



ELTE
EÖTVÖS LORÁND
UNIVERSITY



CRYPTOGRAPHY AND SECURITY

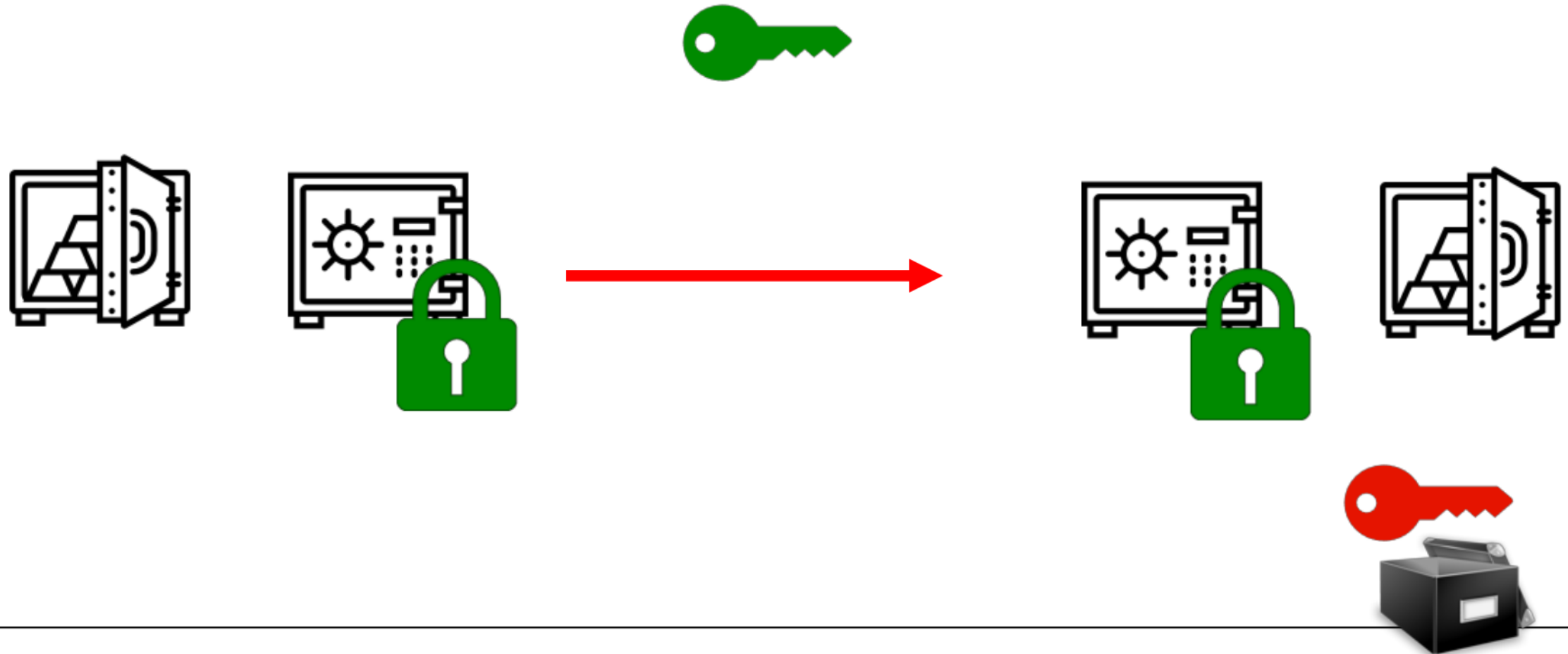
Practice

IP-18FKVKRBG

Lecture 9

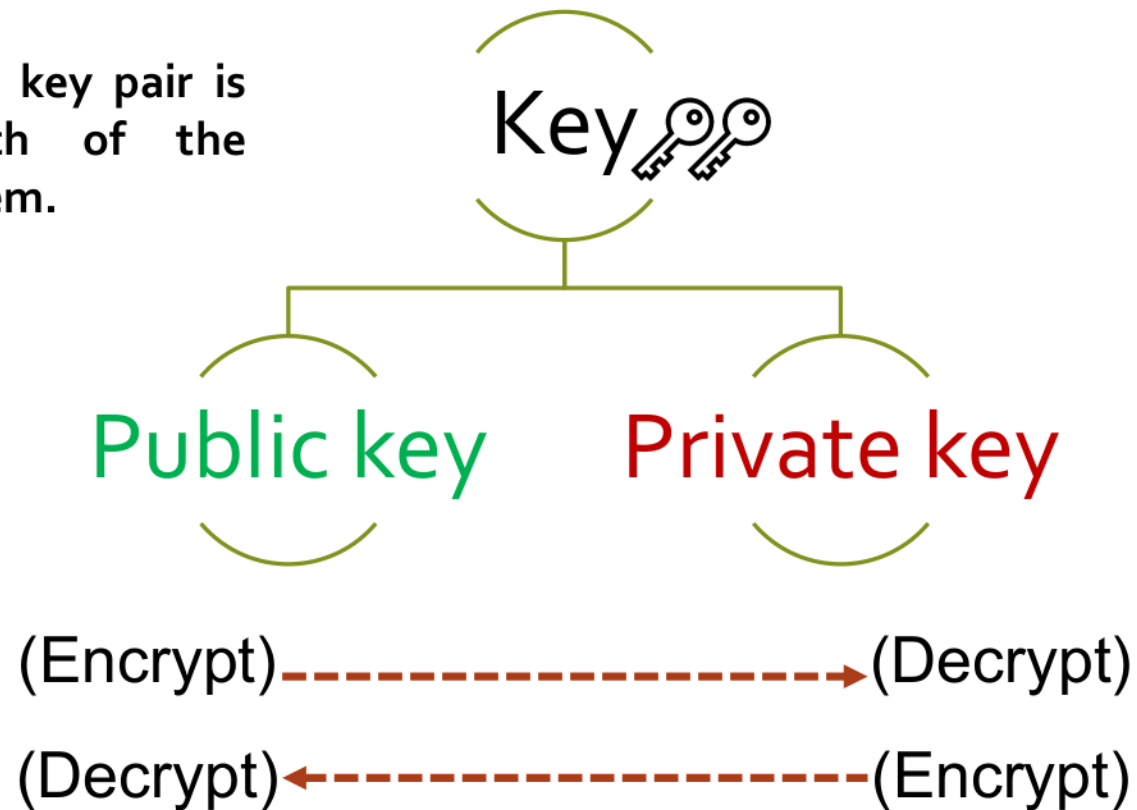
Asymmetric Cryptography

Asymmetric Ciphers



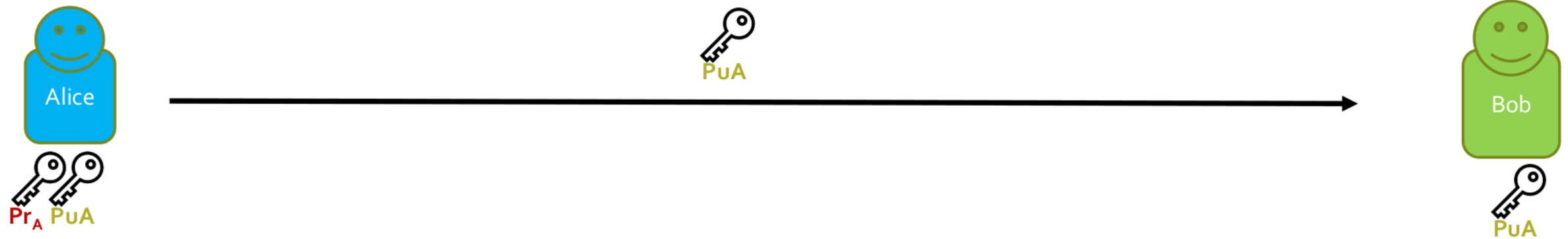
Asymmetric (Public-Key) Cryptography

In key generation, a key pair is generated for each of the members of the system.

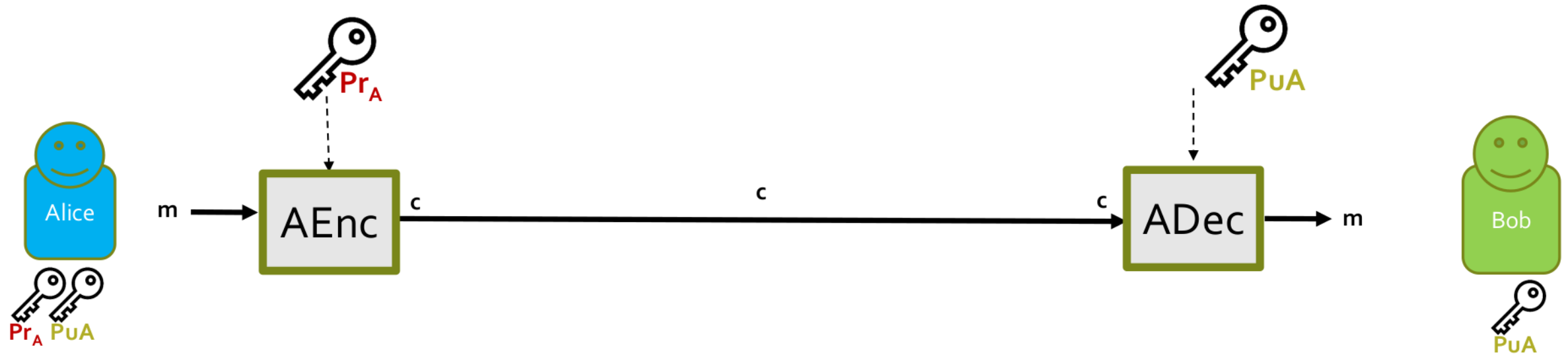


A text Can NOT be encrypted and then decrypted by the SAME key

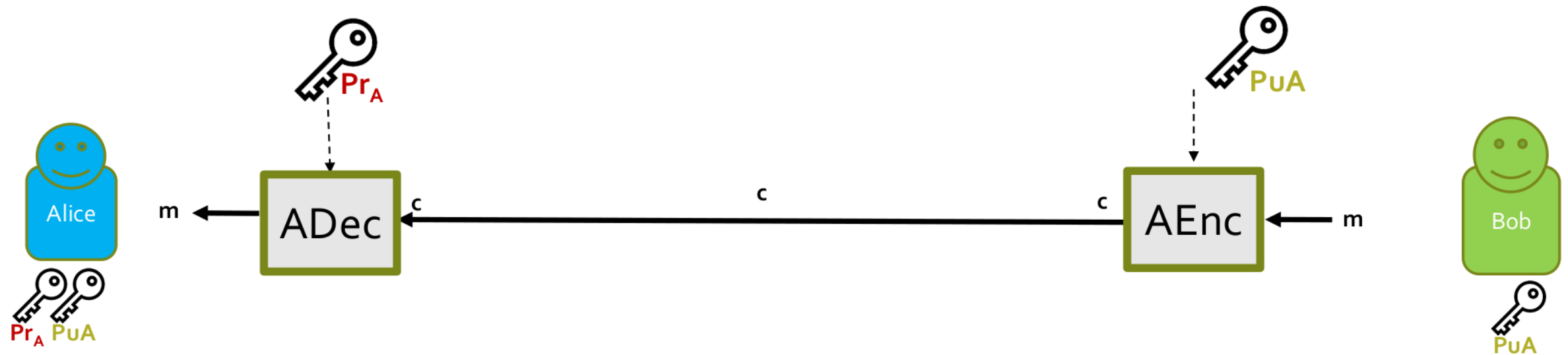
Public-Key Encryption: Sharing public key



Public-Key Encryption: Sender Authentication



Public-Key Encryption: Confidentiality



RSA

RSA: Key generation

- Choose two unequal large prime numbers p and q
- Compute modulo $n = p \cdot q$
- Compute totient function $\varphi(n) = \varphi(p \cdot q) = (p - 1)(q - 1)$
- Generate numbers e, d such that
 - $1 < e < \varphi(n)$ and e relatively prime to $\varphi(n)$
 - Find d such that $ed \equiv 1 \pmod{\varphi(n)}$ \rightarrow it only exist if $\gcd(e, \varphi(n)) = 1$
 - $x^{ed} \equiv x \pmod{n}$

RSA: “schoolbook” Encryption and Decryption

- Public encryption key $\langle n, e \rangle$
 - $\text{Encrypt}(\langle n, e \rangle, x) = x^e \bmod n$
- Private decryption key $\langle n, d \rangle$
 - $\text{Decrypt}(\langle n, d \rangle, y) = y^d \bmod n = x^{ed} \bmod n = x$

RSA: Example

Preparation (Bob Side)

- Bob chooses $p = 41$, $q = 61$, \rightarrow and gets $n = 2501$, $\varphi(n) = 2400$
- He chooses $e = 23$, \rightarrow and gets $d = 2087$
 - By choosing other values of e we get other values of d .
- Alice receives n and e , from Bob

Alice (Sender) Side

- Plaintext = **100**
- Encryption: $100^{23} \bmod 2501 = \mathbf{2306}$

Bob (Receiver) Side

- Decryption: $2306^{2087} \bmod 2501 = \mathbf{100}$

RSA Coding: Necessary functions

```
from sympy import randprime
from math import gcd

def egcd(a, b):
    if a == 0:
        return (b, 0, 1)
    else:
        g, x, y = egcd(b % a, a)
        return (g, y - (b // a) * x, x)

def modinv(b, n):
    g, x, _ = egcd(b, n)
    if g == 1:
        return x % n
```

RSA Coding: Generate two prime numbers

It makes sure that:

- both generated prime numbers are not the same, and
- their product does not exceed the final number of bits (key size)

```
while prime_1 == prime_2 or (prime_1 * prime_2) > 2**key_size:  
    prime_1 = randprime(3, 2**key_size/2)  
    prime_2 = randprime(3, 2**key_size/2)
```

Note: For in lecture testing choose a small key size.

RSA Coding: Generating the public parameter

```
def generate_public_exponents(totient):  
    public_exponent = 0  
    for e in range(3, totient-1):  
        if gcd(e, totient) == 1:  
            public_exponent = e  
            break  
    return public_exponent
```

RSA Coding: Generating the e and d parameters

```
public_exponent = generate_public_exponents(totient)
private_exponent = modinv(public_exponent, totient)
```

RSA Coding: Encryption and decryption

```
def rsa_encrypt(plain_text, rsa_modulus, public_exponent):  
    cipher = ''  
    for ch in plain_text:  
        cipher += chr((ord(ch)**public_exponent) % rsa_modulus)  
    return cipher.encode().hex()  
  
def rsa_decrypt(cipher_text, rsa_modulus, private_exponent):  
    return ''.join(chr((ord(ch)**private_exponent) % rsa_modulus)  
                    for ch in bytearray.fromhex(cipher_text).decode()))
```