

Ported or Relative Oriented Matroids and Electric Circuits

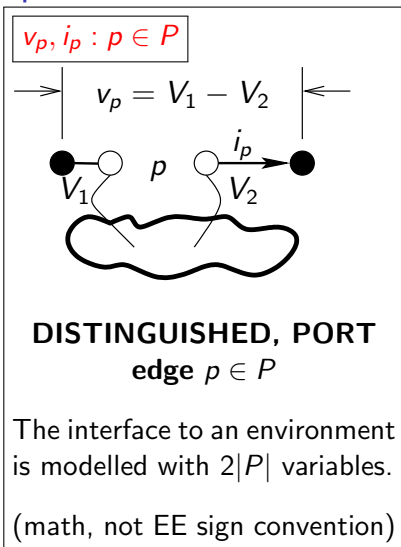
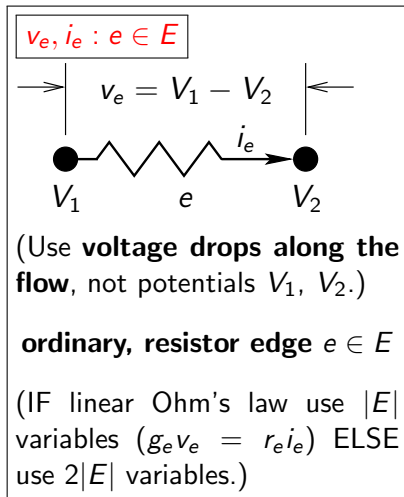
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Why Electricity, EE?

- ▶ Scholarly topic suggested by G.-C. Rota \approx 1980?.
- ▶ \approx 100 yrs. geometry-like intuition of circuit configurations known by engineers, EE books: “Intuitive Analog Circuit Design (2013)” [11]; “Non-linear Circuits” [5] translates to our Oriented Matroid pair model.
- ▶ Geometry of linear spaces and oriented matroids; Tutte decomp. w/ techniques from Barnabi, Brini and Rota’s Exterior Calculus [1])
- ▶ Real behavior \approx ideal plus perturbations, ideal constraints predict intended real behavior,
- ▶ Interesting, accessible, intuitively understandable intentional designs, applicable, easy to both simulate and build physically, dimension \approx 12 or 24, depending on formulation
- ▶ Analogs to chemical (and real algebraic geometry [8]), biological, mechanical, etc., random walks ...
- ▶ Merely one scalar non-linearity can cause chaos.

Kirchhoff (1847) [6] Maxwell (1891) [7] The equivalent resistance problem IS SOLVED by the Matrix Tree Theorem. (1) Let's POSE the problem: the VARIABLES



(2) Let's POSE the problem: EQUATIONS

- ▶ (KCL) $(i_e)_{e \in S}$ is a cycle (a flow).
- ▶ (KVL) $(v_e)_{e \in S}$ is a cocycle
- ▶ (constitutive Law) $i_e = g_e(v_e)$ non-linear, usually monotonic increasing $R \rightarrow R$.
(Sometimes use Ohm's approximation $i_e = g_e v_e$)

Combinatorics!

The signs $\{+, -, 0\}$ have oriented matroid structure (combinatorial, geometric, topological).

SOLUTION: Equiv. Resistance $\equiv -(v_p/i_p)$ observed at a port p by the environment EQUALS a Ratio of Spanning Tree Enumerators! (Port edge p locates the 2 *terminals*.)

$$-(\frac{v_p}{i_p}) = \frac{\text{WTS}(G/p)}{\text{WTS}(G \setminus p)} = \frac{\text{Matrix-Tree Det}(G/p)}{\text{Matrix-Tree Det}(G \setminus p)}$$

- ▶ “Maxwell’s rule” uses MatrTreeT on 2 DIFFERENT GRAPHS
(G/p and $G \setminus p$)

- ▶ Weighted Tree Sum (WTS) is a colored Tutte function:

$$\text{WTS}(G') = g_e \text{WTS}(G'/e) + r_e \text{WTS}(G \setminus e) \text{ for all } e \notin P$$

$$\text{WTS}(\text{coloop}(e)) = g_e$$

$$\text{WTS}(\text{loop}(e)) = r_e$$

Benefits of Multiple Ports

- ▶ One formula expresses $\binom{2|P|}{|P|}$ different Matrix Tree Theorems...
- ▶ ... long vertex-based proofs are shortened; Rayleigh inequalities too.
- ▶ Interesting **non-commutative ranges** of new ORIENTED MATROID Tutte invariants with pattern:

$$TF(N(P \cup E)) = F(N(P \cup E)/E)$$

(They distinguish DIFFERENT ORIENTATIONS of the SAME MATROID.)

- ▶ Formalize composition of systems [9], Tutte poly. splitting formulas.
- ▶ Label variables to observe
- ▶ Model practical devices (transistors, op amps)
- ▶ Align EE applications with knots (Ported = “Relative”) and combinatorial geometry (Ported = “Set Pointed”).

Constraint/Generator Duals and 2 Results.

- ▶ (Part 1) Technique:
Solution Space
 $=$
 \bigcap Constraint Subspaces
- ▶ **Result:** An exterior algebraic Tutte function:
Every Plucker coordinate of it satisfies a Matrix Tree Theorem.
This and det. formulas easily prove Rayleigh inequalities.

- ▶ (Part 2) Combine with:
Solution Space
 $=$
Closure(Set of Generators)
- ▶ **Result:** An oriented matroid pair model for some non-linear problem well-posedness.

Part 1) Coef. Matrix M in CONSTRAINTS $MX = 0$

The Tutte-like function $\mathbf{M}_E(\mathbf{N}) : \text{Extensors} \rightarrow \text{Extensors}$:

Given N (matrix), construct N^\perp with orthog. comp. row space.

Construct: ($G = \text{diag}(g_e)$, $R = \text{diag}(r_e)$)

$$M = \left[\begin{array}{c|c|c} N(P) & 0 & N(E)G \\ \hline 0 & N^\perp(P) & N^\perp(E)R \end{array} \right]$$

with columns labelled by $P_I \cup P_V \cup E$.

Extensor \mathbf{M} over $k[g_e, r_e](P_V \cup P_I \cup E)$ is the product of M 's **row vectors**. The contraction result $\mathbf{M}_E(\mathbf{N}) = \mathbf{M}/E$ appears:

$$\mathbf{M} = \mathbf{M}_E(\mathbf{N})\mathbf{e}_1\mathbf{e}_2 \cdots \mathbf{e}_{|E|} + (\cdots)$$

$\mathbf{M}_E(\mathbf{N})$ is our Tutte function $\mathbf{N} \rightarrow \text{Ext. Alg.}$

(Part 2) Common Covector Model

The cycle space of G_I GENERATES
the covectors of an
oriented matroid over $(E \cup P_I)$.

0

(signs indexed by E) (by P_I)

Non-linear monotone resistors CONSTRAIN SIGNS of
voltage drops (from \downarrow) and flows (from \uparrow)
TO BE EQUAL

(signs indexed by E) (by P_V)

The cocycle space of G_V GENERATES
the covectors of an
oriented matroid over $E \cup P_V$.

G_V
SOMETIMES EQUALS
 G_I

0

Return to the Part 1 Equation $MX = 0$

$$M = \left[\begin{array}{c|c|c} N(P) & 0 & N(E)G \\ \hline 0 & N^\perp(P) & N^\perp(E)R \end{array} \right]; \mathbf{M} = \mathbf{M}_E(\mathbf{N})\mathbf{e}_1\mathbf{e}_2\cdots\mathbf{e}_{|E|}+(\cdots)$$

ELIMINATE the variables indexed by E , leaving $2|P|$ variables labelled by P_I and P_V . ie, CONTRACT E . **Answer** A IS:

$$\mathbf{M}_E = \bigwedge_{\text{JOIN over rows}}^{\text{Exterior}} \left[\begin{array}{c|c} A_{I,I} & A_{I,V} \\ \hline A_{V,I} & A_{V,V} \end{array} \right] [\mathbf{p}_{I_1}, \dots, \mathbf{p}_{I_p}; \mathbf{p}_{V_1}, \dots, \mathbf{p}_{V_p}]^t$$

$$= \dots + C_i \mathbf{XXX} + \dots; \text{Equiv. Resistance} = \text{certain } C_i/C_j$$

All the other C_k 's have similar interpretations.

$\binom{2|P|}{|P|}$ **Matr. Tree Theorems:** Each $C_k(N)$ (a PRINCIPAL MINOR of M ABOVE!) $= g_e C_k(N/e) + r_e C_k(N \setminus e)$ ($e \notin P$, e not (co)loop).

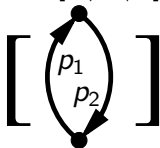
Each C_k is a signed weighted enumerator of forests satisfying **conditions ...**

Conditions (what sets F are enumerated by one det. C_i)

The **conditions** ... are on the rank, nullity of F and, WHAT ORIENTED MINOR is $G/F \setminus (E \setminus F)$, the minor with ONLY PORT EDGES from contracting F and deleting the other resistor edges, leaving the ports.

The conditions for a given C_k *sometimes* make all the signs the same (eg: C_i and C_j in 1-port equivalent resistance $R = C_i/C_j$) *Othertimes*, the oriented **P-minors** in the completed Tutte decomposition of C_k determine some $+$ and some $-$ signs.

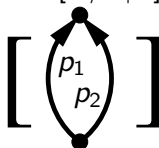
When $[G/F|P]$ is



the term is

$$+ g_F r_{E \setminus F}$$

When $[G/F|P]$ is



the term is

$$- g_F r_{E \setminus F}$$

Application: Rayleigh Identity, “Neg. Spanning Tree Correlation”

$\Gamma_e(G)$ is equivalent conductance across e . Rayleigh: $0 \leq \frac{\partial \Gamma_p}{\partial g_f} = \frac{\partial \frac{T_G}{T_{G/e}}}{\partial g_f}$

is equivalent to

$$0 \leq \frac{\partial T_G}{\partial g_f} T_{G/e} - T_G \frac{\partial T_{G/e}}{\partial g_f} = T_{G/f} T_{G/e} - T_G T_{G/e/f}$$

In fact,

$$T_{G/f} T_{G/e} - T_G T_{G/e/f} = \left(T_{G/e \& G/f}^+ - T_{G/e \& G/f}^- \right)^2$$

$T_{G/e \& G/f}^\pm$ enumerate the \pm common spanning trees.

Known Partial and Full Combinatorial Proofs

$$T_{G/f} T_{G/e} - T_G T_{G/e/f} = \left(T_{G/e \& G/f}^+ - T_{G/e \& G/f}^- \right)^2$$

$T_{G/e \& G/f}^\pm$ enumerate the \pm common spanning trees.

Choe (2004) proved essentially this using the vertex-based all-minors matrix tree theorem, combinatorial cases and Jacobi's theorem relating the minors of a matrix to the minors of its inverse..

Cibulka, Hladky, Lacroix and Wagner (2008) gave a completely bijective proof that utilizes some natural 2:2 and 2:1 correspondences.

Difficulty: Some terms on the left **cancel** and some reduce to terms with coefficients ± 2 .

Linear Alg./Oriented Matroid Proof of Rayleigh's Identity

Let R be the transfer resistance matrix for 2 ports across e and f .
Our result implies that

$$\det R = \begin{vmatrix} R_{ee} & R_{ef} \\ R_{fe} & R_{ff} \end{vmatrix} = + \frac{T_{G/e/f}}{T_G}$$

It and better-known results tell us

$$R_{ee} = \frac{T_{G/e}}{T_G}; \quad R_{ff} = \frac{T_{G/f}}{T_G}; \quad R_{ef} = R_{fe} = \frac{T_{G/e \& G/f}^+ - T_{G/e \& G/f}^-}{T_G}$$

$T_{G/f} T_{G/e} - T_G T_{G/e/f} = \left(T_{G/e \& G/f}^+ - T_{G/e \& G/f}^- \right)^2$ is
immediate after substituting these into

$$\det R = R_{ee} R_{ff} - (R_{ef})^2$$

The $+$ follows from physical grounds if the $g_e, r_e \geq 0$. Our characterization and proof are combinatorial.

New Rayleigh's Identities!

The same method generates identities and inequalities from

$$\begin{vmatrix} R_{ee} & R_{ef} & R_{eg} \\ R_{fe} & R_{ff} & R_{fg} \\ R_{ge} & R_{gf} & R_{gg} \end{vmatrix} = + \frac{T_{G/e/f/g}}{T_G} \geq 0$$

when all $r_{..}, g_{..} \geq 0$, ETC...

(Applications???)

Might the same methods address a much harder problem: The same inequality for forests instead of spanning trees?

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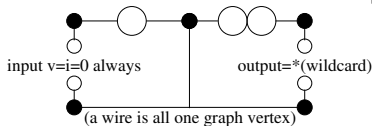
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Voltage and Current graphs N_V , N_I

“Voltage graph” N_V (EE [5, 10], NOT Gross, ...) represents KVL
 $\mathbf{v} \in \text{Cocycles W/ SOME } v_e \equiv 0$

“Current graph” N_I represents KCL $\mathbf{i} \in \text{Cycles}$
 WITH SOME FLOWS $\equiv 0$

- ▶ They are EQUAL GRAPHS for resistor networks.
- ▶ For networks with idealized amplifiers, they are not equal.



The output voltage and current are whatever makes the input voltage and current BOTH BE zero.

- ▶ (More) realistic amp. model = idealized amp. + resistors.

open

$$N_V = N \setminus e$$

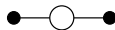
$$N_I = N \setminus e$$

short

$$N_V = N / e$$

$$N_I = N / e$$

nullator



$$N_V = N / e$$

$$N_I = N \setminus e$$

norator



$$N_V = N \setminus e$$

$$N_I = N / e$$

“Colors” are parameters on every Tutte decomposition step

The Bollobos/Riordan/Zaslavsky [2, 12], Traldi-Ellis-Monaghan [4], (sdc unpub) BRZ theory for well-definedness of “Relative Tutte Polynomials for Colored Graphs” ALL GOES THROUGH (Diao and Heteyi [3]): The 3 BRZ conditions on (colors, initial values) GENERALIZE TO 5; activity theory WORKS TOO, when based on linear orders on the non-port-elements.

In a nutshell

The 5 conditions \implies activities define an unambiguous Tutte function from the deletion/contraction and initial value formulas. Additional conditions \implies the Tutte function has a rank-nullity expansion.

(The rank-nullity conditions are satisfied in our application.)

To specify the activity/deletion-contraction linear order GLOBALLY is UNNECESSARY.

The Gordon/McMahon computation-tree-based activity theory also generalizes. (sdc).

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(17)

Our Tutte-like function $\mathbf{M}_E(\mathbf{N}) : \text{Extensors} \rightarrow \text{Extensors}$.

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$\mathbf{M}_E(\mathbf{N})$ is our Tutte function $\mathbf{N} \rightarrow \text{Ext. Alg.}$

What are the conditions like?

$$\text{Extensor} = \sigma(C)(a_1 \vee a_2 \vee \dots \vee a_d)$$

$$C = \dots + t + \dots$$

$t \leftrightarrow$ sets of (“contractible” non-port edges E_1, E_2 for which

$$N_V/E_1 \setminus (E \setminus E_1) = \text{certain OM}s \text{ on } P$$

AND

$$N_I/E_2 \setminus (E \setminus E_2) = \text{certain OM}s \text{ on } P$$

Transfer resistance might be 0 and might be $\neq 0$ iff

$$\exists E_1, E_2 \text{ so (diag) and (diag)}$$

(K_4 Wheatstone bridge diagram)

$$\frac{R_1}{R_2} > \frac{R_3}{R_4} \text{ neg } \frac{R_1}{R_2} < \frac{R_3}{R_4} \text{ pos}$$

Equivalent resistance is a coefficient ratio in an implicitly defined linear function

(diagram)

In other words

$$R_p i_p + v_p = 0$$

or dually,

$$\mathcal{B} = \{(i_p, v_p)\} = \{t(-1, R_p) | t \in R\}$$

The $2 \times d$ port variable constraint space, and its solution spaces, are d -dimensional.

We represent these spaces by carefully defined **extensors**, as Barnabei, Brini and Rota [1] term “decomposable antisymmetric tensors”

The solution extensor (not a ray) satisfies:

$$E(N) = \text{sign}(\dots)(g_e E(N/e) + r_e E(N \setminus e) \text{ for } e \notin P$$

Coefs

$$E(N) = \dots + C_i \mathbf{XXX} + \dots$$

$$R = C_i / C_j$$

All the other C_k 's have similar interpretations.

Each C_k is a determinant.

Each C_k is a signed weighted enumerator of forests satisfying
conditions ...

Each C_k satisfies

$$C_k(N) = g_e C_k(N/e) + r_e C_k(N \setminus e) \text{ for } e \notin P$$

Conditions

What is the nature of the conditions? We state them using the network's graphic oriented matroid.
(diagram-glob w/ ports)