Team notebook

January 13, 2020

Contents

1 FFT 1

1 FFT

```
// computes f (k) = sum(a[x] exp(2 i kx/N)) for all k.
// Useful for convolution: conv(a, b) = c, where c[x] =
   sum(a[i]b[x i])
// convolution of complex numbers or more than two vectors:
   FFT, multiply
// pointwise, divide by n, reverse(start+1, end), FFT back.
   Rounding is safe if
// Otherwise, use long doubles/NTT/FFTMod.
#define IOS
       ios_base::sync_with_stdio(false); \
       cin.tie(0);
       cout.tie(0);
       cin.exceptions(cin.failbit);
#define trav(a, x) for (auto &a : x)
#define rep(i, a, b) for (int i = a; i < (b); ++i)
#define all(x) begin(x), end(x)
```

```
#define sz(x) (int)(x).size()
typedef long long 11;
typedef pair<int, int> pii;
typedef vector<int> vi;
typedef complex<double> C;
typedef vector<double> vd;
void fft(vector<C> &a)
{
       int n = sz(a), L = 31 - __builtin_clz(n);
       static vector<complex<long double>> R(2, 1);
       static vector<C> rt(2, 1);
       for (static int k = 2; k < n; k *= 2)
       {
              R.resize(n);
              rt.resize(n);
               auto x = polar(1.0L, M_PII / k);
              rep(i, k, 2 * k) rt[i] = R[i] = i & 1 ? R[i /
                  2] * x : R[i / 2];
       }
       vi rev(n);
       rep(i, 0, n) rev[i] = (rev[i / 2] | (i & 1) << L) /
       rep(i, 0, n) if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
       for (int k = 1; k < n; k *= 2)</pre>
```

```
for (int i = 0; i < n; i += 2 * k)
                     rep(j, 0, k)
                      {
                             Cz = rt[j + k] * a[i + j + k];
                             a[i + j + k] = a[i + j] - z;
                             a[i + j] += z;
                      }
}
vd conv(const vd &a, const vd &b)
{
       if (a.empty() || b.empty())
              return {};
       vd res(sz(a) + sz(b) - 1);
       int L = 32 - \_builtin\_clz(sz(res)), n = 1 << L;
       vector<C> in(n), out(n);
       copy(all(a), begin(in));
       rep(i, 0, sz(b)) in[i].imag(b[i]);
       fft(in);
       trav(x, in) x *= x;
       rep(i, 0, n) out[i] = in[-i & (n - 1)] - conj(in[i]);
       fft(out):
       rep(i, 0, sz(res)) res[i] = imag(out[i]) / (4 * n);
       return res:
}
// CP ALGO
using cd = complex<double>;
const double PI = acos(-1);
void fft(vector<cd> &a, bool invert)
{
       int n = a.size();
       for (int i = 1, j = 0; i < n; i++)
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{
               int bit = n >> 1;
               for (; j & bit; bit >>= 1)
                       j ^= bit;
               j ^= bit;
               if (i < j)
                       swap(a[i], a[j]);
       }
       for (int len = 2; len <= n; len <<= 1)</pre>
       {
               double ang = 2 * PI / len * (invert ? -1 : 1);
               cd wlen(cos(ang), sin(ang));
               for (int i = 0; i < n; i += len)</pre>
               {
                       cd w(1);
                       for (int j = 0; j < len / 2; j++)</pre>
                              cd u = a[i + j], v = a[i + j +
                                  len / 2] * w;
                              a[i + j] = u + v;
                              a[i + j + len / 2] = u - v;
                              w \neq wlen:
                      }
               }
       }
       if (invert)
       {
               for (cd &x : a)
                       x /= n;
       }
}
vector<int> multiply(vector<int> &a, vector<int> &b)
{
```

```
vector<cd> fa(a.begin(), a.end()), fb(b.begin(),
           b.end());
       int n = 1:
       while (n < a.size() + b.size())</pre>
               n <<= 1:
       fa.resize(n);
       fb.resize(n);
       fft(fa, false);
       fft(fb, false);
       for (int i = 0; i < n; i++)</pre>
               fa[i] *= fb[i];
       fft(fa, true);
       vector<int> result(n);
       for (int i = 0; i < n; i++)</pre>
               result[i] = round(fa[i].real());
       return result;
}
// Description: Can be used for convolutions modulo specific
   nice primes of
// the form 2^a b + 1, where the convolution result has size
   at most 2 a . Inputs
// must be in [0, mod).
// Time: O (N log N )
const 11 mod = (119 \ll 23) + 1, root = 62; // = 998244353
ll modpow(ll n, ll x)
{
       if (x == 0)
               return 1;
       ll z = modpow(n, x / 2);
       z *= z;
       z \% = mod;
       if (x % 2)
```

```
z *= n:
       z \% = mod;
       return z;
// For p < 230 there i s also e . g . 5 << 25, 7 << 26, 479
   << 21
// and 483 << 21 (same root ) . The l a s t two are > 109.
typedef vector<ll> vl;
void ntt(vl &a, vl &rt, vl &rev, int n)
       rep(i, 0, n) if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
       for (int k = 1; k < n; k *= 2)
               for (int i = 0; i < n; i += 2 * k)
                       rep(j, 0, k)
                               ll z = rt[j + k] * a[i + j + k]
                                  % \mod, \&ai = a[i + j];
                               a[i + j + k] = (z > ai ? ai - z
                                  + mod : ai - z):
                               ai += (ai + z >= mod ? z - mod
                                   : z):
                       }
}
vl conv(vl &a, vl &b)
{
       if (a.empty() || b.empty())
               return {};
       int s = sz(a) + sz(b) - 1, B = 32 -
           __builtin_clz(s), n = 1 \ll B;
       vl L(a), R(b), out(n), rt(n, 1), rev(n);
       L.resize(n), R.resize(n);
       rep(i, 0, n) rev[i] = (rev[i / 2] | (i & 1) << B) /
       11 \text{ curL} = \text{mod} / 2, \text{inv} = \text{modpow}(n, \text{mod} - 2);
       for (int k = 2; k < n; k *= 2)</pre>
       {
```

```
ll z[] = \{1, modpow(root, curL /= 2)\};
              rep(i, k, 2 * k) rt[i] = rt[i / 2] * z[i & 1]
                  % mod;
       }
       ntt(L, rt, rev, n);
       ntt(R, rt, rev, n);
       rep(i, 0, n) out[-i & (n - 1)] = L[i] * R[i] % mod *
           inv % mod;
       ntt(out, rt, rev, n);
       return {out.begin(), out.begin() + s};
}
// Transform to a basis with fast convolutions of the form
   c[z] = a[x] b[y], where is one of AND, OR, XOR. The
   size
// of a must be a power of two.
void FST(vi &a, bool inv)
{
       for (int n = sz(a), step = 1; step < n; step *= 2)
              for (int i = 0; i < n; i += 2 * step)</pre>
                      rep(j, i, i + step)
                      {
                             int &u = a[j], &v = a[j + step];
                             tie(u, v) =
                                 inv ? pii(v - u, u) :
                                    pii(v, u + v); // AND
                             inv ? pii(v, u - v) : pii(u +
                                 v, u); // OR
                             pii(u + v, u - v);
                             // XOR
```

```
}
       }
       if (inv)
              trav(x, a) x /= sz(a); // XOR only
}
vi conv(vi a, vi b)
{
       FST(a, 0);
       FST(b, 0);
       rep(i, 0, sz(a)) a[i] *= b[i];
       FST(a, 1);
       return a;
}
// CRT.h
// Description: Chinese Remainder Theorem.
// crt(a, m, b, n) computes x such that x a (mod m), x b
    (mod n). If
// |a| < m and |b| < n, x will obey 0 x < lcm(m, n).
   Assumes mn < 2^62.
// Time: log(n)
ll crt(ll a, ll m, ll b, ll n)
{
       if (n > m)
               swap(a, b), swap(m, n);
       ll x, y, g = euclid(m, n, x, y);
       assert((a - b) \% g == 0); // e l s e no solution
       x = (b - a) \% n * x \% n / g * m + a;
       return x < 0 ? x + m * n / g : x;
}
```