(3.1)

Why is the program counter a pointer and not a counter?

> It saves a step. If it were a counter, an address would still need to be stored some-
> where, and arithmatic on this address would need to be performed using the counter.
> But since the program instructions are stored sequentially in memory, the address
> itself can simply be incremented to get the next instruction.

(3.2)

Explain the function of the following registers in a CPU.

a) PC

> The Program Counter. It is incremented after an instruction is fetched. It
> holds the memory address of the next instruction to be executed.

b) MAR

> The Memory Address Register. It works in tandom with the MDR. It either
> contains the location of the where the MDR data should be written to, or the
> location of data which will be read into the MDR.

c) MDR

> The Memory Data Register. It contains the data referred to in the MAR
> explanation above.

d) IR

> Instruction Register. This contains the instruction which is currently being executed. It is fetched from the location pointed to by the PC.

(3.3)

For each of the following 6-bit operations, calculate the values of the C, Z, V, and N flags.

a.

$$
\begin{array}{r}
001011 \\
+001101 \\
\hline
011000
\end{array}
\quad C = 0, Z = 0, V = 0, N = 0
$$

b.

$$
\begin{array}{r}
111111 \\
+000001 \\
\hline
1000000
\end{array}
\quad C = 1, Z = 1, V = 0, N = 0
$$

c.

$$
\begin{array}{r}
000000 \\
-111111 \\
\hline
111111
\end{array}
\quad C = 0, Z = 0, V = 1, N = 1
$$

d

$$
\begin{array}{r}
101101 \\
+011011 \\
\hline
1001000
\end{array}
\quad C = 1, Z = 0, V = 0, N = 1
$$

e

$$
\begin{array}{r}
000000 \\
-000001 \\
\hline
111110
\end{array}
\quad C = 0,\ Z = 0,\ V = 0,\ N = 1
$$

f.

$$
\begin{array}{r}
111110 \\
+111111 \\
\hline
1111101
\end{array}
\quad C = 1,\ Z = 0,\ V = 1,\ N = 1
$$

(3.10)

Why does the ARM provide a reverse subtract instruction RSB r0,r1,r2 that implements [r0] = [r2] - [r1] when the normal subtraction instruction SUB r0,r2,r1 will do exactly the same job?

These two operand only differ by the order of operation, and can be seen as unnecessary as switching the two operands for SUB should suffice. The problem arises when only the second operand can be constant or other special forms. All operations could be done using the SUB operations but additional registers would be necessary as seen in this example: SUB $R2$, $R4$, $\#8$ if you wanted to do the exact opposite you would need an an additional register to hold the constant $\#8$, but using RSB it would just be RSB $R2$, $R4$, $\#8$

(3.17)

ARM instructions have a 12-bit literal. Instead of permitting a word in the range 0 to $2^12 - 1$, the ARM uses an 8-bit format for the integer and a 4-bit alignment field that allows the integer to be shifted in steps of two. What are the advantages and disadvantages of this mechanism in comparison to a straight 12-bit integer?

> These 8-bit literal that can be scaled by a power of 2. you could even say that arm provides a type of floating point literal.
>
> The four most significant bits of the literal field specify the literal's alignment within a 32-bit word. if the 8 bit immediate value is N and the 4-bit alignment is n in the range 0-12 then the value of the literal is given by $N \times 2^{2n}$. Thus, arm provides an 8-bit literal that is scaled by a power of 2.

(3.18)

Write one or more ARM instructions that will clear bits 20 to 25 inclusive in register r0. All the other bits of r0 should remain unchanged.

> $LDR \quad R3, = 0xFC0F\ FFFF$
> $AND \quad R0, R3, R0$

```
1      AREA clearbits, CODE, READONLY
       ENTRY
3 start
       AND r0,r0,#2_11111111111111111110000001111111
5      END
```

(3.19)

This is a classic problem of assembly language programming. Write a sequence of ARM instructions that swap the contents of registers r0 and r2 without using any additional registers or memory storage (that is, you can't move r1 to a temporary location).

> Our original thought similar the the integer implementation where you add them and together and subtract the other like this
>
> ```
> x=15;
> y=21;
>
> x=x+y;
> y=x-y;
> x=x-y;
> ```
>
> But we had some concern with the S bits. This lead us to some research online where we found an article titled **The Magic of XOR** [1] which introduced us to an extremely fascinating method of swapping two values. This process is basically XORing the three times as the pseudocode below depicts.
>
> ```
> x=1101;
> y=0110;
>
> x=x EOR y;        //1101 EOR 0110 = 1011
> y=x EOR y;        //1011 EOR 0110 = 1101
> x=x EOR y;        //1011 EOR 1101 = 0110
> ```

(3.25)

What is the binary encoding of the following instructions?

	Condition	0 0		Op-code	S	$r_{source1}$	$r_{destination}$	Operand 2
A								

	Condition	0 1	#	P	U	b	W	L	r_{base}	$r_{transfer}$	Operand 2
B											
C											
D											

a) STRB r1, [r2]

E5C21000 → 11100101110000100001000000000000

b) LDR r3, [r4, r5]!

E7B43005 → 11100111101101000011000000000101

c) LDR r3, [r4], r5

E6943005 → 11100110100101000011000000000101

d) LDR r3, [r4,#-6]!

E5343006 → 11100101001101000011000000000110

(3.39)

Write an ARM program that scans a string terminated by the null byte 0x00 and copies the string from a source location pointed by r0 to a destination pointed at by r1.

NullTerminatingString/NullTerminatingString.asm

```
1      AREA NullTerminatingString , CODE , READONLY
       ENTRY
3      ;; copy of code from section 3.9.11

5 Banner   ADR R1 , String
  Bnner1   LDRB   R0 , [R10] , #1
7     CMP    R0 , #0
      SVCNE    WriteC
9     BNE    Bnner1
      MOV    PC , R14
11
  String =  "A String goes here!", &0A , &0D , 0
13 WriteC   EQU    0
      END
```

(3.51)

Write an ARM assembly language program to determine whether a string of characters with an odd length is a palindrome (for example, mom) under the following constraints.

    a. The string of ASIC-encoded characters is stored in memory.

    b. At the start of the program, register r1 contains the address of the first character in the string, and r2 contains the address of the last character. On exit from the program, register r0 contains a 0 if the string is a not a palindrome, and 1 if it is.

../lab2.new/lab2_part2/palindrome.asm

```
      AREA palindrome, CODE, READONLY
2     ENTRY
        ; source: http://www.ing.unlp.edu.ar/electrotecnia/arcom1/09-
          Clements-Chap09-ARM.pdf
4
  start
6     ldr r0, =srcstr
      mov r1,r0              ;copy r0 data
8 loop     ldrb   r2,[r1],#1      ;Double word
      cmp r2,#0
10    bne loop
      sub r1,r1,#2
12    bl   pal

14 pal    mov r10, #0x0
  loop2 ldrb   r3, [r0]
16    ldrb   r4, [r1]
      cmp r3, r4
18    bne notpal

20    cmp    r0, r1
      beq waspal
22    add r2,r0,#1
      cmp r2,r1
24    beq waspal

26    add r0,r0,#1
      sub r1,r1,#1
28    b    loop2

30 waspal     mov r10, #0x1
  notpal     bx lr
32
      AREA   Strings, DATA
34 srcstr  DCB      "First string - source",0
  dststr  DCB      "Second string - destination",0
36    END
      END
```

8

References

[1] Charles Lin. The magic of xor, 2003,2004.