



Machine-Level Programming IV: Data - arrays

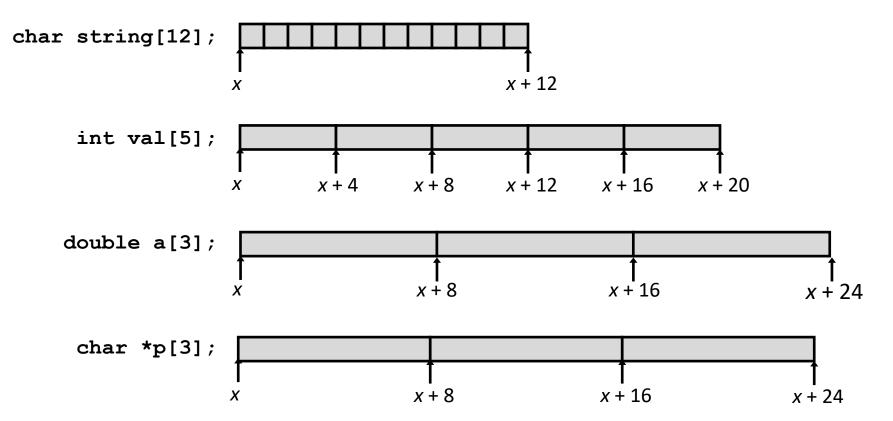
These slides adapted from materials provided by the textbook authors.

Array Allocation

Basic Principle

```
T A[L];
```

- Array of data type T and length L
- Contiguously allocated region of L * sizeof (T) bytes in memory



Array Access

Basic Principle

```
T A[L];
```

- Array of data type T and length L
- Identifier A can be used as a pointer to array element 0: Type T*

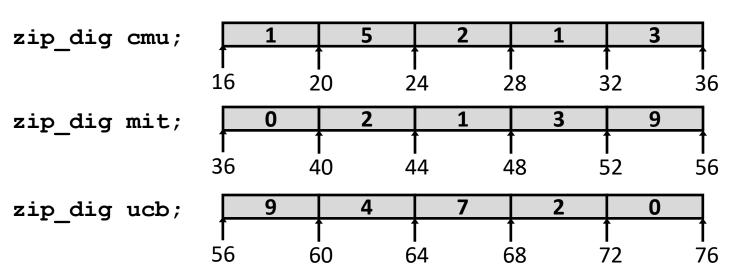
<pre>int val[5];</pre>	1	5	2	1	3	
	Î '	1		1		1
	x x -	+ 4 x -	+ 8 x +	- 12 x +	- 16 <i>x</i> +	- 20

Reference	Type	Value	
val[4]	int	3	
val	int *	X	
val+1	int *	x + 4	
&val[2]	int *	<i>x</i> + 8	
val [5]	int	??	
*(val+1)	int	5	
val + <i>i</i>	int *	x + 4i	

Array Example

```
#define ZLEN 5
typedef int zip_dig[ZLEN];

zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucb = { 9, 4, 7, 2, 0 };
```



- Declaration "zip dig cmu" equivalent to "int cmu[5]"
- Example arrays were allocated in successive 20 byte blocks
 - Not guaranteed to happen in general

Array Accessing Example

```
int get_digit
  (zip_dig z, int digit)
{
  return z[digit];
}
```

IA32

```
# %rdi = z
# %rsi = digit
movl (%rdi, %rsi, 4), %eax # z[digit]
```

- Register %rdi contains starting address of array
- Register %rsi contains array index
- Desired digit at %rdi + 4*%rsi
- Use memory reference
 (%rdi,%rsi,4)

Array Loop Example

```
void zincr(zip_dig z) {
   size_t i;
   for (i = 0; i < ZLEN; i++)
      z[i]++;
}</pre>
```

```
# %rdi = z
                        \# i = 0
 movl $0, %eax
                         # goto middle
 jmp .L3
.L4:
                         # loop:
 addl $1, (%rdi,%rax,4) # z[i]++
 addq $1, %rax
                         # i++
.L3:
                         # middle
 cmpq $4, %rax
                         # i:4
                         # if <=, goto loop</pre>
 jbe .L4
 rep; ret
```

Multidimensional (Nested) Arrays

Declaration

 $T \mathbf{A}[R][C];$

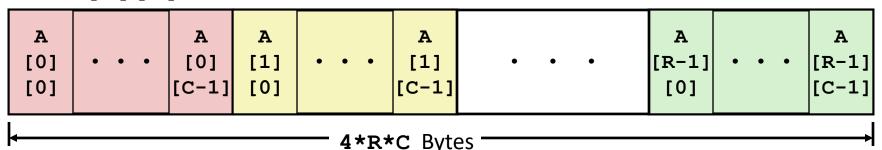
- 2D array of data type T
- R rows, C columns
- Type T element requires K bytes

Array Size

R * C * K bytes

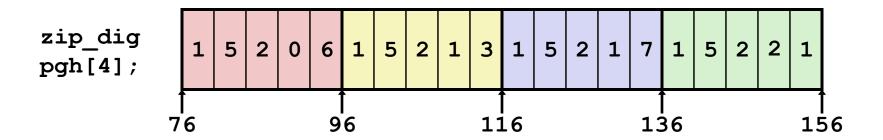
Arrangement

Row-Major Ordering – Going from row to row is Major undertaking int A[R][C];



Nested Array Example

```
#define PCOUNT 4
zip_dig pgh[PCOUNT] =
  {{1, 5, 2, 0, 6},
   {1, 5, 2, 1, 3},
   {1, 5, 2, 1, 7},
   {1, 5, 2, 2, 1 }};
```



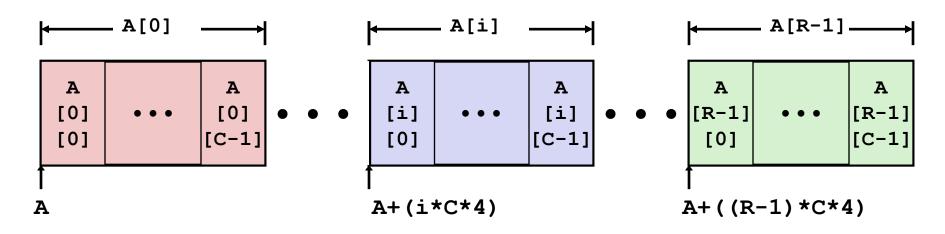
- "zip_dig pgh[4]" equivalent to "int pgh[4][5]"
 - Variable pgh: array of 4 elements, allocated contiguously
 - Each element is an array of 5 int's, allocated contiguously
- "Row-Major" ordering of all elements in memory

Nested Array Row Access

Row Vectors

- A[i] is array of C elements
- Each element of type T requires K bytes
- Starting address A + i * (C * K)

int A[R][C];



Nested Array Row Access Code

```
1 5 2 0 6 1 5 2 1 3 1 5 2 1 7 1 5 2 2 1

pgh

int *get_pgh_zip(int index)
{
    return pgh[index];
}

# %rdi = index
leag (%rdi,%rdi,4),%rax # 5 * index
```

```
leaq pgh(,%rax,4),%rax # pgh + (20 * index)
```

Row Vector

- pgh[index] is array of 5 int's
- Starting address pgh+20*index

Machine Code

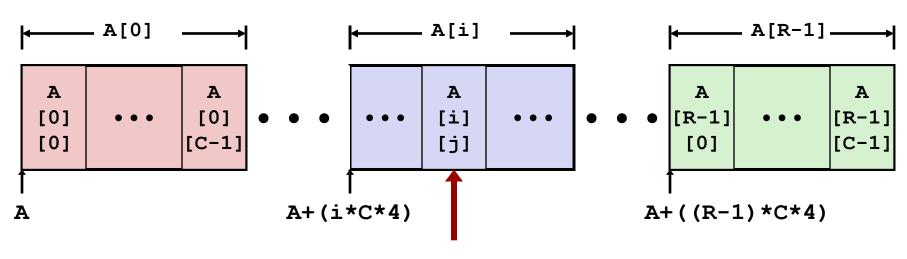
- Computes and returns address
- Compute as pgh + 4* (index+4*index)

Nested Array Element Access

Array Elements

- **A[i][j]** is element of type *T*, which requires *K* bytes
- Address **A** + i * (C * K) + j * K = A + (i * C + j) * K

int A[R][C];



A+(i*C*4)+(j*4)

Nested Array Element Access Code

```
1 5 2 0 6 1 5 2 1 3 1 5 2 1 7 1 5 2 2 1

pgh

int get_pgh_digit
    (int index, int dig)
{
    return pgh[index][dig];
}
```

```
leaq (%rdi,%rdi,4), %rax # 5*index
addl %rax, %rsi # 5*index+dig
movl pgh(,%rsi,4), %eax # M[pgh + 4*(5*index+dig)]
```

Array Elements

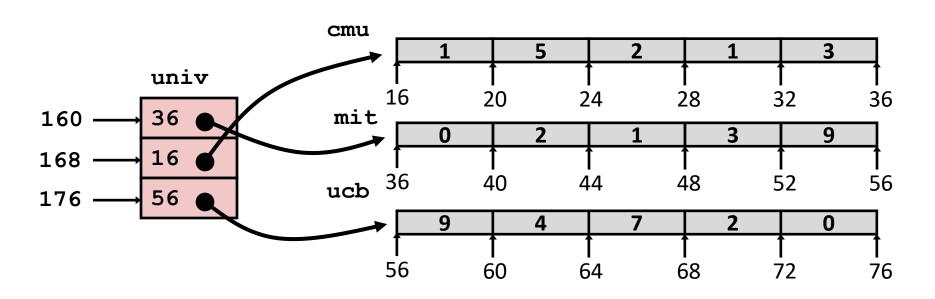
- pgh[index][dig] is int
- Address: pgh + 20*index + 4*dig
 - = pgh + 4*(5*index + dig)

Multi-Level Array Example

```
zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucb = { 9, 4, 7, 2, 0 };
```

```
#define UCOUNT 3
int *univ[UCOUNT] = {mit, cmu, ucb};
```

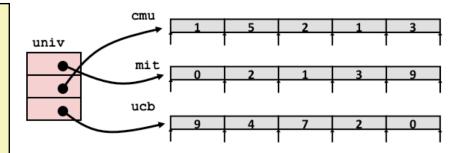
- Variable univ denotes array of 3 elements
- Each element is a pointer
 - 8 bytes
- Each pointer points to array of int's



13

Element Access in Multi-Level Array

```
int get_univ_digit
  (size_t index, size_t digit)
{
  return univ[index][digit];
}
```



```
salq $2, %rsi # 4*digit
addq univ(,%rdi,8), %rsi # p = univ[index] + 4*digit
movl (%rsi), %eax # return *p
ret
```

Computation

- Element access Mem [Mem [univ+8*index]+4*digit]
- Must do two memory reads
 - First get pointer to row array
 - Then access element within array

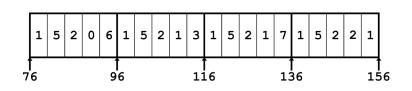
Array Element Accesses

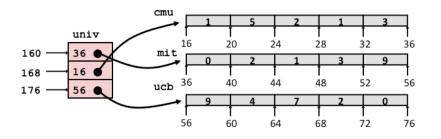
Nested array

```
int get_pgh_digit
   (size_t index, size_t digit)
{
   return pgh[index][digit];
}
```

Multi-level array

```
int get_univ_digit
   (size_t index, size_t digit)
{
   return univ[index][digit];
}
```





Accesses looks similar in C, but address computations very different:

Mem[pgh+20*index+4*digit] Mem[Mem[univ+8*index]+4*digit]

N X N Matrix Code

Fixed dimensions

 Know value of N at compile time

Variable dimensions, explicit indexing

 Traditional way to implement dynamic arrays

Variable dimensions, implicit indexing

Now supported by gcc

16 X 16 Matrix Access

Array Elements

- Address **A** + i * (C * K) + j * K
- C = 16, K = 4

```
/* Get element a[i][j] */
int fix_ele(fix_matrix a, size_t i, size_t j) {
  return a[i][j];
}
```

```
# a in %rdi, i in %rsi, j in %rdx
salq $6, %rsi  # 64*i
addq %rsi, %rdi  # a + 64*i
movl (%rdi,%rdx,4), %eax # M[a + 64*i + 4*j]
ret
```

n X n Matrix Access

Array Elements

- Address **A** + i * (C * K) + j * K
- C = n, K = 4
- Must perform integer multiplication

```
/* Get element a[i][j] */
int var_ele(size_t n, int a[n][n], size_t i, size_t j)
{
   return a[i][j];
}
```

```
# n in %rdi, a in %rsi, i in %rdx, j in %rcx
imulq %rdx, %rdi  # n*i
leaq (%rsi,%rdi,4), %rax # a + 4*n*i
movl (%rax,%rcx,4), %eax # a + 4*n*i + 4*j
ret
```