

Ejercicio 1

a)

- `return x >>= f`
 $\equiv \langle def.return \rangle$
`State (\s -> (x, s)) >>= f`
 $\equiv \langle def.>>= \rangle$
`State (\s -> let (v, s') = runState (State (\s -> (x, s))) s
 in runState (f v) s')`
 $\equiv \langle def.runState \rangle$
`State (\s -> let (v, s') = (\s -> (x, s)) s
 in runState (f v) s')`
 \equiv_{β}
`State (\s -> let (v, s') = (x, s)
 in runState (f v) s')`
 $\equiv \langle def.let \rangle$
`State (\s -> runState (f x) s)`
 \equiv_{η}
`State (runState (f x))`
 $\equiv \langle def.\circ \rangle$
`(State . runState) (f x)`
 $\equiv \langle State \circ runState \equiv id \rangle$
`f x`

- `State g >>= return`
 $\equiv \langle \text{def. } >>= \rangle$
`State (\s -> let (v, s') = runState (State g) s`
`in runState (return v) s')`
 $\equiv \langle \text{def. runState} \rangle$
`State (\s -> let (v, s') = g s`
`in runState (return v) s')`
 $\equiv \langle \text{def. return} \rangle$
`State (\s -> let (v, s') = g s`
`in runState (State (\s -> (v, s))) s')`
 $\equiv \langle \text{def. runState} \rangle$
`State (\s -> let (v, s') = g s`
`in (\s -> (v, s)) s')`
 \equiv_{β}
`State (\s -> let (v, s') = g s`
`in (v, s'))`
 $\equiv \langle \text{def. let} \rangle$
`State (\s -> g s)`
 \equiv_{η}
`State g`

- $((\text{St } t) \gg= f) \gg= g$
 $\equiv \langle \text{def. } \gg= \rangle$
 $\text{St } (\backslash s1 \rightarrow \text{let } (x2, s2) = t \text{ s1}$
 $\quad \text{in runState } (f \ x2) \ s2) \gg= g$
 $\equiv \langle \text{def. } \gg= \rangle$
 $\text{St } (\backslash s3 \rightarrow \text{let } (x4, s4) = (\backslash s1 \rightarrow \text{let } (x2, s2) = t \text{ s1}$
 $\quad \text{in runState } (f \ x2) \ s2) \ s3)$
 $\quad \text{in runState } (g \ x4) \ s4)$
 \equiv_{β}
 $\text{St } (\backslash s3 \rightarrow \text{let } (x4, s4) = \text{let } (x2, s2) = t \text{ s3}$
 $\quad \text{in runState } (f \ x2) \ s2$
 $\quad \text{in runState } (g \ x4) \ s4)$
 $\equiv \langle \text{Lema} \rangle$
 $\text{St } (\backslash s3 \rightarrow \text{let } (x2, s2) = t \text{ s3}$
 $\quad \text{in let } (x4, s4) = \text{runState } (f \ x2) \ s2$
 $\quad \text{in runState } (g \ x4) \ s4)$
 \equiv_{β}
 $\text{St } (\backslash s3 \rightarrow \text{let } (x2, s2) = t \text{ s3}$
 $\quad \text{in } (\backslash s5 \rightarrow \text{let } (x4, s4) = \text{runState } (f \ x2) \ s5$
 $\quad \text{in runState } (g \ x4) \ s4) \ s2)$
 $\equiv \langle \text{runState} \circ \text{St} \equiv \text{id} \rangle$
 $\text{St } (\backslash s3 \rightarrow \text{let } (x2, s2) = t \text{ s3}$
 $\quad \text{in runState } (\text{St } (\backslash s5 \rightarrow \text{let } (x4, s4) = \text{runState } (f \ x2) \ s5$
 $\quad \text{in runState } (g \ x4) \ s4) \ s2)$
 $\equiv \langle \text{def. } \gg= \rangle$
 $\text{St } (\backslash s3 \rightarrow \text{let } (x2, s2) = t \text{ s3}$
 $\quad \text{in runState } ((f \ x2) \gg= g) \ s2)$
 \equiv_{β}
 $\text{St } (\backslash s3 \rightarrow \text{let } (x2, s2) = t \text{ s3}$
 $\quad \text{in runState } ((\backslash w \rightarrow f \ w \gg= g) \ x2) \ s2)$
 $\equiv \langle \text{def. } \gg= \rangle$
 $\text{St } t \gg= (\backslash w \rightarrow f \ w \gg= g)$

- Lema:

```

let x = let y = z
      in f y
in g x
≡ ⟨y ∉ FV (g x)⟩
let y = z
in let x = f y
  in g x

```

b)

- `evalComm :: MonadState m => Comm -> m ()`
`evalComm (Skip) = return ()`
`evalComm (Let v e) = do x <- evalIntExp e;`
`update v x`
`evalComm (Seq com1 com2) = (evalComm com1) >> (evalComm com2)`
`evalComm (Cond expB com1 com2) = do b <- evalBoolExp expB;`
`else evalComm com2`
`evalComm (While expB com) = do b <- evalBoolExp expB;`
`if b`
`then evalComm (Seq com (While expB com))`
`else (evalComm Skip)`
- `evalIntExp :: MonadState m => IntExp -> m Int`
`evalIntExp (Const n) = return n`
`evalIntExp (Var s) = lookfor s`
`evalIntExp (UMinus e) = do x <- evalIntExp e;`
`return (negate x)`
`evalIntExp (Plus e1 e2) = abstract evalIntExp (+) e1 e2`
`evalIntExp (Minus e1 e2) = abstract evalIntExp (-) e1 e2`
`evalIntExp (Times e1 e2) = abstract evalIntExp (*) e1 e2`
`evalIntExp (Div e1 e2) = abstract evalIntExp div e1 e2`

- ```
evalBoolExp :: MonadState m => BoolExp -> m Bool
evalBoolExp BTrue = return True
evalBoolExp BFalse = return False
evalBoolExp (Eq e1 e2) = abstract evalIntExp (==) e1 e2
evalBoolExp (Lt e1 e2) = abstract evalIntExp (<) e1 e2
evalBoolExp (Gt e1 e2) = abstract evalIntExp (>) e1 e2
evalBoolExp (And e1 e2) = abstract evalBoolExp (&&) e1 e2
evalBoolExp (Or e1 e2) = abstract evalBoolExp (||) e1 e2
evalBoolExp (Not e) = do b <- evalBoolExp e;
 return (not b)
```
- ```
abstract :: Monad m => (a -> m b) -> (b -> b -> c) -> a -> a -> m c
abstract ev op e1 e2 = do x <- ev e1;
                          y <- ev e2;
                          return (op x y)
```

Ejercicio 2

- a)

```
instance Monad StateError where
return x = StateError (\s -> Just (x,s))
t >>= f  = StateError (\s ->
                        (runStateError t s >>= \ (x,s') -> (runStateError (f x) s')))
```
- b)

```
instance MonadError StateError where
throw = StateError (\s -> Nothing)
```
- c)

```
instance MonadState StateError where
lookfor v = StateError (\s -> lookfor' v s s) where
    lookfor' _ [] _ = Nothing
    lookfor' v ((u, j):ss) s' | v == u = Just (j,s')
                              | otherwise = lookfor' v ss s'

update v i = StateError (\s -> update' v i s) where
    update' v i [] = Just((),[(v,i)])
    update' v i ((a,b):xs) | v == a = Just ((),(a,i):xs)
                          | v /= a = update' v i xs >>=
                              \ (x,y) -> Just ((),(a,b):y)
```

d)

- `eval :: Comm -> Env`
`eval p = case (runStateError (evalComm p) initState) of`
 `Nothing -> []`
 `Just (_,s) -> s`
- `evalComm :: (MonadState m, MonadError m) => Comm -> m ()`
`evalComm (Skip) = return ()`
`evalComm (Let v e) = do x <- evalIntExp e;`
 `update v x`
`evalComm (Seq com1 com2) = (evalComm com1) >> (evalComm com2)`
`evalComm (Cond expB com1 com2) = do b <- evalBoolExp expB;`
 `else evalComm com2`
`evalComm (While expB com) = do b <- evalBoolExp expB;`
 `if b`
 `then evalComm (Seq com (While expB com))`
 `else (evalComm Skip)`
- `evalIntExp :: (MonadState m, MonadError m) => IntExp -> m Int`
`evalIntExp (Const n) = return n`
`evalIntExp (Var s) = lookfor s`
`evalIntExp (UMinus e) = do x <- evalIntExp e;`
 `return (negate x)`
`evalIntExp (Plus e1 e2) = abstract evalIntExp (+) e1 e2`
`evalIntExp (Minus e1 e2) = abstract evalIntExp (-) e1 e2`
`evalIntExp (Times e1 e2) = abstract evalIntExp (*) e1 e2`
`evalIntExp (Div e1 e2) = do x <- evalIntExp e2`
 `if x == 0 then throw`
 `else abstract evalIntExp div e1 e2`

- `evalBoolExp :: (MonadState m, MonadError m) => BoolExp -> m Bool`
`evalBoolExp BTrue = return True`
`evalBoolExp BFalse = return False`
`evalBoolExp (Eq e1 e2) = abstract evalIntExp (==) e1 e2`
`evalBoolExp (Lt e1 e2) = abstract evalIntExp (<) e1 e2`
`evalBoolExp (Gt e1 e2) = abstract evalIntExp (>) e1 e2`
`evalBoolExp (And e1 e2) = abstract evalBoolExp (&&) e1 e2`
`evalBoolExp (Or e1 e2) = abstract evalBoolExp (||) e1 e2`
`evalBoolExp (Not e) = do b <- evalBoolExp e;`
`return (not b)`
- `abstract :: Monad m => (a -> m b) -> (b -> b -> c) -> a -> a -> m c`
`abstract ev op e1 e2 = do x <- ev e1;`
`y <- ev e2;`
`return (op x y)`

Ejercicio 3

- a) `newtype StateErrorTick a = StateErrorT`
`{ runStateError :: Env -> Maybe (a, Env, Int) }`
- b) `class Monad m => MonadTick m where`
`tick :: m ()`
- c) `instance MonadTick StateErrorTick where`
`tick = StateErrorT (\s -> Just ((),s,1))`
- d) `instance MonadError StateErrorTick where`
`throw = StateErrorT (\s -> Nothing)`
- e) `instance MonadState StateErrorTick where`
`lookfor v = StateErrorT (\s -> lookfor' v s s) where`
`lookfor' _ [] _ = Nothing`
`lookfor' v ((u, j):ss) s' | v == u = Just (j,s',0)`
`| otherwise = lookfor' v ss s'`

`update v i = StateErrorT (\s -> update' v i s) where`
`update' v i [] = Just((),[(v,i)], 0)`
`update' v i ((a,b):xs) | v == a = Just ((), (a,i):xs, 0)`
`| v /= a = update' v i xs >>=`
`\(x,y,n) -> Just ((),(a,b):y, n)`

f)

- `eval :: Comm -> (Env, Int)`
`eval p = case (runStateError (evalComm p) initState) of`
 `Nothing -> ([], 0)`
 `Just (_, s, n) -> (s, n)`
- `evalComm :: (MonadState m, MonadError m, MonadTick m) => Comm -> m ()`
`evalComm (Skip) = return ()`
`evalComm (Let v e) = do x <- evalIntExp e;`
 `update v x`
`evalComm (Seq com1 com2) = (evalComm com1) >> (evalComm com2)`
`evalComm (Cond expB com1 com2) = do b <- evalBoolExp expB;`
 `else evalComm com2`
`evalComm (While expB com) = do b <- evalBoolExp expB;`
 `if b`
 `then evalComm (Seq com (While expB com))`
 `else (evalComm Skip)`
- `evalIntExp :: (MonadState m, MonadError m, MonadTick m) => IntExp -> m Int`
`evalIntExp (Const n) = return n`
`evalIntExp (Var s) = lookfor s`
`evalIntExp (UMinus e) = do x <- evalIntExp e;`
 `return (negate x)`
`evalIntExp (Plus e1 e2) = abstract evalIntExp (+) e1 e2 True`
`evalIntExp (Minus e1 e2) = abstract evalIntExp (-) e1 e2 True`
`evalIntExp (Times e1 e2) = abstract evalIntExp (*) e1 e2 True`
`evalIntExp (Div e1 e2) = do x <- evalIntExp e2`
 `if x == 0 then throw`
 `else abstract evalIntExp div e1 e2 True`
- `evalBoolExp :: (MonadState m, MonadError m, MonadTick m) => BoolExp -> m Bool`
`evalBoolExp BTrue = return True`
`evalBoolExp BFalse = return False`
`evalBoolExp (Eq e1 e2) = abstract evalIntExp (==) e1 e2 False`
`evalBoolExp (Lt e1 e2) = abstract evalIntExp (<) e1 e2 False`
`evalBoolExp (Gt e1 e2) = abstract evalIntExp (>) e1 e2 False`
`evalBoolExp (And e1 e2) = abstract evalBoolExp (&&) e1 e2 False`
`evalBoolExp (Or e1 e2) = abstract evalBoolExp (||) e1 e2 False`
`evalBoolExp (Not e) = do b <- evalBoolExp e;`
 `return (not b)`

- `abstract` `ev op e1 e2 b = do` `x <- ev e1;`
 `y <- ev e2;`
 `if b then do tick;`
 `return (op x y)`
 `else return (op x y)`