

华理小男孩模板

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Chapter 1

数论

1.1 指数降幂公式

$$A^x \equiv A^{x \bmod \phi(p) + \phi(p)} \pmod{p} (x \geq \phi(p))$$

1.2 威尔逊定理

$$(p-1)! \equiv -1 \pmod{p}$$

1.3 费马小定理

$$a^p \equiv a \pmod{p}$$

1.4 欧拉定理

$$a^{\phi(n)} \equiv 1 \pmod{n}$$

1.5 质数表

```
1  const int N = 1000000 + 9;  
2  bool p[N];  
3  int a[N];  
4  int main()  
5  {  
6      int n;  
7      cin >> n;  
8      int cnt = 0;  
9      for (int i = 0; i <= n; i++) p[i] = true;  
10     for (int i = 2; i <= n; i++) {
```

```

11     if (p[i]) a[cnt++] = i;
12     for (int j = 0; j < cnt; j++) {
13         if (i * a[j] > n) break;
14         p[i * a[j]] = false;
15         if (i % a[j] == 0) break;
16     }
17 }
18 cout << cnt << endl;
19 }

```

1.6 素数函数

```

1 #include <bits/stdc++.h>
2 typedef long long ll;
3 using namespace std;
4 ll f[340000], g[340000], n;
5 void init(){
6     ll i, j, m;
7     for(m = 1; m * m <= n; ++m) f[m] = n / m - 1;
8     for(i = 1; i <= m; ++i) g[i] = i - 1;
9     for(i = 2; i <= m; ++i){
10         if(g[i] == g[i - 1]) continue;
11         for(j = 1; j <= min(m - 1, n / i / i); ++j){
12             if(i * j < m) f[j] -= f[i * j] - g[i - 1];
13             else f[j] -= g[n / i / j] - g[i - 1];
14         }
15         for(j = m; j >= i * i; --j) g[j] -= g[j / i] - g[i - 1];
16     }
17 }
18 int main()
19 {
20     while (cin >> n) {
21         init();
22         cout << f[1] << endl;
23     }
24 }

```

1.7 欧拉函数

1.7.1 递推求

```

1 // 傻逼写的
2 int phi[MAXN];

```

```

3 void init() {
4     memset(phi, 0, sizeof(phi));
5     phi[1] = 1;
6     for(int i = 2; i < MAXN; i++) if(!phi[i])
7         for(int j = i; j < MAXN; j += i) {
8             if(!phi[j]) phi[j] = j;
9             phi[j] = phi[j] / i * (i - 1);
10        }
11    }
12    // 聪明人写的
13    void init()
14    {
15        for (int i = 2; i < maxb; i++)
16            if (f[i] == 0)
17                for (int j = i; j < maxb; j += i)
18                {
19                    if (f[j] == 0) f[j] = j;
20                    f[j] = (f[j] / i) * (i - 1);
21                }
22        for (int i = 1; i < maxb; i++)
23            f[i] += f[i - 1];
24    }

```

1.7.2 单个求

```

1 ll phi(ll n){
2     ll ans = n, a = n;
3     for(ll i = 2; i * i <= a; i++){
4         if(a % i == 0){
5             ans = ans / i * (i - 1);
6             while(a % i == 0) a /= i;
7         }
8     }
9     if(a > 1) ans = ans / a * (a - 1);
10    return ans;
11 }

```

1.8 莫比乌斯函数

```

1 int mu[MAXN], prm[MAXN], vis[MAXN];
2 void getmu(int n) {
3     int sz = 0;
4     mu[1] = 1;

```

```

5     for(int i = 2; i <= n; i++) {
6         if(!vis[i]) prm[++sz] = i, mu[i] = -1;
7         for(int j = 1; j <= sz && prm[j] * i <= n; j++) {
8             vis[i * prm[j]] = 1;
9             if(i % prm[j] == 0) {mu[i * prm[j]] = 0;
10                break;}
11             mu[i * prm[j]] = -mu[i];
12         }
13     }

```

1.9 逆元

1.9.1 递推求

```

1  ll mul_mod(ll x, ll y, ll mod = MOD)
2  {
3      return x * y % MOD;
4  }
5  m[1] = 1;
6  inv[1] = 1;
7  for (n = 2; ;n++) {
8      s[n] = s[n - 1] + n;
9      m[n] = mul_mod(m[n - 1], n);
10     inv[n] = (MOD - MOD / n) * inv[MOD % n] % MOD;
11     if (s[n] > MAXX) break;
12 }

```

1.9.2 单个求

用费马小定理

Chapter 2

概率论

2.1 超几何分布

超几何分布是统计学上一种离散概率分布。它描述了由有限个物件中抽出 n 个物件，成功抽出指定种类的物件的个数（不归还）

例如在有 N 个样本，其中 m 个是不及格的。超几何分布描述了在该 N 个样本中抽出 n 个，其中 k 个是不及格的概率：

$$f(k; n, m, N) = \frac{\binom{m}{k} \binom{N-m}{n-k}}{\binom{N}{n}}$$

Chapter 3

数学

3.1 矩阵

3.1.1 矩阵类

```
1 #include <bits/stdc++.h>
2 using namespace std;
3 const int N = 1000, MOD = 1E9 + 7;
4 struct Mat
5 {
6     int a[N][N];
7     int n, m;
8     Mat(int n, int m) : n(n), m(m) {memset(a, 0, sizeof(a));}
9     void eye()
10    {
11        memset(a, 0, sizeof(a));
12        for (int i = 0; i < n; i++) a[i][i] = 1;
13    }
14    void print()
15    {
16        for (int i = 0; i < n; i++) {
17            cout << endl;
18            for (int j = 0; j < m; j++)
19                cout << ' ' << a[i][j];
20        }
21    }
22 };
23 Mat mul(Mat A, Mat B)
24 {
25     Mat t(A.n, B.m);
```

```

26     for (int i = 0; i < A.n; i++)
27         for (int j = 0; j < B.m; j++)
28             for (int k = 0; k < A.m; k++)
29                 t.a[i][j] = (t.a[i][j] + A.a[i][k] * B.a[
                    k][j] % MOD) % MOD;
30     return t;
31 }
32 Mat pow(Mat A, int n)
33 {
34     Mat t(A.n, A.m);
35     t.eye();
36     while (n > 0) {
37         if (n % 2) t = mul(t, A);
38         A = mul(A, A);
39         n /= 2;
40     }
41     return t;
42 }
43 int det(Mat A)
44 {
45     int cnt = 0, ans = 1, n = A.n;
46     for(int i = 0; i < n; i++) {
47         for(int j = i+1; j < n; j++) {
48             int x = i, y = j;
49             while(A.a[y][i]) {
50                 int t = A.a[x][i] / A.a[y][i];
51                 for(int k = 0; k < n; k++) {
52                     A.a[x][k] = A.a[x][k] - A.a[y][k]*t;
53                 }
54                 swap(x, y);
55             }
56             if(x != i) {
57                 for(int k = 0; k < n; k++) {
58                     swap(A.a[x][k], A.a[y][k]);
59                 }
60                 cnt ^= 1;
61             }
62         }
63         if(A.a[i][i] == 0) return 0;
64         else ans *= A.a[i][i];
65     }
66     if(cnt) ans *= -1;
67     return ans;
68 }

```

3.1.2 高斯消元

```
1 double a[MAXN][MAXN];
2 double ans[MAXN];
3 bool f[MAXN]; // 自由变量
4 int sgn(double x) {return (x > eps) - (x < -eps);}
5 // x 0~equ - 1, y 0~var
6 int gauss(int equ, int var) {
7     int k = 0, col = 0;
8     memset(f, true, sizeof(f));
9     for(k = 0; k < equ && col < var; k++, col++) {
10         int r = k;
11         for(int i = k + 1; i < equ; i++)
12             if(fabs(a[i][col]) > fabs(a[r][col])) r = i;
13         if(r != k) for(int j = k; j <= var; j++) swap(a[r
14             ][j], a[k][j]);
15         if(a[k][col] == 0) {k--; continue;}
16         for(int i = k + 1; i < equ; i++) if(a[i][col]) {
17             for(int j = var; j >= col; j--) a[i][j] -= a[
18                 i][col] / a[k][col] * a[k][j];
19         }
20         for(int i = k; i < equ; i++) if(sgn(a[i][col]) != 0)
21             return 0;
22         if(k < var) {
23             for(int i = k - 1; i >= 0; i--) {
24                 int cnt = 0, p;
25                 for(int j = 0; j < var; j++)
26                     if(sgn(a[i][j]) && f[j])
27                         cnt++, p = j;
28                 if(cnt > 1) continue;
29                 double t = a[i][var];
30                 for(int j = 0; j < var; j++)
31                     if(sgn(a[i][j]) && j != p)
32                         t -= a[i][j] * ans[j];
33                 ans[p] = t / a[i][p];
34                 f[p] = 0;
35             }
36         }
37         for(int i = var - 1; i >= 0; i--) {
38             double t = a[i][var];
39             for(int j = i + 1; j < var; j++)
40                 if(sgn(a[i][j]))
41                     t -= a[i][j] * ans[j];
42             ans[i] = t / a[i][i];
43         }
44     }
```

```

42     return 1;
43 }

```

3.2 整除与剩余

3.2.1 扩展欧几里得逆元

```

1  扩展欧几里得
2  //ax + by = d, d = gcd(a, b)
3  void gcd(ll a, ll b, ll &d, ll &x, ll &y){
4      if(!b) {d = a; x = 1; y = 0; return;}
5      else{
6          gcd(b, a % b, d, y, x);
7          y -= x * (a / b);
8          return;
9      }
10 }
11
12 逆元
13 //a在模n下的逆元, a和n互素才有逆元
14 ll inv(ll a, ll n) {
15     ll d, x, y;
16     gcd(a, n, d, x, y);
17     return d == 1 ? (x + n) % n : -1;
18 }

```

3.2.2 中国剩余定理

```

1  // n个方程 x ≡ a[i](mod m[i]) i = 0..n-1
2  ll china(int n, ll *a, ll *m)
3  {
4      ll M = m[0], R = a[0];
5      for (int i = 1; i < n; i++) {
6          ll d = __gcd(M, m[i]);
7          ll c = a[i] - R;
8          if (c % d) return -1;
9          ll k1, k2;
10         extgcd(M, m[i], d, k1, k2);
11         k1 = (c / d * k1) % (m[i] / d);
12         R = R + k1 * M;
13         M = M / d * m[i];
14         R %= M;
15     }
16     if (R < 0) R += M;

```

```

17     return R;
18 }

```

3.3 数值计算

3.4 其他

3.4.1 lucas 定理

```

1 ll qPow (ll a, ll k) {
2     ll ans = 1;
3     while (k) {
4         if (k&1)
5             ans = (ans * a) % p;
6         a = (a * a) % p;
7         k /= 2;
8     }
9     return ans;
10 }
11
12 ll C (ll a, ll b, ll p) {
13     if (a < b)
14         return 0;
15     if (b > a - b)
16         b = a - b;
17
18     ll up = 1, down = 1;
19     for (ll i = 0; i < b; i++) {
20         up = up * (a-i) % p;
21         down = down * (i+1) % p;
22     }
23     return up * qPow(down, p-2) % p; // 逆元
24 }
25 ll lucas (ll a, ll b, ll p) {
26     if (b == 0)
27         return 1;
28     return C(a%p, b%p, p) * lucas(a/p, b/p, p) % p;
29 }

```

3.4.2 递推求组合数

```

1 void calc()
2 {

```

```

3     for (int i = 0; i < N; i++) {
4         c[i][0] = c[i][i] = 1;
5         for (int j = 1; j < i; j++) {
6             c[i][j] = c[i-1][j] + c[i-1][j-1];
7         }
8     }
9 }

```

3.4.3 单个求组合数

```

1 ll C(ll n, ll m)
2 {
3     ll c = 1;
4     for (int i = 1; i <= m; i++) {
5         c *= (n - m + i);
6         if (c % i == 0) c /= i;
7     }
8     return c;
9 }

```

3.4.4 威佐夫博弈

```

1 double gold = (1 + sqrt(5)) / 2;
2 if (m == floor(((n - m) * gold))) cout << 'G' << endl;
   else cout << 'B' << endl;

```

3.4.5 FWT

```

1 //位运算多项式卷积
2 void FWT(int *a, int n) {
3     for(int d = 1; d < n; d<<=1)
4         for(int m = d<<1, i = 0; i < n; i += m)
5             for(int j = 0; j < d; j++) {
6                 int x = a[i + j], y = a[i + j + d];
7                 a[i + j] = (x + y) % MOD, a[i + j + d] =
                    (x - y + MOD) % MOD;
8                 //xor:a[i+j]=x+y,a[i+j+d]=(x-y+mod)%mod;
9                 //and:a[i+j]=x+y;
10                //or:a[i+j+d]=x+y;
11            }
12 }
13 void UFWT(int *a, int n) {
14     for(int d = 1; d < n; d <<= 1)

```

```

15     for(int m = d<<1, i = 0; i < n; i += m)
16         for(int j = 0; j < d; j++) {
17             int x = a[i + j], y = a[i + j + d];
18             a[i + j] = (ll)(x + y) * inv2 % MOD, a[i
                + j + d] = ((ll)(x - y) * inv2 % MOD +
                    MOD) % MOD;
19             //xor:a[i+j]=(x+y)/2,a[i+j+d]=(x-y)/2;
20             //and:a[i+j]=x-y;
21             //or:a[i+j+d]=y-x;
22         }
23     }
24 void conv(int *a, int *b, int n) {
25     FWT(a, n);
26     FWT(b, n);
27     for(int i = 0; i < n; i++) a[i] = (ll)a[i] * b[i] %
        MOD;
28     UFWT(a, n);
29 }

```

3.4.6 FFT

```

1  const double pi = atan(1.0) * 4;
2
3  struct complex {
4      double a, b;
5      complex(double aa = 0.0, double bb = 0.0) { a =
        aa; b = bb; }
6      complex operator +(const complex &e) { return
        complex(a + e.a, b + e.b); }
7      complex operator -(const complex &e) { return
        complex(a - e.a, b - e.b); }
8      complex operator *(const complex &e) { return
        complex(a * e.a - b * e.b, a * e.b + b * e.a);
        }
9  };
10
11 void change(complex * y, long long len) {
12     long long i, j, k;
13     for (i = 1, j = len / 2; i < len - 1; i++) {
14         if (i < j) swap(y[i], y[j]);
15         k = len / 2;
16         while (j >= k) {
17             j -= k;
18             k /= 2;
19         }

```



```

20         if (j < k) j += k;
21     }
22 }
23
24 void fft(complex *y, long long len, long long on) {
25     change(y, len);
26     for (int h = 2; h <= len; h <= 1) {
27         complex wn(cos(-on * 2 * pi / h), sin(-on
28             * 2 * pi / h));
29         for (int j = 0; j < len; j += h) {
30             complex w(1, 0);
31             for (int k = j; k < j + h / 2; k
32                 ++ ) {
33                 complex u = y[k];
34                 complex t = w * y[k + h /
35                     2];
36                 y[k] = u + t;
37                 y[k + h / 2] = u - t;
38                 w = w * wn;
39             }
40         }
41     }
42     if (on == -1)
43         for (int i = 0; i < len; i++)
44             y[i].a /= len;
45 }

```

3.4.7 hell 方程

$$x^2 - ny^2 = 1 \quad x[i+1] = x[1] * x[i] + n * y[1] * y[i]; \quad y[i+1] = x[1] * y[i] + y[1] * x[i]$$

Chapter 4

图论

4.1 最大团

```
1  /*
2  最大团搜索算法
3   $S_i = \{v_i, v_{i+1}, \dots, v_n\}$ 
4   $mc[i] = MC(S_i) =$  点集合  $S_i$  的最大团
5   $MC(V) = mc[1]$ 
6   $mc[i] = mc[i+1]$  //  $mc[i] = mc[i+1] + 1$  因为最多增加一个点
7  后一种情况发生是在  $S_i$  中找到一个包含  $v_i$  的团
8  所以只要搜 是否在  $S_i$  中有一个包含  $v_i$  且比当前最大团还大的团
9  */
10 #include <bits/stdc++.h>
11 #define rep(i, l, r) for (int (i) = (l); (i) <= (r); (i)++)
12 #define log(x) cout << #x << " = " << x << endl
13 #define mem(x, y) memset((x), (y), sizeof(x))
14 using namespace std;
15 int n, ans;
16 bool found;
17 const int N = 50 + 9;
18 int g[N][N];
19 int len[N], mc[N], li[N][N];
20 void dfs(int sz)
21 {
22     if (len[sz] == 0) {
23         if (sz > ans) {
24             ans = sz;
25             found = true;
26         }
27     }
```

```

28     for (int k = 0; k < len[sz] && !found; k++) {
29         if (sz + len[sz] - k <= ans) break;
30         int i = li[sz][k];
31         if (sz + mc[i] <= ans) break;
32         len[sz+1] = 0;
33         for (int j = k + 1; j < len[sz]; j++) {
34             if (g[i][li[sz][j]])
35                 li[sz+1][len[sz+1]++] = li[sz][j];
36         }
37         dfs(sz+1);
38     }
39 }
40 int max_cluster()
41 {
42     mc[n] = ans = 1;
43     for (int i = n - 1; i; i--) {
44         found = false;
45         len[1] = 0;
46         rep(j, i+1, n) {
47             if (g[i][j]) li[1][len[1]++] = j;
48         }
49         dfs(1);
50         mc[i] = ans;
51     }
52     return ans;
53 }
54 int solve()
55 {
56     rep(i, 1, n) rep(j, 1, n) scanf("%d", &g[i][j]);
57     return max_cluster();
58 }
59 int main()
60 {
61     for (;;) {
62         scanf("%d", &n);
63         if (n == 0) break;
64         printf("%d\n", solve());
65     }
66 }

```

4.2 图的遍历和连通性

4.2.1 割点和桥

```

1 int pre[MAXN], iscut[MAXN], dfs_clock = 0;

```

```

2  int dfs(int u, int fa) {
3  int lowu = pre[u] = ++dfs_clock;
4  int child = 0;
5  for(int i = 0; i < g[u].size(); i++) {
6  int v = g[u][i];
7  if(!pre[v]) {
8  child++;
9      int lowv = dfs(v, u);
10     lowu = min(lowu, lowv);
11     if(lowv >= pre[u]) iscut[u] = true;
12     //if lowv > pre[u] (u, v) is bridge
13
14     }
15     else if(pre[v] < pre[u] && v != fa) lowu = min(
        lowu, pre[v]);
16 }
17 if(fa < 0 && child == 1) iscut[u] = false;
18 return lowu;
19 }

```

4.2.2 双连通分量

```

1  int dfs(int u, int fa) {
2  int lowu = pre[u] = ++dfs_clock;
3  int child = 0;
4  for(int i = 0; i < g[u].size(); i++) {
5  int v = g[u][i];
6  if(!pre[v]) {
7  s.push(node(u, v));
8  child++;
9  int lowv = dfs(v, u);
10 lowu = min(lowu, lowv);
11 if(lowv >= pre[u]) {
12     iscut[u] = true;
13     bcc[++bcc_cnt].clear();
14     for(;;) {
15         node temp = s.top(); s.pop();
16         //注意割顶可能包含在多个bcc中,
            bccno不是唯一的标准
17         if(bccno[temp.u] != bcc_cnt)
18         {bcc[bcc_cnt].push_back(temp.u); bccno[
            temp.u] = bcc_cnt;}
19         if(bccno[temp.v] != bcc_cnt)
20         {bcc[bcc_cnt].push_back(temp.v); bccno[
            temp.v] = bcc_cnt;}

```

```

21         if(temp.u == u && temp.v == v) break;
22     }
23 }
24 }
25     else if(pre[v] < pre[u] && v != fa) {
26         s.push(node(u, v));
27         lowu = min(lowu, pre[v]);
28     }
29 }
30 if(fa < 0 && child == 1) iscut[u] = false;
31 return lowu;
32 }
33 void find_bcc(int n) {
34     memset(pre, 0, sizeof(pre));
35     memset(iscut, false, sizeof(iscut));
36     memset(bccno, 0, sizeof(bccno));
37     dfs_clock = bcc_cnt = 0;
38     for(int i = 1; i <= n; i++) {
39         if(!pre[i]) dfs(i, -1);
40     }
41 }

```

4.2.3 强连通分量

```

1 int dfs(int u) {
2     pre[u] = low[u] = ++dfs_clock;
3     s.push(u);
4     for(int i = 0; i < g[u].size(); i++) {
5         int v = g[u][i];
6         if(!pre[v]) {
7             dfs(v);
8             low[u] = min(low[u], low[v]);
9         }
10        else if(!sccno[v]) {
11            low[u] = min(low[u], pre[v]);
12        }
13    }
14    //如果low[u] == pre[u], 那么它就是这个scc的第一个点
15    if(low[u] == pre[u]) {
16        scc_cnt++;
17        for(;;) {
18            int temp = s.top(); s.pop();
19            sccno[temp] = scc_cnt;
20            if(temp == u) break;
21        }

```

```

22     }
23 }
24 void find_scc(int n) {
25     memset(pre, 0, sizeof(pre));
26     memset(low, 0, sizeof(low));
27     memset(sccno, 0, sizeof(sccno));
28     dfs_clock = scc_cnt = 0;
29     for(int i = 1; i <= n; i++) {
30         if(!pre[i]) dfs(i);
31     }
32 }

```

4.2.4 拓扑排序

BFS

判断是否成环

拓扑排序形成的 ans 的 sz != n 则成环

4.2.5 2SAT

```

1 //2-sat中不能走的决策往另一个决策连一条边
2 int dfs(int u) {
3     if(vis[u^1]) return false;
4     if(vis[u]) return true;
5     s[cnt++] = u;
6     vis[u] = 1;
7     for(int i = 0; i < g[u].size(); i++) if(!dfs(g[u][i]))
8         return false;
9     return true;
10 }
11 bool flag = true;
12 //一定记得+=2
13 for(int i = 2; i <= 2 * n; i+= 2) {
14     if(!vis[i] && !vis[i + 1]) {
15         cnt = 0;
16         if(!dfs(i)) {
17             while(cnt) vis[s[--cnt]] = 0;
18             if(!dfs(i + 1)) {flag = false; break;}
19         }
20     }
21 }

```

4.3 路径

4.3.1 非递归欧拉回路

```
1 void euler(int u) {
2     stack<int> st;
3     st.push(u);
4     nxt[st.size()] = -1;
5     while(!st.empty()) {
6         int a = st.top();
7         int i;
8         for(i = last[a]; i < 26; i++) if(!vis[a][i]) {
9             vis[a][i] = 1;
10            st.push(g[a][i]);
11            nxt[st.size()] = i;
12            last[a] = i + 1;
13            break;
14        }
15        if(i == 26) {
16            if(nxt[st.size()] != -1) ans.push_back((char)
17                (nxt[st.size()] + 'a'));
18            st.pop();
19        }
20    }
```

4.4 匹配

4.4.1 二分图最大匹配

最小点覆盖的点数 = 二分图最大匹配

最大独立集的点数 = 总点数 - 二分图最大匹配

```
1 int uN, vN;
2 vector<int> g[MAXN];
3 int linker[MAXN];
4 bool used[MAXN];
5 bool dfs(int u) {
6     for(int i = 0; i < g[u].size(); i++) {
7         int v = g[u][i];
8         if(!used[v]) {
9             used[v] = true;
10            if(linker[v] == -1 || dfs(linker[v])) {
11                linker[v] = u;
12                return true;
13            }
14        }
15    }
```

```

13         }
14     }
15 }
16 return false;
17 }
18 int hungary() {
19     int res=0;
20     memset(linker, -1, sizeof(linker));
21     for(int u=0; u < uN; u++)
22     {
23         memset(used, 0, sizeof(used));
24         if(dfs(u)) res++;
25     }
26     return res;
27 }

```

4.4.2 二分图最优匹配

```

1 struct KM // 二分图最优匹配
2 {
3     int n; //总点数
4     vector<int> g[N];
5     int g2[N][N];
6     void init(int nn)
7     {
8         n = nn;
9         mem(g, 0), mem(g2, 0);
10    }
11    void add_edge(int u, int v, int w)
12    {
13        g[u].push_back(v);
14        g2[u][v] = w;
15    }
16    int lx[N], ly[N], match[N], lcheck[N], rcheck[N];
17    const int INF = INT_MAX;
18    bool dfs(int u)
19    {
20        lcheck[u] = true;
21        for (int v : g[u]) {
22            if (lx[u] + ly[v] == g2[u][v] && !rcheck[v])
23            {
24                rcheck[v] = true;
25                if (match[v] == -1 || dfs(match[v])) {
26                    match[v] = u;

```



```

27         }
28     }
29 }
30     return false;
31 }
32 void update()
33 {
34     int a = INF;
35     rep(u, 1, n) {
36         if (lcheck[u]) {
37             for (int v : g[u]) {
38                 if (!rcheck[v]) {
39                     a = min(a, lx[u] + ly[v] - g2[u][
                        v]);
40             }
41         }
42     }
43 }
44     rep(i, 1, n) {
45         if (lcheck[i]) lx[i] -= a;
46         if (rcheck[i]) ly[i] += a;
47     }
48 }
49 int calc()
50 {
51     rep(i, 1, n) {
52         lx[i] = *max_element(g2[i]+1, g2[i]+n+1);
53         ly[i] = 0;
54         match[i] = -1;
55     }
56     rep(i, 1, n) {
57         for (;;) {
58             mem(lcheck, 0);
59             mem(rcheck, 0);
60             if (dfs(i)) break; else update();
61         }
62     }
63     int ans = 0;
64     rep(i, 1, n) if (~match[i]) ans += lx[match[i]] +
        ly[i];
65     return ans;
66 }
67 };
68 int solve(int n, KM &solver)
69 {
70     solver.init(n);

```

```

71     rep(i, 1, n) rep(j, 1, n) {
72         int x;
73         scanf("%d", &x);
74         solver.add_edge(i, j, x);
75     }
76     return solver.calc();
77 }
78 int main()
79 {
80     KM solver;
81     int n;
82     while (~scanf("%d", &n)) printf("%d\n", solve(n,
83         solver));
84 }

```

4.5 树

4.5.1 prim

```

1  int cnt = 1, ans = 0;
2  priority_queue<edge> pq;
3  for(int i = 0; i < g[1].size(); i++) pq.push(g[1][i]);
4  while(!pq.empty()) {
5      edge t = pq.top(); pq.pop();
6      int v = t.v;
7      if(p[v] != 1) {
8          cnt++; ans += t.w;
9          p[v] = 1;
10     for(int i = 0; i < g[v].size(); i++) if(p[g[v][i].v] !=
11         1) pq.push(g[v][i]);
12     }
13     if(cnt == n) break;
14 }

```

4.5.2 曼哈顿最小距离生成树

```

1  struct Edge {
2      int u, v, w;
3      bool operator < (const Edge& rhs) const {
4          return w < rhs.w;
5      }
6  }edges[8 * MAXN];
7  struct node {
8      int x, y, id;

```

```

9      bool operator < (const node& rhs) const {
10          return x == rhs.x ? y > rhs.y : x > rhs.x;
11      }
12  }nd[MAXN];
13  int ecnt, n, k, sz;
14  int sq[MAXN];
15  int minv[MAXN], pos[MAXN];
16  int p[MAXN];
17  int lowbit(int x) {return x & -x;}
18  void update(int x, int val, int id) {
19      while(x) {
20          if(val < minv[x]) {
21              minv[x] = val; pos[x] = id;
22          }
23          x -= lowbit(x);
24      }
25  }
26  int query(int x) {
27      int ans = 1 << 30, ret = 0;
28      while(x <= sz) {
29          if(minv[x] < ans) {
30              ans = minv[x]; ret = pos[x];
31          }
32          x += lowbit(x);
33      }
34      //cout << ans << endl;
35      return ret;
36  }
37  int fp(int x) {return p[x] == x ? x : p[x] = fp(p[x]);}
38  void addEdge(int u, int v, int w) {
39      edges[ecnt++] = (Edge) {u, v, w};
40  }
41  void solve() {
42      for(int dir = 0; dir < 4; dir++) {
43          if(dir == 1 || dir == 3)
44              for(int i = 1; i <= n; i++) swap(nd[i].x, nd[
45                  i].y);
46          else if(dir == 2)
47              for(int i = 1; i <= n; i++) nd[i].x = -nd[i].
48                  x;
49          sort(nd + 1, nd + n + 1);
50          for(int i = 1; i <= n; i++) sq[i] = nd[i].y - nd[
51              i].x;
52          sort(sq + 1, sq + n + 1);
53          sz = unique(sq + 1, sq + n + 1) - (sq + 1);

```

```

51     for(int i = 1; i <= sz; i++) {minv[i] = (1 << 30)
        ; pos[i] = 0;}
52     for(int i = 1; i <= n; i++) {
53         int p = lower_bound(sq + 1, sq + sz + 1, nd[i]
            ].y - nd[i].x) - sq;
54         int v = query(p);
55         //cout << p << ' ' << v << endl;
56         if(v) addEdge(nd[i].id, nd[v].id, abs(nd[i].x
            - nd[v].x) + abs(nd[i].y - nd[v].y));
57         update(p, nd[i].x + nd[i].y, i);
58     }
59 }
60 sort(edges, edges + ecnt);
61 int cnt = 0;
62 for(int i = 1; i <= n; i++) p[i] = i;
63 for(int i = 0; i < ecnt; i++) {
64     Edge e = edges[i];
65     int x = fp(e.u), y = fp(e.v);
66     if(x != y) {
67         p[x] = y;
68         if(++cnt == n - k) {cout << e.w << endl;
            break;}
69         //cout << cnt << endl;
70     }
71 }
72 }

```

4.6 网络流

4.6.1 最大流 Dinic

```

1  int bfs(int s, int t) {
2      memset(vis, false, sizeof(vis));
3      queue<int> q;
4      q.push(s);
5      vis[s] = true; d[s] = 0;
6      while(!q.empty()) {
7          int u = q.front(); q.pop();
8          for(int i = 0; i < g[u].size(); i++) {
9              Edge e = edges[g[u][i]];
10             if(vis[e.v]) continue; //不能再走访问过的点
11             if(e.cap > e.flow) {
12                 d[e.v] = d[u] + 1;
13                 vis[e.v] = true;

```

```

14         q.push(e.v);
15     }
16 }
17 }
18 return vis[t];
19 }
20 int dfs(int u, int a, int t) {
21     if(a == 0 || u == t) return a;
22     int flow = 0, f;
23     for(int &i = cur[u]; i < g[u].size(); i++) {
24         Edge &e = edges[g[u][i]];
25         if(d[e.v] == d[u] + 1 && (f = dfs(e.v, min(a, e.
26             cap - e.flow), t)) > 0) { //如果从v走还可以增广
27             e.flow += f;
28             edges[g[u][i]^1].flow -= f;
29             flow += f;
30             a -= f;
31             if(a == 0) break;
32         }
33     }
34     return flow;
35 }
36 int dinic(int s, int t) {
37     int flow = 0;
38     while(bfs(s, t)) {
39         memset(cur, 0, sizeof(cur));
40         flow += dfs(s, INF, t);
41     }
42     return flow;
43 }

```

4.6.2 最大流 ISAP

```

1 struct ISAP {
2     struct Edge {
3         int from, to, cap, flow;
4         bool operator < (const Edge& b) {
5             return from < b.from || (from == b.from && to < b
6                 .to);
7         }
8     };
9
10    int n, m, s, t;
11    vector<Edge> edges;

```

```

12     vector<int> g[MAXN];    // 邻接表, g[i][j]表示结点i
                               的第j条边在e数组中的序号
13     bool vis[MAXN];        // BFS使用
14     int d[MAXN];           // 从起点到i的距离
15     int cur[MAXN];         // 当前弧指针
16     int p[MAXN];           // 可增广路上的上一条弧
17     int num[MAXN];         // 距离标号计数
18
19     void add_edge(int from, int to, int cap) {
20         edges.push_back((Edge){from, to, cap, 0});
21         edges.push_back((Edge){to, from, 0, 0});
22         m = edges.size();
23         g[from].push_back(m-2);
24         g[to].push_back(m-1);
25     }
26
27     bool BFS() {
28         memset(vis, 0, sizeof(vis));
29         queue<int> Q;
30         Q.push(t);
31         vis[t] = 1;
32         d[t] = 0;
33         while(!Q.empty()) {
34             int x = Q.front(); Q.pop();
35             for(int i = 0; i < g[x].size(); i++) {
36                 Edge& e = edges[g[x][i]^1];
37                 if(!vis[e.from] && e.cap > e.flow) {
38                     vis[e.from] = 1;
39                     d[e.from] = d[x] + 1;
40                     Q.push(e.from);
41                 }
42             }
43         }
44         return vis[s];
45     }
46
47     void init(int nn) {
48         n = nn;
49         mem(g, 0);
50         edges.clear();
51     }
52
53     void ClearFlow() {
54         for(int i = 0; i < edges.size(); i++) edges[i].
            flow = 0;

```

```

55     }
56
57     int Augment() {
58         int x = t, a = inf;
59         while(x != s) {
60             Edge& e = edges[p[x]];
61             a = min(a, e.cap-e.flow);
62             x = edges[p[x]].from;
63         }
64         x = t;
65         while(x != s) {
66             edges[p[x]].flow += a;
67             edges[p[x]^1].flow -= a;
68             x = edges[p[x]].from;
69         }
70         return a;
71     }
72
73     int max_flow(int s, int t, int need) {
74         this->s = s; this->t = t;
75         int flow = 0;
76         BFS();
77         memset(num, 0, sizeof(num));
78         for(int i = 0; i < n; i++) num[d[i]]++;
79         int x = s;
80         memset(cur, 0, sizeof(cur));
81         while(d[s] < n) {
82             if(x == t) {
83                 flow += Augment();
84                 if(flow >= need) return flow;
85                 x = s;
86             }
87             int ok = 0;
88             for(int i = cur[x]; i < g[x].size(); i++) {
89                 Edge& e = edges[g[x][i]];
90                 if(e.cap > e.flow && d[x] == d[e.to] + 1) {
91                     // Advance
92                     ok = 1;
93                     p[e.to] = g[x][i];
94                     cur[x] = i; // 注意
95                     x = e.to;
96                     break;
97                 }
98             }
99             if(!ok) { // Retreat

```

```

99         int m = n-1; // 初值注意
100         for(int i = 0; i < g[x].size(); i++) {
101             Edge& e = edges[g[x][i]];
102             if(e.cap > e.flow) m = min(m, d[e.to]);
103         }
104         if(--num[d[x]] == 0) break;
105         num[d[x] = m+1]++;
106         cur[x] = 0; // 注意
107         if(x != s) x = edges[p[x]].from;
108     }
109 }
110 return flow;
111 }
112
113 vector<int> Mincut() { // call this after maxflow
114     BFS();
115     vector<int> ans;
116     for(int i = 0; i < edges.size(); i++) {
117         Edge& e = edges[i];
118         if(!vis[e.from] && vis[e.to] && e.cap > 0) ans.
            push_back(i);
119     }
120     return ans;
121 }
122
123 void Reduce() {
124     for(int i = 0; i < edges.size(); i++) edges[i].
        cap -= edges[i].flow;
125 }
126
127 void print() {
128     printf("Graph:\n");
129     for(int i = 0; i < edges.size(); i++)
130         printf("%d->%d, %d, %d\n", edges[i].from, edges
            [i].to, edges[i].cap, edges[i].flow);
131 }
132 };

```

4.6.3 费用流

4.6.4 无源无汇有容量下界网络的可行流

```

1 #include <bits/stdc++.h>

```



```

2  #define log(x) cout << #x << "□=□" << (x) << endl
3  #define mem(x, y) memset((x), (y), sizeof((x)))
4  #define rep(i, l, r) for (int (i) = (l); (i) <= (r); (i)
    ++)
5  using namespace std;
6  typedef long long ll;
7  /*---金牌---*/
8  const int N = 200 + 9;
9  const int INF = 0x3f3f3f3f;
10 const int MAXN = 10*N;
11 int low[MAXN*10];
12 int n, m;
13 struct Dinic
14 {
15     struct Edge
16     {
17         int from, to, cap, flow;
18     };
19     vector<Edge> edges;
20     vector<int> g[MAXN];
21     void init() {
22         mem(g, 0);
23         edges.clear();
24     }
25     void add_edge(int from, int to, int cap) {
26         edges.push_back((Edge){from, to, cap, 0});
27         edges.push_back((Edge){to, from, 0, 0});
28         int m = edges.size();
29         g[from].push_back(m - 2);
30         g[to].push_back(m - 1);
31     }
32     int s, t;
33     bool vis[MAXN];
34     int d[MAXN], cur[MAXN];
35     bool BFS()
36     {
37         memset(vis, 0, sizeof(vis));
38         queue<int> q;
39         q.push(s);
40         d[s] = 0;
41         vis[s] = true;
42         while (!q.empty()) {
43             int x = q.front(); q.pop();
44             for (int i = 0; i < g[x].size(); i++) {
45                 Edge& e = edges[g[x][i]];
46                 if (!vis[e.to] && e.cap > e.flow) {

```

```

47         vis[e.to] = true;
48         d[e.to] = d[x] + 1;
49         q.push(e.to);
50     }
51 }
52 }
53 return vis[t];
54 }
55 int DFS(int x, int a)
56 {
57     if (x == t || a == 0) return a;
58     int flow = 0;
59     int f;
60     for (int& i = cur[x]; i < g[x].size(); i++) {
61         Edge& e = edges[g[x][i]];
62         if (d[x] + 1 == d[e.to] && (f = DFS(e.to, min(a,
63             e.cap - e.flow))) > 0) {
64             e.flow += f;
65             edges[g[x][i]^1].flow -= f;
66             flow += f;
67             a -= f;
68             if (a == 0) break;
69         }
70     }
71     return flow;
72 }
73 int max_flow(int ss, int tt)
74 {
75     s = ss, t = tt;
76     int flow = 0;
77     while (BFS()) {
78         memset(cur, 0, sizeof(cur));
79         flow += DFS(s, INF);
80     }
81     return flow;
82 }
83 void solve(int source, int sink)
84 {
85     max_flow(source, sink);
86     int sz = g[source].size() - 1;
87     bool flag = true;
88     rep(i, 0, sz) {
89         if (edges[g[source][i]].flow < edges[g[source]
90             ][i].cap) {
91             puts("NO");
92             flag = false;
93         }
94     }
95 }

```

```

91         }
92     }
93     if (flag) {
94         puts("YES");
95         rep(i, 1, m) {
96             //log(low[i]);
97             printf("%d\n", edges[i*2+1].cap-edges[i
               *2+1].flow + low[i]);
98         }
99     }
100 }
101 };
102 void solve(Dinic &solver)
103 {
104     solver.init();
105     scanf("%d", &m);
106     int source = n + 1, sink = n + 2;
107     // sink -> source cap : inf
108     solver.add_edge(sink, source, INF);
109     int in[N];
110     mem(in, 0);
111     mem(low, 0);
112     rep(i, 1, m) {
113         int u, v, b, c;
114         scanf("%d%d%d%d", &u, &v, &b, &c);
115         low[i] = b; // 流量下界
116         in[u] -= b;
117         in[v] += b;
118         solver.add_edge(u, v, c - b);
119     }
120     rep(i, 1, n) {
121         if (in[i] < 0) {
122             solver.add_edge(i, sink, -in[i]);
123         } else if (in[i] > 0) {
124             solver.add_edge(source, i, in[i]);
125         }
126     }
127     solver.solve(source, sink);
128 }
129 int main()
130 {
131     Dinic solver;
132     while (~scanf("%d", &n)) solve(solver);
133 }

```

4.7 其他

4.7.1 层次遍历

每次记录队列当前的 sz 然后在一次循环中只出队 sz 次

4.7.2 图解序列

图解序列：一系列非负整数可以构成一个简单图的度序列

Havel 定理：对于 $n > 1$, 长度为 n 的整数序列是图解序列当且仅当 d' 是图解序列, d' 是删除 d 中最大元素 并将紧跟的 $d[1]$ 个元素减 1 后序列

Chapter 5

计算几何

5.1 判断凸包

```
1 #include <cstdio>
2 #include <iostream>
3 using namespace std;
4 const int maxn = 1e8;
5 struct P
6 {
7     int x, y;
8 };
9 int ccw(P a, P b, P c)
10 {
11     return (b.x - a.x) * (c.y - a.y) - (c.x - a.x) * (b.y -
        a.y);
12 }
13 P a[maxn];
14 int main()
15 {
16     for (;;)
17     {
18         int n;
19         cin >> n;
20         if (n == 0) break;
21         for (int i = 0; i < n; i++)
22             cin >> a[i].x >> a[i].y;
23         bool flag = true;
24         a[n++] = a[0];
25         a[n++] = a[1];
26         for (int i = 0; i < n - 3; i++)
```

```

27         if (ccw(a[i], a[i + 1], a[i + 2]) * ccw(a[i + 1], a
28             [i + 2], a[i + 3]) < 0)
29         {
30             flag = false;
31             break;
32         }
33     if (flag) puts("convex"); else puts("concave");
34 }

```

5.2 判断线段是否相交

```

1 struct P
2 {
3     double x, y;
4 };
5 struct Segment
6 {
7     P p, q;
8 };
9 double ccw(P a, P b, P c)
10 {
11     return (b.x - a.x) * (c.y - a.y) - (c.x - a.x) * (b.y
12         - a.y);
13 }
14 bool intersects(Segment a, Segment b)
15 {
16     if (ccw(a.p, a.q, b.p) * ccw(a.p, a.q, b.q) > 0)
17         return false;
18     if (ccw(b.p, b.q, a.p) * ccw(b.p, b.q, a.q) > 0)
19         return false;
20     return true;
21 }

```

5.3 求对称点求交点

```

1 using namespace std;
2 const double eps = 1e-10;
3 double add (double a, double b)
4 {
5     if (abs(a + b) < eps * (abs(a) + abs(b))) return 0;
6     return a + b;
7 }

```

```

8 struct P
9 {
10     double x, y;
11     P(){}
12     P(double x, double y):x(x),y(y){}
13     P operator + (P p)
14     {
15         return P(add(x, p.x), add(y, p.y));
16     }
17     P operator - (P p)
18     {
19         return P(add(x, -p.x), add(y, -p.y));
20     }
21     P operator * (double d)
22     {
23         return P(x * d, y * d);
24     }
25 };
26 double det(P p1, P p2)
27 {
28     return add(p1.x * p2.y, -p1.y * p2.x);
29 }
30 bool g_equal(double a, double b)
31 {
32     if (a > b - eps && a < b + eps) return true;
33     return false;
34 }
35 //求p关于p1, p2所成直线的对称点
36 P symmetric_point(P p, P p1, P p2)
37 {
38     //直线与x轴垂直
39     if (g_equal(p1.x, p2.x)) return P(2 * p1.x - p.x,
40                                         p.y);
41     double k = (p1.y - p2.y) / (p1.x - p2.x);
42     if (g_equal(k, 0)) return P(p.x, 2 * p1.y - p.y);
43     double x = (2*k*k*p1.x + 2*k*p.y - 2*k*p1.y - k*k*p.x +
44                 p.x) / (1 + k * k);
45     double y = p.y - (x - p.x) / k;
46     return P(x, y);
47 }
48 P intersection(P p1, P p2, P q1, P q2)
49 {
50     return p1 + (p2 - p1) * (det(q2 - q1, q1 - p1) / det(q2
51         - q1, p2 - p1));
52 }

```

5.4 终极模板

```
1 #include <bits/stdc++.h>
2 using namespace std;
3
4 struct Point {
5     double x, y;
6     Point(double x = 0, double y = 0) : x(x), y(y) {}
7 };
8
9 typedef Point Vector;
10
11 Vector operator + (Vector A, Vector B) { return Vector(A.
    x + B.x, A.y + B.y); }
12 Vector operator - (Vector A, Vector B) { return Vector(A.
    x - B.x, A.y - B.y); }
13 Vector operator * (Vector A, double p) { return Vector(A.
    x*p, A.y*p); }
14 Vector operator / (Vector A, double p) { return Vector(A.
    x/p, A.y/p); }
15
16 bool operator < (const Point& a, const Point b) {
17     return a.x < b.x || (a.x == b.x && a.y < b.y);
18 }
19
20 const double EPS = 1e-10;
21
22 int dcmp(double x) {
23     if(fabs(x) < EPS) return 0;
24     else return x < 0 ? -1 : 1;
25 }
26
27 bool operator == (const Point& a, const Point& b) {
28     return dcmp(a.x-b.x) == 0 && dcmp(a.y-b.y);
29 }
30
31 //向量a的极角
32 double Angle(const Vector& v) {
33     return atan2(v.y, v.x);\share\CodeBlocks\templates\
    wizard\console\cpp
34 }
35
36 //向量点积
37 double Dot(Vector A, Vector B) { return A.x*B.x + A.y*B.y
    ; }
```



```

38
39 //向量长度\share\CodeBlocks\templates\wizard\console\cpp
40 double Length(Vector A) { return sqrt(Dot(A, A)); }
41
42 //向量夹角
43 double Angle(Vector A, Vector B) { return acos(Dot(A, B)
    / Length(A) / Length(B)); }
44
45 //向量叉积
46 double Cross(Vector A, Vector B) { return A.x*B.y - A.y*B
    .x; }
47
48 //三角形有向面积的二倍
49 double Area2(Point A, Point B, Point C) { return Cross(B-
    A, C-A); }
50
51 //向量逆时针旋转rad度(弧度)
52 Vector Rotate(Vector A, double rad) {
53     return Vector(A.x*cos(rad)-A.y*sin(rad), A.x*sin(rad)
        +A.y*cos(rad));
54 }
55
56 //计算向量A的单位法向量。左转90°, 把长度归一。调用前确保A
    不是零向量。
57 Vector Normal(Vector A) {
58     double L = Length(A);
59     return Vector(-A.y/L, A.x/L);
60 }
61
62 /*
    *****

63 使用复数类实现点及向量的简单操作
64
65 #include <complex>
66 typedef complex<double> Point;
67 typedef Point Vector;
68
69 double Dot(Vector A, Vector B) { return real(conj(A)*B)}
70 double Cross(Vector A, Vector B) { return imag(conj(A)*B)
    ;}
71 Vector Rotate(Vector A, double rad) { return A*exp(Point
    (0, rad)); }
72

```

```

73  *****/
74
75  /*
    *****/

76  * 用直线上的一点 $p_0$ 和方向向量 $v$ 表示一条指向。直线上的所有点
     $P$ 满足 $P = P_0 + t * v$ ;
77  * 如果知道直线上的两个点则方向向量为 $B-A$ ，所以参数方程为 $A$ 
     $+(B-A)*t$ ;
78  * 当 $t$  无限制时， 该参数方程表示直线。
79  * 当 $t > 0$ 时， 该参数方程表示射线。
80  * 当  $0 < t < 1$ 时， 该参数方程表示线段。
81  *****/
    */

82
83  //直线交点,须确保两直线有唯一交点。
84  Point GetLineIntersection(Point P, Vector v, Point Q,
    Vector w) {
85      Vector u = P - Q;
86      double t = Cross(w, u)/Cross(v, w);
87      return P+v*t;
88  }
89
90  //点到直线距离
91  double DistanceToLine(Point P, Point A, Point B) {
92      Vector v1 = B - A, v2 = P - A;
93      return fabs(Cross(v1, v2) / Length(v1)); //不取绝对
    值，得到的是有向距离
94  }
95
96  //点到线段的距离
97  double DistanceToSegmentS(Point P, Point A, Point B) {
98      if(A == B) return Length(P-A);
99      Vector v1 = B-A, v2 = P-A, v3 = P-B;
100     if(dcmp(Dot(v1, v2)) < 0) return Length(v2);
101     else if(dcmp(Dot(v1, v3)) > 0) return Length(v3);
102     else return fabs(Cross(v1, v2)) / Length(v1);
103  }
104
105  //点在直线上的投影
106  Point GetLineProjection(Point P, Point A, Point B) {
107      Vector v = B - A;
108      return A+v*(Dot(v, P-A)/Dot(v, v));
109  }
110

```

```

111 //线段相交判定, 交点不在一条线段的端点
112 bool SegmentProperIntersection(Point a1, Point a2, Point
    b1, Point b2) {
113     double c1 = Cross(a2-a1, b1-a1), c2 = Cross(a2-a1, b2
        -a1);
114     double c3 = Cross(b2-b1, a1-b1), c4 = Cross(b2-b1, a2
        -b1);
115     return dcmp(c1)*dcmp(c2) < 0 && dcmp(c3)*dcmp(c4) <
        0;
116 }
117
118 //判断点是否在点段上, 不包含端点
119 bool OnSegment(Point P, Point a1, Point a2) {
120     return dcmp(Cross(a1-P, a2-P)) == 0 && dcmp((Dot(a1-P,
        a2-P)) < 0));
121 }
122
123 //计算凸多边形面积
124 double ConvexPolygonArea(Point *p, int n) {
125     double area = 0;
126     for(int i = 1; i < n-1; i++)
127         area += Cross(p[i] - p[0], p[i+1] - p[0]);
128     return area/2;
129 }
130
131 //计算多边形的有向面积
132 double PolygonArea(Point *p, int n) {
133     double area = 0;
134     for(int i = 1; i < n-1; i++)
135         area += Cross(p[i] - p[0], p[i+1] - p[0]);
136     return area/2;
137 }
138
139 /*
    *****
140 * Morley定理: 三角形每个内角的三等分线, 相交成的三角形是
    等边三角形。
141 * 欧拉定理: 设平面图的定点数, 边数和面数分别为 $V, E, F$ 。则 $V +
    F - E = 2$ ;
142 *****
    */
143
144 struct Circle {
145     Point c;

```

```

146     double r;
147
148     Circle(Point c, double r) : c(c), r(r) {}
149     //通过圆心角确定圆上坐标
150     Point point(double a) {
151         return Point(c.x + cos(a)*r, c.y + sin(a)*r);
152     }
153 };
154
155 struct Line {
156     Point p;
157     Vector v;
158     double ang;
159     Line() {}
160     Line(Point p, Vector v) : p(p), v(v) {}
161     bool operator < (const Line& L) const {
162         return ang < L.ang;
163     }
164 };
165
166 //直线和圆的交点，返回交点个数，结果存在sol中。
167 //该代码没有清空sol。
168 int getLineCircleInterseccion(Line L, Circle C, double& t1
169 , double& t2, vector<Point>& sol) {
170     double a = L.v.x, b = L.p.x - C.c.x, c = L.v.y, d = L
171         .p.y - C.c.y;
172     double e = a*a + c*c, f = 2*(a*b + c*d), g = b*b + d*
173         d - C.r*C.r;
174     double delta = f*f - 4*e*g;
175     if(dcmp(delta) < 0) return 0; //相离
176     if(dcmp(delta) == 0) { //相切
177         t1 = t2 = -f / (2*e);
178         sol.push_back(C.point(t1));
179         return 1;
180     }
181     //相交
182     t1 = (-f - sqrt(delta)) / (2*e); sol.push_back(C.
183         point(t1));
184     t2 = (-f + sqrt(delta)) / (2*e); sol.push_back(C.
185         point(t2));
186     return 2;
187 }
188
189 //两圆相交

```

```

185 int getCircleCircleIntersection(Circle C1, Circle C2,
    vector<Point>& sol) {
186     double d = Length(C1.c - C2.c);
187     if(dcmp(d) == 0) {
188         if(dcmp(C1.r - C2.r == 0)) return -1;    //两圆完
            全重合
189         return 0;                                //同心
            圆, 半径不一样
190     }
191     if(dcmp(C1.r + C2.r - d) < 0) return 0;
192     if(dcmp(fabs(C1.r - C2.r) == 0)) return -1;
193
194     double a = Angle(C2.c - C1.c);                //向量
            C1C2的极角
195     double da = acos((C1.r*C1.r + d*d - C2.r*C2.r) / (2*
            C1.r*d));
196     //C1C2到C1P1的角
197     Point p1 = C1.point(a-da), p2 = C1.point(a+da);
198     sol.push_back(p1);
199     if(p1 == p2) return 1;
200     sol.push_back(p2);
201     return 2;
202 }
203
204 const double PI = acos(-1);
205 //过定点做圆的切线
206 //过点p做圆C的切线, 返回切线个数。v[i]表示第i条切线
207 int getTangents(Point p, Circle C, Vector* v) {
208     Vector u = C.c - p;
209     double dist = Length(u);
210     if(dist < C.r) return 0;
211     else if(dcmp(dist - C.r) == 0) {
212         v[0] = Rotate(u, PI/2);
213         return 1;
214     } else {
215         double ang = asin(C.r / dist);
216         v[0] = Rotate(u, -ang);
217         v[1] = Rotate(u, +ang);
218         return 2;
219     }
220 }
221
222 //两圆的公切线
223 //返回切线的个数, -1表示有无数条公切线。
224 //a[i], b[i] 表示第i条切线在圆A, 圆B上的切点

```

```

225 int getTangents(Circle A, Circle B, Point *a, Point *b) {
226     int cnt = 0;
227     if(A.r < B.r) {
228         swap(A, B); swap(a, b);
229     }
230     int d2 = (A.c.x - B.c.x)*(A.c.x - B.c.x) + (A.c.y - B
        .c.y)*(A.c.y - B.c.y);
231     int rdiff = A.r - B.r;
232     int rsum = A.r + B.r;
233     if(d2 < rdiff*rdiff) return 0;    //内含
234     double base = atan2(B.c.y - A.c.y, B.c.x - A.c.x);
235     if(d2 == 0 && A.r == B.r) return -1;    //无限多条切线
236     if(d2 == rdiff*rdiff) {           //内切一条切线
237         a[cnt] = A.point(base);
238         b[cnt] = B.point(base);
239         cnt++;
240         return 1;
241     }
242     //有外共切线
243     double ang = acos((A.r-B.r) / sqrt(d2));
244     a[cnt] = A.point(base+ang); b[cnt] = B.point(base+ang
        ); cnt++;
245     a[cnt] = A.point(base-ang); b[cnt] = B.point(base-ang
        ); cnt++;
246     if(d2 == rsum*rsum) { //一条公切线
247         a[cnt] = A.point(base);
248         b[cnt] = B.point(PI+base);
249         cnt++;
250     } else if(d2 > rsum*rsum) { //两条公切线
251         double ang = acos((A.r + B.r) / sqrt(d2));
252         a[cnt] = A.point(base+ang); b[cnt] = B.point(PI+
            base+ang); cnt++;
253         a[cnt] = A.point(base-ang); b[cnt] = B.point(PI+
            base-ang); cnt++;
254     }
255     return cnt;
256 }
257
258 typedef vector<Point> Polygon;
259
260 //点在多边形内的判定
261 int isPointInPolygon(Point p, Polygon poly) {
262     int wn = 0;
263     int n = poly.size();
264     for(int i = 0; i < n; i++) {

```

```

265         if(OnSegment(p, poly[i], poly[(i+1)%n])) return
           -1; //在边界上
266         int k = dcmp(Cross(poly[(i+1)%n]-poly[i], p-poly[
           i]));
267         int d1 = dcmp(poly[i].y - p.y);
268         int d2 = dcmp(poly[(i+1)%n].y - p.y);
269         if(k > 0 && d1 <= 0 && d2 > 0) wn++;
270         if(k < 0 && d2 <= 0 && d1 > 0) wn++;
271     }
272     if(wn != 0) return 1;          //内部
273     return 0;                     //外部
274 }
275
276 //凸包
277 /*
    *****

278 * 输入点数组p, 个数为p, 输出点数组ch。 返回凸包顶点数
279 * 不希望凸包的边上有输入点, 把两个<= 改成 <
280 * 高精度要求时建议用dcmp比较
281 * 输入点不能有重复点。函数执行完以后输入点的顺序被破坏
282 *****
    */
283 int ConvexHull(Point *p, int n, Point* ch) {
284     sort(p, p+n);          //先比较x坐标, 再比较y坐标
285     int m = 0;
286     for(int i = 0; i < n; i++) {
287         while(m > 1 && Cross(ch[m-1] - ch[m-2], p[i]-ch[m-2]) <= 0) m--;
288         ch[m++] = p[i];
289     }
290     int k = m;
291     for(int i = n-2; i >= 0; i++) {
292         while(m > k && Cross(ch[m-1] - ch[m-2], p[i]-ch[m-2]) <= 0) m--;
293         ch[m++] = p[i];
294     }
295     if(n > 1) m--;
296     return m;
297 }
298
299 //用有向直线A->B切割多边形poly, 返回“左侧”。 如果退化, 可能会返回一个单点或者线段
300 //复杂度O(n^2);
301 Polygon CutPolygon(Polygon poly, Point A, Point B) {

```

```

302     Polygon newpoly;
303     int n = poly.size();
304     for(int i = 0; i < n; i++) {
305         Point C = poly[i];
306         Point D = poly[(i+1)%n];
307         if(dcmp(Cross(B-A, C-A)) >= 0) newpoly.push_back(
            C);
308         if(dcmp(Cross(B-A, C-D)) != 0) {
309             Point ip = GetLineIntersection(A, B-A, C, D-C
            );
310             if(OnSegment(ip, C, D)) newpoly.push_back(ip)
            ;
311         }
312     }
313     return newpoly;
314 }
315
316 //半平面交
317
318 //点p再有向直线L的左边。（线上不算）
319 bool Onleft(Line L, Point p) {
320     return Cross(L.v, p-L.p) > 0;
321 }
322
323 //两直线交点，假定交点唯一存在
324 Point GetIntersection(Line a, Line b) {
325     Vector u = a.p - b.p;
326     double t = Cross(b.v, u) / Cross(a.v, b.v);
327     return a.p+a.v*t;
328 }
329
330 int HalfplaneIntersection(Line* L, int n, Point* poly) {
331     sort(L, L+n);           //按极角排序
332
333     int first, last;         //双端队列的第一个元素和
        最后一个元素
334     Point *p = new Point[n]; //p[i]为q[i]和q[i+1]的交
        点
335     Line *q = new Line[n];   //双端队列
336     q[first = last = 0] = L[0]; //队列初始化为只有一个半
        平面L[0]
337     for(int i = 0; i < n; i++) {
338         while(first < last && !Onleft(L[i], p[last-1]))
            last--;

```



```

339         while(first < last && !Onleft(L[i], p[first]))
340             first++;
341         q[++last] = L[i];
342         if(fabs(Cross(q[last].v, q[last-1].v)) < EPS) {
343             last--;
344             if(Onleft(q[last], L[i].p)) q[last] = L[i];
345         }
346         if(first < last) p[last-1] = GetIntersection(q[
347             last-1], q[last]);
348     }
349     while(first < last && !Onleft(q[first], p[last-1]))
350         last--;
351     //删除无用平面
352     if(last-first <= 1) return 0;    //空集
353     p[last] = GetIntersection(q[last], q[first]);
354     //从deque复制到输出中
355     int m = 0;
356     for(int i = first; i <= last; i++) poly[m++] = p[i];
357     return m;
358 }
359 int main() {
360     return 0;
361 }

```

5.5 K 次圆

```

1  #include <bits/stdc++.h>
2  #define rep(i, l, r) for (int(i) = (l); (i) <= (r); (i)
3      ++)
4  #define mem(x, y) memset((x), (y), sizeof(x))
5  using namespace std;
6  const double eps = 1e-8;
7  const int N = 1000 + 9;
8  const double pi = 3.1415926535;
9  double sqr(double x)
10 {
11     return x * x;
12 }
13 int dcmp(double x)
14 {
15     if (fabs(x) < eps)
16         return 0;

```

```

16     else
17         return x < 0 ? -1 : 1;
18 }
19 struct Circle
20 {
21     double x, y, r, angle;
22     int d;
23     Circle() {}
24     Circle(double xx, double yy, double ang = 0, int t =
        0)
25     {
26         x = xx;
27         y = yy;
28         angle = ang;
29         d = t;
30     }
31     void read()
32     {
33         scanf("%lf%lf%lf", &x, &y, &r);
34         d = 1;
35     }
36 };
37
38 Circle cir[N], tp[N * 2];
39 double area[N];
40
41 double dis(Circle a, Circle b)
42 {
43     return sqrt(sqr(a.x - b.x) + sqr(a.y - b.y));
44 }
45
46 double cross(Circle p0, Circle p1, Circle p2)
47 {
48     return (p1.x - p0.x) * (p2.y - p0.y) - (p1.y - p0.y)
        * (p2.x - p0.x);
49 }
50 //圆相交
51 int CirCrossCir(Circle p1, double r1, Circle p2, double
    r2, Circle &cp1, Circle &cp2)
52 {
53     double mx = p2.x - p1.x, sx = p2.x + p1.x, mx2 = mx *
        mx;
54     double my = p2.y - p1.y, sy = p2.y + p1.y, my2 = my *
        my;
55     double sq = mx2 + my2, d = -(sq - sqr(r1 - r2)) * (sq
        - sqr(r1 + r2));

```

```

56     if (d + eps < 0)
57         return 0;
58     if (d < eps)
59         d = 0;
60     else
61         d = sqrt(d);
62     double x = mx * ((r1 + r2) * (r1 - r2) + mx * sx) +
        sx * my2;
63     double y = my * ((r1 + r2) * (r1 - r2) + my * sy) +
        sy * mx2;
64     double dx = mx * d, dy = my * d;
65     sq *= 2;
66     cp1.x = (x - dy) / sq;
67     cp1.y = (y + dx) / sq;
68     cp2.x = (x + dy) / sq;
69     cp2.y = (y - dx) / sq;
70     if (d > eps)
71         return 2;
72     else
73         return 1;
74 }
75 bool circmp(const Circle &u, const Circle &v)
76 {
77     return dcmp(u.r - v.r) < 0;
78 }
79
80 bool cmp(const Circle &u, const Circle &v)
81 {
82     if (dcmp(u.angle - v.angle))
83         return u.angle < v.angle;
84     return u.d > v.d;
85 }
86 //0.5*r*r*(K-sin(K))
87 double calc(Circle cir, Circle cp1, Circle cp2)
88 {
89     double ans = (cp2.angle - cp1.angle) * sqr(cir.r) -
        cross(cir, cp1, cp2) + cross(Circle(0, 0), cp1,
        cp2);
90     return ans / 2;
91 }
92
93 void CirUnion(Circle cir[], int n)
94 {
95     Circle cp1, cp2;
96     sort(cir, cir + n, circmp);
97     for (int i = 0; i < n; ++i)

```

```

198         for (int j = i + 1; j < n; ++j)
199             if (dcmp(dis(cir[i], cir[j]) + cir[i].r - cir
200                     [j].r) <= 0)
201                 cir[i].d++;
202     for (int i = 0; i < n; ++i)
203     {
204         int tn = 0, cnt = 0;
205         for (int j = 0; j < n; ++j)
206         {
207             if (i == j)
208                 continue;
209             if (CirCrossCir(cir[i], cir[i].r, cir[j], cir
210                             [j].r,
211                             cp2, cp1) < 2)
212                 continue;
213             cp1.angle = atan2(cp1.y - cir[i].y, cp1.x -
214                               cir[i].x);
215             cp2.angle = atan2(cp2.y - cir[i].y, cp2.x -
216                               cir[i].x);
217             cp1.d = 1;
218             tp[tn++] = cp1;
219             cp2.d = -1;
220             tp[tn++] = cp2;
221             if (dcmp(cp1.angle - cp2.angle) > 0)
222                 cnt++;
223         }
224         tp[tn++] = Circle(cir[i].x - cir[i].r, cir[i].y,
225                           pi, -cnt);
226         tp[tn++] = Circle(cir[i].x - cir[i].r, cir[i].y,
227                           -pi, cnt);
228         sort(tp, tp + tn, cmp);
229         int p, s = cir[i].d + tp[0].d;
230         for (int j = 1; j < tn; ++j)
231         {
232             p = s;
233             s += tp[j].d;
234             area[p] += calc(cir[i], tp[j - 1], tp[j]);
235         }
236     }
237 }
238
239 int n;
240 void solve()
241 {
242     for (int i = 0; i < n; i++) cir[i].read(); // read x
243                                     y r
244     mem(area, 0);

```

```

137     CirUnion(cir, n);
138     for (int i = 1; i <= n; i++)
139     {
140         area[i] -= area[i + 1];
141         printf("[%d]_=%%.3f\n", i, area[i]);
142     }
143 }
144
145 int main()
146 {
147     while (~scanf("%d", &n)) solve();
148 }

```

Chapter 6

数据结构

6.1 手写堆

```
1 struct Node
2 {
3     int index;
4     int tag;
5     int l;
6     bool operator<(Node ano) const
7     {
8         return l!=ano.l?l<ano.l:index>ano.index;
9     }
10 };
11 Node heap[51111];
12 int heapsz;
13 int a[2<<21]; // 编号为i的在heap的位置
14 void swim(int p)
15 {
16     while (p > 1 && heap[p / 2] < heap[p])
17     {
18         swap(a[heap[p].tag],a[heap[p/2].tag]);
19         swap(heap[p],heap[p/2]);
20         p/=2;
21     }
22 }
23 void sink(int p)
24 {
25     while(p * 2 <= heapsz)
26     {
27         p *= 2;
28         if (p+1<=heapsz) p|=heap[p]<heap[p+1];
```

```

29         swap(a[heap[p].tag],a[heap[p/2].tag]);
30         swap(heap[p],heap[p/2]);
31     }
32 }

```

6.2 左偏树

```

1  #include <bits/stdc++.h>
2  #define rep(i, l, r) for (int (i) = (l); (i) <= (r); (i)
    ++))
3  using namespace std;
4  int n,m,ans;
5  struct Data{int fa,x,dis,l,r;};
6  const int MAXN = 100001;
7  Data q[MAXN];
8  int find(int x)
9  {
10     return x==q[x].fa?x:q[x].fa=find(q[x].fa);
11 }
12 int merge(int x,int y)
13 {
14     if(!x) return y;
15     if(!y) return x;
16     if(q[x].x<q[y].x) swap(x,y);
17     q[x].r=merge(q[x].r,y);
18     q[q[x].r].fa=x;
19     if(q[q[x].l].dis<q[q[x].r].dis) swap(q[x].l,q[x].r);
20     if(!q[x].r) q[x].dis=0;
21     else q[x].dis=q[q[x].r].dis+1;
22     return x;
23 }
24 int pop(int x)
25 {
26     int l=q[x].l,r=q[x].r;
27     q[l].fa=l,q[r].fa=r;
28     q[x].l=q[x].r=q[x].dis=0;
29     return merge(l,r);
30 }
31 void solve()
32 {
33     int x,y,u,v,l1,l2;
34     rep(i, 1, n) {
35         scanf("%d", &q[i].x);
36         q[i].fa=i;

```

```

37         q[i].l=q[i].r=q[i].dis=0;
38     }
39     scanf("%d", &m);
40     rep(i, 1, m) {
41         scanf("%d%d", &x, &y);
42         l1=find(x), l2=find(y);
43         if(l1 == l2){printf("-1\n");continue;}
44         q[l1].x/=2;u=pop(l1);u=merge(u,l1);
45         q[l2].x/=2;v=pop(l2);v=merge(v,l2);
46         ans=merge(u,v);
47         printf("%d\n",q[ans].x);
48     }
49 }
50 int main()
51 {
52     while(~scanf("%d",&n)) slove();
53 }

```

6.3 两优先队列模拟堆

```

1
2 最简分数实数逼近
3 a1 = 100000, a2 = 1, ans = 100000;
4 ll x1 = 0, y1 = 1, xr = 1, yr = 0, xm, ym;
5 while(y1 + yr <= 100000) {
6     xm = x1 + xr;
7     ym = y1 + yr;
8     if(check(xm, ym, a1, a2, kk)) {a1 = xm, a2 = ym;}
9     if(get(xm, ym, kk)) {xr = xm, yr = ym;}
10    else {x1 = xm, y1 = ym;}
11 }
12 struct heap {
13     priority_queue<int> p, q;
14     void init() {
15         while(!p.empty()) p.pop();
16         while(!q.empty()) q.pop();
17     }
18     void add(int x) {p.push(x);}
19     void del(int x) {q.push(x);}
20     int top() {
21         while(1) {
22             if(p.empty()) return -INF;
23             else if(!q.empty() && p.top() == q.top()) {p.
                pop(); q.pop();}

```



```

24         else return p.top();
25     }
26 }
27 int toptwo() {
28     int a = top(); del(a);
29     int b = top(); add(a);
30     if(b == -INF) return a == -INF ? a : 0;
31     else return max(a + b, 0);
32 }
33 }

```

6.4 线段树

```

1  #include <bits/stdc++.h>
2  using namespace std;
3  const int MAXM = 1E6 + 9;
4  const int MAXR = 20 + 9;
5  const int INF = INT_MAX / 10;
6  struct Node
7  {
8      int setv, addv, sumv, minv, maxv;
9      Node()
10     {
11         setv = -1, addv = 0, sumv = 0, minv = 0, maxv =
12         0;
13     }
14 };
15 struct Tree
16 {
17     Node tree[4 * MAXM];
18     int mid(int l, int r)
19     {
20         return l + (r-l) / 2;
21     }
22     void pushdown(int o)
23     {
24         int lc(o*2), rc(o*2+1);
25         if (tree[o].setv >= 0) {
26             tree[lc].setv = tree[rc].setv = tree[o].setv;
27             tree[lc].addv = tree[rc].addv = 0;
28             tree[o].setv = -1;
29         }
30         if (tree[o].addv) {
31             tree[lc].addv += tree[o].addv;

```

```

31         tree[rc].addv += tree[o].addv;
32         tree[o].addv = 0;
33     }
34 }
35 void maintain(int o, int l, int r)
36 {
37     int lc(o*2), rc(o*2+1);
38     if (r > l) {
39         tree[o].sumv = tree[lc].sumv + tree[rc].sumv;
40         tree[o].minv = min(tree[lc].minv, tree[rc].
41             minv);
42         tree[o].maxv = max(tree[lc].maxv, tree[rc].
43             maxv);
44     }
45     if (tree[o].setv >= 0) {
46         tree[o].minv = tree[o].maxv = tree[o].setv;
47         tree[o].sumv = (r-l+1)*tree[o].setv;
48     }
49     if (tree[o].addv) {
50         tree[o].minv += tree[o].addv;
51         tree[o].maxv += tree[o].addv;
52         tree[o].sumv += (r-l+1) * tree[o].addv;
53     }
54 }
55 void update(int o, int l, int r, int op, int ql, int
56     qr, int v)
57 {
58     int lc(o*2), rc(o*2+1);
59     if (ql <= l && qr >= r) {
60         if (op == 1) tree[o].addv += v;
61         else {
62             tree[o].setv = v;
63             tree[o].addv = 0;
64         }
65     } else {
66         pushdown(o);
67         int m = mid(l, r);
68         if (ql <= m) update(lc, l, m, op, ql, qr, v);
69         else maintain(lc, l, m);
70         if (qr > m) update(rc, m+1, r, op, ql, qr, v)
71             ; else maintain(rc, m+1, r);
72     }
73     maintain(o, l, r);
74 }
75 Node query(int o, int l, int r, int ql, int qr)
76 {

```

```

72     Node res;
73     int lc(o*2), rc(o*2+1), m(mid(l, r));
74     maintain(o, l, r);
75     if (ql <= l && qr >= r) {
76         res.sumv = tree[o].sumv;
77         res.minv = tree[o].minv;
78         res.maxv = tree[o].maxv;
79     } else {
80         pushdown(o);
81         Node lres;
82         lres.minv = INF, lres.maxv = -INF;
83         Node rres(lres);
84         if (ql <= m) lres = query(lc, l, m, ql, qr);
85         if (qr > m) rres = query(rc, m+1, r, ql, qr);
86         res.sumv = lres.sumv + rres.sumv;
87         res.minv = min(lres.minv, rres.minv);
88         res.maxv = max(lres.maxv, rres.maxv);
89     }
90     return res;
91 }
92 };
93 Tree tree[MAXR];
94 void solve(int r, int c, int m)
95 {
96     for (int i(1); i <= r; i++) tree[i].tree[1].setv = 0;
97     while (m--) {
98         int op, x1, y1, x2, y2;
99         scanf("%d%d%d%d", &op, &x1, &y1, &x2, &y2);
100         if (op < 3) {
101             int v;
102             scanf("%d", &v);
103             for (int i(x1); i <= x2; i++) tree[i].update
104                 (1, 1, c, op, y1, y2, v);
105         } else {
106             Node ans;
107             ans.sumv = 0, ans.minv = INF, ans.maxv = -INF;
108             for (int i(x1); i <= x2; i++) {
109                 Node res(tree[i].query(1, 1, c, y1, y2));
110                 ans.sumv += res.sumv;
111                 ans.minv = min(ans.minv, res.minv);
112                 ans.maxv = max(ans.maxv, res.maxv);
113             }
114         }
115     }
116 }

```

```

113         printf("%d_%d_%d\n", ans.sumv, ans.minv, ans.
                maxv);
114     }
115 }
116 }
117 int main()
118 {
119     //freopen("in", "r", stdin);
120     int r, c, m;
121     while (scanf("%d%d%d", &r, &c, &m) != EOF) {
122         solve(r, c, m);
123     }
124 }

```

6.5 二维线段树

```

1  #include <bits/stdc++.h>
2  using namespace std;
3  const int MAXN = 800 * 4 + 10;
4  int mx[MAXN][MAXN], mn[MAXN][MAXN], g[810][810];
5  int n;
6  void buildy(int xo, int o, int l, int r, int x) {
7      if(l == r) {
8          if(x != -1) mx[xo][o] = mn[xo][o] = g[x][l];
9          else {
10             int xlch = xo * 2, xrch = xlch + 1;
11             mx[xo][o] = max(mx[xlch][o], mx[xrch][o]);
12             mn[xo][o] = min(mn[xlch][o], mn[xrch][o]);
13         }
14         return;
15     }
16     int mid = (l + r) / 2, lch = o * 2, rch = lch + 1;
17     buildy(xo, lch, l, mid, x);
18     buildy(xo, rch, mid + 1, r, x);
19     mx[xo][o] = max(mx[xo][lch], mx[xo][rch]);
20     mn[xo][o] = min(mn[xo][lch], mn[xo][rch]);
21 }
22 void buildx(int o, int l, int r) {
23     if(l == r) {
24         buildy(o, 1, 1, n, l);
25         return;
26     }
27     int mid = (l + r) / 2, lch = o * 2, rch = lch + 1;
28     buildx(lch, l, mid);

```

```

29     buildx(rch, mid + 1, r);
30     buildy(o, 1, 1, n, -1);
31 }
32 void updatey(int xo, int o, int l, int r, int y, int v) {
33     if(l == r) {
34         if(v != -1) mn[xo][o] = mx[xo][o] = v;
35         else {
36             int xlch = xo * 2, xrch = xlch + 1;
37             mx[xo][o] = max(mx[xlch][o], mx[xrch][o]);
38             mn[xo][o] = min(mn[xlch][o], mn[xrch][o]);
39         }
40         return;
41     }
42     int mid = (l + r) / 2, lch = o * 2, rch = lch + 1;
43     if(y <= mid) updatey(xo, lch, l, mid, y, v);
44     else updatey(xo, rch, mid + 1, r, y, v);
45     mx[xo][o] = max(mx[xo][lch], mx[xo][rch]);
46     mn[xo][o] = min(mn[xo][lch], mn[xo][rch]);
47 }
48 void updatex(int o, int l, int r, int x, int y, int v) {
49     if(l == r) {
50         updatey(o, 1, 1, n, y, v);
51         return;
52     }
53     int mid = (l + r) / 2, lch = o * 2, rch = lch + 1;
54     if(mid >= x) updatex(lch, l, mid, x, y, v);
55     else updatex(rch, mid + 1, r, x, y, v);
56     updatey(o, 1, 1, n, y, -1);
57 }
58 int minv, maxv;
59 void queryy(int xo, int o, int l, int r, int yl, int yr)
60 {
61     if(yl <= l && r <= yr) {
62         minv = min(minv, mn[xo][o]);
63         maxv = max(maxv, mx[xo][o]);
64         return;
65     }
66     int mid = (l + r) / 2, lch = o * 2, rch = lch + 1;
67     if(mid >= yl) queryy(xo, lch, l, mid, yl, yr);
68     if(mid < yr) queryy(xo, rch, mid + 1, r, yl, yr);
69 }
70 void queryx(int o, int l, int r, int xl, int xr, int yl,
71 int yr) {
72     if(xl <= l && r <= xr) {
73         queryy(o, 1, 1, n, yl, yr);
74         return;
75     }

```

```

73     }
74     int mid = (l + r) / 2, lch = o * 2, rch = lch + 1;
75     if(mid >= xl) queryx(lch, l, mid, xl, xr, yl, yr);
76     if(mid < xr) queryx(rch, mid + 1, r, xl, xr, yl, yr);
77 }
78 int main() {
79     int T; scanf("%d", &T);
80     int cas = 1;
81     while(T--) {
82         scanf("%d", &n);
83         for(int i = 1; i <= n; i++)
84             for(int j = 1; j <= n; j++)
85                 scanf("%d", &g[i][j]);
86         buildx(1, 1, n);
87         int q; scanf("%d", &q);
88         printf("Case_#%d:\n", cas++);
89         while(q--) {
90             int x, y, len; scanf("%d%d%d", &x, &y, &len);
91             int xl = max(1, x - len / 2), xr = min(n, x +
                len / 2), yl = max(1, y - len / 2), yr =
                min(n, y + len / 2);
92             minv = 1 << 30, maxv = -minv;
93             queryx(1, 1, n, xl, xr, yl, yr);
94             int ans = (maxv + minv) / 2;
95             updatex(1, 1, n, x, y, ans);
96             printf("%d\n", ans);
97         }
98     }
99 }

```

6.6 Treap

```

1 struct node {
2     node* ch[2];
3     int w, sum, v, k;
4     void maintain() {sum = w + ch[0]->sum + ch[1]->
        sum;}
5     int cmp(int vv) {return vv == v ? -1 : vv > v;}
6     bool operator < (const node& rhs) {return k < rhs
        .k;}
7     }nd[MAXN * 10];
8     node* null, *rt;
9     int ncnt;
10    void rot(node*& o, int d) {

```

```

11     node* k = o->ch[d^1]; o->ch[d^1] = k->ch[d]; k->ch[d]
        = o;
12     //一定先维护o, 因为o是k的子节点
13     o->maintain(); k->maintain();
14     o = k;
15 }
16 void ins(node*& o, int x) {
17     if(o == null) {
18         o = &nd[ncnt++];
19         *o = *null;
20         o->v = x; o->k = rand();
21         o->w = 1; o->maintain();
22     }
23     else {
24         int d = o->cmp(x);
25         if(d == -1) {o->w++; o->sum++;}
26         else {
27             ins(o->ch[d], x);
28             o->maintain();
29             if(o < o->ch[d]) rot(o, d^1);
30         }
31     }
32 }
33 void del(node*& o, int x) {
34     if(o == null) return;
35     int d = o->cmp(x);
36     if(d == -1) {
37         if(o->w > 1) {o->w--; o->sum--;}
38         else {
39             if(o->ch[0] != null && o->ch[1] != null)
40             {
41                 int dd = o->ch[0] > o->ch[1];
42                 rot(o, dd);
43                 del(o->ch[dd], x);
44             }
45             else o = o->ch[0] == null ? o->ch[1] : o
46                 ->ch[0];
47         }
48     }
49     else del(o->ch[d], x);
50     o->maintain();
51 }
52 int getrk(node* o, int x) {
53     if(o == null) return 0;
54     int d = o->cmp(x);
55     if(d == -1) return o->ch[0]->sum + 1;

```

```

54     else {
55         if(d == 1) return o->ch[0]->sum + o->w +
            getrk(o->ch[1], x);
56         return getrk(o->ch[0], x);
57     }
58 }
59 int getkth(node* o, int x) {
60     if(o == null) return 0;
61     int ls = o->ch[0]->sum, cs = o->w;
62     if(x <= ls) return getkth(o->ch[0], x);
63     else if(x > ls && x <= ls + cs) return o->v;
64     else return getkth(o->ch[1], x - ls - cs);
65 }
66 int ans;
67 //找第一个小于x的数
68 void getpre(node* o, int x) {
69     if(o == null) return;
70     if(x > o->v) {
71         ans = o->v;
72         getpre(o->ch[1], x);
73     }
74     else getpre(o->ch[0], x);
75 }
76 //第一个大于x的数
77 void getnxt(node* o, int x) {
78     if(o == null) return;
79     if(x < o->v) {
80         ans = o->v;
81         getnxt(o->ch[0], x);
82     }
83     else getnxt(o->ch[1], x);
84 }

```

6.7 splay

```

1  int ch[MAXN][2], fa[MAXN], val[MAXN], sz[MAXN], cnt[MAXN]
    ];
2  int rt, ncnt;
3  void update(int r) {sz[r] = sz[ch[r][0]] + sz[ch[r][1]] +
    cnt[r];}
4  void rot(int x) {
5      int y = fa[x], z = fa[y], d = ch[y][1] == x;
6      ch[y][d] = ch[x][d^1], fa[ch[y][d]] = y;
7      ch[x][d^1] = y; fa[y] = x;

```



```

8     fa[x] = z;
9     if(z) ch[z][ch[z][1] == y] = x;
10    update(y);
11 }
12 void splay(int r, int tp) {
13     for(int y, z; (y = fa[r]) != tp; rot(r)) {
14         z = fa[y];
15         if(z == tp) continue;
16         if((ch[z][0] == y) == (ch[y][0] == r)) rot(y);
17         else rot(r);
18     }
19     if(!tp) rt = r;
20     update(r);
21 }
22 void ins(int r, int x) {
23     int y = 0;
24     while(r && val[r] != x) {y = r; r = ch[r][val[r] < x
25         ];}
26     if(r) ++cnt[r];
27     else {
28         r = ++ncnt;
29         sz[r] = cnt[r] = 1;
30         val[r] = x; fa[r] = y;
31         ch[r][0] = ch[r][1] = 0;
32         if(y) ch[y][val[y] < x] = r;
33     }
34     splay(r, 0);
35 }
36 void get(int v) {
37     int x = rt; if(!x) return;
38     while(ch[x][val[x] < v] && val[x] != v) x = ch[x][val
39         [x] < v];
40     splay(x, 0);
41 }
42 int getrk(int v) {
43     get(v);
44     return sz[ch[rt][0]];
45 }
46 int getkth(int x) {
47     int y = rt, p;
48     if(x > sz[rt]) return 0;
49     while(1) {
50         p = ch[y][0];
51         if(sz[p] + cnt[y] < x) {
52             x -= sz[p] + cnt[y];
53             y = ch[y][1];

```

```

52     }
53     else if(sz[p] >= x) y = p;
54     else return val[y];
55 }
56 }
57 int nxt(int x, bool op) {
58     get(x);
59     if((val[rt] > x && op) || (val[rt] < x && !op))
60         return rt;
61     int p = ch[rt][op];
62     while(ch[p][op^1]) p = ch[p][op^1];
63     return p;
64 }
65 void del(int v) {
66     int p = nxt(v, 0), s = nxt(v, 1);
67     splay(p, 0);
68     splay(s, p);
69     p = ch[s][0];
70     if(cnt[p] > 1) -- cnt[p], splay(p, 0);
71     else ch[s][0] = 0;
72 }

```

6.8 倍增 LCA

```

1 void initp() {
2     for(int j = 1; (1 << j) <= n; j++)
3         for(int i = 1; i <= n; i++) if(p[i][j - 1])//一定
4             要有这个if
5             p[i][j] = p[p[i][j - 1]][j - 1];
6 }
7 int LCA(int u, int v) {
8     if(d[u] < d[v]) swap(u, v);
9     int lim;
10    //确定最大的2^lim不超过d[u]
11    for(lim = 0; (1 << lim) <= d[u]; lim++); lim--;
12    int ret = 0;
13    //把u上升到v相同的高度
14    for(int i = lim; i >= 0; i--) if(d[u] - (1 << i) >= d
15        [v]) u = p[u][i];
16    if(u == v) return u; //一定要有这个判断
17    for(int i = lim; i >= 0; i--) if(p[u][i] != p[v][i])
18        {u = p[u][i]; v = p[v][i];}
19    return p[u][0];
20 }

```

6.9 主席树

```
1 void build(node* &now, node* &pre, int l, int r, int x) {
2     now = &tn[ncnt++];
3     *now = *null;
4     int mid = (l + r) / 2;
5     if(l == r) {
6         *now = *pre;
7         now->val++;
8         return;
9     }
10    if(x <= sq[mid]) {
11        build(now->ch[0], pre->ch[0], l, mid, x);
12        now->ch[1] = pre->ch[1];
13        now->maintain();
14    }
15    else {
16        build(now->ch[1], pre->ch[1], mid + 1, r, x);
17        now->ch[0] = pre->ch[0];
18        now->maintain();
19    }
20 }
```

6.10 树剖

```
1 void dfs1(int u, int fa, int dep) {
2     sz[u] = 1; d[u] = dep; ch[u] = 0; p[u] = fa;
3     for(int i = head[u]; i != -1; i = ed[i].next) {
4         int v = ed[i].v;
5         if(v == fa) continue;
6         dfs1(v, u, dep + 1);
7         sz[u] += sz[v];
8         if(sz[v] > sz[ch[u]]) ch[u] = v;
9     }
10 }
11 void dfs2(int u, int rt) {
12     idx[u] = id++;
13     top[u] = rt;
14     if(ch[u]) dfs2(ch[u], rt);
15     for(int i = head[u]; i != -1; i = ed[i].next) {
16         int v = ed[i].v;
17         if(v == p[u] || v == ch[u]) continue;
18         dfs2(v, v);
19     }
20 }
```

```

20 }
21 int ask(int u, int v) {
22     int ret = 0;
23     while(top[u] != top[v]) {
24         //一定是top[u]的深度大于等于top[v]
25         if(d[top[u]] < d[top[v]]) swap(u, v);
26         ret = max(ret, query(1, 1, n, idx[top[u]], idx[u]));
27         u = p[top[u]];
28     }
29     if(d[u] < d[v]) swap(u, v);
30     if(u != v) ret = max(ret, query(1, 1, n, idx[ch[v]], idx[u]));
31     return ret;
32 }

```

6.11 点分治

```

1 void getsz(int u, int fa) {
2     sz[u] = 1; f[u] = 0;
3     for(int i = head[u]; i != -1; i = ed[i].next) {
4         int v = ed[i].v;
5         if(v == fa || vis[v]) continue;
6         getsz(v, u);
7         sz[u] += sz[v];
8         f[u] = max(f[u], sz[v]);
9     }
10 }
11 //找到最大子树最小的点作为分治的中心
12 void getrt(int r, int u, int fa) {
13     //用父边所连的子树更新f
14     f[u] = max(f[u], sz[r] - sz[u]);
15     if(f[u] < minf) {minf = f[u]; rt = u;}
16     for(int i = head[u]; i != -1; i = ed[i].next) {
17         int v = ed[i].v;
18         if(v == fa || vis[v]) continue;
19         getrt(r, v, u);
20     }
21 }
22 int solve(int u) {
23     minf = n;
24     getsz(u, 0);
25     getrt(u, u, 0);
26     vis[rt] = 1;

```

```

27     int ret = getdp(rt); //分治的结果
28     for(int i = head[rt]; i != -1; i = ed[i].next) {
29         int v = ed[i].v;
30         if(!vis[v]) ret = max(ret, solve(v));
31     }
32     return ret;
33 }

```

6.12 RMQ

6.13 整体二分

```

1  Divide_Conquer(Q, AL, AR)
2  //Q是当前处理的操作序列
3  //WANT是要求的贡献, CURRENT为已经累计的贡献(记录的是1~AL
   -1内所有修改的贡献)
4  //[AL, AR]是询问的答案范围区间
5  if AL = AR then
6      将Q中所有是询问操作的答案设为AL
7  end if
8  //我们二分答案, AM为当前的判定答案
9  AM = (AL+AR) / 2
10 //Solve是主处理函数, 只考虑参数满足判定标准[AL, AM]的修改
   的贡献, 因为CURRENT域中已经记录了[1, AL-1]的修改的贡献了
   , 这一步是保证时间复杂度的关键, 因为SOLVE只于当前Q的长度
   有关, 而不与整个操作序列的长度有线性关系, 这保证了主定理
   解出来只多一个log
11 Solve(Q, AL, AM)
12 //Solve之后Q中各个参数满足判定标准的修改对询问的贡献被存
   储在ANS数组
13 //Q1, Q2为了两个临时数组, 用于划分操作序列
14 for i = 1 to Length(Q) do
15     if (Q[i].WANT <= Q[i].CURRENT + ANS[i]) then
16         //当前已有贡献不小于要求贡献, 说明最终答案应当不大
           于判定答案
17         向数组Q1末尾添加Q[i]
18     else
19         //当前已有贡献小于要求贡献, 说明最终答案应当大于判
           定答案
20         //这里是整体二分的关键, 把当前贡献累计入总贡献, 以
           后不再重复统计!

```

```

21         Q[i].CURRENT = Q[i].CURRENT + ANS[i]
22         向数组Q2末尾添加Q[i]
23     end if
24 end for
25 //分治,递归处理
26 Divide_Conquer(Q1, AL, AM)
27 Divide_Conquer(Q2, AM+1, AR)
28 以上别人的伪代码 非常清楚
29
30 带修改的整体二分
31 void cal(int ql, int qr, int l, int mid) {
32     for(int i = ql; i <= qr; i++) {
33         if(nd[i].k) nd[i].cnt = query(nd[i].r) - query(nd
34             [i].l - 1);
35         else if(nd[i].r <= mid) update(nd[i].l, nd[i].cnt
36             );
37     }
38     for(int i = ql; i <= qr; i++) if(nd[i].r <= mid && (!
39         nd[i].k)) update(nd[i].l, -nd[i].cnt);
40 }
41 void divide(int ql, int qr, int l, int r) {
42     if(l == r) {
43         for(int i = ql; i <= qr; i++) if(nd[i].k) ans[nd[
44             i].id] = l;
45         return;
46     }
47     int mid = (l + r) / 2;
48     cal(ql, qr, l, mid);
49     int p1 = 0, p2 = 0;
50     for(int i = ql; i <= qr; i++) {
51         if(nd[i].k) {
52             if(nd[i].cnt >= nd[i].k) t1[++p1] = nd[i];
53             else {
54                 nd[i].k -= nd[i].cnt;
55                 t2[++p2] = nd[i];
56             }
57         }
58         else {
59             if(nd[i].r <= mid) t1[++p1] = nd[i];
60             else t2[++p2] = nd[i];
61         }
62     }
63     for(int i = 1; i <= p1; i++) nd[ql - 1 + i] = t1[i];
64     for(int i = 1; i <= p2; i++) nd[ql + p1 - 1 + i] = t2
65         [i];
66     if(p1) divide(ql, ql + p1 - 1, l, mid);

```

```

62         if(p2) divide(q1 + p1, qr, mid + 1, r);
63     }

```

6.14 莫队

```

1  struct qnode {
2      int l, r, id;
3      ll a, b;
4      bool operator < (const qnode& rhs) const {
5          return dv[l] == dv[rhs.l] ? dv[r] < dv[rhs.r] :
              dv[l] < dv[rhs.l];
6      }
7  }qn[MAXN];
8  int bd = (int)sqrt(n);
9  for(int i = 1; i <=n; i++) {
10     dv[i] = (i - 1) / bd + 1;
11 }

```

6.15 KDtree

给 N 个 K 维点, 找这 K 维点里面最近的 M 个点

```

1  #include <bits/stdc++.h>
2  using namespace std;
3  const int N = 50080;
4  const int K = 5;
5  typedef pair<int, int> Pair;
6  struct Node
7  {
8      int x[K], d;
9      bool f;
10     static int k, cd;
11     void read() {
12         for (int i = 0; i < k; i++) {
13             cin >> x[i];
14         }
15         f = 0;
16     }
17     bool operator < (Node node) const {
18         return x[cd] < node.x[cd];
19     }
20 };
21 int Node::k = 0, Node::cd = 0;
22 int mid(int l, int r)

```

```

23 {
24     return l + (r - 1) / 2;
25 }
26 int dist(Node a, Node b) //算k维的距离
27 {
28     int res = 0;
29     for(int i = 0; i < Node::k; i++)
30         res += (a.x[i] - b.x[i])*(a.x[i] - b.x[i]);
31     return res;
32 }
33 Node a[N];
34 priority_queue<Pair> pq; // 距离 下标
35 void build(int l, int r, int d)
36 {
37     if (l > r) return;
38     int m = mid(l, r);
39     Node::cd = d;
40     nth_element(a+l, a+m, a+r+1); //保证第n大的在第n的位置,
    类似快排的基准
41     //这样就可以保证a[m]是中间的 这样树就可以建的非常平均
42     a[m].d = d; //这个点的维度是d
43     if(l == r) //到了叶节点了
44     {
45         a[m].f = true; //应是不是叶节点的意思
46         return;
47     }
48     build(l, m-1, (d+1)%Node::k); //递归建树
49     build(m+1, r, (d+1)%Node::k);
50 }
51 int num;
52 void fd(int l, int r, Node tar)
53 {
54     if (l > r) return;
55     int m = mid(l, r);
56     int d = dist(a[m], tar);
57     if(a[m].f) { //如果是叶子
58         if(pq.size() < num) { // 还每找满m个
59             pq.push(make_pair(d, m));
60         } else if(d < pq.top().first) { // 已经找了m个了,要
        删了 并且是比当前最远点小
61             pq.pop();
62             pq.push(make_pair(d, m));
63         }
64         return;
65     }

```



```

66     int t = tar.x[a[m].d] - a[m].x[a[m].d];
67     // 要调查的点和 这个节点的距离
68     if(t > 0) { // 右子树
69         fd(m+1, r, tar); //先在右子树
70         if(pq.size() < num) {
71             pq.push(make_pair(d, m));
72             fd(l,m-1,tar); // 再找左子树
73         } else {
74             if(d < pq.top().first) {
75                 pq.pop();
76                 pq.push(make_pair(d, m));
77             }
78             if(pq.top().first > t*t) fd(l,m-1,tar);
79         }
80     } else { //
81         fd(l,m-1,tar);
82         if(pq.size() < num) {
83             pq.push(make_pair(d, m));
84             fd(m+1,r,tar);
85         } else {
86             if(pq.top().first > d) {
87                 pq.pop();
88                 pq.push(make_pair(d, m));
89             }
90             if (pq.top().first > t*t) fd(m+1,r,tar);
91         }
92     }
93 }
94 void solve(int n, int k)
95 {
96     Node::k = k;
97     for (int i = 0; i < n; i++) a[i].read();
98     build(0, n-1, 0);
99     int t;
100    cin >> t;
101    while (t--) {
102        Node q;
103        q.read();
104        while (!pq.empty()) pq.pop();
105        cin >> num; // 题目中需要找与q距离最近的num个点
106        fd(0, n-1, q);
107        vector<int> ans;
108        while (!pq.empty()) {
109            ans.push_back(pq.top().second);
110            pq.pop();
111        }

```

```

112     printf("the closest %d points are:\n", num);
113     for(int j = num - 1; j >= 0; j--) {
114         for(int kk = 0; kk < k; kk++) {
115             kk == 0 ? cout << a[ans[j]].x[kk] : cout << ' '
                        << a[ans[j]].x[kk];
116         }
117         cout << endl;
118     }
119 }
120 }
121 int main()
122 {
123     int n, k;
124     while(cin >> n >> k) solve(n, k);
125 }

```

6.16 优先队列可删除任意元素的堆

6.17 动态树带

6.18 动态树维护后缀自动机 parent 树

6.19 虚树

Chapter 7

字符串

7.1 最小表示法

```
1 #include <cstdio>
2 #include <algorithm>
3 #include <iostream>
4 using namespace std;
5 const int maxn = 100010;
6 int main()
7 {
8     int n;
9     while (scanf("%d", &n) != EOF)
10    {
11        string s[maxn];
12        for (int ii = 0; ii < n; ii++)
13        {
14            string ss;
15            cin >> ss;
16            ss = ss + ss;
17            bool flag = false;
18            int i = 0, j = 1, k = 0, l = ss.size() / 2, p = 0;
19            while (i < l && j < l)
20            {
21                k = 0;
22                while (ss[i + k] == ss[j + k] && k < l) k++;
23                if (k == l)
24                {
25                    p = i;
26                    flag = true;
27                    break;
28                }
29            }
30        }
31    }
```

```

29         if (ss[i + k] > ss[j + k])
30             if (i + k + 1 > j) i = i + k + 1; else i = j +
                1;
31         else if (j + k + 1 > i) j = j + k + 1;
32         else j = i + 1;
33     }
34     if (!flag)
35         if (i < j) p = i; else p = j;
36     s[ii] = ss.substr(p, l);
37 }
38 sort(s, s + n);
39 int ans = 1;
40 for (int i = 1; i < n; i++)
41     if (s[i] != s[i - 1]) ans++;
42 printf("%d\n", ans);
43 }
44 }

```

7.2 KMP

```

1 void getfail() {
2     f[0] = f[1] = 0;
3     for(int i = 1; i < n; i++) {
4         int j = f[i];
5         //找到与当前后缀匹配的最靠右的前缀的位置
6         while(j && c[i] != c[j]) j = f[j];
7         f[i + 1] = c[i] == c[j] ? j + 1 : 0;
8     }
9 }
10 void findp(char *T, char *P, int* f) {
11     int n = strlen(T), m = strlen(P);
12     int j = 0;
13     for(int i = 0; i < n; i++) {
14         //如果不匹配就往前找
15         while(j && P[j] != T[i]) j = f[j];
16         if(P[j] == T[i]) j++;
17         if(j == m) printf("%d\n", i - m + 1);
18     }
19 }

```

7.3 Manacher

```

1 int len = strlen(s);

```

```

2  int n = 1, pre = 0, ans = 0;
3  ss[0] = '$';
4  for(int i = 0; i < len; i++) {
5      ss[n++] = '#';
6      ss[n++] = s[i];
7  }
8  ss[n] = '#';
9  for(int i = 1; i < n; i++) {
10 //pre是在i之前半径延伸地最远的点
11 //如果pre的区间包含了i, 那么i延伸的区间是i向右和i关于pre
    的对称点向左
12 延伸的较小值
13  if(i < pre + p[pre]) p[i] = min(p[pre - (i - pre)], p[
        pre] + pre - i);
14  else p[i] = 1;
15  //继续扩展i的回文区间
16  while(ss[i - p[i]] == ss[i + p[i]]) ++p[i];
17  //用i更新pre
18      if(pre + p[pre] < i + p[i]) pre = i;
19  }

```

7.4 AC 自动机

```

1  void getfail() {
2      queue<int> q;
3      for(int i = 0; i < 26; i++) if(ch[0][i]) q.push(ch
        [0][i]);
4      while(!q.empty()) {
5          int rt = q.front(); q.pop();
6          for(int i = 0; i < 26; i++) {
7              int u = ch[rt][i];
8              if(!u) {ch[rt][i] = ch[fail[rt]][i]; continue
                ;}
9                  //如果本来没有u节点 那么就走失配
                    边 从而形成了图
10             q.push(u);
11             int v = fail[rt];
12             while(v && !ch[v][i]) v = fail[v];
13             fail[u] = ch[v][i];
14                 //Trie中一个节点可能匹配了多个 所
                    以还要记录上一个匹配的后缀
15             last[u] = val[fail[u]] ? fail[u] : last[fail[
                u]];
16         }

```

```

17     }
18 }
19 void print(int j) {
20     if(j) {
21         printf("%d:_%d", j, val[j]);
22         print(last[j]);
23     }
24 }
25 void findp(char * T) {
26     int len = strlen(T);
27     int j = 0;
28     for(int i = 0; i < n; i++) {
29         int x = T[i] - 'a';
30         while(j && !ch[j][x]) j = fail[j];
31         j = ch[j][x];
32         if(val[j]) print(j);
33         else if(last[j]) print(last[j]);
34     }
35 }

```

7.5 后缀数组

```

1  int sa[MAXN], rk[MAXN], ht[MAXN], cnt[MAXN], t1[MAXN], t2
   [MAXN];
2  void getsa(int m) {
3      int *x = t1, *y = t2;
4      for(int i = 0; i < m; i++) cnt[i] = 0;
5      for(int i = 0; i < n; i++) cnt[x[i]] = s[i]++;
6      for(int i = 1; i < m; i++) cnt[i] += cnt[i - 1];
7      for(int i = n - 1; i >= 0; i--) sa[--cnt[x[i]]] = i;
8      //sa数组从0到n-1
9      for(int k = 1; ; k *= 2) {
10         int p = 0;
11         //2nd
12         for(int i = n - k; i < n; i++) y[p++] = i;
13         for(int i = 0; i < n; i++) if(sa[i] >= k) y[p++]
           = sa[i] - k;
14         //1st
15         for(int i = 0; i < m; i++) cnt[i] = 0;
16         for(int i = 0; i < n; i++) cnt[x[y[i]]]++;
17         for(int i = 1; i < m; i++) cnt[i] += cnt[i - 1];
18         for(int i = n - 1; i >= 0; i--) sa[--cnt[x[y[i]
           ]]] = y[i];
19         swap(x, y);

```

```

20     p = 1; x[sa[0]] = 0;
21     for(int i = 1; i < n; i++)
22         x[sa[i]] = (y[sa[i]] == y[sa[i - 1]] && y[sa[
                i] + k] == y[sa[i - 1] + k] && sa[i] + k <
                n && sa[i - 1] + k < n) ? p - 1 : p++;
23     if(p >= n) break;
24     m = p;
25 }
26 }
27 void getheight() {
28     int k = 0;
29     for(int i = 0; i < n; i++) rk[sa[i]] = i;
30     for(int i = 0; i < n; i++) {
31         if(rk[i] == 0) continue;
32         if(k) k--;
33         int j = sa[rk[i] - 1];
34         while(s[i + k] == s[j + k] && i + k < n && j + k
                < n) k++;
35         ht[rk[i]] = k;
36     }
37 }

```

7.6 后缀自动机

```

1 //rt = 1
2 int ncnt, last, ch[MAXN][26], val[MAXN], par[MAXN];
3 int c[MAXN], rk[MAXN];
4 void init(int x, int v) {
5     memset(ch[x], 0, sizeof(ch[x]));
6     par[x] = 0; val[x] = v;
7 }
8 void add(int x) {
9     int p = last, np = ++ncnt;
10    memset(ch[np], 0, sizeof(ch[np]));
11    init(np, val[p] + 1);
12    while(p && !ch[p][x]) {
13        ch[p][x] = np;
14        p = par[p];
15    }
16    if(p == 0) par[np] = 1;
17    else {
18        int q = ch[p][x];
19        if(val[p] + 1 == val[q]) par[np] = q;
20        else {

```

```

21         int nq = ++ncnt;
22         memcpy(ch[nq], ch[q], sizeof(ch[q]));
23         val[nq] = val[p] + 1;
24         par[nq] = par[q];
25         par[q] = nq;
26         par[np] = nq;
27         while(p && ch[p][x] == q) {
28             ch[p][x] = nq;
29             p = par[p];
30         }
31     }
32 }
33 last = np;
34 }
35 void tsort() {
36     memset(c, 0, sizeof(c));
37     for(int i = 1; i <= ncnt; i++) c[val[i]]++;
38     for(int i = 1; i <= ncnt; i++) c[i] += c[i - 1];
39     for(int i = 1; i <= ncnt; i++) rk[c[val[i]]--] = i;
40 }

```

7.7 shift-and

Chapter 8

其他

8.1 蔡勒公式

```
1 int calc(int y,int m,int d)
2 {
3     if (m < 3) m += 12, y--;
4     int w = (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4
5         - y / 100 + y / 400) % 7;
6     return w;
7 }
8 bool isRun(int YYYY)
9 {
10     return (YYYY % 4 == 0 && YYYY % 100 != 0) || (YYYY %
11         400 == 0);
12 }
```

8.2 斜率 DP

8.3 最长子序列

```
1 bool solve()
2 {
3     int n, k;
4     scanf("%d%d", &n, &k);
5     int a[N];
6     int dp[N];
7     int g[N];
```

```

8     rep(i, 1, n) scanf("%d", &a[i]), g[i] = inf;
9     int ans1 = 0;
10    rep(i, 1, n) {
11        int k = upper_bound(g+1, g+1+n, a[i]) - g;
12        dp[i] = k;
13        g[k] = a[i];
14        ans1 = max(ans1, k);
15    }

```

8.4 四边形不等式

```

1  for(int i=2;i<=m+1;++i)
2  {
3      s[i][n+1]=n;
4      for(int j=n;j>i;--j)
5      {
6          for(int k=s[i-1][j];k<=s[i][j+1];++k)
7          {
8              LL tmp=dp[i-1][k]+w[k+1][j];
9              // cout<<i<<' '<<j<<' '<<k<<' '<<tmp<<endl;
10             if(tmp<dp[i][j])
11             {
12                 dp[i][j]=tmp;
13                 s[i][j]=k;
14             }
15         }
16     }
17 }

```

8.5 数位 DP

```

1 //基本上都是这个模板，但是有点慢
2 int dp(int pos, int sum, int lz, int lim) {
3     if(!lim && !lz && d[pos][sum] != -1) return d[pos][sum];
4     if(pos == 0) return sum >= 33;
5     int bd = lim ? s[pos] : 1;
6     int res = 0;
7     for(int i = 0; i <= bd; i++) {
8         if(i == 0) {
9             if(lz) res += dp(pos - 1, sum, lz, lim && i == bd);

```

```

10         else res += dp(pos - 1, sum + 1, 0, lim && i
11             == bd);
12     }
13     else res += dp(pos - 1, sum - 1, 0, lim && i ==
14         bd);
15 }
16 if(!lim && !lz) d[pos][sum] = res;
17 return res;
18 }
19 int solve(int x) {
20     int cnt = 0;
21     while(x) {
22         if(x & 1) s[++cnt] = 1;
23         else s[++cnt] = 0;
24         x /= 2;
25     }
26     return dp(cnt, 33, 1, 1);
27 }

```

8.6 大数

```

1  #include <cstdio>
2  #include <cstring>
3  #include <cstdlib>
4  #include <iostream>
5  #include <algorithm>
6  using namespace std;
7
8  #define MAXN 9999
9  #define MAXSIZE 10
10 #define DLEN 4
11
12 class BigNum
13 {
14 private:
15     int a[2000];    //可以控制大数的位数
16     int len;        //大数长度
17 public:
18     BigNum(){ len = 1;memset(a,0,sizeof(a)); }    //构造函数
19     BigNum(const int);    //将一个int类型的变量转化为
20     BigNum(const char*);    //将一个字符串类型的变量转化
    为大数

```

```

21     BigNum(const BigNum &); //拷贝构造函数
22     BigNum &operator=(const BigNum &); //重载赋值运算
        符，大数之间进行赋值运算
23
24     friend istream& operator>>(istream&, BigNum&); //
        重载输入运算符
25     friend ostream& operator<<(ostream&, BigNum&); //
        重载输出运算符
26
27     BigNum operator+(const BigNum &) const; //重载加法
        运算符，两个大数之间的相加运算
28     BigNum operator-(const BigNum &) const; //重载减法
        运算符，两个大数之间的相减运算
29     BigNum operator*(const BigNum &) const; //重载乘法
        运算符，两个大数之间的相乘运算
30     BigNum operator/(const int &) const; //重载除法
        运算符，大数对一个整数进行相除运算
31
32     BigNum operator^(const int &) const; //大数的n次
        方运算
33     int operator%(const int &) const; //大数对一个
        int类型的变量进行取模运算
34     bool operator>(const BigNum & T) const; //大数和另
        一个大数的大小比较
35     bool operator<(const BigNum & T) const;
36     bool operator==(const BigNum & T) const;
37     bool operator>(const int & t) const; //大数和一
        个int类型的变量的大小比较
38     bool operator<(const int & t) const;
39     bool operator==(const int & t) const;
40
41     void print(); //输出大数
42 };
43
44 bool BigNum::operator==(const BigNum & T) const {
45     return !(*this > T) && !(T > *this);
46 }
47 bool BigNum::operator==(const int & t) const {
48     BigNum T = BigNum(t);
49     return *this == T;
50 }
51 bool BigNum::operator<(const BigNum & T) const {
52     return T > *this;
53 }
54 bool BigNum::operator<(const int & t) const {
55     return BigNum(t) > *this;

```

```

56 }
57 BigNum::BigNum(const int b)      //将一个int类型的变量转化
    为大数
58 {
59     int c,d = b;
60     len = 0;
61     memset(a,0,sizeof(a));
62     while(d > MAXN)
63     {
64         c = d - (d / (MAXN + 1)) * (MAXN + 1);
65         d = d / (MAXN + 1);
66         a[len++] = c;
67     }
68     a[len++] = d;
69 }
70 BigNum::BigNum(const char*s)    //将一个字符串类型的变量
    转化为大数
71 {
72     int t,k,index,l,i;
73     memset(a,0,sizeof(a));
74     l=strlen(s);
75     len=l/DLEN;
76     if(l%DLEN)
77         len++;
78     index=0;
79     for(i=l-1;i>=0;i-=DLEN)
80     {
81         t=0;
82         k=i-DLEN+1;
83         if(k<0)
84             k=0;
85         for(int j=k;j<=i;j++)
86             t=t*10+s[j]-'0';
87         a[index++]=t;
88     }
89 }
90 BigNum::BigNum(const BigNum & T) : len(T.len) //拷贝构造
    函数
91 {
92     int i;
93     memset(a,0,sizeof(a));
94     for(i = 0 ; i < len ; i++)
95         a[i] = T.a[i];
96 }

```

```

97  BigNum & BigNum::operator=(const BigNum & n)    //重载赋值
    运算符，大数之间进行赋值运算
98  {
99      int i;
100     len = n.len;
101     memset(a,0,sizeof(a));
102     for(i = 0 ; i < len ; i++)
103         a[i] = n.a[i];
104     return *this;
105 }
106 ostream& operator>>(ostream & in,  BigNum & b)    //重载输
    入运算符
107 {
108     char ch[MAXSIZE*4];
109     int i = -1;
110     in>>ch;
111     int l=strlen(ch);
112     int count=0,sum=0;
113     for(i=l-1;i>=0;)
114     {
115         sum = 0;
116         int t=1;
117         for(int j=0;j<4&& i>=0;j++,i--,t*=10)
118         {
119             sum+=(ch[i]-'0')*t;
120         }
121         b.a[count]=sum;
122         count++;
123     }
124     b.len =count++;
125     return in;
126 }
127 }
128 ostream& operator<<(ostream& out,  BigNum& b)    //重载输
    出运算符
129 {
130     int i;
131     cout << b.a[b.len - 1];
132     for(i = b.len - 2 ; i >= 0 ; i--)
133     {
134         cout.width(DLEN);
135         cout.fill('0');
136         cout << b.a[i];
137     }
138     return out;
139 }

```

```

140
141 BigNum BigNum::operator+(const BigNum & T) const //两个
    大数之间的相加运算
142 {
143     BigNum t(*this);
144     int i, big; //位数
145     big = T.len > len ? T.len : len;
146     for(i = 0 ; i < big ; i++)
147     {
148         t.a[i] += T.a[i];
149         if(t.a[i] > MAXN)
150         {
151             t.a[i + 1]++;
152             t.a[i] -= MAXN + 1;
153         }
154     }
155     if(t.a[big] != 0)
156         t.len = big + 1;
157     else
158         t.len = big;
159     return t;
160 }
161 BigNum BigNum::operator-(const BigNum & T) const //两个
    大数之间的相减运算
162 {
163     int i, j, big;
164     bool flag;
165     BigNum t1, t2;
166     if(*this > T)
167     {
168         t1 = *this;
169         t2 = T;
170         flag = 0;
171     }
172     else
173     {
174         t1 = T;
175         t2 = *this;
176         flag = 1;
177     }
178     big = t1.len;
179     for(i = 0 ; i < big ; i++)
180     {
181         if(t1.a[i] < t2.a[i])
182         {
183             j = i + 1;

```

```

184         while(t1.a[j] == 0)
185             j++;
186         t1.a[j--]--;
187         while(j > i)
188             t1.a[j--] += MAXN;
189         t1.a[i] += MAXN + 1 - t2.a[i];
190     }
191     else
192         t1.a[i] -= t2.a[i];
193 }
194 t1.len = big;
195 while(t1.a[t1.len - 1] == 0 && t1.len > 1)
196 {
197     t1.len--;
198     big--;
199 }
200 if(flag)
201     t1.a[big-1] = 0 - t1.a[big-1];
202 return t1;
203 }
204
205 BigNum BigNum::operator*(const BigNum & T) const    //两个
大数之间的相乘运算
206 {
207     BigNum ret;
208     int i,j,up;
209     int temp,temp1;
210     for(i = 0 ; i < len ; i++)
211     {
212         up = 0;
213         for(j = 0 ; j < T.len ; j++)
214         {
215             temp = a[i] * T.a[j] + ret.a[i + j] + up;
216             if(temp > MAXN)
217             {
218                 temp1 = temp - temp / (MAXN + 1) * (MAXN
+ 1);
219                 up = temp / (MAXN + 1);
220                 ret.a[i + j] = temp1;
221             }
222             else
223             {
224                 up = 0;
225                 ret.a[i + j] = temp;
226             }
227         }

```



```

228         if(up != 0)
229             ret.a[i + j] = up;
230     }
231     ret.len = i + j;
232     while(ret.a[ret.len - 1] == 0 && ret.len > 1)
233         ret.len--;
234     return ret;
235 }
236 BigNum BigNum::operator/(const int & b) const    //大数对
        一个整数进行相除运算
237 {
238     BigNum ret;
239     int i,down = 0;
240     for(i = len - 1 ; i >= 0 ; i--)
241     {
242         ret.a[i] = (a[i] + down * (MAXN + 1)) / b;
243         down = a[i] + down * (MAXN + 1) - ret.a[i] * b;
244     }
245     ret.len = len;
246     while(ret.a[ret.len - 1] == 0 && ret.len > 1)
247         ret.len--;
248     return ret;
249 }
250 int BigNum::operator %(const int & b) const    //大数对一
        个int类型的变量进行取模运算
251 {
252     int i,d=0;
253     for (i = len-1; i>=0; i--)
254     {
255         d = ((d * (MAXN+1))% b + a[i])% b;
256     }
257     return d;
258 }
259 BigNum BigNum::operator^(const int & n) const    //大数的
        n次方运算
260 {
261     BigNum t,ret(1);
262     int i;
263     if(n<0)
264         exit(-1);
265     if(n==0)
266         return 1;
267     if(n==1)
268         return *this;
269     int m=n;
270     while(m>1)

```

```

271     {
272         t=*this;
273         for( i=1;i<<1<=m;i<=1)
274         {
275             t=t*t;
276         }
277         m-=i;
278         ret=ret*t;
279         if(m==1)
280             ret=ret*(*this);
281     }
282     return ret;
283 }
284 bool BigNum::operator>(const BigNum & T) const    //大数和
        另一个大数的大小比较
285 {
286     int ln;
287     if(len > T.len)
288         return true;
289     else if(len == T.len)
290     {
291         ln = len - 1;
292         while(a[ln] == T.a[ln] && ln >= 0)
293             ln--;
294         if(ln >= 0 && a[ln] > T.a[ln])
295             return true;
296         else
297             return false;
298     }
299     else
300         return false;
301 }
302 bool BigNum::operator >(const int & t) const    //大数和
        一个int类型的变量的大小比较
303 {
304     BigNum b(t);
305     return *this>b;
306 }
307
308 void BigNum::print()    //输出大数
309 {
310     int i;
311     printf("%d", a[len-1]);
312     for (int i = len-2; i >= 0; --i) {
313         printf("%04d", a[i]);
314     }

```

```

315     puts("");
316 }

```

8.7 可以重复走的异或路径

```

1  #include <bits/stdc++.h>
2  using namespace std;
3  typedef long long ll;
4  const int MAXN = 1e5 + 10;
5  vector<pair<int, ll> > g[MAXN];
6  vector<ll> a;
7  int vis[MAXN];
8  ll dis[MAXN];
9  ll b[60];
10 int n, m;
11 void dfs(int u, ll d) {
12     vis[u] = 1;
13     dis[u] = d;
14     for(int i = 0; i < g[u].size(); i++) {
15         int v = g[u][i].first;
16         if(vis[v]) a.push_back(d^g[u][i].second^dis[v]);
17         else dfs(v, dis[u]^g[u][i].second);
18     }
19 }
20 int main() {
21     cin >> n >> m;
22     for(int i = 1; i <= m; i++) {
23         int u, v;
24         ll w;
25         scanf("%d%d%lld", &u, &v, &w);
26         g[u].push_back(make_pair(v, w));
27         g[v].push_back(make_pair(u, w));
28     }
29     dfs(1, 0);
30     for(int i = 0; i < a.size(); i++) {
31         for(int j = 40; j >= 0; j--) {
32             if(a[i] & (1LL << j)) {
33                 if(!b[j]) {b[j] = a[i]; break;}
34                 else a[i]^=b[j];
35             }
36         }
37     }
38     ll ans = dis[n];

```

```
39     for(int i = 40; i >= 0; i--) ans = min(ans, ans^b[i])  
        ;  
40     cout << ans << endl;  
41 }
```

8.8 最简分数逼近