

# L4: OOP, First Class Functions, Continuation and Closure

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# In-class Exercise 6

- `def unzip(xs: List[(Int,Int)]): (List[Int], List[Int])` reverses what `zip` does. Make it so that it's polymorphic. The input can be any `List[(A, B)]`.
- `def countWhile[T](xs: List[T], key: T): Int` that counts the number of times `key` repeats itself in the prefix of `xs`.
- `def topK(xs: List[Int], k: Int): List[Int]` that tallies the elements of `xs` and return elements with the top `k` frequencies (if there are ties, break ties in any way you like). Look at Scaladoc for inspiration.
- Make a sum type called `Dessert`, which can be one of the following: – `Pie(kind: String)`, – `Smoothie(fruits: List[String])`, – `Cake(toppings: String)`
- Then, write a function `def isLiquid(what: Dessert): Boolean`

# Let's Create a Rational Number

- Idea 1: Use a pair

- `type Rational = (Int, Int)`
- `def add(p: Rational, q: Rational) = (p, q) match {  
 case ((np, dp), (nq, dq)) => (np*dq + nq*dp, dp*dq)  
}`
- `def toString(p: Rational) = (p, q) match {  
 p.toString + "/" + q.toString  
}`

- Idea 2: Use a record

- `case class Rational(n: Int, d: Int)`
- `def add(p: Rational, q: Rational) = Rational(p.n*q.d + q.n*p.d,  
 p.d*q.d)`
- `def toString(p: Rational) = p.n.toString + "/" + p.d.toString`

# Using the Class

- Let's use a class to define a rational number
- ```
class Rational(n: Int, d: Int) {  
    def numer = n  
    def denom = d  
}
```
- Instantiation in Scala
  - `new`
- `val r = new Rational(3, 4)`
- Accessing `r` can be done by using *r.numer* and *r.denom*

# Implementing an Add

- This can be done so it becomes a class method
  - `def add(that: Rational) = new Rational(this.numer*that.denom + that.numer*this.denom, this.denom*that.denom)`
- `def mult(that: Rational) = new Rational(numer*that.numer, denom*that.denom)`
- `override def toString = numer + "/" + denom`

# Public and Private

- ```
private def gcd(a: Int, b: Int): Int =  
  if (b == 0) a else gcd(b, a % b)  
private val g = gcd(n, d)  
def numer = n/g  
def denom = d/g  
require(d>0, "denominator must be positive")
```
- By default, def is public

# Constructor

- We can define a constructor by adding aux. constructors
- `def this(n:Int) = this(n, 1)`
- Notice how we use “this” to self-reference
- Example:
- `def less(that: Rational) =  
 numer * that.denom < that.numer * denom`
- `// This could have been: this.numer * that.denom <  
 that.numer * this.denom`

# Overloading Operators

- Unlike an Integer, adding a rational class is different
  - You cannot just call `x+y`
  - You ended up having to define `r.add`
- Alternative: overloading the “+” operator
  - Operators are treated like a function, you can define it
- `def + (r: Rational) = ...`
- `def - (r: Rational) = ...`
- `def unary_- = ...`
- Keep in mind that operators have precedence rule
  - Overloaded operators keep the same rule



# Abstract Class

- What if we want to make an abstract class?
  - Q: What is an abstract class?
- Let's say we want to create the following things:
- An IntSet, where it collects a set of Integers
  - `add(x: Int): IntSet` – produce a new set taking the union of this set and `{x}`.
  - `has(x: Int): Boolean` — ask if `x` is a member of this set
- How can we specify the interface?

# Abstract Class: Interface

- Create an abstract class
- abstract class IntSet {  
    def add(x: Int): IntSet  
    def has(x: Int): Boolean  
}
- Then, we can implement this class later

# Abstract Class: Implementation

- Example: Implement this IntSet using a linked list
- class Empty extends IntSet {  
    def has(x: Int) = false  
    def add(x: Int) = ... // new NonEmpty(x, new Empty)  
}
- class NonEmpty(elt: Int, other: IntSet) extends IntSet {  
    def has(x: Int) = (x==elt) || (other has x)  
    def add(x: Int) = new NonEmpty(x, this)  
}

# Abstract Class: Implementation

- Empty and NonEmpty extend the class IntSet
- Both conforms to IntSet
- IntSet is the superclass to Empty and NonEmpty
  - Vice versa, Empty and NonEmpty are the subclasses
- Everything has an Object as a superclass
  - This includes your REPL statements
  - This means you can override the implementation
    - val on top on existing variables

# Limit Copies to One

- From our example, the Empty set should really have one copy, right?
  - This is pretty easy to fix using static class in other languages
- In Scala, this is also an easy fix using a singleton ***object***
- ```
object Empty extends IntSet {  
    def has(x: Int) = false  
    def add(x: Int) = new NonEmpty(x, Empty)  
}
```
- This define an object called Empty, no other instances of this object can be created
- Empty evaluate to itself (it is a value)

# In-Class Exercise 7

- Recreate an object for the Expression type with we been using, with the following traits
  - trait Expr {
    - def +(that: Expr) = [Fill in this blank]
    - def \*(that: Expr) = [Fill in this blank]
    - def unary\_- = [Fill in this blank]
    - def toVal(implicit ctx: Map[String, Double]): Double
  - }
- It should have the following methods
  - case class Var(name: String) extends Expr
  - case class Constant(n: Double) extends Expr
  - case class Negate(e: Expr) extends Expr
  - case class Sum(e1: Expr, e2: Expr) extends Expr
  - case class Prod(e1: Expr, e2: Expr) extends Expr
  - Each of these should implement its version of **toVal**

# First-class Functions

# Recap: First Class Function

- Functions become values
- Conceptually, this allows you to pass functions in, and return a function
- Example: Repeat a function n times
  - `def nTimes[A](f: A => A, n: Int, x: A): A =  
 if (n==0) x else f(nTimes(f, n-1, x))`



# Examples: Functions as Inputs

- Let's define:  
def triple(x: Int) = 3\*x  
def addTwo(x: Int) = x+2  
def doTail[T](xs: List[T]) = xs.tail
- What does these do?  
nTimes(triple, 7, 11)  
nTimes(addTwo, 4, 9)  
nTimes(doTail, 2, List(3,5,2,4,9,7))  
nTimes(doTail[Int], 2, List(3,5,2,4,9,7))

# Examples: Functions as Outputs

- ```
def tripleNTimes(n: Int, x: Int) = {  
  def triple(x: Int) = 3*x  
  nTimes(triple, n, x)  
}
```
- ```
// use the shorthand form for defining a function  
def tripleNTimes(n: Int, x: Int) =  
  nTimes((x: Int) => 3*x, n, x)
```

# Scala: Methods vs. Functions

- When we write `def inc(x: Int) = x+1`
  - This is not really a function
  - `def` with parameters is a method
- In Scala, method can be polymorphic
- Also in Scala, functions are never polymorphic
  - They will have a type
- `inc _` gives a functional form, it takes an `Int`, and will return an `Int`

# Types

- For now, let's assume functions are polymorphic
- `def nTimes[A](f: A => A, n: Int, x: A): A =  
 if (n==0) x else f(nTimes(f, n-1, x))`
  - This has the type `((A => A), Int, A) => A`
    - What does this mean?
- In this same example, A is a placeholder for a type
- But, these functions does not have to be polymorphic
  - `def timesUntilZero(f: Int => Int, x: Int): Int =  
 if (x==0) 0 else 1 + timesUntilZero(f, f(x))`

# Reducing the Function

- Consider this example
  - if  $((x*y+2 < 10) == \text{true})$  true else false
- Rewrite once
  - if  $(x*y+2 < 10)$  true else false
- Rewrite again to
  - $(x*y+2 < 10)$

# Reducing the Function: Example 2

- Can I rewrite the following?
  - `nTimes(doTail[Int], 2, List(3,2,1))`
- `nTimes((xs: List[Int]) => xs.tail, 2, List(3,2,1))`
- `nTimes[List[Int]](_.tail, 2, List(3,2,1))`

# More Abstraction

- Consider this example
  - ```
def sillyLottery(f: Int => Int, n: Int) =  
  if (f(n)%2 == 0) {  
    (x: Int) => x/2  
  } else {  
    (x: Int) => 2*x+1 }
```
- What is the type?
  - ```
((Int => Int), Int) => (Int => Int)
```
  - If we give `Int => Int` and one `Int`, we will get `Int => Int`
  - Which we can bind to a variable
- Let's consider `val magic = sillyLottery(x=>3*x-9, 25)`
- What is `magic(21)`?

# Scala with I/O

- You can import `scala.io.Source` to deal with I/O



# Example

- What if I want to count the number of word in a file?

```
import scala.io.Source
```

```
object SimpleWordCount extends App {
```

```
  def countPerLine(line: String): Int =
```

```
    line.split("\\W+") .length
```

```
  val wordsPerLine =
```

```
    Source.stdin .getLines.map(countPerLine).toSeq
```

```
  val lineCount = wordsPerLine.length
```

```
  val wordCount = wordsPerLine.sum
```

```
  println(s"lineCount: $lineCount")
```

```
  println(s"wordCount: $wordCount")
```

```
}
```

# Continuation

# Continuations

- So far, all our functions return something
- You can also call a new function at the end
  - This is called the continuation passing style (CPS)
- This allows you to make every function a tail call
- Let's use a sum of all integer as an example

```
def sum(L: List[Int]): Int = L match {  
  case Nil => 0 case x::xs => sum(xs) + x
```

# Continuations

- Tail call version:
- `sum_tc(L: List[Int]): Int = {  
 def sumHelper(L: List[Int], a: Int): Int = L match {  
 case Nil => a  
 case x::xs => sumHelper(xs, a + x) }  
 sumHelper(L, 0) }`

# Continuations

- Continuation version
- ```
def sum_cont(L: List[Int]): Int = {  
  def sumHelper(L: List[Int], K: Int => Int): Int = L match {  
    case Nil => K(0)  
    case x::xs => sumHelper(xs, (r: Int) => K(r + x)) }  
  sumHelper(L, (x: Int) => x) }
```

# Example 2: Binary Tree

- Fundamentally, a binary tree can be defined recursively
  - A node can be
    - Empty
    - A node with two sub-tree
- Let's make the trait Tree
- sealed trait Tree
  - case object Empty extends Tree
  - case class Node(left: Tree, key: Int, right: Tree) extends Tree

# Example 2: Binary Tree

- Making a tree can be done by
- `val t = Node(Node(Node(Empty, 2, Empty), 5, Node(Node(Empty, 6, Empty), 7, Node(Empty, 9, Empty)))`
- What does this tree looks like?

# Example 2: Binary Tree Traversal

- ```
def walkInorder(t: Tree): List[Int] = t match {  
  case Empty => Nil  
  case Node(l, k, r) => walkInorder(l)::(k::walkInorder(r))  
}
```
- What if we want to use continuation?
  - We need to pass down the functions to call at the end
  - What should that function do?



# Example 2: Binary Tree Traversal

- ```
def walkInorder(t: Tree): List[Int] = {  
  def contWalk(t: Tree, K: List[Int] => List[Int]): List[Int] = t  
  match {  
    case Empty => K(Nil)  
    case Node(l, k, r) => contWalk(l, leftList => {  
      contWalk(r, rightList => K(leftList:::(k::rightList)))  
    })  
  }  
  contWalk(t, (r: List[Int]) => r)  
}
```

# Currying

# Currying

- Instead of accepting parameters normally, accept through a sequence of functions
- `def sortedUncurry(x: Int, y: Int, z: Int) = x <= y && y <= z`
- `val sortedTriple = { (x: Int) => (y: Int) => (z: Int) => x <= y && y <= z }`
- Currying version:
  - `def sortedTriple (x: Int) (y: Int) (z: Int) = x <= y && y <= z`

# Currying

- Benefits:
  - You can stage the function
    - Parts of the execution can run as soon as the values are ready
  - Maps well with dataflow model
  - This can allow the compiler and the hardware to be faster
    - Eliminate data dependency as soon as possible
- Actual efficiency: It depends
  - Compiler is very smart nowadays
  - Run a profiler to test the two formats

# States and Mutable Variables

- We assumed variables are immutable
  - This is annoying in some cases
  - What if we need to store a state
- **State:** the intermediate steps that need to be stored
  - Real hardware also needs the concept of state
- So, many functional languages have mutable variables
- Benefit of mutable variables
  - Allow programmer to keep the state

# Declaring Mutable Variables

- `var x = value`
- Example: implementing a while loop
  - Using currying and mutable variables
- ```
def my_while (condition: => Boolean) (block: => Unit):  
  Unit = {  
    if (condition) {  
      block  
      my_while (condition) (block)  
    } else ()  
  }
```

**Before We Leave Today**

# In-class Exercise 8

- Implement fibonacci recursively in the continuation-passing style
- Preorder traversal: visit root first, then left, then right. Write preorderWalk in CPS style