

## Particle methods

CS B553  
Spring 2013

## Announcements

- A3 posted
  - Due Friday March 8, 11:59PM

### Basic graph cut construction

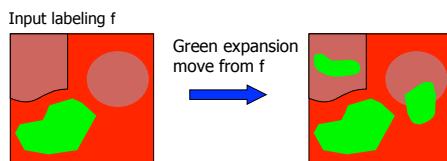
- One non-terminal vertex per pixel
    - Each pixel has edge to s, t, and neighbors
    - Edge p-s has weight  $D_p(0)$ , edge p-t has weight  $D_p(1)$
    - Edge (p,q) has weight  $V_{pq}(0,1)$
  - Run graph cuts to find a min cut
    - Label pixel p 0 if connected to t, and 1 if connected to s
  - Cost of cut is the cost of the entire MRF labeling
    - So min cut means we've found min-cost labeling!
- $$E(x_1, \dots, x_n) = \sum_p D_p(x_p) + \sum_{p,q} V_{pq}(x_p, x_q)$$
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- Adapted from R. Zabih's slide

### Can this be generalized for multi-label problems?

- Not easily.
  - NP-hard for even the Potts model [K/BVZ 01]
- Two main approaches
  1. Exact solution [Ishikawa 03]
    - Large graph, convex  $V$  (arbitrary  $D$ )
  2. Approximate solutions [BVZ 01]
    - Solve a binary labeling problem, repeatedly
    - Expansion move algorithm

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### Expansion move algorithm



- Make green expansion move that most decreases cost
  - Then make the best blue expansion move, etc
  - Done when no  $\alpha$ -expansion move decreases the energy, for any label  $\alpha$
  - See [BVZ 01] for details

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Adapted from R. Zabih's slide

### Binary sub-problem



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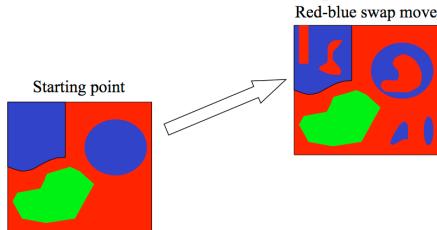
Adapted from R. Zabih's slide

## The expansion move algorithm

1. Start with an arbitrary labeling
2. Cycle through every label  $A$  in some order
  - 2.1 Find the lowest cost labeling that involves an  $A$ -expansion move – this is a binary subproblem!
  - 2.2 Make the move if its cost is lower than current labeling
3. If cost did not decrease in the cycle, we're done  
Otherwise, go to step 2

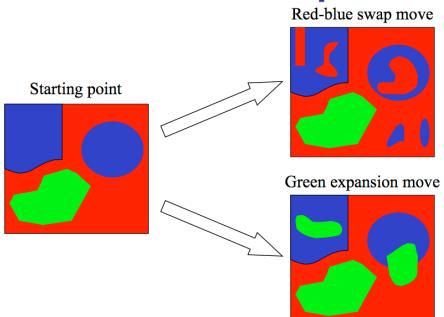
Adapted from R. Zabih's slide

## Move examples



Adapted from R. Zabih's slide

## Move examples



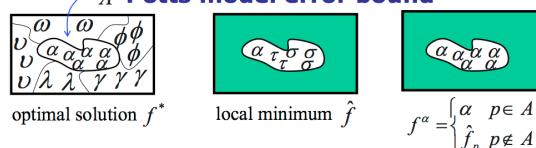
Adapted from R. Zabih's slide

## The swap move algorithm

1. Start with an arbitrary labeling
2. Cycle through every label pair  $(A, B)$  in some order
  - 2.1 Find the lowest cost labeling within a single  $AB$ -swap
  - 2.2 Go there if its cost is lower than the current labeling
3. If cost did not decrease in the cycle, we're done  
Otherwise, go to step 2

Adapted from R. Zabih's slide

## Potts model error bound



Summing up over all labels:

$$E(\hat{f}) \leq E(f^*) + E_{\partial}(f^*) \leq 2E(f^*)$$

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## Multi-label graph cuts

- The approximate algorithm works for:
  - $D$  of any form
  - $V$  must satisfy a (generalized) submodularity constraint:

$$\begin{aligned} V(\alpha, \alpha) + V(f(p), f(q)) &\leq \\ V(f(p), \alpha) + V(\alpha, f(q)) &\end{aligned}$$

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Adapted from R. Zabih's slide



