

Chris Grimes

CS 46101

HW 6

1. (10 points) A company named RT&T has a network of n switching stations connected by m high-speed communication links. Each customer's phone is directly connected to one station in his or her area. The engineers of RT&T have developed a prototype video-phone system that allows two customers to see each other during a phone call. In order to have acceptable image quality, however, the number of links used to transmit video signals between the two parties cannot exceed 4. Suppose that RT&T network is represented by a graph. Design and give the pseudo-code for an efficient algorithm that computes, for each station, the set of stations it can reach using no more than 4 links. Analyze its running time.

Algorithm span4Video(a vertex A, distance from source):

input: a vertex A and an int I which is the distance from the source of the video

// I should begin at 0 the first time this method is called

output: a set of vertices within 3 edges of vertex A

set rtnval=null

for each(edge in the graph that contains vertex A)//runs $O(n+m)$

 edge.label()=-1

 ++I

 if(I < 5)

 for each(incident edge of A)// runs a max of $O(n+m)$

 if(edge.label()==-1)

 edge.label()=I

 rtnval.add(other vertex connected to edge)

 span4video(other vertex connected to edge, I)

return rtnval

Assuming that links means edges the above algorithm is correct, if links means vertices the "if(I<5)" should be "if(I<4)"

Do to the two for loops in the above algorithm, the span4video algorithm will have a worst case running time of $O(n+m)$.

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2. (10 points) An Eulerian circuit of a directed graph G with n vertices and m edges is a cycle that traverses each edge of G exactly once according to its direction. Such a cycle always exists if the in-degree is equal to the out-degree for each vertex in G . Describe a $O(n + m)$ time algorithm for finding an Eulerian circuit of such a graph G . Analyze its running time.

Algorithm eulerianCircuit(a graph G):

input: a graph G

output: a list containing the vertices of a eulerian circuit in G

an empty list `rtval`

for each(`edge` in graph g)

`edge.color()`=red

v = an edge contained in graph G from which we'll start our search

`rtval.add(v)`

while(there is a red incident edge of V)

`edge.color()`=white

`rtval.add(other vertex connected to edge)`

v = other vertex connected to edge

return `rtval`

Do to the for loop and the while loop in the above algorithm, the eulerianCircuit algorithm will have a worst case running time of $O(n+m)$.

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3. (10 points) Consider the graph G depicted below, and perform the following graph search algorithms. Whenever faced with a decision of which vertex to pick from a set of vertices, pick the vertex whose label occurs earliest in the alphabet.

(a) Trace the execution of BFS beginning at vertex A, labeling each edge as a discovery or cross edge.

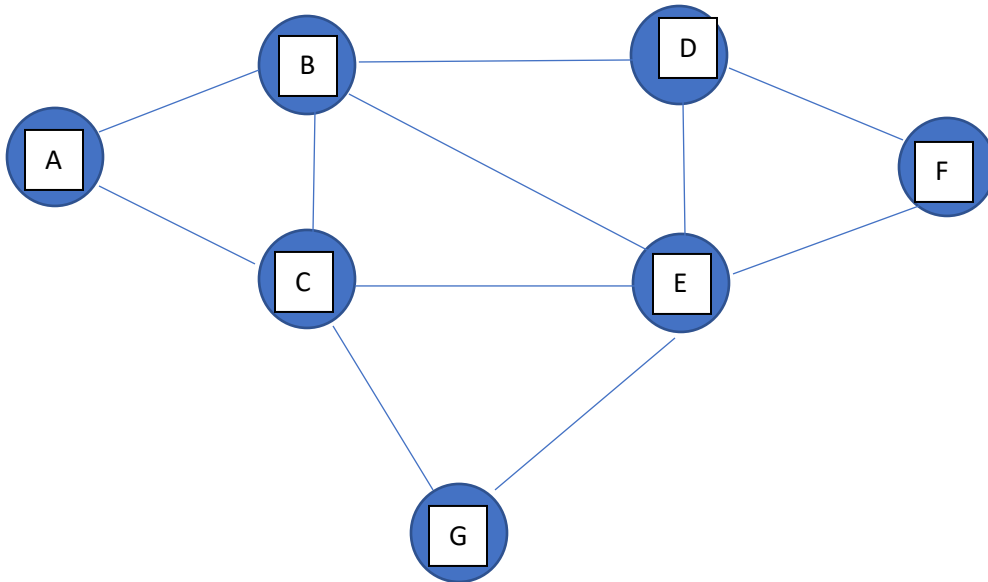
Blue circle= unexplored vertex

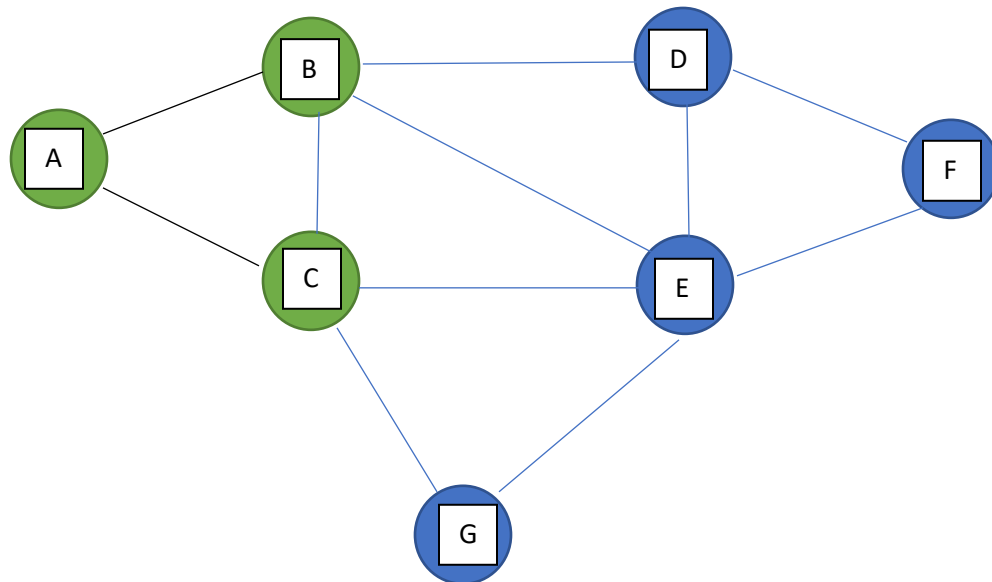
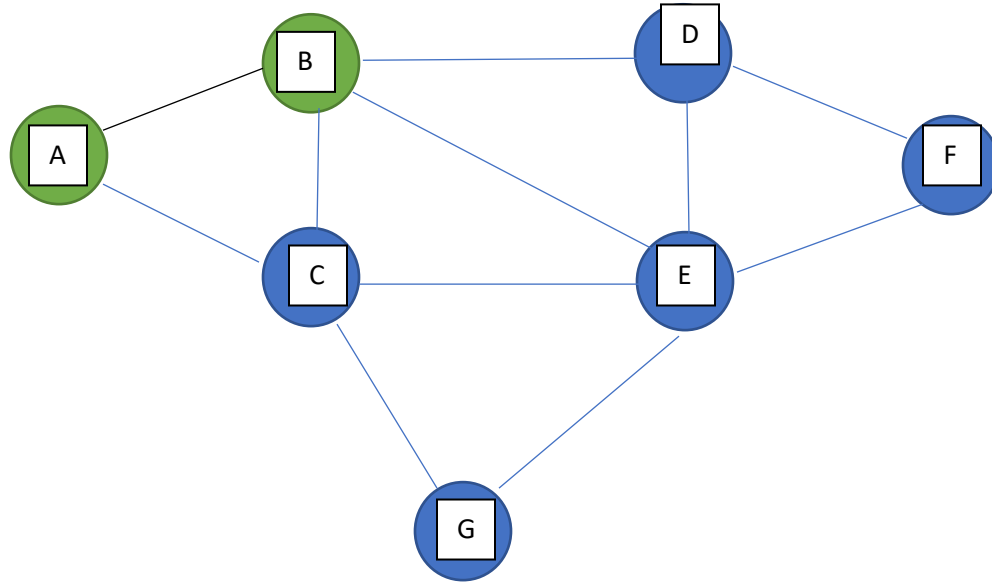
blue line= unexplored edge

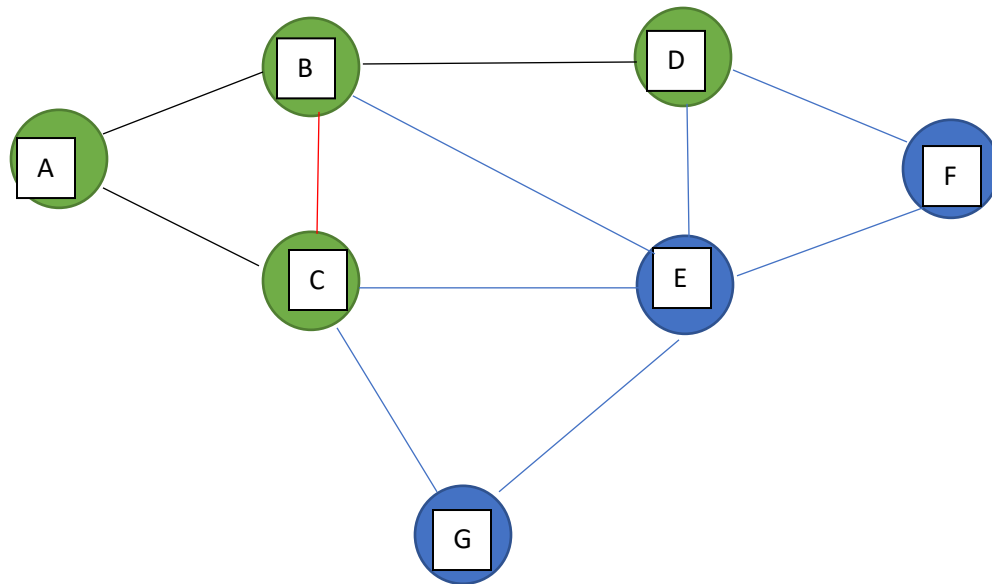
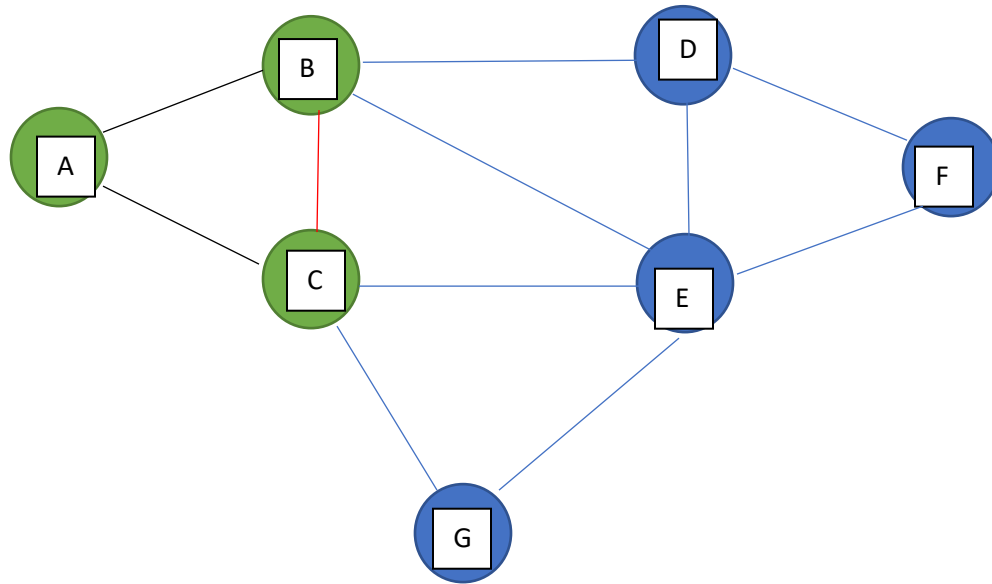
green circle= visited vertex

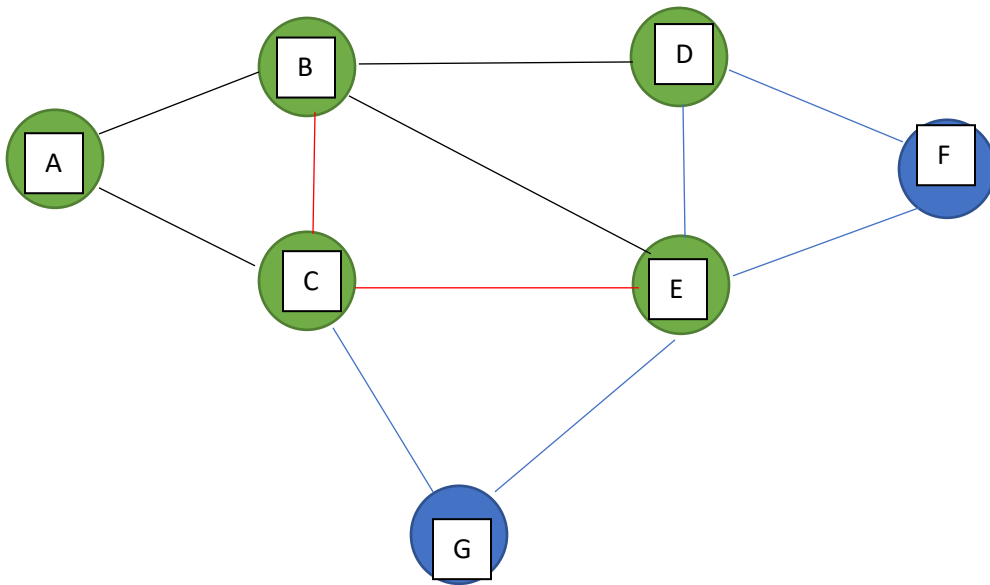
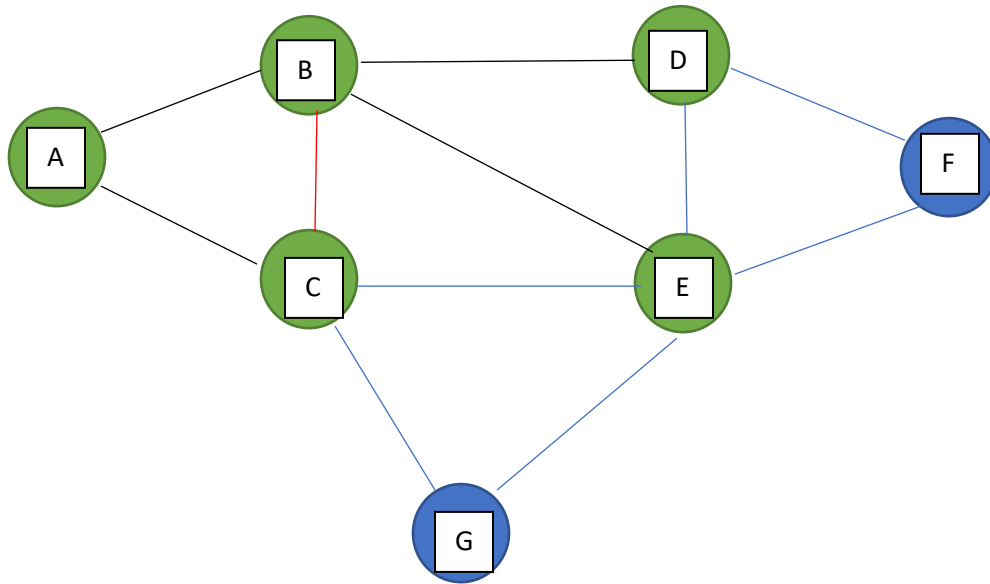
black line= discovery edge

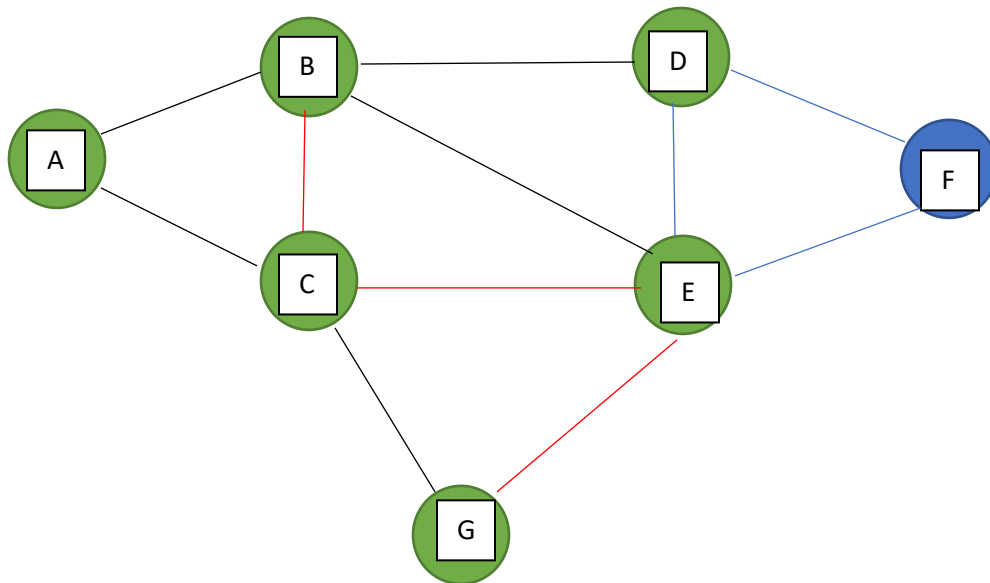
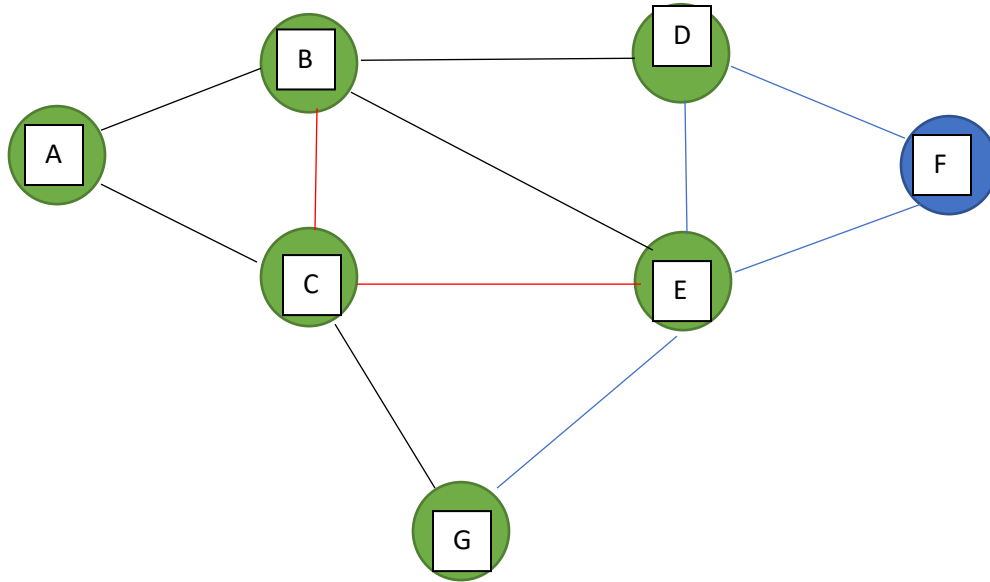
red line= cross edge

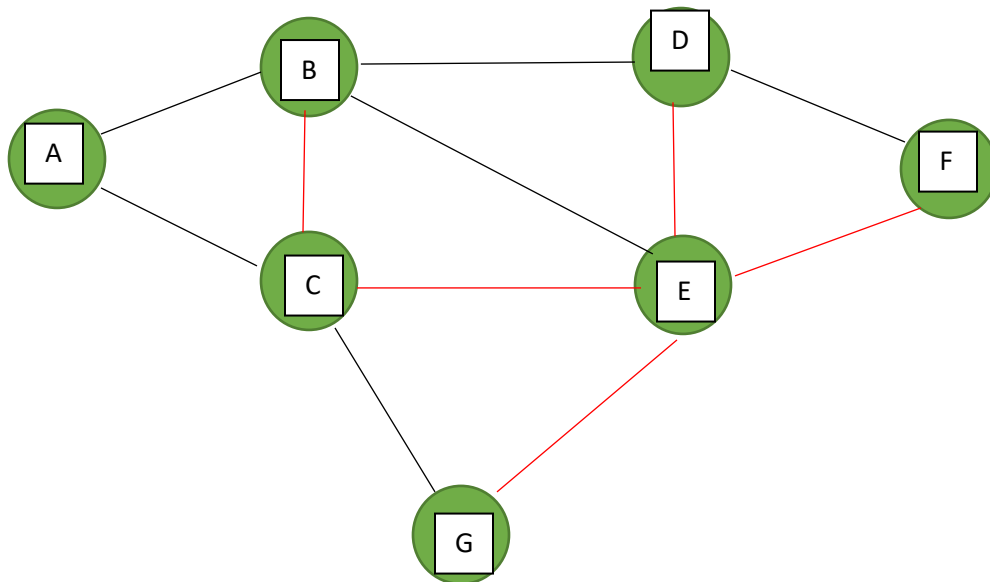
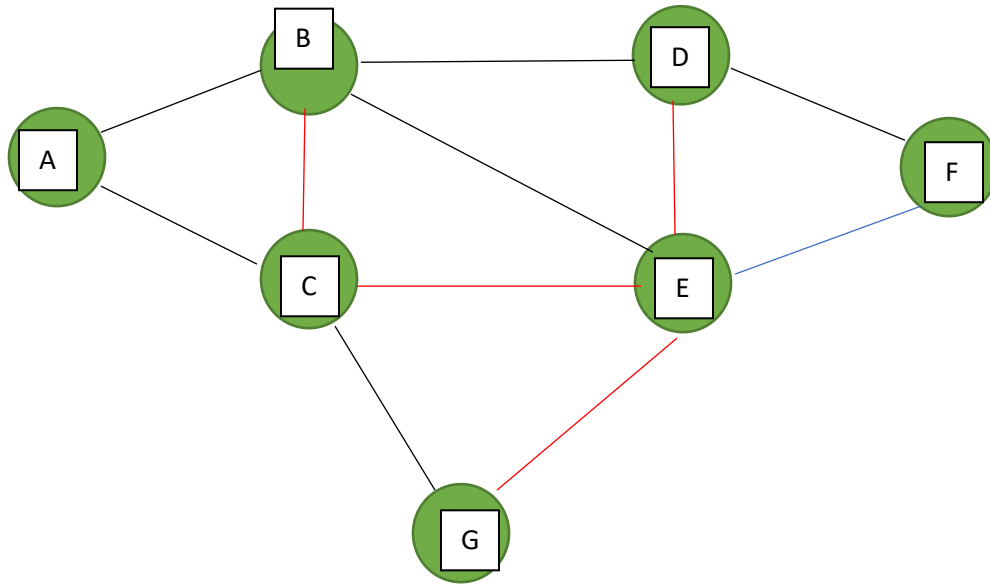












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(b) Trace the execution of DFS beginning at vertex A, labeling each edge as a discovery or back edge.

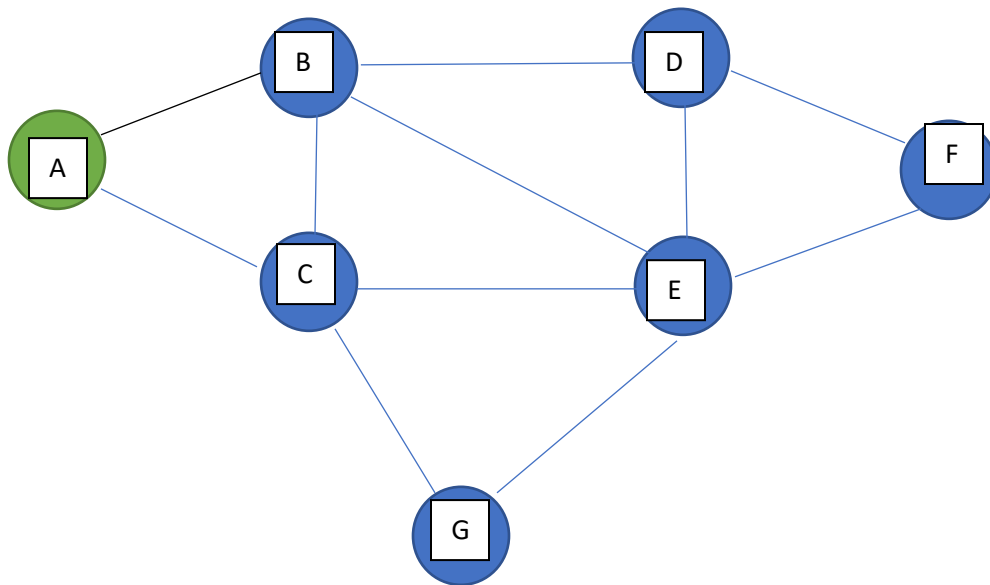
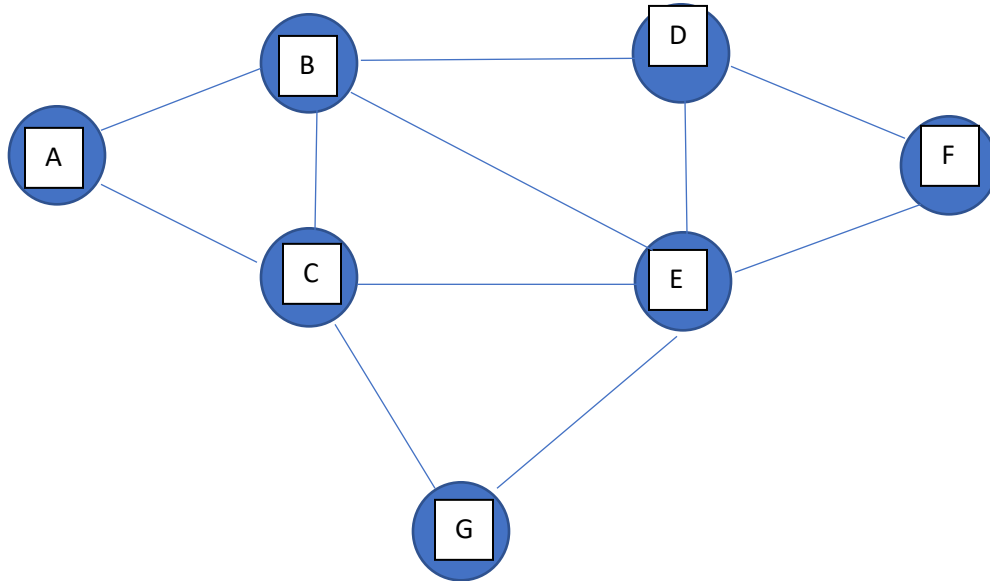
Blue circle= unexplored vertex

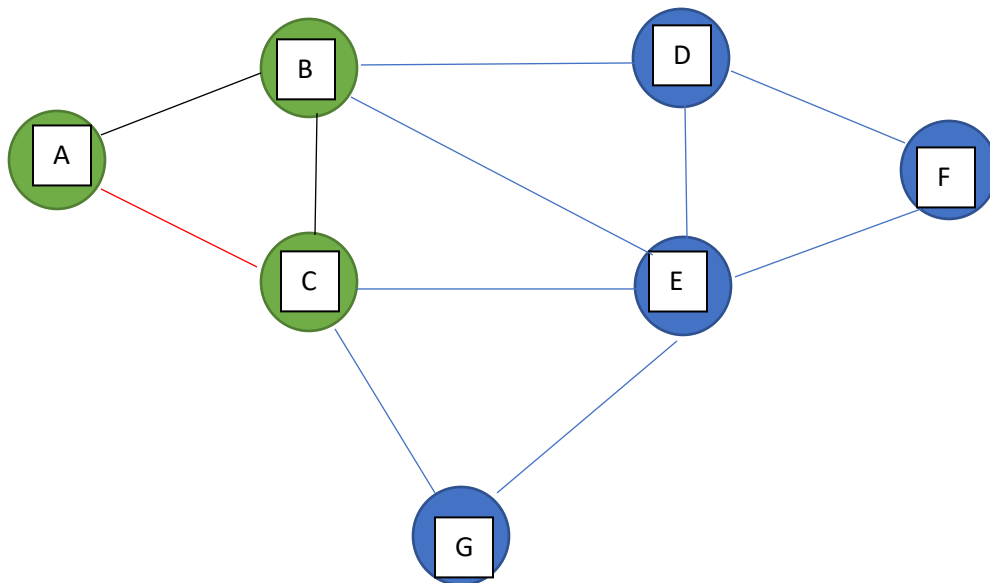
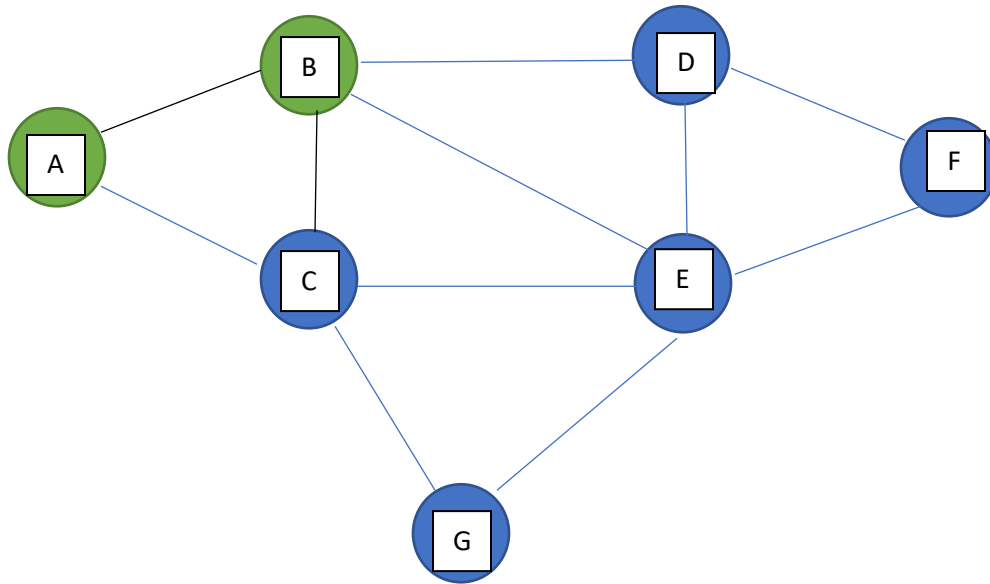
blue line= unexplored edge

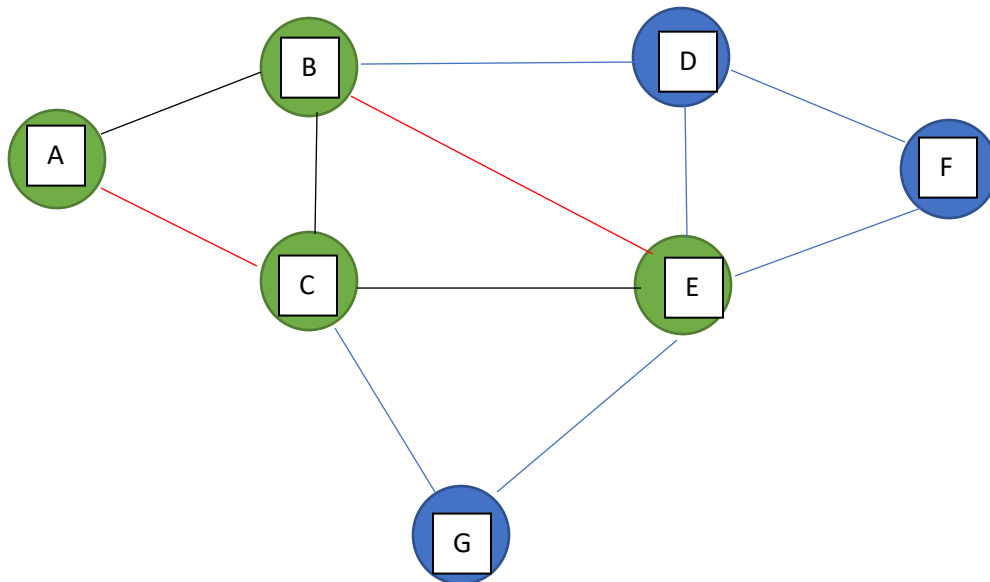
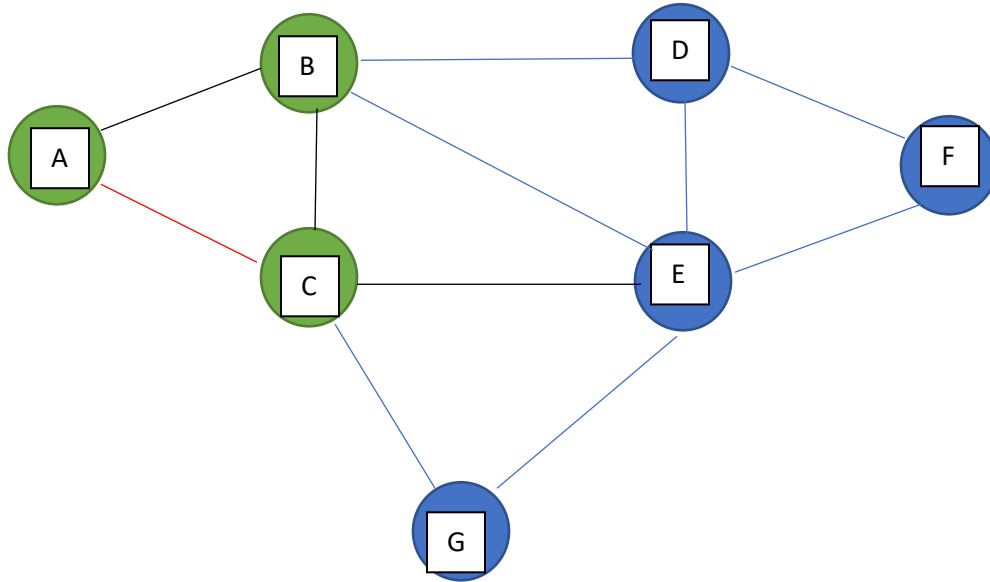
green circle= visited vertex

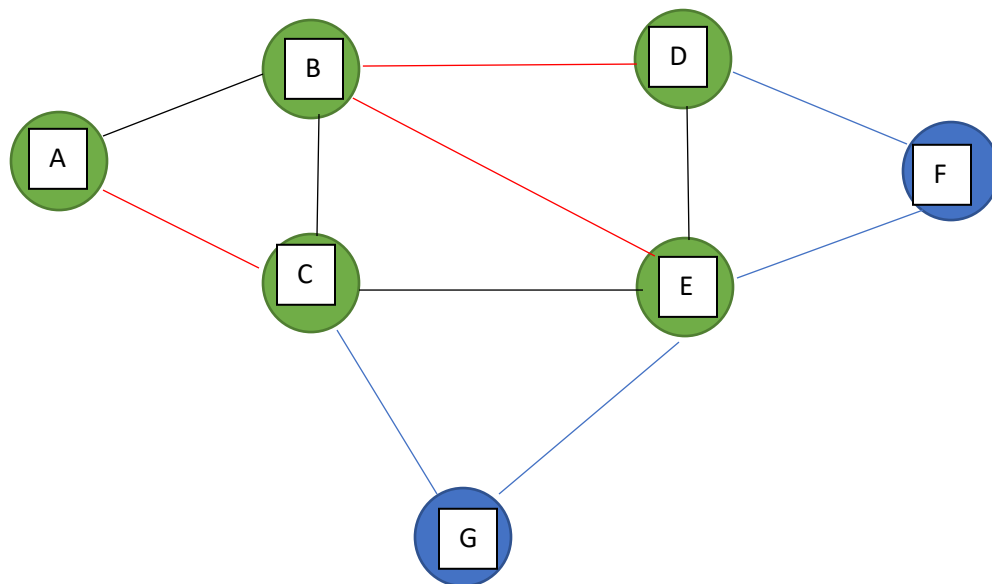
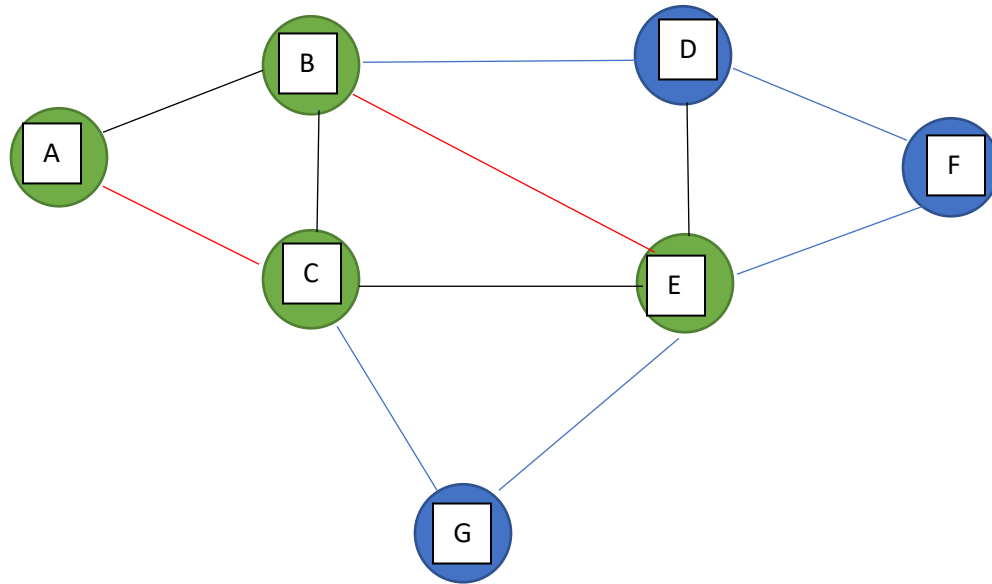
black line= discovery edge

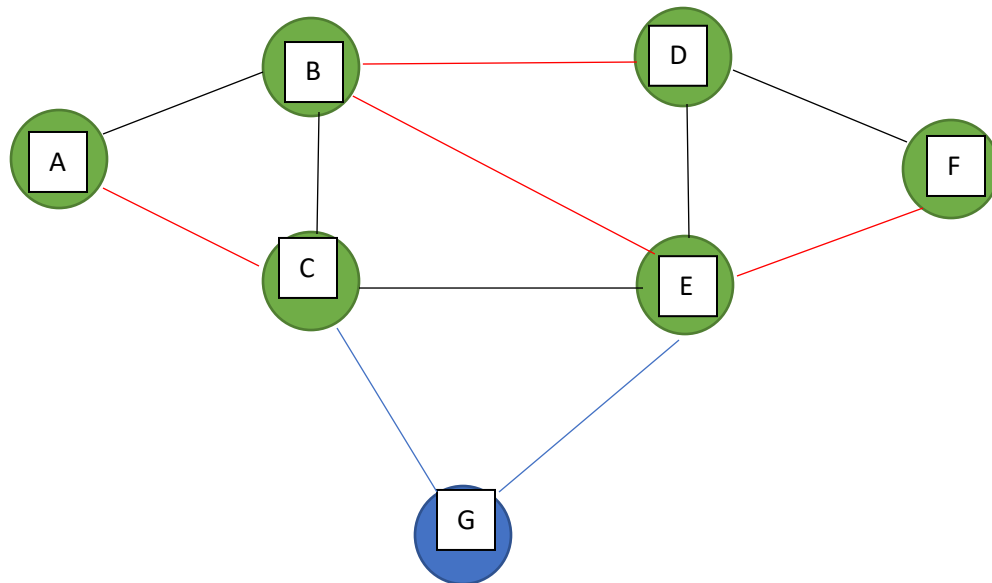
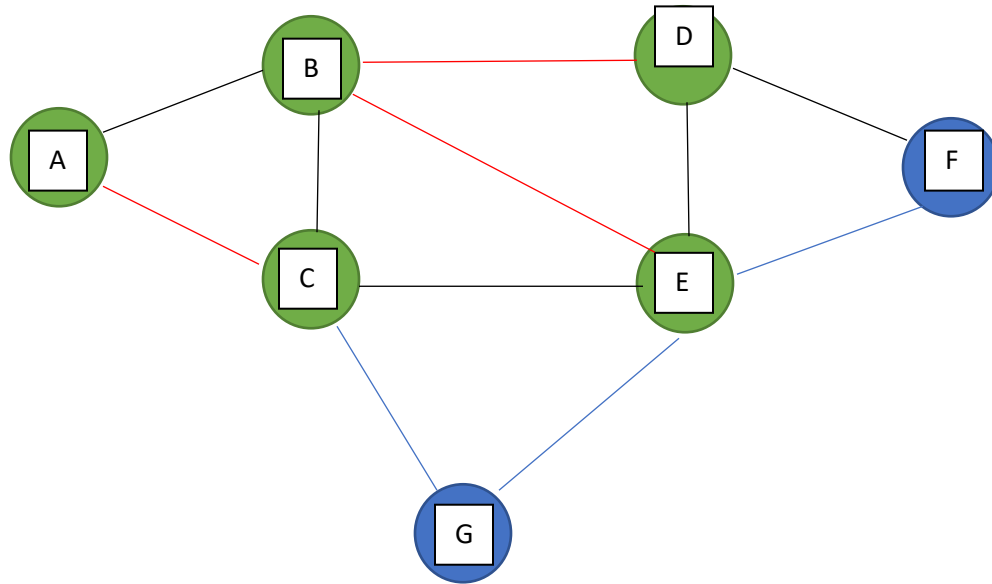
red line= back edge

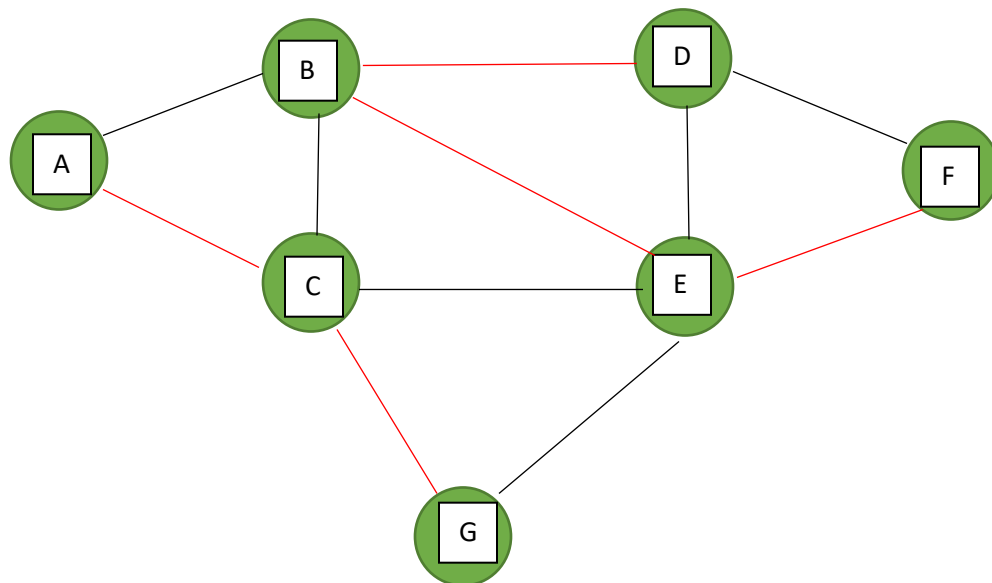
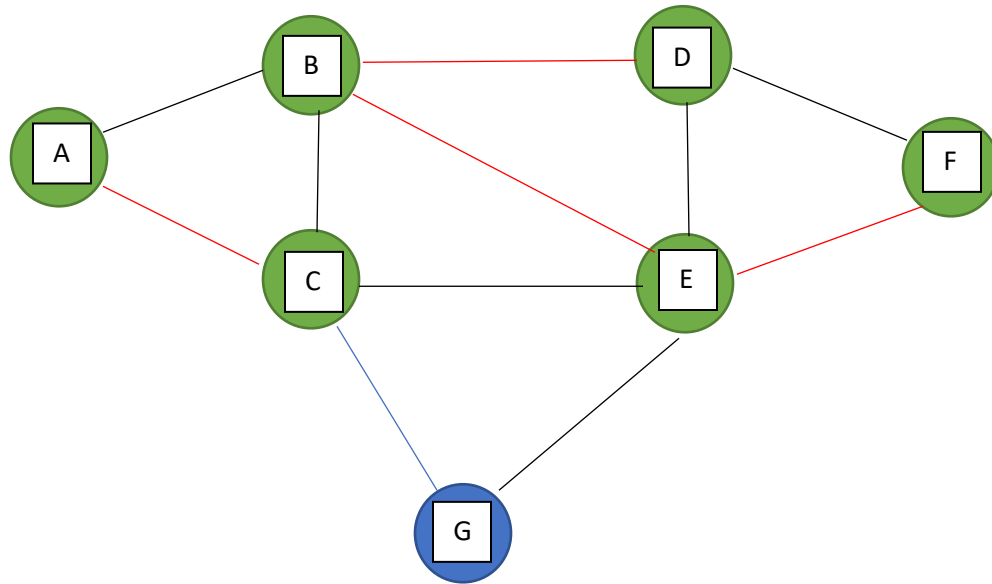












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4. (10 points) In each of the following problems, describe in a few sentences how the task is accomplished via DFS or BFS using only the discovery-edges or back/cross-edges (depending on the searching algorithm).

(a) Find a spanning tree of G , assuming G is connected.

Use dfs on G and remove any back edges from the tree that results. After removing the back edges you should be left with a spanning tree.

(b) Determine if G is acyclic.

Use dfs on G , if there exists no back edges, then there are no cycles and G is acyclic.

(c) Find a path from a vertex u to a vertex v .

Use dfs on (G,u) , find the cycle that includes both u and v . Remove any and all back edges from said cycle, the resulting path should lead you from u to v .

(d) Find a shortest path from a vertex u to a vertex v .

Use bfs on (G,u) , remove to cross edges from the resulting tree. After removing the cross edges you should be left with the shortest path from u to v .

(e) Find all connected components of G .

Use dfs on G , remove any and all back edges from the resulting tree. After removing any and all back edges you should be left with a tree with edges connecting all of the connected components.

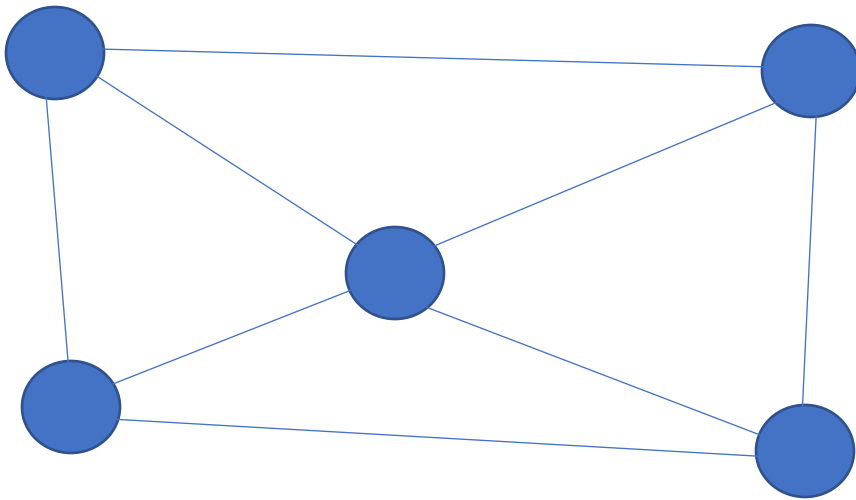
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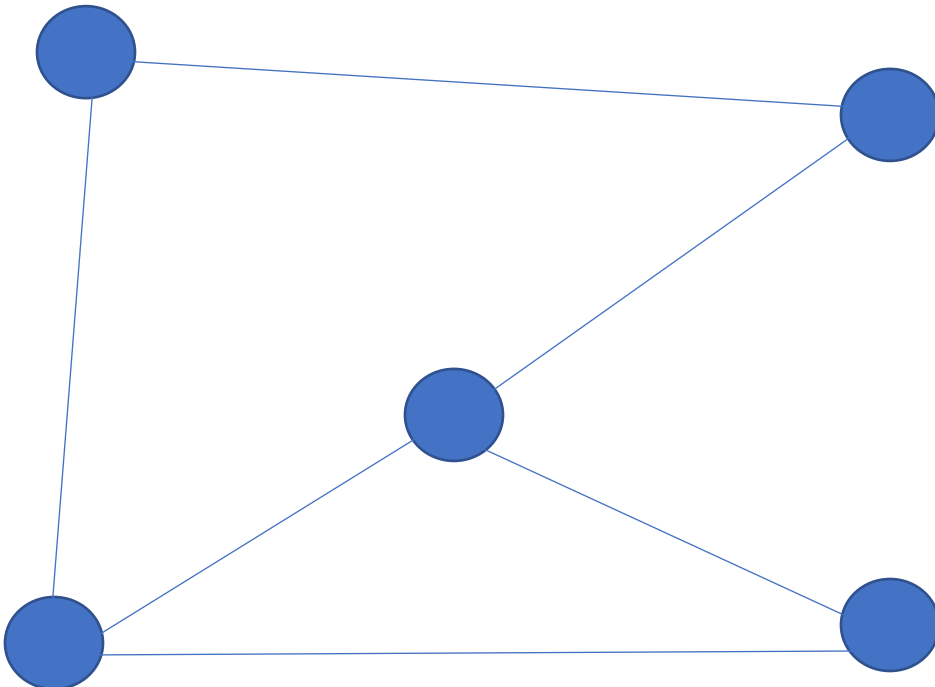
5. (10 points) A graph is triconnected if one has to remove at least 3 vertices from the graph to disconnect it. Construct examples of the following graphs or explain why it cannot be done. Assume the graph is undirected.

(a) A triconnected graph with exactly 5 vertices and 8 edges.

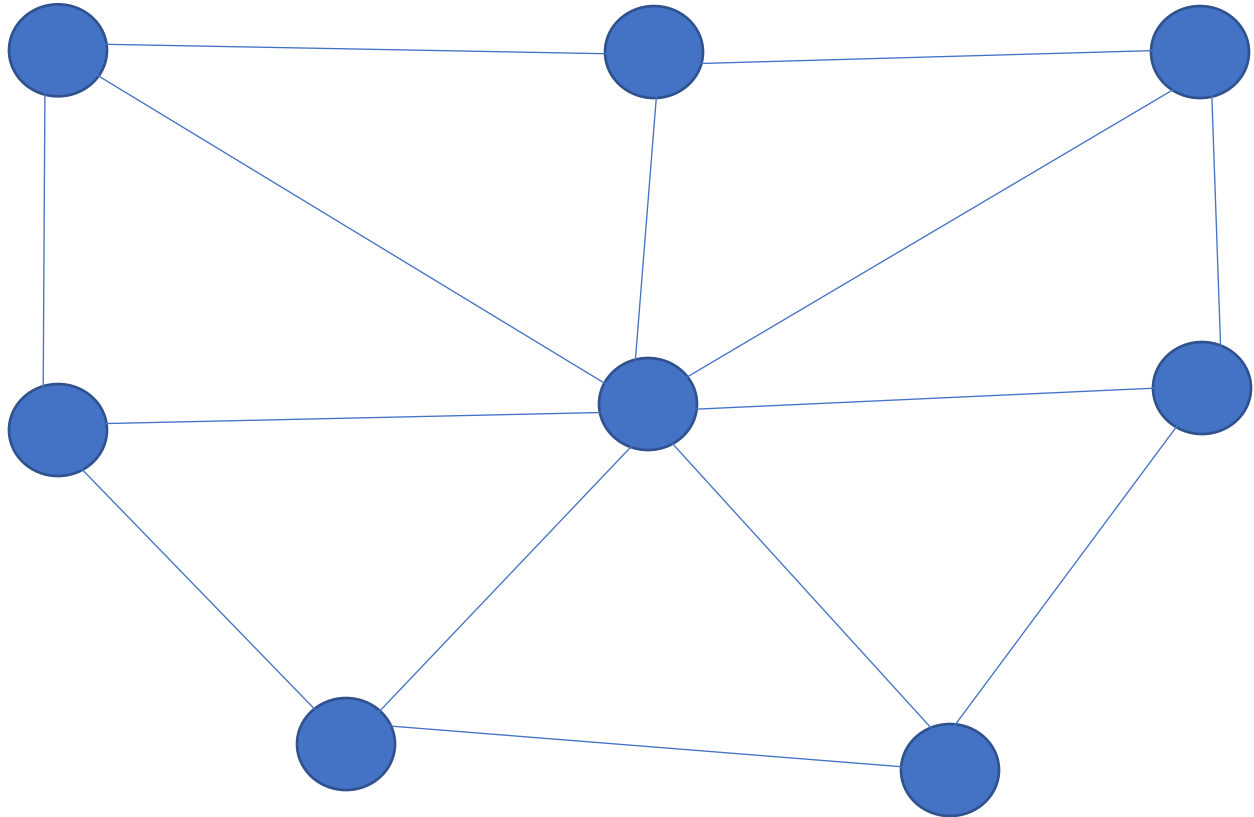


(b) A triconnected graph with exactly 5 vertices and 6 edges.

Such a graph does not exist as the removal of only two vertices would disconnect said graph.



(c) A triconnected graph with exactly 8 vertices and 14 edges.



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6. (5 points) Bob loves foreign languages and wants to plan his course schedule to take the following nine language courses: LA15, LA16, LA22, LA31, LA32, LA126, LA127, LA141, and LA169. The course prerequisites are:

- LA15: (none)
- LA16: LA15
- LA22: (none)
- LA31: LA15
- LA32: LA16, LA31
- LA126: LA22, LA32
- LA127: LA16
- LA141: LA22, LA16
- LA169: LA32

Find a sequence of courses that allows Bob to satisfy all the prerequisites

LA15-> LA22-> LA16-> LA31-> LA32-> LA126->LA169->LA141->LA127

