# Chapter 7

**Entity Relationship Model** 

## Entity-Relationship Model

- Design Process
- Modeling
- Constraints
- E-R Diagram
- Design Issues
- Weak Entity Sets
- Extended E-R Features
- Design of the Bank Database
- Reduction to Relation Schemas
- Database Design
- UML

#### Modeling

- A database can be modeled as:
  - A collection of entities,
  - Relationship among entities.
- An entity is an object that exists and is distinguishable from other objects.
  - Example: specific person, company, event, plant
- Entities have attributes
  - Example: people have names and addresses
- An entity set is a set of entities of the same type that share the same properties.
  - Example: set of all persons, companies, trees, holidays

## Entity Sets instructor and student

| Instructor_ID | Instructor_Name |
|---------------|-----------------|
| 76766         | Crick           |
| 45565         | Katz            |
| 10101         | Srinivasan      |
| 98345         | Kim             |
| 76543         | Singh           |
| 22222         | Einstein        |

instructor

| Student_Name |
|--------------|
| Tanaka       |
| Shankar      |
| Zhang        |
| Brown        |
| Aoi          |
| Chavez       |
| Peltier      |
|              |

student

### Relationship Sets

 A relationship is an association among several entities Example:

| 44553 (Peltier) | Advisor          | 22222 (Einstein)  |
|-----------------|------------------|-------------------|
| student entity  | relationship set | instructor entity |

 A relationship set is a mathematical relation among n ≥ 2 entities, each taken from entity sets

$$\{(e_1, e_2, ..., e_n) \mid e_1 \in E_1, e_2 \in E_2, ..., e_n \in E_n\}$$
  
where  $(e_1, e_2, ..., e_n)$  is a relationship

• Example : (44553, 22222) ∈ advisor

## Relationship Set advisor

| Instructor_ID | Instructor_Name |
|---------------|-----------------|
| 76766         | Crick           |
| 45565         | Katz            |
| 10101         | Srinivasan      |
| 98345         | Kim             |
| 76543         | Singh           |
| 22222         | Einstein        |

| Student_ID | Student_Name |
|------------|--------------|
| 98988      | Tanaka       |
| 12345      | Shankar      |
| 00128      | Zhang        |
| 76543      | Brown        |
| 76653      | Aoi          |
| 23121      | Chavez       |
| 44553      | Peltier      |
|            |              |

instructor

student

### Relationship Sets (cont)

• At attribute can also be a property of a relationship set

• For instance, the advisor relationship set between entity sets instructor and student may have the attribute date which tracks when the student started being

associated with the advisor

|               |                 |     | Student_ID | Student_Name |
|---------------|-----------------|-----|------------|--------------|
| Instructor_ID | Instructor_Name |     | 98988      | Tanaka       |
| 76766         | Crick           | 9 / | 12345      | Shankar      |
| 45565         | Katz            |     | 00128      | Zhang        |
| 10101         | Srinivasan      |     | 76543      | Brown        |
| 98345         | Kim             |     | 76653      | Aoi          |
| 76543         | Singh           |     | 23121      | Chavez       |
| 22222         | Einstein        |     | 44553      | Peltier      |
| inst          | ructor          |     |            | student      |

## Degree of a Relationship Set

- Binary relationship
  - Involve two entity sets (or degree two).
  - Most relationship sets in a database system are binary.
- Relationships between more than two entity sets are rare. Most relationships are binary. (More on this later.)
  - Example: *students* work on research *projects* under the guidance of an *instructor*.
  - Relationship proj\_guide is a ternary relationship between instructor, student, and project

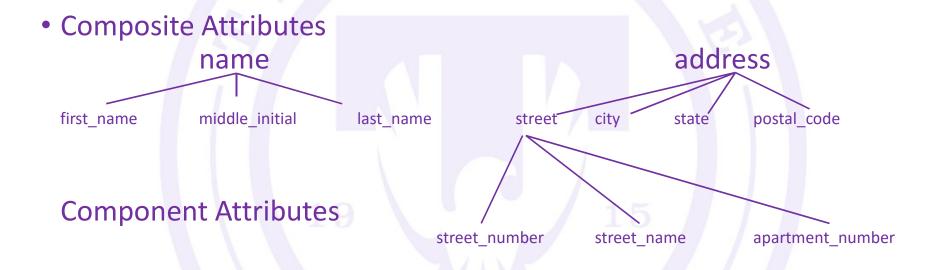
#### Attributes

- An entity is represented by a set of attributes, that is descriptive properties possessed by all members of an entity set.
- Example:

```
instructor = (ID, name, street, city, salary )
course= (course_id, title, credits)
```

- Domain the set of permitted values for each attribute
- Attribute types:
  - Simple and composite attributes.
  - Single-valued and multivalued attributes
    - Example: multivalued attribute: phone\_numbers
  - Derived attributes
    - Can be computed from other attributes
    - Example: age, given date\_of\_birth

## Composite Attributes

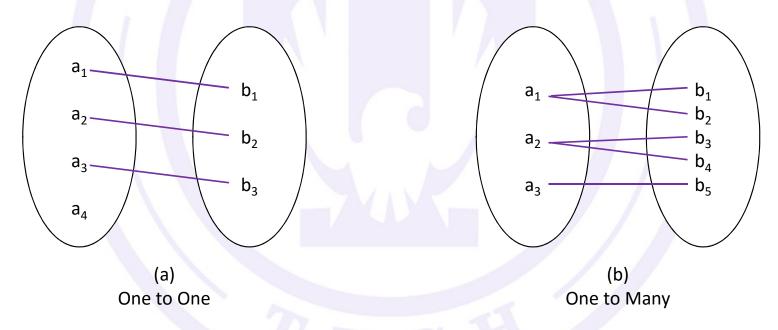


### Mapping Cardinality Constraints

- Express the number of entities to which another entity can be associated via a relationship set.
- Most useful in describing binary relationship sets.
- For a binary relationship set the mapping cardinality must be one of the following types:
  - One to one
  - One to many
  - Many to one
  - Many to many

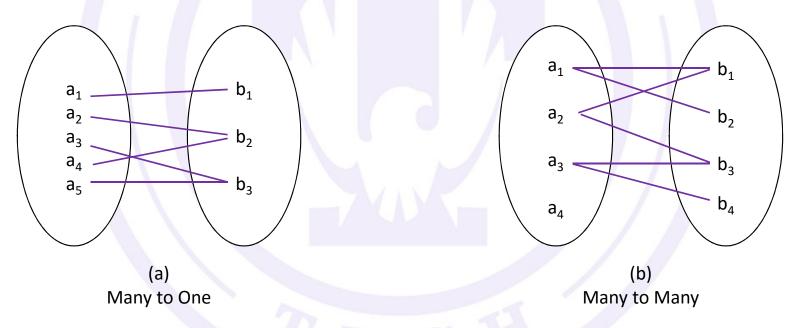
## Mapping Cardinalities

 Note: Some elements in A and B may not be mapped to any elements in the other set



## Mapping Cardinalities

 Note: Some elements in A and B may not be mapped to any elements in the other set



#### Keys

- A super key of an entity set is a set of one or more attributes whose values uniquely determine each entity.
- A candidate key of an entity set is a minimal super key
  - *ID* is candidate key of *instructor*
  - course\_id is candidate key of course
- Although several candidate keys may exist, one of the candidate keys is selected to be the primary key.

#### Keys for Relationship Sets

- The combination of primary keys of the participating entity sets forms a super key of a relationship set.
  - (s\_id, i\_id) is the super key of advisor
  - NOTE: this means a pair of entity sets can have at most one relationship in a particular relationship set.
    - Example: if we wish to track multiple meeting dates between a student and her advisor, we cannot assume a relationship for each meeting. We can use a multivalued attribute though
- Must consider the mapping cardinality of the relationship set when deciding what are the candidate keys
- Need to consider semantics of relationship set in selecting the primary key in case of more than one candidate key

#### Redundant Attributes

- Suppose we have entity sets
  - instructor, with attributes including dept\_name
  - department
    - and a relationship
  - inst\_dept relating instructor and department
- Attribute dept\_name in entity instructor is redundant since there is an explicit relationship inst\_dept which relates instructors to departments
  - The attribute replicates information present in the relationship, and should be removed from instructor
  - BUT: when converting back to tables, in some cases the attribute gets reintroduced, as we will see.

#### E-R Diagrams



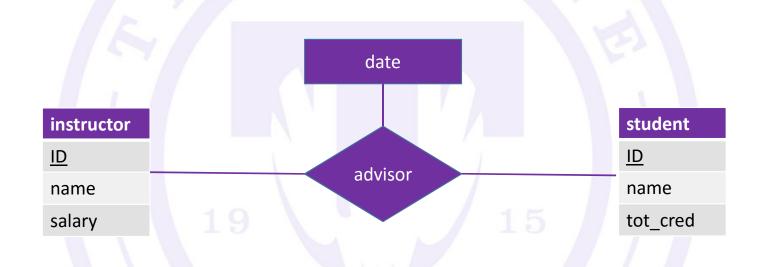
- Rectangles represent entity sets
- Diamonds represent relationship sets
- Attributes listed inside entity rectangle
- Underline indicates primary key attributes

Entity with Composite, Multivalued and Derived Attributes

Instructor

ID name first\_name middle\_initial last\_name address street street\_number street\_name apt number city state zip { phone\_number } date\_of\_birth age()

# Relationship Sets with Attributes



#### Roles

- Entity sets of a relationship need not be distinct
  - Each occurrence of an entity set plays a "role" in the relationship
- The labels "course\_id" and "prereq\_id" are called roles



### **Cardinality Constraints**

- We express cardinality constraints by drawing either a directed line (→), signifying "one," or an undirected line (—), signifying "many," between the relationship set and the entity set.
- One-to-one relationship:
- A student is associated with at most one instructor via the relationship advisor
- A student is associated with at most one department via stud\_dept

## One-to-One Relationship

- One-to-one relationship between an instructor and a student
  - An instructor is associated with at most one student via advisor
  - A student is associated with at most one instructor via advisor



## One-to-Many Relationship

- One-to-many relationship between an instructor and a student
  - An instructor is associated with several (including 0) students via advisor
  - A student is associated with at most one instructor via advisor



## Many-to-One Relationship

- Many-to-one relationship between an instructor and a student
  - An instructor is associated with at most one students via advisor
  - A student is associated with several (including 0) instructor via advisor



## Many-to-Many Relationship

- Many-to-many relationship between an instructor and a student
  - An instructor is associated with several (including 0) students via advisor
  - A student is associated with several (including 0) instructor via advisor



# Participation of an Entity Set in a Relationship Set

- Total participation (indicated by double line): every entity in the entity set participates in at least one relationship in the relationship set
  - E.g., participation of section in sec\_course is total
    - every section must have an associated course
- Partial participation: some entities may not participate in any relationship in the relationship set
  - Example: participation of instructor in advisor is partial

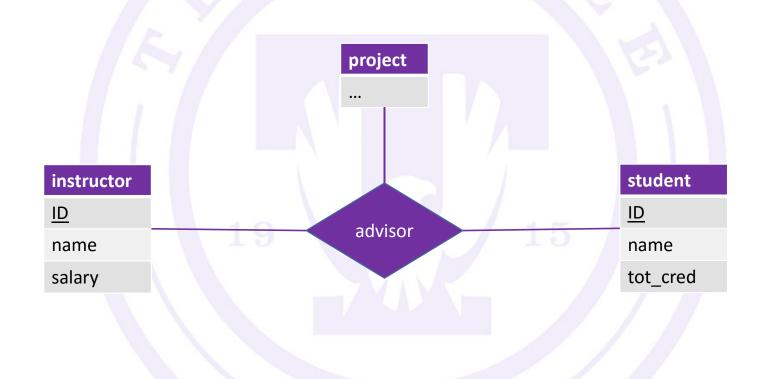


## Alternative Notation for Cardinality Limits

Cardinality limits can also express participation constraints



## E-R Diagram with a Ternary Relationship



# Cardinality Constraints on Ternary Relationship

- We allow at most one arrow out of a ternary (or greater degree) relationship to indicate a cardinality constraint
- E.g., an arrow from proj\_guide to instructor indicates each student has at most one guide for a project
- If there is more than one arrow, there are two ways of defining the meaning.
  - E.g., a ternary relationship R between A, B and C with arrows to B and C could mean
    - 1. each A entity is associated with a unique entity from B and C or
- 2. each pair of entities from (A, B) is associated with a unique C entity, and each pair

(A, C) is associated with a unique B

- Each alternative has been used in different formalisms
- To avoid confusion we outlaw more than one arrow

#### Weak Entity Sets

- An entity set that does not have a primary key is referred to as a weak entity set.
- The existence of a weak entity set depends on the existence of a identifying entity set
  - It must relate to the identifying entity set via a total, one-to-many relationship set from the identifying to the weak entity set
  - Identifying relationship depicted using a double diamond
- The discriminator (or partial key) of a weak entity set is the set of attributes that distinguishes among all the entities of a weak entity set.
- The primary key of a weak entity set is formed by the primary key of the strong entity set on which the weak entity set is existence dependent, plus the weak entity set's discriminator.

#### Weak Entity Sets

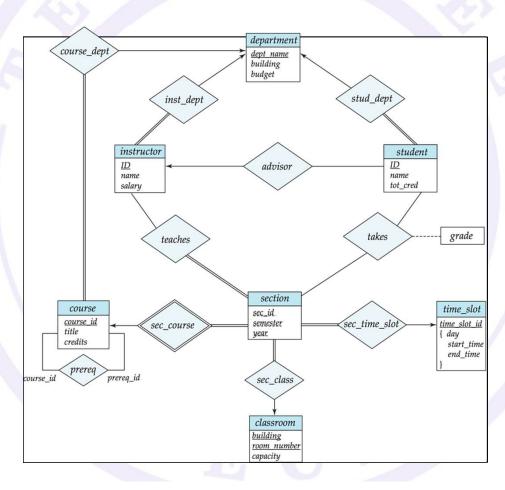
- We underline the discriminator of a weak entity set with a dashed line.
- We put the identifying relationship of a weak entity in a double diamond.
- Primary key for section (course\_id, sec\_id, semester, year)



#### Weak Entity Sets

- Note: the primary key of the strong entity set is not explicitly stored with the weak entity set, since it is implicit in the identifying relationship.
- If course\_id were explicitly stored, section could be made a strong entity, but then the relationship between section and course would be duplicated by an implicit relationship defined by the attribute course\_id common to course and section

## E-R Diagram for a University Enterprise



#### Reduction to Relation Schemas

- Entity sets and relationship sets can be expressed uniformly as relation schemas that represent the contents of the database.
- A database which conforms to an E-R diagram can be represented by a collection of schemas.
- For each entity set and relationship set there is a unique schema that is assigned the name of the corresponding entity set or relationship set.
- Each schema has a number of columns (generally corresponding to attributes), which have unique names.

# Representing Entity Sets with Simple Attributes

- A strong entity set reduces to a schema with the same attributes student(<u>ID</u>, name, tot\_cred)
- A weak entity set becomes a table that includes a column for the primary key of the identifying strong entity set section (<u>course id</u>, <u>sec id</u>, <u>sem</u>, <u>year</u>)



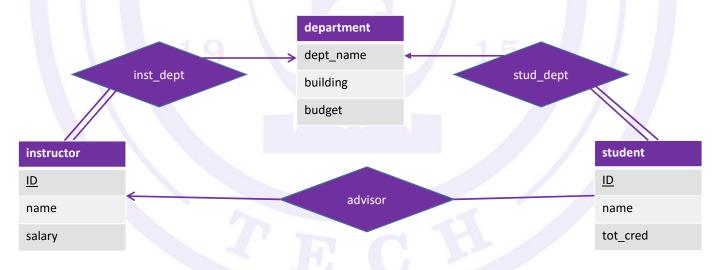
## Representing Relationship Sets

- A many-to-many relationship set is represented as a schema with attributes for the primary keys of the two participating entity sets, and any descriptive attributes of the relationship set.
- Example: schema for relationship set advisor advisor = (s\_id, i\_id)



# Redundancy of Schemas

- Many-to-one and one-to-many relationship sets that are total on the many-side can be represented by adding an extra attribute to the "many" side, containing the primary key of the "one" side
- Example: Instead of creating a schema for relationship set inst\_dept, add an attribute dept\_name to the schema arising from entity set instructor



## Redundancy of Schemas

- For one-to-one relationship sets, either side can be chosen to act as the "many" side
  - That is, extra attribute can be added to either of the tables corresponding to the two entity sets
- If participation is partial on the "many" side, replacing a schema by an extra attribute in the schema corresponding to the "many" side could result in null values
- The schema corresponding to a relationship set linking a weak entity set to its identifying strong entity set is redundant.
  - Example: The section schema already contains the attributes that would appear in the sec\_course schema

# Composite and Multivalued Attributes

#### instructor ID name first\_name middle initial last name address street street number street name apt number city state zip { phone\_number } date\_of\_birth age ()

- Composite attributes are flattened out by creating a separate attribute for each component attribute
  - Example: given entity set instructor with composite attribute name with component attributes first\_name and last\_name the schema corresponding to the entity set has two attributes name\_first\_name and name\_last\_name
    - Prefix omitted if there is no ambiguity
- Ignoring multivalued attributes, extended instructor schema is

### Composite and Multivalued Attributes

- A multivalued attribute M of an entity E is represented by a separate schema EM
  - Schema EM has attributes corresponding to the primary key of E and an attribute corresponding to multivalued attribute M
  - Example: Multivalued attribute phone\_number of instructor is represented by a schema:
    - inst\_phone= (ID, phone\_number)
  - Each value of the multivalued attribute maps to a separate tuple of the relation on schema EM
    - For example, an instructor entity with primary key 22222 and phone numbers 456-7890 and 123-4567 maps to two tuples:
      - (22222, 456-7890) and (22222, 123-4567)

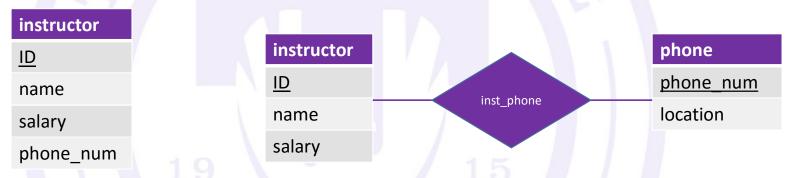
#### Multivalued Attributes

- Special case:entity time\_slot has only one attribute other than the primary-key attribute, and that attribute is multivalued
  - Optimization: Don't create the relation corresponding to the entity, just create the one corresponding to the multivalued attribute
  - time\_slot(time\_slot\_id, day, start\_time, end\_time)
  - Caveat: time\_slot attribute of section (from sec\_time\_slot) cannot be a foreign key due to this optimization



### Design Issues

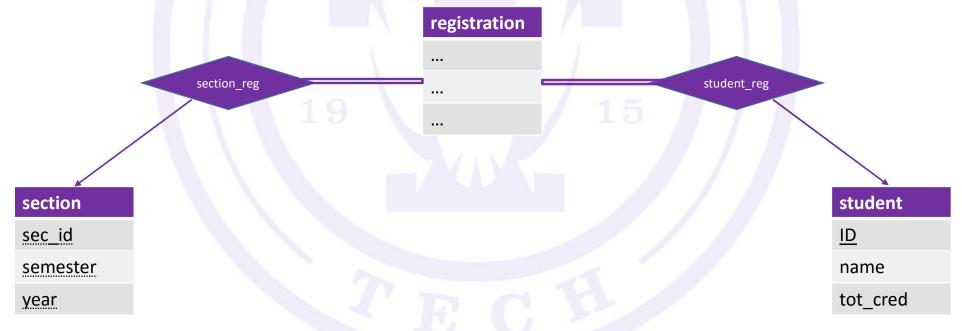
Use of entity sets v attributes



• Use of phone as an entity allows extra information about phone numbers (plus multiple phone numbers)

#### Design Issues

Use of entity sets vs. relationship sets
 Possible guideline is to designate a relationship set to describe an action that occurs between entities



#### Design Issues

- Binary versus n-ary relationship sets
  - Although it is possible to replace any nonbinary (n-ary, for n > 2) relationship set by a number of distinct binary relationship sets, a n-ary relationship set shows more clearly that several entities participate in a single relationship.
- Placement of relationship attributes
   e.g., attribute date as attribute of advisor or as attribute of student

## Binary v Non-Binary Relationships

- Some relationships that appear to be non-binary may be better represented using binary relationships
  - E.g., A ternary relationship parents, relating a child to his/her father and mother, is best replaced by two binary relationships, father and mother
    - Using two binary relationships allows partial information (e.g., only mother being know)
  - But there are some relationships that are naturally non-binary
    - Example: proj guide

# Converting Non-Binary Relationships to Binary Form

- In general, any non-binary relationship can be represented using binary relationships by creating an artificial entity set.
  - Replace R between entity sets A, B and C by an entity set E, and three relationship sets:
  - R<sub>A</sub>, relating E and A
     R<sub>B</sub>, relating E and B
     R<sub>C</sub>, relating E and C
  - Create a special identifying attribute for E
  - Add any attributes of R to E
  - For each relationship (a<sub>i</sub> , b<sub>i</sub> , c<sub>i</sub>) in R, create
    1. a new entity e<sub>i</sub> in the entity set E
    2. add (e<sub>i</sub> , a<sub>i</sub> ) to R<sub>A</sub>
    3. add (e<sub>i</sub> , b<sub>i</sub> ) to R<sub>B</sub>

    A

    B
    R
    C

    B
    R
    C

    C

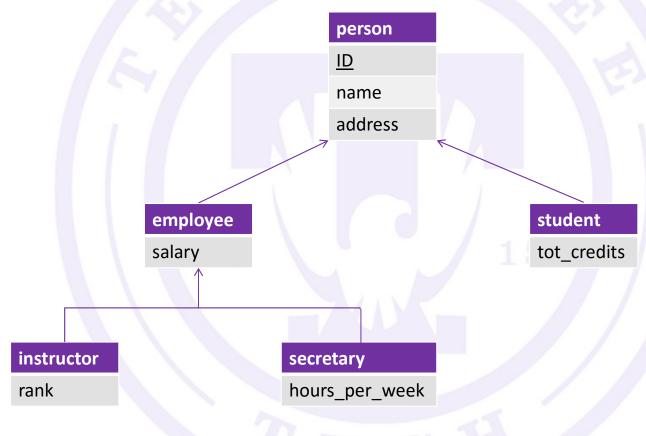
# Converting Non-Binary Relationships

- Also need to translate constraints
  - Translating all constraints may not be possible
  - There may be instances in the translated schema that cannot correspond to any instance of R
    - Exercise: add constraints to the relationships R<sub>A</sub>, R<sub>B</sub> and R<sub>C</sub> to ensure that a newly created entity corresponds to exactly one entity in each of entity sets A, B and C
  - We can avoid creating an identifying attribute by making E a weak entity set (described shortly) identified by the three relationship sets

#### Extended E-R Features: Specialization

- Top-down design process; we designate subgroupings within an entity set that are distinctive from other entities in the set.
- These subgroupings become lower-level entity sets that have attributes or participate in relationships that do not apply to the higher-level entity set.
- Depicted by a triangle component labeled ISA (E.g., instructor "is a" person).
- Attribute inheritance a lower-level entity set inherits all the attributes and relationship participation of the higher-level entity set to which it is linked.

# Specialization Example



#### Extended ER Features: Generalization

- A bottom-up design process combine a number of entity sets that share the same features into a higher-level entity set.
- Specialization and generalization are simple inversions of each other; they are represented in an E-R diagram in the same way.
- The terms specialization and generalization are used interchangeably.

#### Specialization and Generalization

- Can have multiple specializations of an entity set based on different features.
- E.g., permanent\_employee vs. temporary\_employee, in addition to instructor vs. secretary
- Each particular employee would be
  - A member of one of permanent\_employee or temporary\_employee,
  - And also a member of one of instructor, secretary
- The ISA relationship also referred to as superclass subclass relationship

# Design Constraints on Specialization and Generalization

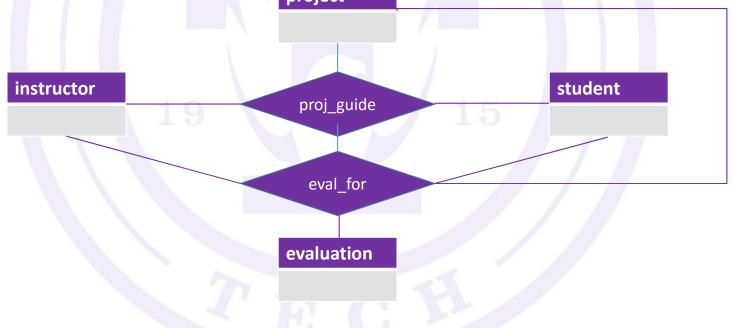
- Constraint on which entities can be members of a given lower-level entity set.
  - Condition-defined
    - Example: all customers over 65 years are members of senior-citizen entity set; senior-citizen ISA person.
  - User-defined
- Constraint on whether or not entities may belong to more than one lowerlevel entity set within a single generalization.
  - Disjoint
    - An entity can belong to only one lower-level entity set
    - Noted in E-R diagram by having multiple lower-level entity sets link to the same triangle
  - Overlapping
    - An entity can belong to more than one lower-level entity set

# Design Constraints on Specialization and Generalization

- Completeness constraint specifies whether or not an entity in the higher-level entity set must belong to at least one of the lower-level entity sets within a generalization.
  - Total: an entity must belong to one of the lower-level entity sets
  - Partial: an entity need not belong to one of the lower-level entity sets

#### Aggregation

- Consider the ternary relationship proj\_guide, which we saw earlier
- Suppose we want to record evaluations of a student by a guide on a project



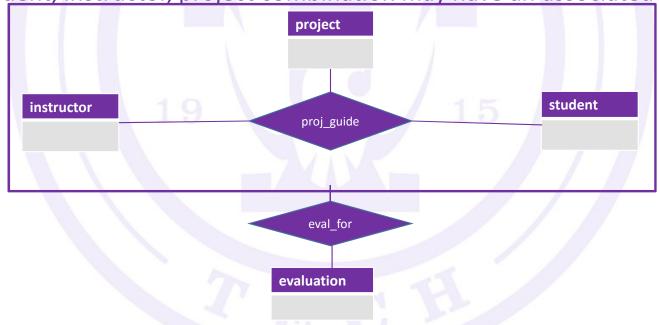
#### Aggregation

- Relationship sets eval\_for and proj\_guide represent overlapping information
  - Every eval\_for relationship corresponds to a proj\_guide relationship
  - However, some proj\_guide relationships may not correspond to any eval\_for relationships
    - So we can't discard the proj\_guide relationship
- Eliminate this redundancy via aggregation
  - Treat relationship as an abstract entity
  - Allows relationships between relationships
  - Abstraction of relationship into new entity

#### Aggregation

- Without introducing redundancy, the following diagram represents:
  - A student is guided by a particular instructor on a particular project

- A student, instructor, project combination may have an associated evaluation



### Representing Specialization via Schemas

#### Method 1:

- Form a schema for the higher-level entity
- Form a schema for each lower-level entity set, include primary key of higher-level entity set and local attributes

| schema   | attributes                       |
|----------|----------------------------------|
| person   | ID, name, street, city           |
| student  | ID, name, street, city, tot_cred |
| employee | ID, name, street, city, salary   |

 Drawback: getting information about, an employee requires accessing two relations, the one corresponding to the low-level schema and the one corresponding to the high-level schema

### Representing Specialization as Schemas

- Method 2:
  - Form a schema for each entity set with all local and inherited attributes

| schema   | attributes             |
|----------|------------------------|
| person   | ID, name, street, city |
| student  | ID, tot_cred           |
| employee | ID, salary             |

- If specialization is total, the schema for the generalized entity set (*person*) not required to store information
  - Can be defined as a "view" relation containing union of specialization relations
  - But explicit schema may still be needed for foreign key constraints
- Drawback: name, street and city may be stored redundantly for people who are both students and employees

# Schemas Corresponding to Aggregation

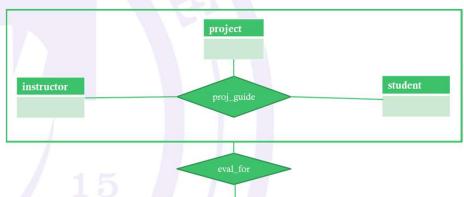
- To represent aggregation, create a schema containing
  - Primary key of the aggregated relationship,
  - The primary key of the associated entity set
  - Any descriptive attributes

# Schemas Corresponding to Aggregation

 For example, to represent aggregation manages between relationship works\_on and entity set manager, create a schema eval\_for (s\_ID, project\_id, i\_ID,

 Schema proj\_guide is redundant provided we are willing to store null values for attribute manager\_name in relation on schema manages

evaluation id)

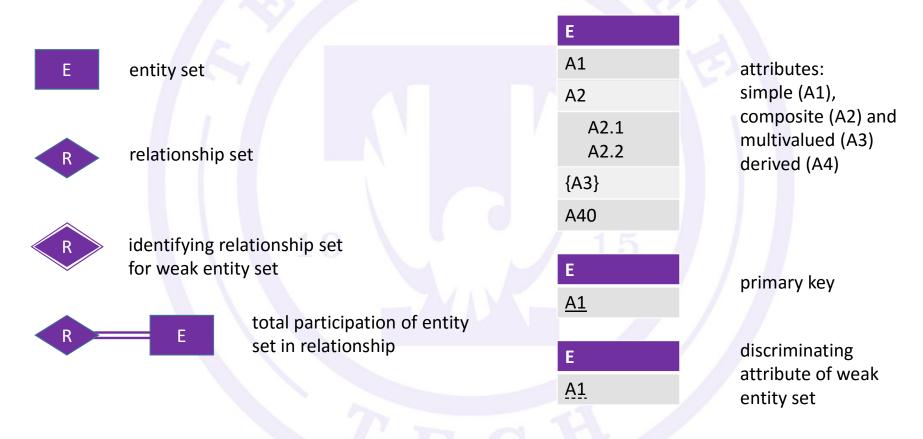


evaluation

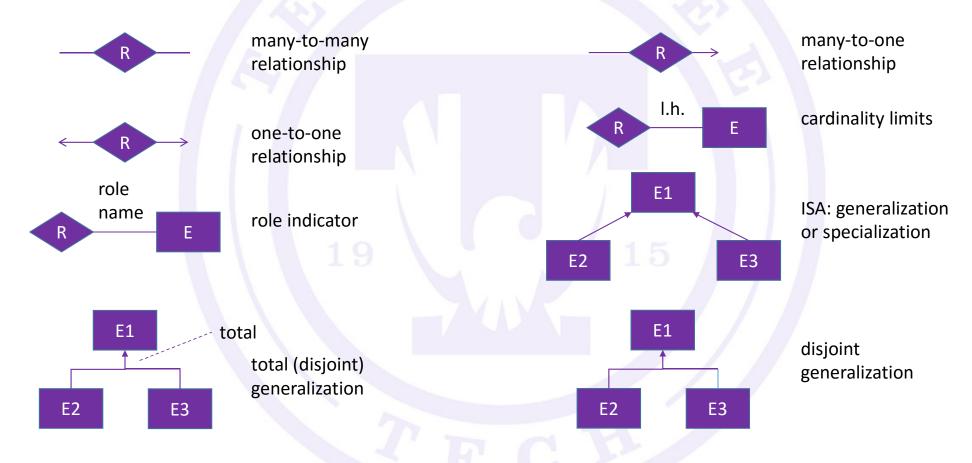
#### E-R Design Decisions

- The use of an attribute or entity set to represent an object.
- Whether a real-world concept is best expressed by an entity set or a relationship set.
- The use of a ternary relationship versus a pair of binary relationships.
- The use of a strong or weak entity set.
- The use of specialization/generalization contributes to modularity in the design.
- The use of aggregation can treat the aggregate entity set as a single unit without concern for the details of its internal structure.

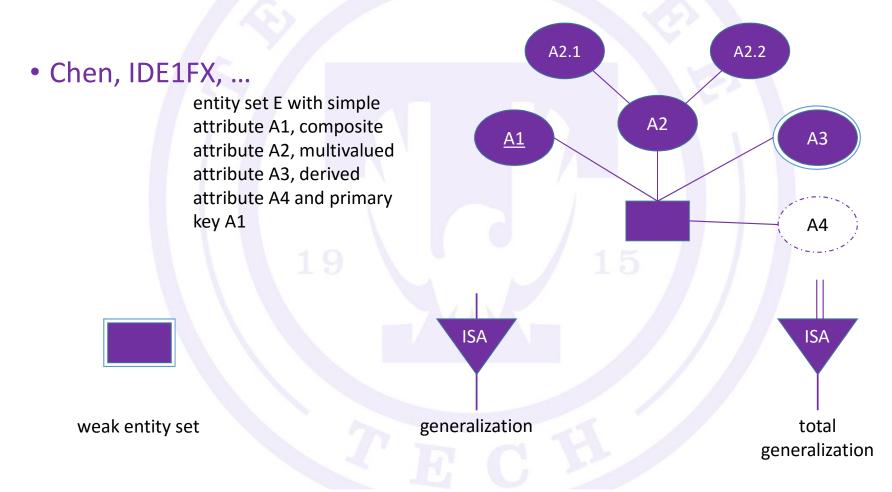
# Summary of Symbols Used in E-R Notation



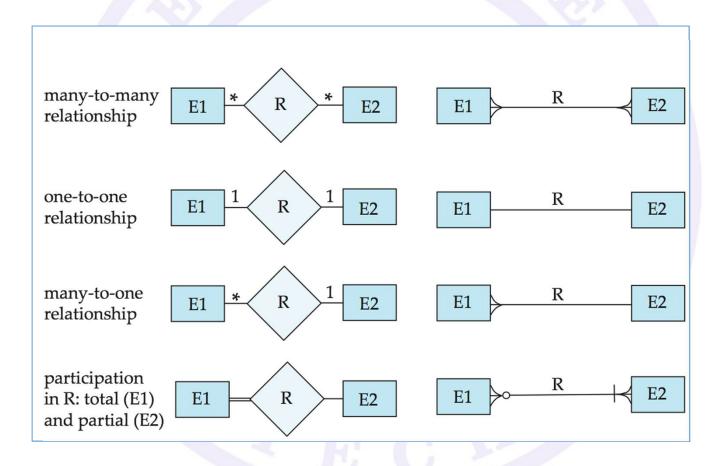
## Summary of Symbols Used in E-R Notation



#### Alternative ER Notations



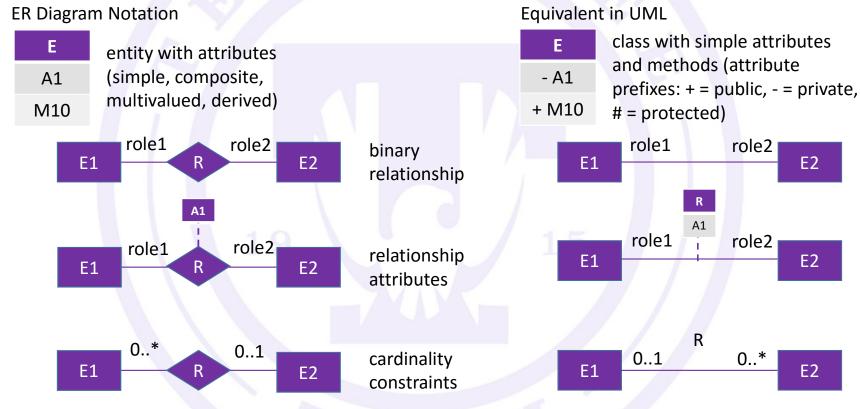
#### Alternative ER Notations



#### **UML**

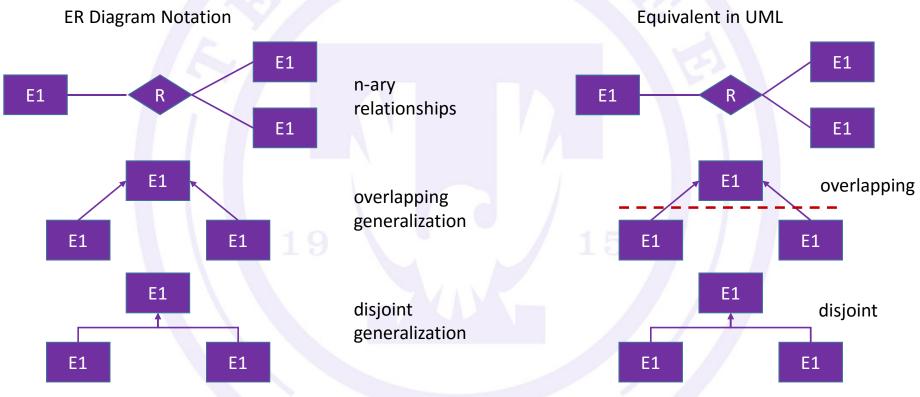
- UML: Unified Modeling Language
  - UML has many components to graphically model different aspects of an entire software system
  - UML Class Diagrams correspond to E-R Diagram, but several differences.

#### ER v UML Class Diagrams



<sup>\*</sup> Note the reversal of position in cardinality constraints depiction

#### ER v UML Class Diagrams



<sup>\*</sup>Generalization can use merged or separate arrows independent of disjoint/overlapping

#### ER v UML Class Diagrams

- Binary relationship sets are represented in UML by just drawing a line connecting the entity sets. The relationship set name is written adjacent to the line.
- The role played by an entity set in a relationship set may also be specified by writing the role name on the line, adjacent to the entity set.
- The relationship set name may alternatively be written in a box, along with attributes of the relationship set, and the box is connected, using a dotted line, to the line depicting the relationship set.

End of Chapter 7