



These slides adapted from materials provided by the textbook

Mechanisms in Procedures

Passing control

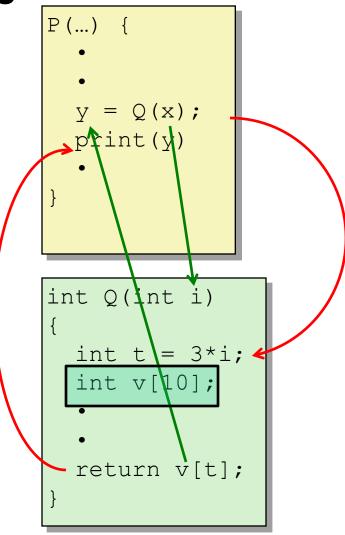
- To beginning of procedure code
- Back to return point

Passing data

- Procedure arguments
- Return value

Memory management

- Allocate during procedure execution
- Deallocate upon return
- Mechanisms all implemented with machine instructions
- x86-64 implementation of a procedure uses only those mechanisms required

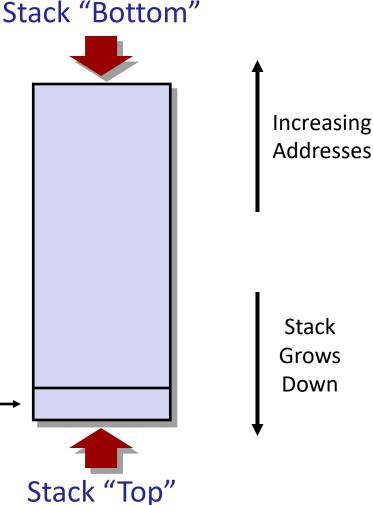


- Procedures
 - Stack Structure
 - Calling Conventions
 - Passing control
 - Passing data
 - Managing local data

x86-64 Stack

- Region of memory managed with stack discipline
- Grows toward lower addresses
- Register %rsp contains lowest stack address
 - address of "top" element

Stack Pointer: %rsp → Stack "Tou



x86-64 Stack: Push

pushq Src

- Fetch operand at Src
- Decrement %rsp by 8
- Write operand at address given by %rsp

Addresses Stack **Grows** Down Stack Pointer: %rsp Stack "Top"

Stack "Bottom"

Increasing

x86-64 Stack: Pop

■ popq Dest

- Read value at address given by %rsp
- Increment %rsp by 8
- Store value at Dest (must be register)

Addresses Stack **Grows** Down Stack Pointer: %rsp Stack "Top"

Stack "Bottom"

Increasing

- Procedures
 - Stack Structure
 - Calling Conventions
 - Passing control
 - Passing data
 - Managing local data

Code Examples

```
void multstore
  (long x, long y, long *dest)
{
    long t = mult2(x, y);
    *dest = t;
}
```

```
      0000000000000400540
      <multstore>:

      400540: push %rbx # Save %rbx

      400541: mov %rdx,%rbx # Save dest

      400544: callq 400550 <mult2> # mult2(x,y)

      400549: mov %rax,(%rbx) # Save at dest

      40054c: pop %rbx # Restore %rbx

      40054d: retq # Return
```

```
long mult2
  (long a, long b)
{
  long s = a * b;
  return s;
}
```

```
0000000000400550 <mult2>:
   400550: mov %rdi,%rax # a
   400553: imul %rsi,%rax # a * b
   400557: retq # Return
```

Procedure Control Flow

- Use stack to support procedure call and return
- Procedure call: call label
 - Push return address on stack
 - Jump to label
- Return address:
 - Address of the next instruction right after call
 - Example from disassembly
- Procedure return: ret
 - Pop address from stack
 - Jump to address

Control Flow Example #1

```
0x130

0x128

0x120

0x128

0x120

0x120

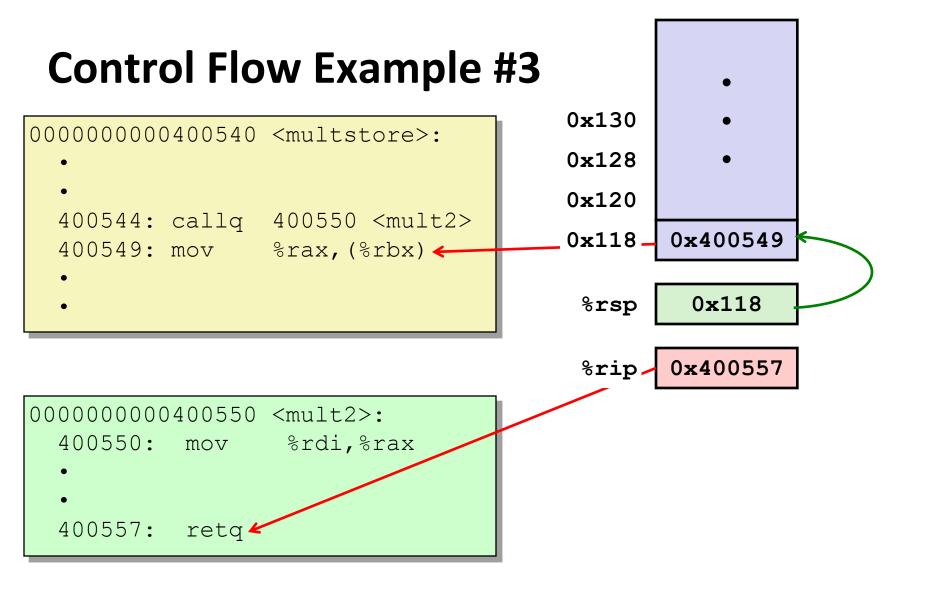
0x120

0x120

0x120
```

```
000000000400550 <mult2>:
   400550: mov %rdi,%rax
   •
   400557: retq
```

Control Flow Example #2 0x1300000000000400540 <multstore>: 0x1280x120400544: callq 400550 <mult2> 0x4005490x118 400549: mov %rax, (%rbx) ← 0x118%rsp %rip. 0×400550 0000000000400550 <mult2>: 400550: mov %rdi,%rax 400557: retq



Control Flow Example #4

```
0000000000400550 <mult2>:
   400550: mov %rdi,%rax
   •
   400557: retq
```

- Procedures
 - Stack Structure
 - Calling Conventions
 - Passing control
 - Passing data
 - Managing local data

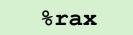
Procedure Data Flow

Registers

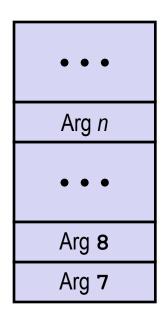
First 6 arguments



Return value



Stack



Only allocate stack space when needed

Data Flow Examples

```
void multstore
  (long x, long y, long *dest)
{
    long t = mult2(x, y);
    *dest = t;
}
```

```
long mult2
  (long a, long b)
{
  long s = a * b;
  return s;
}
```

```
0000000000000400550 <mult2>:
    # a in %rdi, b in %rsi
400550: mov %rdi,%rax # a
400553: imul %rsi,%rax # a * b
# s in %rax
400557: retq # Return
```

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