

Name - CHAKSH

Section - Ds 1

Roll No : 37

1 - Linear Search (pseudocode)

int lin-search (int * ar, int n, int key)

{ for (i = 0 to n-1)

{ if (ar[i] == key)

return i;

} return -1;

2 - Iterative insertion sort

void insert sort (int ar[], int n)

int i, temp, j;

for i = 1 to n

temp = ar[i]

j = i-1

while (j >= 0 AND ar[j] > temp)

ar[j+1] = ar[j]

ar[j+1] = temp

Recursive Insertion sort

void insert = sort (int ar[], int n)

if (n <= 1)

return

insert = sort (ar, n-1)

last = ar[n-1]

j = n-2

while (j >= 0 && ar[j] > last)

ar[j+1] = ar[j]

j--

arr[j+1] = last :-

Why insertion sort is called online sorting?
Because it doesn't need to know anything about what values it will sort and the information is requested while the algo. is running.

3- (i) Selection sort :

T.C = Best case = $O(n^2)$; worst case = $O(n^2)$

S.C = $O(1)$

(ii) Insertion sort :

T.C = Best case = $O(n)$; worst case = $O(n^2)$

S.C = $O(1)$

(iii) Merge sort :

~~the~~ T.C = Best case = $O(n \log n)$; worst case = $O(n \log n)$

S.C = $O(n)$

(iv) Quick sort :

T.C = Best case = $O(n \log n)$; worst case = $O(n^2)$

S.C = $O(1)$

(v) Bubble sort :

T.C = Best case = $O(n^2)$; worst case = $O(n^2)$

S.C = $O(1)$

4- Sorting	Inplace	Stable	Online
selection	✓		
Insertion	✓	✓	✓
Merge		✓	
Quick	✓		
Heap	✓		
Bubble	✓	✓	

5- Iterative Binary Search

```
int bin_search (int ar[], int l, int r, int x)
```

```
{ while (l <= r) {  
    int = (l + r) / 2;
```

```
    if (ar[m] == x)
```

```
        return m;
```

```
    if (ar[m] < x)
```

```
        l = m + 1;
```

```
    else
```

```
        r = m - 1;
```

```
}  
return -1;
```

```
}
```

Recursive Binary search

```
int bin_search (int ar[], int l, int r, int x)
```

```
{ if (x >= 1) {
```

```
    int mid = (l + r) / 2;
```

```
    if (ar[mid] == x)
```

```
        return mid;
```

```
    else if (ar[mid] > x)
```

```
        return bin_search(ar, l, mid - 1, x);
```

```
    else
```

```
        return bin_search(ar, mid + 1, r, x);
```

```
}  
return -1;
```

```
}
```

T.C

Best case = $O(1)$

Avg. case = $O(\log n)$

Worst case = $O(\log n)$

T.C

Best case = $O(1)$

Avg. case = $O(\log n)$

Worst case = $O(\log n)$

6 - Recurrence Relation for binary recursive search

$$T(n) = T(n/2) + C$$

8 - Which sort is used practically?

Quick sort is the fastest general purpose sort. In most practical situation, quick sort is the method of choice. If stability is important and space is available, merge sort might be best.

9 - What is Inversion count for an array?

How far (or close) the array is from being sorted. If the array is already sorted, then the inversion count is 0, but if array is sorted in reverse order, the inversion count is max...

For following array $ar[] = \{7, 21, 31, 8, 10, 1, 20, 6, 4, 5\}$

```
#include <bits/stdc++.h>
```

```
using namespace std;
```

```
int mergesort (int ar[], int temp[], int left, int right);
```

```
int merge (int ar[], int temp[], int left, int mid, int right);
```

```
int mergesort (int ar[], int temp[], int left, int right)
```

```
{ int mid, inv-count = 0;
```

```
  if (right > left) {
```

```
    mid = (right + left) / 2;
```

```
    inv-count += mergesort (ar, temp, left, mid);
```

```
    inv-count += mergesort (ar, temp, mid + 1, right);
```

```
    inv-count += merge (ar, temp, left, mid + 1, right);
```

```
  }
```

```
  return inv-count;
```

```

int merge (int ar [], int temp [], int left, int mid, int right);
{
    int i, j, k;
    int inv-count = 0;
    i = left;
    j = mid;
    k = left;
    while (i <= mid-1) & (j <= right)
    {
        if (ar[i] < ar[j])
            temp[k++] = ar[i++];
        else
        {
            temp[k++] = ar[j++];
            inv-count = inv-count + (mid-i);
        }
    }
    while (i <= mid-1)
        temp[k++] = ar[i++];
    while (j <= right)
        temp[k++] = ar[j++];
    for (i = left; i <= right; i++)
        ar[i] = temp[i];
    return inv-count;
}

int main ()
{
    int ar [] = { 7, 21, 31, 8, 10, 1, 20, 6, 4, 5 };
    int n = size of (ar) (size of (ar[0]));
    int ans = merge sort (ar, n);
    cout << "No. of inversion are" << ans;
    return 0;
}

```

10 - The worst case time complexity of quick sort is $O(n^2)$ the worst case occurs when the picked pivot is always an extreme (smallest or largest) element. This happens when input array is sorted and either first or last element is picked as pivot. The best case of quick sort is when we will select pivot as a mean element.

11) Recurrence relation of:

① Merge sort $\rightarrow T(n) = 2T(n/2) + n$

② Quick sort $\rightarrow T(n) = \alpha T(n/2) + n$

\rightarrow Merge sort is more efficient & works faster than quick sort in case of larger array.

\rightarrow Worst case complexity for quick sort is $O(n^2)$ where $O(n \log n)$ for merge sort.

12) Stable Selection Sort

using namespace;

```
void stab_sel_sort (int a[], int n) {
```

```
    for (int i = 0; i < n - 1; i++) {
```

```
        int min = i;
```

```
        for (int j = i + 1; j < n; j++)
```

```
            if (a[min] > a[j])
```

```
                min = j;
```

```
        int key = a[min];
```

```
        while (min > i)
```

```
            { a[min] = a[min - 1];
```

```
            }
    }
```

```
    a[i] = key;
```

```
}
```

```

int main(){
    int a[] = {4, 5, 3, 4, 18};
    int n = size of (a) / size of (a[0]);
    stab - setsort - (a, n);
    for (int i=0; i<n; i++)
        cout << a[i] << " ";
    cout << endl;
    return 0;
}

```

- 13 - The easiest way to do this is to use external sorting. We divide our source file into temporary files of size equal to the size of RAM and first sort these files -
-) External Sorting : If the input data is such that it cannot adjust in the memory entirely at once, it needs to be stored in hard disk, floppy disk etc. This is external sorting.
 -) Internal Sorting : If the input data is such that it can adjust in the main memory at once, it is called internal sorting.