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MATH 308

Fall 2021

HW 2: Due 09/24

“Some people are immune to good advice.”

– Saul Goodman, Breaking Bad

Problem 1. (10pt) Determine if the following sentences are predicates. If the sentence is a predicate, mark it ‘T’; otherwise, mark the sentence ‘F.’ [Let the universal set be \mathbb{R} .]

- (a) T : x is odd.
- (b) F : $x^2 + x + 1$
- (c) T : $P(x): x^2 + 1 < 0$
- (d) T : $Q(x): x$ is an integer
- (e) T : $R(x, y): x^2 < y^3$

Solution. Recall that predicate becomes a proposition when the variables are substituted for their values.

- (a) Given an x , either x is odd or it is not.
- (b) Given an x , the result is simply an expression and will not hold a true or false value.
- (c) Given an x , the resulting inequality will either be true or false.
- (d) Given an x , either x is an integer or not.
- (e) Given a x, y , either the resulting inequality will either be true or false.

Problem 2. (10pt) Give an original example of a predicate having more than one variable.

Solution. Remember that to be a predicate, once the variables have been substituted, the result should be a proposition, i.e. something which is unambiguously true or false. There are infinitely many examples, e.g. $P(x, y): x = y$, $Q(x, y, z): x^2 + y^2 = z^2$, etc.

Problem 3. (10pt) Let $P(x)$ be the predicate $P(x): 1 \leq 2^x \leq 100$. Suppose that the domain is the nonnegative integers. What is the truth set for $P(x)$? What is the truth set if the domain were instead the set of real numbers?

Solution. Suppose the domain is the nonnegative integers, i.e. $\mathbb{Z}_{\geq 0} = \{0, 1, 2, 3, \dots\}$. We need $1 \leq 2^x \leq 100$. For the left inequality to hold, we need $1 \leq 2^x$, which implies that $0 \leq x$. For the right inequality to hold, we require $2^x \leq 100$. But this implies that $x \leq \log_2 100 \approx 6.644$. Because $x \in \mathbb{Z}$, this implies that $x \leq 6$. Then $x \geq 0$ and $x \leq 6$, i.e. $0 \leq x \leq 6$. Therefore, the truth set is $\{0, 1, 2, 3, 4, 5, 6\}$.

Suppose now that the domain is the set of real numbers. For the left inequality to hold, we need $1 \leq 2^x$, which implies that $0 \leq x$. For the right inequality to hold, we require $2^x \leq 100$. But this implies that $x \leq \log_2 100 \approx 6.644$. Then $0 \leq x$ and $x \leq \log_2 100$. Therefore, the truth set is $[0, \log_2 100]$.

Problem 4. (10pt) Defining appropriate propositional functions and variables, write the following English sentences using logical symbols and the functions/variables that you defined.

- (a) Every cloud has a silver lining.
- (b) All that glitters is not gold.
- (c) Every human is guilty of all the good that they did not do.
- (d) None but the brave deserve the fair.

Solution. There are many ways to do this. . .

- (a) Let $P(x)$ be the predicate that x is a cloud and $Q(x)$ denote the predicate that x has a silver lining. Then 'Every cloud has a silver lining' could be translated as $\forall x, P(x) \rightarrow Q(x)$.
- (b) Let $P(x)$ be the predicate that x glitters and let $Q(x)$ denote the predicate that x is gold. Then 'All that glitters is not gold' could be translated as $\exists x, P(x) \wedge \neg Q(x)$.
- (c) Let $P(x, y)$ denote the predicate that x did y , $Q(x)$ denote the predicate that x is good, and $R(x, y)$ denote the predicate that x is guilty of y . Then 'Every human is guilty of all the good that they did not do' could be translated as $(\forall x)(\forall y)[Q(y) \wedge \neg P(x, y) \rightarrow R(x, y)]$.
- (d) Let $P(x)$ be the predicate that x is brave and let $Q(x)$ denote the predicate that x deserves the fair. Then 'None but the brave deserve the fair' could be translated as $\forall x, Q(x) \rightarrow P(x)$.

Problem 5. (10pt) Define the following predicates:

- (i) $P(x) : x > 0$
- (ii) $Q(x) : x$ is even
- (iii) $R(x) : x$ is a perfect square
- (iv) $S(x) : x$ is divisible by 4
- (v) $T(x) : x$ is divisible by 5

Write the following in symbolic form:

- (a) Any perfect square is positive.
- (b) If an integer is divisible by 4, then the integer is even.
- (c) No even integer is divisible by 5.

Write the following in the form of an English sentence:

- (d) $\forall x (S(x) \rightarrow Q(x))$
- (e) $\exists x (S(x) \wedge \neg R(x))$

Solution.

- (a) $\forall x, R(x) \rightarrow P(x)$
- (b) $\forall x, S(x) \rightarrow Q(x)$
- (c) $\forall x, Q(x) \rightarrow \neg T(x)$
- (d) All integers divisible by 4 are even.
- (e) There exists an integer divisible by four that is not a perfect square.

Problem 6. (10pt) Let $P(x)$ be the predicate $x^2 = x$. Determine if the following statements are true or false:

- (a) T : $P(0)$
- (b) F : $P(-1)$
- (c) F : $\forall x, P(x)$
- (d) T : $\exists x, P(x)$
- (e) F : $\exists!x, P(x)$

Solution.

- (a) $P(0)$: $0^2 = 0$, which is equivalent to $0 = 0$, which is true.
- (b) $P(-1)$: $(-1)^2 = -1$, which is equivalent to $1 = -1$, which is false.
- (c) From (b), we already know that -1 is a counterexample to this statement.
- (d) From part (a), we already know that $x = 0$ is such an example.
- (e) $P(1)$: $1^2 = 1$, which is equivalent to $1 = 1$, which is true. Combining this with (a), we know there are at least two values for x , namely 0 and 1, for which $P(x)$ is true.

Problem 7. (12pt) Let $P(x)$, $Q(x)$, $R(x)$ denote the predicates $1 - 2x = 7$, $x^2 = 9$, and $x^2 > 9$, respectively. Determine whether the following propositions are true or false. If the statement is true, explain why. If the statement is false, give a counterexample.

- (a) $(\forall x)(P(x) \wedge Q(x))$
- (b) $(\exists x)(P(x) \wedge Q(x))$
- (c) $(\forall x)(P(x) \rightarrow Q(x))$
- (d) $(\forall x)(P(x) \rightarrow R(x))$
- (e) $(\exists x)(P(x) \vee R(x))$
- (f) $(\exists!x)(P(x) \wedge Q(x))$

Solution.

- (a) The statement is *false*. For the statement to be true, we need both $P(x)$ and $Q(x)$ to be true. However, for ‘most’ x in the domain, both $P(x)$ and $Q(x)$ are false. As a counterexample, if $x = 0$ then both $P(x)$ and $Q(x)$ are false so that $P(x) \wedge Q(x)$ are false. Then the truth set is not the entire domain.
- (b) The statement is *true*. For the statement to be true, there must be an x such that $P(x)$ and $Q(x)$ are true. If $P(x)$ is true, then $1 - 2x = 7$, which implies $x = -3$. But if $x = -3$, then $(-3)^2 = 9$ so that $Q(x)$ is true. But then if $x = -3$, both $P(x)$ and $Q(x)$ are true so that $P(x) \wedge Q(x)$ is true. But then the truth set is non-empty.
- (c) The statement is *true*. We know that $P(x)$ is true if and only if $1 - 2x = 7$ is true, which is true if and only if $x = -3$. From (b), we know that if $x = -3$ that $Q(x)$ is true. But then for $x = -3$, $P(x) \rightarrow Q(x)$ is true. If $x \neq -3$, then $P(x)$ is false so that $P(x) \rightarrow Q(x)$ is true. But then the truth set is the entire domain.
- (d) The statement is *false*. From (c), we know that if $x = -3$, then $P(x)$ is true. But if $x = -3$, then $R(x): (-3)^2 > 9$ is false. Then $P(x) \rightarrow R(x)$ is false. Therefore, $x = -3$ is a counterexample. Then the truth set is not the entire domain.
- (e) The statement is *true*. For $P(x) \vee R(x)$ to be true, either $P(x)$ or $R(x)$ is true. From (b), we know that if $x = -3$, then $P(x)$ is true. But then $P(x) \vee R(x)$ is true. Therefore, the truth set is non-empty.
- (f) The statement is *true*. For $P(x) \wedge Q(x)$ to be true, $P(x)$ and $Q(x)$ need to be true. From (c), we know the only x in the domain such that $P(x)$ is true is $x = -3$. But from (c), we know that $Q(x)$ is true if $x = -3$. Therefore, if $x = -3$, then $P(x)$ and $Q(x)$ are true so that $P(x) \wedge Q(x)$ is true. But if $x \neq -3$, then $P(x)$ is false so that $P(x) \wedge Q(x)$ is false. But then $x = -3$ is the unique value in the truth set.

Problem 8. (10pt) What well-known property does the following proposition represent: $\forall x \forall y \forall z [x + (y + z) = (x + y) + z]$.

Solution. This is quantified version of the associative property for addition for the real numbers.

Problem 9. (10pt) Determine if the following statements are true or false:

- (a) T : $\exists x \exists y (xy = 1)$
- (b) F : $\exists x \forall y (xy = 1)$
- (c) F : $\forall x \exists y (xy = 1)$
- (d) F : $\forall x \forall y (x^2 + y = 1)$
- (e) T : $\forall x \exists y (x^2 + y = 1)$

Solution.

- (a) If $x = y = 1$, then $xy = 1$.
- (b) If $y = 0$, then $xy = 0 \neq 1$ for all $x \in \mathbb{R}$.
- (c) If $x = 0$, then $xy = 0 \neq 1$ for all $y \in \mathbb{R}$. Note that if we require $x \neq 0$, the statement is true: if $x \neq 0$, define $y := 1/x$. But then $xy = x \cdot \frac{1}{x} = 1$.
- (d) If $x = y = 0$, then $x^2 + y = 0 \neq 1$.
- (e) Given $x \in \mathbb{R}$, define $y := 1 - x^2$. But then $x^2 + y = x^2 + (1 - x^2) = 1$.

Problem 10. (10pt) Negate the following proposition:

$$\forall x \exists y (P(x, y) \wedge Q(x, y) \rightarrow R(x, y))$$

Solution.

$$\neg[\forall x \exists y (P(x, y) \wedge Q(x, y) \rightarrow R(x, y))]$$

$$\exists x \neg[\exists y (P(x, y) \wedge Q(x, y) \rightarrow R(x, y))]$$

$$\exists x \forall y \neg (P(x, y) \wedge Q(x, y) \rightarrow R(x, y))$$

$$\exists x \forall y (P(x, y) \wedge Q(x, y) \wedge \neg R(x, y))$$

Problem 11. (10pt) One way of stating the definition for a function $f(x)$ to have limit L at x , i.e. $\lim_{x \rightarrow a} f(x) = L$, is as follows:

$$(\forall \epsilon > 0)(\exists \delta > 0)(\forall x \in \mathbb{R})[0 < |x - a| < \delta \Rightarrow |f(x) - L| < \epsilon]$$

Give a definition for a function $f(x)$ to not have a limit at $x = a$ by negating the statement above. Your answer should not contain any negations.

Solution.

$$\neg[(\forall \epsilon > 0)(\exists \delta > 0)(\forall x \in \mathbb{R})[0 < |x - a| < \delta \Rightarrow |f(x) - L| < \epsilon]]$$

$$(\exists \epsilon > 0)\neg[(\exists \delta > 0)(\forall x \in \mathbb{R})[0 < |x - a| < \delta \Rightarrow |f(x) - L| < \epsilon]]$$

$$(\exists \epsilon > 0)(\forall \delta > 0)\neg[(\forall x \in \mathbb{R})[0 < |x - a| < \delta \Rightarrow |f(x) - L| < \epsilon]]$$

$$(\exists \epsilon > 0)(\forall \delta > 0)(\exists x \in \mathbb{R})\neg[0 < |x - a| < \delta \Rightarrow |f(x) - L| < \epsilon]$$

$$(\exists \epsilon > 0)(\forall \delta > 0)(\exists x \in \mathbb{R})[0 < |x - a| < \delta \wedge \neg(|f(x) - L| < \epsilon)]$$

$$(\exists \epsilon > 0)(\forall \delta > 0)(\exists x \in \mathbb{R})[0 < |x - a| < \delta \wedge (|f(x) - L| \geq \epsilon)]$$

Problem 12. (10pt) The universal and existential quantifier do not necessarily ‘distribute’ over \wedge and \vee . One of the following ‘equivalences’ is not correct:

$$\begin{aligned}\forall x (P(x) \wedge Q(x)) &\iff \forall x P(x) \wedge \forall x Q(x) \\ \forall x (P(x) \vee Q(x)) &\iff \forall x P(x) \vee \forall x Q(x)\end{aligned}$$

Determine which one is always true and state it. For the one that is false, give an example to show that it is false.

Solution. The proposition $\forall x (P(x) \wedge Q(x)) \iff \forall x P(x) \wedge \forall x Q(x)$ is true. If the left side is true, then for all x in the domain, $P(x) \wedge Q(x)$ is true, i.e. $P(x)$ is true and $Q(x)$ is true. But then $\forall x P(x)$ is true and $\forall x, Q(x)$ is true. But then $\forall x P(x) \wedge \forall x Q(x)$ is true. If the left side is false, then there is an x in the domain such that $P(x) \wedge Q(x)$ is false, i.e. either $P(x)$ is false or $Q(x)$ is false. Without loss of generality, assume $P(x)$ is false. But then $\forall x P(x)$ is false so that $\forall x P(x) \wedge \forall x Q(x)$ is false. Now suppose the right hand side is true. Then $\forall x P(x) \wedge \forall x Q(x)$ is true. But then $\forall x P(x)$ is true and $\forall x, Q(x)$ is true. But then $P(x) \wedge Q(x)$ is true for all x in the domain so that $\forall x (P(x) \wedge Q(x))$ is true. Suppose that the right hand side is false. Then $\forall x P(x) \wedge \forall x Q(x)$ is false so that either $\forall x, P(x)$ is false or $\forall x, Q(x)$ is false. Without loss of generality, assume that $\forall x, P(x)$ is false. Then there is an x in the domain such that $P(x)$ is false. But then $P(x) \wedge Q(x)$ is false for this x so that $\forall x (P(x) \wedge Q(x))$ is false.

To see that $\forall x (P(x) \vee Q(x))$ and $\forall x P(x) \vee \forall x Q(x)$ are not equivalent, consider predicates the following example: let the domain be the set of chess pieces and define $P(x)$ to be the predicate that x is black and define $Q(x)$ be the predicate that x is white. It is true that for any x in the domain, that either x is black or x is white, i.e. $\forall x (P(x) \vee Q(x))$. However, $\forall x P(x) \vee \forall x Q(x)$ is false because it is not true that every chess piece is black or that every chess piece is white. However, it is true that $\forall x P(x) \vee \forall x Q(x) \Rightarrow \forall x (P(x) \vee Q(x))$.