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Abstract

Final polynomials have been used to prove non-representability for oriented matroids, i.e. to decide whether geometric embeddings of combinatorial structures exist. They received more attention when Dress and Sturmfels, independently, pointed out that non-representable oriented matroids always possess a final polynomial as a consequence of an appropriate real version of Hilbert's Nullstellensatz. We discuss the more difficult problem of determining such final polynomials algorithmically. We introduce the notion of bi-quadratic final polynomials, and we show that finding them is equivalent to solving an LP-problem. We apply a new theorem about symmetric oriented matroids to a series of cases of geometrical interest.

1. Introduction

The method of using final polynomials, as a tool for proving non-representability of oriented matroids or for deciding geometric embeddings of combinatorial structures, was introduced when the first author studied 3-spheres and their polytopality. Later on the notion final polynomial was used, and a variety of applications of this method of proof has grown. Sturmfels and Dress have independently pointed out that through the use of a real version of Hilbert's Nullstellensatz, a final polynomial exists for every non-representable oriented matroid. This gave new impetus to the more difficult task of finding such polynomials algorithmically. A detailed description of the final polynomial method was given in [5]. Any further progress in this direction would be a great advance, and this article is devoted precisely to this question.

Oriented matroids (and matroids) are general discrete structures which can perhaps be compared with graphs or lattices. Their general structure is an appropriate tool for tackling mathematical problems when linear dependencies and independencies play a substantial role. According to these very general properties, the rather young theory of oriented matroids has been used and applied in many different fields of mathematics.

A fundamental problem in the theory of oriented matroids is deciding representability. Final polynomials were suggested as a general tool for proving non-representability. We emphasize the

Proposed running head: Finding final polynomials.

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advantage of this method: once a final polynomial is known, the calculations for finding it are no longer needed in the proof. So far, only a rather small number of final polynomials have been found. In every case, these polynomials were determined by using rather long computer-algebra-like calculations, and sometimes even greater effort was needed in solving a specific problem, e.g. in [3]. This more difficult problem of finding a particular final polynomial for specific applications will be discussed in Section 2. We provide a polynomial algorithmic method, using essentially a suitable variant of linear programming for finding classes of final polynomials. It is very likely that there are final polynomials which we cannot find in this way. On the other hand, the set we do get at least includes all final polynomials known to us so far. Thus our method, deciding first a suitable LP-problem, can at least be applied as a good pre-processing.

Section 3 is devoted to symmetrical chirotopes. For symmetric d-chirotopes it is natural to decide whether there is a symmetrical realization with respect to a given finite subgroup of the orthogonal group O(d). Our main result in Section 3 asserts that, once the non-realizability of our chirotope along with (geometric) symmetry assumptions has been achieved by our method (!), we can solve the general case as well, i.e., the symmetry assumption can be dropped. This is a substantial shortcut in attaining the final polynomial, algorithmically.

In Section 4 we apply our method of finding new non-embeddability proofs to a series of examples, for instance, we apply it to an example in the class of *star-chirotopes*, cf. [7], which were key-examples in establishing the Non-Steinitz-theorem by Sturmfels, [13] and in which it seemed to be difficult to find corresponding final polynomials. Still another application of our method will be given in a forthcoming paper [4], in which we determine all (simplicial) embeddings of a certain class of matroid manifolds.

In conclusion, we would like to mention that our approach to the problem has been coordinate free. This new development has also been presented in the collection of papers written by Crapo, Havel, Sturmfels, White, and Whiteley in Minnesota in 1987, [8].

2. Finding final polynomials through linear programming

It is generally considered that the algorithmic search for final polynomials is very difficult. Sometimes the effort seems to be beyond our reach when the number of points increase. In this section, we provide an algorithmic method for finding final polynomials in a reasonable length of time. All final polynomials known to us so far, can be achieved by our method. In describing the details, we restrict ourselves to the simplicial case, thus simplifying the main idea.

Let χ be a given simplicial d-chirotope of n points $\chi: \Lambda(n,d) \to \{-1,1\}$. Here $\Lambda(n,d)$ denotes the set of all ordered d-tupels of n elements $\lambda = (\lambda_1, \ldots, \lambda_d)$ with $\lambda_1 \leq \ldots \leq \lambda_d$. If neccessary, we refer to the tupels λ of $\Lambda(n,d)$ as

$$\lambda^1 = (1, 2, ..., d), \quad \lambda^2 = (1, 2, ..., d-1, d+1), \quad ..., \quad \lambda^{\binom{n}{2}} = (n-d-1, ..., n).$$

The map χ is considered to be extended alternatingly onto all d-tupels $\lambda \in \{1, \ldots, n\}^d$. To every $\lambda = (\lambda_1, \lambda_2, \ldots, \lambda_d) \in \{1, \ldots, n\}^d$, we assign a formal variable $[\lambda_1, \lambda_2, \ldots, \lambda_d]$, called bracket, which in case of a realizable chirotope with points $x_1, \ldots, x_n \in R^d, x_i = (x_i^1, x_i^2, \ldots, x_i^d)$ will be equal to the determinant

$$[\lambda_1, \lambda_2, \dots, \lambda_d] = \det egin{pmatrix} x_{\lambda_1}^1 & \dots & x_{\lambda_1}^d \\ \vdots & & \vdots \\ x_{\lambda_d}^1 & \dots & x_{\lambda_d}^d \end{pmatrix}.$$

But in general, we consider these brackets as only being formal variables. We abbreviate $\tau := \tau_1, \ldots, \tau_{d-2}$. The following term will be called Graßmann-Plücker polynomial,

$$\{\tau \mid \lambda_1, \lambda_2, \lambda_3, \lambda_4\} := [\tau, \lambda_1, \lambda_2][\tau, \lambda_3, \lambda_4] - [\tau, \lambda_1, \lambda_3][\tau, \lambda_2, \lambda_4] + [\tau, \lambda_1, \lambda_4][\tau, \lambda_2, \lambda_3].$$

Again, in the case of the brackets being determinants, this Graßmann-Plücker polynomial is equal to zero, since in this case we have precisely a Graßmann-Plücker syzygy.

The definition of a chirotope requires one of the following relations

$$\{-1,+1\} \subset P$$
 or $\{0\} = P$ with

$$P := \{\chi((\tau, \lambda_1, \lambda_2))\chi((\tau, \lambda_3, \lambda_4)), -\chi((\tau, \lambda_1, \lambda_3))\chi((\tau, \lambda_2, \lambda_4)), \chi((\tau, \lambda_1, \lambda_4))\chi((\tau, \lambda_2, \lambda_3))\}.$$

This is a neccessary condition for χ being realizable, otherwise no choice of coordinates for the points x_1, \ldots, x_n would fulfill both, the requirement

$$\chi(\lambda) = \chi((\lambda_1, \dots, \lambda_d)) = \text{sign } [\lambda_1, \dots, \lambda_d]$$

and the above single Graßmann-Plücker syzygy.

The set of vectors $\alpha = (\alpha_1, \ldots, \alpha_{\binom{n}{2}}) \in R^{\binom{n}{2}}$ with

$$\alpha_1 = [1, 2, ..., d], \quad \alpha_2 = [1, 2, ..., d-1, d+1], \quad ..., \quad \alpha_{\binom{n}{d}} = [n-d-1, ..., n],$$

such that all Graßmann-Plücker syzygies are fulfilled, define the Grassmannian $G(n,d) \subset R^{\binom{n}{2}}$. The chirotope χ is realizable, if and only if there is a point

$$\alpha = (\alpha_1, \ldots, \alpha_{\binom{d}{n}}) \in G(n, d)$$

in the Grassmannian G(n,d), which also lies in the cone $C(\chi) \subset R^{\binom{n}{2}}$, defined as

$$C(\chi) := \{\alpha = (\alpha_1, \ldots, \alpha_{\binom{d}{n}}) \mid \operatorname{sign}(\alpha_j) = \chi(\lambda^j), \ 1 \leq j \leq \binom{n}{d}\}.$$

Remember that the change in the above Graßmann-Plücker polynomial, caused by a permutation of the points on the right hand side of the vertical line in

$$\{\tau \mid \lambda_1, \lambda_2, \lambda_3, \lambda_4\},\$$

keeps the relation property (the chirotope condition) mentioned above invariant. The notation was accordingly chosen in order to avoid multiple checkings of this condition. The polynomial remains equal up to choosing the actual brackets

$$[\lambda_1, \ldots, \lambda_d]$$
 or $\operatorname{sign} \pi [\lambda_{\pi(1)}, \ldots, \lambda_{\pi(d)}]$

for a permutation $\pi \in S_d$, where S_d denotes the symmetric group, or it is equal up to ordering the bracket products. By a suitable permutation of the points $\lambda_1, \lambda_2, \lambda_3, \lambda_4$, we can attain that all three bracket products are positive

$$\underbrace{[\tau,\lambda_1,\lambda_2][\tau,\lambda_3,\lambda_4]}_{>0} - \underbrace{[\tau,\lambda_1,\lambda_3][\tau,\lambda_2,\lambda_4]}_{>0} + \underbrace{[\tau,\lambda_1,\lambda_4][\tau,\lambda_2,\lambda_3]}_{>0}.$$

In this case, we say, we have normalized our Grasmann-Plücker polynomial. In case χ is realizable and in the normalized case, we can deduce that the inequality

$$[\tau, \lambda_1, \lambda_2][\tau, \lambda_3, \lambda_4] < [\tau, \lambda_1, \lambda_3][\tau, \lambda_2, \lambda_4]$$

must hold. In what follows, it will be useful to have a special notation for these formal inequalities, arising from normalized Graßmann-Plücker polynomials.

Definition 2.1. Let χ be a simplicial chirotope and let

$$\{\tau \mid \lambda_1, \lambda_2, \lambda_3, \lambda_4\} := [\tau, \lambda_1, \lambda_2][\tau, \lambda_3, \lambda_4] - [\tau, \lambda_1, \lambda_3][\tau, \lambda_2, \lambda_4] + [\tau, \lambda_1, \lambda_4][\tau, \lambda_2, \lambda_3]$$

be a Graßmann-Plücker polynomial such that

$$\chi((\tau,\lambda_1,\lambda_2))\cdot\chi((\tau,\lambda_3,\lambda_4))=\chi((\tau,\lambda_1,\lambda_3))\cdot\chi((\tau,\lambda_2,\lambda_4))=\chi((\tau,\lambda_1,\lambda_4))\cdot\chi((\tau,\lambda_2,\lambda_3))=1,$$

then we call the 4-tupel B of formal variables,

$$B = \langle [\tau, \lambda_1, \lambda_2], [\tau, \lambda_3, \lambda_4] \mid [\tau, \lambda_1, \lambda_3], [\tau, \lambda_2, \lambda_4] \rangle,$$

a bi-quadratic inequality of χ . The set of all bi-quadratic inequalities of χ will be denoted by \mathcal{B}_{χ} .

Remark 2.1. Whenever χ is realizable, we have

$$A \cdot B < C \cdot D$$
 for all $\langle A, B \mid C, D \rangle \in \mathcal{B}_{\chi}$.

In order to prepare a definition for bi-quadratic final polynomials, we must identify those brackets, which in the realizable case would change according to the alternating determinant rules, i.e., for any permutation $\pi \in S_d$ we set

$$[\lambda_1, \lambda_2, \ldots, \lambda_d] - \operatorname{sign} \pi \cdot [\lambda_{\pi(1)}, \lambda_{\pi(2)}, \ldots, \lambda_{\pi(d)}] = 0.$$

In other words, if R is the integer polynomial ring generated by the formal brackets $\{[\lambda] \mid \lambda \in \{1, ..., n\}^d\}$ and if I is the ideal generated by polynomials of the form

$$[\lambda_1, \lambda_2, \ldots, \lambda_d] = \operatorname{sign} \pi \cdot [\lambda_{\pi(1)}, \lambda_{\pi(2)}, \ldots, \lambda_{\pi(d)}],$$

we calculate within the quotient ring R/I. Note that R/I is similar, but not identical, to the bracket ring introduced by White, see [8], where I contains also the Graßmann-Plücker polynomials. Now we are ready to define bi-quadratic final polynomials as follows

Definition 2.2. A simplicial chirotope χ admits a bi-quadratic final polynomial, whenever, there is a collection of bi-quadratic inequalities

$$\langle A_i, B_i \mid C_i, D_i \rangle \in \mathcal{B}_{\gamma}; \quad 1 \leq i \leq k$$

such that the following equality holds within the ring R/I.

$$\prod_{i=1}^k A_i \cdot B_i \stackrel{I}{=} \prod_{i=1}^k C_i \cdot D_i.$$

Here $a \stackrel{J}{=} b$ expresses that the ring variables a and b are equal modulo the set J.

The above Definition 2.2. allows the following claim.

Lemma 2.1. If χ admits a bi-quadratic final polynomial, then χ is not realizable.

Proof

 χ admits a bi-quadratic final polynomial, i.e., there is a collection of bi-quadratic inequalities

$$\langle A_i, B_i \mid C_i, D_i \rangle \in \mathcal{B}_{\chi}; \quad 1 \leq i \leq k$$

such that

$$\prod_{i=1}^k A_i \cdot B_i \stackrel{\underline{I}}{=} \prod_{i=1}^k C_i \cdot D_i.$$

In assuming that χ is realizable, we have by Remark 1:

 $A_i \cdot B_i < C_i \cdot D_i$ for all $i \in \{1, ..., k\}$, and since both sides in the above equation are by definition positive, we also have

$$\prod_{i=1}^k A_i \cdot B_i < \prod_{i=1}^k C_i \cdot D_i.$$

In using

$$[\lambda_1, \lambda_2, \ldots, \lambda_d] - \operatorname{sign} \pi \cdot [\lambda_{\pi(1)}, \lambda_{\pi(2)}, \ldots, \lambda_{\pi(d)}] = 0,$$

whenever neccessary, we arrive at a contradiction.

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Remark 2.2. Now, it is a straight-forward task of translating a bi-quadratic final polynomial into an ordinary final polynomial as introduced in [5].

Of course, it seems as if one looses quite a bit of information in searching for bi-quadratic final polynomials rather then looking for final polynomials in general. But there are two reasons for doing this.

- 1. The solvability of the remaining system of inequalities can now be seen to be a LP-problem, yielding a polynomial algorithm by the ellipsoid method, compare [10].
- 2. Again all known examples of non-representable oriented matroids, where non-representability was proven by means of final polynomials, can be treated automatically and more easily than before.

And above all, this might be the only accessible method of solving a more difficult problem under consideration. In any case, we suggest solving this LP-problem as a pre-processing method. How do we get the LP-problem?

We consider the set of all formal inequalities \mathcal{B}_{χ} . We replace each variable $X = [\lambda_1, \ldots, \lambda_d]$ by their formal absolute value $X \cdot \chi((\lambda_1, \ldots, \lambda_d))$. It is clear how one gets new formal inequalities $\langle A, B \mid C, D \rangle$ for the formal absolute values. Assuming realizability, we can take the logarithm on both sides, and we get a linear problem (*) with integer coefficients and strict inequalities for the new variables $Y := log(X \cdot \chi((\lambda_1, \ldots, \lambda_d)))$ chosen appropriately.

In principle, it is now possible to use the special structure of this inequality system when solving this system of inequalities. Consider the inequalities with positive coefficients

$$\sum \ldots + \alpha_l Y + \ldots < \sum \ldots + \ldots$$

or

$$\sum \ldots + \ldots < \sum \ldots + \beta_{\tau} Y + \ldots ,$$

and multiply those inequalities, which contain the variable Y, with suitable factors in order to get mY, where m denotes the smallest common multiple of all (positive) factors α_l , β_r of Y. We rewrite all inequalities in the form

$$\sum \ldots < mY < \sum \ldots$$

The decisive inductive step of our process for solving all variables is replacing this system of inequalities by all pairs of inequalities, and comparing all possible parts, left and right that one has solved for Y together with the remaining system of inequalities. Now, either one gets a solution this way, or one finally arrives at a contradiction 0 < 0, which then can be traced back in order to find the bi-quadratic final polynomial. We formulate this result as follows.

Theorem 1.

A chirotope χ admits a bi-quadratic final polynomial, if and only if the dual of the above LP-problem (*) is admissable (LP-Phase I). Moreover, the solution of the LP-problem can be used for constructing the bi-quadratic final polynomial.

Remark 3. We don't recommend the procedure used above for algorithmic purposes in general. For practical solutions we applied the MINOS-software system. When judging the above theorem it is the algorithmic point of view which one has to keep in mind, We were surprised ourselves about the practical significance as will be shown in concrete applications in Section 4.

3. Final polynomials with symmetries

This section is devoted to symmetric chirotopes and corresponding final polynomials with respect to a given symmetry. For symmetric d-chirotopes it is natural, not only to look for a realization at all, but also to decide whether there is a symmetrical realization with respect to a given finite subgroup of the orthogonal group O(d).

In a straight forward way, this leads to the tool of bi-quadratic final polynomials with respect to a symmetry group for disproving symmetrical embeddings. Our main result in this section asserts that once the non-realizability of our chirotope according to (geometric) symmetry assumptions has been established, by using this tool (!), we can then derive a final polynomial even for the general case, by dropping the (geometric) symmetry assumption. The result being that our chirotope is not realizable at all. The advantage of our theorem is immediate as in these cases, it is an essential shortcut in finding a final polynomial, algorithmically.

In order to formulate this theorem, we first have to fix our notation. We abbreviate for any $\lambda = (\lambda_1, \lambda_2, \dots, \lambda_d) \in E^d$ and any $\sigma \in S_n$:

$$\sigma\lambda = (\sigma(\lambda_1), \sigma(\lambda_2), \ldots, \sigma(\lambda_d)).$$

Let χ be a simplicial d-chirotope. An element $\sigma \in S_n$ of the permutation group S_n is called a rotation of χ if

$$\chi(\lambda) = \chi(\sigma\lambda)$$
 for all $\lambda = (\lambda_1, \ldots, \lambda_d) \in \{1, \ldots, n\}^d$,

denoted by $\sigma\chi = \chi$. Similarly, an element $\sigma \in S_n$ is called a reflection of χ if

$$\chi(\lambda) = -\chi(\sigma\lambda)$$
 for all $\lambda = (\lambda_1, \ldots, \lambda_d) \in \{1, \ldots, n\}^d$.

In this case, we write $\sigma \chi = -\chi$. We call

$$R_{\chi} = \{ \sigma \in S_n \mid \sigma \chi = \chi \}$$

the set of rotations of χ and

$$M_{\chi} = \{ \sigma \in S_n \mid \sigma \chi = -\chi \}$$

the set of reflections of χ . The union of these sets forms the group of automorphisms $\operatorname{Aut}(\chi)$ of χ denoted by $G_{\chi} := R_{\chi} \cup M_{\chi} = \operatorname{Aut}(\chi)$.

The group G_{χ} (resp. any subgroup of G_{χ}) acts on the polynomial ring R as follows. For any bracket and any $\sigma \in S_n$, we define

$$\sigma * [\lambda] = \begin{cases} [\sigma \lambda] & \text{if } \sigma \in R_{\chi}, \\ -[\sigma \lambda] & \text{if } \sigma \in M_{\chi}. \end{cases}$$

For any polynomial $P([\lambda], [\mu], \ldots)$ in the bracket variables $[\lambda], [\mu], \ldots$, we define

$$\sigma * P([\lambda], [\mu], \ldots) = P(\sigma * [\lambda], \sigma * [\mu], \ldots).$$

The action of the group can be extended to any quotient R/J, whenever the ideal J remains fixed under the group, in particular, we define for the above polynomial ring R/I: for any $P+I \in R/I$; $P \in R$ and any $\sigma \in G_{\chi}$:

$$\sigma * (P+I) := (\sigma * P) + I.$$

Our next claim is that according to

$$\sigma * \langle A, B | C, D \rangle = \langle \sigma * A, \sigma * B | \sigma * C, \sigma * D \rangle,$$

 G_{χ} also acts on the set of bi-quadratic inequalities. The only non-trivial fact, we have to prove, is the following

Lemma 3.1. The property of $\langle A, B \mid C, D \rangle$, being a bi-quadratic inequality, remains fixed under the action of G_{χ} , i.e., for any $\sigma \in G_{\chi}$, $\sigma * \langle A, B \mid C, D \rangle$ is again a bi-quadratic inequality.

Proof. If $\langle A, B \mid C, D \rangle$ is a bi-quadratic inequality, then there are brackets E and F such that AB - CD + EF is a Graßmann-Plücker polynomial, and AB > 0, CD > 0, EF > 0. Since Graßmann-Plücker polynomials in R/I do not change by renumbering the vertices,

$$(\sigma * A)(\sigma * B) - (\sigma * C)(\sigma * D) + (\sigma * E)(\sigma * F)$$

is a Graßmann-Plücker polynomial as well. For any rotation or reflection σ , we have:

$$(\sigma * A)(\sigma * B) > 0$$
, $(\sigma * C)(\sigma * D) > 0$, $(\sigma * E)(\sigma * F) > 0$.

Thus, by Definition 2.1, $\langle \sigma * A, \sigma * B | \sigma * C, \sigma * D \rangle$ yields a bi-quadratic inequality again, in other words, G acts on the set of bi-quadratic inequalities. \triangle

Now, we introduce the notion of a symmetric realization of χ with respect to G_R , where $G_R < G_{\chi}$ and $G_R < O(d)$, is any subgroup of G_{χ} and O(d) respectively. By that, we mean a realization R of χ with a geometric automorphism group isomorphic to G_R .

Given such a symmetric realization, we have for any $\lambda \in E^d$ and any $\sigma \in G_{\mathcal{R}}$

$$[\lambda] - \sigma * [\lambda] = 0.$$

Now, we want to factorize our ring R/I, such that elements together with their symmetric images are identified. Therefore, we define the ideal $I_{G_{\mathcal{R}}}$ in R which is generated by polynomials of the following type

 $[\lambda] - \sigma * [\lambda]$ for any $\sigma \in G_{\mathcal{R}}$ and any $\lambda \in E^d$. We also consider the polynomial ring $R/\overline{I_{G_{\mathcal{R}}}}$, where $\overline{I_{G_{\mathcal{R}}}}$ denotes the ideal generated by $I \cup I_{G_{\mathcal{R}}}$.

Here we have factored out the symmetry of χ . A symmetric bi-quadratic final polynomial with respect to $G_{\mathcal{R}}$ will now be defined as follows.

Definition 3.1. A simplicial chirotope χ admits a symmetric bi-quadratic final polynomial with respect to $G_{\mathcal{R}}$, whenever, there is a collection of bi-quadratic inequalities

$$\langle A_i, B_i \mid C_i, D_i \rangle \in \mathcal{B}_{\chi}; \quad 1 \leq i \leq k,$$

such that

$$\prod_{i=1}^k A_i \cdot B_i \stackrel{\overline{I_{G_R}}}{=} \prod_{i=1}^k C_i \cdot D_i.$$

With this definition we obtain the following result:

Theorem 2. For any subgroup $G_{\mathcal{R}} < G_{\chi}$ we have: A (symmetric) chirotope admits a bi-quadratic final polynomial, if and only if it admits a symmetric bi-quadratic final polynomial with respect to $G_{\mathcal{R}}$.

Remark: The result of Theorem 2 heavily and positively influences the computational part of our applications. Whenever we want to prove that a symmetric chirotope is not realizable and we are looking for a bi-quadratic final polynomial, we can reduce the computational time tremendously. Especially, if the inequality system being considered generally has N inequalities, we only have to solve a system of approximately $N/|G_\chi|$ inequalities. Our proof will show how to construct the bi-quadratic final polynomial whenever the symmetric one is given.

Proof of Theorem 2.

One part of this theorem is very easily proven. Since for any $P, Q \in R$,

$$P \stackrel{I}{=} Q$$
 also implies $P \stackrel{\overline{I_{G_R}}}{=} Q$,

the existence of a bi-quadratic final polynomial implies the existence of a symmetric one. To prove the reverse, we first provide the following two lemmas.

Lemma 3.1. For any $\lambda_1, \lambda_2 \in E^d$, we have

$$[\lambda_1] \stackrel{\overline{I_{G_R}}}{=} [\lambda_2] \Rightarrow [\lambda_1] \in G_R * ([\lambda_2] + I).$$

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Proof. $[\lambda_1] = \overline{I_{\sigma_{\mathcal{R}}}} [\lambda_2]$ implies, there exists a $\sigma \in G_{\mathcal{R}}$ such that $[\lambda_1] = \overline{I_{\sigma_{\mathcal{R}}}} [\lambda_2]$ holds, and this in turn can be written as

$$[\lambda_1] \in \sigma * [\lambda_2] + I = \sigma * ([\lambda_2] + I) \in G_{\mathcal{R}}([\lambda_2] + I).$$

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Lemma 3.2. For any $\lambda_1, \lambda_2 \in E^d$, we have

$$[\lambda_1] \stackrel{\overline{I_{G_R}}}{=} [\lambda_2] \quad \Rightarrow \quad \prod_{\sigma \in G_R} (\sigma * [\lambda_1]) \stackrel{\underline{I}}{=} \prod_{\sigma \in G_R} (\sigma * [\lambda_2]) .$$

Proof. For $\lambda_1, \lambda_2 \in E^d$, we assume

$$[\lambda_1] \stackrel{\overline{I_{G_R}}}{=} [\lambda_2].$$

We consider the orbits $\Omega([\lambda_i]) := \prod_{\sigma \in G_{\mathcal{R}}} (\sigma * [\lambda_i])$, i = 1, 2, of $G_{\mathcal{R}}$ within R generated by $[\lambda_1]$, and $[\lambda_2]$, respectively. They induce corresponding orbits

$$\Omega([\lambda_i]+I) = \prod_{\sigma \in G_{\mathcal{R}}} \sigma * \{[\lambda_i]+I\} = \prod_{\sigma \in G_{\mathcal{R}}} (\sigma * [\lambda_i]) + I, \ i=1,2.$$

To pursue the proof of Theorem 2, let us assume, that χ admits a symmetric bi-quadratic final polynomial with respect to $G_{\mathcal{R}}$. Then, there is a collection of bi-quadratic inequalities

$$\langle A_i, B_i | C_i, D_i \rangle \in \mathcal{B}_{\chi}; \quad 1 \leq i \leq k,$$

such that

$$(**) \qquad \prod_{i=1}^k A_i \cdot B_i \stackrel{\overline{I_{\sigma_n}}}{=} \prod_{i=1}^k C_i \cdot D_i.$$

We will show that in replacing any bi-quadratic inequality by its orbit under $G_{\mathcal{R}}$, we get a biquadratic final polynomial in R/I. In other words, the collection

$$\sigma * \langle A_i, B_i | C_i, D_i \rangle \in \mathcal{B}_{\chi}; \quad 1 \leq i \leq k; \quad \sigma \in G_{\mathcal{R}}$$

is a bi-quadratic final polynomial. Finally, we have to show that

$$\prod_{\substack{1 \leq i \leq k \\ \sigma \in G_{\mathcal{R}}}} (\sigma * A_i) \cdot (\sigma * B_i) \stackrel{I}{=} \prod_{\substack{1 \leq i \leq k \\ \sigma \in G_{\mathcal{R}}}} (\sigma * C_i) \cdot (\sigma * D_i).$$

This is equivalent to

$$\prod_{1 \leq i \leq k} \left(\prod_{\sigma \in G_{\mathcal{R}}} \sigma * A_i \right) \cdot \left(\prod_{\sigma \in G_{\mathcal{R}}} \sigma * B_i \right) \stackrel{I}{=} \prod_{1 \leq i \leq k} \left(\prod_{\sigma \in G_{\mathcal{R}}} \sigma * C_i \right) \cdot \left(\prod_{\sigma \in G_{\mathcal{R}}} \sigma * D_i \right)$$

Since (**) is only valid, if for any bracket $[\lambda]$ on the left of (**), there is a bracket $[\lambda']$ on the right of (**) with $[\lambda] = \overline{I_{G_R}} [\lambda']$, the result follows according to Lemma 2.2. \triangle

4. Applications

In this section we provide a collection of non-representable oriented matroids together with their bi-quadratic final polynomials. We first give some examples of known non-representable proofs and even known final polynomials by using our theorems as discussed in Section 3, we can provide extremely short non-realizability proofs.

Example 1. The Vamos matroid

As a first example, we give a non-realizability proof for an oriented version of the well-known Vamos matroid of 8 points in rank 4. To give a combinatorial description of the oriented matroid, we start with the chirotope χ which corresponds to the following rank 4 configuration.

Notice that this configuration has a symmetry group isomorphic to $\mathbb{Z}_2 \times S_4$, since the points 1, 2, 3 and 4 form a regular tetraedron in the subspace $\{(w, x, y, z) \mid w = 1\}$ and the points 5, 6, 7 and 8 form a similar one in the subspace $\{(w, x, y, z) \mid w = -1\}$. Notice furthermore that we have

$$[1, 2, 5, 6] = [1, 3, 5, 7] = [1, 4, 5, 8] = [2, 3, 6, 7] = [2, 4, 6, 8] = [3, 4, 7, 8] = 0.$$

We modify the chirotope χ in order to get a new map χ_V by replacing exactly these zero-orientations by 1, i.e., we require

$$\chi_{V}(1,2,5,6) = \chi_{V}(1,3,5,7) = \chi_{V}(1,4,5,8) = \chi_{V}(2,3,6,7) = \chi_{V}(2,4,6,8) = \chi_{V}(3,4,7,8) = +1$$

and $\chi = \chi_V$ for all other 4-tupels. We will show that χ_V is again a chirotope. We consider any Graßmann-Plücker polynomial $\{a,b \mid \ldots\}$ where a,b are vertices of the same tetrahedron, as mentioned above. In this polynomial $\{a,b \mid \ldots\}$, there occurs no other basis λ with $\chi(\lambda) = 0$, since all 4-tupels contain a and b. But if a and b lie in different tetrahedra, and if there would be another λ in the polynomial with $\chi(\lambda) = 0$, we see from the structure of the polynomial $\{a,b \mid \ldots\}$ and from the six 4-tupels above that this zero-valued $\chi(\lambda)$ is the second factor in the same summand of the polynomial which is already equal to zero. Two other summands of opposite signs must exist, showing that the chirotope condition for χ_V is fulfilled.

Notice that χ_V has a combinatorial symmetry isomorphic to the alternating group A_4 which can be generated by the permutations

So far, we have only given a description of our chirotope χ_V . Now we claim that χ_V is not realizable.

Lemma 4.1. The Vamos matroid χ_V is not realizable, and χ_V admits a bi-quadratic final polynomial.

Proof. Consider the Graßmann-Plücker polynomial

$$\{1, 2 \mid 3, 6, 5, 4\} = \underbrace{[1, 2, 3, 6][1, 2, 5, 4]}_{+} - \underbrace{[1, 2, 3, 5]}_{+} \underbrace{[1, 2, 6, 4]}_{+} + \underbrace{[1, 2, 3, 4]}_{-} \underbrace{[1, 2, 6, 5]}_{-} = 0.$$

 $\langle [1,2,3,6], [1,2,5,4] | [1,2,3,5], [1,2,6,4] \rangle$ is a bi-quadratic inequality of χ_V . We assume that χ_V is symmetrically (A_4) realizable, therefore, we have permutations $\sigma_1 := (123)(567), \quad \sigma_2 := (124)(568) \in G_{\chi_V}$ such that

$$\sigma_1^{-1} * [1, 2, 3, 6] = [1, 2, 3, 5] \text{ and } \sigma_2 * [1, 2, 5, 4] = [1, 2, 6, 4].$$

Thus, we have $[1, 2, 3, 6][1, 2, 5, 4] = \tilde{I}_{G\times V}[1, 2, 3, 5][1, 2, 6, 4]$. In conclusion, Theorem 2 proves the lemma, i.e., χ_V is not realizable at all.

Example 2. Desargues Theorem.

Consider the configuration D_{10}^3 given in Figure 1. This configuration is an oriented version of Desargues Theorem. Notice that in the corresponding chirotope χ_D , we have the mutations

$$[1, 0, 6], [2, 0, 4], [3, 0, 5], [1, 3, 7], [2, 1, 8], [3, 2, 9], [7, 5, 6], [8, 6, 4], [9, 4, 5], [7, 8, 9].$$

When we require $\chi(\lambda) = 0$ for the corresponding 3-tupels λ , we get exactly a chirotope corresponding to Desargues configuration. χ_D has a symmetry isomorphic to Z_3 generated by $\sigma = (123)(456)(789)$.

insert Figure 1.

Lemma 4.2.

The Desargue chirotope χ_D is not realizable, and χ_D admits a bi-quadratic final polynomial. Proof. Consider the Grassmann-Plücker relations

The underlined brackets have to be negative in any realization of χ_D . This yields the bi-quadratic inequalities

$$\langle [1, 2, 0], [1, 6, 3] | [1, 2, 6], [1, 0, 3] \rangle,$$

 $\langle [1, 7, 6], [1, 5, 3] | [1, 7, 5], [1, 6, 3] \rangle,$
 $\langle [7, 5, 1], [7, 8, 6] | [7, 5, 8], [7, 1, 6] \rangle,$
 $\langle [7, 8, 5], [7, 0, 9] | [7, 8, 0], [7, 5, 9] \rangle.$

Now the symmetry of χ_D yields the following identities

$$\sigma^2 * [1, 2, 0] = [3, 1, 0], \sigma * [1, 5, 3] = [2, 6, 1], \sigma^2 * [7, 8, 6] = [9, 7, 5], \sigma * [7, 0, 9] = [8, 0, 7],$$

and this in turn leads to the identity

$$[1, 2, 0][1, 6, 3][1, 7, 6][1, 5, 3][7, 5, 1][7, 8, 6][7, 5, 8][7, 0, 9] \stackrel{\overline{I}_{G_{XD}}}{=} D$$

$$[1, 2, 6][1, 0, 3][1, 7, 5][1, 6, 3][7, 5, 8][7, 1, 6][7, 8, 0][7, 5, 9].$$

Now Theorem 2 gives the desired result.

Example 3. A star-chirotope

In [7] an example of an infinite class of nonrealizable chirotopes was given, such that any proper minor of each of those chirotopes is realizable. The non-realizability proof was given by standard geometric arguments, and finding corresponding final polynomials turned out to be difficult, when specific coordinates were chosen. Here, we give a bi-quadratic final polynomial for the smallest example of this class, see Figure 2.

insert Figure 2.

Lemma 4.3.

The star-chirotope χ_5 corresponding to the configuration in Figure 2 is not realizable, and χ_5 admits a bi-quadratic final polynomial.

Proof. We give a list of the used Graßmann-Plücker relations and the corresponding bi-quadratic inequalities.

Any involved bracket appears exactly once, both on the right and on the left. In conclusion, the product of all brackets on the right equals in the ring R/I the product of all brackets on the left, proving our lemma.

Example 4. The non-polytopal Altshuler sphere M_{425}^{10} .

We give the combinatorial neighborly description of Altshuler's 3-sphere M_{425}^{10} with 10 vertices in terms of its facets.

```
1048
1045
           4579
                 8275
3267
     1062
           4531
                 0497
5489
     0135
           7604
                 2619
7901
     0197 7628 4831
9823
     3260 6791
                 6053
0132
     3284
           6753
     2357
           9826
2354
     2319
4576
           9840
6798
     5482
           8913
8910
     5406 8975
```

The cyclic symmetry-group G generated by $\sigma = (13579)(24680)$ is a subgroup of the automorphism group of this sphere. Since the sphere M_{425}^{10} is neighborly, it is also rigid by Shemer's result,[12], i.e., there are only two chirotopes χ and $-\chi$ which are compatible in the sense of [6] to the polyhedral structure of the polytope. M_{425}^{10} is realizable, if and only if these chirotopes are realizable. Since the chirotopes are unique up to a multiplication by -1, they must have the same automorphism group as in the sphere M_{425}^{10} . The facial structure of M_{425}^{10} forces the following identities (for details see [3]):

$$\chi((2,6,0,1,8)) = \chi((2,6,0,1,4)) = \chi((2,6,0,1,3))$$
$$= \chi((2,6,0,4,3)) = \chi((2,6,0,8,3)) = \chi((8,6,0,1,7))$$

On the other hand, the symmetry requires:

$$\chi((2,6,0,1,4)) = \chi(\sigma^{3}(2,6,0,1,4)) = \chi((8,2,6,7,0)),$$

$$\chi((2,6,0,4,3)) = \chi(\sigma^{2}(2,6,0,4,3)) = \chi((6,0,4,8,7)),$$

$$\chi((2,6,0,8,3)) = \chi(\sigma^{4}(2,6,0,8,3)) = \chi((0,4,8,6,1)).$$

We claim that M_{425}^{10} is not realizable. To prove this, we have simply to consider two Grasmann-Plücker polynomials $\{6, 8, 0 \mid 1, 7, 4, 2\}$ and $\{2, 6, 0 \mid 1, 3, 8, 4\}$. With these preparations, we leave it to the reader to work out the conclusion by using Theorem 2.

Example 5. A projective theorem.

Our examples chosen above were oriented versions of projective theorems of the form whenever the brackets $[a], [b], [c], \ldots$ are zero, then [x] has to be zero, too.

Another type of projective theorem is of the form

whenever the brackets $[a], [b], [c], \ldots$ are zero and $[r], [s], [t], \ldots$ are non-zero, then [x] has to be nonzero, too.

The following theorem of this type has been proven in [8] by using a computer algebra system. Theorem. For any 10 points given in R^3 , we have, whenever

$$[1,2,3] = [1,6,0] = [1,7,9] = [2,4,0] = [2,5,6] = [3,4,9] = [3,5,7] = [4,5,9] = [6,7,8] = 0$$

and whenever all other brackets except [8, 9, 0] are nonzero, then [8, 9, 0] is nonzero. Especially, we have $[8, 9, 0] = [2, 3, 4] \cdot [1, 6, 7]^2 \cdot [3, 4, 7]$.

One oriented version of this theorem is that the configuration in Figure 3. is not realizable. This fact can again be proven by using a bi-quadratic final polynomial. The smallest final polynomial we know of was found on a computer algebra program that solves LP-problems with integer

coefficients. This final polynomial is horribly large and consists of 135 bi-quadratic inequalities. A symmetric version with respect to the symmetry-group generated by (176)(253)(908) with 45 bi-quadratic inequalities can be sent by the first author on request.

insert Figure 3.

Example 6. Classification of matroid manifolds.

What actually stimulated the investigations which led to our results in this paper was a classification problem for 31 matroid manifolds. This was recently completed but will be described in a forthcoming paper [4].

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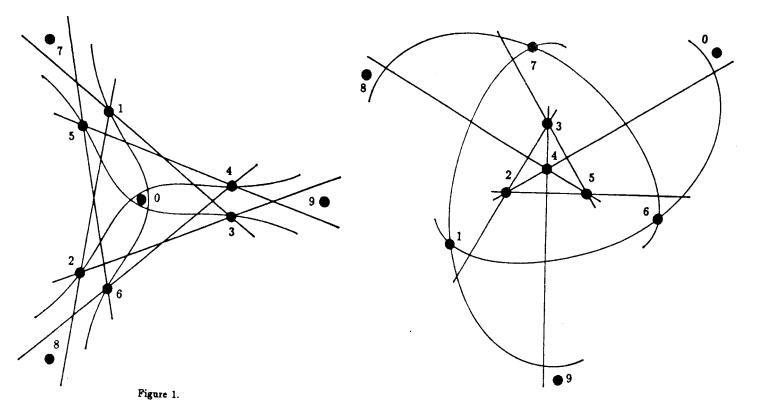


Figure 3.

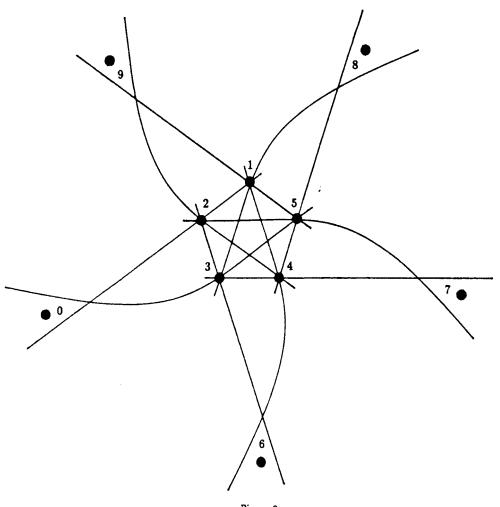


Figure 2.