

Chapter 3

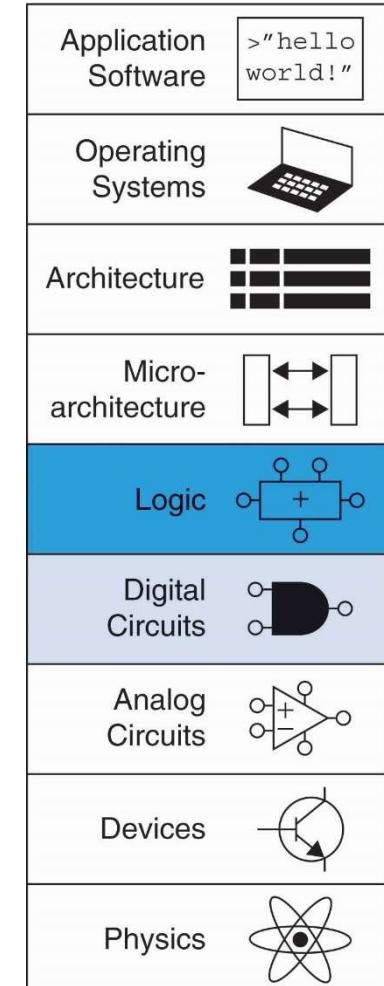
Digital Design and Computer Architecture, 2nd Edition

David Money Harris and Sarah L. Harris



Chapter 3 :: Topics

- **Introduction**
- **Latches and Flip-Flops**
- **Synchronous Logic Design**
- **Finite State Machines**
- **Timing of Sequential Logic**
- **Parallelism**



Introduction

- Outputs of sequential logic depend on current *and* prior input values – it has ***memory***.
- Some definitions:
 - **State:** all the information about a circuit necessary to explain its future behavior
 - **Latches and flip-flops:** state elements that store one bit of state
 - **Synchronous sequential circuits:** combinational logic followed by a bank of flip-flops

Sequential Circuits

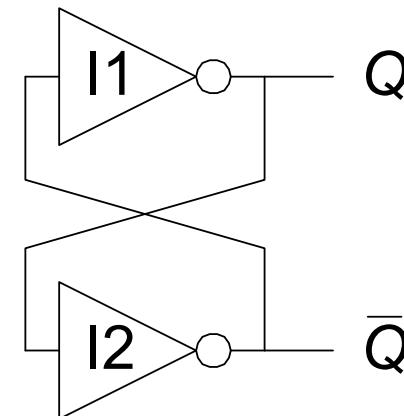
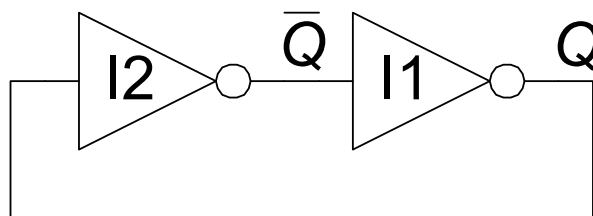
- Give sequence to events
- Have memory (short-term)
- Use feedback from output to input to store information

State Elements

- The state of a circuit influences its future behavior
- State elements store state
 - Bistable circuit
 - SR Latch
 - D Latch
 - D Flip-flop

Bistable Circuit

- Fundamental building block of other state elements
- Two outputs: Q, \bar{Q}
- No inputs

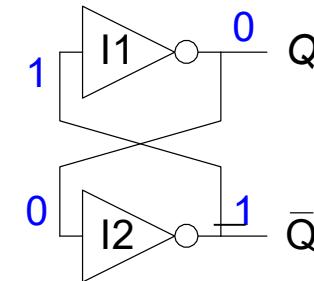


Bistable Circuit Analysis

- Consider the two possible cases:

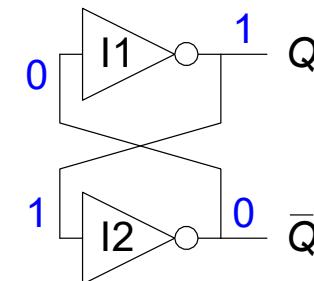
- $- Q = 0:$

then $\bar{Q} = 1, Q = 0$ (consistent)



- $- Q = 1:$

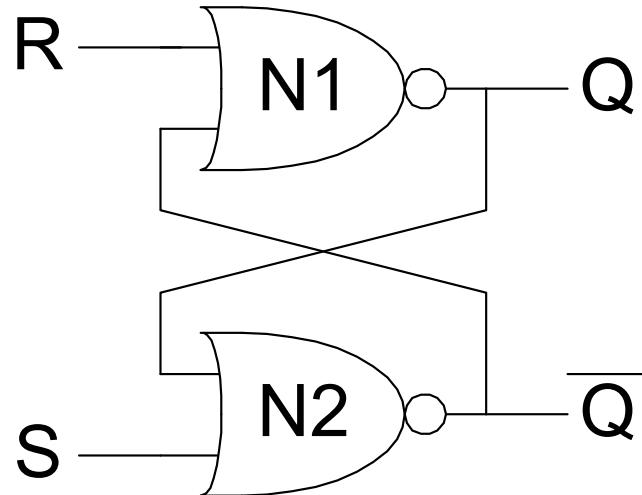
then $\bar{Q} = 0, Q = 1$ (consistent)



- Stores 1 bit of state in the state variable, Q (or \bar{Q})
- But there are **no inputs to control the state**

SR (Set/Reset) Latch

- SR Latch

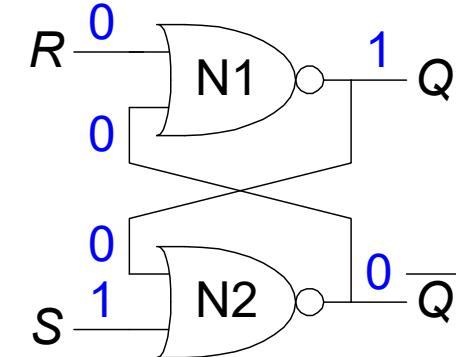


- Consider the four possible cases:
 - $S = 1, R = 0$
 - $S = 0, R = 1$
 - $S = 0, R = 0$
 - $S = 1, R = 1$

SR Latch Analysis

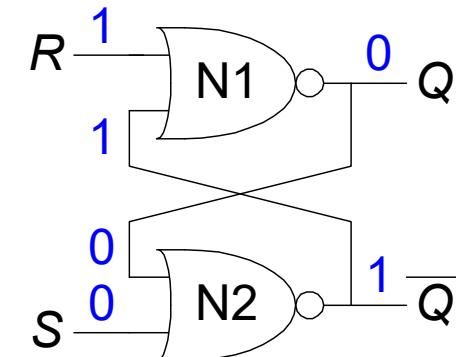
– $S = 1, R = 0$:

then $Q = 1$ and $\bar{Q} = 0$



– $S = 0, R = 1$:

then $Q = \bar{1}$ and $\bar{Q} = 0$

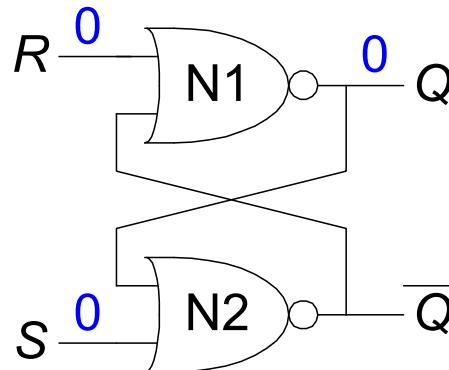


SR Latch Analysis

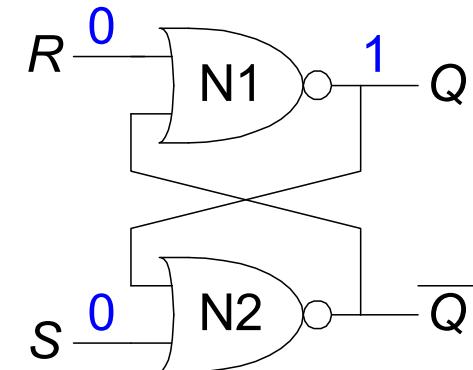
$- S = 0, R = 0:$

then $Q = Q_{prev}$

$$Q_{prev} = 0$$

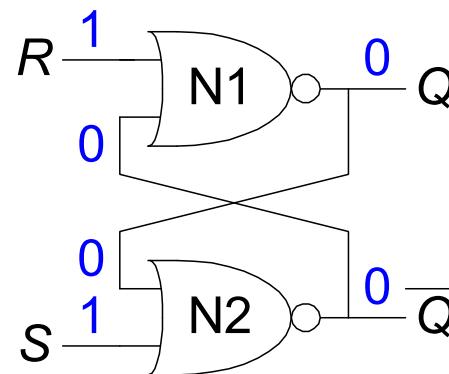


$$Q_{prev} = 1$$



$- S = 1, R = 1:$

then $Q = 0, \bar{Q} = 0$



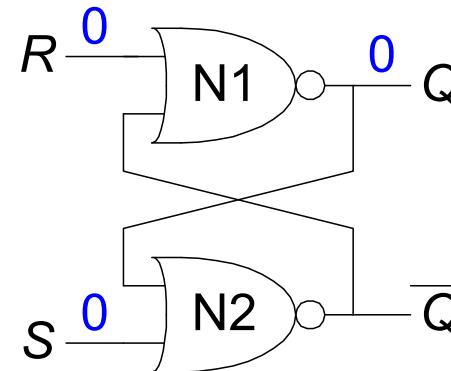
SR Latch Analysis

– $S = 0, R = 0$:

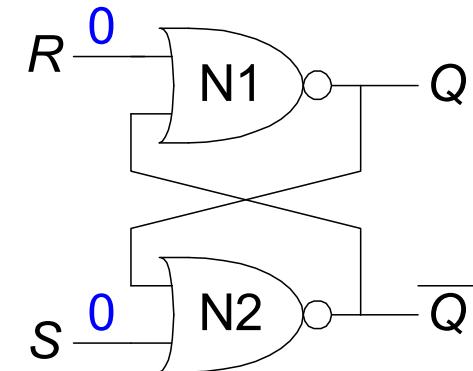
then $Q = Q_{prev}$

– **Memory!**

$Q_{prev} = 0$



$Q_{prev} = 1$

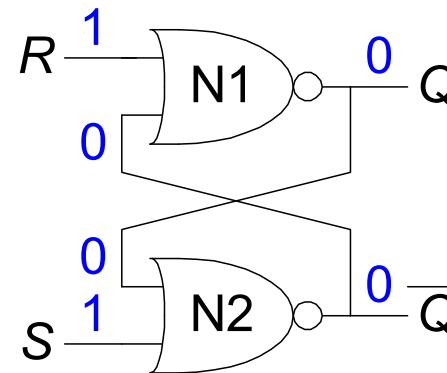


– $S = 1, R = 1$:

then $Q = 0, \bar{Q} = 0$

– **Invalid State**

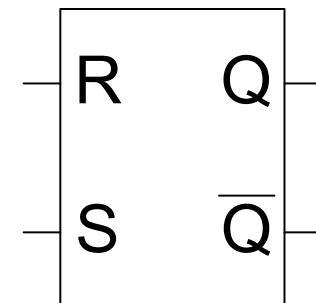
$\bar{Q} \neq \text{NOT } Q$



SR Latch Symbol

- SR stands for Set/Reset Latch
 - Stores one bit of state (Q)
- Control what value is being stored with S , R inputs
 - **Set:** Make the output 1
($S = 1, R = 0, Q = 1$)
 - **Reset:** Make the output 0
($S = 0, R = 1, Q = 0$)
- **Must do something to avoid invalid state (when $S = R = 1$)**

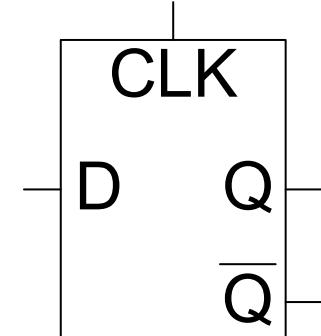
SR Latch
Symbol



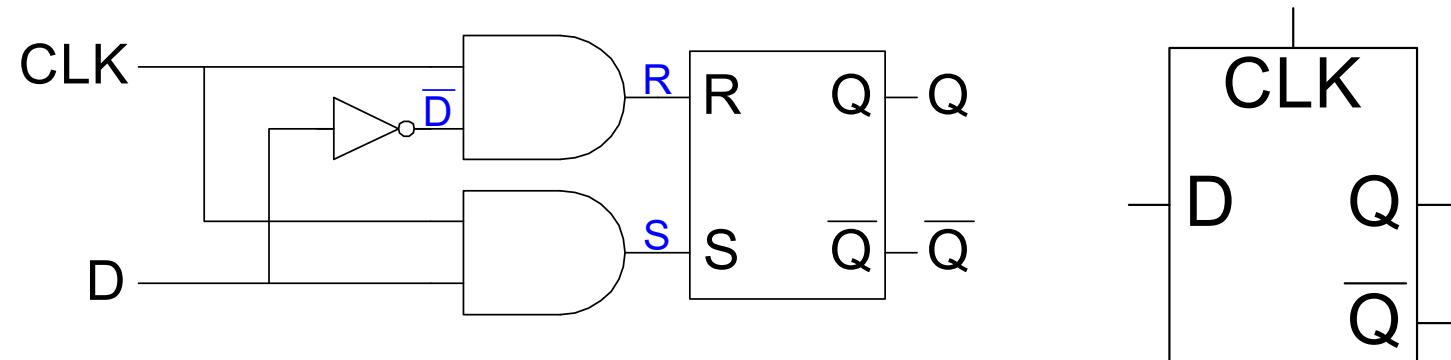
D Latch

- Two inputs: CLK , D
 - CLK : controls *when* the output changes
 - D (the data input): controls *what* the output changes to
- Function
 - When $CLK = 1$,
 D passes through to Q (*transparent*)
 - When $CLK = 0$,
 Q holds its previous value (*opaque*)
- Avoids invalid case when
 $Q \neq \text{NOT } \bar{Q}$

D Latch
Symbol

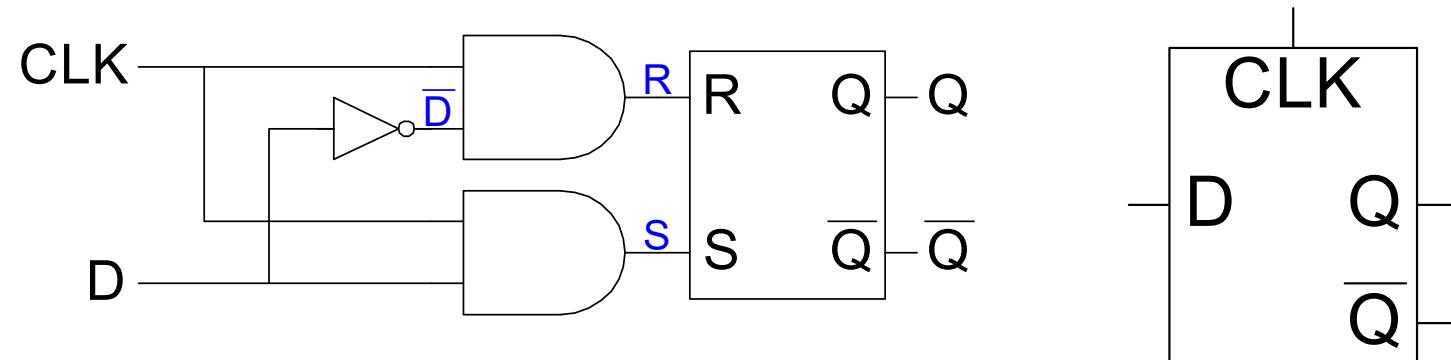


D Latch Internal Circuit



CLK	D	\bar{D}	S	R	Q	\bar{Q}
0	X					
1	0					
1	1					

D Latch Internal Circuit

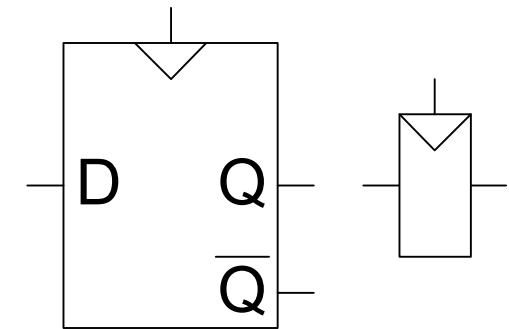


CLK	D	\bar{D}	S	R	Q	\bar{Q}
0	X	\bar{X}	0	0	Q_{prev}	\bar{Q}_{prev}
1	0	1	0	1	0	1
1	1	0	1	0	1	0

D Flip-Flop

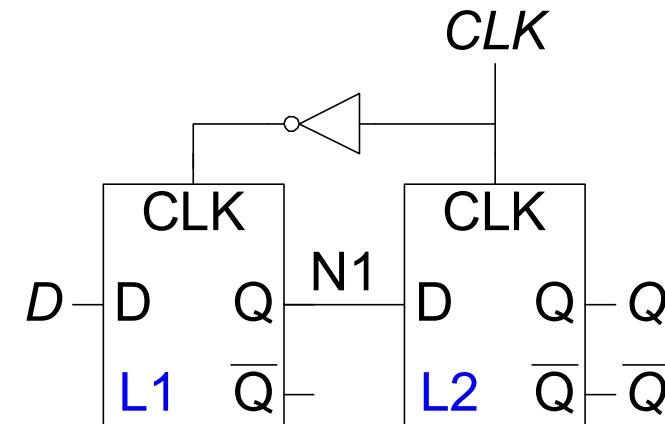
- **Inputs:** CLK, D
 - **Function**
 - Samples D on rising edge of CLK
 - When CLK rises from 0 to 1, D passes through to Q
 - Otherwise, Q holds its previous value
 - Q changes only on rising edge of CLK
 - Called *edge-triggered*
 - Activated on the clock edge

D Flip-Flop Symbols

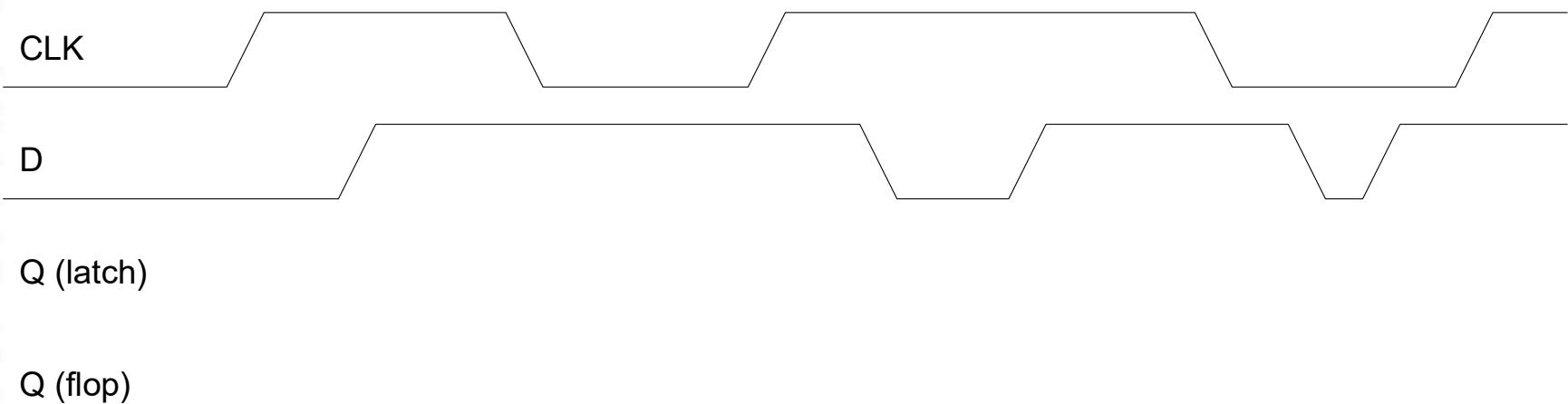
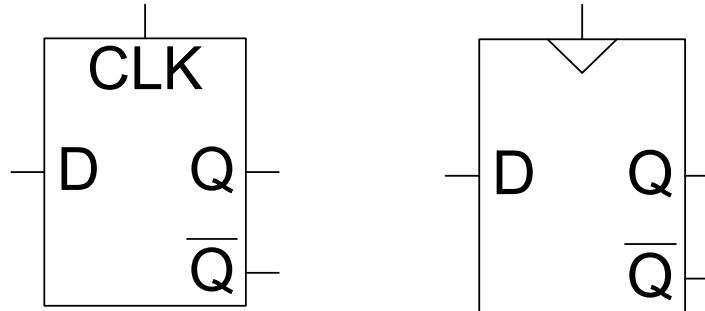


D Flip-Flop Internal Circuit

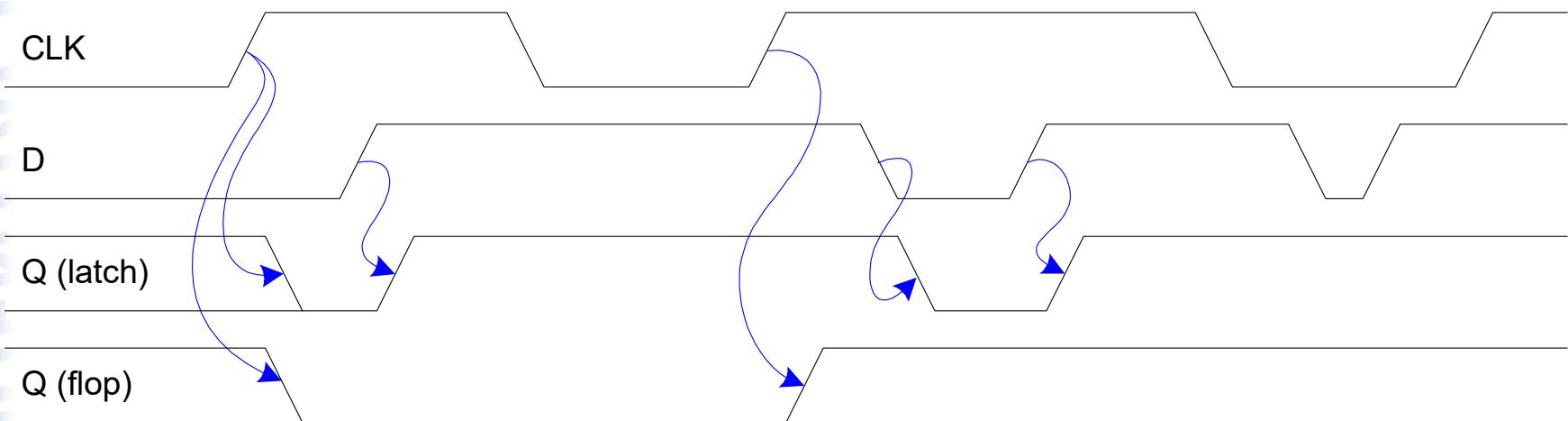
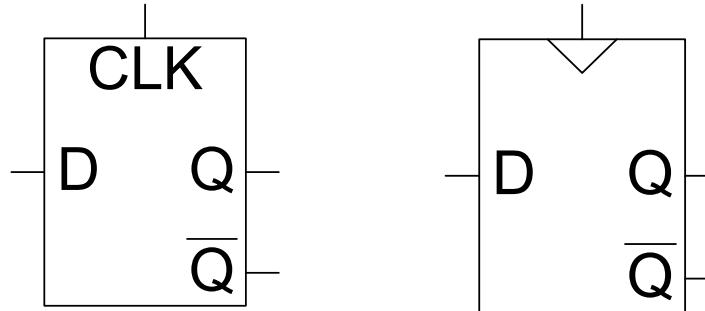
- Two back-to-back latches (L1 and L2) controlled by complementary clocks
- When $\text{CLK} = 0$
 - L1 is transparent
 - L2 is opaque
 - D passes through to N1
- When $\text{CLK} = 1$
 - L2 is transparent
 - L1 is opaque
 - N1 passes through to Q
- Thus, on the edge of the clock (when CLK rises from 0 → 1)
 - D passes through to Q



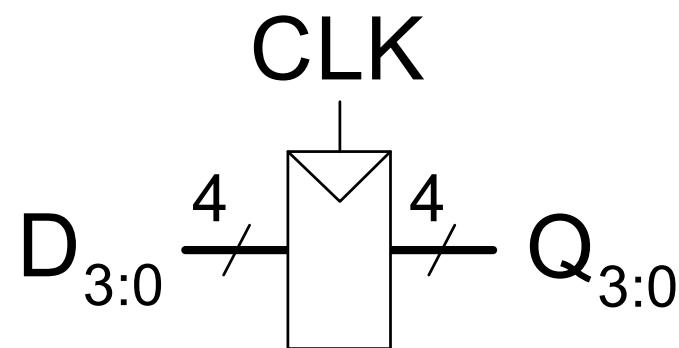
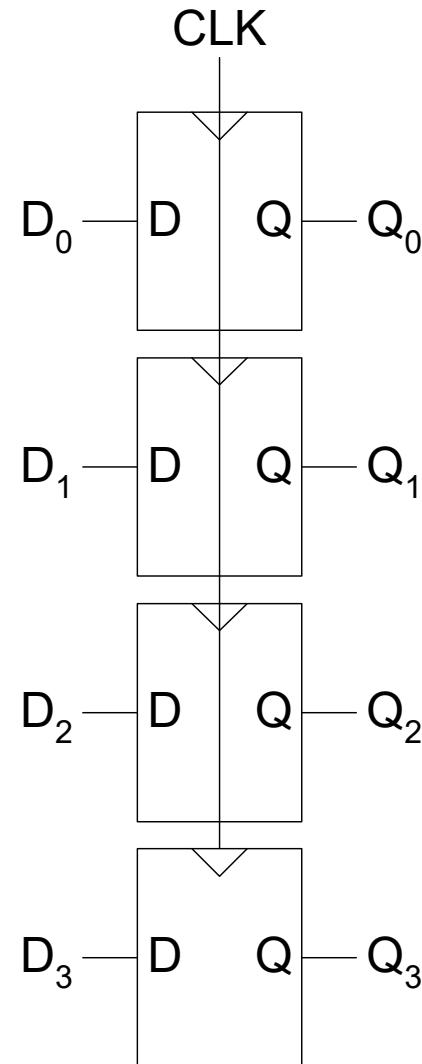
D Latch vs. D Flip-Flop



D Latch vs. D Flip-Flop

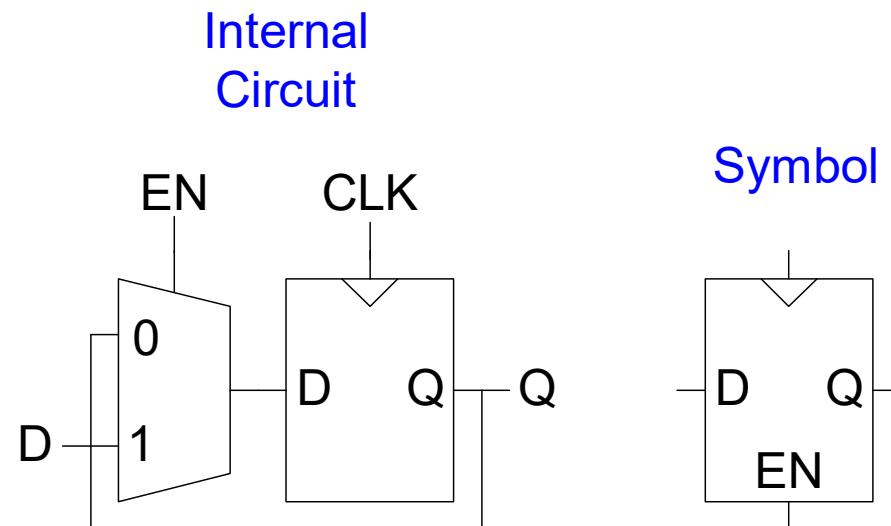


Registers



Enabled Flip-Flops

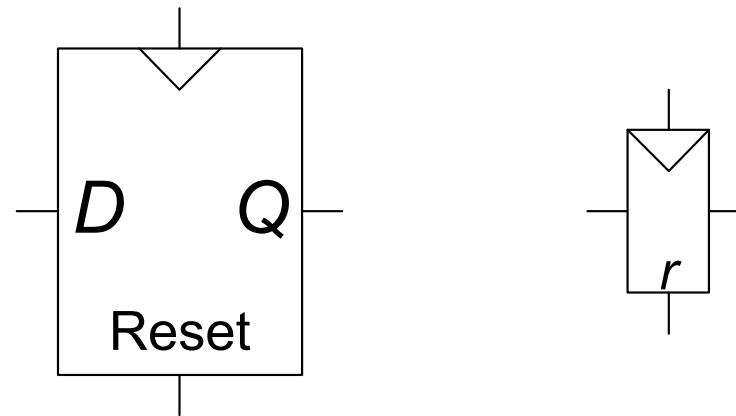
- **Inputs:** CLK , D , EN
 - The enable input (EN) controls when new data (D) is stored
- **Function**
 - $EN = 1$: D passes through to Q on the clock edge
 - $EN = 0$: the flip-flop retains its previous state



Resettable Flip-Flops

- **Inputs:** $CLK, D, Reset$
- **Function:**
 - $Reset = 1$: Q is forced to 0
 - $Reset = 0$: flip-flop behaves as ordinary D flip-flop

Symbols

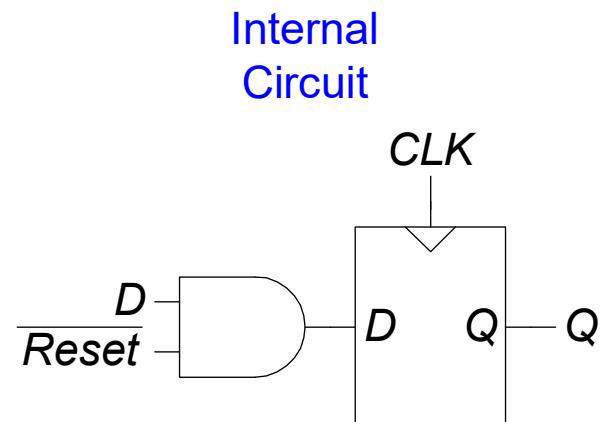


Resettable Flip-Flops

- Two types:
 - **Synchronous**: resets at the clock edge only
 - **Asynchronous**: resets immediately when $Reset = 1$
- Asynchronously resettable flip-flop requires changing the internal circuitry of the flip-flop
- Synchronously resettable flip-flop?

Resettable Flip-Flops

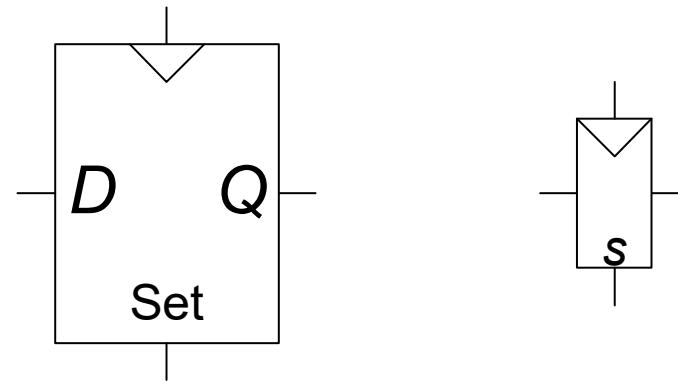
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Settable Flip-Flops

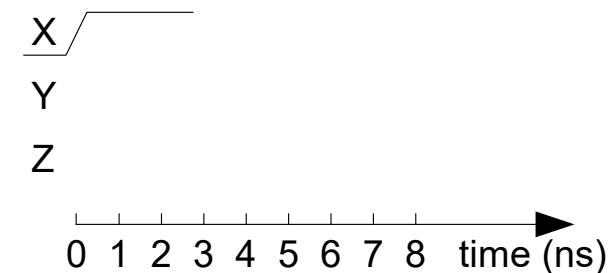
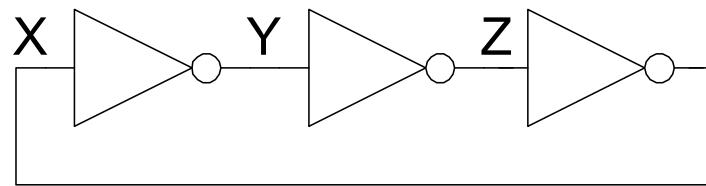
- Inputs: CLK , D , Set
- Function:
 - $Set = 1$: Q is set to 1
 - $Set = 0$: the flip-flop behaves as ordinary D flip-flop

Symbols



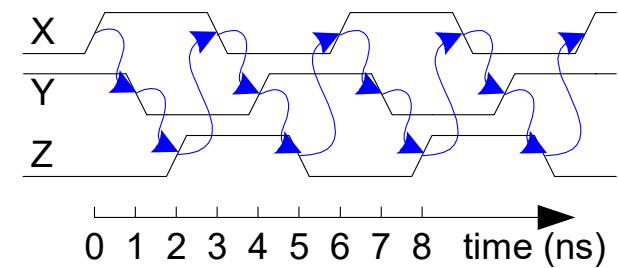
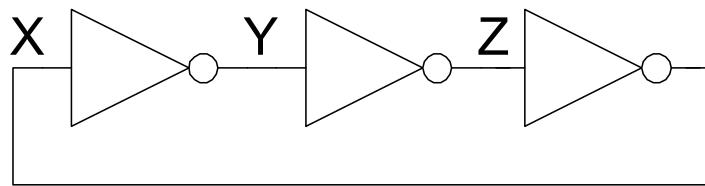
Sequential Logic

- Sequential circuits: all circuits that aren't combinational
- A problematic circuit:



Sequential Logic

- Sequential circuits: all circuits that aren't combinational
- A problematic circuit:



- No inputs and 1-3 outputs
- Astable circuit, oscillates
- Period depends on inverter delay
- It has a *cyclic path*: output fed back to input

Synchronous Sequential Logic Design

- Breaks cyclic paths by **inserting registers**
- Registers contain **state** of the system
- State changes at clock edge: system **synchronized** to the clock
- **Rules** of synchronous sequential circuit composition:
 - Every circuit element is either a register or a combinational circuit
 - At least one circuit element is a register
 - All registers receive the same clock signal
 - Every cyclic path contains at least one register
- Two common synchronous sequential circuits
 - Finite State Machines (FSMs)
 - Pipelines