"Optimizing Logistics Through Data" - Courier Nexus

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I. PROBLEM STATEMENT

As the logistics company grows, managing shipment flow becomes increasingly complex. Current tools, mainly Excel, can't handle the rising volume and complexity, leading to inefficiencies in tracking shipments, managing customer feedback, and resolving claims. This results in delayed deliveries, poor customer satisfaction, and operational bottlenecks.

Additionally, Excel's basic password protection is vulnerable, and manual backups risk data loss, highlighting the need for a secure, automated, and scalable database solution.

We propose implementing a scalable logistics management system that provides real-time tracking, efficient delivery monitoring, and streamlined claims processing. The new database will handle large datasets, support multiple users, and generate reports to optimize operations, improving data accuracy, decision-making, and overall efficiency.

II. TARGET USERS

The data-driven platform, Courier Nexus, targets four key user groups for logistics optimization.

Logistics Operations Team: They manage shipments by adding details, selecting shipment methods, and updating delivery statuses. For example, Sarah, a coordinator, inputs package info and tracks its progress.

Customer Service Representatives: They resolve inquiries and issues, using the system to track shipments and initiate claims. John, a representative, checks shipment status in real-time to assist customers.

Warehouse and Delivery Teams: They log delivery attempts and update package statuses. Mike, a delivery person, records failed delivery attempts, helping the logistics team reschedule.

Administrators and IT Team: They ensure security, performance, and system upkeep, managing permissions and backups. Maria, a database admin, protects sensitive data and ensures system compliance.

III. ENTITY - RELATIONSHIP DIAGRAM

A. Entities

1) Claims: claim_id (Primary Key, INTEGER): Unique identifier for each claim. This cannot be NULL. claim_status (TEXT): Represents the status of the claim (e.g., "Pending", "Resolved"). Can be NULL.

 $claim_date$ (DATE): The date when the claim was filed. Can be NULL.

resolution_date (DATE): The date when the claim was resolved. Can be NULL.

amount_claimed (REAL): The amount claimed in the process. Default value is 0.0.

shipment_id (Foreign Key, INTEGER): References the shipment_id in the Shipments table to associate a claim with a specific shipment. Cannot be NULL.

Primary Key Justification: The claim_id serves as the primary key to ensure that each claim is uniquely identifiable within the system. No duplicate claims are allowed, and every claim needs a distinct ID for traceability.

Foreign Key Action: When a shipment (referenced by shipment_id) is deleted, a cascade delete action can be applied, meaning the associated claims would also be deleted.

Query Query History									
1	1 SELECT * FROM Claims;								
Data	Data Output Messages Notifications								
=+ F V C V F SQL									
	claim_id [PK] integer	shipment_id /	claim_status text	claim_date /	resolution_date /	amount_claimed /			
1	1	1897	Resolved	2023-12-21	2024-01-29	259.59564			
2	2	387	Open	2022-03-09	2022-11-02	307.9183			
3	3	1030	Resolved	2023-06-21	2024-01-27	487.45978			
4	4	237	Open	2022-11-18	2024-02-16	753.4093			
5	5	212	Open	2024-06-15	2022-10-31	637.2745			
6	6	758	Open	2024-02-06	2023-03-15	814.48846			
7	7	2180	Open	2022-07-12	2024-09-23	515.70807			
8	8	144	Resolved	2023-01-14	2023-10-02	11.549548			
9	9	793	Resolved	2023-07-27	2022-09-18	839.8145			
10	10	148	Open	2023-02-21	2023-07-30	496.20346			

Fig. 1. Table Claims

2) Customer Feedback: feedback_id (Primary Key, INTEGER): Unique identifier for customer feedback. Cannot be NULL.

customer_id (INTEGER): Represents the customer providing feedback. Can be NULL.

rating (INTEGER): The rating given by the customer (e.g., 1–5). Can be NULL.

comments (TEXT): Feedback or comments from the customer. Can be NULL.

feedback_date (DATE): The date when the feedback was provided. Can be NULL.

shipment_id (Foreign Key, INTEGER): References the shipment_id in the Shipments table to tie feedback to a

specific shipment. Cannot be NULL.

Primary Key Justification: The feedback_id uniquely identifies each feedback entry, ensuring no duplicates and allowing feedback to be tracked individually.

Foreign Key Action: On deletion of a referenced shipment, the associated customer feedback can be set to NULL to maintain feedback records without shipment information.

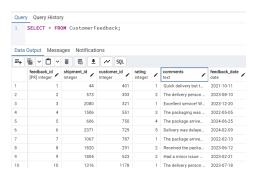


Fig. 2. Table Customer Feedback

3) **Delivery Attempts**: attempt_id (Primary Key, INTEGER): Unique identifier for each delivery attempt. Cannot be NULL.

attempt_date (DATE): The date of the delivery attempt. Can be NULL.

attempt_status (TEXT): The result or status of the delivery attempt (e.g., "Failed", "Successful"). Can be NULL.

notes (TEXT): Additional information about the attempt. Can be NULL.

shipment_id (Foreign Key, INTEGER): References
shipment_id in the Shipments table. Cannot be NULL.

Primary Key Justification: The attempt_id is necessary to uniquely identify each attempt made to deliver a shipment, ensuring that all records are traceable.

Foreign Key Action: On deletion of a shipment, a cascade delete can be applied, removing all associated delivery attempts for that shipment.

Quer	y Query Histor	у						
1	SELECT * FROM DeliveryAttempts;							
Data	Output Messa	iges Notificati	ons					
=+ (a) ∨ (b) ∨ (a) (b) (b) (c) (c)								
	attempt_id [PK] integer	shipment_id /	attempt_date /	attempt_status /	notes text			
1	1	1228	2022-07-14	Success	Receiver una			
2	2	1733	2022-03-01	Failed	Address not			
3	3	498	2024-07-01	Failed	Unable to del			
4	4	2089	2023-05-06	Failed	Package han			
5	5	1179	2022-10-23	Success	Unable to del			
6	6	1777	2022-10-14	Success	Delivery succ			
7	7	887	2023-01-29	Failed	Address not			
8	8	1813	2022-10-24	Success	Package han			
9	9	565	2023-01-27	Failed	Delivery succ			
10	10	592	2024-03-10	Success	Delivery succ			

Fig. 3. Table Delivery Attempts

4) **Delivery Locations**: location_id (Primary Key, INTEGER): Unique identifier for each delivery location.

Cannot be NULL.

longitude (REAL): The longitude of the delivery location. Can be NULL.

latitude (REAL): The latitude of the delivery location. Can be NULL.

delivery_date (DATE): The date of delivery at this location. Can be NULL.

shipment_id (Foreign Key, INTEGER): References shipment id in the Shipments table. Cannot be NULL.

Primary Key Justification: The location_id is used to ensure that every delivery location is unique, which is crucial for tracking where shipments were delivered.

Foreign Key Action: On deletion of a shipment, all associated delivery locations should also be deleted, enforcing a cascade delete policy.

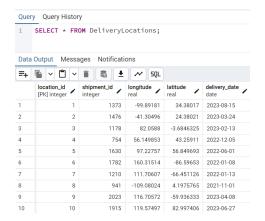


Fig. 4. Table Delivery Locations

5) Shipping Rates: rate_id (Primary Key, INTEGER): Unique identifier for each shipping rate. Cannot be NULL. shipping_method (TEXT): The shipping method used (e.g., "Air", "Ground"). Can be NULL.

weight_limit (REAL): The maximum weight allowed for this rate. Default value is 0.0.

base_price (REAL): The base price for this shipping rate. Default value is 0.0.

additional_cost_per_kg (REAL): Additional cost per kilogram over the weight limit. Default value is 0.0.

Primary Key Justification: The rate_id ensures each shipping rate is unique, as different rates may apply based on method and weight.

Foreign Key Action: No foreign keys are present, so no specific action is required for deletions.

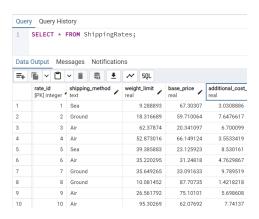


Fig. 5. Table Shipping Rates

6) Packages: package_id (Primary Key, INTEGER): Unique identifier for each package. Cannot be NULL. weight (REAL): Weight of the package. Default value is 0.0.

length (REAL): Length of the package. Default value is 0.0.

width (REAL): Width of the package. Default value is 0.0. height (REAL): Height of the package. Default value is 0.0.

content_description (TEXT): Description of the contents of the package. Can be NULL.

Primary Key Justification: The package_id uniquely identifies each physical package, allowing shipments to be tied to specific packages.

Foreign Key Action: No foreign keys are present, so no specific action is required for deletions.

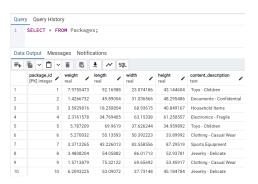


Fig. 6. Table Packages

7) Senders: sender_id (Primary Key, INTEGER): Unique identifier for each sender. Cannot be NULL. name (TEXT): Name of the sender. Can be NULL. address (TEXT): Address of the sender. Can be NULL. contact (TEXT): Contact information for the sender. Can be NULL.

Primary Key Justification: The sender_id uniquely identifies each sender, allowing packages to be tracked by who sent them.

Foreign Key Action: No foreign keys are present, so no specific action is required for deletions.

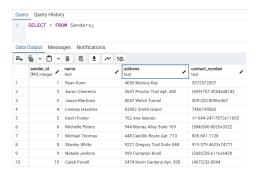


Fig. 7. Table Senders

8) Receivers: receiver_id (Primary Key, INTEGER): Unique identifier for each receiver. Cannot be NULL. name (TEXT): Name of the receiver. Can be NULL. address (TEXT): Address of the receiver. Can be NULL. contact (TEXT): Contact information for the receiver. Can be NULL.

Primary Key Justification: The receiver_id ensures that each receiver is unique, allowing the system to track who receives which shipment.

Foreign Key Action: No foreign keys are present, so no specific action is required for deletions.

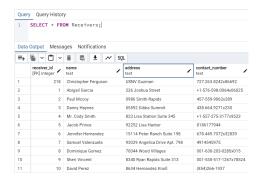


Fig. 8. Table Receivers

9) Service Areas: area_id (Primary Key, INTEGER): Unique identifier for each service area. Cannot be NULL. service_area_name (TEXT): Name of the service area. Can be NULL.

coverage_description (TEXT): Description of the area covered by the service. Can be NULL.

is_active (BOOLEAN): Indicates whether the service area is active. Default value is TRUE.

Primary Key Justification: The area_id uniquely identifies each service area, ensuring distinct coverage zones. **Foreign Key Action:** No foreign keys are present, so no specific action is required for deletions.

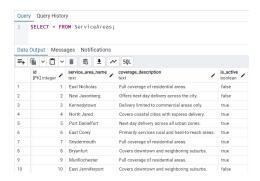


Fig. 9. Table Service Areas

10) Shipments: shipment_id (Primary Key, INTEGER): Unique identifier for each shipment. Cannot be NULL. package_id (Foreign Key, INTEGER): References package_id in the Packages table. Cannot be NULL. sender_id (Foreign Key, INTEGER): References sender_id in the Senders table. Cannot be NULL. receiver_id (Foreign Key, INTEGER): References receiver_id in the Receivers table. Cannot be NULL. shipment_method (TEXT): The method used for shipment. Can be NULL.

tracking_number (TEXT): Tracking number for the shipment. Can be NULL.

status (TEXT): The status of the shipment (e.g., "In Transit", "Delivered"). Can be NULL.

Primary Key Justification: The shipment_id ensures that each shipment is uniquely identified, allowing detailed tracking of every shipment in the system.

Foreign Key Action: For package_id, sender_id, and receiver_id, when the primary key is deleted, all associated shipments should be deleted using cascade delete.

Query Query History									
1	1 SELECT * FROM Shipments;								
Data	Output Me:	ssages Notifi	cations						
=+	· ·	v ii 💲	. ₩ 5	QL					
	shipment_id [PK] integer	package_id /	sender_id /	receiver_id /	shipping_met text	tracking_number text	status text		
1	1	84	1095	842	Sea	adbd0029-19ac-493e-bfa3	Pending		
2	2	1397	645	642	Ground	99a901c7-b8b0-41b5-b11c	Pending		
3	3	661	602	396	Air	8bb56f5a-6b8a-4d35-9b5e	Pending		
4	4	1325	275	615	Air	873a17ba-84cd-41bd-804f	Delivered		
5	5	201	719	226	Sea	8c322a82-8334-4f84-8c15	Pending		
6	6	430	1098	94	Ground	9e9e16cd-eb71-4b32-a284	Delivered		
7	7	641	682	405	Air	34f15b1b-6180-4a25-80ad	Delivered		
8	8	530	891	360	Ground	fc8c32ac-53d5-4f42-9341-f	Pending		
9	9	454	593	931	Sea	349d20dd-ce93-4303-a2ea	Delivered		
10	10	528	1054	1166	Air	df0065dc-05e6-4068-b547	Delivered		

Fig. 10. Table Shipments

B. Relations and Attributes

Table I below gives an overview of some of the important entities that describe the relationships and attributes of the logistics enterprise, including claims, shipments, and customer feedback. It provides full elaboration, featuring specific details for each entity on its main attributes and how those entities relate to one another through foreign keys. The following table shows the structure of this system.

TABLE I ENTITY RELATIONS AND ATTRIBUTES

Entity	Attributes
Claims	claim_id (primary key),
	claim_status, claim_date,
	resolution_date,
	amount_claimed, shipment_id
	(foreign key)
Customer Feedback	feedback_id (primary key),
	customer_id, rating, comments,
	feedback_date, shipment_id
	(foreign key)
Delivery Attempts	attempt_id (primary key),
	attempt_date, attempt_status,
	notes, shipment_id (foreign key)
Delivery Locations	location_id (primary key),
	longitude, latitude, delivery_date, shipment_id
	delivery_date, shipment_id
	(foreign key)
Shipping Rates	rate_id (primary key),
	shipping_method,
	weight_limit, base_price,
Packages	additional_cost_per_kg package_id (primary key), weight, length, width, height,
	weight, length, width, height,
	content description
Senders	sender_id (primary key), name,
	address, contact
Receivers	receiver_id (primary key), name,
	address, contact
Service Areas	area_id (primary key),
	service_area_name,
	coverage_description,
	is_active (Boolean)
Shipments	shipment_id (primary key),
•	package_id (foreign key), sender_id
	(foreign key), receiver_id (foreign key), shipment_method,
	tracking_number, status

C. Relationships Between Entities

Fig. 11 illustrates the relationships among the logistics system entities. Key relationships include:

- Shipments → Packages: A one-to-one relationship where each shipment corresponds to a single package.
- Shipments → Senders / Receivers: A many-to-one relationship; one sender or receiver can be linked to multiple shipments.
- Shipments → Claims: A one-to-one relationship; each shipment is associated with at most one claim.
- Shipments → Customer Feedback: A one-to-many relationship; a shipment can receive multiple feedback entries.
- Shipments → Delivery Attempts: A one-to-many relationship; multiple delivery attempts may be logged for a shipment.
- Shipments → Delivery Locations: A one-to-many relationship; several delivery locations may be recorded for a shipment.
- Service Areas → Shipments: Service Areas define zones for shipment deliveries, influencing shipping rates based on package weight and delivery method.

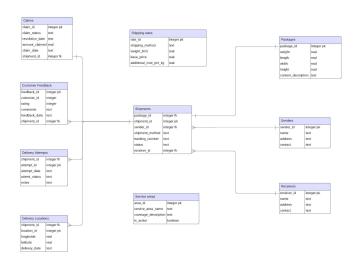


Fig. 11. Entity Relationship Diagram for Logistics Enterprise

IV. DATABASE GENERATION

In this project, Python scripts generated a synthetic dataset tailored to an imaginary domain. This flexible dataset supports SQL-based applications, enabling queries, trend analysis, and updates to meet specific project needs effectively.

A. Generating the table Packages

```
import sqlite3
from faker import Faker
import random
# Initialize Faker
fake = Faker()
# Connect to SQLite database (or create it)
conn = sqlite3.connect('courier_services5.db#) Initialize Faker
cursor = conn.cursor()
# Create Tables
cursor.execute('''CREATE TABLE IF NOT
EXISTS Packages (
    package id INTEGER PRIMARY KEY,
    weight REAL,
    length REAL,
    width REAL,
    height REAL,
    content_description TEXT
)''')
# Commit table creation
conn.commit()
# Helper functions to generate random data
def create_packages(n):
    for _ in range(n):
        cursor.execute('''INSERT INTO Packages Senders(sender_id),
```

```
(weight, length, width, height,
        content_description)
             VALUES (?, ?, ?, ?, ?)''',
            (random.uniform(1, 10),
            random.uniform(10, 100),
            random.uniform(10, 100),
            random.uniform(10, 100),
            fake.text(max nb chars=50)))
    conn.commit()
# Generate Data for a larger set
num_packages = 1500
create packages (num packages)
# Close the connection
conn.close()
```

```
17 V CREATE TABLE Packages (
18 package_id SERIAL PRIMARY KEY,
19 weight REAL,
20
21
             length REAL,
width REAL,
             height REAL
             content description TEXT
 Data Output Messages Notifications
 Query returned successfully in 179 msec.
```

Fig. 12. Creating the table Packages

B. Generating the table Shipments

```
import sqlite3
from faker import Faker
import random
fake = Faker()
# Connect to SOLite database
conn = sqlite3.connect('courier_services5.db')
cursor = conn.cursor()
# Create Tables
cursor.execute('''CREATE TABLE IF NOT
EXISTS Shipments (
    shipment_id INTEGER PRIMARY KEY,
    package_id INTEGER,
    sender_id INTEGER,
    receiver_id INTEGER,
    shipping_method TEXT,
    tracking_number TEXT,
    status TEXT,
    FOREIGN KEY (package_id) REFERENCES
    Packages (package_id),
    FOREIGN KEY (sender_id) REFERENCES
```

```
FOREIGN KEY (receiver_id) REFERENCES
   Receivers (receiver_id)
)''')
# Helper functions to generate random data
def create_shipments(n):
   for _ in range(n):
        cursor.execute('''INSERT INTO Shipments
        (package_id, sender_id, receiver_id,
        shipping_method, tracking_number, status)
            VALUES (?, ?, ?, ?, ?)''',
            (random.randint(1, num_packages),
            random.randint(1, num senders),
            random.randint(1, num_receivers),
            random.choice(['Air', 'Ground',
            'Sea'])
            ,fake.uuid4(),
            random.choice(['Pending',
            'Delivered'])))
   conn.commit()
# Generate Data for a larger set
```

```
do CREATE TABLE Shipments (
shipment_id SERIAL PRIMARY KEY,
package_id INTEGER,
sassen_id INTEGER,
receiver_id INTEGER,
tracking_number TEXT,
tracking_number TEXT,
FORETON KEY (package_id) REFERENCES Packages(package_id) ON DELETE CASCADE,
FORETON KEY (package_id) REFERENCES Packages(package_id) ON DELETE CASCADE,
FORETON KEY (receiver_id) REFERENCES Senders(sender_id) ON DELETE CASCADE,
FORETON KEY (receiver_id) REFERENCES Receivers(receiver_id) ON DELETE CASCADE,
Data Output Messages
Notifications
CREATE TABLE

Output Additional Control of the Control of the Cascade o
```

 $num_shipments = 2500$

Close the connection

conn.close()

create shipments (num shipments)

Fig. 13. Creating the table Shipments

As generated above, we have created the remaining tables in the database.

V. QUERY EXECUTION

SQL queries on the logistics database enable efficient analysis of shipments, claims, feedback, and delivery attempts using commands like GROUP BY, JOIN, sub-queries, and aggregate functions. These queries answer key questions, such as shipment counts by status or identifying claims. Below are examples:

A. Using GROUP BY to Get the Count of Shipments by Status

```
/*This query counts how many shipments are in each status
(e.g., "Pending", "Shipped", "Delivered").*/
SELECT status, COUNT(*) AS shipment_count
FROM Shipments
GROUP BY status;
```

Fig. 14. Using Group by function

Below is the output to the query in Fig. 14:

	status text	shipment_count bigint	
1	Delivered	1212	
2	Pending	1288	

Fig. 15. Output of the Group by function

B. Using a Sub-query to Find Shipments with Claims

```
/*This query retrieves details of all shipments.]
that have had claims filed against them using a sub-query to filter shipments.*/
SELECT shipment_id, package_id, sender_id, receiver_id, status
FROM Shipments
MMERE shipment_id ITM (

SELECT shipment_id ITM (

SELECT shipment_id FROM Claims ):
```

Fig. 16. Using Sub-query

Below is the output to the above query in Fig. 16:

	shipment_id [PK] integer	package_id /	sender_id /	receiver_id /	status text
1	12	460	632	267	Pending
2	19	1476	157	734	Pending
3	21	654	932	1082	Pending
4	22	941	280	428	Delivered
5	33	413	660	763	Delivered
6	37	110	97	889	Delivered
7	38	1024	483	19	Delivered
8	42	1412	589	330	Pending
9	43	675	519	995	Pending
10	57	1471	7	763	Delivered
Total rows: 459 of 459 Query complete 00:00:00.057 Ln					735, Col 3

Fig. 17. Output of the Sub-query

C. Using JOIN to Get Detailed Information About Shipments

```
/*This query uses an inner 30IH to get a detailed view of the shipments, including the sender's name and receiver's name.r/
SELECT s.nihpment (d. p.package_id, s.status, senume AS sender_name, r.name AS receiver_name
FROM Shipments a
30IM Receivers of 0s.package_id = p.package_id
30IM Senders so 0M s.package_id = se.sender_id
30IM Receivers 0 0s.receiver_id * r.receiver_id;
```

Fig. 18. Using Join function

Below is the output to the above query in Fig. 18:

	shipment_id a	package_id integer	status text	sender_name text	receiver_name text	
1	1	84	Pending	Joseph Pennington	Shane Williams	
2	2	1397	Pending	Christopher Nash Jr.	Kimberly Reynolds	
3	3	661	Pending	Brian Alvarez	Scott Booth	
4	4	1325	Delivered	Michael Leon	Michael Smith	
5	5	201	Pending	Brianna Cooper	Robert Soto	
6	6	430	Delivered	Omar Donovan	Brandon Schwartz	
7	7	641	Delivered	Ashlee Hamilton	Krista Jackson	
8	8	530	Pending	Crystal Blake	Nicole Bailey	
9	9	454	Delivered	Heather Mcgee	Vanessa Campbell	
10	10	528	Delivered	Seth Rivera	Mr. Corey Smith	
Tota	Total rows: 1000 of 2500					

Fig. 19. Output of the Join function

D. Using Aggregate Functions with a Sub-query to get Total Claims to specified Status

```
/*This query calculates the total value of all claims for shipments that have a specific status (c.g., "Delivered") using an aggregate function and sub-query.*/
SELECT SUMC.count_claime() AS total_claimed
FROM claims c
MEREE c.shipment_id IN (
SELECT shipment_id FROM Shipments MHERE status = 'Delivered');
;
;
;
```

Fig. 20. Using Aggregate function on Sub-query

Below is the output to the above query in Fig. 20:



Fig. 21. Output of the Aggregate function on Sub-query

VI. DECOMPOSITIONS

Decomposition ensures compliance with Boyce-Codd Normal Form (BCNF). After evaluating functional dependencies and applying the Chase Test, the schema was confirmed to already satisfy BCNF, with every determinant being a candidate key. This compliance applies to all relations, including packages, senders, recipients, and delivery-related entities.

VII. FUNCTIONAL DEPENDENCIES

Each table was evaluated to ensure that all determinants are a candidate key, thus satisfying the BCNF criteria. The following is the list of Functional Dependencies (FDs) identified for all tables:

```
package_id → {weight, length, width,
                     height, content description}
      receiver_id \rightarrow {name, address, contact_number}
        sender_id → {name, address, contact_number}
          area id \rightarrow {service area name,
                     coverage_description, is_active}
     shipment_id → {package_id, sender_id,
                     receiver id, shipping method,
                     tracking number, status}
tracking_number → {shipment_id, package_id,
                     sender id, receiver id,
                     shipping method, status}
         claim id → {shipment_id, claim_status,
                     claim_date, resolution_date,
                     amount_claimed}
      feedback id \rightarrow {shipment id,
                     customer_id, rating, comments,
                     feedback_date}
      attempt_id \rightarrow {shipment_id,
                     attempt_date, attempt_status, notes}
      location_id \rightarrow {shipment_id,
                     longitude, latitude, delivery date}
```

Each of these dependencies was analyzed for compliance with BCNF, ensuring that there are no partial or transitive dependencies. In every relation, the determinant is a superkey, making all tables BCNF-compliant. Below is the analysis of each table in terms of its superkeys:

A. Packages

The **Packages** table contains the attributes {package_id, weight, length, width, height, content_description}. The functional dependency package_id \rightarrow {weight, length, width, height, content_description} shows that package_id is a superkey, as it uniquely identifies all attributes in the table. Thus, the table satisfies BCNF.

B. Receivers

The **Receivers** table contains the attributes {receiver_id, name, address, contact_number}. The functional dependency receiver_id → {name, address, contact_number} demonstrates that receiver_id is a superkey, as it uniquely identifies all attributes in the table. Therefore, this table is BCNF-compliant.

C. Senders

The **Senders** table contains the attributes {sender_id, name, address, contact_number}. The functional dependency sender_id → {name, address, contact_number}

ensures that sender_id is a superkey. This guarantees that the table is in BCNF.

D. Service Areas

The Service Areas table contains the attributes {area_id, service_area_name, coverage_description, is_active}. The functional dependency area_id → {service_area_name, coverage_description, is_active} ensures that area_id is a superkey. As a result, the table satisfies BCNF.

E. Shipments

The **Shipments** table contains the attributes

{shipment_id, package_id, sender_id, receiver_id, shipping_method, tracking_number, status }.

The functional dependency

shipment $id \rightarrow all$ fields and tracking number \rightarrow all fields, where all fields represents all other attributes in the table, confirm that both shipment id and tracking number are superkeys.

Since all dependencies involve superkeys, the table satisfies BCNF.

F. Claims

The Claims table contains the attributes {claim id, shipment id, claim status, claim date, resolution_date, amount_claimed \}. The functional dependency claim_id → all_fields, where all fields represents all other attributes in the table, confirms that claim_id is a superkey. Thus, this table satisfies BCNF.

G. Customer Feedback

The Customer Feedback table contains the attributes {feedback_id, shipment_id, customer_id, rating, comments, feedback_date \}.

The functional dependency feedback id \rightarrow {shipment_id, customer_id, rating, comments, feedback_date} confirms that feedback_id is a superkey. This makes the table compliant with BCNF.

H. Delivery Attempts

The **Delivery Attempts** table contains the attributes {attempt_id, shipment_id, attempt_date, attempt_status, notes}. Applying the functional dependency unifies b, c, d with a, The functional dependency attempt_id → {shipment_id, attempt_date, attempt_status, notes} confirms that attempt_id is a superkey. Therefore, this table satisfies BCNF.

I. Delivery Locations

The **Delivery Locations** table contains the attributes {location_id, shipment_id, longitude, latitude, delivery_date}. The functional dependency location id \rightarrow {shipment_id, longitude, latitude, delivery_date} ensures that location_id is a superkey. As such, this table satisfies BCNF.

VIII. CHASE TEST

The Chase test was applied to verify whether the decomposition of each table ensures lossless joins and preserves functional dependencies. Below is the detailed analysis for all 10 tables.

A. Packages

The Packages table has the schema {package_id, weight, length, width, height, content_description} and the functional dependency

$$package_id \rightarrow \left\{ \begin{array}{l} weight, \ length, \ width, \\ height, \ content_description \end{array} \right.$$

Starting with:

$$\begin{split} t_1[\text{package_id}] &= a, \ t_1[\text{weight}] = b, \\ t_1[\text{length}] &= c, \ t_1[\text{width}] = d, \\ t_1[\text{height}] &= e, \ t_1[\text{content_description}] = f. \end{split}$$

Applying the functional dependency unifies b, c, d, e, f with a, resulting in one row.

Lossless Join: Yes.

Dependency Preservation: Yes.

B. Receivers

The Receivers table has the schema {receiver_id, name, address, contact_number} and the functional dependency receiver id \rightarrow {name, address, contact number}. Starting with

$$t_2[\text{receiver_id}] = a, \ t_2[\text{name}] = b, \ t_2[\text{address}] = c, \\ t_2[\text{contact_number}] = d.$$

Applying the functional dependency unifies b, c, d with a, resulting in one row.

Lossless Join: Yes. Dependency Preservation: Yes.

C. Senders

The **Senders** table has the schema {sender_id, name, address, contact_number} and the functional dependency $sender_id \rightarrow \{name, address, contact_number\}$. Starting with

$$t_3[{\rm sender_id}] = a, \ t_3[{\rm name}] = b, \ t_3[{\rm address}] = c, \\ t_3[{\rm contact_number}] = d.$$

resulting in one row.

Lossless Join: Yes. Dependency Preservation: Yes.

D. Service Areas

The Service Areas table schema is {area_id, service_area_name, coverage_description, is_active}, with the functional dependency: area_id → {other_fields}, where other_fields includes service_area_name, coverage_description, and is_active. Given: noitemsep

- $t_4[area_id] = a$
- Other fields unify with $a: b, c, d \rightarrow a$

This results in a single row for each unique area_id. Lossless Join: Yes. Dependency Preservation: Yes.

E. Shipments

The Shipments table schema is {shipment_id, package_id, sender_id, receiver_id, shipping_method, tracking_number, status}, with simplified functional dependencies: noitemsep

- shipment_id →
 {tracking_number, other_fields}
- tracking_number → {shipment_id, other_fields}

Given: noitemsep

• $t_5[{\rm shipment_id}] = a, \, t_5[{\rm package_id}] = b, \ t_5[{\rm sender_id}] = c, \, t_5[{\rm receiver_id}] = d, \ t_5[{\rm shipping_method}] = e, \ t_5[{\rm tracking_number}] = f, \, t_5[{\rm status}] = g.$

Applying FD1 unifies b, c, d, e, f, g with a, and FD2 unifies a, b, c, d, e, g with f.

Lossless Join: Yes. Dependency Preservation: Yes.

IX. SUMMARY OF 3NF AND BCNF EVALUATION

Third Normal Form (3NF) and Boyce-Codd Normal Form (BCNF) are designed to reduce redundancy and dependency anomalies, but BCNF imposes stricter rules by requiring all determinants to be superkeys. Each relation in the schema was evaluated and found to satisfy both BCNF and 3NF. Key examples include the Packages, Receivers, Senders, and **Service Areas** tables, where the functional dependencies have superkeys as determinants, ensuring compliance with BCNF. Similarly, the **Shipments** table, with candidate keys shipment_id and tracking_number, and the Claims table, with claim_id as a superkey, satisfy BCNF. The Customer Feedback, Delivery Attempts, and Delivery Locations tables also exhibit superkey-based dependencies. As all relations meet BCNF, the stricter form ensures no redundancy or dependency anomalies, making it unnecessary to relax the schema to 3NF.

X. REVISED ENTITY - RELATIONSHIP DIAGRAM

A. Entities

The schema meets BCNF requirements, requiring no further decomposition. Key entities—Packages, Receivers, Senders, Service Areas, Shipments, Claims, Customer Feedback, Delivery Attempts, and Delivery Locations—are robust and efficient, free of partial or transitive dependencies.

B. Relations and Attributes

The relations and their attributes also remain consistent with the previous schema, as the design meets BCNF requirements. Each table is normalized with its respective primary keys and attributes. These relations are mentioned in Table I.

C. Relationship Between Entities

Relationships remain intact, with Shipments linking Packages, Senders, and Receivers. Claims, Feedback, Delivery Attempts, and Locations are tied to Shipments via foreign keys, ensuring consistency and referential integrity.

D. Constraints

- 1) Primary Keys: Primary key constraints are defined in the table(s) to ensure unique identification of each record.
- 2) Foreign Keys: Foreign key constraints are defined to enforce relationships between table(s), ensuring referential integrity.
- 3) Not Null Constraints: Not null constraints are applied to specific fields to ensure data completeness and maintain data integrity.



Fig. 22. Indexing

E. Redundancy Analysis

The BCNF-compliant schema minimizes redundancy, avoids data anomalies, and optimizes storage efficiency and consistency by ensuring functional dependencies depend only on superkeys.

XI. CHALLENGES AND OPTIMIZATION TECHNIQUES FOR MANAGING LARGE DATASETS

Handling larger datasets led to performance issues, particularly slower query execution times due to the table size and complex queries. Operations involving joins and filters on non-indexed attributes, like $tracking_number$, status, and

To optimize performance, indexes were added on frequently queried fields, such as a unique index on $tracking_number and non-clustered indexes on$

Indexing reduced query execution time from seconds to milliseconds, minimized full table scans, and optimized resource usage, improving system responsiveness. Indexing increased storage requirements and slightly slowed write operations, but the benefits in read-heavy operations outweighed these trade-offs. This experience highlighted the importance of optimizing query patterns, selecting appropriate indexes, and periodic index maintenance.

XII. SQL OPERATIONS AND QUERY EXECUTION RESULTS

A. Insertion Queries

For the Insert Query - I, we inserted a row in the Packages table as seen in Fig . 22.



Fig. 23. Insert Query - I

For the Insert Query - II, we inserted a row in the Senders table as seen in Fig. 23.

```
S VINSERT INTO Senders (id, name, address, contact_number)

6 VALUES (1201, 'Alice Johnson', '123 Elm Street', '123-456-7890');

7

Data Output Messages Notifications

INSERT 0 1

Query returned successfully in 84 msec.
```

Fig. 24. Insert Query - II

B. Update Queries

For the first update query, we updated the content description of the package id '125' to 'Electronics' from 'Books - Hardcover'.



Fig. 25. Update Query - I (Pre update)

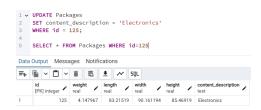


Fig. 26. Update Query - I (Post update)

For the second update query, we increased the base price present in the table 'ShippingRates' by 10% for all the rows where the given weight limit was more than the average weight limit of all the shipping methods.



Fig. 27. Update Query - II

C. Select Queries

In query 1, we retrieve detailed information about shipments, including sender and receiver details, where the shipment status is not "Delivered."

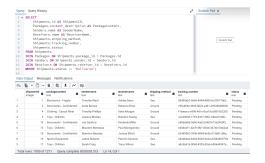


Fig. 28. Query - I

In query 2, we get the total number of shipments delivered and pending, grouped by shipping method.

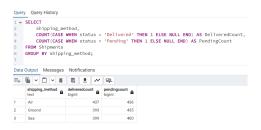


Fig. 29. Query - II

In query 3, we get the top 5 highest-rated feedback comments for shipments, including sender names and feedback details.

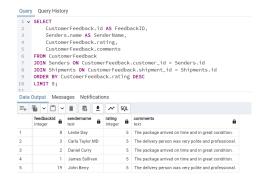


Fig. 30. Query - III

In query 4, we find the two most recent claims filed for each shipment and rank them by claim date.

Query Query History 1 v WITH RankedClaims AS (SELECT c.shipment_id, c.id AS claim_id, c.claim_status, c.claim_date, c.amount_claimed, RANK() OVER PARTITION BY c.shipment_id ORDER BY c.claim_date DESC) AS rank_position SELECT claim_id, claim_status, claim_date, amount claimed FROM RankedClaims
WHERE rank_position <= 2 ORDER BY shipment_id, rank_position LIMIT 20: Data Output Messages Notifications =+ **1** ∨ 1 ∨ 1 8 ± **×** SQL shipment_id a claim_id claim_id claim_status a claim_date a amount_claimed a text 242 Open 2023-09-22 493.3209 707.30774 2024-05-11 357 Resolved 2024-02-18 885.921 2023-11-23 699.79095 229 Open 177 Open 2024-09-17 990,67206 927.9941 2022-02-21 13 Open 180 Resolved 2024-08-17 668.5279 276 Resolved 2023-08-29 417.74557 157 Resolved 2024-06-06 784.66425 396 Resolved 2024-02-27 790.14307 11 79 190 Open 2022-05-22 Total rows: 20 of 20 Query complete 00:00:00.236 Ln 24, Col 1 483.8005

Fig. 31. Query - IV

D. Deletion Queries

In the first deletion query, we removed entries from the Claims table for claims that are older than one year and have been resolved.

```
Query Query History

1 v DELETE FROM Claims
2 WHERE claim_status = 'Resolved'
3 AND claim_date < CURRENT_DATE - INTERVAL '1 year';
4

Data Output Messages Notifications

DELETE 177

Query returned successfully in 130 msec.
```

Fig. 32. Deletion Query - I

In the second deletion query, we delete claims where the claimed amount exceeds the shipment's base price.

```
Ouery Overy History

Vollette FROM Claims

WHERE Shipments of IN (

SELECT Shipments.id IN (

SELECT Shipments.id IN (

SOUTH Shipments SOUTH Shipments.shipping_method = ShippingRates.shipping_method = Shipping_method = Shipping_m
```

Fig. 33. Deletion Query - II

XIII. QUERY EXECUTION ANALYSIS

Targeted Indexing: Indexes were added to key columns like sender_id, receiver_id, and shipment_id in the Shipments table, and claim_status and feedback_date in Claims and CustomerFeedback tables.

Faster Queries: Indexes reduced table scans, speeding up data retrieval and cutting execution time.

Smoother Joins: Indexes on foreign keys improved join efficiency between tables like Shipments, Claims, and CustomerFeedback.

Better Filtering and Aggregations: Indexes on claim_status and feedback_date sped up filtering and aggregation queries.

Less System Strain: Faster queries reduced system load, enabling better performance as data scales.

A. Query 1



Fig. 34. Query - I (Pre Improvement)

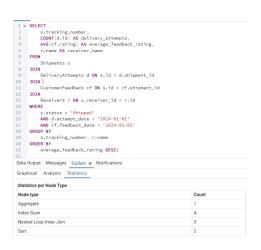


Fig. 35. Query - I (Post Improvement)

B. Query 2



Fig. 36. Query - II (Pre Improvement)

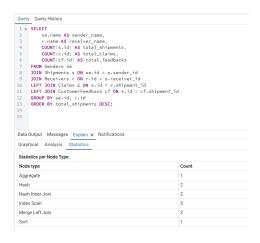


Fig. 37. Query - II (Post Improvement)

C. Query 3

Fig. 38. Deletion Query - III (Pre Improvement



Fig. 39. Deletion Query - III (Post Improvement

XIV. BONUS TASK: WEB APPLICATION FOR DATABASE QUERY VISUALIZATION

Below are a few screenshots of the running website required for the bonus task.



Fig. 40.



Fig. 41.

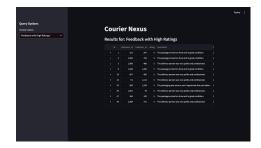


Fig. 42.

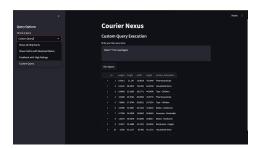


Fig. 43.

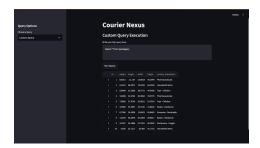


Fig. 44.