

# CS/CE 6378: Advanced Operating Systems

## Section 001

### Project 1

Instructor: Neeraj Mittal

Assigned on: Tuesday September 5, 2023

This is an individual project. *Code sharing among students as well as copying code from the Internet is strictly prohibited and will result in disciplinary action being taken.* Each student is expected to write their *own code* and demonstrate the operation of this project to the instructor or the TA. You will be graded not only on the correct behavior of your project, but also on your responses to questions concerning your implementation which the instructor or TA may ask during your demo.

You can do this project in C, C++ or Java. Since the project involves socket programming, you can only use machines `dcXX.utdallas.edu`, where  $XX \in \{01, 02, \dots, 45\}$ , for running the program. Although you may develop the project on any platform, the demonstration has to be on `dcXX` machines; otherwise you will be assessed a penalty of 20%.

## 1 Project Description

This project consists of four parts.

### 1.1 Part 1

Implement a distributed system consisting of  $n$  nodes, numbered 0 to  $n - 1$ , arranged in a certain topology. The topology and information about other parameters will be provided in a configuration file.

All channels in the system are bidirectional, reliable and satisfy the first-in-first-out (FIFO) property. You can implement a channel using a reliable socket connection (with TCP or SCTP). For each channel, the socket connection should be created at the beginning of the program and should stay intact until the end of the program. All messages between neighboring nodes are exchanged over these connections.

All nodes execute the following protocol:

- Initially, each node in the system is either *active* or *passive*. **At least one node must be active at the beginning of the protocol.**
- While a node is *active*, it sends anywhere from `minPerActive` to `maxPerActive` messages, and then turns *passive*. For each message, it makes a uniformly random selection of one of its neighbors as the destination. Also, if the node stays *active* after sending a message, then it waits for at least `minSendDelay` time units before sending the next message.

- Only an *active* node can send a message.
- A *passive* node, on receiving a message, becomes *active* if it has sent fewer than `maxNumber` messages (summed over all *active* intervals). Otherwise, it stays *passive*.

We refer to the protocol described above as the MAP protocol.

## 1.2 Part 2

Implement the Chandy and Lamport's protocol for recording a consistent global snapshot as discussed in the class. Assume that the snapshot protocol is always initiated by node 0 and all channels in the topology are bidirectional. Use the snapshot protocol to detect the termination of the MAP protocol described in Part 1. The MAP protocol terminates when all nodes are *passive* and all channels are *empty*. To detect termination of the MAP protocol, augment the Chandy and Lamport's snapshot protocol to collect the information recorded at each node at node 0 using a converge-cast operation over a spanning tree. The tree can be built once in the beginning or on-the-fly for an instance using `MARKER` messages.

Note that, in this project, the messages exchanged by the MAP protocol are *application* messages and the messages exchanged by the snapshot protocol are *control* messages. The rules of the MAP protocol (described in Part 1) only apply to application messages. They do not apply to control messages.

## 1.3 Part 3

To test that your implementation of the Chandy and Lamport's snapshot protocol is correct, implement Fidge/Mattern's vector clock protocol described in the class. The vector clock of a node is part of the local state of the node and its value is also recorded whenever a node records its local state. Node 0, on receiving the information recorded by all the nodes, uses these vector timestamps to verify that the snapshot is indeed consistent. Note that only application messages will carry vector timestamps.

## 1.4 Part 4

Design and implement a protocol for bringing all nodes to a halt after node 0 has detected termination of the MAP protocol.

# 2 Submission Information

To ensure that you are able to complete the project on time, we are asking you to submit your code as per the following timeline:

- (i) Thursday, September 14, 2023: Part 1
- (ii) Thursday, September 21, 2023: Part 1 and Part 3
- (iii) Tuesday, October 3, 2023: Complete Project

All submissions will be through eLearning. Submit all the source files necessary to compile the program and run it. Also, submit a `README` file that contains instructions to compile and run your program. You should also include launcher (`launcher.sh`) and cleanup (`cleanup.sh`) scripts

which run your program on the machines as configured and cleanup your processes. You may need to modify the launcher and cleanup scripts that are on eLearning to work for this project. If you have questions about the script, please contact the TA.

### 3 Configuration Format

Your program should run using a configuration file in the following format:

The configuration file will be a plain-text formatted file no more than 100kB in size. Only lines which begin with an unsigned integer are considered to be valid. Lines which are not valid should be ignored. The configuration file will contain  $2n + 1$  valid lines. The first valid line of the configuration file contains **six** tokens. The first token is the number of nodes in the system. The second and third tokens are values of **minPerActive**, and **maxPerActive** respectively. The fourth and fifth tokens are values of **minSendDelay** and **snapshotDelay**, in milliseconds. **snapshotDelay** is the amount of time to wait between initiating snapshots in the Chandy and Lamport algorithm. The sixth token is the value of **maxNumber**. After the first valid line, the next  $n$  lines consist of three tokens. The first token is the node ID. The second token is the host-name of the machine on which the node runs. The third token is the port on which the node listens for incoming connections. After the first  $n + 1$  valid lines, the next  $n$  lines consist of a space delimited list of at most  $n - 1$  tokens. The  $k^{th}$  valid line after the first line is a space delimited list of node IDs which are the neighbor of node  $k$ . Your parser should be written so as to be robust concerning leading and trailing white space or extra lines at the beginning or end of file, as well as interleaved with valid lines. The **#** character will denote a comment. On any valid line, any characters after a **#** character should be ignored.

You are responsible for ensuring that your program runs correctly when given a valid configuration file. Make no additional assumptions concerning the configuration format. If you have any questions about the configuration format, please ask the TA.

Listing 1: Example configuration file

```
# six global parameters (see above)
5 6 10 100 2000 15

0 dc02 1234 # nodeID hostName listenPort
1 dc03 1233
2 dc04 1233
3 dc05 1232
4 dc06 1233

1 4 # space delimited list of neighbors for node 0
0 2 3 # space delimited list of neighbors for node 1
1 3 # ... node 2
1 2 4 # ... node 3
0 3 # ... node 4
```

## 4 Output Format

If the configuration file is named `<config_name>.txt` and is configured to use  $n$  nodes, then your program should output  $n$  output files, named in according to the following format:

`<config_name>-<node_id>.out`, where `node_id`  $\in \{0, \dots, n - 1\}$ .

The output file for process  $j$  should be named `<config_name>-j.out` and should contain the following: If your program took  $m$  snapshots, then each output file should contain  $m$  lines. The  $i^{th}$  line should contain the vector timestamp of the  $i^{th}$  snapshot *as seen by process  $j$* . Each line of the output file should contain  $n$  space delimited tokens, each of which should be a non-negative integer. In each line, the timestamps must appear in increasing order of process id. That is, the  $k^{th}$  number in the  $i^{th}$  line should be the timestamp value for process  $k$ . An example file is described below.

Listing 2: Example output file for 7 nodes

```
0 4 3 6 0 2 3
3 7 6 7 2 4 4
6 9 11 10 5 7 5
8 12 14 23 8 10 7
```

In this example, the first snapshot has vector clock value  $\langle 0, 4, 3, 6, 0, 2, 3 \rangle$ ; the value of node 0's clock is 0, and the value of node 3's clock is 6.