DBMS - no concrawench fill system - redundancy MBMS as soln. - Inconsistency - harder access data abstraction: hiding irrelevant details from users. 1 topy physical level (acroll store in disk). 10 logical ierel (define schema) @ view derel. physical -> logical - view 2 schema - lagical structure of db (blueprint). instance - actual data stored in db (CRUD on instance):

data ind models (tools)

- relational model
- ER model
- object based model
- network and hierarchical model.

DMI - data manipulation language.

- access and manipulate

- CRUD opelouhions.

disk sto

(1)

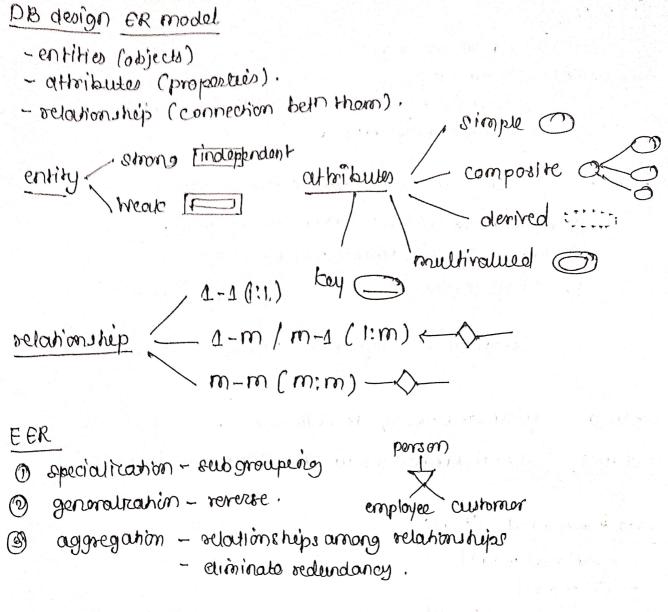
DDL - data defination language

manage schema and structure.

create, other, drop and bruncate.

Architeculure

- 1. tier architecutice (db and uz insame system)
- 2. Hir architecture client -> server + db
- 3. tein architecture client apr (ur) -> serrer + db.



(X) Relational DB Design

- represent data un tables (relations).
- pitfalls redundant data
 - Insertion anomalies
 - delation anomalies (unintended loss of data)
 - updation anomalies (single update req multi vou update)

Normalisation: decompase large selations into small, to eaduce eadendary

1NF - atomic value (anique)

- no multi valued atteibutes.

rame	branch	rollno	language	name branch rolling lang.
charen	AIM L	40	python, C	(h) 1 m n m 1 90
Bharesh	Electrical	4 Q	Rust, Ctt	

2NF - should 1NF - no partial dependencies (PK → ® duplicate value).

JNF =		name	(plc)	(milno).	(FK) rallno language	
	\Rightarrow	cheran	40	AIML	MO	python
	1	* bharest) pla	e/echrical	40	RUST
					40	C++·

BNF - should 2NF

- no transitive dependency

- non prime atteibute should depend on primary key (PK).

studentid	name d	OI toler	dept nam	le	M-CHEST MERCHANISM	itive dependen
101	cherom	1	AIML			nol·
102	bhareth	2	66	shidi	entid <	dept Id & dept name
	shidentid	name	dept Id		dept id	dept nome
	101	cheroun			1	AIML
	102	bhorest	7 2	black by Mills	2	66

BCNF - should GNF.

- dependency A > B then A should super key

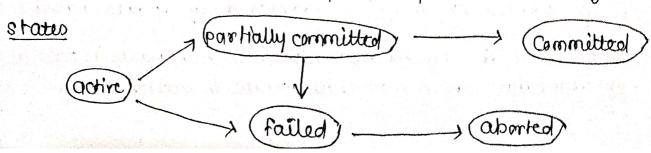
* Transactions
- transaction is unit of work performed in db.

ACED - atomicity Call reflect or none)

- consistency (after execution consistency in db)

- Prolation C seperate execution)

- durability (beickup after rystem failure, changes soved).



0

schedules	
~ CI Carrette	

- sequence of instructions that specify chronlogical order of concurrent transactions are executed.

1) serial scheduling - one transaction completed when 71 another is executed.

concurrent scheduling - in betto execution.

serializibility - preserve db consistency.

final result align with requential execution.

1 conflict serializability - transform in serial schedule by swap

@ view reprodicibility - view equivalent to resid schedule.

Read - compatible instruction Weite - non-compatible instruction.

Lock based Protocal

- Lock: mechanism to control concurrent access to data item

100k-x

@ shared - read only - lock-s

S D F E(K) F F

deadlock -> rollback and lock released.

T1-5 then T2 gets 8

soln: Two phase locking protocol

- 1) growing phase obtain locks, not release.
- 1 Shrinking phase release locks, not obtain

Recovery (Log based recovery) - log records of update activities in db.

A derected db - record all modification but update in partial commit

13) immediate ab - immediate update as withen / change