#### CMSC 25025 / STAT 37601

# **Machine Learning and Large Scale Data Analysis**

### Assignment 3

Due: Thursday, April 25, 2013

This assignment consists of three problems. Al of them are "pencil and paper" problems. They are intended to be short, and to help you with the programming project that you are currently working on.

### 1. Feeling the Heat (30 points)

The Laplacian is the differential operator given by the sum of the second derivatives

$$\Delta f(x) = \frac{\partial^2 f(x)}{\partial x_1^2} + \frac{\partial^2 f(x)}{\partial x_2^2} + \cdots + \frac{\partial^2 f(x)}{\partial x_d^2}$$

for a function  $f: \mathbb{R}^d \longrightarrow \mathbb{R}$ . One of the many uses of the Laplacian is to model the diffusion of heat through a substance.

Recall that for a graph on n nodes, we suppose we have a weight  $w_{ij}$  assigned to edge (i, j). Assume the weights are symmetric, so that  $w_{ij} = w_{ji}$ , and let  $W = [w_{ij}]$  be the  $n \times n$  weight matrix. The graph Laplacian is the  $n \times n$  matrix defined by

$$\Delta = D - W. \tag{1}$$

Show that

$$\sum_{ij} w_{ij} (f_i - f_j)^2 = f^T \Delta f$$

for any vector  $f = (f_1, f_2, \dots, f_d)$ .

#### 2. Working in Harmony (30 points)

Now recall our harmonic function semi-supervised learning algorithm. We have some points  $x_i$  labeled  $y_i = 1$ , some points labeled  $y_i = 0$ , and many points without any label. We form a weighted graph with graph Laplacian  $\Delta = D - W$ . We then estimate a value  $f_j$  for each node j by minimizing

$$\sum_{ij} w_{ij} (f_i - f_j)^2 = f^T \Delta f$$

subject to  $f_i = y_i$  if  $x_i$  is labeled  $y_i$ . Show that the solution is *harmonic*; namely,

$$f_i = \frac{\sum_j w_{ij} f_j}{\sum_j w_{ij}}$$

for each node i.

## 3. Loving Your Neighbors (40 points)

neighbors of  $x_i$ .

Suppose we have n points  $\{x_i\}_{i=1}^n$  and a distance  $d_{ij} = d_{ji} = \operatorname{dist}(x_i, x_j)$  between each pair. The k nearest neighbors of a point  $x_i$  are the k points closest to  $x_i$ , in terms of the distance  $d_{ij}$ . The k-nearest neighbor graph  $G_k$  has n nodes, one for each point  $x_i$ , and an edge between i and j if either  $x_j$  is one of the k nearest neighbors of  $x_i$ , or if  $x_i$  is one of the k nearest

Derive an algorithm that computes the smallest k such that  $G_k$  is a connected graph. Sketch the algorithm in pseudo-code. What is the complexity of the algorithm?