

# SOME PORTED, RELATIVE, OR SET POINTED PARAMETRIZED TUTTE FUNCTIONS

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ABSTRACT. Tutte decompositions with deletion and contraction not done of elements in a fixed set of ports  $P$ , and the resulting polynomial expressions and functions for matroids, oriented matroids and labelled graphs are investigated. We give conditions on parameters  $x_e, y_e, X_e, Y_e$  for  $e \notin P$ , and on initial values  $I(Q)$  for indecomposables, that are necessary and sufficient for the following equations to have a well-defined solution:  $T(M) = X_e T(M/e)$  for coloops  $e \notin P$ ,  $T(M) = Y_e T(M \setminus e)$  for loops  $e \notin P$ , and  $T(M) = x_e T(M/e) + y_e T(M \setminus e)$  for other  $e \notin P$ . They generalize similar conditions given by Bollobás and Riordan, Zaslavsky, and Ellis-Monaghan and Traldi for Tutte functions defined without the  $e \notin P$  restriction. We complete the generalization to matroids given Diao and Heteyi which was motivated by invariants for the virtual knots studied by Kauffman. Our motivations include electrical network analysis, oriented matroids, and negative correlation of edge appearances in spanning trees; our  $P$ -ported Tutte polynomials of oriented matroids express orientation information that ordinary Tutte polynomials cannot. We also apply the computation tree formalism of Gordon and McMahon to give activities expansions for  $P$ -ported parametrized Tutte polynomials that are more general than those just determined by the linear element orderings which originated in Tutte's dichromate.

The polynomials expressing conditions for above Tutte equation to have a solution all have one factor  $I(Q_i)$ . Since the elements are labelled, the methodology also applies to some other situations with ports including graphs with labelled vertices for which similar theorems can be proven. Examples also demonstrate cases situations where conditions similar these polynomials are not sufficient for Tutte equations to be solvable.

## 1. INTRODUCTION

The Tutte polynomial is a well-known invariant of graphs and matroids, and various useful generalizations of its theory, defined by attaching parameters to elements, have been found. We generalize by restricting the Tutte equations so that only the elements not in a fixed subset  $P$  can be deleted and/or contracted; see (TA) and (TSSM) below. Our work is largely based on [11], which addresses parametrization in detail and which unifies the main results of Zaslavsky [24] and Bollobás and Riordan [2]. Diao and Heteyi [9] give special cases of our generalization together with its application to virtual knot invariants [17]. Recent work on a different generalization, weak Tutte functions (see sec. 1.1), has been done by Ellis-Monaghan and Zaslavsky [25].

The separator-strong Tutte equations (TA), (TSSM) equations reduce to the well-known identities for the two-variable Tutte polynomial when the parameters are  $x_e = y_e = 1$ ,  $X_e = x$  and  $Y_e = y$  for all  $e$  and the set  $P = \emptyset$ . In other words, we study the effects of distinguishing certain elements, those in  $P$ , so they are never deleted or contracted in the course of Tutte decompositions that carry the  $x_e, y_e, X_e, Y_e$  parameters into the resulting

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polynomials. We generalize some basic theorems about parametrized Tutte functions and about Tutte polynomials, both of matroids and of graphs, to this situation.

We call  $P$  the *set of ports*. The resulting Tutte polynomials or functions have been called *set-pointed* [20, 19], *ported* [5, 7] and *relative* [9]. We prefer the terms “port” and “ $P$ -ported equations or functions” because of our applications [6, 7] and because the  $P$  can be specified.

When parameters  $x_e$  and  $y_e$  are introduced into the additive Tutte equation, the route toward generalizations of the Tutte polynomial becomes complicated. The next subsection (1.1) elaborates on the particular destination—*separator-strong Tutte functions*, to which we add the port restrictions. Here is the main definition with the restriction for  $P$ :

**Definition 1.** *Let  $P$  be a set and  $\mathcal{C}$  be a class of matroids or oriented matroids closed under taking minors or oriented minors by, for  $e \notin P$ , deleting  $e$  if  $e$  is not a coloop or contracting  $e$  if  $e$  is not a loop. Such a  $\mathcal{C}$  is called a  $P$ -family.  $\mathcal{C}$  is given with four parameters in a commutative ring  $R$ ,  $(x_e, y_e, X_e, Y_e)$ , for each  $e \notin P$  that is an element in some  $M \in \mathcal{C}$ .*

*A separator-strong  $P$ -ported Tutte function  $T$  maps  $\mathcal{C}$  to  $R$  or to an  $R$ -module and satisfies conditions (TA) and (TSSM) below for all  $M \in \mathcal{C}$  and all  $e$  in  $M$ .*

$$(TA) \quad T(M) = x_e T(M/e) + y_e T(M \setminus e) \\ \text{if } e \notin P \text{ and } e \text{ is a non-separator, i.e., neither a loop nor a coloop.}$$

$$(TSSM) \quad \text{If } e \notin P \text{ is a coloop in } M \text{ then } T(M) = X_e T(M/e). \\ \text{If } e \notin P \text{ is a loop in } M \text{ then } T(M) = Y_e T(M \setminus e).$$

*The  $M \in \mathcal{C}$  for which neither (TA) nor (TSSM) apply are called indecomposables and the  $P$ -quotients in  $\mathcal{C}$ . The analogous definitions are used for graphs and directed graphs.*

Except in remarks about peripheral work, all Tutte polynomials, functions and equations in this paper refer to the separator-strong kind, except when they are qualified to be strong (see TSM and sec. 5-6).

Evidently, (TA) and (TSSM) specify how  $T(M)$  can be recursively calculated from the values of  $T$  on indecomposables. The definition that  $T$  is a function means that all calculations of  $T(M)$  for  $M \in \mathcal{C}$  using (TA) and (TSSM) yield the same result. A simple induction on  $|E(M)|$ , the number of elements in  $M$  that are not in  $P$ , shows that if a Tutte function with specified values on the indecomposable matroids or oriented matroids exists for the given  $P$ -family and parameters, then the function is unique.

Restricting Tutte decomposition operations to  $e \notin P$  gives some new information about oriented matroids and oriented (i.e., directed) graphs. The fundamental consequence of the restriction is that the indecomposable matroids (or graphs) are minors of  $M$  that have all their elements (or edges) in  $P$ . We will therefore also call them  $P$ -quotients. If  $M$  is an oriented matroid, each indecomposable is an oriented matroid because the oriented minor  $M/A \setminus B$  is well-defined when  $A, B$  partitions  $E$ . Hence, when  $P \neq \emptyset$ ,  $P$ -ported Tutte functions can have different values on different orientations of the same orientable matroid. On the other hand, we can define Tutte functions of matroids obtained by forgetting the orientations of oriented matroids. Many of our results will be stated for “matroids or oriented matroids” because there are different indecomposables and Tutte functions depending on whether or not the matroids carry an orientation. Analogous statements of course apply to graphs and directed graphs.

We presented in [7] a new kind of strong Tutte-like function on  $P$ -ported oriented graphic matroids (more generally, unimodular, i.e. regular oriented matroids) where the values vary with the orientation. Each function value  $F(G)$  is in the exterior algebra over  $R^{2p}$ , where  $R$  is the real polynomial fraction field with the  $x_e, y_e$  parameters as variables, and  $|P| = p$ . So formulated, the function obeys an *anti-commutative* variant of (TSM) below with exterior multiplication  $\wedge$ . (When  $P = \emptyset$ , the function reduces to the reduced Laplacian determinant in the famous Matrix Tree Theorem [16, 4].) It is the first example we know of “the possibility of making use of a noncommutative generalization of the Tutte polynomial at some point in the future.” mentioned by Bollobás and Riordan in [2]. We won’t say more beyond that each of the  $\binom{2p}{p}$  Plücker coordinates of  $F(G)$  is a  $P$ -ported Tutte function of the kind we cover here.

**1.1. Complications from Parametrization.** It is known that, even when  $P = \emptyset$ , (TA) fails to have a solution  $T$  except if certain algebraic relations are true about the parameters and the so-called initial values of  $T$ .

For example, if  $M = U_1^{ef}$  is  $e, f$  in parallel, then applying (TA) for  $e$  gives the polynomial ( $U_0^e, U_1^e$  are the loop, coloop matroids on  $\{e\}$ , etc.)

$$x_e T(U_0^f) + y_e T(U_1^f)$$

whence applying it for  $f$  gives

$$x_f T(U_0^e) + y_f T(U_1^e).$$

These are different polynomials in the parameters and initial values. The equation that says they are equal is an example of a relation that is necessary for a solution to exist. The example demonstrates that when the Tutte identities are parametrized, it is important to carefully distinguish between a solution value  $T(M)$ , where a solution  $T$  is a function that satisfies all the relevant identities, and a formal polynomial that results from using a subset of the identities to calculate  $T(M)$  for one  $M$ . We will in section 4 investigate using *Tutte computation tree* [15] to express the activities expansions for all such polynomials obtained by recursion, with any port set  $P$ .

Generally speaking, given a system of equations that, like Tutte equations, specifies a function value  $T(M)$  with recursive reductions, for the system to “have a solution” means that all recursive calculations of  $T(M)$  from the equations have the same result. We then say that  $T(M)$  is *well-defined*. The above example shows calculations of two different polynomials in the parameters and initial values. For this reason, we follow [24, 11] and use the term “Tutte function” for a solution to (separator-strong) Tutte identities, rather than Tutte polynomial. The term “Tutte polynomial” is reserved for certain polynomial expressions that equal the Tutte function value if additional conditions are satisfied. We follow [2, 11] to say  $T(M)$  is well-defined when the parameters are in ring  $R$  and the initial values are in  $R$  or an  $R$ -module for which the polynomial expressions for  $T(M)$  obtained by all the recursive applications of the relevant Tutte equations are equal. The conditions are conveniently expressed as generators for the ideal  $I$  such that the universal Tutte function is into a quotient ring or module modulo  $I$ . See Corollary 7.

Besides (TA), classical Tutte polynomials and other so-called *strong Tutte functions* [24] satisfy the multiplicative identity

$$(\text{TSM}) \quad T(M_1 \oplus M_2) = T(M_1)T(M_2).$$

Confusion between Tutte polynomials and Tutte functions arises because the classical two-variable Tutte polynomial in  $x, y$  expresses a strong, non-parametrized Tutte function for every assignment of values to  $x$  and  $y$ .

Different combinations of variations of Tutte equations determine different kinds of Tutte functions. For example, the additive identity (TA) alone (with  $P = \emptyset$ ) characterizes *weak Tutte functions* [24, 25].

Our generalization is the kind of Tutte function determined by (TA) and (TSSM); the latter is the weakening of (TSM) so it applies only for cases where one of  $M_1$  or  $M_2$  is a *separator*, that is, a loop or coloop. We generalize by adding the restriction on port elements. The result, Theorem 4 (sec. 3), is a straightforward generalization of Theorem 2 below paraphrased from [11] about the existence and universal form for the family of Tutte functions characterized by (TA) and (TSSM). This family was subsequently named *separator-strong Tutte functions* by [12] because it is wider than the strong Tutte functions characterized by (TA) and (TSM), which are the subject of [24]. We follow these authors' terminology when we extend the family to matroids, oriented matroids and graphs with distinguished port elements, and then to ported objects representing matroids or oriented matroids. The conclusions for  $P$ -ported parametrized strong Tutte functions easily follow from those for the separator-strong ones. See 30 in sec. 5.

We build upon [11] which reconciles the results of [24] and [2] with a common generalization. It generalizes the fields and strong Tutte functions of [24] to the commutative rings and separator-strong Tutte functions of [2], and the definedness on all matroids in [2] to definedness on a minor-closed class in [24]. Further, neither the matroid or graph Tutte functions need to be 0 or 1 on  $\emptyset$ . A main result of [11] is called the Zaslavsky-Bollobás-Riordan (ZBR) theorem for matroids:

**Theorem 2.** *Let  $R$  be a commutative ring, let  $\mathcal{C}$  be a minor-closed class of matroids defined on subsets of an  $R$ -parametrized class  $U$ , and let  $\alpha \in R$ . Then there is a parametrized Tutte polynomial on  $\mathcal{C}$  with  $T(\emptyset) = \alpha$  if and only if the following identities are satisfied.*

- a *Whenever  $e$  and  $f$  are digonal in  $M \in \mathcal{C}$  (i.e., they constitute a two element circuit),*

$$\alpha \cdot (x_e Y_f + y_e X_f) = \alpha \cdot (x_f Y_e + y_f X_e).$$

- b *Whenever  $e, f$  and  $g$  are triangular in  $M \in \mathcal{C}$  (i.e. they constitute a three element circuit),*

$$\alpha \cdot X_g \cdot (x_e Y_f + y_e x_f) = \alpha \cdot X_g \cdot (Y_e x_f + x_e y_f).$$

- c *Whenever  $e, f$  and  $g$  are triadic in  $M \in \mathcal{C}$  (i.e. they constitute a three element cocircuit),*

$$\alpha \cdot Y_g \cdot (x_e Y_f + y_e x_f) = \alpha \cdot Y_g \cdot (Y_e x_f + x_e y_g).$$

The terminology “ $R$ -parametrized class  $U$ ” (of matroid elements) was used in [11] to emphasize that the parameters  $(x_e, y_e, X_e, Y_e)$  are attached to elements  $e$ , not equations. The assumption that  $\mathcal{C}$  is minor-closed implies that the three conditions can be restricted to the pair or triple being the only elements in  $M$ . Our generalization, Theorem 4 (sec. 3) is expressed that way.

The one indecomposable for the ZBR theorem is the empty matroid  $\emptyset$ . Loop or coloop matroids on  $e$  are decomposable; the values of the Tutte function on them equal respectively  $Y_e T(\emptyset)$  and  $X_e T(\emptyset)$ .

**1.2. Summary.** We first generalize to  $P$ -ported matroids and oriented matroids the Bollobás-Riordan-Zaslavsky theorem as synthesized by Ellis-Monaghan and Traldi in [11]. Parameters are only attached to  $e \notin P$ .

In our generalization of Theorem 2, it is easy to see that the indecomposables are those matroids  $Q_i$  in  $\mathcal{C}$  whose ground set  $S(Q_i) \subseteq P$ . When  $P \neq \emptyset$ , instead of every monomial resulting from a Tutte decomposition containing the factor  $T(\emptyset)$ , every monomial resulting from a Tutte decomposition of  $M$  contains a factor  $T(Q_i)$  where  $Q_i$  is some minor of  $M$  obtained by contracting or deleting every  $e \notin P$ . These *initial* values must be given for those  $Q_i$  in order to define a particular Tutte function. They generalize the  $\alpha = T(\emptyset)$  value.

Our generalization of the ZBR theorem first replaces  $T(\emptyset) = \alpha$  in its three equations with  $I(Q_i)$ . It also adds two more equations pertaining to  $M \in \mathcal{C}$  in which series and parallel pairs  $\{e, f\}$ ,  $e \notin P$ ,  $f \notin P$  are connected to one or more elements of  $P$ . Again, only one  $Q_i$  appears in each equation. See Theorem 4. We let some Tutte functions take values in modules to facilitate both defining universal Tutte functions with quotient modules and expressing formulas for Tutte functions of direct sums and related combinations.

An interesting consequence is that no relationships are required between  $T(Q_i)$  and  $T(Q_j)$  for different indecomposables  $Q_i$  and  $Q_j$  for the Tutte function to be well-defined. This answers a question we raised in [7].

Second, we extend to  $P \neq \emptyset$  the activities Tutte polynomial expressions and corresponding interval partitions of the Boolean lattice  $2^{S(M) \setminus P}$ . Like the activities expressions given by [15] for greedoid Tutte polynomials, the expressions we present are not just those obtained when a fixed linear ordering on the elements is used to determine for which element  $e$  to apply (TA) when two or more elements are eligible. Each of our activities expressions is based on a formal (Tutte) *computation tree* as defined in [15].

Third, we describe the relationship between the relevant separator-strong Tutte function values when a matroid is a direct sum. It is more complicated than the product formula  $T(M^1 \oplus M^2)T(\emptyset) = T(M^1)T(M^2)$  for matroids and related formulas for graphs [11] because a  $P$ -ported Tutte function value can involve more than one indecomposable. When the module that contains the universal Tutte polynomials is extended to an algebra by defining direct sum and similar compositions of indecomposables by the new algebra's multiplication  $*$ , the relationship can be expressed by  $T(M^1 \oplus M^2) = T(M^1) * T(M^2)$ . Families of objects with matroids, where deletion and contraction is consistent with the matroids and restricted to  $e \notin P$ , are introduced to facilitate studying  $P$ -ported strong Tutte functions of graphs.

Finally, we specialize to graphs and so extend some of the graph results of [11]. Additional hypotheses are given for  $P$ -ported Tutte functions of graphs, including those with labelled vertices, for characterizations like Theorems 2 and 4 to be true.

**1.3. Background and Other Related Work.** The topic of Tutte decomposition restricted so some elements are not deleted or contracted probably made its first appearance in the early work of Thomas Brylawski[3]. He formulated a Tutte polynomial for “pregeometries with basepoint  $p_0$ ” with the four variables  $z$ ,  $x$ ,  $z'$  and  $x'$ . His polynomial is obtained from (PAE) in sec. 4 with the substitutions  $P = \{p_0\}$ ,  $x_e = y_e = 1$ ,  $X_e = z$ ,  $Y_e = x$  for all  $e \neq p_0$ , and  $[U_1^{p_0}] = z'$  and  $[U_0^{p_0}] = x'$ . Another early appearance is [21].

Ellis-Monaghan and Traldi [11] explain that by leaving the reduction by  $e_0$  to last so  $e_0$  is always contracted as a coloop or deleted as a loop, the Tutte function value can be expressed by  $T(M) = (rX_{e_0} + sY_{e_0})T(\emptyset)$  where  $r, s$  are not-necessarily-unique elements in  $R$ . As one application, they give a formula for the parametrized Tutte polynomial for the

parallel connection across  $e_0$  which generalized Brylawski's work. These  $r, s$  appear in the  $P$ -ported Tutte function expression  $rT(U_1^{e_0}) + sT(U_0^{e_0})$  when  $P = \{e_0\}$ . They are parametrized generalizations of the coefficients of  $z'$  and  $x'$  in Brylawski's four variable Tutte polynomial.

Las Vergnas defined and gave basic properties of “set-pointed” Tutte polynomials (with no parameters) and used them to study matroid perspectives. The polynomial given in [18, 20] has a variable  $\xi_l$  for each subset in a collection of  $k$  subsets  $P_l \subseteq P$ ,  $l = 1, \dots, k$ . Each term in (PAE) had  $\prod \xi_l^{r_i(P_l)}$  for  $[Q_i]$  where  $r_i$  is the rank function for  $Q_i$ . Therefore (TSM) was satisfied and the association of the term to (non-oriented) matroid  $Q_i$  could be assured by taking all  $2^{|P|}$  subsets for the  $P_l$ . The matroid perspective is the strong map  $M \setminus E(M) \rightarrow M/E(M)$  given by the identity on  $P$ .

In [5], we reproduced Las Vergnes' theory with explicit  $P$ -quotient (matroid) variables (see the  $[Q_i]$  symbols in Corollary 7 in sec. 3.3) in place of  $\prod \xi_l^{r_i(P_l)}$ . We then gave formulas for the  $P$ -ported Tutte polynomial for the union and its dual of matroids whose common elements are in  $P$ . These formulas work in a way similar to what appears in sec. 3.3. We extend to algebras the  $\mathbb{Z}[u, w]$ -module generated by the  $[Q_i]$  by defining multiplications  $\tilde{*}$  with the rules  $[Q_i]\tilde{*}[Q_j] = r_{i,j}[Q_{i,j}]$ , with  $r_{i,j} \in \mathbb{Z}[u, w]$  and  $Q_{i,j} = Q_i * Q_j$  depending on  $(Q_i, Q_j)$  and whether  $*$  represents union or its dual. It is not often recognized that series and parallel connection of matroids across basepoint  $p$  is equivalent to matroid union and its dual on matroids with only element  $p$  in common. We plan to investigate whether the formulas for parametrized Tutte polynomials of parallel connections in [11] can be generalized to the dual of union when  $|P| > 1$ , and to detail the relationship when  $|P| = 1$ .

We introduced  $P$ -ported parametrized Tutte polynomial for normal Tutte functions in [7]. Most of the results in the current paper, when restricted to *normal*  $P$ -ported parametrized Tutte functions (those with corank-nullity polynomial expressions), appeared in [7] or can be derived by adding parameters and oriented matroid considerations to material in [5]. These include computation tree[15] based activities expansions with terms corresponding to  $P$ -subbases (which are called “contracting sets” in [9]). In the normal case, the indecomposables can be assigned arbitrarily. We used this to show the our extensor-valued Tutte-like function [5] is expressible by assigning extensors as the initial values. In this electrical network application, the indecomposables are oriented matroids and different values *are* assigned to different orientations of the same underlying matroid.

We had left open questions about when non-normal  $P$ -ported parametrized Tutte functions are well-defined. They are whether arbitrary values can be assigned to the indecomposables and what is the appropriate generalization of conditions on the parameters given by Zaslavsky [24], Bollobás and Riordan [2] and Ellis-Monaghan and Traldi [11]. Theorem 4 resolves these questions: The new conditions are obvious revisions of those for  $P = \emptyset$ . The indecomposables can be assigned arbitrarily so long as for each one separately, the conditions of Theorem 4 are satisfied. In addition, if the function is to be strong, Theorem 30 tells us that it is sufficient for the assignment to indecomposables be strong.

Diao and Heteyi gave conditions on the parameters, similar to ours, for the Tutte polynomial to be well-defined for every assignment of values on the indecomposables that obeys a symmetry condition. That condition is motivated by the graph specialization. They gave the very natural application to invariants of virtual knots calculated from their diagrams. Zero edges (what we call port edges) correspond to the virtual crossings, and the sign parameters of the other  $\pm 1$  edges derive from left-over or right-over sense of the regular crossings. Their preprint in fact motivated us to pursue the current topic[9]. This topic leads us to

ask if the  $P$ -ported objects with matroids abstraction, and its related Tutte computation tree expansions for parametrized Tutte functions (sec. 5), can be usefully applied to objects besides graphs or directed graphs, such as various kinds of knot diagrams.

Another open project is to classify the solutions to the conditions of Theorem 4 (about separator-strong  $P$ -ported Tutte functions) and Theorem 30 (about the strong ones) for rings and for fields along the lines of [24] and [2].

## 2. PRELIMINARIES

For a matroid or oriented matroid  $M$ , the ground set is denoted by  $S(M)$  and the rank function is denoted by  $r$ . Given a set  $P$ ,  $S(M) \setminus P = \{e \in S(M) \mid e \notin P\}$  is denoted by  $E(M)$ . The elements of  $P$  are called *ports*, and a matroid or oriented matroid given with a set of ports  $P$  is called  *$P$ -ported*. Analogous terminology is used with graphs or directed graphs; then, the edges are the elements.

A  *$P$ -family* is a collection  $\mathcal{C}$  of matroids or oriented matroids such that given  $M \in \mathcal{C}$  and  $e \in E(M)$ , the contraction  $M/e \in \mathcal{C}$  if  $e$  is not a loop in  $M$  and the deletion  $M \setminus e \in \mathcal{C}$  if  $e$  is not a coloop in  $M$ . The set of non-port elements that occur in  $\mathcal{C}$  is denoted by  $E(\mathcal{C})$ . Specifically  $E(\mathcal{C}) = \{e \mid e \in E(M) \text{ for some } M \in \mathcal{C}\}$ . It is straightforward to extend these definitions to families of objects, such as graphs, where each member object has an associated matroid on appropriate elements of the object, such as edges, and those elements can be deleted and/or contracted consistently with the matroid.

All the minors of  $M$  obtained by deleting or contracting zero or more non-port elements are called  *$P$ -minors*. Thus, a  *$P$ -family* is a  *$P$ -minor closed* collection of matroids or oriented matroids. The  *$P$ -minors*  $Q_i$  for which  $S(Q_i) \subseteq P$ , i.e., those with no non-port elements, are called the  *$P$ -quotients* of  $M$ . We say a  *$P$ -quotient* belongs to  $\mathcal{C}$  if it is a  *$P$ -quotient* of some  $M \in \mathcal{C}$ . Note that if  $P$  is finite, (and the objects are just matroids or oriented matroids) there are only a finite number of  *$P$ -quotients* because there are only a finite number of matroids or oriented matroids over subsets of  $P$ .

As usual, a *separator* is an element that is a loop, or is a coloop, i.e., an isthmus in a graph.

Let  $R$  be a commutative ring with multiplicative and additive identities 1 and 0 respectively. We sometimes assume that a  *$P$ -family*  $\mathcal{C}$  comes equipped with four parameters  $x_e, y_e, X_e, Y_e \in R$  for each  $e \in E(\mathcal{C})$  and one initial value  $I(Q_i) \in R$  or in an  $R$ -module for each  *$P$ -quotient*  $Q_i$  in  $\mathcal{C}$ . Note that whether or not the empty matroid  $\emptyset$  is a  *$P$ -quotient* depends on  $\mathcal{C}$ . For example, if  $M \in \mathcal{C}$  with  $S(M) \cap P = \emptyset$  then  $\emptyset$  certainly is a  *$P$ -quotient*. In that case, (TSSM) specifies that  $T(M) = X_e I(\emptyset)$  or  $T(M) = Y_e I(\emptyset)$  if  $e \notin P$  is a separator. Therefore, we consider the  $X_e$  and  $Y_e$  to be parameters because  $X_e$  and  $Y_e$  are values of the Tutte function only if  $\emptyset \in \mathcal{C}$  and  $I(\emptyset) = 1$ . The Tutte equations (TA) and (TSSM) justify calling the  *$P$ -quotients* indecomposables.

The uniform matroid with elements  $\{e, f, \dots\}$  and rank  $r$  is denoted by  $U_r^{ef\dots}$ .

In the remainder, we consider only elements  $e, f, g$  none of which are in  $P$ .

Two distinct elements  $e, f$  in matroid  $M$  are *parallel* when every cocircuit that contains one of them also contains the other. This is equivalent to  $\{e, f\}$  being a two-element circuit. They are *series* when every circuit that contains one of them also contains the other. This is equivalent to  $\{e, f\}$  being a two-element cocircuit. They are called a *dyad* when they are both parallel and series. Note that every dyad is a connected component of  $M$ .

Two distinct elements  $\{e, f\}$  are a *parallel pair connected to  $P$*  when they are parallel and there is a cocircuit of the form  $\{e, f\} \cup P'$  with  $\emptyset \neq P' \subseteq P$ .

Two distinct elements  $e, f$  are a *series pair connected to  $P$*  when they are series and there is a circuit of the form  $\{e, f\} \cup P'$  with  $\emptyset \neq P' \subseteq P$ .

Three distinct elements  $e, f, g$  are called a *triangle* when they comprise a 3 element circuit  $U_2^{efg}$  that is a connected component of  $M$ .

Three distinct elements  $e, f, g$  are called a *triad* when they comprise a 3 element cocircuit  $U_1^{efg}$  that is a connected component of  $M$ .

The following is critical to the proof that the generalizations of identities in Theorem 2 all have the form  $I(Q_i) \cdot r = 0$  where  $r$  is a polynomial in the  $x_e, y_e, X_e, Y_e$  parameters and  $Q_i$  is one  $P$ -quotient.

**Proposition 3.** *Suppose  $e, f$  are in series, or are in parallel, in matroid or oriented matroid  $M$ .*

- (1) *The minors  $M/e \setminus f = M/f \setminus e$  are equal as matroids.*
- (2) *If  $M$  is oriented, the oriented minors  $M/e \setminus f = M/f \setminus e$  are equal as oriented matroids.*

*Proof.* In the following, take all matroids as oriented or not depending on how  $M$  is given.

If  $e, f$  are in series, note that  $M/e \setminus f = M \setminus f/e$ .  $e$  is a coloop in  $M \setminus f$ , so  $M \setminus f/e = M \setminus \{e, f\}$ , which is clearly the same matroid or oriented matroid if  $e, f$  are interchanged. The relevant theory for oriented matroids can be found in XXXXXXXXXXXXXXXX of [1].

If  $e, f$  are in parallel,  $e, f$  are in series in the matroid or oriented matroid dual  $M^*$  of  $M$ . By the first case,  $M^* \setminus e/f = M^* \setminus f/e$  as matroids or as oriented matroids. Thus  $M/e \setminus f = (M^* \setminus e/f)^* = (M^* \setminus f/e)^* = M/f \setminus e$  as matroids or as oriented matroids.  $\square$

### 3. PARAMETRIZED PORTED TUTTE FUNCTIONS

Let  $P$  be a set and  $\mathcal{C}$  be a  $P$ -family of matroids or oriented matroids. In this section, we state, discuss and prove this generalization of Theorem 2 of [11]:

**Theorem 4.** *The following two statements are equivalent.*

- (1)  *$T$  from  $\mathcal{C}$  to  $R$  or an  $R$ -module is a  $P$ -ported separator-strong parametrized Tutte function with  $R$ -parameters  $(x, y, X, Y)$  whose values  $T(Q_i)$  on  $P$ -quotients  $Q_i \in \mathcal{C}$  are the initial values  $I(Q_i)$ .*

- (2) (a) *For every  $M = U_1^{ef} \oplus Q_j \in \mathcal{C}$  with  $P$ -quotient  $Q_j$  ( $U_1^{ef}$  is a dyad),*

$$I(Q_j)(x_e Y_f + y_e X_f) = I(Q_j)(x_f Y_e + y_f X_e).$$

- (b) *For every  $M = U_2^{efg} \oplus Q_j \in \mathcal{C}$  with  $P$ -quotient  $Q_j$  ( $U_2^{efg}$  is a triangle),*

$$I(Q_j)X_g(x_e y_f + y_e X_f) = I(Q_j)X_g(x_f y_e + y_f X_e).$$

- (c) *For every  $M = U_1^{efg} \oplus Q_j \in \mathcal{C}$  with  $P$ -quotient  $Q_j$  ( $U_1^{efg}$  is a triad),*

$$I(Q_j)Y_g(x_e Y_f + y_e x_f) = I(Q_j)Y_g(x_f Y_e + y_f x_e).$$

- (d) *If  $\{e, f\} = E(M)$  is a parallel pair connected to  $P$ ,*

$$I(Q_j)(x_e Y_f + y_e x_f) = I(Q_j)(x_f Y_e + y_f x_e)$$

*where  $P$ -quotient  $Q_j = M/e \setminus f = M/f \setminus e$ .*



(e) If  $\{e, f\} = E(M)$  is a series pair connected to  $P$ ,

$$I(Q_j)(x_e y_f + y_e X_f) = I(Q_j)(x_f y_e + y_f X_e)$$

where  $P$ -quotient  $Q_j = M/e \setminus f = M/f \setminus e$ .

**3.1. Remarks.** Proposition 3 assures that the different expressions for  $P$ -quotients  $Q_i$  in Theorem 4 are in fact equal as matroid or as oriented matroids, depending on how  $M$  was given.

The first three cases are trivial extensions of the conditions in Theorem 2 [11]. The only difference is that our conditions have the factor  $I(Q_j)$  in place of  $\alpha = T(\emptyset)$ . Just two new conditions are required by  $P \neq \emptyset$ . They are vacuous when  $P = \emptyset$ .

Our proof is the immediate result of adding considerations of ports to the proof in [11], there described as “a straightforward adaption of the proof of Theorem 3.3 of [24].”

The idea is to use induction on  $n = |E(M)|$  rather than on  $|S(M)|$ . As in [24, 2, 11], the necessity of each condition is verified for families of examples with  $n = 2$  and 3, where the family members correspond to different  $P$ -quotients and different choices of  $e$ ,  $f$  and  $g$ . When  $P = \emptyset$ , the only  $P$ -quotient is  $\emptyset$ .

As in [11], we rely on the hypothesis the  $\mathcal{C}$  is closed under our  $P$ -minors in order to verify (1) that the conditions imply  $T(M)$  is well-defined for  $n = 0, 1$  and 2 and (2) that in a larger minimum  $n$  counterexample, the elements of  $E(M)$  are either all in series or all in parallel, and then the conditions imply that all calculation orders give the same result. All the cases involve two different combinations of deleting and contracting of several elements in  $E(M)$  where both combinations produce the same  $P$ -quotients.

The empty matroid  $\emptyset$  is clearly the only indecomposable for the separator strong Tutte identities with  $P = \emptyset$ . Then  $\emptyset \in \mathcal{C}$  is required, provided  $\mathcal{C} \neq \emptyset$ . When we generalize to the  $P$ -ported separator strong Tutte identities with  $P \neq \emptyset$ , the indecomposables depend on  $\mathcal{C}$  and  $\emptyset \notin \mathcal{C}$  is possible.

**3.2. Proof.** We sketch the proof with a few details, pointing out differences from [11].

As in [11], the necessary relations are easy to deduce by applying (TA) and (TSSM) to the particular matroids or oriented matroids of  $\mathcal{C}$  to which they apply. Now on to the converse.

Let  $M \in \mathcal{C}$  be a counterexample with minimum  $n = |E(M)|$ . Therefore, whenever  $M'$  is a proper  $P$ -minor of  $M$ ,  $T(M')$  is well-defined. The Tutte conditions (TA) and (TSSM) have the property that given  $M$  and  $e \in M$ , exactly one equation applies. Therefore, the induction hypothesis entails that calculations that yield different values for  $T(M)$  must start with reducing by different elements of  $E(M)$ . Since  $T(M)$  is given unambiguously by the initial value  $I(M)$  when  $n = 0$ , we can assume  $n \geq 2$ .

$M$  cannot contain a separator  $e \in E(M)$ . This is a consequence of the fact, applied to  $P$ -minors, that a separator  $e \in E(M)$  is a separator in every minor of  $M$  containing  $e$ . Therefore, as observed in [11], every computation has the same result  $X_e T(M/e)$  or  $Y_e T(M \setminus e)$  depending on whether  $e$  is a coloop or a loop.

Let  $e$  be one element in  $E(M)$ . Since no element in  $E(M)$  is a separator,  $V = x_e T(M/e) + y_e T(M \setminus e)$  is well-defined, and so is  $x_{e'} T(M/e') + y_{e'} T(M \setminus e')$  for each other  $e' \in E$ . We follow [11] and define  $D = \{e \in E(M) \mid V = x_e T(M/e) + y_e T(M \setminus e)\}$ . The induction hypothesis then tells us that there is at least one element  $f \in E(M) \setminus D$ . (Our notation means of course that  $e, e', f \notin P$ .)

Suppose that  $e$  is a separator in both  $M \setminus f$  and  $M/f$  and  $f$  is a separator in both  $M \setminus e$  and  $M/e$ . Then,  $T$  would be well-defined for all four of these  $P$ -minors and so we can write

$$T(M) = x_e x_f T(M/\{e, f\}) + x_e y_f T(M/e \setminus f) + y_e x_f T(M \setminus e/f) + y_e y_f T(M \setminus \{e, f\}).$$

Both computations give the same value because in this situation the reductions by  $e$  and  $f$  commute. So, for  $M$  to be a counterexample, there must be  $e \in D$  and  $f \notin D$  ( $e, f \notin P$ ) to which one of the following two lemmas apply:

**Lemma 5.** [24] *Let  $e, f$  be nonseparators in a matroid  $M$ . The following statements are equivalent:*

- (1)  $e$  is a separator in  $M \setminus f$ .
- (2)  $e$  is a coloop in  $M \setminus f$ .
- (3)  $e$  and  $f$  are in series in  $M$ .
- (4)  $f$  is a separator in  $M \setminus e$ .

**Lemma 6.** [24] *Let  $e, f$  be nonseparators in a matroid  $M$ . The following statements are equivalent:*

- (1)  $e$  is a separator in  $M/f$ .
- (2)  $e$  is a loop in  $M/f$ .
- (3)  $e$  and  $f$  are in parallel in  $M$ .
- (4)  $f$  is a separator in  $M/e$ .

We claim that one of the following five cases must be satisfied:

- (1)  $n = 2$  and  $E(M) = \{e, f\}$  is a dyad.
- (2)  $n \geq 3$  and  $E(M)$  is a circuit not connected to  $P$ .
- (3)  $n \geq 3$  and  $E(M)$  is a cocircuit not connected to  $P$ .
- (4)  $n \geq 2$  and for some  $\emptyset \neq P' \subseteq P$ ,  $P' \cup E(M)$  is a circuit.
- (5)  $n \geq 2$  and for some  $\emptyset \neq P' \subseteq P$ ,  $P' \cup E(M)$  is a cocircuit.

As in [11], we draw the conclusion that if  $e \in D$  and  $f \notin D$  then  $e, f$  are either series or parallel. It was further proven that a series pair and a parallel pair cannot have exactly one element in common. Therefore, the pairs  $e, f$  satisfying the conditions are either all series pairs or all parallel pairs. By minimality of  $n$ ,  $E(M)$  is either an  $n$ -element parallel class or an  $n$ -element series class. The last two cases are distinguished from the first three according to whether or not  $E(M)$  is disconnected or not from elements of  $P$  in matroid  $M$ . We now use (TA) and (TSSM) to show that, in each case, the calculations that start with  $e$  and those that start with  $f$  have the same result, which contradicts  $e \in D$  and  $f \notin D$ .

We give the details for case 4. By hypothesis, each of  $T(M/e)$ ,  $T(M \setminus e)$ ,  $T(M/f)$ ,  $T(M \setminus f)$ ,  $T(M/f/e) = T(M/e/f)$ ,  $T(M/e \setminus f)$  and  $T(M/f \setminus e)$  is well-defined. Furthermore, by Proposition 3,  $M/e \setminus f = M/f \setminus e$  as matroids or oriented matroids depending on how  $M$  was given.

Starting with  $e$  and with  $f$  give the two expressions:

$$\begin{aligned} V &= x_e x_f T(M/e/f) + x_e y_f T(M/e \setminus f) + y_e T(M \setminus e) \\ V &\neq x_f x_e T(M/f/e) + x_f y_e T(M/f \setminus e) + y_f T(M \setminus f) \end{aligned}$$

Let  $M'$  be the  $P$ -minor obtained by contracting each element in  $E(M)$  except for  $e$  and  $f$  ( $M' = M$  if  $n = 2$ .) Since  $E(M') = \{e, f\}$ , (TA) tells us that

$$I(Q)(x_e y_f + y_e x_f) = I(Q)(x_f y_e + y_f x_e),$$

where  $Q = M'/e \setminus f = M'/f \setminus e$ . The latter two matroids or oriented matroids are equal because  $e, f$  are in series in  $M'$  and so Proposition 3 applies. Since  $A = E(M) \setminus \{e, f\}$  is a set of coloops ( $\emptyset$  if  $n = 2$ ) in  $M/e \setminus f = M/f \setminus e$ , we write  $X_A = \prod_{a \in A} X_a$  (1 if  $A = \emptyset$ ) by  $X_A$  and use (TSSM) to write

$$T(M/e \setminus f) = X_A I(Q).$$

and

$$T(M/f \setminus e) = X_A I(Q).$$

$$T(M \setminus e) = Y_f X_A I(Q).$$

$$T(M \setminus f) = Y_e X_A I(Q).$$

So

$$x_e y_f T(M/e \setminus f) + y_e T(M \setminus e) = x_f y_e T(M/f \setminus e) + y_f T(M \setminus f)$$

which contradicts  $V \neq x_f x_e T(M/\{f, e\}) + x_f y_e T(M/f \setminus e) + y_f T(M \setminus f)$ .

The remaining cases can be completed analogously. It might be noted that our proof differs slightly from [11] in that the cases of  $n = 3$  and  $n \geq 4$  are not distinguished.

**3.3. Universal Tutte Polynomial.** It is easy to follow [2, 11] to define a universal, i.e., most general  $P$ -ported parametrized Tutte function  $T^{\mathcal{C}}$  for the  $P$ -minor closed class  $\mathcal{C}$  given without parameters or initial values. To do this, we take indeterminates  $x_e, y_e, X_e, Y_e$  for each  $e \in E(\mathcal{C})$  and an indeterminate  $[Q_i]$  for each  $P$ -quotient  $Q_i \in \mathcal{C}$ . Let  $\mathbb{Z}[x, y, X, Y]$  denote the integer polynomial ring generated by the  $x_e, y_e, X_e, Y_e$  indeterminates, define  $\tilde{\mathbb{Z}}$  to be the  $\mathbb{Z}[x, y, X, Y]$ -module generated by the  $[Q_i]$ . Let  $I^{\mathcal{C}}$  denote the ideal of  $\tilde{\mathbb{Z}}$  generated by the identities of Theorem 4, comprising for example  $[Q_i](x_e Y_f + y_e X_f - x_f Y_e - y_f X_e)$  for each subcase of case (a), etc. The universal Tutte function has values in the quotient module  $\tilde{\mathbb{Z}}/I^{\mathcal{C}}$ . Finally, observe that the range of Tutte function  $T$  can be considered to be a module generated by the values  $T(Q_i)$  over a ring  $S$  containing the  $x, y, X, Y$  parameters; if the  $T(Q_i)$  are in  $S$ , consider  $S$  to be the  $S$ -module. We follow [11] to write the corresponding consequence of Theorem 4:

**Corollary 7.** *Let  $\mathcal{C}$  be a  $P$ -minor closed class of matroids or oriented matroids. Then there is a  $\tilde{\mathbb{Z}}/I^{\mathcal{C}}$ -valued function  $T_P$  on  $\mathcal{C}$  with  $T_P(Q_i) = [Q_i]$  for each  $P$ -quotient  $Q_i \in \mathcal{C}$  that is a  $P$ -ported parametrized Tutte function on  $\mathcal{C}$  where the parameters are the  $x, y, X, Y$  indeterminates. Moreover, if  $T$  is any  $R$ -parametrized Tutte function with parameters  $x'_e, y'_e, X'_e, Y'_e$ , then  $T$  is the composition of  $T_P$  with the homomorphism determined by  $[Q_i] \rightarrow T(Q_i)$  for  $P$ -quotient and  $x_e \rightarrow x'_e$ , etc., for each  $e \in E(\mathcal{C})$ .*

In the next section we define particular expressions for  $T_P(M)$  which will be called Tutte polynomials. Our purpose for allowing Tutte function values and polynomials to be in  $R$ -modules is that it facilitates giving formulas for Tutte functions of combinations such as direct sum in terms of multiplication rules that make the module into algebra.

#### 4. TUTTE COMPUTATION TREES AND ACTIVITIES

Several authors [10, 11] surveyed the two ways that the two-variable Tutte polynomial can be defined: It may be defined either as a universal solution to the recursive strong Tutte equations, or as a generating function. Further, two kinds of generating function definitions have been given. The first is what Tutte originally used for graphs [22, 23] and is called the *basis* or *activities* expansion. It enumerates each basis  $B \subseteq E$  by a term  $x^{i(B)} y^{e(B)}$ , where

the *numbers of internally and externally active elements*  $i(B)$  and  $e(B)$  are determined from a given linear order on the elements of  $E$  (see Definition 11). It was shown that even though  $(i(B), e(B))$  for particular  $B$  might vary with the order, the resulting polynomial is independent of this order, and that it satisfies the Tutte equations. The second, called the *rank-nullity* generating function, is well-defined automatically because it enumerates each subset  $A \subseteq E$  with the term  $(x - 1)^{r(E) - r(A)}(y - 1)^{(|A| - r(A))}$ . This generating function is then shown to satisfy the Tutte equations.

A remarkable conclusion about adding parameters, given by Zaslavsky [24], is that the activities expansion generalizes to a universal form for Tutte functions whereas the rank-nullity function generalization can only express the *normal* Tutte functions. The latter are characterized by (CNF). See sec. 4.2.

Ellis-Monaghan and Traldi [11] remarked that the Tutte equation approach appears to give a shorter proof of the ZBR theorem than the activities expansion approach. Diao and Heteyi [9] proved specializations of Theorem 4 by means of the activities expansion approach. The inductive proofs on  $|E|$  that we and [11] give demonstrate that every calculation of  $T(M)$  from Tutte equations produces the same result when the conditions on the parameters and initial values are satisfied. It is then almost a tautology that the polynomial expression resulting from a particular calculation will equal the Tutte function value in the ring  $R$ . We show that every recursive calculation (see below) gives rise to an activities expansion, when the activities are defined in the more general way given by McMahon and Gordon [15]. We suggest a heuristic reason why the inductive Tutte equation approach is more succinct: The induction assures that *every* computation for a matroid with smaller  $|E|$  gives the same result, not just those computations that are determined by linear orders on  $E$ .

We call any computation that uses Tutte equations to find  $T(M)$  in terms of the Tutte function of indecomposables and/or of  $M'$  with  $|E(M')| < |E(M)|$  *recursive*. All the recursive computations of  $T(M)$  are expressible by “computation trees,” formally defined by McMahon and Gordon [15]. Their motivation was to generalize activities expansions and the corresponding interval partitions of the subset lattice from matroids to greedoids. Unlike matroids, some greedoids do not have an activities expansion for their Tutte polynomial that derives from an element ordering.

Proofs of activities expansions for matroids, and their generalizations for  $P$ -ported matroids, seem more informative and certainly no harder when the expansions are derived from a general Tutte computation tree, than when the expansions are only those that result from an element order. From the retrospective that the Tutte equations specify a non-deterministic recursive computation [13], it seems artificial to start with element-ordered computations and then prove first that all linear orders give the same result and second that it satisfies the Tutte equations, in order prove that all recursions give the same result. We therefore take advantage of the Tutte computation tree formalism and the more general expansions it enables.

**4.1. Computation Tree Expansion.** We begin with the definition of what generalizes the matroid bases in the activities expansion when  $P \neq \emptyset$ . In the following,  $\mathcal{B}(M)$  denotes the set of bases in  $M$ .

**Definition 8.** *Given  $P$ -ported matroid or oriented matroid  $M$ , a  **$P$ -subbasis**  $F \in \mathcal{B}_P(M)$  is an independent set with  $F \subseteq E(M)$  (so  $F \cap P = \emptyset$ ) for which  $F \cup P$  is a spanning set for  $M$  (in other words,  $F$  spans  $M/P$ ).*

An equivalent definition was given in [20]. The following proposition shows our definition is equivalent to that given in [9].  $C$  and  $D$  below are called “contracting and deleting sets” in that paper.

**Proposition 9.**  *$C$  is a  $P$ -subbasis if and only if  $C \subseteq E((M))$  has no circuits and  $D = E(M) \setminus C$  has no cocircuits.*

*Proof.*  $C$  has no circuits means  $C$  is an independent set in  $M$ .  $D$  has no cocircuits means  $D$  is independent in the dual of  $M$ , i.e.,  $D$  is coindependent.  $D$  is coindependent if and only if  $S(M) \setminus D = P \cup C$  spans  $M$ .  $\square$

**Proposition 10.** *For every  $P$ -subbasis  $F$  there exists an independent set  $Q \subseteq P$  that extends  $F$  to a basis  $F \cup Q \in \mathcal{B}(M)$ . Conversely, if  $B \in \mathcal{B}(M)$  then  $F = B \cap E = B \setminus P$  is a  $P$ -subbasis.*

*Proof.* Immediate.  $\square$

The next definition is also equivalent to one in [9]. It generalizes Tutte’s definitions based on element orderings [23, 22] extended to matroids [8]. We will see that expansions based on computation trees generalize these further.

**Definition 11** (Activities with respect to a  $P$ -subbasis and an element ordering  $O$ ). *Let ordering  $O$  have every  $p \in P$  before every  $e \in E$ . Let  $F$  be a  $P$ -subbasis. Let  $B$  be any basis for  $M$  with  $F \subseteq B$ .*

- *Element  $e \in F$  is internally active if  $e$  is the least element within its principal cocircuit with respect to  $B$ . Thus, this principal cocircuit contains no ports. The reader can verify this definition is independent of the  $B$  chosen to extend  $F$ . Elements  $e \in F$  that are not internally active are called internally inactive.*
- *Dually, element  $e \in E$  with  $e \notin F$  is externally active if  $e$  is the least element within its principal circuit with respect to  $B$ . Thus, each externally active element is spanned by  $F$ . Elements  $e \in E \setminus F$  that are not externally active are called externally inactive.*

**Definition 12** (Computation Tree, following [15]). *A  $P$ -ported (Tutte) computation tree for  $M$  is a binary tree whose root is labeled by  $M$  and which satisfies:*

- (1) *If  $M$  has non-separating elements not in  $P$ , then the root has two subtrees and there exists one such element  $e$  for which one subtree is a computation tree for  $M/e$  and the other subtree is a computation tree for  $M \setminus e$ .*

*The branch to  $M/e$  is labeled with “ $e$  contracted” and the other branch is labeled “ $e$  deleted”.*

- (2) *Otherwise (i.e., every element in  $E(M)$  is separating) the root is a leaf.*

An immediate consequence is

**Proposition 13.** *Each leaf of a  $P$ -ported Tutte computation tree for  $M$  is labeled by the direct sum of some  $P$ -quotient (oriented if  $M$  is oriented) summed with loop and/or coloop matroids with ground sets  $\{e\}$  for various distinct  $e \in E$  (possibly none).*

**Definition 14** (Activities with respect to a leaf). *For a  $P$ -ported Tutte computation tree for  $M$ , a given leaf, and the path from the root to this leaf:*

- *Each  $e \in E(M)$  labeled “contracted” along this path is called **internally passive**.*
- *Each coloop  $e \in E(M)$  in the leaf’s matroid is called **internally active**.*

- Each  $e \in E(M)$  labeled “deleted” along this path is called **externally passive**.
- Each loop  $e \in E(M)$  in the leaf’s matroid is called **externally active**.

**Proposition 15.** *Given a leaf of a  $P$ -ported Tutte computation tree for  $M$ , the set of internally active or internally passive elements constitutes a  $P$ -subbasis of  $M$  which we say **belongs to the leaf**. Furthermore, every  $P$ -subbasis  $F$  of  $N$  belongs to a unique leaf.*

*Proof.* For the purpose of this proof, let us extend Definition 14 so that, given a computation tree with a given node  $i$  labeled by matroid  $M_i$ ,  $e \in E$  is called internally passive when  $e$  is labeled “contracted” along the path from root  $M$  to node  $i$ . Let  $IP_i$  denote the set of such internally passive elements.

It is easy to prove by induction on the length of the root to node  $i$  path that (1)  $IP_i \cup S(M_i)$  spans  $M$  and (2)  $IP_i$  is an independent set in  $M$ . The proof of (1) uses the fact that elements labeled deleted are non-separators. The proof of (2) uses the fact that for each non-separator  $f \in M/IP_i$ ,  $f \cup IP_i$  is independent in  $M$ .

These properties applied to a leaf demonstrate the first conclusion, since each  $e \in E$  in the leaf’s matroid must be a separator by Definition 12.

Given a  $P$ -subbasis  $F$ , we can find the unique leaf with the algorithm below. Note that it also operates on arbitrary subsets of  $E$ .

**Tree Search Algorithm:** Beginning at the root, descend the tree according to the rule: At each branch node, descend along the edge labeled “ $e$ -contracted” if  $e \in F$  and along the edge labeled “ $e$ -deleted” otherwise (when  $e \notin F$ ).  $\square$

The above definitions and properties lead us to reproduce element order based activities:

**Proposition 16.** *Given element ordering  $O$  in which every  $p \in P$  is ordered before each  $e \notin P$ , suppose we construct the unique  $P$ -ported computation tree  $\mathcal{T}$  in which the greatest non-separator  $e \in E$  is deleted and contracted in the matroid at each tree node.*

*The activity of each  $e \in E$  relative to ordering  $O$  and  $P$ -subbasis  $F \subseteq E$  is the same as the activity of  $e$  defined with respect to the leaf belonging to  $F$  in  $\mathcal{T}$ .*

**Definition 17.** *Given a computation tree for  $P$ -ported (oriented) matroid  $M$ , each  $P$ -subbasis  $F \subseteq E$  is associated with the following subsets of non-port elements defined according to Definition 14 from the unique leaf determined by the algorithm given above.*

- $IA(F) \subseteq F$  denotes the set of internally active elements,
- $IP(F) \subseteq F$  denotes the set of internally passive elements,
- $EA(F) \subseteq E \setminus F$  denotes the set of externally active elements, and
- $EP(F) \subseteq E \setminus F$  denotes the set of externally passive elements.
- $A(F) = IA(F) \cup EA(F)$  denotes the set of active elements.

**Proposition 18.** *Given a  $P$ -ported Tutte computation tree for  $M$ , the boolean lattice of subsets of  $E = E(M)$  is partitioned by the collection of intervals  $[IP(F), F \cup EA(F)]$  (note  $F \cup EA(F) = IP(F) \cup A(F)$ ) determined from the collection of  $P$ -subbases  $F$ , which correspond to the leaves.*

*Proof.* Every subset  $F' \subseteq E = E(M) \setminus P$  belongs to the unique interval corresponding to the unique leaf found by the tree search algorithm given at the end of the previous proof.  $\square$

Dualizing, we obtain:

**Proposition 19.** *Given a  $P$ -ported Tutte computation tree for  $M$  the boolean lattice of subsets of  $E = E(M)$  is also partitioned by the collection of intervals  $[EP(F), E \setminus F \cup IA(F)]$  (note  $E \setminus F \cup IA(F) = EP(F) \cup A(F)$ ).*

*Proof.* The dual of the tree search algorithm is to descend along the edge labeled “ $e$ -deleted” if  $e \in F$ .  $\square$

The following generalizes the activities expansion expression given in [24] to ported (oriented) matroids, as well as Theorem 8.1 of [20].

**Definition 20.** *Given parameters  $x_e, y_e, X_e, Y_e$ , and  $P$ -ported matroid or oriented matroid  $M$  the Tutte polynomial expression determined by the sets in Definition 17 from a computation tree is given by*

$$(PAE) \quad \sum_{F \in \mathcal{B}_P} I(M/F|P) X_{IA(F)} x_{IP(F)} Y_{EA(F)} y_{EP(F)}.$$

Each Tutte polynomial expression is constructed by applying some of the Tutte equations. Each expression  $M/F|P$  denotes a  $P$ -quotient of  $M$ , so the expression is a polynomial in the parameters and in the initial values  $I(M/F|P)$ . We can consider each  $P$ -quotient to be a variable, which we denote  $[M/F|P]$  as in [5] (except we recognize now that  $M$  and  $M/F|P$  might be oriented matroids). Therefore, if  $P$ -ported Tutte function  $F$  is in the domain  $\mathcal{C}$  Tutte function  $F$ , then  $F(M)$  is given by any Tutte polynomial expression with  $F(M/F|P)$  substituted for each oriented or unoriented matroid variable  $[M/F|P]$ . We have proven:

**Theorem 21.** *For every  $P$ -ported parametrized Tutte function  $T$  into  $R$  on  $\mathcal{C}$ , for every computation tree for  $M \in \mathcal{C}$  (and so for every ordering of  $E(M)$ ), the polynomial expression (PAE) equals  $T(M)$  in  $R$  when the initial values are assigned so  $[Q_i] = T(Q_i)$  for every  $P$ -quotient  $Q_i$ .*

#### 4.2. Expansions of Normal Tutte Functions.

Zaslavsky [24] defined *normal* Tutte functions to be those for which  $T(\emptyset) = 1$  **CHECK!!**, and for which there exist  $u, w \in R$  so that for each  $e \in E(M)$ ,

$$(CNF) \quad X_e = x_e + uy_e \text{ and } Y_e = y_e + wx_e.$$

It is easy to generalize the well-known rank-nullity expansion for the classical Tutte polynomial to the following expansion for a  $P$ -ported parametrized normal Tutte function. A proof that this is a universal solution for (TA) and (TSM) is easily obtained by including the  $x_e, y_e$  parameters, and allowing  $M$  and its minors to sometimes be oriented, into the proof we detailed in [5].

$$(PGF) \quad T(M) = \sum_{A \subseteq E(M)} T(M/A | P) x_A y_{E \setminus A} u^{r(M) - r(M/A|P) - r(A)} w^{|A| - r(A)}.$$

We can now also verify (PGF) expresses a normal Tutte function by deriving it from the Boolean interval expansion obtained when (CNF) parameters are substituted into the activities expansion of Theorem 21.

4.2.1. *Boolean Interval Expansion.* The following proposition expresses the ported corank-nullity polynomial in terms of a  $P$ -ported activities expansion. It is obtained by substituting binomials  $X_e = x_e + y_e u$ ,  $Y_e = y_e + x_e v$  and leaving the matroid or oriented matroid variables unchanged in Definition 20.

**Proposition 22.** *The polynomial  $R_P(M)$  is given by the following activities and boolean interval expansion formula:*

$$(1) \quad R_P(M) = \sum_{F \in \mathcal{B}_P} [M/F|P] \left( \sum_{\substack{IP(F) \subseteq K \subseteq F \\ EP(F) \subseteq L \subseteq E \setminus F}} x_{K \cup (E \setminus F \setminus L)} v^{|E \setminus F \setminus L|} y_{L \cup (F \setminus K)} u^{|F \setminus K|} \right)$$

*Proof.* Let  $A = K \cup (E \setminus F \setminus L)$  within (PAE). We can verify  $\bar{A} = E \setminus A = L \cup (F \setminus K)$ . For each  $A \subseteq E$  a unique  $P$ -subbasis  $F$ , and two tree leaves are determined, one by the tree descending algorithm and the other leaf by the dual algorithm. Thus  $A$  and  $\bar{A}$  respectively belong to intervals within the boolean lattice partitions of Propositions 18 and 19. In particular,  $A \in [IP(F), F \cup EA(F)]$  and  $\bar{A} \in [EP(F), E \setminus F \cup IA(F)]$ . Therefore the terms in the above sum are equal one by one to the terms in the corank-nullity polynomial's subset expansion (Definition PGF) because the terms are in correspondence with the subsets of  $E(M)$ .  $\square$

4.2.2. *Geometric Lattice Flat Expansion.* A formula for the unparametrized ported Tutte (or corank-nullity) polynomials of non-oriented matroids in terms of the lattice of flats (closed sets) and its Möbius function was given in [5]. We generalize: (1) The expansion's monomials  $[Q_i]$  can signify either oriented matroid minors, when  $M$  is oriented, or non-oriented minors when  $M$  is not oriented. (2) The polynomial is parametrized with  $x_e, y_e$  for each  $e \in E$ . The derivation relies on the fact that the oriented or non-oriented matroid minor  $[M/A|P]$  (according to whether  $M$  is oriented or not) depends only on the flat spanned by  $A \subseteq E$ .

**Proposition 23.** *Let  $N$  be an oriented or unoriented. Let  $R_P(M)$  be given by (PGF). In the formula below,  $F$  and  $G$  range over the geometric lattice of flats contained in  $M$  restricted to  $E$ .*

$$(2) \quad R_P(M)(u, v) = \sum_{Q_i} [Q_i] \sum_{\substack{F \subseteq E \\ [M/F|P] = [Q_i]}} u^{r(M) - r(Q_i) - r(F)} v^{-r(F)F} \sum_{G \subseteq F} \mu(G, F) \prod_{e \in G} (y_e + x_e v)$$

*Proof.* It follows the steps for theorem 8 in [5].  $\square$

## 5. DIRECT SUMS AND OBJECTS WITH MATROIDS

The introduction of ports  $P$  into the theory of separator-strong Tutte functions of direct sums suffers from two complications. First, a  $P$ -family of matroids  $\mathcal{C}$  can have more than the one indecomposable  $\emptyset$ . Second,  $\mathcal{C}$  might be partitioned into non-empty disjoint  $P$ -minor closed classes. Each class has different indecomposables. The second situation also occurs when  $P = \emptyset$  with minor closed subclasses of graphs [11] where deletion of isthmuses and isolated vertices is forbidden. When  $P \neq \emptyset$ , each subclass can have multiple indecomposables.

We define an abstraction of graphs called  *$P$ -families of objects with matroids or oriented matroids*. It facilitates a discussion of how the known notions of separator-strong and strong Tutte functions of graphs can be extended when deletion and contraction are restricted to  $e \notin P$ . Proofs [11] based on the one matroid of an indecomposable being  $\emptyset$  and on just one indecomposable graph (the edgeless graph with  $k$  vertices) in each subclass also seem easier



to generalize with Tutte computation trees. Our abstraction and proof of Theorem 2 lead us to formulate new ZBR-like conditions, see Theorem 36 below.

It is useful to think that a  $P$ -ported Tutte computation tree may have objects  $N$  for its node labels, such as graphs, that have matroids or oriented matroids  $M(N)$  associated with them. Each  $N$  has elements  $S(N) = S(M(N))$ , each  $p \in S(N) \cap P$  is called a port, and  $E(N) = S(N) \setminus P$ . Loops, coloops and non-separators of  $N$  are characterized by their status in  $M(N)$ . So we say  $N$  is an *object with a matroid or an oriented matroid*.

Contraction  $N/e$  and deletion  $N \setminus e$  of object  $N$  is defined when  $e \in E(N)$ ; and  $e$  is not a coloop in  $M(N)$  and  $e$  is not a loop in  $M(N)$ , respectively. Under those conditions,  $M(N/e) = M(N)/e$  and  $M(N \setminus e) = M(N) \setminus e$  (as matroids or oriented matroids). Thus  $P$ -minors are defined, and an indecomposable or  $P$ -quotient is a  $P$ -minor  $Q$  for which  $S(Q) = S(M(Q)) \subseteq P$ .

A  $P$ -family  $\mathcal{N}$  is a  $P$ -minor closed class of objects with matroid or oriented matroids. Tutte computation trees are then defined for such  $N \in \mathcal{N}$ . The matroids of course constrain the structure of these trees. It is possible (as when edgeless graphs  $G_k$  have different vertex sets but all  $M(G_k) = \emptyset$ , see sec. 6) for different objects, even different indecomposables, to have the same matroid or oriented matroid. It also naturally occurs (see 6) that  $N/e \setminus f \neq N/f \setminus e$  (as objects) even though  $M(N)/e \setminus f = M(N)/f \setminus e$ . The latter equation when  $e, f$  are in parallel or in series (see Proposition 3) is critical to the above ZBR theorems. It is also conceivable that  $N/e/f \neq N/f/e$  or  $N \setminus e \setminus f \neq N \setminus f \setminus e$ .

We can say that  $T$  on  $\mathcal{N}$  is a  $P$ -ported separator-strong Tutte function on  $\mathcal{N}$  into the ring  $R$  containing parameters  $x_e, y_e, X_e, Y_e$  for each  $e \in E(N)$  for some  $N \in \mathcal{N}$ , or into an  $R$ -module containing the values on the indecomposables, if  $T(N)$  satisfies (TA) and (TSSM) for all  $N \in \mathcal{N}$ . It follows:

**Proposition 24.**  *$T$  is a  $P$ -ported separator-strong Tutte function on  $\mathcal{N}$  if and only if for each  $N \in \mathcal{N}$ , all Tutte computation trees for  $T(N)$  yield polynomial expressions that are equal in the range ring or  $R$ -module.*

We can still talk about Tutte decompositions and a Tutte computation tree for  $N$  even without a Tutte function. If we are only given values  $I(Q_i)$  for the indecomposables, each Tutte computation tree for  $N$  yields a value in the  $R$ -module generated by the  $I(Q_i)$ . The Tutte decompositions, and the universal Tutte polynomial (if it exists!) of each  $N \in \mathcal{N}$  are determined by  $M(N)$  and the indecomposables, i.e.,  $P$ -quotients  $Q_i$  in  $\mathcal{N}$ , which of course satisfy  $Q_i \in \mathcal{N}$ .

See sec. 6.2 for graph examples.

The following easy observations formalize how Tutte computation trees for  $N$  and  $N'$  have the same structure when  $M(N) = M(N')$  and how the Tutte polynomials are related.

**Lemma 25.** *If  $M(N) = M(N')$  and  $\sum_{i \in Z} I(N/B_i|P)r_i$  is the value from Tutte computation tree  $\mathcal{T}$  for  $N$ , then  $\mathcal{T}$  can be relabelled to construct a Tutte computation tree for  $N'$  that yields the value  $\sum_{i \in Z} I(N'/B_i|P)r_i$  with the same index set  $Z$  and  $r_i \in R$  for  $i \in Z$ .*

**Proposition 26.** *If  $T$  is a separator-strong Tutte function on  $\mathcal{N}$  and  $T(N) = \sum_{i \in Z} I(N/B_i|P)r_i$  is the Tutte polynomial expression from some computation tree, then  $T(N') = \sum_{i \in Z} I(N'/B_i|P)r_i$  with the same index set  $Z$  and  $r_i \in R$  for  $i \in Z$ .*

It is a common situation that  $\{N^1, N^2, N\} \subseteq \mathcal{N}$  and their matroids or oriented matroids  $M(N^1) \oplus M(N^2) = M(N)$ . Tutte computation trees help.

**Definition 27.** If  $\mathcal{T}_1$  and  $\mathcal{T}_2$  are Tutte computation trees then  $\mathcal{T}_1 \cdot \mathcal{T}_2$  is the tree obtained by appending a separate copy of  $\mathcal{T}_2$  at each leaf of  $\mathcal{T}_1$ . The root is the root of the expanded  $\mathcal{T}_1$ .

**Proposition 28.** Suppose  $N$ ,  $N^1$  and  $N^2$  are all in  $\mathcal{N}$  and  $M(N^1) \oplus M(N^2) = M(N)$ . Then if  $\mathcal{T}_1$  and  $\mathcal{T}_2$  are Tutte computation trees for  $N^1$  and  $N^2$  respectively with values given by (DS1) and (DS2), then there is a Tutte computation tree for  $N$  that yields the value given by (DS).

$$(DS1) \quad \sum_{Q_i^1} I(Q_i^1) c_1(Q_i^1) \text{ where } Q_i^1 = N/B_i^1|P.$$

$$(DS2) \quad \sum_{Q_j^2} I(Q_j^2) c_2(Q_j^2) \text{ where } Q_j^2 = N/B_j^2|P.$$

$$(DS) \quad \sum_{Q_i^1, Q_j^2} I(Q_{i,j}) c_1(Q_i^1) c_2(Q_j^2) \text{ where } Q_{i,j} = N/B^1/B^2|P.$$

Furthermore, if  $T$  is a Tutte function on  $\mathcal{N}$  and  $T(N^1)$  and  $T(N^2)$  equal the Tutte polynomials given by (DS1) and (DS2) then  $T(N)$  equals the polynomial given by (DS).

*Proof.* We show how to relabel  $\mathcal{T}_1 \cdot \mathcal{T}_2$  to obtain a Tutte computation tree for  $N$ .  $M(N^1) \oplus M(N^2) = M(N)$  is defined means  $S(M(N^1)) \cap S(M(N^2)) = \emptyset$  and  $S(M(N)) = S(M(N^1)) \cup S(M(N^2))$ . Each node of  $\mathcal{T}_1 \cdot \mathcal{T}_2$  is determined by deleting and/or contracting some elements of  $E(M(N^1)) \cup E(M(N^2))$ . Relabel that node with the  $P$ -minor of  $N$  obtaining deleting and/or contracting the same elements respectively. The result is a computation tree for  $N$  because  $M(N^1) \oplus M(N^2) = M(N)$ . Assume  $P \subseteq S(M(N))$  (otherwise, take a smaller  $P$ ) and let  $P^1 = S(M(N^1)) \cap P$  and  $P^2 = S(M(N^2)) \cap P$ . At a leaf of the relabelled tree, there will be the  $P$ -quotient  $N/B_1/B_2|P$  where  $B_1$  is a  $P^1$ -subbasis of  $M(N^1)$  and  $B_2$  is a  $P^2$ -subbasis of  $M(N^2)$ . □

**5.1. Strong Tutte Functions.** Zaslavsky defined a parametrized Tutte function  $T$  on a minor-closed family of matroids  $\mathcal{C}$  to be strong if  $T(M^1 \oplus M^2) = T(M^1)T(M^2)$  whenever  $M^1, M^2$  and direct sum  $M^1 \oplus M^2 \in \mathcal{C}$ . It follows from the formula [11]  $T(M^1 \oplus M^2)T(\emptyset) = T(M^1)T(M^2)$  for separator-strong Tutte functions and  $\emptyset = \emptyset \oplus \emptyset$  that a separator-strong Tutte function is strong if and only if  $T(\emptyset)$  is idempotent, i.e.,  $T(\emptyset) = T(\emptyset)^2$  (Corollary 2.3 of [11]).

Let us extend the definition of strong parametrized Tutte function to  $P$ -families  $\mathcal{N}$  of objects with matroids and oriented matroids, in the way that abstracts the known notion of strong Tutte functions on minor closed families of graphs[11]. Of course, taking  $\mathcal{N}$  to be a  $P$ -family  $\mathcal{C}$  of matroids or oriented matroids gives us the extension to such  $\mathcal{C}$ . Then, there might still be indecomposables besides or instead of  $\emptyset$ .

**Definition 29.** A  $P$ -ported separator-strong Tutte function  $T$  on a  $P$ -family of objects  $\mathcal{N}$  with matroids is called strong if whenever  $\{N^1, N^2, N\} \subseteq \mathcal{N}$  and  $M(N^1) \oplus M(N^2) = M(N)$ , then  $T(N^1)T(N^2) = T(N)$ .

We can use Proposition 28 to prove the generalization of the  $T(\emptyset)T(\emptyset) = T(\emptyset)$  characterization of strong Tutte functions.

**Theorem 30.** *A  $P$ -ported separator-strong Tutte function  $T$  on a  $P$ -family of objects with matroids or oriented matroids  $\mathcal{N}$  is strong if and only if  $T$  restricted to the indecomposables of  $\mathcal{N}$ , is strong; i.e., whenever  $Q^1$ ,  $Q^2$  and  $Q$  are indecomposables and  $M(Q^1) \oplus M(Q^2) = M(Q)$  then  $T(Q^1)T(Q^2) = T(Q)$ .*

*Proof.* Every  $P$ -quotient is in  $\mathcal{N}$ , so clearly  $T$  restricted to the  $P$ -quotients is strong.

Conversely, suppose  $N^1$ ,  $N^2$  and  $N$  are in  $\mathcal{N}$  and  $M(N^1) \oplus M(N^2) = M(N)$ , so Proposition 28 applies.

Since  $M(N^1) \oplus M(N^2) = M(N)$ ,  $M(N/(B_1 \cup B_2)|P) = (M(N^1)/B_1|P) \oplus (M(N^2)/B_2|P) = M(N^1/B_1|P) \oplus M(N^2/B_2|P)$ . We now use the fact that  $Q_{ij} = N/B_1/B_2|P$ ,  $Q_i^1 = N^1/B_1|P$  and  $Q_j^2 = N^2/B_2|P$  are  $P$ -quotients and the hypothesis to write  $T(Q_{ij}) = T(Q_i^1)T(Q_j^2)$ .

We therefore conclude  $T(N) = T(N^1)T(N^2)$  from (DS1), (DS2) and (DS).  $\square$

**5.2. Matroidal Direct Sums and Multiplicative Tutte Functions.** Often  $\mathcal{N}$  comes equipped with one operation “ $*$ ”, or more, that satisfy the following definition. Examples for  $P = \emptyset$  are disjoint union  $\amalg$  and one-point unions of graphs [11], see sec. 6.

**Definition 31.** *A partially defined binary operation “ $*$ ” on a  $P$ -family of objects with matroids or oriented matroids  $\mathcal{N}$  is a matroidal direct sum if whenever  $N^1 * N^2 \in \mathcal{N}$  is defined for  $\{N^1, N^2\} \subseteq \mathcal{N}$ , the matroids or oriented matroids satisfy  $M(N^1) \oplus M(N^2) = M(N^1 * N^2)$ .*

Proposition 28 applies when  $N^1 * N^2 = N$  is defined. It gives a general recipe for  $T(N^1 * N^2)$  which generalizes the identity [11]  $T(M^1 \oplus M^2)T(\emptyset) = T(M^1)T(M^2)$  for separator-strong Tutte functions of matroids. The  $P$ -ported generalization is more complicated and generally cannot be expressed by a product in the domain ring of  $T$ .

**Proposition 32.** *Suppose  $*$  is a matroidal direct sum and  $N^1$ ,  $N^2$  and  $N^1 * N^2$  are each members of a  $P$ -family for which  $T$  is a Tutte function.*

*If for  $P$ -quotients  $Q_i^j$  and  $R$ -coefficients  $c_j(Q_i^j)$ ,  $j = 1$  and  $2$ ,*

$$(MD1) \quad T(N^1) = \sum_{Q_i^1} T(Q_i^1) c_1(Q_i^1)$$

*and*

$$(MD2) \quad T(N^2) = \sum_{Q_j^2} T(Q_j^2) c_2(Q_j^2)$$

*then*

$$(MD) \quad T(N^1 * N^2) = \sum_{Q_i^1, Q_j^2} T(Q_i^1 * Q_j^2) c_1(Q_i^1) c_2(Q_j^2).$$

*Proof.* Substitute  $Q_{i,j} = Q_i^1 * Q_j^2$  in (DS) of Proposition 28.  $\square$

When  $\mathcal{N}$  is a  $P$ -family of matroids or oriented matroids, direct matroid or oriented matroid sum is obviously a matroidal direct sum operation, and so Proposition 32 is applicable.

**Corollary 33.** [11] *Let  $P = \emptyset$ .  $T(M^1 \oplus M^2)T(\emptyset) = T(M^1)T(M^2)$  for Tutte function  $T$  of matroids.*

*Proof.* Our proof demonstrates how Proposition 32 generalizes this formula to  $P$ -families. The expansions DS1 and DS2 take the one-term form  $T(M^j) = T(\emptyset)c_j(\emptyset)$ ,  $j = 1, 2$ , so  $T(M^1)T(M^2) = T(\emptyset)^2c_1(\emptyset)c_2(\emptyset)$ . Expansion DS is then  $T(M^1 \oplus T^2) = T(\emptyset)c_1(\emptyset)c_2(\emptyset)$ .  $\square$

Following the definitions for graphs in [11], we write:

**Definition 34.** *Given a matroidal direct sum  $*$  on  $\mathcal{N}$ , a Tutte function  $T$  on  $\mathcal{N}$  is multiplicative (with respect to “ $*$ ”) if whenever  $N^1 * N^2$  is defined for  $\{N^1, N^2\} \subseteq \mathcal{N}$ , the Tutte function values satisfy  $T(N^1)T(N^2) = T(N^1 * N^2)$ .*

A strong Tutte function is certainly multiplicative for any “ $*$ ”, but not conversely. A consequence of Proposition 28 is that, like strong Tutte functions, multiplicative Tutte functions are characterized by being that way on the indecomposables.

**Corollary 35.** *A  $P$ -ported Tutte function  $T$  on  $P$ -family  $\mathcal{N}$  is multiplicative with respect to matroidal direct product “ $*$ ” if and only if for every pair of indecomposables  $\{Q_i, Q_j\} \in \mathcal{N}$  for which  $Q_i * Q_j \in \mathcal{N}$  is defined,  $T(Q_i)T(Q_j) = T(Q_i * Q_j)$ .*

*Proof.* When  $N^1 * N^2$  is defined, Proposition 28 applies because  $M(N^1) \oplus M(N^2) = M(N^1 * N^2)$ .  $T(N^1 * N^2) = T(N^1)T(N^2)$  is then a consequence of  $T(Q_i * Q_j) = T(Q_i)T(Q_j)$ .  $\square$

**5.3. ZBR-type Theorems.** Examples in sec. 6 below force us to recognize that for  $N$  an object with a matroid or oriented matroid  $M(N)$  with  $e, f \in E(N)$  in series or in parallel, it might happen that  $N/e \setminus f \neq N/f \setminus e$  even though, by Proposition 3,  $M(N)/e \setminus f \neq M(N)N/f \setminus e$ . Note that the additional hypotheses in the theorem below do not pertain to cases of Proposition 3.

**Theorem 36** (ZBR Theorem for  $P$ -families of objects with matroids or oriented matroids). *Let  $\mathcal{N}$  be a  $P$ -family of objects with matroids or oriented matroids.*

*Suppose, in addition, the  $P$ -minor operations are commutative in the way they are for matroids. Specifically, for all  $N \in \mathcal{N}$  and  $e, f \in E(N)$ ,  $N/e/f = N/f/e$ ,  $N \setminus e \setminus f = N \setminus f \setminus e$ ,  $N/e \setminus f = N \setminus f/e$  whenever neither  $e$  nor  $f$  is a loop in any of these contractions and neither is a coloop in any of these deletions.*

*The following two statements are equivalent.*

(1)  *$T$  from  $\mathcal{N}$  to  $R$  or an  $R$ -module is a  $P$ -ported separator-strong parametrized  $P$ -ported Tutte function with  $R$ -parameters  $(x, y, X, Y)$  whose values  $T(Q_i)$  on  $P$ -quotients  $Q_i \in \mathcal{N}$  are the initial values  $I(Q_i)$ .*

(2) *For every  $N \in \mathcal{N}$ :*

(a) *If  $M(N) = U_1^{ef} \oplus M(Q_j) = U_1^{ef} \oplus M(Q'_j)$  with  $P$ -quotients  $Q_j = N/e \setminus f$  and  $Q'_j = N/f \setminus e$ ,*

$$I(Q_j)(x_e Y_f - y_f X_e) = I(Q'_j)(x_f Y_e - y_e X_f).$$

(b) *If  $M(N) = U_2^{efg} \oplus M(Q_j) = U_2^{efg} \oplus M(Q'_j)$  with  $P$ -quotients  $Q_j = N/e \setminus f/g$  and  $Q'_j = N/f \setminus e/g$ ,*

$$I(Q_j)X_g(x_e y_f - y_f X_e) = I(Q'_j)X_g(x_f y_e - y_e X_f).$$

(c) *If  $M(N) = U_1^{efg} \oplus M(Q_j) = U_1^{efg} \oplus M(Q'_j)$  with  $P$ -quotients  $Q_j = N/e \setminus f \setminus g$  and  $Q'_j = N/f \setminus e \setminus g$ ,*

$$I(Q_j)Y_g(x_e Y_f - y_f X_e) = I(Q'_j)Y_g(x_f Y_e - y_e X_f).$$

(d) If  $\{e, f\} = E(M(N))$  is a parallel pair connected to  $P$ ,

$$I(Q_j)(x_e Y_f - y_f x_e) = I(Q'_j)(x_f Y_e - y_e x_f)$$

where  $P$ -quotients  $Q_j = N/e \setminus f$  and  $Q'_j = N/f \setminus e$ .

(e) If  $\{e, f\} = E(M(N))$  is a series pair connected to  $P$ ,

$$I(Q_j)(x_e y_f - y_f x_e) = I(Q'_j)(x_f y_e - y_e x_f)$$

where  $P$ -quotients  $Q_j = N/e \setminus f$  and  $Q'_j = N/f \setminus e$ .

*Proof.* Analogous to the proof we gave for Theorem 4.  $\square$

**Corollary 37.** *A  $P$ -family of objects with matroids satisfies a ZBR-type theorem with the identities given in Theorem 4 if, in addition to the  $P$ -minor operations being commutative as with matroids, the object  $P$ -quotients  $N/e \setminus f = N/f \setminus e$  when  $\{e, f\} = E(N)$  is a series or parallel pair,  $N/e \setminus f/g = N/f \setminus e/g$  when  $\{e, f, g\} = E(N)$  is a triangle and  $N/e \setminus f \setminus g = N/f \setminus e \setminus g$  when  $\{e, f, g\} = E(N)$  is a triad.*

For convenience in subsequent discussion, we will call a theorem with identities like those in Theorem 4 or in the ZBR theorem for graphs given by Ellis-Monaghan and Traldi a ZBR-matroid type theorem, and one with identities like those just above a ZBR-object type theorem. The common feature of these ZBR-type theorems is that the identities are each associated with a minor having only two or three non-port elements.

## 6. APPLICATIONS TO GRAPHS

The  $P$ -family of objects with matroids or oriented matroids abstraction helps us explore multiple kinds of graphs and their  $P$ -ported Tutte functions. We consider only separator-strong and strong Tutte functions. We first define graphs with labelled vertices and extend the Ellis-Monaghan and Traldi's ZBR theorem for graphs to them.

That is followed by an example whose Tutte functions do satisfy ported ZBR-matroid type theorems, and one where only a ported ZBR-object type theorem is satisfied.

Indeed, the immediate generalization of that theorem fails when  $P \neq \emptyset$ . That is because, even when  $P = \emptyset$  so the theorem holds, the Tutte function depends on the indecomposable graphs, not just their matroids. (In the  $P = \emptyset$  case, the indecomposables are  $E_k$ , the graphs with no edges and  $k$  vertices, so the matroids are all  $\emptyset$ .) [[[SHOULD THIS BE IN SUM SECTION?? Next, we prove that a separator-strong Tutte function is strong if and only if it is strong on the indecomposables. This is independent of whether the graphs satisfy a ZBR-type theorem or not]]] The section concludes with a generalization of the ZBR theorem for graphs with labelled vertices. The Tutte function must be equal on all indecomposables  $Q_i$  with the same partition  $\pi$  of vertex labels, assignment of port edges in  $Q_i$  to parts of  $\pi$ , and matroid or oriented matroid on the port edges. When  $P = \emptyset$ , all the matroids become  $\emptyset$  as in Ellis-Monaghan and Traldi's ZBR theorem for unlabelled graphs.

**Definition 38.** *A  $P$ -minor closed class of graphs  $\mathcal{G}$  is closed under deletion of loops not in  $P$ , contraction of isthmuses (i.e. coloops in the graph's matroid) not in  $P$ , and deletion and contraction of non-separating edges not in  $P$ .*

*A  $P$ -ported separator-strong Tutte function of graphs is defined on  $\mathcal{G}$  given with parameters in ring  $R$  and obeys the conditions (TA) and (TSSM) interpreted for (directed) graphs with their graphic, i.e. circuit (oriented) matroids.*

**6.1. Labelled Vertices.** Since our attitude is to retain the labels of the entities not removed by deletion or contraction, we begin with a definition of data preserved in  $P$ -minors that is more detailed than the number of graph components because it includes vertex labels. We remind the reader that our definition [11] of graph minor excludes deletion of isthmuses or of isolated vertices, and considers the removal of a loop to be a deletion.

**Definition 39** (Graphs with set-labelled vertices). *The elements of such a graph  $S(G) = E(G) \cup (P \cap S(G))$  are edges. The vertices are labelled with non-empty finite sets so the two sets labelling distinct vertices in one graph are disjoint. Only non-loop edges  $e \notin P$  can be contracted; when an edge is contracted, its two endpoints are replaced by one vertex whose label is the union of the labels of the two endpoints. Only non-isthmus edges  $e \notin P$  can be deleted; deletion doesn't change labels. The graph has its graphic matroid if it is undirected and its oriented graphic matroid if it is directed.*

**Definition 40** (Graphs with set-labelled components). *The elements of such a graph  $S(G) = E(G) \cup (P \cap S(G))$  are edges. The path-connected components are labelled by non-empty finite sets so two components in the same graph always have disjoint labels. In other words, the set labels of the components are a partition  $\pi_V$ . Only non-loop edges  $e \notin P$  can be contracted and only non-isthmus edges  $e \notin P$  can be deleted. The component labels are unchanged by these minor operations. Definition 39 specifies the matroids or oriented matroids.*

*Two such graphs  $G_1$  and  $G_2$  are equivalent means they have the same component label partition, the same subset of port edges, and if  $h \in S(G_1) \cap S(G_2)$  then the components of  $G_1$  and  $G_2$  that contain  $h$  have the same component labels.*

It is clear how to construct  $P$ -families of graphs with set-labelled vertices, or with set-labelled components, result, from closing a collection of vertex labelled graphs under  $P$ -minors.

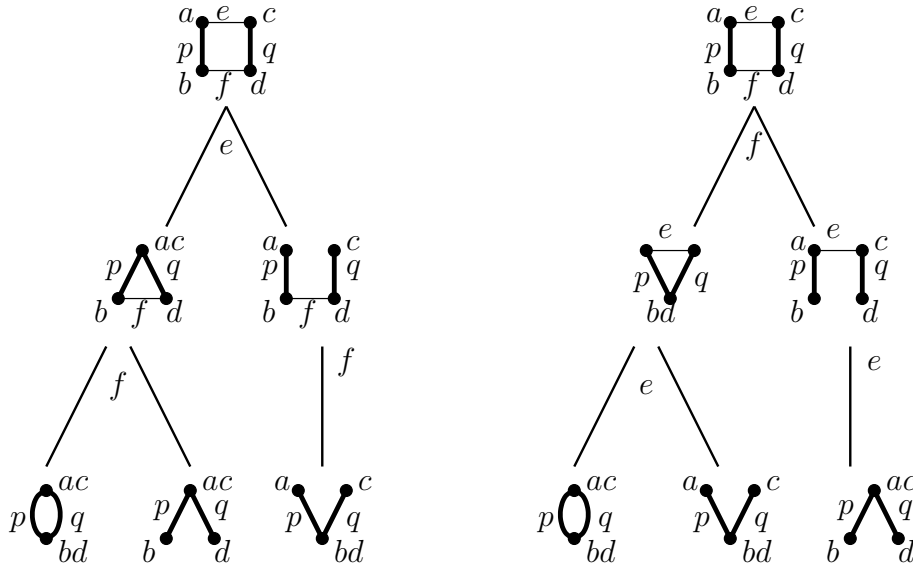
**Lemma 41.** *Let  $\mathcal{G}$  be a  $P$ -family of graphs with set-labelled components. Each part in the partition of  $\mathcal{G}$  into equivalent graphs is  $P$ -minor closed.*

It is clear that the Ellis-Monaghan and Traldi's arguments can be extended to conclude:

**Proposition 42.** *A class of graphs with set-labelled components  $\mathcal{G}$  is  $P$ -minor closed if and only if  $\mathcal{G}_\pi$  is  $P$ -minor closed for each non-empty equivalence class  $\mathcal{G}_\pi \subseteq \mathcal{G}$ .*

*A separator-strong  $P$ -ported Tutte function is well-defined on a  $P$ -minor closed class of graphs with set-labelled components if and only if it is well-defined on each non-empty equivalence  $\mathcal{G}_\pi \subseteq \mathcal{G}$ .*

**6.2. Examples.** The first example shows that graphs with vertices labelled by sets satisfy a ZBR-type theorem only when  $P = \emptyset$ .



The next example shows that additional conditions are required for a ZBR-type theorem for  $P$ -ported graphs even when the vertices are unlabelled.

The major complication due to port edges is illustrated with the graph  $G$  composed of five edges in the circle ordered  $(e, p, f, q, r)$  where the ports  $P = \{p, q, r\}$ .  $e$  and  $f$  are a series pair connected to  $P$ , but the  $P$ -quotient graphs  $Q_1 = G/e \setminus f$  and  $Q_2 = G/f \setminus e$  are different graphs, even though they have the same matroid.  $Q_1$  is the path  $qrp$  and  $Q_2$  is the path  $pqr$ . So, if this  $G \in \mathcal{G}$ , a necessary condition would be

$$(3) \quad T(Q_1)(x_e y_f + y_e X_f) = T(Q_2)(x_f y_e + y_f X_e).$$

**6.3. Strong Tutte Functions.** Strong parametrized Tutte functions for graphs are immediately generalized to  $P$ -ported graphs, and Theorem 30 applies directly. It generalizes part of the concluding proposition about strong Tutte functions for graphs in [11].

In the case when  $P = \emptyset$  and the vertices are unlabelled, we can prove a strengthening of part of Ellis-Monaghan and Traldi's Corollary 3.13. It is stronger because it does not require any additional hypotheses on  $\mathcal{G}$  to prove that all initial values that occur are the same idempotent.

**Corollary 43.** *Suppose  $T$  is parametrized Tutte function on a minor-closed class of graphs  $\mathcal{G}$  (note  $P = \emptyset$ .)  $T$  is strong if and only if there is an idempotent  $\alpha = \alpha^2 \in R$  and  $T(E_k) = \alpha$  whenever  $\mathcal{G}_k \neq \emptyset$ .*

*Proof.*  $M(E_k) = \emptyset$  for all  $k \geq 1$  and  $\emptyset \oplus \emptyset = \emptyset$ , so  $T(E_i)T(E_j) = T(E_k)$  whenever  $\mathcal{G}_i, \mathcal{G}_j$  and  $\mathcal{G}_k$  are all non-empty. Hence, if  $\mathcal{G}_k \neq \emptyset$  then  $T(E_k)T(E_k) = T(E_k) = \alpha$ . Further, if  $\mathcal{G}_j \neq \emptyset$  with  $j \neq k$ ,  $\alpha = T(E_k) = T(E_k)T(E_k) = T(E_j)$ .  $\square$

The other case of Corollary 3.13 [11] requires additional conditions for a Tutte function that is multiplicative on both disjoint union  $\amalg$  and one-point union  $\cdot$  to always be strong. Consider  $\mathcal{N} = \{E_3, E_4, E_5, \dots\}$ ,  $T(E_5) = 0$  and  $T(E_k) = 1$  for  $k \geq 3, k \neq 5$ . If  $T$  were strong then  $1 = T(E_3)T(E_4) = T(E_5) = 0$ , but this  $T$  is multiplicative on disjoint and one-point union because  $E_5$  cannot be expressed as either kind of union of graphs in  $\mathcal{G}$ . The other conditions are that  $\mathcal{G}_k \neq \emptyset$  for all  $k$  and that  $\mathcal{G}$  is closed under one-point unions and removal of isolated vertices.

**6.4. Matroidal Direct Sums in Graphs.** Ellis-Monaghan and Traldi derived relationships between the parametrized Tutte functions of two graphs  $G^1, G^2 \mathcal{G}$  and their disjoint union  $G^1 \amalg G^2$  and one-point union  $G^1 \cdot G^2$ , where  $\mathcal{G}$  was a minor closed class. Vertices were unlabelled and  $\mathcal{G}$  was partitioned into the minor-closed subclasses  $\mathcal{G}_k$  of graphs with  $k$  graph components. In this situation, each non-empty  $\mathcal{G}_k$  has exactly one indecomposable  $E_k$ , the graph with no edges and  $k$  vertices. The resulting ZBR theorem for graphs had  $\alpha_k = T(E_k)$  in place of  $\alpha = T(\emptyset)$  in the ZBR theorem for matroids, Theorem 2. Every Tutte function value had the form  $T(G) = \alpha_k r$ ,  $r \in R$  where  $G \in \mathcal{G}_k$ .

**Proposition 44.** (1) *If  $M(G) = M(G')$  then there exist  $r_i \in R$  such that  $T(G) = \sum T(Q_i)r_i$  and  $T(G') = \sum T(Q'_i)r_i$ , where  $Q'_i$  is the  $P$ -quotient of obtained by deleting and contracting respectively the same set of edges from  $G'$  as  $Q_i$  was obtained from  $G$ .*

*For  $P = \emptyset$ , this specializes to  $T(G) = \alpha_k r$  and  $T(G') = \alpha_{k'} r$  when  $G \in \mathcal{G}_k$  and  $G' \in \mathcal{G}_{k'}$  and*

(2) *Given any  $e \in E(G)$ , one expression for  $T(G)$  is  $\sum (r_i X_e + s_i Y_e) T(Q_i)$ .*

*For  $P = \emptyset$ , this specializes to  $T(G) = (r X_e + s Y_e) \alpha_k$ , when  $G \in \mathcal{G}_k$ .*

(3) *If  $G^1, G^2$  and  $G^1 \amalg G^2 \in \mathcal{G}$  and  $T(G^j) = \sum T(Q_i^j) r_i^j$  for  $i = 1, 2$ , then*

$$T(G^1 \amalg G^2) = \sum T(Q_i^1 \amalg T(Q_j^2) r_i^1 r_j^2$$

*and*

$$T(G^1)T(G^2) = \sum T(Q_i^1)T(Q_j^2) r_i^1 r_j^2.$$

*When  $P = \emptyset$ , this specializes to  $T(G^1 \amalg G^2) = \alpha_{k_1+k_2} r$  and  $T(G^1)T(G^2) = \alpha_{k_1} \alpha_{k_2} r$ .*

(4)  *$T$  is multiplicative with respect to  $\amalg$ , that is, if  $G^1, G^2$  and  $G^1 \amalg G^2$  are all in  $\mathcal{G}$ , then  $T(G^1 \amalg G^2) = T(G^1)T(G^2)$ , if and only if whenever  $Q_i, Q_j$  and  $Q_i \amalg Q_j$  are all  $P$ -quotients in  $\mathcal{G}$ , then  $T(Q_i)T(Q_j) = T(Q_i \amalg Q_j)$ .*

*When  $P = \emptyset$ , this specializes to  $T$  is multiplicative with respect to  $\amalg$  if and only if  $\alpha_{k_1} \alpha_{k_2} = \alpha_{k_1+k_2}$  whenever  $\mathcal{G}_{k_1}, \mathcal{G}_{k_2}$  and  $\mathcal{G}_{k_1+k_2}$  are all non-empty.*

*Proof.* The first two can be proved as in [11], or by analyzing Tutte computation trees

The last two follow from Proposition 32. □

[11] also gave analogous results for one-point union, which can be generalized in the same way.

**6.5. A Natural ZBR Theorem for Labelled Ported Graphs.** So, to generalize the ZBR theorem to graphs with port edges in a way that the edge deletion and contraction operations remain compatible with both the graphic matroid (or oriented matroid) structure and vertex structure, we have to add a hypothesis. To do that and get the exact generalization of Corollary 3.4, the generalized ZaslavskyBollobásRiordan theorem for graphs in [11] we require that the function values  $T(G_1) = T(G_2)$  whenever graphs  $G_1$  and  $G_2$  are path-connected and have the same matroids (or oriented matroids).

The separator-strong reductions applied to a graph clearly preserve the vertex set  $V$  and the partition of  $V$  induced by path-connectedness. The graphic 2-edge-connectivity structure is of course identical to the matroid connectivity structure. It is therefore natural, and motivated by Diao and Heteyi's applications to virtual knots [9], to define a separator-strong Tutte function of graphs as follows:



**Definition 45.** *Two graphs  $G_1, G_2$  with the same vertex set  $V$  and edge set  $S$  have the same connectivity structure if and only if they have the same graphic matroids (or oriented matroids) on  $S$  and the same partition of  $V$  induced by path-connectedness.*

*A separator-strong Tutte function  $T$  of graphs is a function from a  $P$ -family of graphs to ring  $R$  which satisfies (TA), (TSSM) applied to their graphic matroids and for which  $T(G_1) = T(G_2)$  if  $G_1$  and  $G_2$  have the same connectivity structure.*

Let us say two  $P$ -quotients  $Q_i$  and  $Q_j$  are equivalent if they have the same matroid (or oriented matroid) and have the same connectivity structure.

**Lemma 46.**  *$F$  is a Tutte function on a  $P$ -family of graphs if and only if  $F$  satisfies (TA) and (TSSM), and  $F(Q_i) = F(Q_j)$  for all pairs of equivalent  $P$ -quotients  $Q_i$  and  $Q_j$ .*

## 7. ACKNOWLEDGEMENTS

I wish to thank the Newton Institute for Mathematical Sciences of Cambridge University for hospitality and support of my participation in the Combinatorics and Statistical Mechanics Programme, January to July 2008, during which some of this work and much related subjects were reviewed and discussed.

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## 8. OLD GRAPH SECTION

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Strong parametrized Tutte functions for graphs are immediately generalized to  $P$ -ported graphs, and Theorem 30 applies directly. It generalizes part of the concluding proposition about strong Tutte functions for graphs in [11].

In the case when  $P = \emptyset$  and the vertices are unlabelled, we can prove a strengthening of part of Ellis-Monaghan and Traldi's Corollary 3.13. It is stronger because it does not require any additional hypotheses on  $\mathcal{G}$  to prove that all initial values that occur are the same idempotent.

**Corollary 47.** *Suppose  $T$  is parametrized Tutte function on a minor-closed class of graphs  $\mathcal{G}$  (note  $P = \emptyset$ .)  $T$  is strong if and only if there is an idempotent  $\alpha = \alpha^2 \in R$  and  $T(E_k) = \alpha$  whenever  $\mathcal{G}_k \neq \emptyset$ .*

*Proof.*  $M(E_k) = \emptyset$  for all  $k \geq 1$  and  $\emptyset \oplus \emptyset = \emptyset$ , so  $T(E_i)T(E_j) = T(E_k)$  whenever  $\mathcal{G}_i, \mathcal{G}_j$  and  $\mathcal{G}_k$  are all non-empty. Hence, if  $\mathcal{G}_k \neq \emptyset$  then  $T(E_k)T(E_k) = T(E_k) = \alpha$ . Further, if  $\mathcal{G}_j \neq \emptyset$  with  $j \neq k$ ,  $\alpha = T(E_k) = T(E_k)T(E_k) = T(E_j)$ .  $\square$

The other case of Corollary 3.13 [11] requires additional conditions for a Tutte function that is multiplicative on both disjoint union  $\amalg$  and one-point union  $\cdot$  to always be strong. Consider  $\mathcal{N} = \{E_3, E_4, E_5, \dots\}$ ,  $T(E_5) = 0$  and  $T(E_k) = 1$  for  $k \geq 3, k \neq 5$ . If  $T$  were strong then  $1 = T(E_3)T(E_4) = T(E_5) = 0$ , but this  $T$  is multiplicative on disjoint and one-point union because  $E_5$  cannot be expressed as either kind of union of graphs in  $\mathcal{G}$ . The other conditions are that  $\mathcal{G}_k \neq \emptyset$  for all  $k$  and that  $\mathcal{G}$  is closed under one-point unions and removal of isolated vertices.

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It is natural to study separator-strong Tutte functions on graphs. See [11] for details. Our aim is to extend Ellis-Monaghan and Traldi's discussion and results when a set of port edges  $P$  is given so deletion and contraction is done only of edges  $e \notin P$ . The extension of their definitions are

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**Definition 48.** *A  $P$ -minor closed class of graphs  $\mathcal{G}$  is closed under deletion of loops not in  $P$ , contraction of isthmuses (i.e., coloops in the graph's matroid) not in  $P$ , and deletion and contraction of non-separating edges not in  $P$ .*

*A  $P$ -ported separator-strong Tutte function of graphs is defined on  $\mathcal{C}$  given with parameters in ring  $R$  and obeys the conditions (TA) and (TSSM) interpreted for (directed) graphs with their graphic (oriented) matroids.*

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Some useful features of graphs are preserved in minors restricted as above. Any such feature partitions  $\mathcal{G}$  into  $P$ -minor closed classes, which of course have different indecomposables. Therefore, a Tutte function on  $\mathcal{G}$  is equivalent to Tutte functions on separate subclasses, and these Tutte functions have no relations between each other. For example, when  $P = \emptyset$  and vertices are unlabelled, [11] considered  $\mathcal{G}$  partitioned into the subclasses  $\mathcal{G}_k$  of graphs in  $\mathcal{G}$  with  $k$  graph components. The one indecomposable in a non-empty  $\mathcal{G}_k$  is  $E_k$ , the graph with  $k$  vertices and no edges. Their ZBR theorem for graphs therefore had  $\alpha_k = T(E_k)$  in place of  $\alpha = T(\emptyset)$  in the matroid version. Furthermore, one can express  $T(G_i) = \alpha_{k_i} R_i$  with some  $R_i \in R$  when  $G_i \in \mathcal{G}_{k_i}$ . The values on disjoint unions  $G_1 \amalg G_2$  and on one-point unions  $G_1 \cdot G_2$

(the graphs share one vertex) can then be written as  $T(G_1 \amalg G_2) = \alpha_{k_1+k_2} T(G_1) T(G_2)$  and  $T(G_1 \cdot G_2) = \alpha_{k_1+k_2-1} T(G_1) T(G_2)$ . These product formulas result from  $E(G_1) \cap E(G_2) = \emptyset$  and that the graphic matroid  $M(G_1 \amalg G_2) = M(G_1 \cdot G_2) = M(G_1) \oplus M(G_2)$  is a direct sum.

Generalizing remarks from [11], graphs  $G$  and  $G'$  might belong to different  $P$ -minor closed classes yet have the same matroid. We can formalize how the calculations of  $T(G)$  and of  $T(G')$  are identical by saying that they have isomorphic collections of Tutte computation trees that differ only at the leaves, and corresponding leaves have corresponding  $P$ -quotients. Combined with Proposition 32 this enables us to generalize the conclusions from [11] to graphs with port edges. We will take one step further to consider the vertices to be labelled. We end up with certain equivalence classes of graphs with among other things, the same matroids or oriented matroids, as our indecomposables.

---

The key fact about separator-strong Tutte decomposition of graphs (with  $P = \emptyset$ ) is that the indecomposables are the graphs  $E_k$  consisting of  $k$  isolated vertices. In [11], graph vertices are assumed to be not labelled. The identities and conditions in the ZBR theorem for graphs are identical to those in the theorem 2 for matroids except that the common factor in each identity is  $\alpha_k = T(E_k)$  instead of  $\alpha = T(\emptyset)$ , where  $k$  is the number of (graph) components in the graph to which the identity applies.

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The major complication due to port edges is illustrated with the graph  $G$  composed of five edges in the circle ordered  $(e, p, f, q, r)$  where the ports  $P = \{p, q, r\}$ .  $e$  and  $f$  are a series pair connected to  $P$ , but the  $P$ -quotient graphs  $Q_1 = G/e \setminus f$  and  $Q_2 = G/f \setminus e$  are different graphs, even though they have the same matroid.  $Q_1$  is the path  $qrp$  and  $Q_2$  is the path  $pqr$ . So, if this  $G \in \mathcal{G}$ , a necessary condition would be

$$(4) \quad T(Q_1)(x_e y_f + y_e X_f) = T(Q_2)(x_f y_e + y_f X_e).$$

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### 8.1. Natural $P$ -ported ZBR Theorem for Graphs. MOVED TO ABOVE..

## 9. OTHER WRITING

We will see that sometimes conditions adapted from those in the ZBR theorem for graphs and our  $P$ -ported ZBR theorem for matroids are equivalent to the well-definedness of Tutte functions, and sometimes they are not.

... Application to express properties of negative correlation of occurrences of graph edges in spanning trees and conjectures for forests.

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Sagan and Gessel demonstrated [14] how various interesting identities involving interval partitions can be proven by expanding the Tutte polynomial of a graph by following various schemes for picking which edge to delete or contract at each stage.

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NOTES: Matroid facts used: (1) A set in which each pair of elements is either a series pair or a parallel pair has the property that the whole set is either a series set or a parallel set. (2) Commutation of deletion and contraction on series pairs and on parallel pairs. (I.E., if  $e, f$  is series, then  $M/e \setminus f = M/f \setminus e$ , and dually). Commutation of ordinary pairs. (I.E.,  $M \setminus e/f = M/f \setminus e$ ,  $M/e/f = M/f/e$ , and  $M \setminus e \setminus f = M \setminus f \setminus e$ .)

The basis expansion originated by Tutte[22] for graphs and Crapo[8] for matroids depends on a particular but arbitrary ground set element order  $O$ . Each basis determines a term from the internal and external activities of elements with respect to that basis according to the ordering  $O$ . Our way to generalize is to restrict  $O$  to orders in which every port element is ordered before each  $e \in E$ . (We use the convention that the deleted/contracted element is the last, i.e., greatest element under order  $O$  eligible for reduction.)

Gordan and McMahon define[15] a “computation tree” to formalize the application of a subset of Tutte equations to a matroid and some of its minors. Each (Tutte) computation tree determines a polynomial in the parameters and point values. Therefore, when  $\mathcal{N}$  is in the domain of a Tutte function, each of these computation trees determine the same value. Computation trees are a way to give a basis expansion expression in terms of a more general definition of internal and external activities of elements with respect to a basis. The expansion is more general because it is based on any Tutte equation computation rather than on an element order  $O$ . We will extend to ported computation trees the classification[15] of elements as internally or externally, active or passive with respect to each path down the tree. In each case, the result is an interval partition of the boolean subset lattice of  $E$ .

the Tutte function is given by the following generalization of what is often called the corank-nullity polynomial (see section 4)

Tutte functions of this form are called “normal” in their extension for  $P \neq \emptyset$  is obvious. Indeed, the ZBR theorem is interesting because Tutte functions for some parameter choices might fail to be normal. Equation PGF generalizes to  $P$ -ported normal Tutte functions the corank-nullity expansion given in Corollary 5.2 of [11].

After verifying that the function given by this equation satisfies the  $P$ -ported Tutte equations (following our calculations in [5]), we see that this  $P$ -ported Tutte function of every matroid or oriented matroid exists for arbitrary assignments of initial values  $T(Q_i)$ . In [7] we utilized this fact to express an extensor-valued function  $F$  of  $P$ -ported graphic (and regular) oriented matroids by the assignments  $T(Q_i) = F(Q_i)$  and  $u = w = 0$  in (CNF). In that paper, we left open the question of whether for all parametrizations, not just those for which (CNF) is satisfied, can the initial values  $T(Q_i)$  be assigned arbitrarily. Theorem 4 answers that question: For the Tutte function to exist on  $\mathcal{C}$ , the  $T(Q_i)$  can be assigned arbitrarily so long as the  $(x, y, X, Y)$  parameters satisfy identities each of the form  $T(Q_i)p(x, y, X, Y) = 0$  where  $p$  is a 2nd or 3rd degree polynomial in the  $x, y, X$  and  $Y$ . In particular, no relation involves more than one  $T(Q_i)$ .

---

We can however take advantage of the fact that every  $P$ -quotient  $Q_i$  of  $M = M^1 \oplus M^2$  has a unique decomposition  $Q_i = Q_{i_1}^1 \oplus Q_{i_2}^2$  where  $Q_{i_1}^1$  is a  $P$ -quotient of  $M_1$  and  $Q_{i_2}^2$  is a  $P$ -quotient of  $M_2$ .

## FIGURES TO PLACE

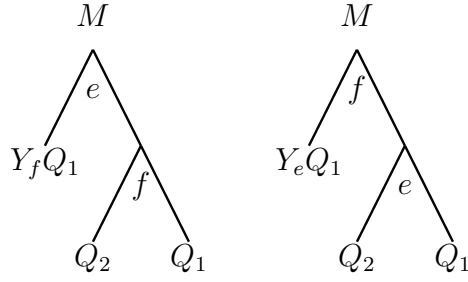


FIGURE 1.  $\{e, f\} = E(M)$  are parallel connected to  $P$ .  $Q_1 = M/e \setminus f = M/f \setminus e$ ,  $Q_2 = M \setminus \{e, f\}$ .

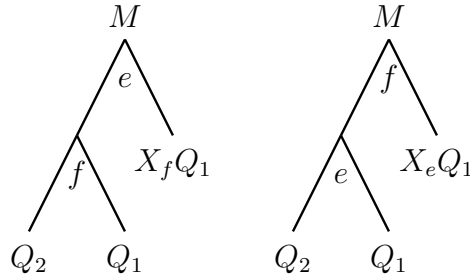
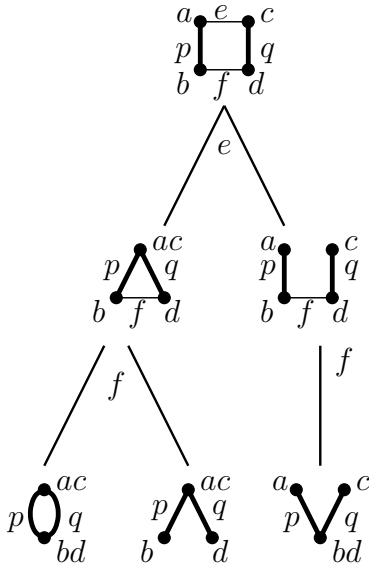


FIGURE 2.  $\{e, f\} = E(M)$  are series connected to  $P$ .  $Q_1 = M/e \setminus f = M/f \setminus e$ ,  $Q_2 = M/\{e, f\}$ .



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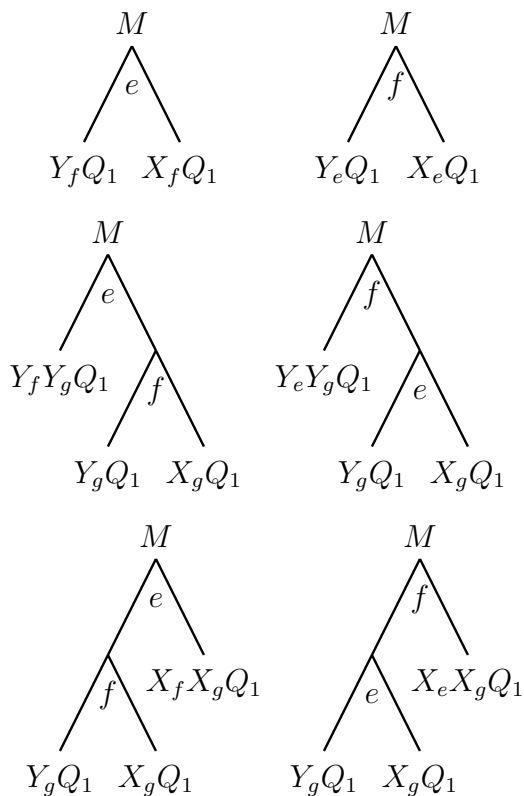


FIGURE 3. Cases from ZBR

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