building effective language parsers for penetration testers

columbus owasp 28 march 2019

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i like breaking software. i've done it for these organizations.









analyzing source code for bugs is hard for many reasons.

lexical ambiguity variable state considerations grammar complexity variable scope knowledge of source / sinks effects of sanitization

before we jump into code parsing, let's explore some fundamentals with english.

"I shot the elephant in my pajamas."

"I shot the elephant in my pajamas."

- The shooter wore pajamas while shooting the elephant.
- 2. The elephant was wearing the shooter's pajamas.

"He drove down the street in the car."

"He drove down the street in the car."

- 1. He was driving a car on the street.
- 2. The street he drove down was actually *in his car*.

"The complex houses married and single soldiers and their families."

"The complex houses married and single soldiers and their families."

1. The housing complex contains both married and single soldiers, as well as their families.

"Buffalo buffalo buffalo buffalo buffalo buffalo buffalo buffalo."

"Buffalo buffalo buffalo buffalo buffalo buffalo buffalo buffalo."

1. Bison from Buffalo, New York, who are intimidated by other bison in Buffalo, New York also happen to intimidate other bison in Buffalo New York. Kids make nutritious snacks Milk drinkers are turning to powder Drunk musician gets nine months in violin case Man eating piranha mistakenly sold as pet fish Include your children when baking cookies Red tape holds up new bridge

linguistic principles: context-free grammars

formal definition of a context-free grammar

A context-free grammar can be defined as $G = (N, \Sigma, R, S)$ where:

- N is a set of non-terminal symbols
- Σ is a set of terminal symbols
- R is a set of rules
- S ∈ N is a distinguished start symbol

A Simplified Context-free Grammar for English

 $G = (N, \Sigma, R, S)$ where:

- **N** = { <u>S</u>, NP, VP, PP, DT, Vi, Vt, NN, IN, PR }
- Σ = { sleeps, saw, man, woman, telescope, the, with, in }
- R is the set of rules (or "derivations") shown below
- S = S

s	\rightarrow	NP	VP
VP	\rightarrow	Vi	
VP	\rightarrow	Vt	NP
VP	→	VP	NP
NP	\rightarrow	PR	
NP	\rightarrow	DT	NN
NP	→	NP	PP
PP	\rightarrow	IN	NP

Vi	\rightarrow	sleeps
Vt	→	saw
NN	→	man
NN	→	woman
NN	→	telescope
DT	→	the
IN	→	with
IN	→	in
PR	→	he

Σ

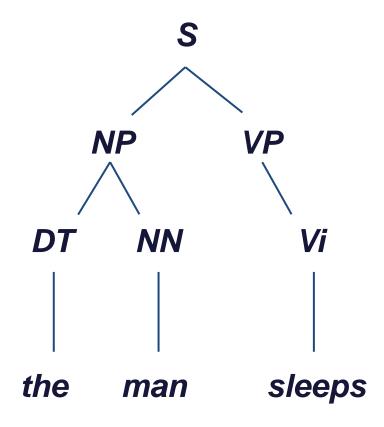
S	sentence
VP	verb phrase
NP	noun phrase
PP	prepositional phrase
DT	determiner
Vi	intransitive verb
Vt	transitive verb
NN	noun
IN	preposition
PR	pronoun

N

Parse Tree for this Simplified Grammar

S	→	NP	VP	1
VP	→	Vi		
VP	→	Vt	NP	
VP	→	VP	NP	
NP	→	PR		
NP	→	DT	NN	
NP	→	NP	PP	
PP	→	IN	NP	

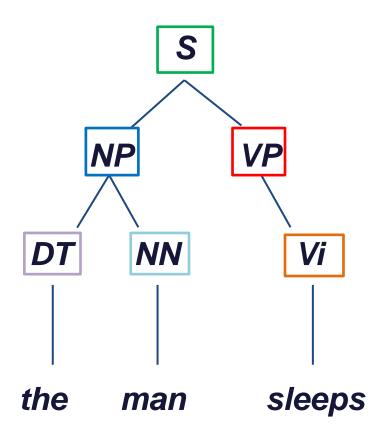
Vi	\rightarrow	sleeps
Vt	\rightarrow	saw
NN	\rightarrow	man
NN	→	woman
NN	→	telescope
DT	→	the
IN	→	with
IN	→	in
PR	→	he

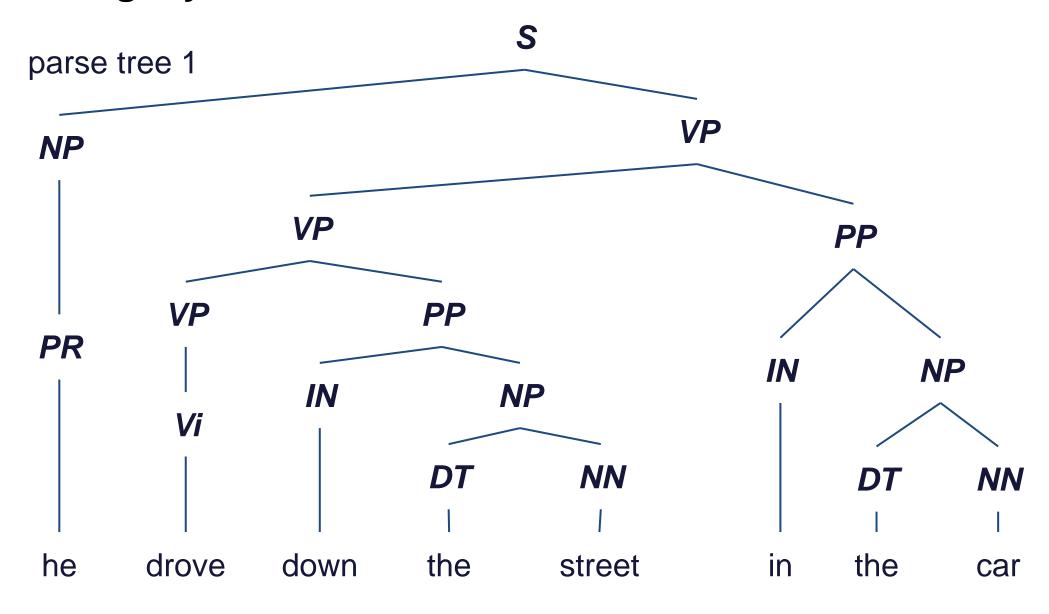


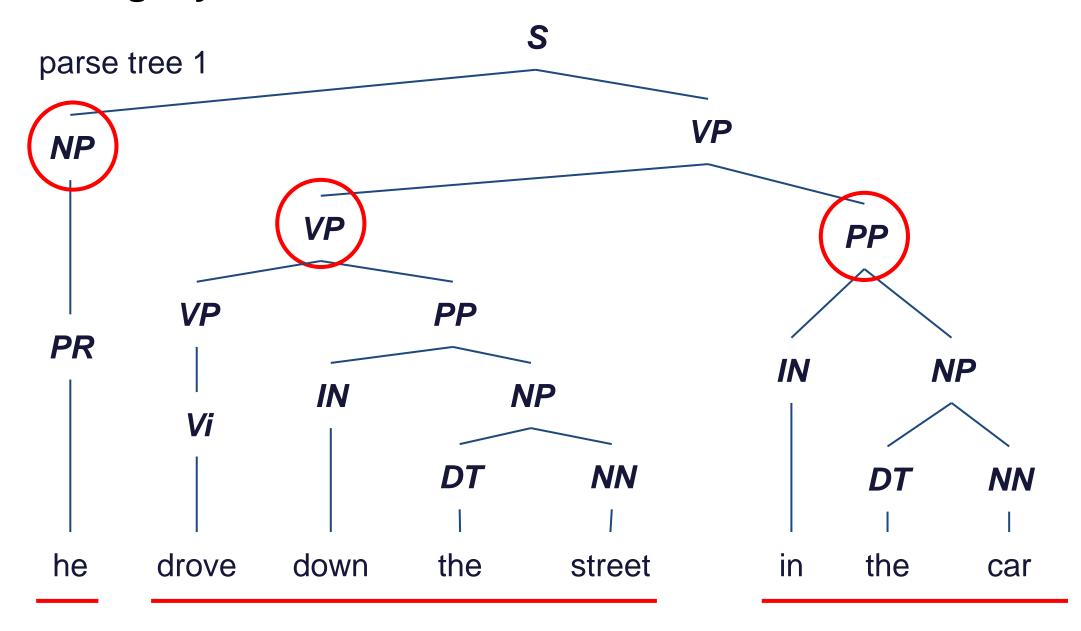
Parse Tree for this Simplified Grammar

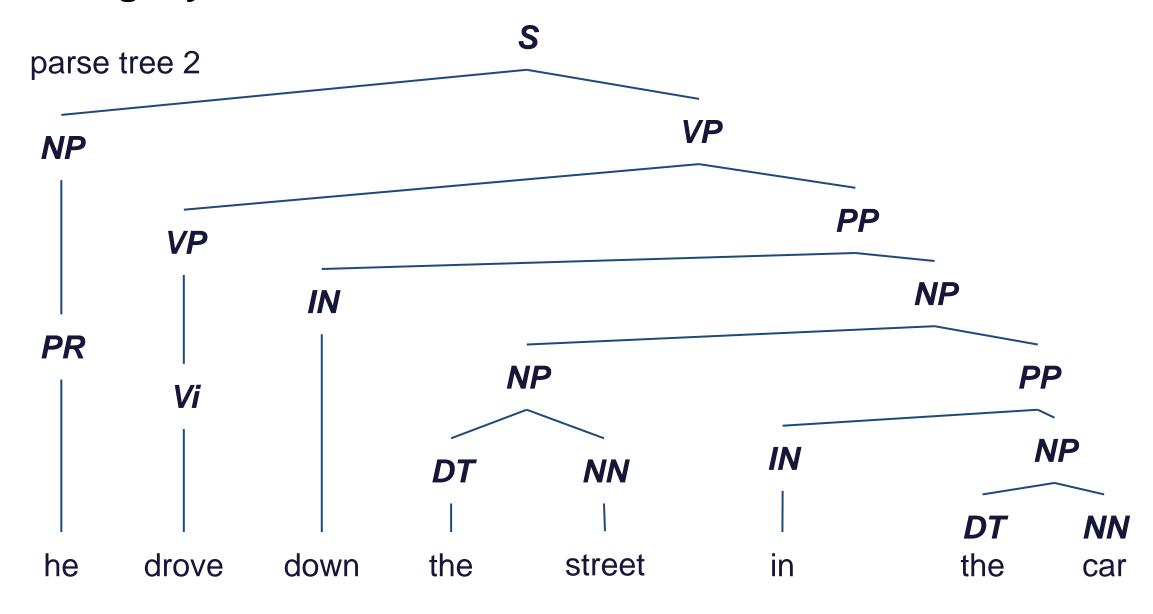
s	\rightarrow	NP	VP	F
VP	\rightarrow	Vi		
VP	→	Vt	NP	
VP	→	VP	NP	
NP	\rightarrow	PR		
NP NP	→ →	PR DT	NN	

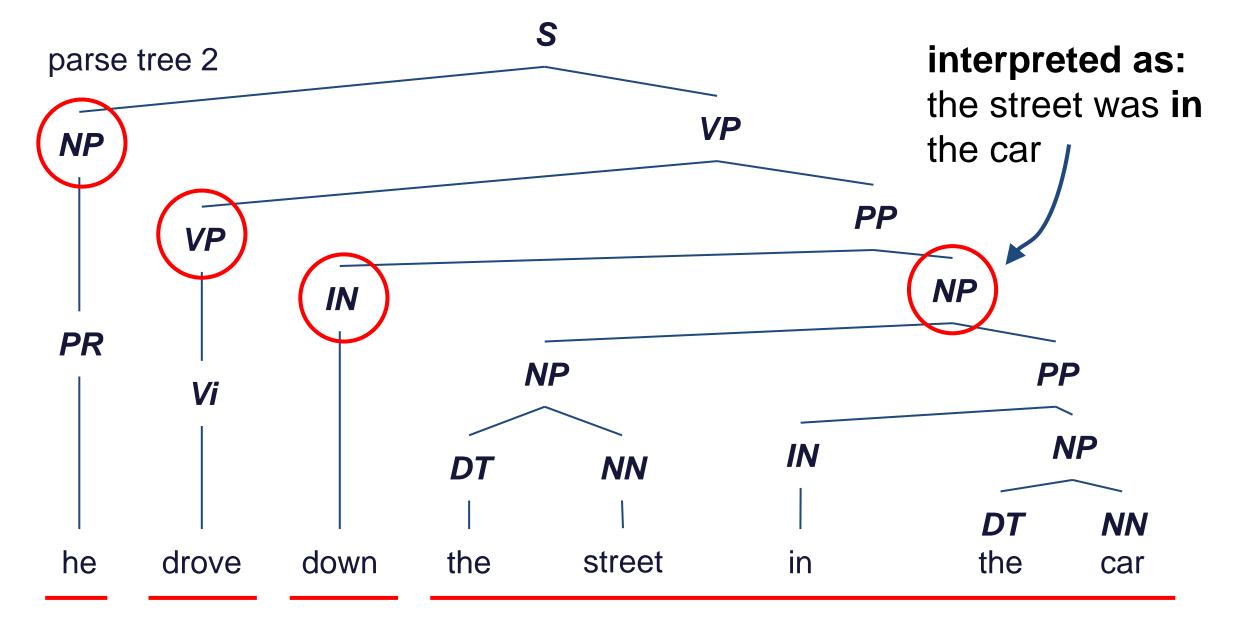
Vi	\rightarrow	sleeps	Σ
Vt	\rightarrow	saw	
NN	→	man	
NN	→	woman	
NN	→	telescope	
DT	→	the	
IN	→	with	
IN	→	in	
PR	\rightarrow	he	









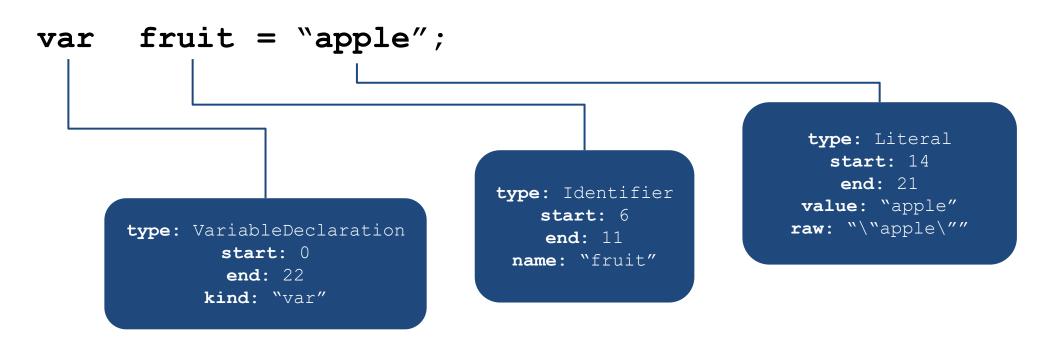


lexing & parsing using context-free grammars

lexing / tokenization

lexing is the process of breaking an input stream into discrete components (lexemes) and applying defining characteristic information to them.

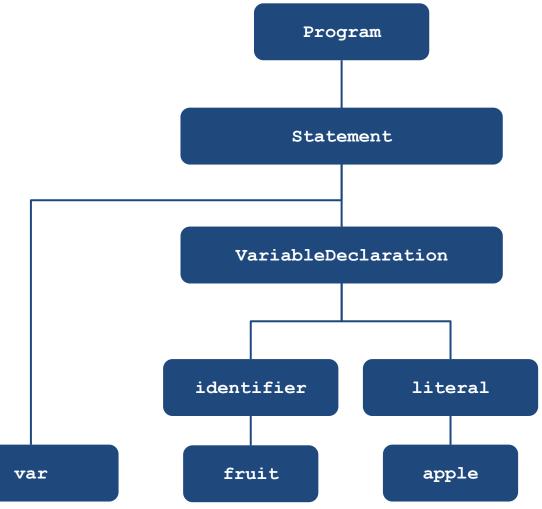
lexing is a fancy word for tokenization, lexeme is a fancy word for token



parsing

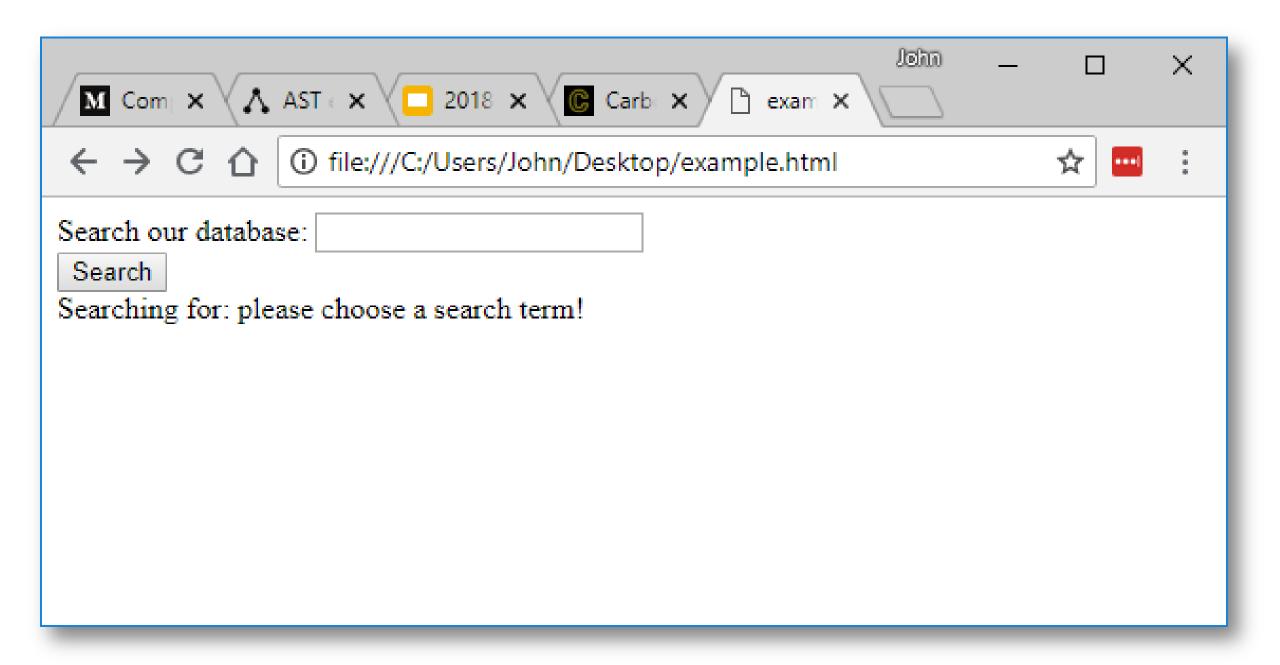
parsing is the process of applying structure to an input token stream.

```
var fruit = "apple";
```

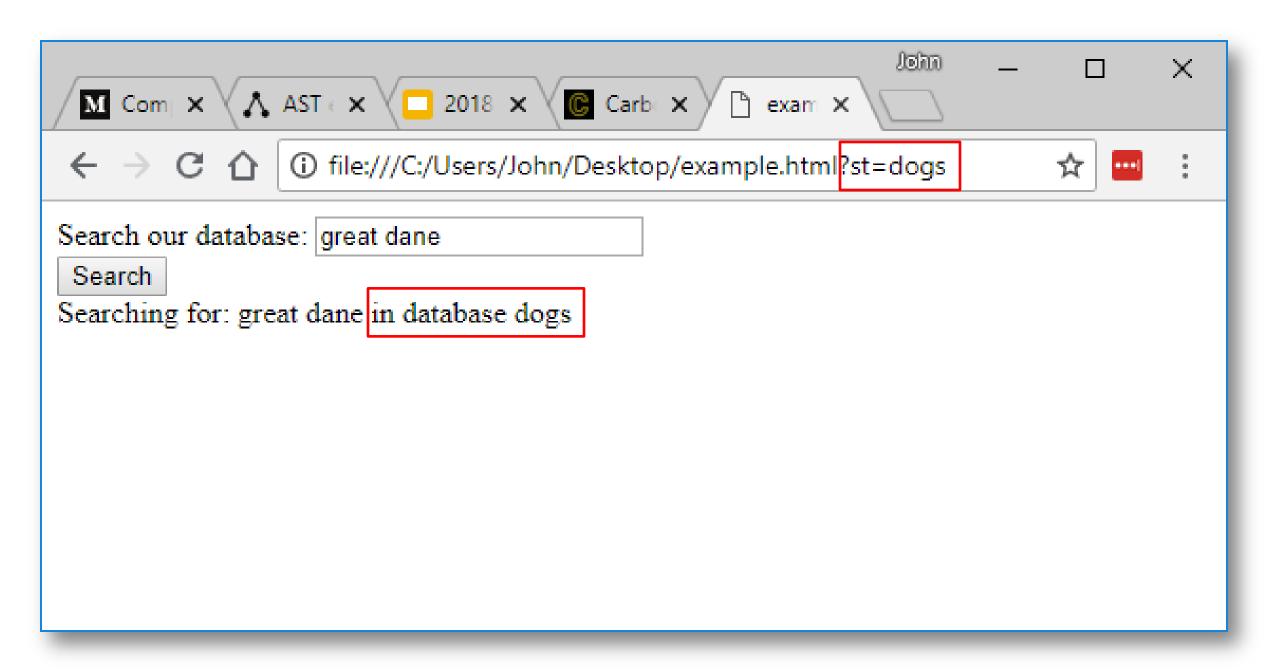


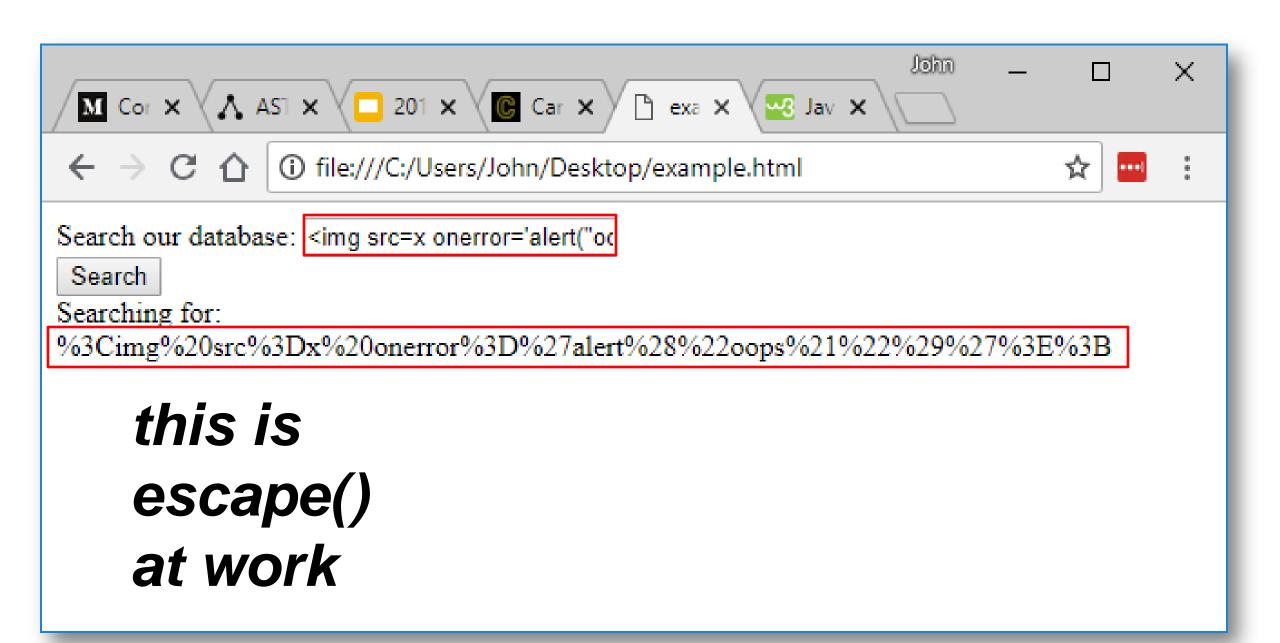
question: can you spot the vulnerability in this javascript code?

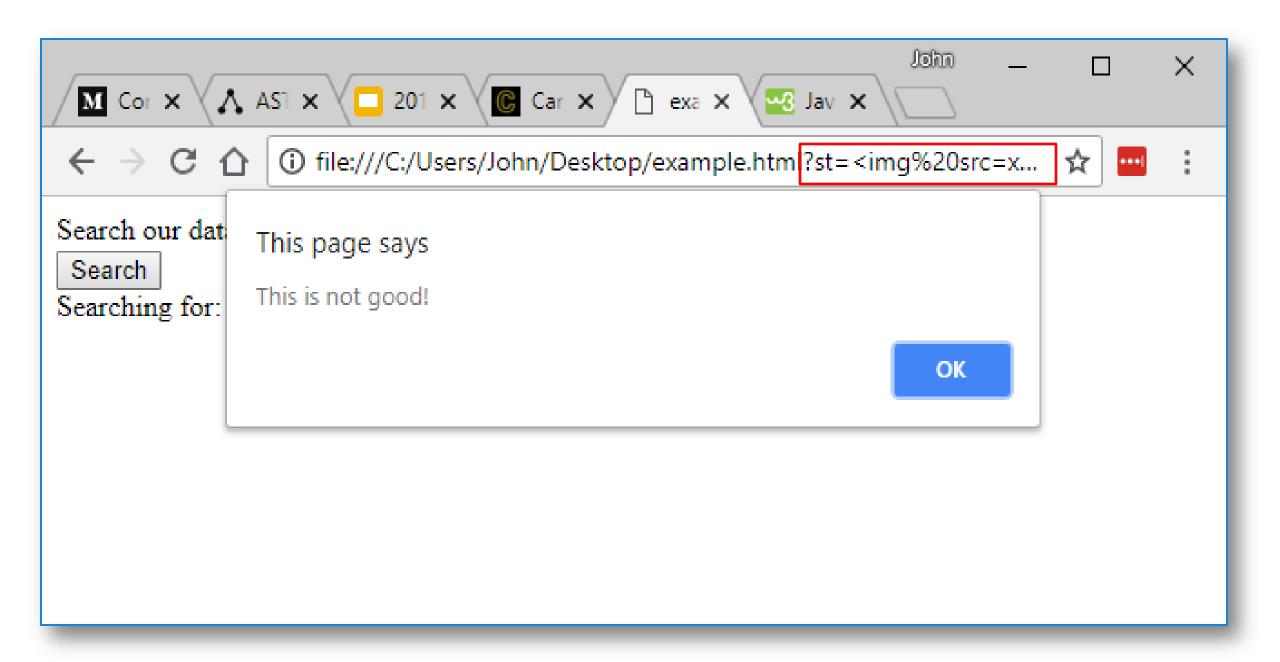
```
<!DOCTYPE html>
<html>
<body>
<form>
  Search our database: <input type="text" name="sterm" id="sterm"><br>
  <input type="button" onclick="search()" value="Search">
</form>
<div id="msq"></div>
<script>
    var urlParams = new URLSearchParams(window.location.search);
    function search() {
        document.getElementById("msg").innerHTML = "Searching for: " +
        (escape(document.getElementById("sterm").value) || "please choose a search term!") +
        (urlParams.get('st') ? " in database " + urlParams.get('st') : "");
</script>
</body>
</html>
```











```
source
<!DOCTYPE html>
                                                                               sink
<html>
<body>
                                                                             sanitize
<form>
 Search our database: <input type="text" name="sterm" id="sterm"><br>
 <input type="button" onclick="search()" value="Search">
</form>
<div id="msq"></div>
<script>
   var urlParams = new URLSearchParams(window.location.search);
    function search() {
       document.getElementById("msg").innerHTML = "Searching for: " +
        (escape(document.getElementById("sterm").value) || "please choose a search term!") +
        (urlParams.get('st') ? " in database " + urlParams.get('st') : "");
</script>
</body>
</html>
```

• • •

```
var urlParams = new URLSearchParams(window.location.search);

function search() {
    // 'st' is the 'search term'
    document.getElementById("msg").innerHTML = "Searching for: " +
        (escape(document.getElementById("sterm").value) || "please
        choose a search term!") + (urlParams.get('st') ? " in database "
        + urlParams.get('st') : "");
}
```

source sink sanitize

can you automate this process using grep or regex?

can you automate this process using grep or regex?

you can:

- find sources of user-controlled data
- find vulnerable sinks
- look for sanitization functions

you cannot:

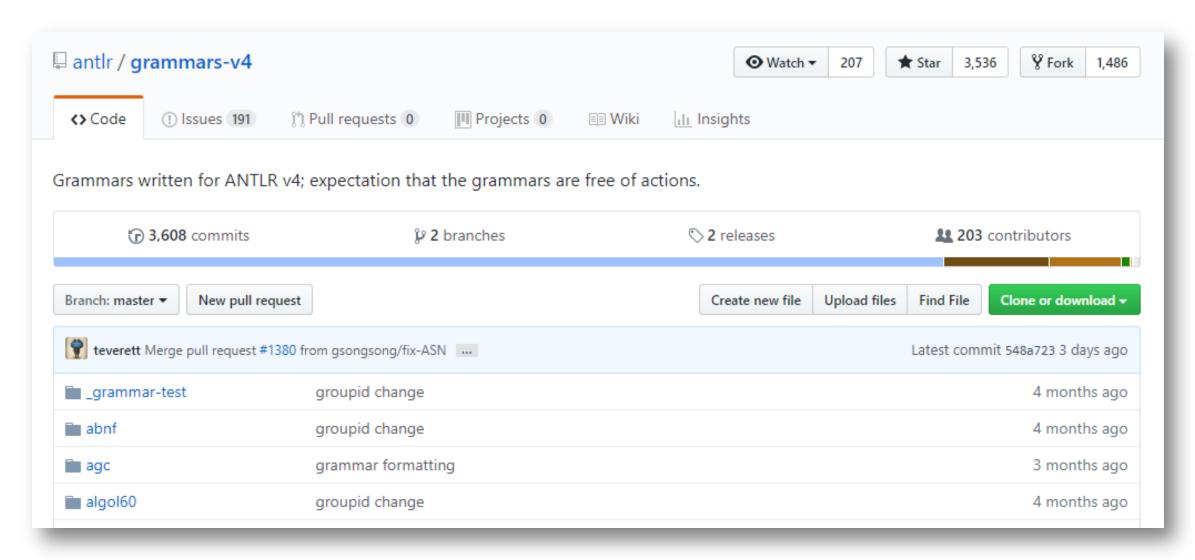
- determine the relationships between sources, sinks, and sanitizers
- resolve variable scope easily
- deconflict ambiguous variable names

automating this process using parsing strategies.

let's use ANTLR as our lexer / parser generator, to build parse trees and walk them looking for issues.

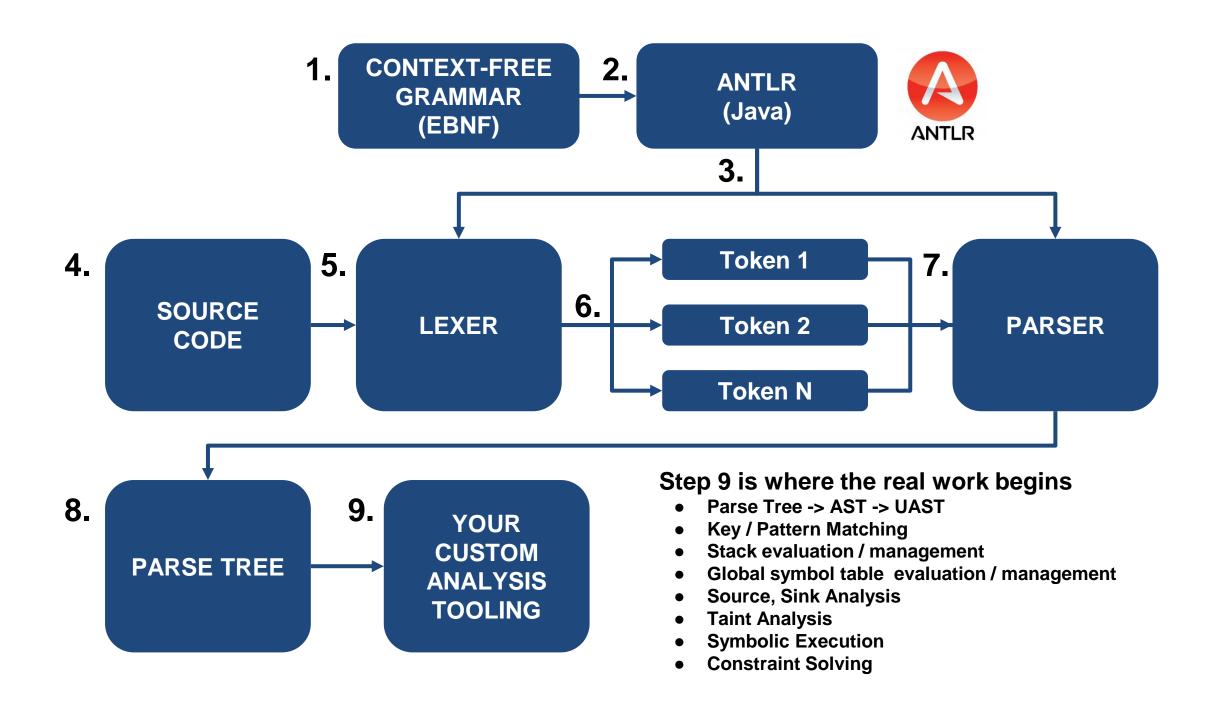


currently 190 grammars available



extended backus-naur form (EBNF)

```
arrayLiteral
    : '[' ','* elementList? ','* ']'
228
     elementList
229
          : singleExpression (','+ singleExpression)* (','+ lastElement)?
          | lastElement
232
233
     lastElement
                                  // ECMAScript 6: Spread Operator
         : Ellipsis Identifier
235
237
     objectLiteral
         : '{' (propertyAssignment (',' propertyAssignment)*)? '.'? '}'
239
240
```



demo: fun with improvisational parsers (fwip)





github.com/cetfor/fwip

using fwip: help

```
PS C:\Users\John\Desktop\fwip> node .\fwip.js -h
Usage: fwip [options]

Options:

-a, --analyze [file] Analyze a target file or directory of files
-s, --scrape [url] Scrape a target URL
-d, --debug Print debug strings (analyzers will not run in this mode)
-v, --version output the version number
-h, --help output usage information
PS C:\Users\John\Desktop\fwip>
```

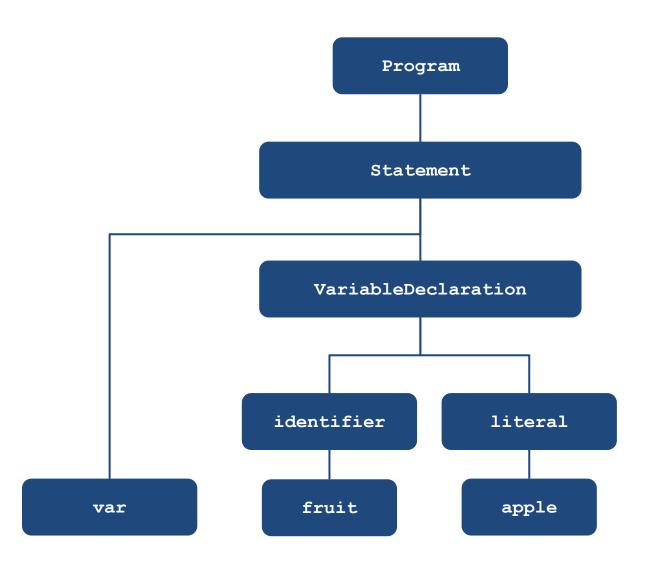
using fwip: analyzing a file

```
PS C:\Users\John\Desktop\fwip> node .\fwip.js -a .\examples\test.js
Analyzing file .\examples\test.js

Finished analyzing 1 file(s) with 0 error(s) and 0 skipped file(s).

PS C:\Users\John\Desktop\fwip>
```

remember this example parse tree?



```
enterProgram: varfruit="apple";<EOF>
 enterSourceElements: varfruit="apple";
    enterSourceElement: varfruit="apple";
      enterStatement: varfruit="apple";
        enterVariableStatement: varfruit="apple";
         enterVarModifier: var
          exitVarModifier: var
          enterVariableDeclarationList: fruit="apple"
            enterVariableDeclaration: fruit="apple"
            enterIdentifierName: fruit
            exitIdentifierName: fruit
            enterLiteralExpression: "apple"
              enterLiteral: "apple"
              exitLiteral: "apple"
            exitLiteralExpression: "apple"
            exitVariableDeclaration: fruit="apple"
          exitVariableDeclarationList: fruit="apple"
          enterEos: ;
          exitEos: ;
        exitVariableStatement: varfruit="apple";
      exitStatement: varfruit="apple";
    exitSourceElement: varfruit="apple";
 exitSourceElements: varfruit="apple";
exitProgram: varfruit="apple";<EOF>
```

10

11

12

13

14

15

16

17

18

19

20

21

22

23

24

this is the result of

node .\fwip.js -a test.js --debug

where test.js is just

var fruit = "apple";

let's revisit our example vulnerable code and build an analyzer.

```
<!DOCTYPE html>
<html>
<body>
 Search our database: <input type="text" name="sterm" id="sterm"><br>
 <input type="button" onclick="search()" value="Search">
</form>
<div id="msq"></div>
<script>
   var urlParams = new URLSearchParams(window.location.search);
   function search() {
       document.getElementById("msg").innerHTML = "Searching for: " +
       (escape(document.getElementById("sterm").value) || "please choose a search term!") +
       (urlParams.get('st') ? " in database " + urlParams.get('st') : "");
</script>
</body>
</html>
```

what we want to check for:

- 1. is a dangerous sink used?
- 2. is data from a usercontrolled <u>source</u> passed to the sink?
- 3. iff, is a <u>sanitize</u> function used on the source data?

```
<!DOCTYPE html>
<html>
<body>
 Search our database: <input type="text" name="sterm" id="sterm"><br>
 <input type="button" onclick="search()" value="Search">
</form>
<div id="msq"></div>
<script>
   var urlParams = new URLSearchParams(window.location.search);
   function search() {
       document.getElementById("msg").innerHTML = "Searching for: " +
       (escape(document.getElementById("sterm").value) || "please choose a search term!") +
       (urlParams.get('st') ? " in database " + urlParams.get('st') : "");
</script>
</body>
</html>
```

we'll use the fwip --debug switch to see what elements we should check.

```
// variable source
     enterVariableStatement: varurlParams=newURLSearchParams(window.location.search);
     enterVarModifier: var
     enterVariableName: urlParams
4
     enterArgumentsExpression: URLSearchParams(window.location.search)
     enterArguments: (window.location.search)
6
     // sink
8
     enterMemberDotExpression: document.getElementById("msg").innerHTML
9
     exitMemberDotExpression: document.getElementById
10
     enterArguments: ("msg")
11
12
     enterIdentifierName: innerHTML
13
14
     // sanitized source
15
     enterArgumentsExpression: escape(document.getElementById("sterm").value)
     enterIdentifierExpression: escape
16
                                                                                 select
17
     enterArguments: (document.getElementById("sterm").value)
18
19
     // non-sanitized, variable source
                                                                      results of
     enterArgumentsExpression: urlParams.get('st')
20
     enterMemberDotExpression: urlParams.get
21
     enterIdentifierExpression: urlParams
22
                                                         node .\fwip.js -a owasp.html --debug
23
     enterIdentifierName: get
     enterArguments: ('st')
24
```

there are many ways we could analyze this code. we'll go really simple by using a stack (an array)

```
const stack = [
   // interesting stuff goes here
]
```

whenever we enter interesting conditions, we'll push select leaf nodes to the stack.

```
const stack = [
   // interesting stuff goes here
enterMemberDotExpression:
        document.getElementById("msg").innerHTML
enterIdentifierName:
        innerHTML
exitMemberDotExpression:
        document.getElementById("msg").innerHTML
stack.push('innerHTML')
stack = [
   'innerHTML',
```

when parsing completes we'll call our analyzers on the program "state" we've built using stacks.

now onto demos and code.

```
PS C:\Users\John\Documents\GitHub\fwip> node fwip.js -a .\examples\owasp.html
Analyzing file .\examples\owasp.html
>> Potential vulnerability on line: 5
    Line 5: document.getElementById("msg").innerHTML="Searching for: "+(escap...
    Description: This line contains 3 source(s), 2 sink(s), and 1 sanitizer(s).
    Sources: value,URLSearchParams,URLSearchParams
    Sinks: innerHTML,value
    Sanitizers: escape

Finished analyzing 1 file(s) with 0 error(s) and 0 skipped file(s).
```