

Resistive Networks, Linear Spaces and Tutte Polynomials

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The ± 1 vertex-edge incidence matrix of a graph fixes an arbitrary direction of each edge.

$$Mi = 0 \iff \{i_e\}_{e \in E} \text{ is a } \mathbf{flow} \text{ (of conserved current)}$$

Let's make M full row rank by removing redundant rows.

Let $T \subset E$ be any spanning tree. $E \setminus T$ is a cotree.

Each $\{i_e\}_{e \notin T} \in \mathbb{R}^{T^c}$ extends to a unique flow $\{i_e\}_{e \in E}$.

How? Take the $\{i_e\}_{e \in T}$ to **balance the excess at each vertex**.

That is: Row operations transform M to

$$\left[I \mid 0, \pm 1s \right]$$

if and only if T is a spanning tree.

M **represents** the (graphic) matroid whose bases are the spanning trees.

Electrical network problem: Solve for $i, v \in \mathbb{R}^n$

$M_V i = 0 \Leftarrow \text{equivalently} \Rightarrow i \in \text{Row space}(M_I), \text{ flows.}$

$M_I v = 0 \Leftarrow \text{equivalently} \Rightarrow v \in \text{Row space}(M_V), \text{ bonds.}$

$$F(i, v) = 0 \text{ locally rank } n.$$

First two are Kirchhoff's two laws: Combinatorial, assumed exact in electrical network applications.

Second are the **constitutive** constraints.

Linear one-port network: For all but one edge, **Ohm's law** is written

$$r_e i_e - g_e v_e = 0$$

For the one **port edge** p , demarking a pair of **terminal vertices**, either $i_p = 1$ then solve for $v_p = \text{equivalent resistance}$ by say eliminating the v_e, i_e for the resistors.

or $v_p = 1$ then solve for $i_p = \text{equivalent conductance} \dots$

Equiv. Resistance $\equiv -(v_p/i_p)$ observed at a port p by the environment

Theorem

(Kirchhoff, called “Maxwell’s rule”) Let g_T denote $\prod_{e \in T} g_e$, etc.
 G/p is G with edge p **contracted** (vertices identified).
 $G \setminus p$ is G with edge p **deleted**.

$$-\left(\frac{v_p}{i_p}\right) = \frac{\sum_{T: \text{spanning tree of } (G/p)} g_T r_{T^c}}{\sum_{T: \text{spanning tree of } (G \setminus p)} g_T r_{T^c}} = \frac{\text{Matrix-Tree Det}(G/p)}{\text{Matrix-Tree Det}(G \setminus p)}$$

It’s usually proved via the Matrix Tree Thm. on 2 DIFFERENT GRAPHS G/p and $G \setminus p$.

Pick any (resistor) edge e ; factor the sums: ($T^c = E \setminus e \setminus T$ below)

$$\frac{\text{WTS}(G/p)}{\text{WTS}(G \setminus p)} =$$

$$\frac{r_e \sum_{T: \text{spanning tree of } ((G \setminus e)/p)} g_T r_{T^c} + g_e \sum_{T: \text{spanning tree of } ((G/e)/p)} g_T r_{T^c}}{r_e \sum_{T: \text{spanning tree of } ((G \setminus e) \setminus p)} g_T r_{T^c} + g_e \sum_{T: \text{spanning tree of } ((G/e) \setminus p)} g_T r_{T^c}}$$

So, when we express our ratio by **the line** of all non-zero \mathbb{R} multiples of $(\text{WTS}(G/p), \text{WTS}(G \setminus p))$, carefully picked generators satisfy **Tutte decomposition**: For “ordinary” $e \neq p$

$$\begin{aligned} &(\text{WTS}(G/p), \text{WTS}(G \setminus p)) = \\ &r_e(\text{WTS}((G \setminus e)/p), \text{WTS}((G \setminus e) \setminus p)) + \\ &g_e(\text{WTS}((G/e)/p), \text{WTS}((G/e) \setminus p)) \end{aligned}$$

Interesting combinatorial theory emerges when quantities or relations of linear electrical network analysis are expressed as lines or more generally affine linear subspaces.

Our starting point is Thevenin's and Norton's theorems. They conclude that the voltage v_p and current i_p at a pair of terminals are characterized by the affine constraint $av_p + bi_p + c = 0$.

How the load line is used...

Outline

1. Spanning trees and equivalent (linear) resistance.
2. An exterior algebra (extensor) Tutte function and a (linear) resistance network's behavior projected on distinguished coordinates.
3. Rayleigh's inequalities.
4. Tutte polynomials on pairs and (linear) amplifier networks.
5. Distinguished graph vertices and splitting formulas.

Next steps

1. One (terminal-pair) port \rightarrow set of ports P .
2. 1-dim subspace of homogeneous coordinates of solutions $((v_p, i_p)) \rightarrow$ p-dim subspace of $k^{2|P|}$.
3. p-dim subspace \rightarrow EXTENSOR (decomposable exterior algebra, i.e., anti-symmetric tensor) with $\binom{2p}{p}$ Plucker coordinates (determinants).

Theorem

(After careful definitions...) For fixed P ,

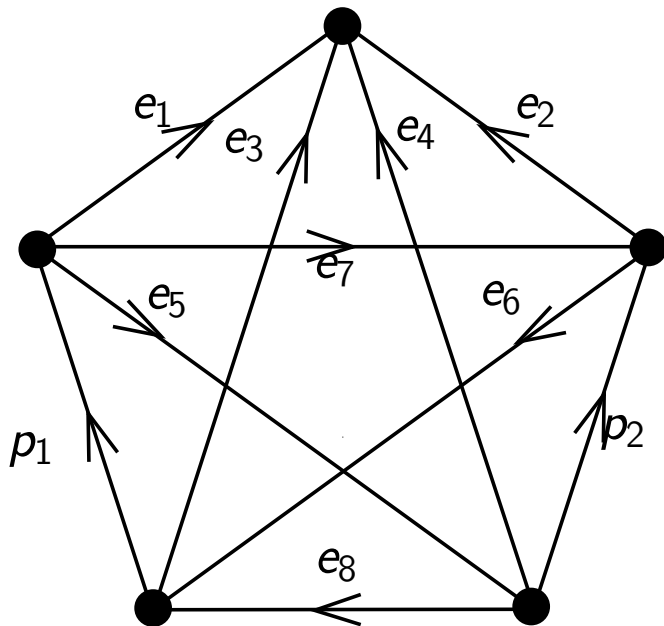
*each Plucker
coordinate*

and

*this extensor in the exterior
algebra*

satisfy weighed Tutte recursion, when $/e$ and $\backslash e$ are restricted to $e \notin P$.

Example



Example

Here's M of the electrical network equations $Mx = 0$. Kirchhoff's laws apply to all cycles and cocycles with $r_i x_{e_i}$ as voltage and $g_i x_{e_i}$ as current of resistor (not port) edges. TWO SEPARATE voltage and current variables are used for each port edge.

ip_1	ip_2	e_1	e_2	e_3	e_4	e_5	e_6	e_7	e_8	vp_1	vp_2
1	0	0	0	$+g_3$	0	0	$-g_6$	0	$-g_8$		
-1	0	$-g_1$	0	0	0	$+g_5$	0	$+g_7$	0		
0	+1	0	0	0	$+g_4$	$-g_5$	0	0	$+g_8$		
0	-1	0	$-g_2$	0	0	0	g_6	$+g_7$	g_8		
		$+r_1$	0	$-r_3$	0	0	0	0	0	1	0
		0	$+r_2$	0	$-r_4$	0	0	0	0	0	1
		$-r_1$	0	0	$+r_4$	$+r_5$	0	0	0	0	0
		0	$-r_2$	$+r_3$	0	0	$+r_6$	0	0	0	0
		$-r_1$	$+r_2$	0	0	0	0	$+r_7$	0	0	0
		0	0	$+r_3$	$+r_4$	0	0	0	$+r_8$	0	0

Top 4 rows: Basis for cocycle space. Represents graphic matroid.

Bot 6 rows: Basis for cycle space. Represents cographic matroid.

The Tutte Decomposition

For all choices denoted by ?? of the $\binom{2|P|}{|P|}$ size $|P|$ subsets of the $2|P|$ columns $\{ip_k, vp_k\}$, the matrices in the equation below are square.

So the elementary multilinearity of determinants means Tutte decomposition holds for all $e_i \notin P$:

$$\begin{array}{c} \text{??} \quad e_i \quad \text{??} \\ \left| \begin{array}{|c|c|c|} \hline & g_i \begin{array}{c} \vdots \\ \vdots \end{array} & \\ \hline & r_i \begin{array}{c} \vdots \\ \vdots \end{array} & \\ \hline \end{array} \right| = g_i \begin{array}{c} \text{??} \quad e_i \quad \text{??} \\ \left| \begin{array}{|c|c|c|} \hline & 1 \begin{array}{c} \vdots \\ \vdots \end{array} & \\ \hline & 0 \begin{array}{c} \vdots \\ \vdots \end{array} & \\ \hline \end{array} \right| + r_i \begin{array}{c} \text{??} \quad e_i \quad \text{??} \\ \left| \begin{array}{|c|c|c|} \hline & 0 \begin{array}{c} \vdots \\ \vdots \end{array} & \\ \hline & 1 \begin{array}{c} \vdots \\ \vdots \end{array} & \\ \hline \end{array} \right| \end{array}$$

(Technical detail: Define the Tutte function on all graphs with distinguished or port subset P so the det. signs are consistent with the decomposition.)

Rayleigh Identity which \Rightarrow inequality, “Neg. Spanning Tree Correlation”

$\Gamma_e(G)$ is equivalent conductance across e . Rayleigh: $0 \leq \frac{\partial \Gamma_p}{\partial g_f} = \frac{\partial \frac{T_G}{T_{G/e}}}{\partial g_f}$

is equivalent to

$$0 \leq \frac{\partial T_G}{\partial g_f} T_{G/e} - T_G \frac{\partial T_{G/e}}{\partial g_f} = T_{G/f} T_{G/e} - T_G T_{G/e/f}$$

Theorem

$$T_{G/f} T_{G/e} - T_G T_{G/e/f} = \left(T_{G/e \& G/f}^+ - T_{G/e \& G/f}^- \right)^2$$

$T_{G/e \& G/f}^\pm$ enumerate the \pm common spanning trees.

Choe, Cibulka, Hladky, Lacroix and Wagner gave bijective proofs;
we give det. based proofs and generalizations.

Linear Alg./Oriented Matroid Proof of Rayleigh's Identity

Let R be the transfer resistance matrix for 2 ports across e and f .
Our result implies that

$$\det R = \begin{vmatrix} R_{ee} & R_{ef} \\ R_{fe} & R_{ff} \end{vmatrix} = + \frac{T_{G/e/f}}{T_G}$$

It and better-known results tell us

$$R_{ee} = \frac{T_{G/e}}{T_G}; \quad R_{ff} = \frac{T_{G/f}}{T_G}; \quad R_{ef} = R_{fe} = \frac{T_{G/e \& G/f}^+ - T_{G/e \& G/f}^-}{T_G}$$

$T_{G/f} T_{G/e} - T_G T_{G/e/f} = \left(T_{G/e \& G/f}^+ - T_{G/e \& G/f}^- \right)^2$ is
immediate after substituting these into

$$\det R = R_{ee} R_{ff} - (R_{ef})^2$$

The $+$ follows from physical grounds if the $g_e, r_e \geq 0$. Our characterization and proof are combinatorial.

New Rayleigh's Identities!

The same method generates identities and inequalities from

$$\begin{vmatrix} R_{ee} & R_{ef} & R_{eg} \\ R_{fe} & R_{ff} & R_{fg} \\ R_{ge} & R_{gf} & R_{gg} \end{vmatrix} = + \frac{T_{G/e/f/g}}{T_G} \geq 0$$

when all $r_{..}, g_{..} \geq 0$, ETC...

(Applications???)

Might the same methods address a much harder problem: The same inequality for forests instead of spanning trees?

Pairs: The Common Covector Model

The cycle space of G_I GENERATES
the covectors of an
oriented matroid over $(E \cup P_I)$.

0

(signs indexed by E) (by P_I)

Non-linear monotone resistors CONSTRAIN SIGNS of
voltage drops (from \downarrow) and flows (from \uparrow)
TO BE EQUAL

(signs indexed by E) (by P_V)

The cocycle space of G_V GENERATES
the covectors of an
oriented matroid over $E \cup P_V$.

G_V
SOMETIMES EQUALS
 G_I

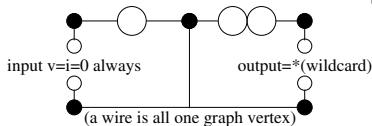
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Voltage and Current graphs G_V , G_I

“Voltage graph” G_V (EE [?, ?], NOT Gross, ...) represents KVL
 $\mathbf{v} \in \text{Cocycles W/ SOME } v_e \equiv 0$

“Current graph” G_I represents KCL $\mathbf{i} \in \text{Cycles}$
 WITH SOME FLOWS $\equiv 0$

- ▶ They are EQUAL GRAPHS for resistor networks.
- ▶ For networks with idealized amplifiers, they are not equal.



The output voltage and current are whatever makes the input voltage and current BOTH BE zero.

- ▶ (More) realistic amp. model = idealized amp. + resistors.

open

$$G_V = G \setminus e$$

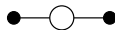
$$G_I = G \setminus e$$

short

$$G_V = G / e$$

$$G_I = G / e$$

nullator



$$G_V = G / e$$

$$G_I = G \setminus e$$

norator



$$G_V = G \setminus e$$

$$G_I = G / e$$

Distinguished graph vertices and splitting formulas

Let Q be a set of distinguished, labelled graph VERTICES, analogous to the distinguished port edges P

Theorem

Given graph $G(V \cup Q, E \cup P)$ let $T(G, P, Q)$ be the Tutte polynomial determined by restricting $/e$ and $\backslash e$ to $e \notin P$ AND carrying along the partition of Q defined by the components of the contracted edges.

Construct $G^Q(V \cup Q, E \cup P \cup P_Q)$ by adding to G a new vertex Z and the $|Q|$ new port edges from Z to each vertex in Q .

Then $T(G, P, Q)$ and $T(G^Q, P \cup P_Q)$ (the ported Tutte polynomial) determine each other by substitutions.

So we can use ported Tutte polys to express splitting formulas for Tutte polynomials of graph, beginning with Crapo [?] and continuing with Andrzejak [?], Bonin and de Meir [?], and Narayanan [?, ?].

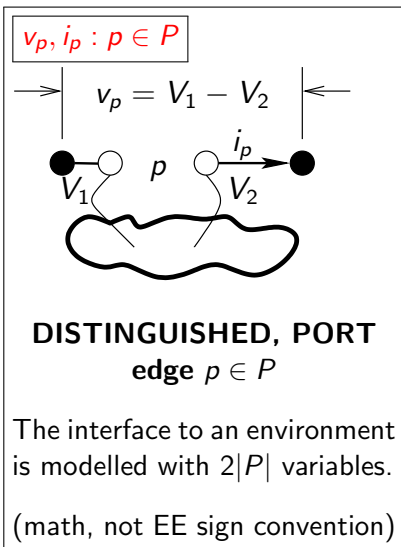
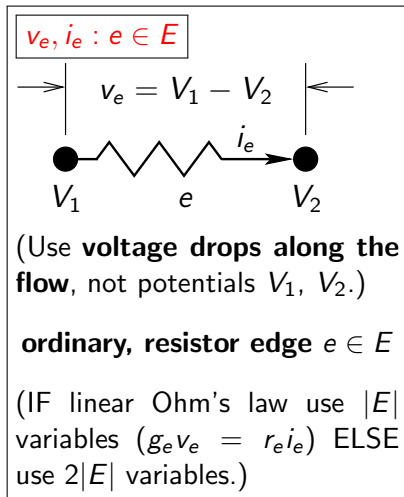
etc

Extra slides...

Why Electricity, EE?

- ▶ Scholarly topic suggested by G.-C. Rota \approx 1980?.
- ▶ \approx 100 yrs. geometry-like intuition of circuit configurations known by engineers, EE books: “Intuitive Analog Circuit Design (2013)” [?]; “Non-linear Circuits” [?] translates to our Oriented Matroid pair model.
- ▶ Geometry of linear spaces and oriented matroids; Tutte decomp. w/ techniques from Barnabi, Brini and Rota’s Exterior Calculus [?])
- ▶ Real behavior \approx ideal plus perturbations, ideal constraints predict intended real behavior,
- ▶ Interesting, accessible, intuitively understandable intentional designs, applicable, easy to both simulate and build physically, dimension \approx 12 or 24, depending on formulation
- ▶ Analogs to chemical (and real algebraic geometry [?]), biological, elastic/tensegrity strs. etc., random walks ...
- ▶ Merely one scalar non-linearity can cause chaos.

Kirchhoff (1847) [?] Maxwell (1891) [?] The equivalent resistance PROBLEM IS SOLVED by the Matrix Tree Theorem. (1) POSE! the **VARIABLES** or **COORDINATES**



(2) POSE: EQUATIONS. Preview the consequences.

- ▶ (KCL) $(i_e)_{e \in S}$ is a cycle (a flow).
- ▶ (KVL) $(v_e)_{e \in S}$ is a cocycle
- ▶ (constitutive Law) $i_e = g_e(v_e)$ non-linear, usually monotonic increasing $R \rightarrow R$.
(Sometimes use Ohm's approximation $i_e = g_e v_e$)

Combinatorics!

The signs $\{+, -, 0\}$ have a **DUAL-PAIR ORIENTED MATROID** structure (combinatorial, geometric, topological).

Engineering with amplifiers!

There's good unique solvability due to STRUCTURE, when the **NON-DUAL PAIR** (for voltages and currents) is ALMOST DUAL:
No common covectors.

Multiple Ports. (your stereo: 3=power plug & 2 speakers)

- ▶ One formula expresses $\binom{2|P|}{|P|}$ different Matrix Tree Theorems...
- ▶ ... long vertex-based proofs are shortened; Rayleigh inequalities too.
- ▶ Interesting **non-commutative ranges** of new ORIENTED MATROID Tutte invariants with pattern:

$$TF(N(P \cup E)) = F(N(P \cup E)/E)$$

(They distinguish DIFFERENT ORIENTATIONS of the SAME MATROID.)

- ▶ Ported/Relative OM Tutte Poly. terms **embed SPECIFIC MINORS** as **variables**, making proofs just with $\partial T / \partial x_e$ easier.
- ▶ Formalize composition of systems [?], Tutte poly. splitting formulas.
- ▶ Model practical devices (transistors, op amps); Label variables to observe.
- ▶ Align EE applications with knots [?] (Ported = “Relative”) and combinatorial geometry [?] (Ported = “Set Pointed”).

Constraint/Generator Duals and 2 Results.

- ▶ (Part 1) Technique:
Solution Space
=
 \bigcap Constraint Subspaces
- ▶ **Result:** An exterior algebraic Tutte function: Each of its $\binom{2|P|}{|P|}$ Plücker coordinates satisfies a Matrix Tree Theorem.
This and det. formulas easily prove Rayleigh inequalities.
- ▶ (Part 2) Combine with:
Solution Space
=
Closure(Set of Generators)
- ▶ To apply: An oriented matroid's COVECTOR SET encodes ALL POSSIBLE $(+, -, 0)$ coordinate behaviors or δ s.
- ▶ **Result:** An oriented matroid pair model for some non-linear problem (**AMPLIFIER!**) well-posedness. (How? Sign contradictions \Rightarrow a KERNEL= $\{(0)\}$.)

Part 1) Use Matrix M in CONSTRAINTS $MX = 0$ to get...

The Tutte-like function $\mathbf{M}_E() : \text{Extensor } \mathbf{N} \rightarrow \text{Extensor } \mathbf{M}_E(\mathbf{N})$.

(**STUDENT NOTE:** An EXTENSOR represents the row-space of an $r \times s$ r -rank matrix M by the $\binom{s}{r}$ -TUPLE of the DETERMINANTS of M 's $r \times r$ submatrices. **Plücker coords.**)

Given N (matrix), construct N^\perp with orthog. comp. row space.

Construct: $(G = \text{diag}(g_e), R = \text{diag}(r_e))$

$$M = \left[\begin{array}{c|c|c} N(P) & 0 & N(E)G \\ \hline 0 & N^\perp(P) & N^\perp(E)R \end{array} \right]$$

with columns labelled by $P_I \cup P_V \cup E$.

Extensor \mathbf{M} over $k[g_e, r_e](P_V \cup P_I \cup E)$ is the **\wedge -product** of M 's **row vectors**. The contraction result $\mathbf{M}_E(\mathbf{N}) = \mathbf{M}/E$ appears:

$$\mathbf{M} = \mathbf{M}_E(\mathbf{N})\mathbf{e}_1\mathbf{e}_2 \cdots \mathbf{e}_{|E|} + (\cdots)$$

$\mathbf{M}_E(\mathbf{N})$ is our Tutte function $\mathbf{N} \rightarrow \text{Ext. Alg.}$

Contracting means “Eliminate variables”

ELIMINATE the variables indexed by E , leaving $2|P|$ variables labelled by P_I and P_V . ie, CONTRACT E . **Answer M_E** IS:

$$M_E = \bigwedge_{\text{JOIN over rows}}^{\text{Exterior}} \left[\frac{A_{I,I} \mid A_{I,V}}{A_{V,I} \mid A_{V,V}} \right] [p_{I_1}, \dots, p_{I_p}; p_{V_1}, \dots, p_{V_p}]^t$$

$$= \dots + C_i \mathbf{XXX} + \dots; \text{Equiv. Resistance} = \text{certain } C_i/C_j$$

All the other C_k 's have similar interpretations.

$\binom{2|P|}{|P|}$ **Matr. Tree Theorems:** Each $C_k(N)$ (a PRINCIPAL MINOR of MATRIX **A** ABOVE!) = $g_e C_k(N/e) + r_e C_k(N \setminus e)$ ($e \notin P$, e not (co)loop).

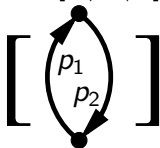
Each C_k is a signed weighted enumerator of forests satisfying **conditions ...**

Conditions (what sets F are enumerated by one det. C_i)

The **conditions** ... are on the rank, nullity of F and, WHAT ORIENTED MINOR is $G/F \setminus (E \setminus F)$, the minor with ONLY PORT EDGES from contracting F and deleting the other resistor edges, leaving the ports.

The conditions for a given C_k *sometimes* make all the signs the same (eg: C_i and C_j in 1-port equivalent resistance $R = C_i/C_j$) *Othertimes*, the oriented **P-minors** in the completed Tutte decomposition of C_k determine some $+$ and some $-$ signs.

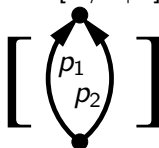
When $[G/F|P]$ is



the term is

$$+ g_F r_{E \setminus F}$$

When $[G/F|P]$ is



the term is

$$- g_F r_{E \setminus F}$$

Application: Rayleigh Identity, “Neg. Spanning Tree Correlation”

$\Gamma_e(G)$ is equivalent conductance across e . Rayleigh: $0 \leq \frac{\partial \Gamma_p}{\partial g_f} = \frac{\partial \frac{T_G}{T_{G/e}}}{\partial g_f}$

is equivalent to

$$0 \leq \frac{\partial T_G}{\partial g_f} T_{G/e} - T_G \frac{\partial T_{G/e}}{\partial g_f} = T_{G/f} T_{G/e} - T_G T_{G/e/f}$$

In fact,

$$T_{G/f} T_{G/e} - T_G T_{G/e/f} = \left(T_{G/e \& G/f}^+ - T_{G/e \& G/f}^- \right)^2$$

$T_{G/e \& G/f}^\pm$ enumerate the \pm common spanning trees.

Known Partial and Full Combinatorial Proofs

$$T_{G/f} T_{G/e} - T_G T_{G/e/f} = \left(T_{G/e \& G/f}^+ - T_{G/e \& G/f}^- \right)^2$$

$T_{G/e \& G/f}^\pm$ enumerate the \pm common spanning trees.

Choe (2004) proved essentially this using the vertex-based all-minors matrix tree theorem, combinatorial cases and Jacobi's theorem relating the minors of a matrix to the minors of its inverse..

Cibulka, Hladky, Lacroix and Wagner (2008) gave a completely bijective proof that utilizes some natural 2:2 and 2:1 correspondences.

Difficulty: Some terms on the left **cancel** and some reduce to terms with coefficients ± 2 .

“Colors” are parameters on every Tutte decomposition step

The Bollobos/Riordan/Zaslavsky [?, ?],
Traldi-Ellis-Monaghan [?, ?], (sdc unpub) BRZ theory for
well-definedness of “Relative Tutte Polynomials for Colored
Graphs” ALL GOES THROUGH (Diao and Heteyi [?]): The 3 BRZ
conditions on (colors, initial values) GENERALIZE TO 5; activity
theory WORKS TOO, when based on linear orders on the
non-port-elements.

In a nutshell

The 5 conditions \implies activities define an unambiguous Tutte
function from the deletion/contraction and initial value formulas.
Additional conditions \implies the Tutte function has a rank-nullity
expansion.

(The rank-nullity conditions are satisfied in our application.)

To specify the activity/deletion-contraction linear order
GLOBALLY is UNNECESSARY.

The Gordon/McMahon computation-tree-based activity theory also
generalizes. (sdc).

References I