# Biblioteca de Gabriel Pessoa:)

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# 1 String Algorithms

# 1.1 String Alignment

```
int pd[ms][ms];
int edit_distance(string &a, string &b) {
    int n = a.size(), m = b.size();
    for(int i = 0; i <= n; i++) pd[i][0] = i;
    for(int j = 0; j <= m; j++) pd[0][j] = j;
    for(int i = 1; i <= n; i++) {
        for(int j = 1; j <= m; j++) {
            int del = pd[i][j-1] + 1;
            int ins = pd[i-1][j] + 1;
            int mod = pd[i-1][j-1] + (a[i-1] != b[j-1]);
            pd[i][j] = min(del, min(ins, mod));
    }
    return pd[n][m];
}</pre>
```

### 1.2 KMP

```
string p, t;
int b[ms], n, m;
void kmpPreprocess() {
    int i = 0, j = -1;
   b[0] = -1;
    while(i < m) {</pre>
        while (j \ge 0 \&\& p[i] != p[j]) j = b[j];
        b[++i] = ++j;
void kmpSearch() {
   int i = 0, j = 0, ans = 0;
    while(i < n) {</pre>
        while(j \ge 0 \&\& t[i] != p[j]) j = b[j];
        i++; j++;
        if(j == m) {
            //ocorrencia aqui comecando em i - j
            ans++;
            j = borda[j];
```

# 1.3 Trie

 $\frac{2}{2}$ 

2

3

3

3

```
int trie[ms][sigma], terminal[ms], z;
void init() {
    memset(trie[0], -1, sizeof trie[0]);
    z = 1;
int get_id(char c) {
    return c - 'a';
void insert(string &p) {
    int cur = 0;
    for(int i = 0; i < p.size(); i++) {</pre>
       int id = get_id(p[i]);
        if(trie[cur][id] == -1) {
            memset(trie[z], -1, sizeof trie[z]);
            trie[cur][id] = z++;
        cur = trie[cur][id];
    terminal[cur]++;
int count(string &p) {
    int cur = 0;
    for(int i = 0; i < p.size(); i++) {</pre>
        int id = get_id(p[i]);
        if(trie[cur][id] == -1) {
            return false;
        cur = trie[cur][id];
    return terminal[cur];
```

## 1.4 Aho-Corasick

// Construa a Trie do seu dicionario com o codigo acima  $\,$ 

```
int fail[ms], q[ms], front, rear;
void buildFailure() {
    front = 0; rear = 0; q[rear++] = 0;
    while(front < rear) {</pre>
       int node = q[front++];
        for(int pos = 0; pos < sigma; pos++) {</pre>
            int &v = trie[node][pos];
            int f = node == 0 ? 0 : trie[fail[node]][pos];
            if(v == -1) {
               v = f;
            } else {
                fail[v] = f;
                q.push(v);
                // juntar as informacoes da borda para o V ja q um match em V implica um
                     match na borda
                terminal[v] += terminal[f];
int search(string &txt) {
   int node = 0;
    int ans = 0;
   for(int i = 0; i < txt.length(); i++) {</pre>
       int pos = get_id(txt[i]);
       node = trie[node][pos];
        // processar informacoes no no atual
      ans += terminal[node];
    return ans;
```

### 2 Data Structures

### 2.1 BIT - Binary Indexed Tree

```
int arr[ms], bit[ms], n;

void update(int v, int idx) {
   while(idx <= n) {
      bit[idx] += v;
      idx += idx & -idx;
   }
}

int query(int idx) {
   int r = 0;
   while(idx > 0) {
      r += bit[idx];
      idx -= idx & -idx;
   }
   return r;
}
```

# 2.2 Iterative Segment Tree

```
int n, t[2 * ms];

void build() {
    for(int i = n - 1; i > 0; --i) t[i] = t[i<<1] + t[i<<1|1];
}

void update(int p, int value) { // set value at position p
    for(t[p += n] = value; p > 1; p >>= 1) t[p>>1] = t[p] + t[p^1];
}

int query(int 1, int r) {
    int res = 0;
```

```
for(l += n, r += n; l < r; l >>= 1, r >>= 1) {
        if(l&l) res += t[l++];
        if(r&l) res += t[--r];
    }
    return res;
}

// If is non-commutative
S query(int l, int r) {
S resl, resr;
    for (l += n, r += n; l < r; l >>= 1, r >>= 1) {
        if (l&l) resl = combine(resl, t[l++]);
        if (r&l) resr = combine(t[--r], resr);
    }
    return combine(resl, resr);
```

# 2.3 Iterative Segment Tree with Interval Updates

```
int n, t[2 * ms];

void build() {
    for(int i = n - 1; i > 0; --i) t[i] = t[i<<1] + t[i<<1|1];
}

void update(int 1, int r, int value) {
    for(1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
        if(1&1) t[1++] += value;
        if(r&1) t[--r] += value;
    }
}

int query(int p) {
    int res = 0;
    for(p += n; p > 0; p >>= 1) res += t[p];
    return res;
}

void push() { // push modifications to leafs
    for(int i = 1; i < n; i++) {
        t[i<<1] += t[i];
        t[i<<1|1] += t[i];
        t[i] = 0;
}</pre>
```

# 2.4 Recursive Segment Tree

```
int arr[4 * ms], seg[4 * ms], n;
void build(int idx = 0, int l = 0, int r = n - 1) {
    int mid = (1+r)/2, left = 2 * idx + 1, right = 2 * idx + 2;
    if(1 == r) {
        seg[idx] = arr[1];
   build(left, 1, mid); build(right, mid + 1, r);
    seg[idx] = seg[left] + seg[right];
int query(int L, int R, int idx = 0, int l = 0, int r = n - 1) {
    int mid = (1+r)/2, left = 2 * idx + 1, right = 2 * idx + 2;
    if(R < 1 || L > r) return 0;
   if(L <= 1 && r <= R) return seg[idx];</pre>
    return query(L, R, left, l, mid) + query(L, R, right, mid + 1, r);
void update(int V, int I, int idx = 0, int l = 0, int r = n - 1) {
    int mid = (1+r)/2, left = 2 * idx + 1, right = 2 * idx + 2;
    if(1 > I || r < I) return;
   if(1 == r) {
        arr[I] = V;
```

```
seg[idx] = V;
    return;
}
update(V, I, left, 1, mid); update(V, I, right, mid + 1, r);
seg[idx] = seg[left] + seg[right];
}
```

# 2.5 Segment Tree with Lazy Propagation

```
int arr[4 \star ms], seg[4 \star ms], lazy[4 \star ms], n;
void build(int idx = 0, int l = 0, int r = n - 1) {
    int mid = (1+r)/2, left = 2 * idx + 1, right = 2 * idx + 2;
    if(1 == r) {
        seg[idx] = arr[1];
        return:
    build(left, 1, mid); build(right, mid + 1, r);
    seg[idx] = seg[left] + seg[right];
int query (int L, int R, int idx = 0, int 1 = 0, int r = n - 1) {
    int mid = (1+r)/2, left = 2 * idx + 1, right = 2 * idx + 2;
    if(lazy[idx]) {
        seg[idx] += lazy[idx] * (r - l + 1);
        if(1 < r) {
            lazv[left] += lazv[idx];
            lazy[right] += lazy[idx];
        lazy[idx] = 0;
    if(R < 1 || L > r) return 0;
   if(L <= 1 && r <= R) return seg[idx];</pre>
    return query(L, R, left, 1, mid) + query(L, R, right, mid + 1, r);
void update(int V, int L, int R, int idx = 0, int l = 0, int r = n - 1) {
    int mid = (1+r)/2, left = 2 * idx + 1, right = 2 * idx + 2;
    if(lazy[idx]) {
        seg[idx] += lazy[idx] * (r - l + 1);
        if(1 < r) {
            lazy[left] += lazy[idx];
            lazy[right] += lazy[idx];
        lazy[idx] = 0;
    if(1 > R \mid \mid r < L) return;
    if(L <= 1 && r <= R) {
        lazv[idx] += V;
        return:
    update(V, L, R, left, l, mid); update(V, L, R, right, mid + 1, r);
    seg[idx] = seg[left] + seg[right];
```

# 2.6 Color Updates Structure

```
struct range {
    int 1, r;
    int v;

    range(int 1 = 0, int r = 0, int v = 0) : 1(1), r(r), v(v) {}

    bool operator < (const range &a) const {
        return 1 < a.1;
    }
};

set<range> ranges;

vector<range> update(int 1, int r, int v) { // [1, r)
        vector<range> ans;
```

```
if(l >= r) return ans;
auto it = ranges.lower_bound(1);
if(it != ranges.begin()) {
        it--:
        if(it->r>1) {
               auto cur = *it;
                ranges.erase(it);
                ranges.insert(range(cur.1, 1, cur.v));
               ranges.insert(range(l, cur.r, cur.v));
it = ranges.lower_bound(r);
if(it != ranges.begin()) {
       it--;
        if(it->r>r) {
               auto cur = *it;
                ranges.erase(it);
                ranges.insert(range(cur.l, r, cur.v));
                ranges.insert(range(r, cur.r, cur.v));
for(it = ranges.lower_bound(1); it != ranges.end() && it->1 < r; it++) {</pre>
        ans.push_back(*it);
ranges.erase(ranges.lower_bound(l), ranges.lower_bound(r));
ranges.insert(range(l, r, v));
return ans;
```

# 3 Graph Algorithms

#### 3.1 Dinic Max Flow

```
const int ms = 1e3; // Quantidade maxima de vertices
const int me = 1e5; // Quantidade maxima de arestas
int adj[ms], to[me], ant[me], wt[me], z, n;
int copy_adj[ms], fila[ms], level[ms];
void clear() {
   memset(adj, -1, sizeof adj);
    z = 0:
int add(int u, int v, int k) {
   to[z] = v;
    ant[z] = adj[u];
    wt[z] = k;
   adj[u] = z++;
    swap(u, v);
    to[z] = v;
    ant[z] = adj[u];
    wt[z] = 0;
    adj[u] = z++;
int bfs(int source, int sink) {
        memset(level, -1, sizeof level);
        level[source] = 0;
        int front = 0, size = 0, v;
        fila[size++] = source;
        while(front < size) {</pre>
                v = fila[front++];
                for(int i = adj[v]; i != -1; i = ant[i]) {
                        if(wt[i] && level[to[i]] == -1) {
                                level[to[i]] = level[v] + 1;
                                fila[size++] = to[i];
        return level[sink] != -1;
int dfs(int v, int sink, int flow) {
```

```
if(v == sink) return flow;
        int f:
        for(int &i = copy_adj[v]; i != -1; i = ant[i]) {
                if(wt[i] && level[to[i]] == level[v] + 1 && (f = dfs(to[i], sink, min(flow,
                     wt[i])))) {
                        wt[i] -= f;
                        wt[i ^ 1] += f;
                        return f;
        return 0:
int maxflow(int source, int sink) {
        int ret = 0, flow;
        while(bfs(source, sink)) {
                memcpy(copy_adj, adj, sizeof adj);
                while((flow = dfs(source, sink, 1 << 30))) {</pre>
                        ret += flow:
        return ret:
```

## 3.2 Articulations Points and Bridges

```
int idx[ms], art[ms], bridge[me], ind, child;
int dfs(int v, int par = -1) {
   int low = idx[v] = ind++;
    for(int i = adj[v]; i > -1; i = ant[i]) {
        if(idx[to[i]] == -1) {
           if(par == -1) child++;
            int temp = dfs(to[i], v);
            if(temp >= idx[v]) art[v] = true;
           if(temp > idx[v]) bridge[i] = true;
            low = min(low, temp);
        } else if(to[i] != par) low = min(low, idx[to[i]]);
    return low:
void artPointAndBridge() {
   ind = 0;
    memset(idx, -1, sizeof idx);
    memset(art, 0, sizeof art);
    for(int i = 0; i < n; i++) if(idx[i] == -1) {</pre>
        child = 0;
        dfs(i);
        art[i] = child > 1;
```

# 3.3 Biconnected Components

```
int idx[ms], bc[me], ind, nbc, child, st[me], top;

void generateBc(int edge) {
    while(st[--top] != edge) {
        bc[st[top]] = nbc;
    }
    bc[edge] = nbc++;
}

int dfs(int v, int par = -1) {
    int low = idx[v] = ind++;
    for(int i = adj[v]; i > -1; i = ant[i]) {
        if(idx[to[i]] == -1) {
            if(par == -1) child++;
            st[top++] = i;
            int temp = dfs(to[i], v);
            if(par == -1 && child > 1 || ~par && temp >= idx[v]) generateBc(i);
```

```
if(temp >= idx[v]) art[v] = true;
    if(temp > idx[v]) bridge[i] = true;
    low = min(low, temp);
} else if(to[i]! = par && idx[to[i]] < low) {
    low = idx[to[i]];
    st[top++] = i;
}

return low;
}

void biconnected() {
    ind = 0;
    nbc = 0;
    top = -1;
    memset(idx, -1, sizeof idx);
    for(int i = 0; i < n; i++) if(idx[i] == -1) {
        child = 0;
        dfs(i);
}</pre>
```

# 3.4 SCC - Strongly Connected Components / 2SAT

```
int idx[ms], low[ms], ind, comp[ms], ncomp, st[ms], top;
int dfs(int v) {
    if(~idx[v]) return idx[v] ? idx[v] : ind;
    low[v] = idx[v] = idx++;
    st[top++] = v;
    for(int w = adj[v]; ~w; w = ant[w]) {
        low[v] = min(low[v], dfs(to[w]));
    if(low[v] == idx[v]) {
        idx[st[--top]] = 0;
        while(st[top] != v)
            int w = st[--top];
            idx[w] = 0;
            low[w] = low[v];
            comp[w] = ncomp;
        comp[v] = ncomp++;
    return low[v];
bool solveSat() {
   memset(idx, -1, sizeof idx);
    ind = 1; top = -1;
    for(int i = 0; i < n; i++) dfs(i);</pre>
    for(int i = 0; i < n; i++) if(comp[i] == comp[i^1]) return false;</pre>
// Operacoes comuns de 2-sat
// ~v = "nao v"
#define trad(v) (v<0?((~v)*2)^1:v*2)
void addImp(int a, int b) { add(trad(a), trad(b)); }
void addOr(int a, int b) { addImp(~a, b); addImp(~b, a); }
void addEqual(int a, int b) { addOr(a, ~b); addOr(~a, b); }
void addDiff(int a, int b) { addEqual(a, ~b); }
// valoracao: value[v] = comp[trad(v)] < comp[trad(~v)]</pre>
```

#### 3.5 LCA - Lowest Common Ancestor

```
int par[ms] [mlg], lvl[ms];

void dfs(int v, int p, int l = 0) {
    lvl[v] = 1;
    par[v][0] = p;
    for(int i = adj[v]; i > - 1; i = ant[i]) {
        if(to[i] != p) dfs(to[i], v, l + 1);
}
```

```
}

void processAncestors(int root = 0) {
    dfs(root, root);
    for(int k = 1; k <= mlg; k++) {
        for(int i = 0; i < n; i++) {
            par[i][k] = par[par[i][k-1]][k-1];
        }
    }
}

int lca(int a, int b) {
    if(lv1[b] > lv1[a]) swap(a, b);
    for(int i = mlg; i >= 0; i--) {
        if(lv1[a] - (1 << i) >= lv1[b]) a = par[a][i];
    }
    if(a == b) return a;
    for(int i = mlg; i >= 0; i--) {
        if(par[a][i] != par[b][i]) a = par[a][i], b = par[b][i];
    }
    return par[a][0];
}
```

### 4 Math

# 4.1 Discrete Logarithm

```
11 discreteLog(ll a, ll b, ll m) {
        // a^ans == b mod m
        // ou -1 se nao existir
        11 \text{ cur} = a, \text{ on } = 1;
        for (int i = 0; i < 100; i++) {
                 cur = cur * a % m;
        while (on \star on \leftarrow m) {
                 cur = cur * a % m;
                 on++:
        map<11, 11> position;
        for (ll i = 0, x = 1; i * i <= m; i++) {
                 position[x] = i * on;
                 x = x * cur % m;
        for(ll i = 0; i <= on + 20; i++) {
                 if(position.count(b)) {
                          return position[b] - i;
                 b = b * a % m;
        return -1;
```

#### 4.2 GCD - Greatest Common Divisor

```
11 gcd(11 a, 11 b) {
     while(b) a %= b, swap(a, b);
     return a;
}
```

#### 4.3 Extended Euclides

```
// euclides estendido: acha u e v da equacao: 
// u * x + v * y = gcd(x, y); 
// u eh inverso modular de x no modulo y 
// v eh inverso modular de y no modulo x
```

### 4.4 Fast Exponentiation

# 4.5 Matrix Fast Exponentiation

```
const 11 \mod = 1e9+7;
const int m = 2; // size of matrix
struct Matrix {
        11 mat[m][m];
        Matrix operator * (const Matrix &p) {
                 Matrix ans;
                 for (int i = 0; i < m; i++)</pre>
                        for(int j = 0; j < m; j++)</pre>
                                 for(int k = ans.mat[i][j] = 0; k < m; k++)</pre>
                                         ans.mat[i][j] = (ans.mat[i][j] + mat[i][k] * p.mat[k]
                                                ][j]) % mod;
                 return ans;
        }
};
Matrix fExp(Matrix a, ll b) {
        Matrix ans;
        for (int i = 0; i < m; i++) for (int j = 0; j < m; j++)
                 ans.mat[i][j] = i == j;
        while(b) {
                if(b \& 1) ans = ans * a;
                 a = a * a;
                 b >>= 1;
        return ans;
```

#### 4.6 FFT - Fast Fourier Transform

```
typedef complex<double> Complex;
const double pi = acosl(-1.0);
int rbit[1 << 23];

void calcReversedBits(int n) {
   int lg = 0;
   while(1 << (lg + 1) < n) {</pre>
```

```
lg++;
    for(int i = 1; i < n; i++) {</pre>
        bits[i] = (bits[i >> 1] >> 1) | ((i & 1) << lq);
void fft(Complex a[], int n, bool inv = false) {
    calcReversedBits(n):
    for(int i = 0; i < n; i++) {</pre>
        if(rbit[i] > i) swap(a[i], a[rbit[i]]);
    double ang = inv ? -pi : pi;
    for(int m = 1; m < n; m += m) {</pre>
        Complex d(cosl(ang/m), sinl(ang/m));
        for (int i = 0; i < n; i += m+m) {
            Complex cur = 1;
            for(int j = 0; j < m; j++) {
                Complex u = a[i + j], v = a[i + j + m] * cur;
                a[i + j] = u + v;
                a[i + j + m] = u - v;
                cur *= d;
   if(inv) {
        for(int i = 0; i < n; i++) a[i] /= n;</pre>
```

# 5 Geometry

# 5.1 Geometry

```
const double inf = 1e100, eps = 1e-9;
struct PT {
   double x, y;
    PT(double x, double y) : x(x), y(y) {}
    PT operator + (const PT &p) { return PT(x + p.x, y + p.y); }
    PT operator - (const PT &p) { return PT(x - p.x, y - p.y); }
    PT operator * (double c) { return PT(x * c, y * c); }
    PT operator / (double c) { return PT(x / c, y / c); }
    bool operator < (const PT &p) const {
                return x < p.x | | (x == p.x && y < p.y);
};
double dot(PT p, PT q) { return p.x * q.x + p.y * q.y; }
double dist2(PT p, PT q) { return dot(p - q, p - q); }
double cross(PT p, PT q) { return p.x * q.y - p.y * q.x; }
// Rotaciona o ponto CCW ou CW ao redor da origem
PT rotateCCW90 (PT p) { return PT(-p.y, p.x); }
PT rotateCW90(PT p) { return PT(p.y, -p.x); }
PT rotateCCW(PT p, double d) {
    return PT(p.x * cos(t) - p.y * sin(t), p.x * sin(t) + p.y * cos(t));
// Projeta ponto c na linha a - b assumindo a != b
PT projectPointLine(PT a, PT b, PT c) {
    return a + (b - a) * dot(c - a, b - a) / dot(b - a, b - a);
// Projeta ponto c no segmento a - b
PT projectPointSegment (PT a, PT b, PT c) {
    double r = dot(b - a, b - a);
    if(abs(r) < eps) return a;</pre>
    r = dot(c - a, b - a) / r;
    if(r < 0) return a;</pre>
    if(r > 1) return b;
    return a + (b - a) * r;
```

```
// Calcula distancia entre o ponto c e o segmento a - b
double distancePointSegment (PT a, PT b, PT c) {
    return sqrt(dist2(c, projectPointSegment(a, b, c)));
// Calcula distancia entre o ponto (x, y, z) e o plano ax + by + cz = d
double distancePointPlane(double x, double y, double z, double a, double b, double c, double
    return abs(a * x + b * y + c * z - d) / sqrt(a * a + b * b + c * c);
// Determina se as linhas a - b e c - d sao paralelas ou colineares
bool linesParallel(PT a, PT b, PT c, PT d) {
    return abs(cross(b - a, c - d)) < eps;</pre>
bool linesCollinear(PT a, PT b, PT c, PT d) {
    return linesParallel(a, b, c, d) && abs(cross(a - b, a - c)) < eps && abs(cross(c - d, c
         - a)) < eps;</pre>
// Determina se o segmento a - b intersecta com o segmento c - d
bool segmentsIntersect(PT a, PT b, PT c, PT d) {
    if(linesCollinear(a, b, c, d)) {
        if(dist2(a, c) < eps || dist2(a, d) < eps || dist2(b, c) < eps || dist2(b, d) < eps)</pre>
             return true;
        if(dot(c - a, c - b) > 0 && dot(d - a, d - b) > 0 && dot(c - b, d - b) > 0) return
             false:
        return true;
    if(cross(d - a, b - a) * cross(c - a, b - a) > 0) return false;
    if(cross(a - c, d - c) * cross(b - c, d - c) > 0) return false;
    return true:
// Calcula a intersecao entre as linhas a - b e c - d assumindo que uma unica intersecao
// Para intersecao de segmentos, cheque primeiro se os segmentos se intersectam
PT computeLineIntersection(PT a, PT b, PT c, PT d) {
   b = b - a; d = c - d; c = c - a;
    return a + b * cross(c, d) / cross(b, d);
// Calcula centro do circulo dado tres pontos
PT computeCircleCenter(PT a, PT b, PT c) {
   b = (a + b) / 2;
    c = (a + c) / 2;
    return computeLineIntersection(b, b + rotateCW90(a - b), c, c + rotateCW90(a - c));
// Determina se o ponto esta num poligno possivelmente nao-convexo
// Retorna 1 para pontos estritamente dentro, 0 para pontos estritamente fora do poligno
// e 0 ou 1 para os pontos restantes
// Eh possivel converter num teste exato usando inteiros e tomando cuidado com a divisao
// e entao usar testes exatos para checar se esta na borda do poligno
bool pointInPolygon(const vector<PT> &p, PT q) {
  bool c = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int j = (i + 1) % p.size();
     \textbf{if} ( (p[i].y \mathrel{<=} q.y \&\& q.y \lessdot p[j].y \mid\mid p[j].y \mathrel{<=} q.y \&\& q.y \lessdot p[i].y ) \&\& \\
      q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) / (p[j].y - p[i].y))
  return c:
// Determina se o ponto esta na borda do poligno
bool pointOnPolygon(const vector<PT> &p, PT q) {
  for(int i = 0; i < p.size(); i++)</pre>
    if(dist2(projectPointSegment(p[i], p[(i + 1) % p.size()], q), q) < eps)
      return true;
    return false;
// Calcula intersecao da linha a - b com o circulo centrado em c com raio r > 0
vector<PT> circleLineIntersection(PT a, PT b, PT c, double r) {
 vector<PT> ans;
```

b = b - a;

```
a = a - c;
  double x = dot(b, b);
  double y = dot(a, b);
  double z = dot(a, a) - r * r;
  double w = y * y - x * z;
  if (w < -eps) return ans;</pre>
  ans.push_back(c + a + b * (-y + sqrt(w + eps)) / x);
  if (w > eps)
   ans.push_back(c + a + b \star (-y - sqrt(w)) / x);
  return ans;
// Calcula intersecao do circulo centrado em a com rajo r e o centrado em b com rajo R
vector<PT> circleCircleIntersection(PT a, PT b, double r, double R) {
  vector<PT> ans:
  double d = sqrt(dist2(a, b));
  if (d > r + R \mid \mid d + min(r, R) < max(r, R)) return ans;
  double x = (d * d - R * R + r * r)/(2 * d);
  double y = sqrt(r * r - x * x);
  PT v = (b - a) / d;
  ans.push_back(a + v * x + rotateCCW90(v) * y);
   ans.push_back(a + v * x - RotateCCW90(v) * y);
// Calcula a area ou o centroide de um poligono (possivelmente nao-convexo)
// assumindo que as coordenadas estao listada em ordem horaria ou anti-horaria
// O centroide eh equivalente a o centro de massa ou centro de gravidade
double computeSignedArea(const vector<PT> &p) {
  double area = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
   int j = (i + 1) % p.size();
   area += p[i].x * p[j].y - p[j].x * p[i].y;
  return area / 2.0;
double computeArea(const vector<PT> &p) {
  return abs(computeSignedArea(p));
PT computeCentroid(const vector<PT> &p) {
  PT c(0,0);
  double scale = 6.0 * ComputeSignedArea(p);
  for(int i = 0; i < p.size(); i++) {</pre>
   int j = (i + 1) % p.size();
    c = c + (p[i] + p[j]) * (p[i].x * p[j].y - p[j].x * p[i].y);
  return c / scale;
// Testa se o poligno listada em ordem CW ou CCW eh simples (nenhuma linha se intersecta)
bool isSimple(const vector<PT> &p) {
  for(int i = 0; i < p.size(); i++) {</pre>
    for(int k = i + 1; k < p.size(); k++) {</pre>
     int j = (i + 1) % p.size();
      int 1 = (k + 1) % p.size();
      if (i == 1 \mid \mid j == k) continue;
      if (segmentsIntersect(p[i], p[j], p[k], p[l]))
        return false:
  return true;
```

#### 5.2 Convex Hull

```
vector<PT> convexHull(vector<PT> p)) {
   int n = p.size(), k = 0;
   vector<PT> h(2 * n);
   sort(p.begin(), p.end());
   for(int i = 0; i < n; i++) {
      while(k >= 2 && cross(h[k - 1] - h[k - 2], p[i] - h[k - 2]) <= 0) k--;
      h[k++] = p[i];
   }</pre>
```

```
for(int i = n - 2, t = k + 1; i >= 0; i--) {
    while(k >= t && cross(h[k - 1] - h[k - 2], p[i] - h[k - 2]) <= 0) k--;
    h[k++] = p[i];
}
h.resize(k);
return h;</pre>
```

#### 5.3 ClosestPair

```
double closestPair(vector<PT> p) {
   int n = p.size(), k = 0;
   sort(p.begin(), p.end());
   double d = inf;
   set<PT> ptsInv;
   for(int i = 0; i < n; i++) {
        while(k < i && p[k].x < p[i].x - d) {
            ptsInv.erase(swapCoord(p[k++]));
        }
        for(auto it = ptsInv.lower_bound(PT(p[i].y - d, p[i].x - d));
            it!= ptsInv.end() && it->x <= p[i].y + d; it++) {
                 d = min(d, !(p[i] - swapCoord(*it)));
            }
            ptsInv.insert(swapCoord(p[i]));
        return d;
}</pre>
```

#### 5.4 Intersection Points

```
// Utiliza uma seg tree
int intersectionPoints(vector<pair<PT, PT>> v) {
    int n = v.size();
    vector<pair<int, int>> events, vertInt;
    for(int i = 0; i < n; i++) {
        if(v.first.x == v.second.x) { // Segmento Vertical
            int y0 = min(v.first.y, v.second.y), y1 = max(v.first.y, v.second.y);
            events.push_back({v.first.x, vertInt.size()}); // Tipo = Indice no array
            vertInt.push_back({y0, y1});
        } else { // Segmento Horizontal
            int x0 = min(v.first.x, v.second.x), x1 = max(v.first.x, v.second.x);
            events.push_back({x0, -1}); // Inicio de Segmento
            events.push_back({x1, inf}); // Final de Segmento
    sort(events.begin(), events.end());
    int ans = 0;
    for(int i = 0; i < events.size(); i++) {</pre>
        int t = events[i].second;
        if(t == -1) {
            segUpdate(events[i].first, 1);
        } else if(t == inf)
           segUpdate(events[i].first, 0);
        } else {
           ans += segQuery(vertInt[t].first, vertInt[t].second);
    return ans;
```

# 5.5 Maximum Segments Overlap (Sweep Line)

```
int maxSegOverlap(vector<pair<int, int>> v) {
   int n = v.size();
   vector<pair<int, int>> events(n * 2);
   // tipo 1 = Inicio de segmento
   // tipo 0 = Fim de segmento
   for(int i = 0; i < n; i++) {
      events[i*2] = {v[i].first, 1};
}</pre>
```

# 6 Miscellaneous

# 6.1 LIS - Longest Increasing Subsequence

```
}
return len;
}
vi getLis() {
   int len = lis();
   vi ans;
   for(int i = pos[lisArr[len - 1]]; i >= 0; i = bef[i]) {
        ans.push_back(arr[i]);
   }
   reverse(ans.begin(), ans.end());
   return ans;
}
```

# 6.2 Binary Search

```
int smallestSolution() {
    int x = -1;
    for(int b = z; b >= 1; b /= 2) {
        while(!ok(x+b)) x += b;
    }
    return x + 1;
}

int maximumValue() {
    int x = -1;
    for(int b = z; b >= 1; b /= 2) {
        while(f(x+b) < f(x+b+1)) x += b;
    }
    return x + 1;
}</pre>
```