

# Integer & network Linear Programming Assignment

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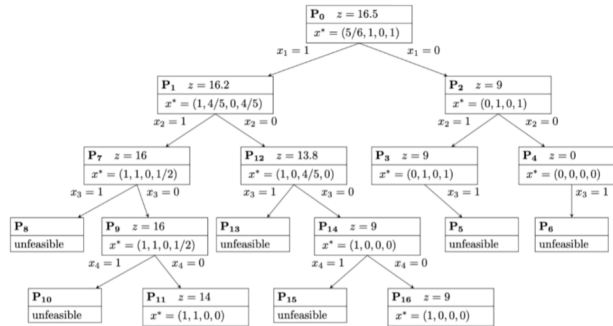
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## Problem 1

Consider the following ILP:

$$\begin{aligned}
 \max \quad & 9x_1 + 5x_2 + 6x_3 + 4x_4 \\
 \text{s.t.} \quad & 6x_1 + 3x_2 + 5x_3 + 2x_4 \leq 10, \\
 & x_3 + x_4 \leq 1, \\
 & -x_1 + x_3 \leq 0, \\
 & -x_2 + x_4 \leq 0 \\
 & x_1, x_2, x_3, x_4 \in \{0, 1\}
 \end{aligned}$$

The following tree represents the solutions of all possible relaxations of the problem in which no sub-problem has been excluded (fathoming).



Suppose that the Branch and Bound (BB) algorithm applies to this problem. Also, let's suppose that the algorithm visits the sub-problems in the following order  $P_0, P_1, \dots, P_{16}$ . Clearly, the algorithm does not visit all nodes.

- 1) Determine the nodes that will be visited by the BB algorithm and for each of them get the upper and lower limit deduced by the algorithm in the execution.

The algorithm starts from  $P_0$  where the only non integer variable is  $X_1$ , since in an integer problem we can not accept variables solutions unless all of them are integer, the Branch and Bound algorithm (BB) will split the starting problem in two branches:  $P_1$  by considering  $X_1 = 1$  and  $P_2$  in the case where  $X_1 = 0$ .

The Upper bound (Ub) and the Lower bound (Lb) for each node are computed in the following way: the Ub corresponds to the value of  $z$  that is the relaxed solution of the problem, instead the Lb is computed by solving  $z$  rounding down the non integer values for the decision variables to the closest integer value (0 in this case), therefore with the Lb is presented

the pessimistic minimal solution for the problem.

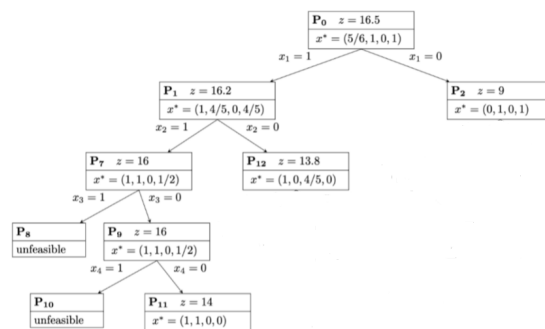
As presented above the first node explore by the  $BB$  is  $P_0$  where  $Ub = 16.5$  and  $Lb = 9$ .

Let's consider now  $P_2$ , where all the coefficients for the decision variables are now integer, since  $Ub = Lb = 9$  meaning that all the sub-problem starting from  $P_2$  will not find a better solution and it is completely worthless for the algorithm to proceed further, then 9 can be considered as a possible candidate for the optimal solution of the problem.

Considering  $P_1$ , in the left side of the graph, it present  $Ub = 16.2$  and  $Lb = 9$ , right after the  $BB$  will visit  $P_7$  having the same  $Lb$  computed in the previous node. At this point  $P_8$  is unfeasible then the algorithm will look for another optimal solution in  $P_9$ : this node present  $Ub = 16$  and  $Lb = 14$ , therefore being  $Lb_{(P_9)} > Lb_{(P_2)}$  the algorithm will no longer consider  $P_2 = 9$  as a candidate optimal solution for the problem because the algorithm will surely find a better solution for the problem.  $P_{10}$  is unfeasible too then it will not be considered by  $BB$ , while  $P_{11}$  with  $Ub = Lb = 14$  and all the coefficients found are integer so for now this can be considered as a candidate for an optimal solution.

Looking at  $P_{12}$  the  $Ub = 13.8$ , so being  $Ub_{(P_{12})} < Ub_{(P_{11})}$  the algorithm will stop here because there will not be a better solution for the problem. So the algorithm will find the optimal solution in the sub-problem  $P_{11}$  having  $z^* = 14$  and  $x^* = (1, 1, 0, 0)$

In the Immage below is represented the path chosen by the  $BB$  algorithm.



2) Solve the problem with an ILP solver and check the value of the objective function matches the one found at point 1.

```
# Import packages.
```

```
library(lpSolve)
```

```
library(lpSolveAPI)
```

```
library(DiagrammerR)
```

```
# Initializing the lp model with 0 constraints and 4 decision variables.
```

```
model <- make.lp(0, 4)
```

```
# Maximization problem.
```

```
lp.control(model, sense = 'max')
```

```
# Objective function.
```

```
set.objfn(model, c(9, 5, 6, 4))
```

```
# Constrains.
```

```
add.constraint(model, c(6, 3, 5, 2), "<=", 10)
```

```
add.constraint(model, c(1, 1), "<=", 1, c(3, 4))
```

```
add.constraint(model, c(-1, 1), "<=", 0, c(1, 3))
```

```
add.constraint(model, c(-1, 1), "<=", 0, c(2, 4))
```

```
# Set decision variables value to be binary integer {0, 1}.
```

```
set.type(model, c(1:4), "binary")
```

```
# 0 means that the model was solved correctly.
solve(model)
```

```
## [1] 0
```

```
get.objective(model)
```

```
## [1] 14
```

```
get.variables(model)
```

```
## [1] 1 1 0 0
```

As expected the maximized value for the problem is 14 and the coefficients found are the same seen in the above graph.

## Problem 2

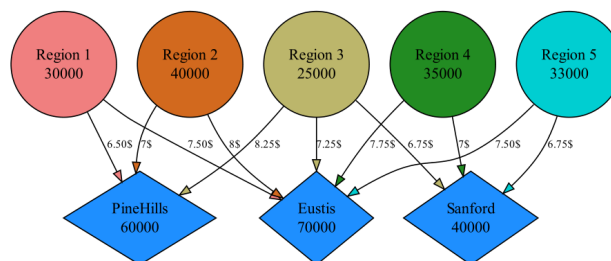
SunNet is a residential Internet Service Provider (ISP) in the central Florida area. Presently, the company operates one centralized facility that all of its clients call into for Internet access.

To improve service, the company is planning to open three satellite offices in the cities of Pine Hills, Eustis, and Sanford. The company has identified five different regions to be serviced by these three offices. The following table summarizes the number of customers in each region, the service capacity at each office, and the monthly average cost per customer for providing the service to each region from each office. Table entries of “n.a.” indicate infeasible region-to-service center combinations.

SunNet would like to determine how many customers from each region to assign to each service center to minimize the total cost.

Region	Pine Hills	Eustis	Sanford	Customers
1	\$6.50	\$7.50	n.a.	30,000
2	\$7.00	\$8.00	n.a.	40,000
3	\$8.25	\$7.25	\$6.75	25,000
4	n.a.	\$7.75	\$7.00	35,000
5	n.a.	\$7.50	\$6.75	33,000
Capacity	60,000	70,000	40,000	

- 1) Draw a network flow model to represent this problem.



The goal about the problem is to find the optimal way to assign the customers between the  $i$  –  $th$  office and the  $j$  –  $th$  region. The decision variable can be explained in this way.

Region	Pine Hills	Eustis	Sanford	Customers
1	$x_{1,1}$	$x_{2,1}$	n.a.	30,000
2	$x_{1,2}$	$x_{2,2}$	n.a.	40,000
3	$x_{1,3}$	$x_{2,3}$	$x_{3,3}$	25,000
4	n.a.	$x_{2,4}$	$x_{3,4}$	35,000
5	n.a.	$x_{2,5}$	$x_{3,5}$	33,000
Capacity	60,000	70,000	40,000	

The goal is to minimize the average cost of the services for each customer. The objective function is presented below since the average costs coefficients  $c_{i,j}$  are given in the first table:

$$MIN : \sum_{i=1}^3 \sum_{j=1}^5 c_{i,j} x_{i,j}$$

2) Implement your model and solve it.

```
model <- make.lp(0, 11)
lp.control(model, sense = 'min')

# Objective function.
set.objfn(model, c(6.5, 7.5, 7, 8, 8.25, 7.25, 6.75, 7.75, 7, 7.5, 6.75))

# Capacity constraints.
add.constraint(model, c(1, 1, 1), "<=", 60000, c(1, 3, 5))
add.constraint(model, c(1, 1, 1, 1, 1), "<=", 70000, c(2, 4, 6, 8, 10))
add.constraint(model, c(1, 1, 1), "<=", 40000, c(7, 9, 11))

# Customer constraints.
add.constraint(model, c(-1, -1), "<=", -30000, c(1, 2))
add.constraint(model, c(-1, -1), "<=", -40000, c(3, 4))
add.constraint(model, c(-1, -1, -1), "<=", -25000, c(5, 6, 7))
add.constraint(model, c(-1, -1), "<=", -35000, c(8, 9))
add.constraint(model, c(-1, -1), "<=", -33000, c(10, 11))

# Integer programming
set.type(model, c(1:11), "integer")
solve(model)
```

```
## [1] 0
```

3) What is the optimal solution?

```
get.objective(model)
```

```
## [1] 1155000
```

```
get.variables(model)
```

```
## [1] 20000 10000 40000 0 0 25000 0 0 35000 28000 5000
```

```
get.primal.solution(model)[2:8]
```

```
## [1] 60000 63000 40000 -30000 -40000 -25000 -35000
```

From the outputs it is possible to discover that the minimized cost is equal to 1155000\$, and the values for each variable of the objective function can be seen below.

$x_{1,1}$	$x_{2,1}$	$x_{1,2}$	$x_{2,2}$	$x_{1,3}$	$x_{2,3}$	$x_{3,3}$	$x_{2,4}$	$x_{3,4}$	$x_{2,5}$	$x_{3,5}$
20000	10000	40000	0	0	25000	0	0	35000	28000	5000

In the last output it's easy to notice that all the right hand side of the constraints are maximized except for the second capacity constrain where over 70000 only 63000 ( $x_{2,1} + x_{2,2} + x_{2,3} + x_{2,4} + x_{2,5}$ ) were selected.