COMBINATIONAL LOGIC DESIGN

Chapter 2

***Digital Design and Computer Architecture*, 2nd Edition**

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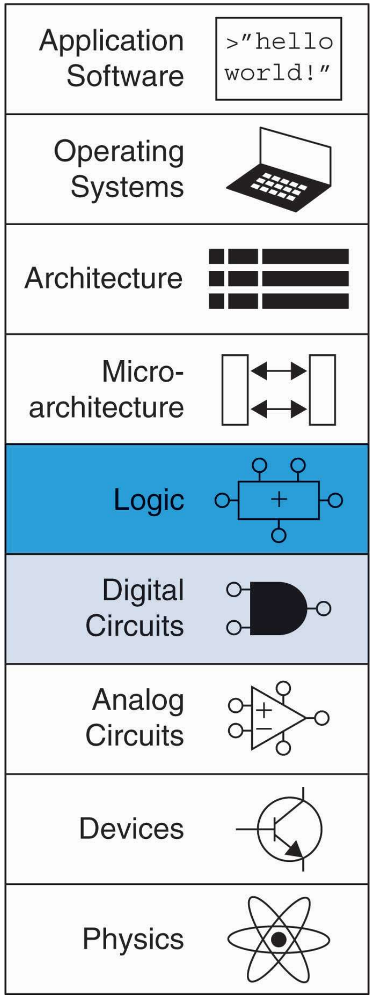
COMBINATIONAL LOGIC DESIGN

Chapter 2 :: Topics

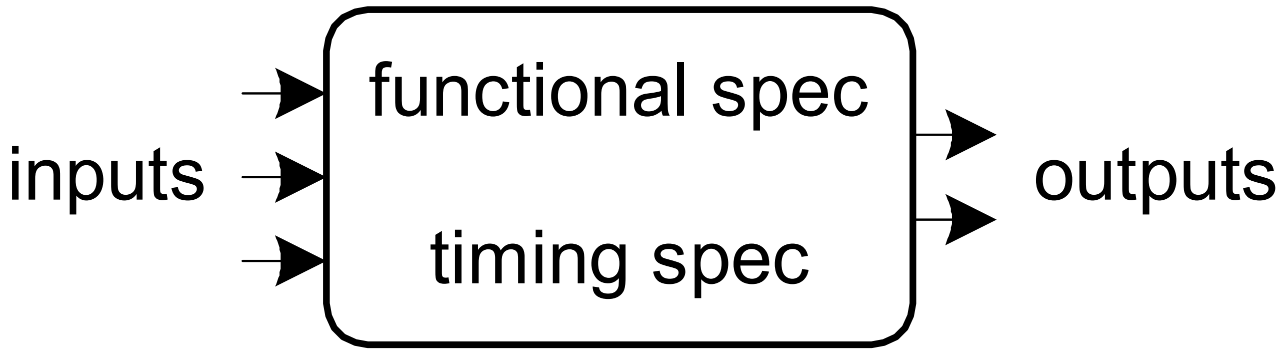
- **Introduction**
- **Boolean Equations**
- **Boolean Algebra**
- **From Logic to Gates**
- **Multilevel Combinational Logic**
- **X’s and Z’s, Oh My**
- **Karnaugh Maps**
- **Combinational Building Blocks**
- **Timing**



Introduction

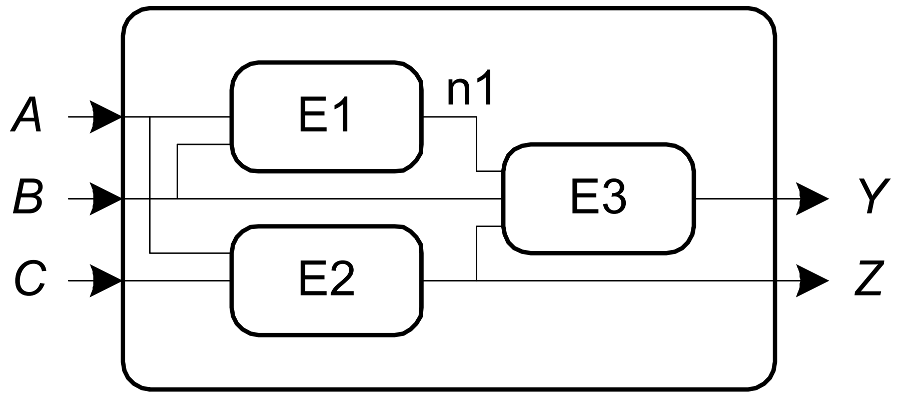
A logic circuit is composed of:

- Inputs
- Outputs
- Functional specification
- Timing specification



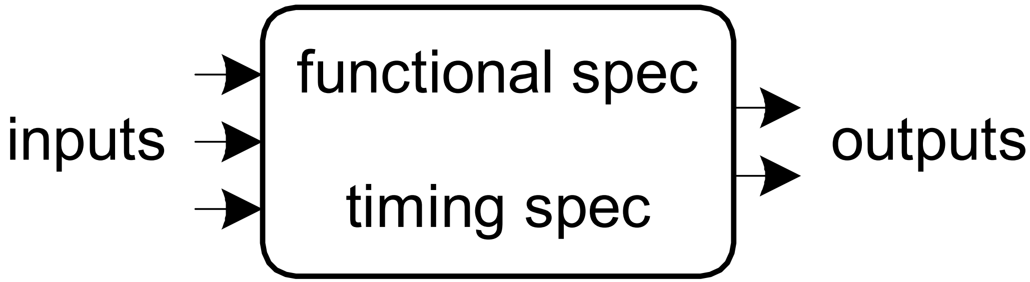
Circuits

- Nodes
  - Inputs: *A*, *B*, *C*
  - Outputs: *Y*, *Z*
  - Internal: n1
- Circuit elements
  - E1, E2, E3
  - Each a circuit



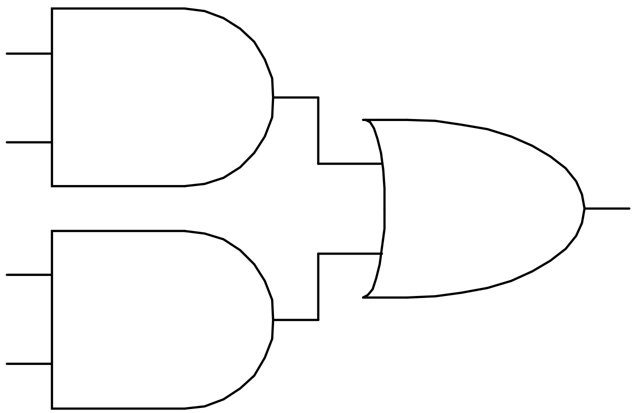
Types of Logic Circuits

- **Combinational Logic**
  - Memoryless
  - Outputs determined by current values of inputs
- **Sequential Logic**
  - Has memory
  - Outputs determined by previous and current values of inputs



Rules of Combinational Composition

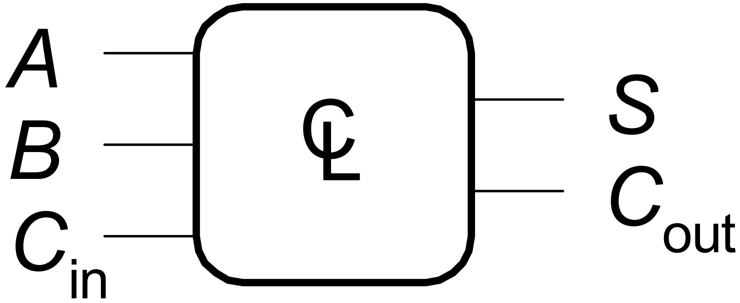
- Every element is combinational
- Every node is either an input or connects to *exactly one* output
- The circuit contains no cyclic paths
- **Example:**



Boolean Equations

- Functional specification of outputs in terms of inputs
- **Example:** $S = F(A, B, C_{in})$

  $C_{out} = F(A, B, C_{in})$



$$S = A \oplus B \oplus C_{in}$$
$$C_{out} = AB + AC_{in} + BC_{in}$$

Some Definitions

- Complement: variable with a bar over it
  $\bar{A}, \bar{B}, \bar{C}$
- Literal: variable or its complement
  $A, \bar{A}, B, \bar{B}, C, \bar{C}$
- Implicant: product of literals
  $AB\bar{C}, \bar{A}C, BC$
- Minterm: product that includes all input variables
  $AB\bar{C}, A\bar{B}\bar{C}, ABC$
- Maxterm: sum that includes all input variables
  $(A+\bar{B}+C), (\bar{A}+B+\bar{C}), (\bar{A}+\bar{B}+C)$

Sum-of-Products (SOP) Form

- All equations can be written in SOP form
- Each row has a **minterm**
- A minterm is a product (AND) of literals
- Each minterm is TRUE for that row (and only that row)
- Form function by ORing minterms where the output is TRUE
- Thus, a sum (OR) of products (AND terms)

A	B	Y	minterm	minterm name
0	0	0	$\overline{A}\ \overline{B}$	m_0
0	1	1	$\overline{A}\ B$	m_1
1	0	0	$A\ \overline{B}$	m_2
1	1	1	$A\ B$	m_3

$$Y = \mathrm{F}(A, B) =$$

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1	0	0	$A\ \overline{B}$	m_2
1	1	1	$A\ B$	m_3

$$Y = F(A, B) = \overline{A}B + AB = \Sigma(1, 3)$$

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Product-of-Sums (POS) Form

- All Boolean equations can be written in POS form
- Each row has a **maxterm**
- A maxterm is a sum (OR) of literals
- Each maxterm is FALSE for that row (and only that row)
- Form function by ANDing the maxterms for which the output is FALSE
- Thus, a product (AND) of sums (OR terms)

A	B	Y	maxterm	maxterm name
0	0	0	$A + B$	M_0
0	1	1	$A + \overline{B}$	M_1
1	0	0	$\overline{A} + B$	M_2
1	1	1	$\overline{A} + \overline{B}$	M_3

$$Y = \mathbf{F}(A, B) = (A + B)(\overline{A} + B) = \Pi(0, 2)$$

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Boolean Equations Example

- You are going to the cafeteria for lunch
  - You won't eat lunch (E)
  - If it's not open (O) or
  - If they only serve corndogs (C)
- Write a truth table for determining if you will eat lunch (E).

O	C	E
0	0	
0	1	
1	0	
1	1	

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O	C	E
0	0	0
0	1	0
1	0	1
1	1	0

SOP & POS Form

- SOP – sum-of-products

O	C	E	minterm
0	0		$\overline{O}\ \overline{C}$
0	1		$\overline{O}\ C$
1	0		$O\ \overline{C}$
1	1		$O\ C$

- POS – product-of-sums

O	C	Y	maxterm
0	0		$O + C$
0	1		$O + \overline{C}$
1	0		$\overline{O} + C$
1	1		$\overline{O} + \overline{C}$

SOP & POS Form

- SOP – sum-of-products

O	C	E	minterm
0	0	0	$\overline{O}\,\overline{C}$
0	1	0	$\overline{O}\,C$
1	0	1	$O\,\overline{C}$
1	1	0	$O\,C$

$$Y = O\overline{C}$$
$$= \Sigma(2)$$

- POS – product-of-sums

O	C	E	maxterm
0	0	0	$O + C$
0	1	0	$O + \overline{C}$
1	0	1	$\overline{O} + C$
1	1	0	$\overline{O} + \overline{C}$

$$Y = (O + C)(O + \overline{C})(\overline{O} + \overline{C})$$
$$= \Pi(0, 1, 3)$$

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Boolean Algebra

- Axioms and theorems to **simplify** Boolean equations
- Like regular algebra, but simpler: variables have only two values (1 or 0)
- **Duality** in axioms and theorems:
  - ANDs and ORs, 0's and 1's interchanged

Boolean Axioms

	Axiom		Dual	Name
A1	$B = 0$ if $B \neq 1$	A1′	$B = 1$ if $B \neq 0$	Binary field
A2	$\overline{0} = 1$	A2′	$\overline{1} = 0$	NOT
A3	$0 \bullet 0 = 0$	A3′	$1 + 1 = 1$	AND/OR
A4	$1 \bullet 1 = 1$	A4′	$0 + 0 = 0$	AND/OR
A5	$0 \bullet 1 = 1 \bullet 0 = 0$	A5′	$1 + 0 = 0 + 1 = 1$	AND/OR

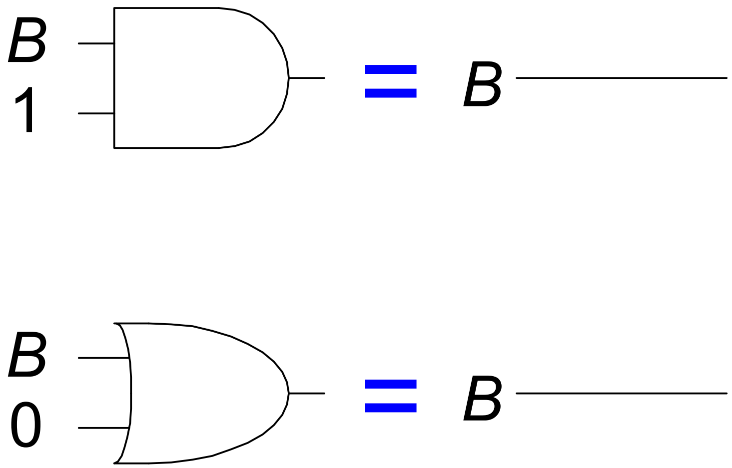
	Theorem		Dual	Name
T1	$B \bullet 1 = B$	T1′	$B + 0 = B$	Identity
T2	$B \bullet 0 = 0$	T2′	$B + 1 = 1$	Null Element
T3	$B \bullet B = B$	T3′	$B + B = B$	Idempotency
T4		$\overline{\overline{B}} = B$		Involution
T5	$B \bullet \overline{B} = 0$	T5′	$B + \overline{B} = 1$	Complements

T1: Identity Theorem

- $B \cdot 1 = B$
- $B + 0 = B$

T1: Identity Theorem

- $B \cdot 1 = B$
- $B + 0 = B$



T2: Null Element Theorem

- $B \cdot 0 = 0$
- $B + 1 = 1$

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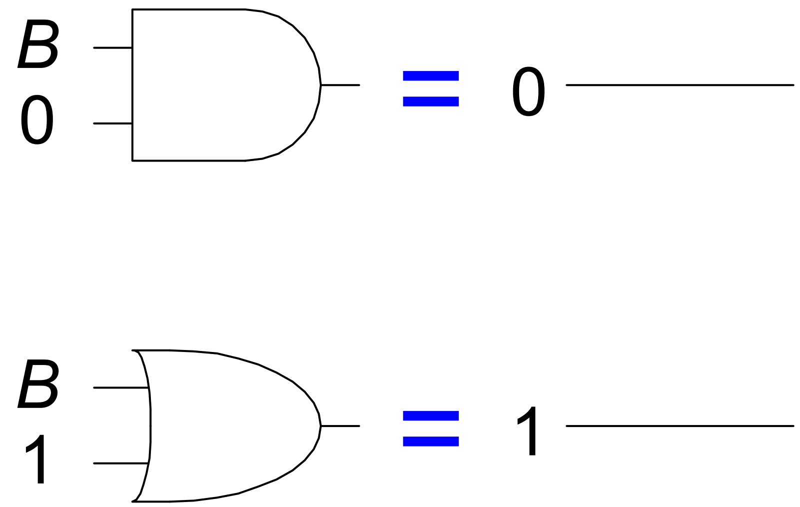
- $B \cdot 0 = 0$
- $B + 1 = 1$



T3: Idempotency Theorem

- $B \cdot B = B$
- $B + B = B$

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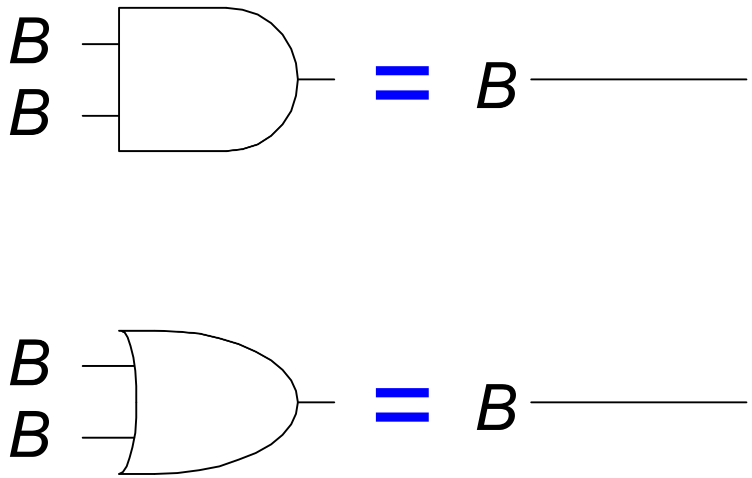
T3: Idempotency Theorem

- $B \cdot B = B$
- $B + B = B$



T4: Identity Theorem

- $\overline{\overline{B}} = B$

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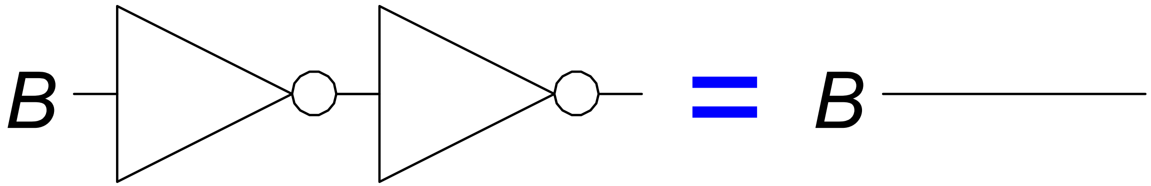
- $\overline{\overline{B}} = B$



T5: Complement Theorem

- $B \bullet \overline{B} = 0$
- $B + \overline{B} = 1$

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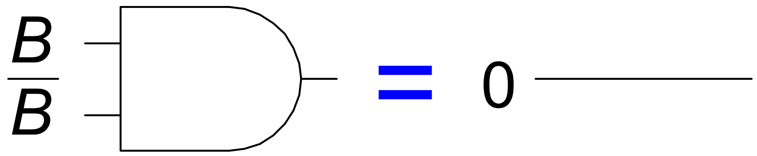

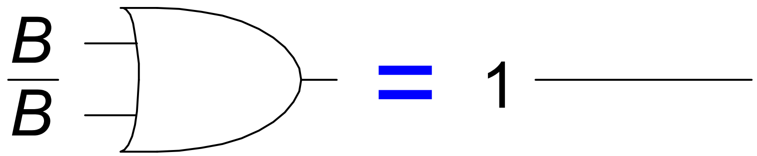

Boolean Theorems Summary

	Theorem		Dual	Name
T1	$B \bullet 1 = B$	T1′	$B + 0 = B$	Identity
T2	$B \bullet 0 = 0$	T2′	$B + 1 = 1$	Null Element
T3	$B \bullet B = B$	T3′	$B + B = B$	Idempotency
T4		$\bar{\bar{B}} = B$		Involution
T5	$B \bullet \bar{B} = 0$	T5′	$B + \bar{B} = 1$	Complements

Boolean Theorems of Several Vars

	Theorem		Dual	Name
T6	$B \bullet C = C \bullet B$	T6′	$B + C = C + B$	Commutativity
T7	$(B \bullet C) \bullet D = B \bullet (C \bullet D)$	T7′	$(B + C) + D = B + (C + D)$	Associativity
T8	$(B \bullet C) + B \bullet D = B \bullet (C + D)$	T8′	$(B + C) \bullet (B + D) = B + (C \bullet D)$	Distributivity
T9	$B \bullet (B + C) = B$	T9′	$B + (B \bullet C) = B$	Covering
T10	$(B \bullet C) + (B \bullet \overline{C}) = B$	T10′	$(B + C) \bullet (B + \overline{C}) = B$	Combining
T11	$(B \bullet C) + (\overline{B} \bullet D) + (C \bullet D)$ $= B \bullet C + \overline{B} \bullet D$	T11′	$(B + C) \bullet (\overline{B} + D) \bullet (C + D)$ $= (B + C) \bullet (\overline{B} + D)$	Consensus
T12	$\overline{B_0 \bullet B_1 \bullet B_2 \ldots}$ $= (\overline{B_0} + \overline{B_1} + \overline{B_2} \ldots)$	T12′	$\overline{B_0 + B_1 + B_2 \ldots}$ $= (\overline{B_0} \bullet \overline{B_1} \bullet \overline{B_2})$	De Morgan's Theorem

Simplifying Boolean Equations

Example 1:

- $Y = AB + \overline{A}B$

Simplifying Boolean Equations

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- $Y = AB + \overline{A}B$

$$= B(A + \overline{A}) \quad \text{T8}$$

$$= B(1) \quad \text{T5'}$$

$$= B \quad \text{T1}$$

Simplifying Boolean Equations

Example 2:

- $Y = A(AB + ABC)$

Simplifying Boolean Equations

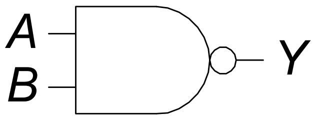
Example 2:

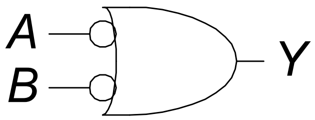
- $Y = A(AB + ABC)$

$= A(AB(1 + C))$	T8
$= A(AB(1))$	T2’
$= A(AB)$	T1
$= (AA)B$	T7
$= AB$	T3

DeMorgan's Theorem

- $Y = \overline{AB} = \overline{A} + \overline{B}$





- $Y = \overline{A + B} = \overline{A} \bullet \overline{B}$