Graphical Representation for Global Protocols

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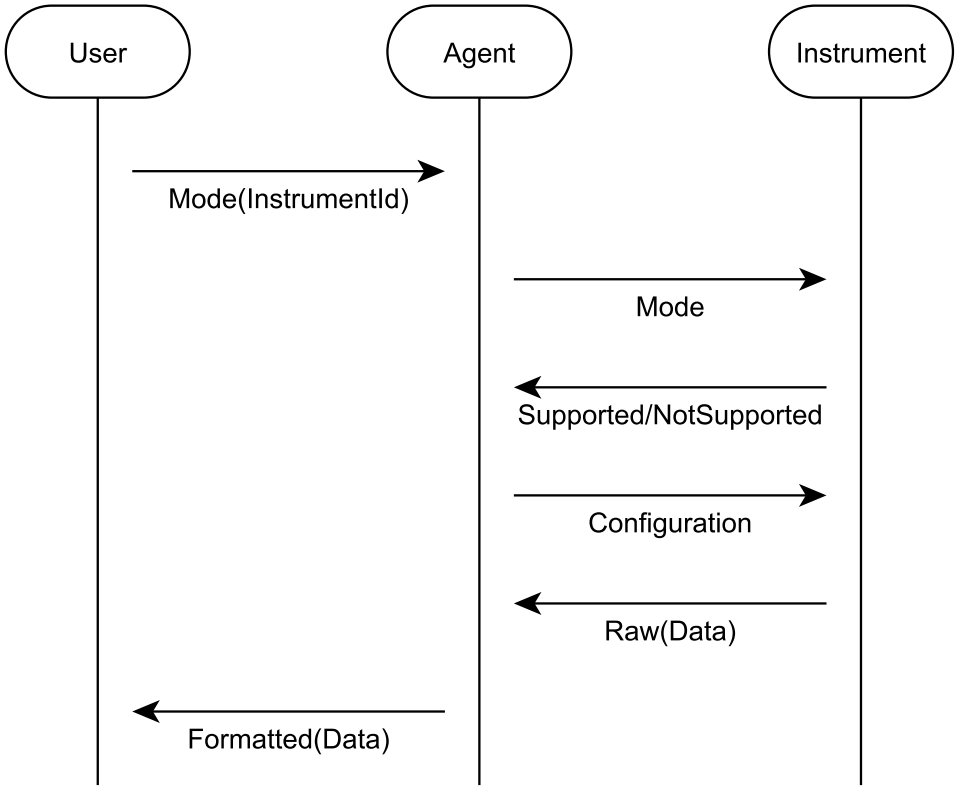
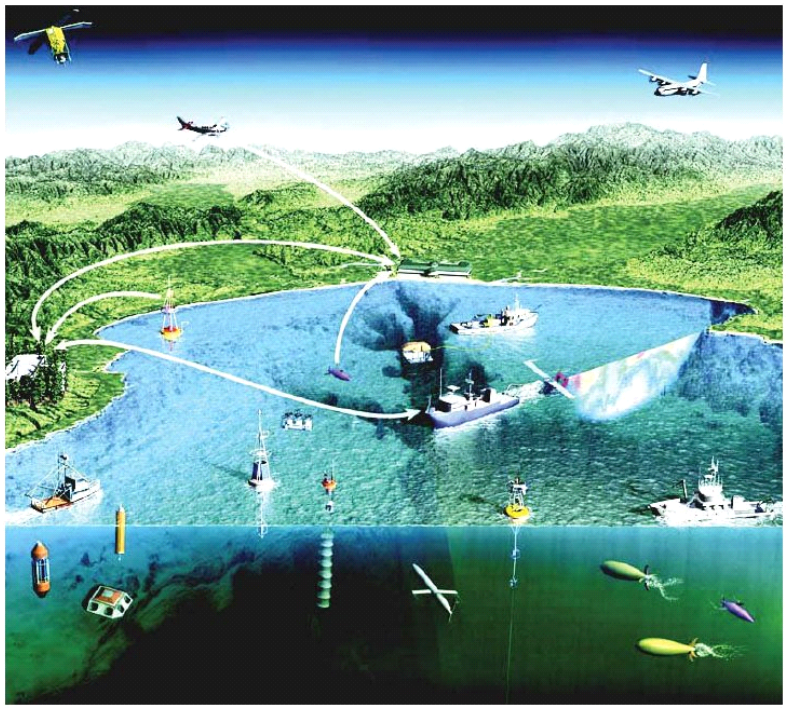
Ocean Observatories Initiative Use Case : DataAcquisition





Global Protocol in SCRIBBLE

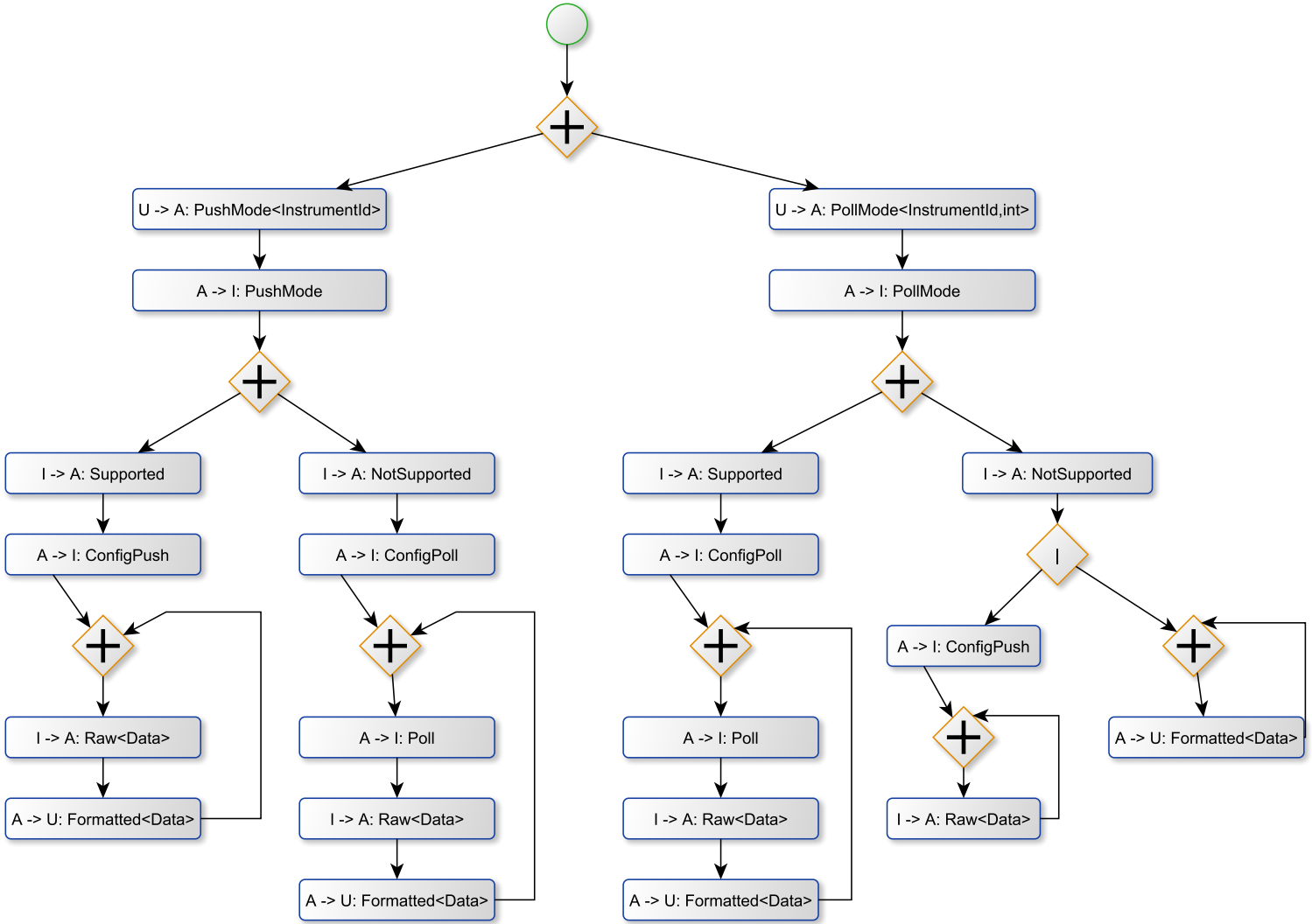
```
0 // U is User, A is ION Agent (Integrated
1 // Observatories Network), I is Instrument

2 global protocol DataAcquisition (role U, role A, role I) {

3 interruptible { choice at U {
4     PushMode(InstrumentId) from U to A;
5     PushMode from A to I;
6     choice at I {
7           Supported from I to A;
8           ConfigPush from A to I;
9           rec PUSH {
10                Raw(Data) from I to A;
11                Formatted(Data) from A to U;
12                continue  PUSH;}
13    } or {NotSupported from I to A;
14          ConfigPoll from A to I;
15          rec POLL {
16                Poll from A to I;
17                Raw(Data) from I to A;
18                Formatted(Data) from A to U;
19                continue  POLL;} }
20 } or {   PollMode(InstrumentId, int) from U to A;
21    PollMode from A to I;
22    choice at I {
23          Supported from I to A;
24          ConfigPoll from A to I;
25          rec POLL {
26                Poll from A to I;
27                Raw(Data) from I to A;
28                Formatted(Data) from A to U;
29                continue POLL;}
30    } or {NotSupported from I to A;
31          parallel {
32                ConfigPush from A to I;
33                rec PUSH {
34                      Raw(Data) from I to A;
35                      continue PUSH;}
36          } and {     rec POLL {
37                      Formatted(Data) from A to U;
38                      continue POLL; }} }}

39 } by U with Stop
40 }
```

Graphical Representation



Generalised Multiparty Session Types: Global Types

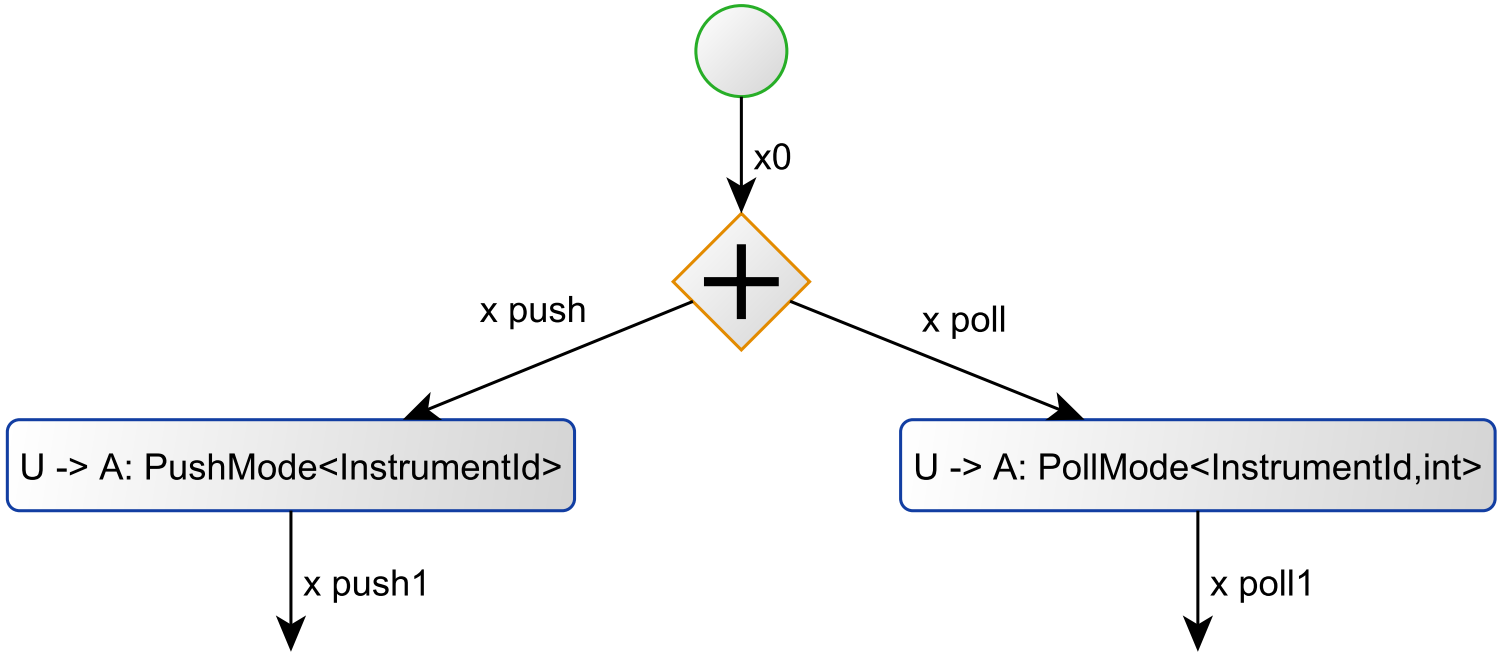
G	::=	def G in x	Global type
G	::=	x = p $\rightarrow$ q : $l\langle U\rangle$; x'	Labelled messages
	|	x = x' | x"	Fork
	|	x = x' + x"	Choice
	|	x | x' = x"	Join
	|	x + x' = x"	Merge
	|	x = end	End
U	::=	$\langle G\rangle$ | bool | nat | ...	Sorts

G = def

$$
\begin{aligned}
x_0 &= x_{push} + x_{poll} \\
x_{push} &= \text{U} \rightarrow \text{A: PushMode}\langle\ \text{InstrumentId}\ \rangle; x_{push1} \\
x_{push1} &= \text{A} \rightarrow \text{I: PushMode}\ ; x_{push2} \\
x_{push2} &= x_{ps} + x_{pns} \\
x_{ps} &= \text{I} \rightarrow \text{A: Supported}; x_{ps1} \\
x_{ps1} &= \text{A} \rightarrow \text{I: ConfigPush}; x_{ps2} \\
x_{ps2} + x_{ps3} &= x_{ps4} \\
x_{ps4} &= \text{I} \rightarrow \text{A: Raw}\langle\ \text{Data}\ \rangle; x_{ps5} \\
x_{ps5} &= \text{A} \rightarrow \text{U: Formatted}\langle\ \text{Data}\ \rangle; x_{ps3} \\
x_{pns} &= \text{I} \rightarrow \text{A: NotSupported}; x_{pns1} \\
x_{pns1} &= \text{A} \rightarrow \text{I: ConfigPoll}; x_{pns2}\ \ldots
\end{aligned}
$$

in x_0

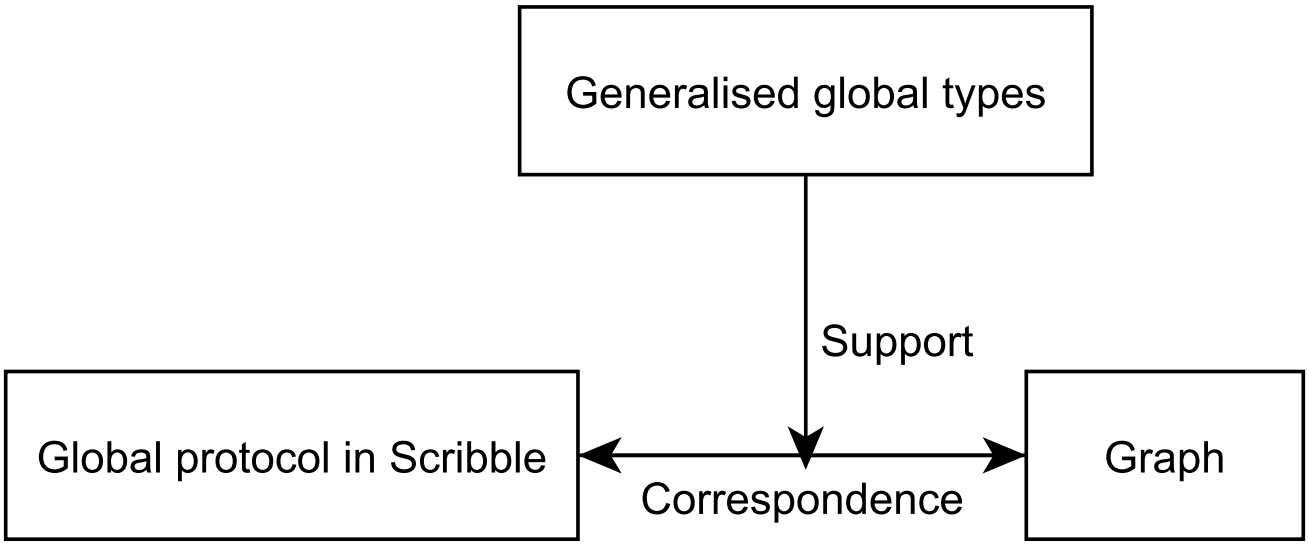
Generalised Multiparty Session Types: Graph syntax



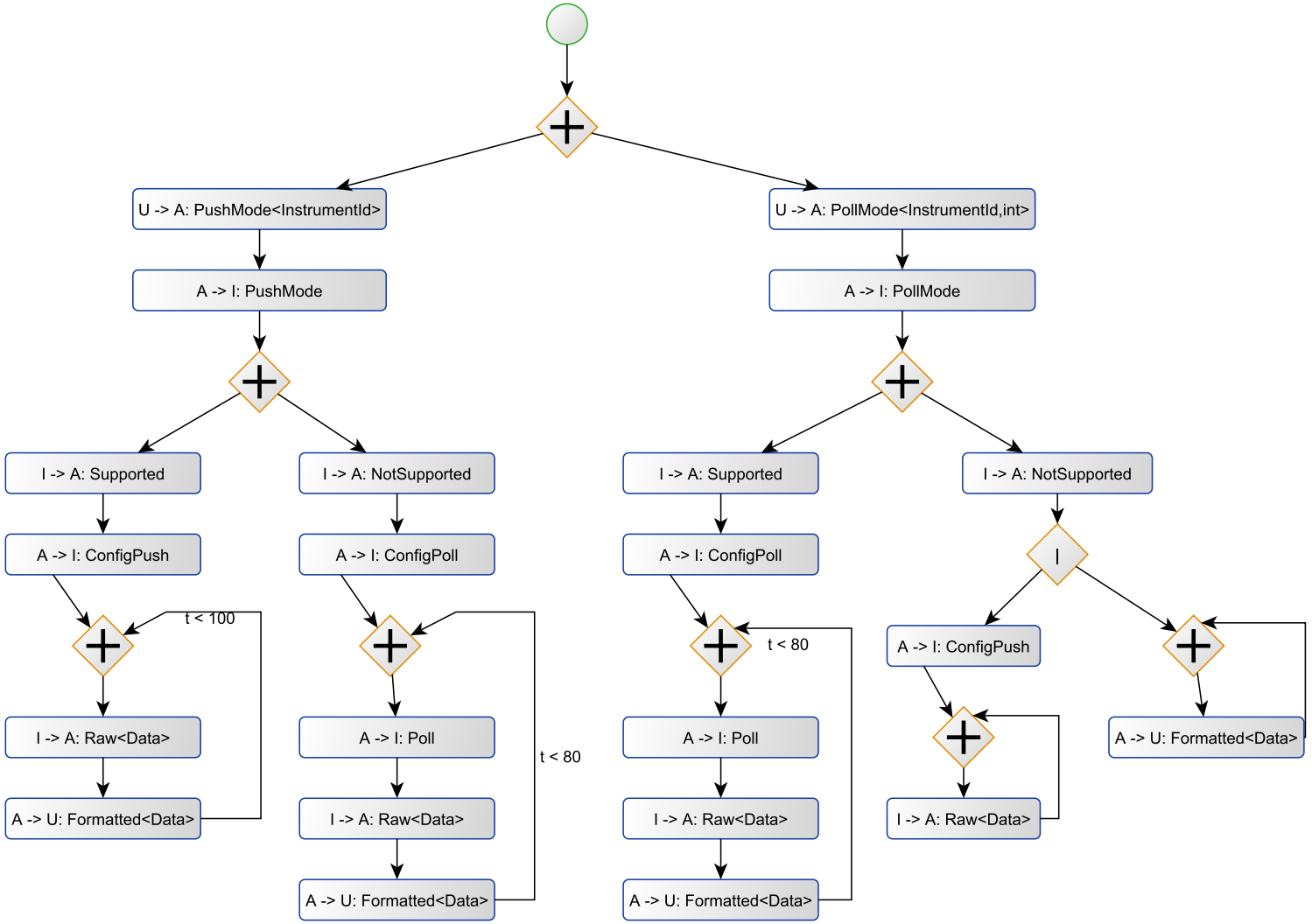

```
choice at U {
  PushMode(InstrumentId) from U to A;
} or {
  PollMode(InstrumentId,int) from U to A;
}
```

$$
\begin{array}{rcl}
x_0 & = & x_{push} + x_{poll} \\
x_{push} & = & \mathrm{U} \rightarrow \mathrm{A}\colon \mathrm{PushMode}\langle\ \mathrm{InstrumentId}\ \rangle; x_{push1} \\
x_{poll} & = & \mathrm{U} \rightarrow \mathrm{A}\colon \mathrm{PollMode}\langle\ \mathrm{InstrumentId,int}\ \rangle; x_{poll1}
\end{array}
$$

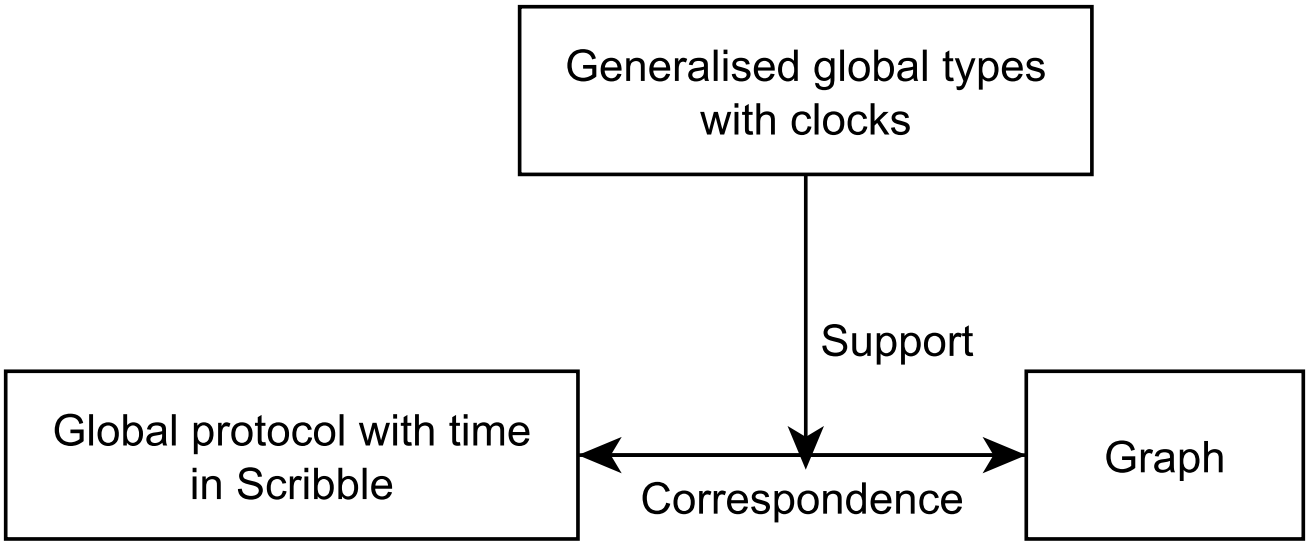
Overview of the project



Graphical Representation with time constraints



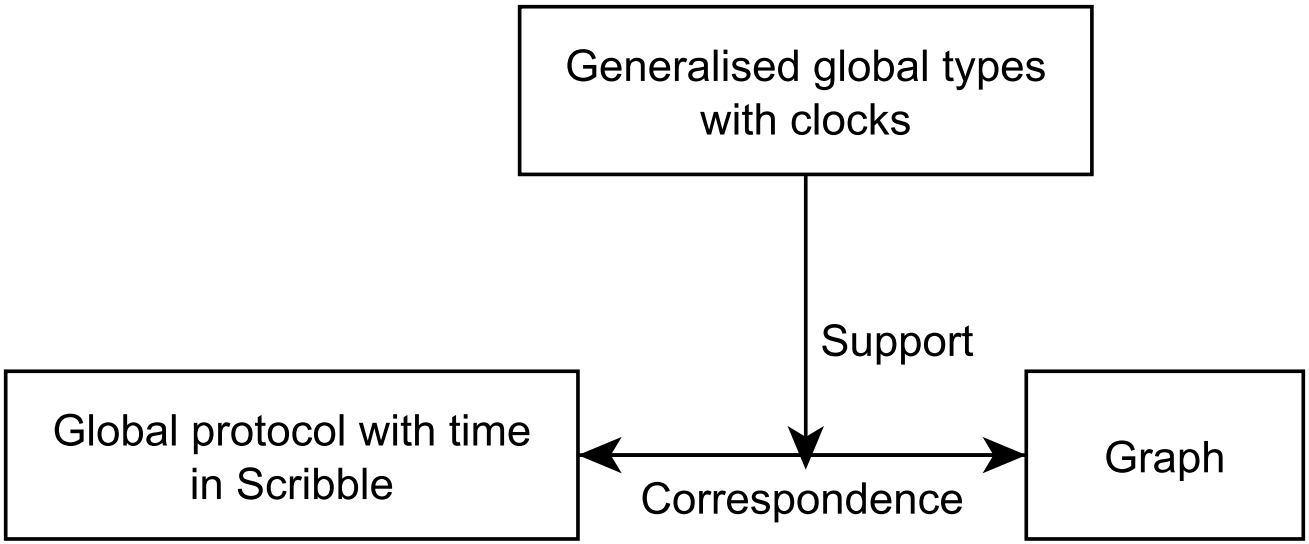
Overview of the project



Contributions:

- Design of the graph
- Extension with clocks
- Implementation of the correspondence

Overview of the project



Contributions:

- Design of the graph
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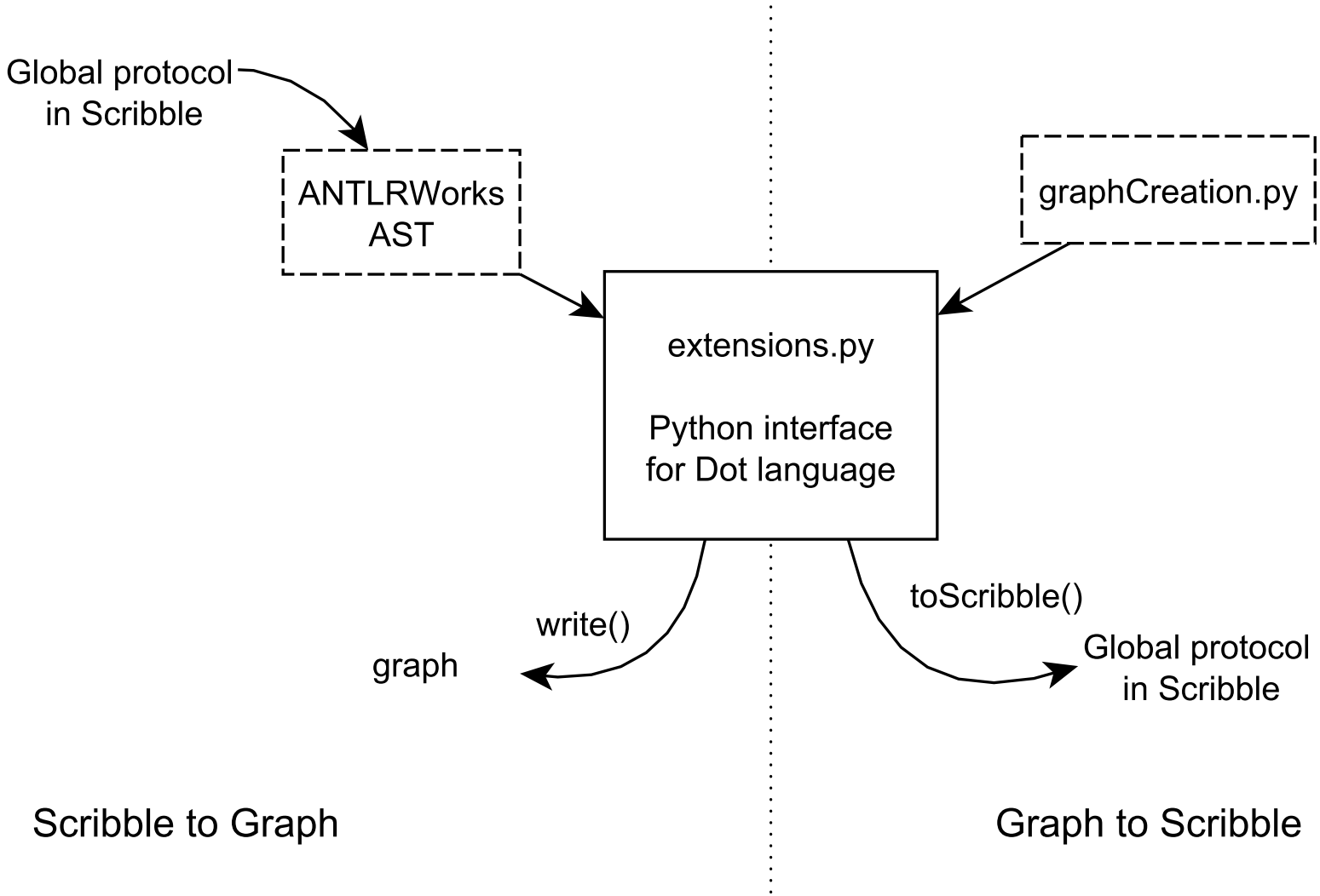
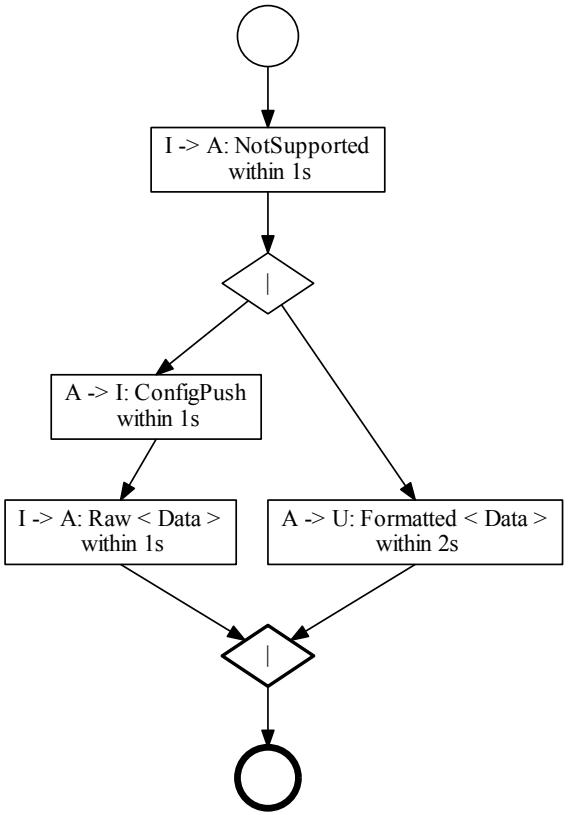
Overview of the implementation



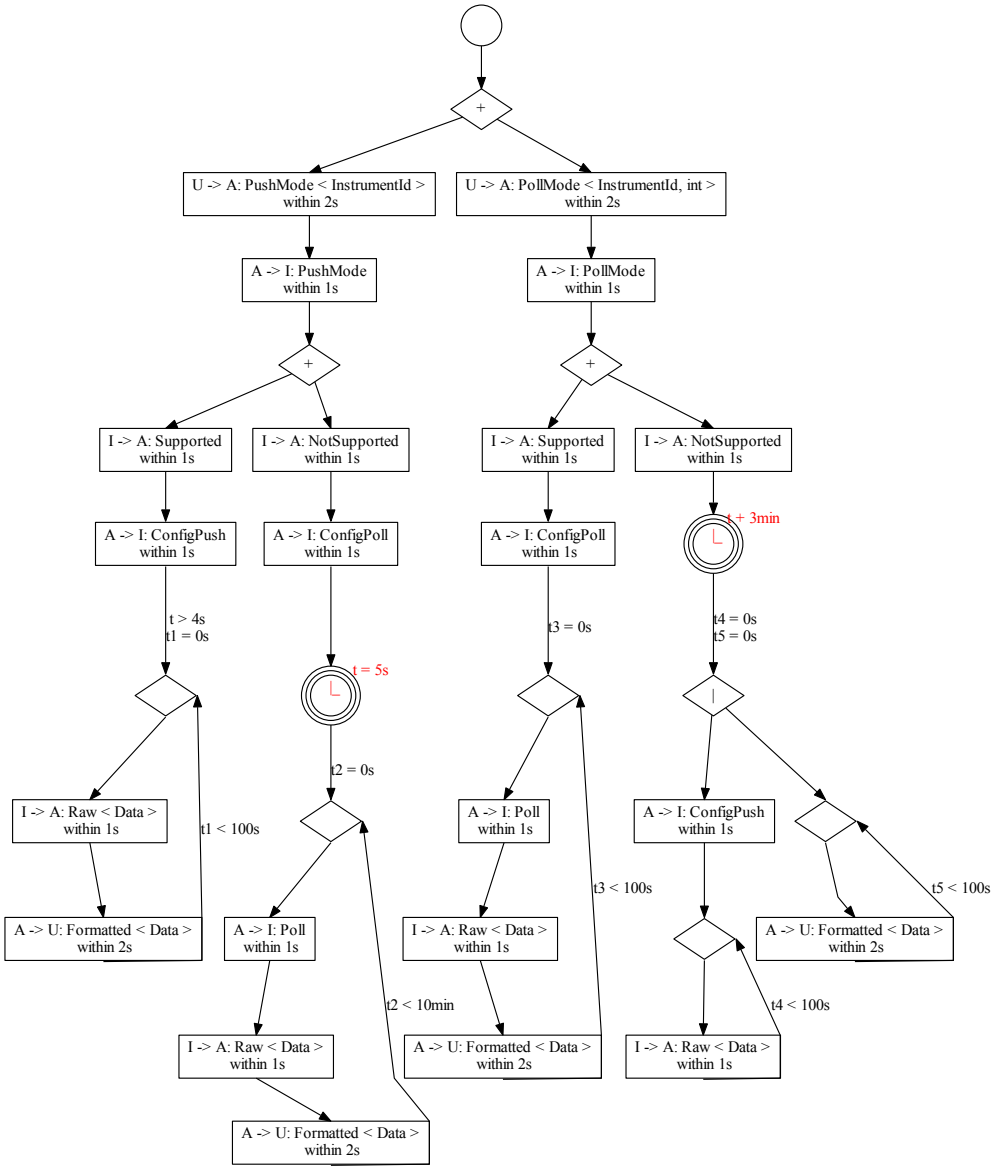
Contents

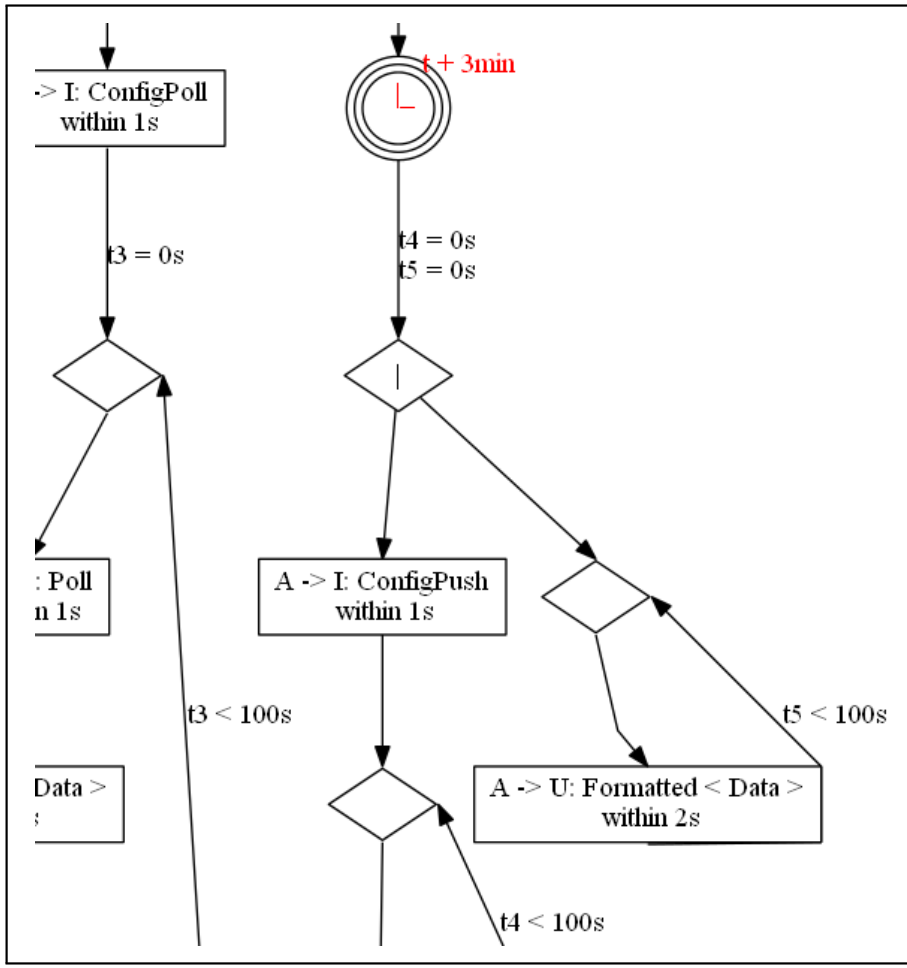
Graphical notations




```
global protocol FirstParallel (role U, role A, role I) {
  NotSupported from I to A within 1s;
  parallel {
      ConfigPush from A to I within 1s;
      Raw(Data) from A to I within 1s;
  } and {
      Formatted(Data) from A to U within 2s;
      }
}
```

Graph with clocks





Zoom

Timed global types

Let C be a set of clocks. C = $\{t, t_1, t_2, ..., t_x, ...\}$

Definition:

$\delta := t \leq v \mid t \geq v \mid \neg\delta \mid \delta_1 \wedge \delta_2 \mid \varepsilon$
where t is a clock in C and v is a constant in $\mathbb{Q}$.

Abbreviations:

t = v	means	t $\leq$ v and t $\geq$ v
t < v	means	$\neg$ t $\geq$ v
t > v	means	$\neg$ t $\leq$ v

G	::=	def $\tilde{G}$ in x	Global type
G	::=	x = p $\rightarrow p'$: $l\langle U\rangle, \lambda_O, \delta_O, \lambda_I, \delta_I; x'$	Labelled messages
	|	x = x' | x"	Fork
	|	x = x' + x"	Choice
	|	x | x' = x"	Join
	|	x + x' = x"	Merge
	|	x = end	End
U	::=	$\langle G\rangle$ | bool | nat | ...	Sorts

Timed local types and projection

T	::=	def $\tilde{T}$ in x	Local type
T	::=	$x = !\langle p, l\langle U\rangle, \lambda, \delta\rangle.x'$	Message sending
	|	$x = ?\langle p, l\langle U\rangle, \lambda, \delta\rangle.x'$	Message receiving
	|	x = x' | x"	Fork
	|	x = x' $\oplus$ x"	Internal choice
	|	x = x' & x"	External choice
	|	x | x' = x"	Join
	|	x + x' = x"	Merge
	|	x = x'	Inaction
	|	x = end	End

Projection algorithm:

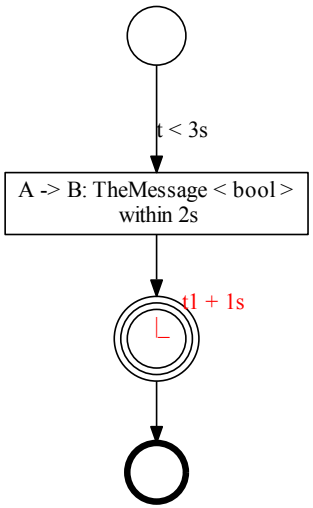
$$\text{def } \tilde{G} \text{ in x} \upharpoonright \text{p} = \text{def } \tilde{G} \upharpoonright_{\tilde{G}} \text{p in x}$$

$x = p \to p' : l\langle U\rangle, \lambda_O, \delta_O, \lambda_I, \delta_I; x' \upharpoonright_{\tilde{G}} p$	=	$x = !\langle p', l\langle U\rangle, \lambda_O, \delta_O\rangle.x'$	
$x = p \to p' : l\langle U\rangle, \lambda_O, \delta_O, \lambda_I, \delta_I; x' \upharpoonright_{\tilde{G}} p'$	=	$x = ?\langle p, l\langle U\rangle, \lambda_I, \delta_I\rangle.x'$	
$x = p \to p' : l\langle U\rangle, \lambda_O, \delta_O, \lambda_I, \delta_I; x' \upharpoonright_{\tilde{G}} p''$	=	x = x'	(p" $\notin$ {p,p'})
x = x' | x" $\upharpoonright_{\tilde{G}}$ p	=	x = x' | x"	
x = x' + x" $\upharpoonright_{\tilde{G}}$ p	=	x = x' $\oplus$ x"	(if p = ASend($\tilde{G}$)(x))
x = x' + x" $\upharpoonright_{\tilde{G}}$ p	=	x = x' & x"	(otherwise)
x | x' = x" $\upharpoonright_{\tilde{G}}$ p	=	x | x' = x"	
x + x' = x" $\upharpoonright_{\tilde{G}}$ p	=	x + x' = x"	
x = end $\upharpoonright_{\tilde{G}}$ p	=	x = end	

Timed global types: Example

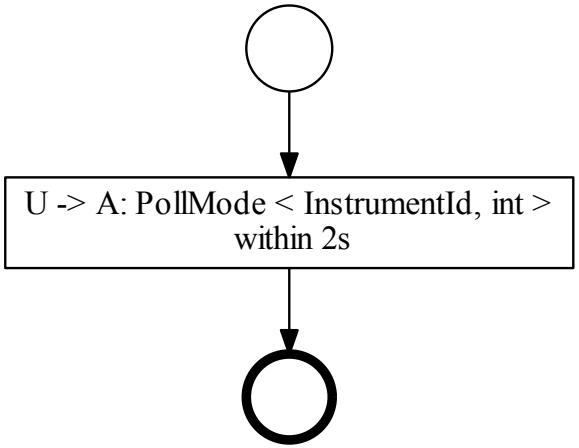
$$x = p \rightarrow p' : l\langle U \rangle, \lambda_O, \delta_O, \lambda_I, \delta_I \text{ ;x'}$$

$$x = A \rightarrow B : \text{TheMessage <bool> }, \{t_x\}, t < 3, \{t_1\}, t_x < 2 \text{ ; x'}$$




```
global protocol Example (role A, role B) {
  t before 3s
  TheMessage(bool) from A to B within 2s;
  wait for t1+ 1s
}
```

The *within* statement




```
global protocol TestMessage (role A, role U) {
  PollMode(InstrumentId,int) from U to A within 2s;
}
```

To support this protocol we define the global type as follows:

$G =$ def x = U $\rightarrow$ A : PollMode <InstrumentId, int> , $\{t_x\}, \varepsilon, \varnothing, t_x < 2$; x'
x' = end
in x

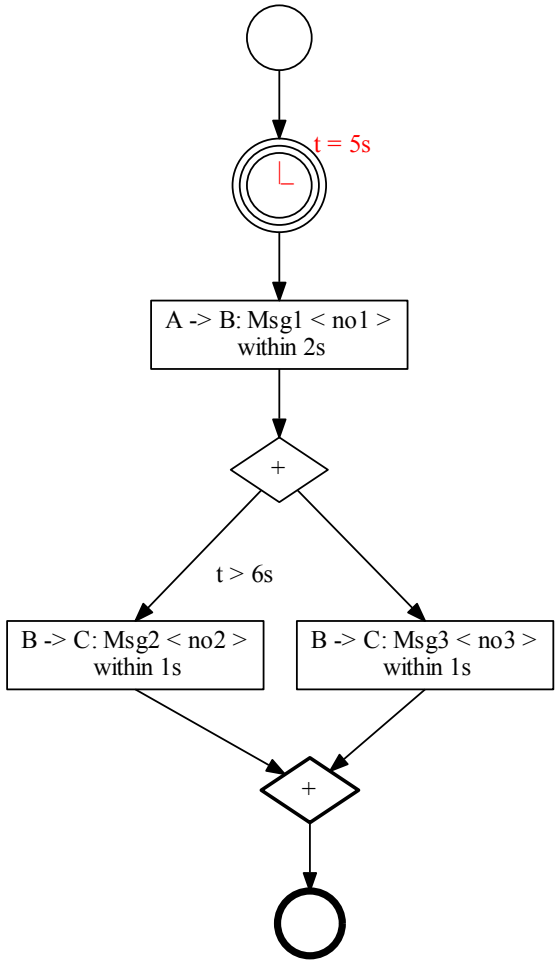
$T_U =$ def x = $!\langle$ A, PollMode <InstrumentId, int> ,$\{t_x\}, \varepsilon\rangle$. x'
x' = end
in x

$T_A =$ def x = $?\langle$ U,PollMode <InstrumentId, int> , $\varnothing, t_x < 2\rangle$. x'
x' = end
in x

Temporal satisfiability for clocks conditions

Temporal satisfiability *If δ, the clock condition of a given transition, is satisfiable at some point, for each constraint δ', appearing in a later transition, it is eventually possible to satisfy δ'.*

```
global protocol TestTemporalSatisfiability (role A, role
B, role C) {
  wait for t is 5s
  Msg1(no1) from A to B within 2s;
     choice at B {
        t after 6s
        Msg2(no2) from B to C within 1s;
     } or {
        Msg3(no3) from B to C within 1s;
     } }
```



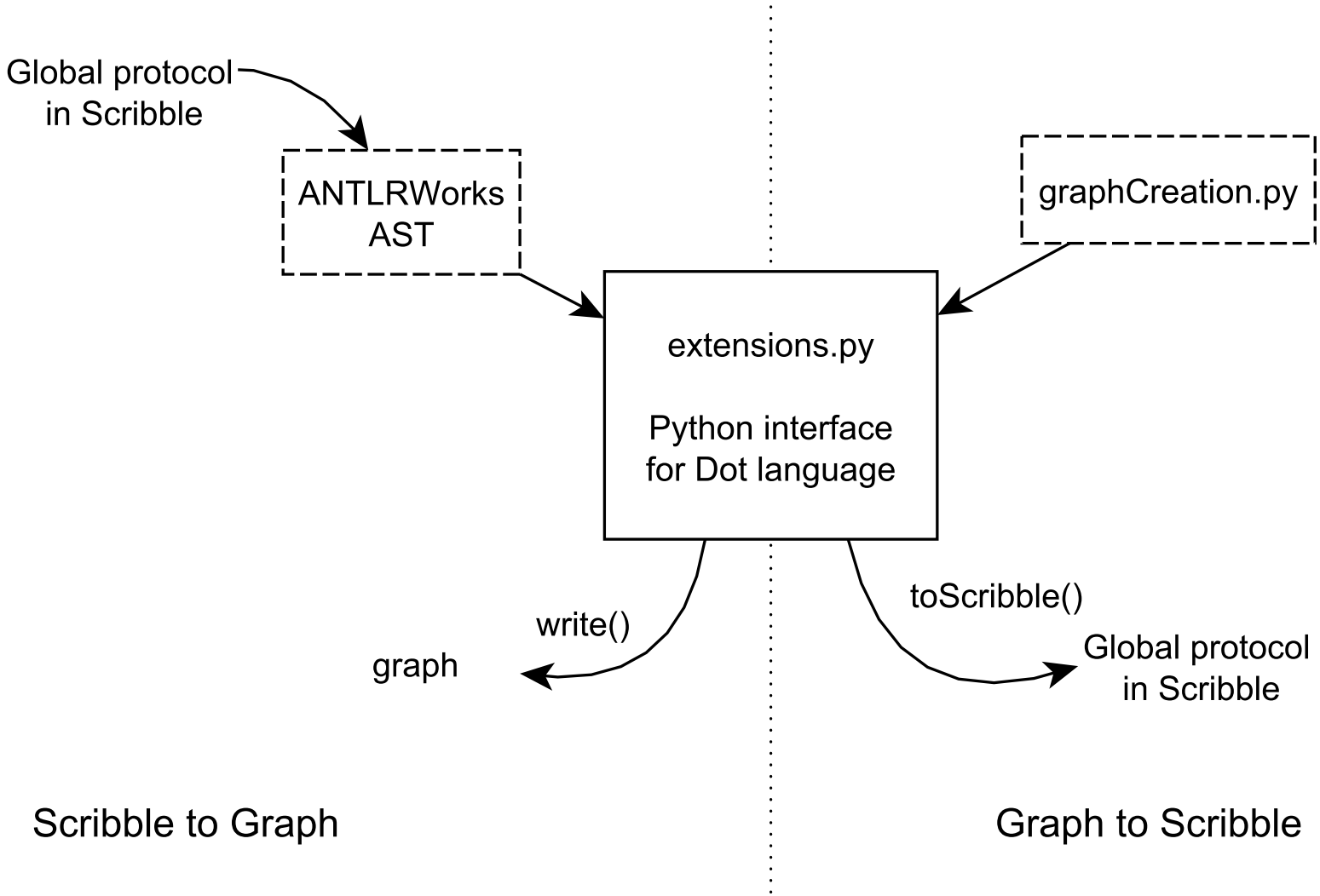
Timed processes

P	::=	def $\tilde{P}$ in X	definition
P	::=	$x(\tilde{x}) = x\langle G, C\rangle.x'(\tilde{e})$	init
	|	$x(\tilde{x}) = x[p](y): x'(\tilde{e})$	accept
	|	$x(\tilde{x}) = x!\langle p; l < e >, \lambda_O, \delta_O\rangle: x'(\tilde{e})$	send
	|	$x(\tilde{x}) = x?\langle p; l(y), \lambda_I, \delta_I\rangle: x'(\tilde{e})$	receive
	|	$x(\tilde{x}) = x'(\tilde{y}) \mid x''(\tilde{z})$	parallel
	|	$x(\tilde{x}) \mid x'(\tilde{y}) = x''(\tilde{z})$	join
	|	$x(\tilde{x}) + x'(\tilde{x}) = x''(\tilde{x})$	merge
	|	$x(\tilde{x}) = x'(\tilde{x})\ \&\ x''(\tilde{x})$	external choice
	|	if e then $x'(\tilde{e}')$ else $x''(\tilde{e}'')$	conditional
	|	$x(\tilde{x}) = 0$	null
	|	$x(\tilde{x}) = (\nu a)x'(a\tilde{x})$	new name
X	::=	$x(\tilde{v})$ | X | X	thread, parallel
	|	$x(\nu a)X$ | 0	restriction, null
e	::=	v | x | $e \wedge \delta$ | $e \wedge e$ | ...	expressions
v	::=	a | s[p] | true | false | ...	values
α, β	::=	$s[p,q]!\langle p; l < e >, \lambda, \delta\rangle$	labels
	|	$s[p,q]?\langle p; l(y), \lambda, \delta\rangle$	
	|	$a\langle G, C\rangle$ | $a\langle p\rangle$ [s] | $\langle \tau, \lambda, \delta\rangle$	

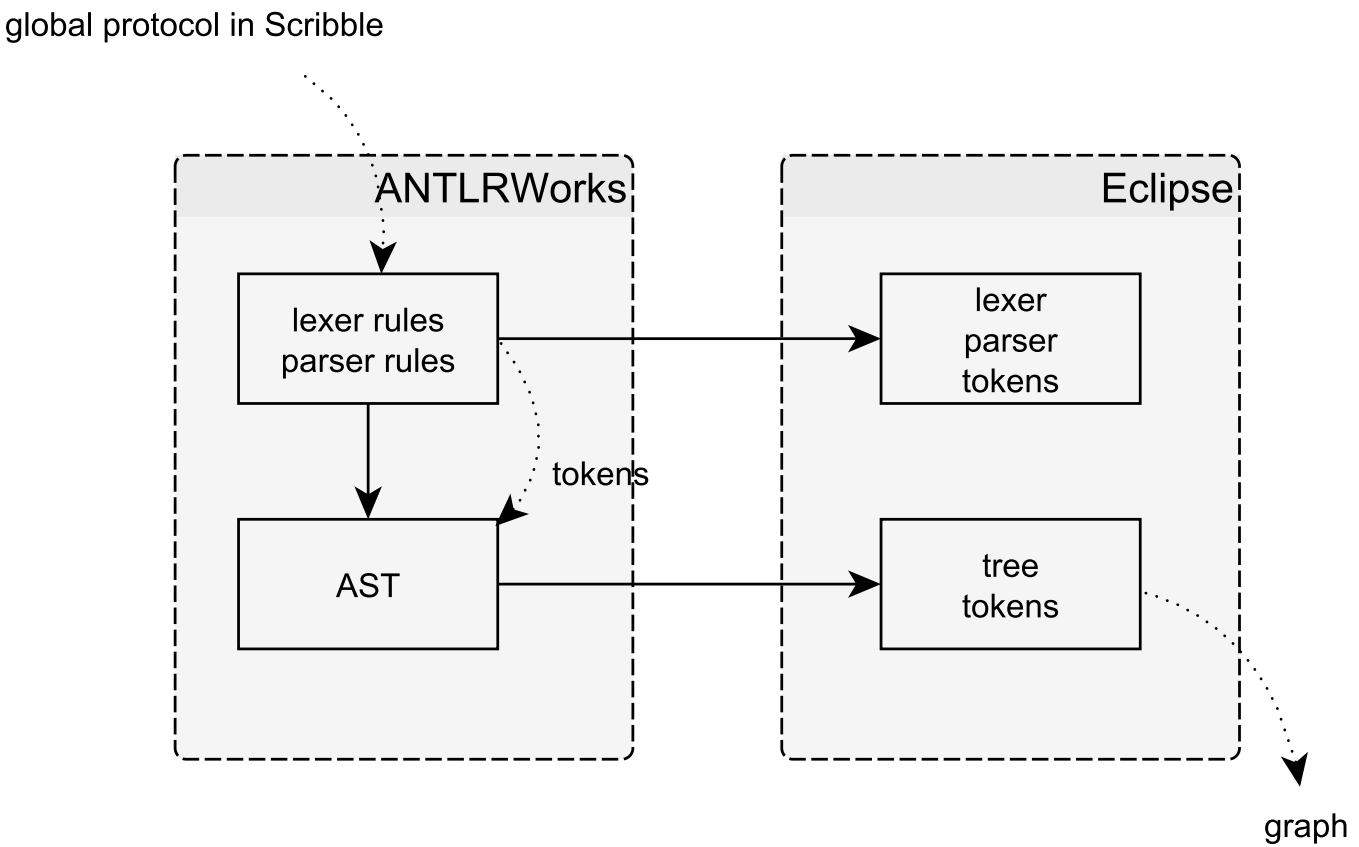
Table: Syntax for timed processes

Contents

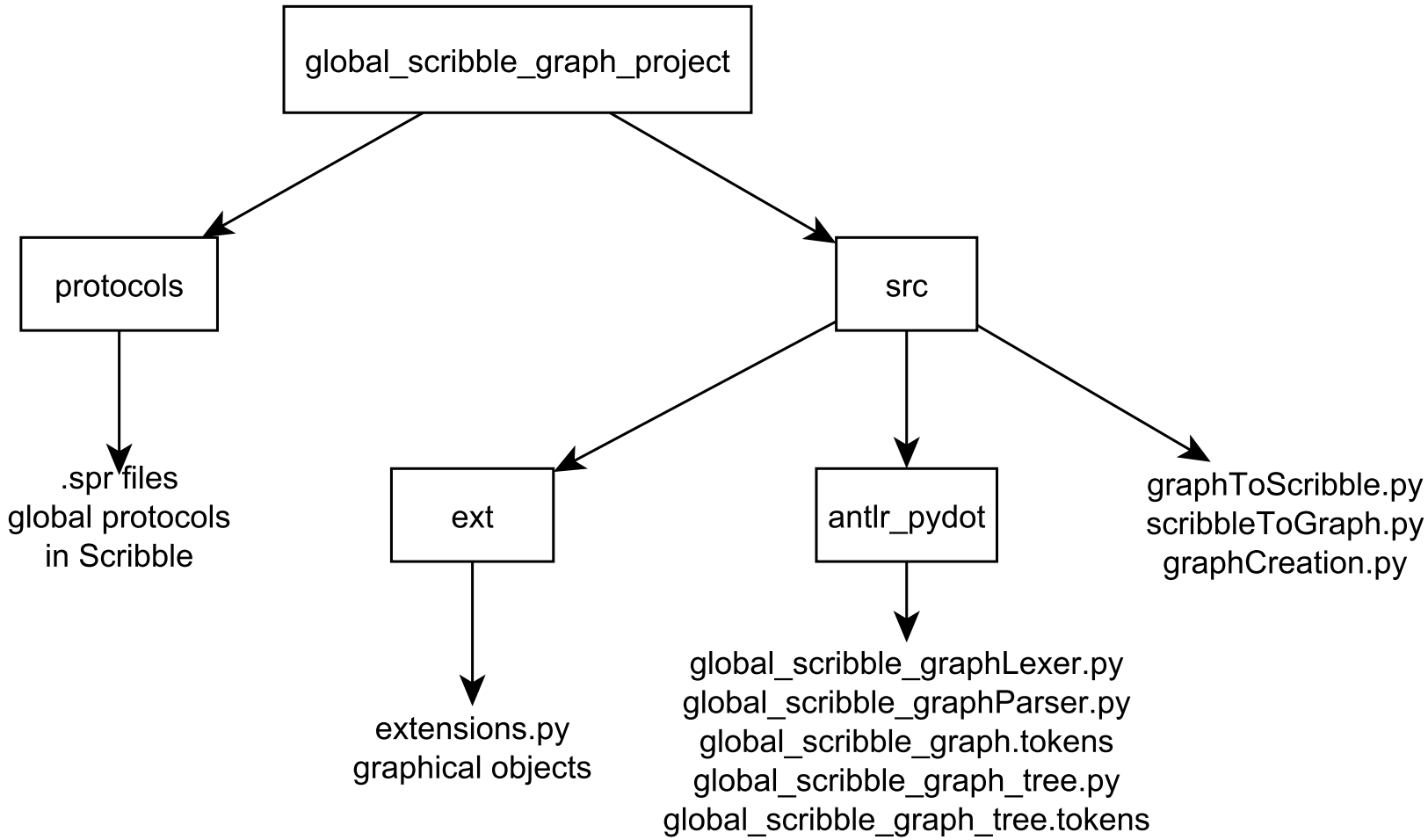
Overview


Global protocol
in Scribble
ANTLRWorks
AST
graphCreation.py
extensions.py
Python interface
for Dot language
write()
graph
toScribble()
Global protocol
in Scribble
Scribble to Graph
Graph to Scribble

Link between ANTLRWorks and Eclipse



Demonstration



Future work

- Well-formedness verification
- Further extensions of Scribble: merge, join, etc.
- Merge with Guillaume's project about the plugin Eclipse
- Proofs of the properties for clocks condition and timed processes

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- Well-formedness verification
- Further extensions of Scribble: merge, join, etc.
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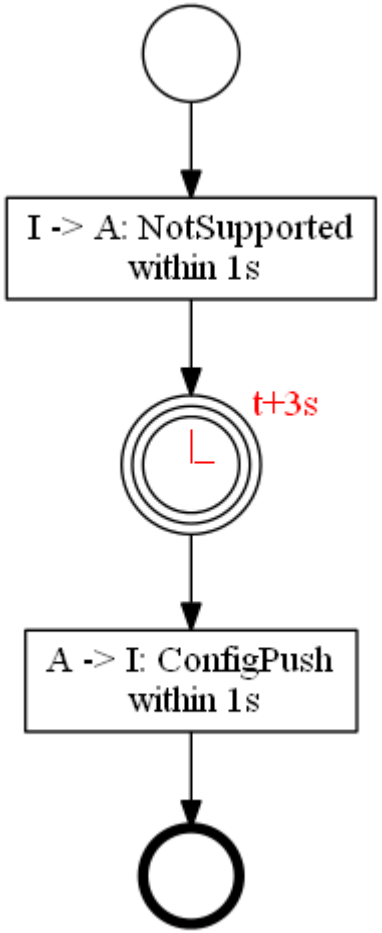
Questions 

Extended Scribble language

global-protocol-body	::=	global-interaction-block
global-interaction-block	::=	{ global-interaction-sequence }
global-interaction-sequence	::=	(global-interaction)*
global-interaction	::=	[time-constraints] message
	|	[time-constraints] choice
	|	[time-constraints] parallel
	|	[time-constraints] recursion
	|	[time-constraints] continue
	|	[time-constraints] delay
message	::=	(message-signature | identifier) `from` role-name `to` role-name `within` time;
delay	::=	`wait for` time-identifier symbol time ;
	|	`wait for` time-identifier `is` time ;
time-constraints	::=	constraint (`and` constraint)*
constraint	::=	time-identifier `after` time
	|	time-identifier `before` time
	|	time-identifier `is` time
time-identifier	::=	identifier
time	::=	(digit)* identifier
symbol	::=	('+' | '*')

Delay

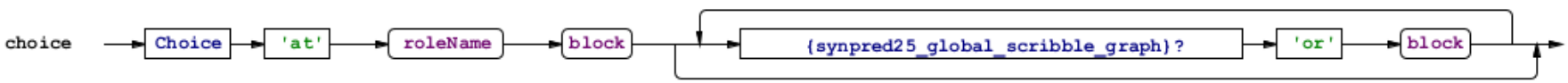



```
global protocol FirstDelay (role A, role I) {
  NotSupported from I to A within 1s;
  wait for t+ 3s
  ConfigPush from A to I within 1s;
}
```

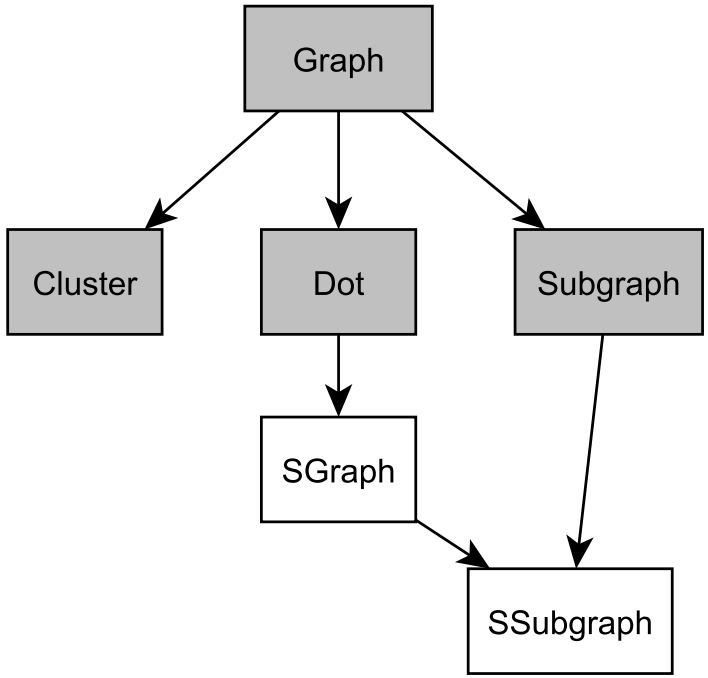
How to write grammar with ANTLR?

```
choice at U {
  PushMode(InstrumentId) from U to A;
} or {
  PollMode(InstrumentId,int) from U to A;
}
```

```
choice
        :       Choice 'at' roleName block ( 'or'  block )* -> ^(Choice roleName block+ );
```

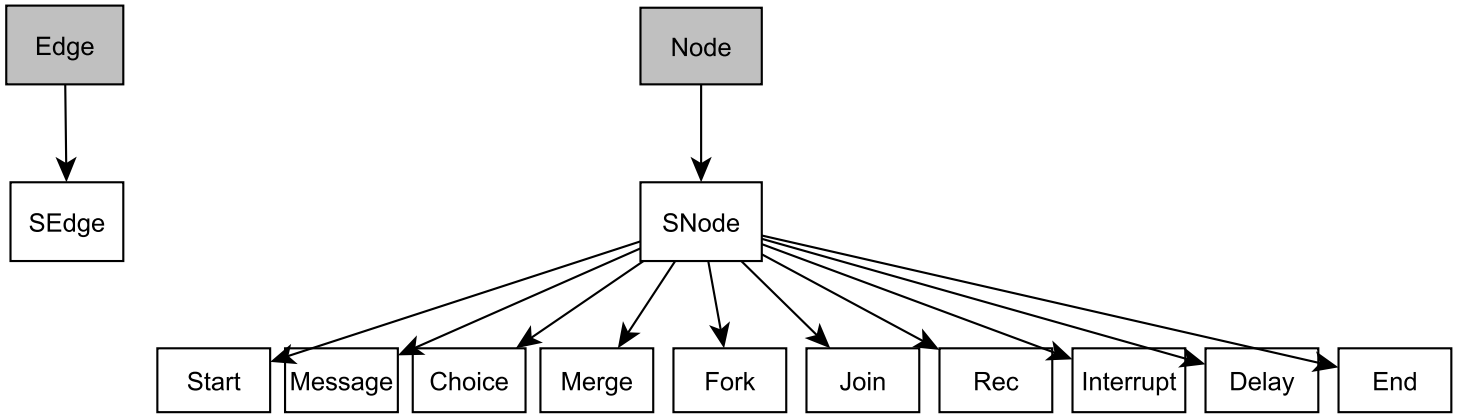


Class diagrams



Pydot

Extensions



Pydot

Extensions