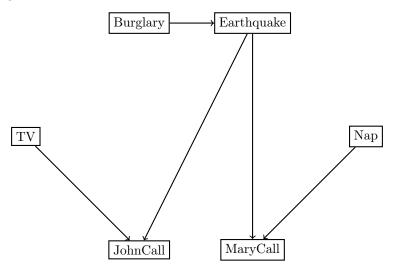
HW 3 3.10,3.11a,3.14,3.15

3.10. Let X be a node in a graph \mathscr{G} and let $Pa_X\mathscr{G}$ be the parents of X. Let P be a probability distribution and $\mathscr{I}(\mathscr{G})$ be the set of global independencies over P. Let Y be a non-descendant of X. Then $X \perp Y \mid Z$ locally, where $Z \in Pa_X\mathscr{G}$. We prove that X and Y are d-separated. By way of contradiction, assume that X and Y are not d-separated. Then there exists an active trail between X and Y given Z.

By definition 3.6, this is only possible if for the trail $X \rightleftharpoons X_0 \rightleftharpoons X_n \rightleftharpoons Y$ where each X_i is a possible random variable in the trail, then whenever there is a v-structure, $X_{i-1} \to X_i \leftarrow X_{i+1}$, X_i or one of its descendants is Z and no other nodes are in Z. First assume that there are no v-structures. Since Z is a parent of X, then $X \to X_0$. However, since Y is a non-descendant of X, then there exists some $X_k \leftarrow X_{k+1}$. That means there must exist at least one v-structure. Thus, Z is in the trail between X and Y. This is a contradiction. So assume that there is at least one v-structure. Then $X \leftarrow Z$. This is a contradiction since if Z is in the trail, then all surrounding arrows must point towards Z. So there does not exist an active trail between X and Y, hence X and Y are d-separated given Z. Since X,Y and Z are general, we see every locally independency is represented in $\mathscr{I}(\mathscr{G})$. So global independencies imply local independencies.

3.11a. Let \mathscr{B} be the burglary alarm network from fig 3.15. Consider the marginal distribution $P_{\mathscr{B}}(B, E, T, N, J, M)$. We construct the diagram



3.14. Let \mathscr{G} be the Bayesian network, X the source variable, and \mathbf{Z} be the observations on the network. We want to find all the nodes reachable from X given \mathbf{Z} via active trails. This means we want all the trails from X to Y (where Y is arbitrary) which only have a v-structure $X_{i-1} \to X_i \leftarrow X_{i+1}$ if X_i or a descendant of X_i is in Z. This means there are 3 cases:

- 1. Y and X do not have Z in their trail.
- 2. Z is a descendant of both X and Y.
- 3. Y is Z.

We show that the algorithm finds all trails in these three cases.

- 1. If Y is an ancestor of X, then we look at the trail beginning with (X,\uparrow) . We assume that any node on this trail is not in Z. We add this to our set L and then select it from L, add it to V and add it to R since $X \notin Z$. Since $d = \uparrow$, we add the parents of X to L and repeat, starting with the first parent of L. As we continue this, we will pass through all the ancestors of X, so we will find the trail between Y and X.
 - If Y is not a direct ancestor of X, then after travelling through X, we check if Y is a direct ancestor of the new node we select. If it is, we follow the above proof. If it is not, add all the children and all the parents as nodes we could travel to (as to check these active trails connect). Then as we pass through nodes, since Z is not in the trail, we will always continue to check the children and parent trails. Eventually, we will arrive at the node with ancestor Y (or is Y). At this point, we will find the trail by visting trails up through the node.
- 2. Assume that Z is in the trail from X to Y. We continue through all possible up and down active trails as in part (i) until we arrive at Z. Once we arrive at Z, then we know we are travelling down (which must be true because it is a v-structure. We then must travel to the parents of Z (which are in L from phase 1). We then continue the part (i) until we arrive at Y. If there is more than one node from Z in the trail, each time we arrive at the node, we know we must travel upwards, which happens in the last if statement of the algorithm. We do this until we have passed through all observed nodes and arrive Y.
- 3. If Z is the node we are travelling to, we can apply part (i).

Hence we can find every trail with this algorithm.

3.15.

- 1. There is no Bayesian network I-equivalent to (a).
- 2. The graph below is equivalent to (b).

