

Pseudocode for the main thread:

1. Initialize:

- Parse input arguments: teamA, teamB
- Validate inputs: check if teamA and teamB are even and their sum is divisible by 4
- If invalid, print error and exit

2. Initialize synchronization mechanisms:

- Initialize custom semaphores for Team A and Team B
- Initialize a mutex lock
- Initialize a barrier with a count equal to the sum of teamA and teamB

3. Create fan threads:

- For each fan in teamA and teamB:
 - Create a thread executing 'threads' function
 - Pass a 'personal' struct with group ('A' or 'B') and driver status ('-')

4. Wait for fan threads:

- For each created thread:
 - Wait (join) until the thread completes

5. Clean up:

- Destroy custom semaphores for Team A and Team B
- Destroy the mutex lock and barrier

6. Print "The main terminates" and exit

Pseudocode for the fan thread:

1. Lock the mutex
2. Determine the team (A or B) of the fan
3. Determine if the fan becomes a driver:
 - If fan's team count reaches 4:
 - Set the fan as a driver
 - Reset the team count
 - Post (signal) the semaphore of the fan's team 3 times
 - Else if both teams have at least 2 fans each:
 - Set the fan as a driver
 - Reset both team counts
 - Post the semaphore of the fan's team once
 - Post the semaphore of the other team twice
 - Else:
 - Wait on the semaphore of the fan's team
 - (When signaled, continue)
4. Unlock the mutex
5. Wait at the barrier until all fans find their spots
6. If the fan is a driver:
 - Lock the mutex
 - Increment the global captainID
 - Print "end" message with the captainID
 - Unlock the mutex
7. Exit the thread

SYNCHRONIZATION:

Custom Semaphore:

- Pivotal in managing access to shared resources, specifically the allocation of fans to teams.
- The custom semaphore, incorporating a condition variable and a mutex, is manipulated through `custom_sem_post` and `custom_sem_wait` to regulate thread access and ensure orderly progression.

Mutex (lock):

- A critical component to ensure exclusive access to shared variables such as the team counts and the captainID. This mutex is vital in preventing data races and ensuring consistent state changes within the program.

Barrier:

- A key mechanism to align thread execution at a specific point in the program.
- It is particularly important in ensuring all threads representing fans reach a common stage (finding a spot in a car) before any thread (captain) proceeds to drive the car.

Condition Variable & Custom Semaphores:

- These are used to efficiently block threads when a resource is unavailable and wake them up when it becomes available, thereby minimizing busy-waiting and enhancing the program's overall efficiency.

CORRECTNESS:

Main Thread Correctness:

- The main thread is meticulously designed to first validate the input arguments for the required numerical conditions.
- Post-validation, it initiates the creation of child threads and waits for their completion, thus ensuring no premature termination of the program.
- The final action of the main thread is to print "The main terminates," confirming the orderly conclusion of all operations.

Fan Thread Correctness:

- Each fan thread goes through an ordered sequence of operations (init, mid, end).
- Semaphore operations and mutex locks/unlocks ensure that no more than 4 fans (2 from each team) find spots in a car between two end prints.
- The use of a global captainID with mutex protection guarantees unique and sequential captain IDs.
- Barrier synchronization ensures that all threads reach the mid state before any end is printed.

Valid Combinations:

- The program guarantees that for every fan thread, the init, mid, and end messages are printed in the correct sequence. The total count of these messages aligns with the specified criteria: init and mid messages are equal to the sum of numA and numB, while end messages are a quarter of this sum.
- The implementation ensures that there are exactly four mid messages between each end message, and prior to the first end message. The teams represented in these messages are in valid combinations as per the team count checks. This aspect is critical for demonstrating that the team assembly is not only correct but also adheres to the problem's logical constraints.
- The mechanism for assigning unique car IDs to captains (incrementing a global counter) is managed within a mutually exclusive block. This approach guarantees that each captain is assigned a distinct and sequentially increasing ID, thereby fulfilling a key requirement of the problem statement.