Course: CS 836

Name: Chao Zhang

Student#: 200383834

1. Give the computational formulation of rough set approximations:

a). Give a definition of information table.

 $T=(OB,AT,\{V_a|a\in AT\},\{I_a|a\in AT\})$ 

where OB is a finite non-empty set of objects which is the universe. AT is a finite non-empty set of attributes.  $V_a$  is the domain of attributes and  $I_a$  is an information function from OB to  $V_a$  which means each object is mapped to exactly one value from objects to attributes.

b). Define an equivalence relation in an information table. In an information table, one can construct an equivalence relation by using a subset of attributes. For a subset of attributes  $A \subseteq AT$ , we define an equivalence relation  $E_A$ :  $xE_Ay \iff \forall a \in A \ (I_a(x) = I_a(y))$ .

c). Define the partition induced by the equivalence relation.

Consider the equivalence relation  $E_A$  induced by a subset of attributes  $A \subseteq AT$ . The equivalence relation  $E_A$  induces a partition  $OB/E_A$  of OB for instance a family of nonempty and pair-wise disjoint subsets whose union is the universe. For an object  $x \in OB$ , its equivalence class is given by:  $[x]_A = \{y \in OB \mid xE_Ay\}$ .

d). Define the atomic Boolean algebra by using the partition as the set of atoms. One can construct an atomic sub-Boolean algebra  $B(U/E_A)$  of  $2^{OB}$  with  $OB/E_A$  as the set of atoms:  $B(OB/E_A)=\{U \ F | F\subseteq OB/E_A\}$ . Each element in  $B(OB/E_A)$  is the union of a family of equivalence classes. The Boolean algebra  $B(OB/E_A)$  contains the empty set  $\emptyset$ , the whole set OB, and is closed under set complement, intersection, and union.

e). Define the rough set approximations.

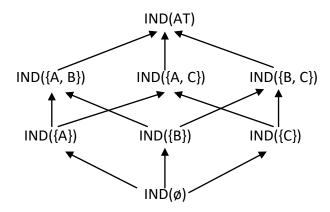
For a subset  $X\subseteq OB$ , the lower and upper approximations of X are defined by:

Granule-based definition:

2. Consider the following information table:

Object	A	В	С
1	$a_1$	$b_1$	$c_1$
2	$a_1$	$b_1$	$c_2$
3	$a_2$	$b_2$	$c_1$
4	$a_2$	$b_3$	$c_1$
5	$a_2$	$b_3$	$c_1$
6	$a_2$	$b_1$	$c_1$
7	$a_2$	$b_3$	$c_2$
8	$a_1$	$b_1$	$c_2$

- a). Construct the family of all equivalence relations defined by subsets of attributes.  $IND(\emptyset)$ ,  $IND(\{A\})$ ,  $IND(\{B\})$ ,  $IND(\{C\})$ ,  $IND(\{A,B\})$ ,  $IND(\{A,C\})$ ,  $IND(\{B,C\})$ , IND(AT)
- b). Draw a diagram to show the connections between those equivalence relations.



IND(AT) includes IND({A, B}), IND({A, C}), IND({B, C})

IND({A, B}) includes IND({A}) and IND({B})

IND({A, C}) includes IND({A}) and IND({C})

IND({B, C}) includes IND({B}) and IND({C})

IND({A}), IND({B}), IND({C}) include IND(Ø)

c). Construct an equivalence relation on the power set of the set of attributes: two subsets of attributes are equivalent if and only if they produce the same equivalence relation on U. Give the partition of the power set of the set of attributes.

```
OB/ IND(ø)={{1,2,3,4,5,6,7,8}}

OB/ IND({A})={{1,2,8},{3,4,5,6,7}}

OB/ IND({B})={{1,2,6,8},{3},{4,5,7}}

OB/ IND({C})={{1,3,4,5,6},{2,7,8}}

OB/ IND({A, B})={{1,2,8},{3},{6},{4,5,7}}

OB/ IND({A, C})={{1},{2,8},{3,4,5,6},{7}}

OB/ IND({B, C})={{1,6},{2,8},{3},{4,5},{7}}

OB/ IND(AT)={{1},{2,8},{3},{4,5},{7}}
```

From above, we can get none of them produce the same equivalence relation on U so there do not exist the equivalence relation on attributes. So the partition of the power set of the set of attributes is  $\{\{\emptyset\},\{A\},\{B\},\{C\},\{A,B\},\{A,C\},\{B,C\},\{A,B,C\}\}\}$ 

3. Show that the conceptual formulation and the computational formulation are equivalent.

In an information table T, for a subset  $X \subseteq OB$ , the lower and upper approximations of X are defined by the following pair of definable sets,  $\overline{aprA}(X)$  is the least definable set in  $DEF_A(T)$  containing X and  $\underline{aprA}(X)$  is the greatest definable set in  $DEF_A(T)$  contained by X. We can easily get  $\underline{aprA}(X) \subseteq X \subseteq \overline{aprA}(X)$ . From the fact that  $[x]_A$  is a definable set and  $\underline{aprA}(X)$  is the greatest definable set contained in X, we can also conclude that  $[x]_A \subseteq X \Longrightarrow [x]_A \subseteq \underline{aprA}(X)$ , so then we combine these two and get  $\underline{aprA}(X) = U\{[x]_A \in OB/E_A \mid [x]_A \in X\}$  and  $\overline{aprA}(X) = U\{[x]_A \in OB/E_A \mid [x]_A \cap X \neq \emptyset\}$ . For  $\underline{aprA}(X) = U\{[x]_A \in OB/E_A \mid [x]_A \in X\}$  which has the same meaning with the greatest definable set in  $DEF_A(T)$  contained by X. And for  $\overline{aprA}(X) = U\{[x]_A \in OB/E_A \mid [x]_A \cap X \neq \emptyset\}$ , where  $[x]_A \cap X \neq \emptyset$  is equal to  $[x]_A \cap \overline{aprA}(X) \neq \emptyset$ , it has the same meaning with the least definable set in  $DEF_A(T)$  containing X. Since the conceptual formulation and the computational formulation have the same meaning, they are equivalent.