

Assignment 9

R.3.11,

5, 16, 22, 45, 2, 19, 18, 30, 50, 12, 13, 33

1) 5, 16, 22

2) 5, 16, 22, 45

3) 5, 16, 22, 45, 10, 18, 22, 45

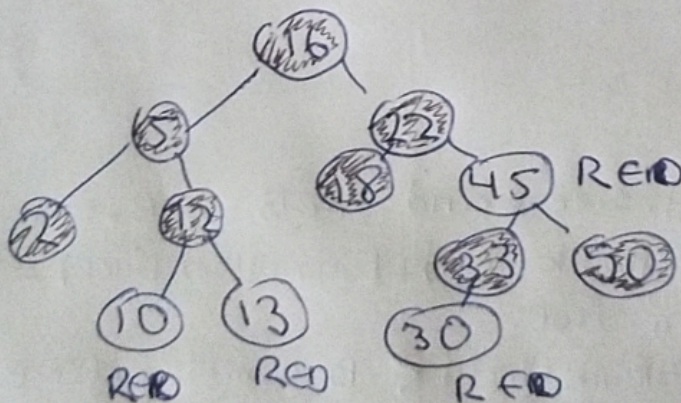
4) 16, 22, 2, 5, 10, 18, 30, 45, 50

5) 5, 16, 22, 2, 10, 12, 18, 30, 45, 50

5, 16, 22, 2, 10, 12, 13, 18, 30, 45, 50

16, 22, 45, 50, 5, 2, 10, 12, 13, 18, 30, 33

• RED - BLACK



R.3.14.

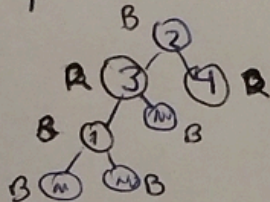
- a Subtree of a red-black tree is itself a red-black tree.

True, because satisfy this

- Every Node is either red or Black
- Every leaf (NIL) is Black
- If a node is red, then both its children are black
- For each node, all simple paths from the node to any descendant leaves contain the same number of black nodes

- the sibling of an external node is either external or it is red

False, because an external node can be either external or black



- Given a red-black tree T , there is a unique $(2,4)$ tree T' associated with T , because $(2,4)$ tree is a binary search tree that satisfy:

- Every node has at least two children and at most four children.
- All leaves are at the same level.
- The keys in all left subtrees are less than the key in the root.
- The keys in all right subtree are greater than the key in the root.

- Given $(2,4)$ tree T , there is a unique red-black tree T' associated with T .

False, because there are multiple ways to color the nodes of T red or black such that the resulting tree is a red-black tree,

Algorithm is ValidAVL(Root)

Input ~~tree~~ is a node Root

if ~~root~~ == null then

return true

leftHeight := getHeight(root.left)

rightHeight := getHeight(root.right)

if (Abs(leftHeight - rightHeight) > 1) then

return false

return is ValidAVL(root.left) and is ValidAVL(root.right)

Algorithm getHeight(root)

Input root is a node

output Integer

if root == null then

return 0

return 1 + Math.max(getHeight(root.left), getHeight(root.right))

Algorithm is Permutation(A, B)

Input A, B are Sequences.

1 If A.size \neq B.size then
return false

1 D := new Dictionary

n for i := 0 to A.size do
n if D.findItem(A.elemAtRank(i)) == null then
n D.insertItem(A.elemAtRank(i), 1)

else

n D.insertItem(A.elemAtRank(i), D.findValue(A.elemAtRank(i)) + 1)

n for i := 0 to B.size do

n count = D.findValue(B.elemAtRank(i))

~~if count == null~~

n If count = null or count = NO-SUCH-KEY then
return false

n ~~if~~ D.insertItem(B.elemAtRank(i), count - 1)

1 return true

Algorithm FindAllInRange (k_1, k_2) •

Input k_1, k_2 are range

$S :=$ New Sequence

findHelper($T, k_1, k_2, T.\text{root}(), S$)

return S .

Algorithm findHelper(T, k_1, k_2, N, S)

if T is External(P) then

$e := P.\text{element}$ return

if $k_1 > e$ then

findHelper($T, k_1, k_2, T.\text{leftChild}(P), S$)

$e := P.\text{element}()$

if $k_1 \leq e \wedge e \leq k_2$ then

$S.\text{insertLast}(e)$

if $e > k_2$ then

findHelper($T, k_1, k_2, T.\text{rightChild}(P), S$)