

# Reconfigurable Broadcast Networks and Asynchronous Shared-Memory Systems are Equivalent

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# Two Models

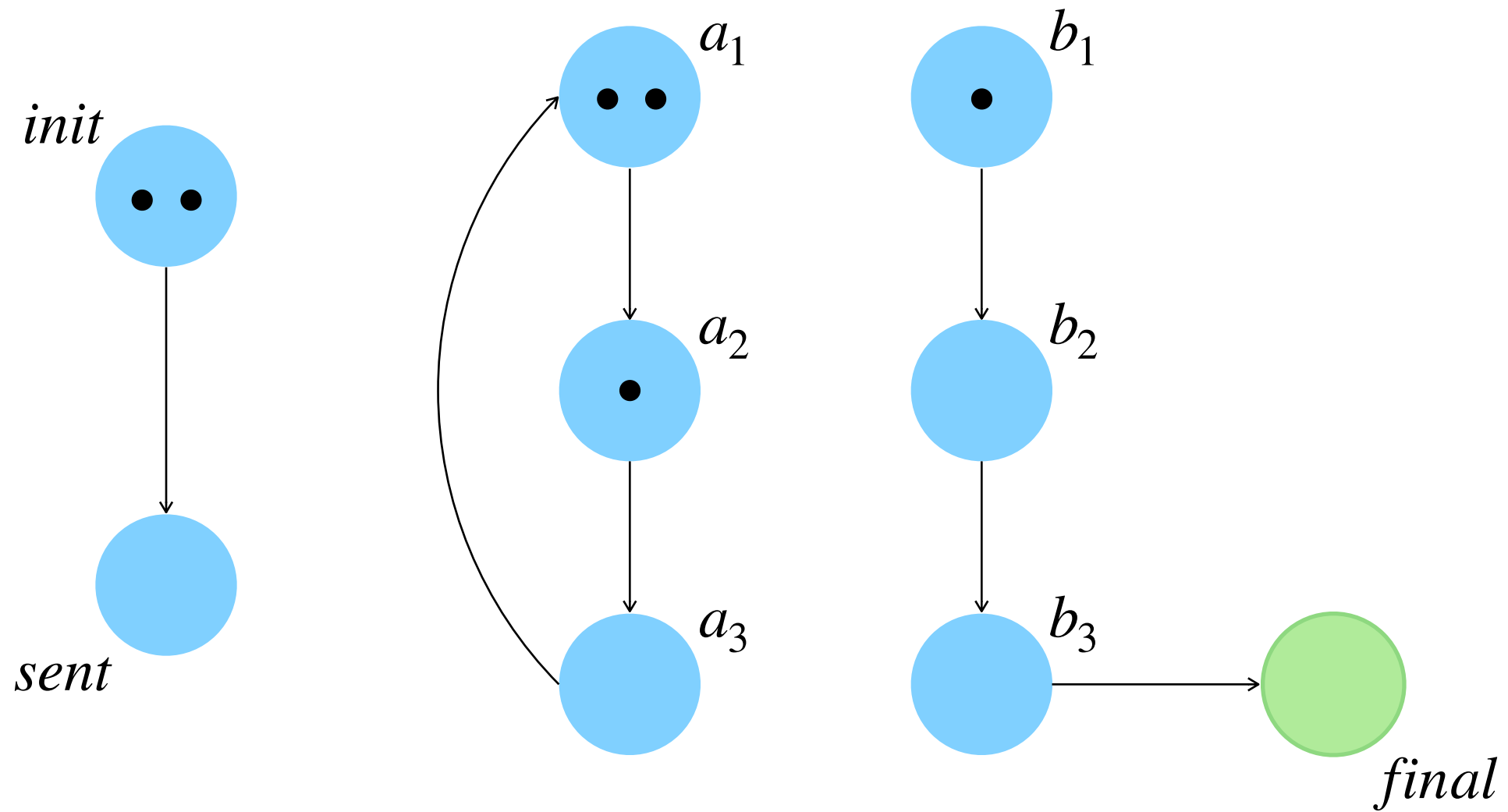
## **Reconfigurable Broadcast Network (RBN)**

- introduced in [Delzanno, Sangnier & Zavattaro, CONCUR '10]
- anonymous, identical processes which can communicate by selective broadcast.

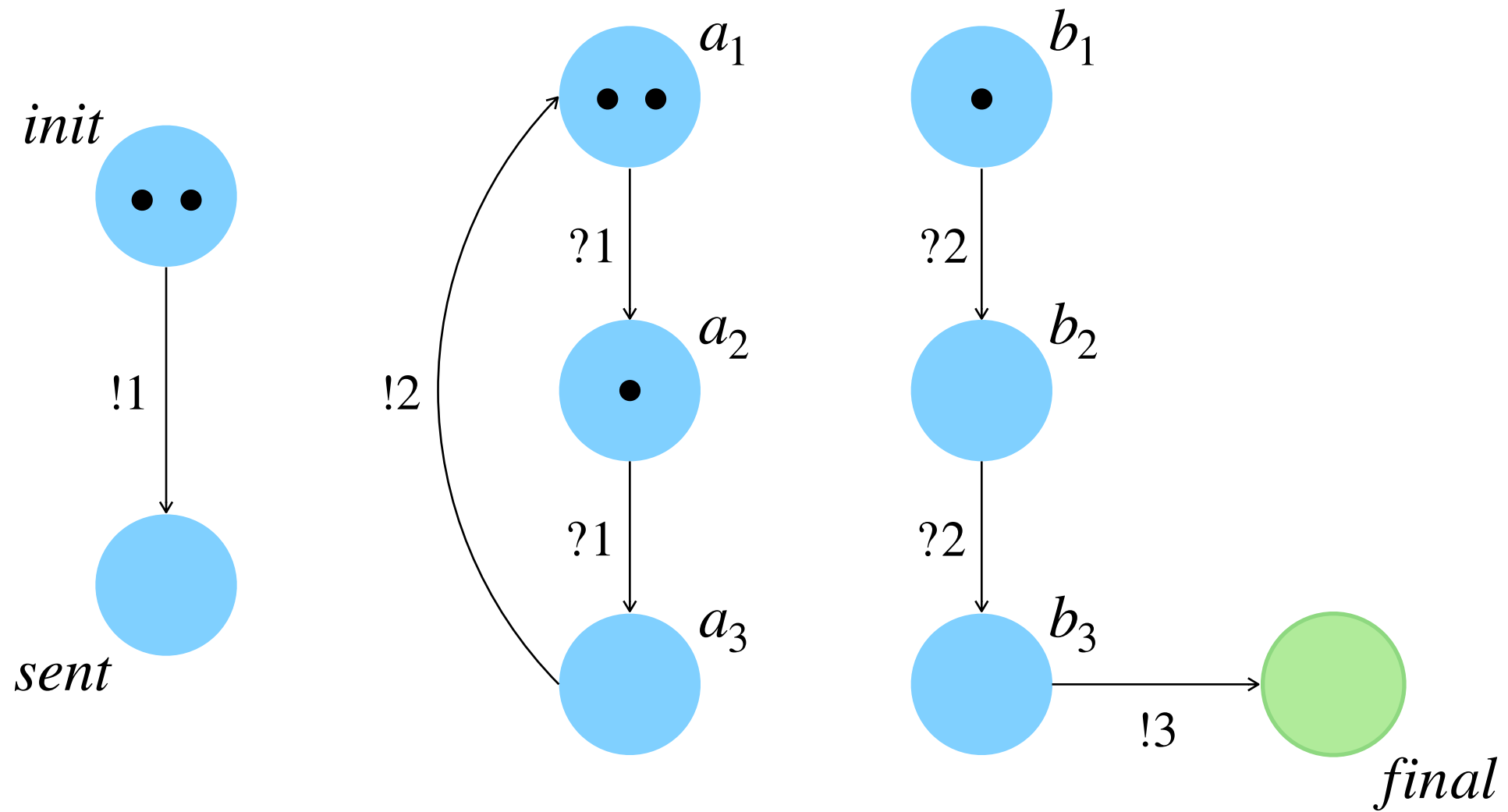
## **Asynchronous Shared Memory System (ASMS)**

- introduced in [Esparza, Ganty & Majumdar, CAV '13]
- anonymous, identical processes which can communicate by writing to a shared register.

# RBN

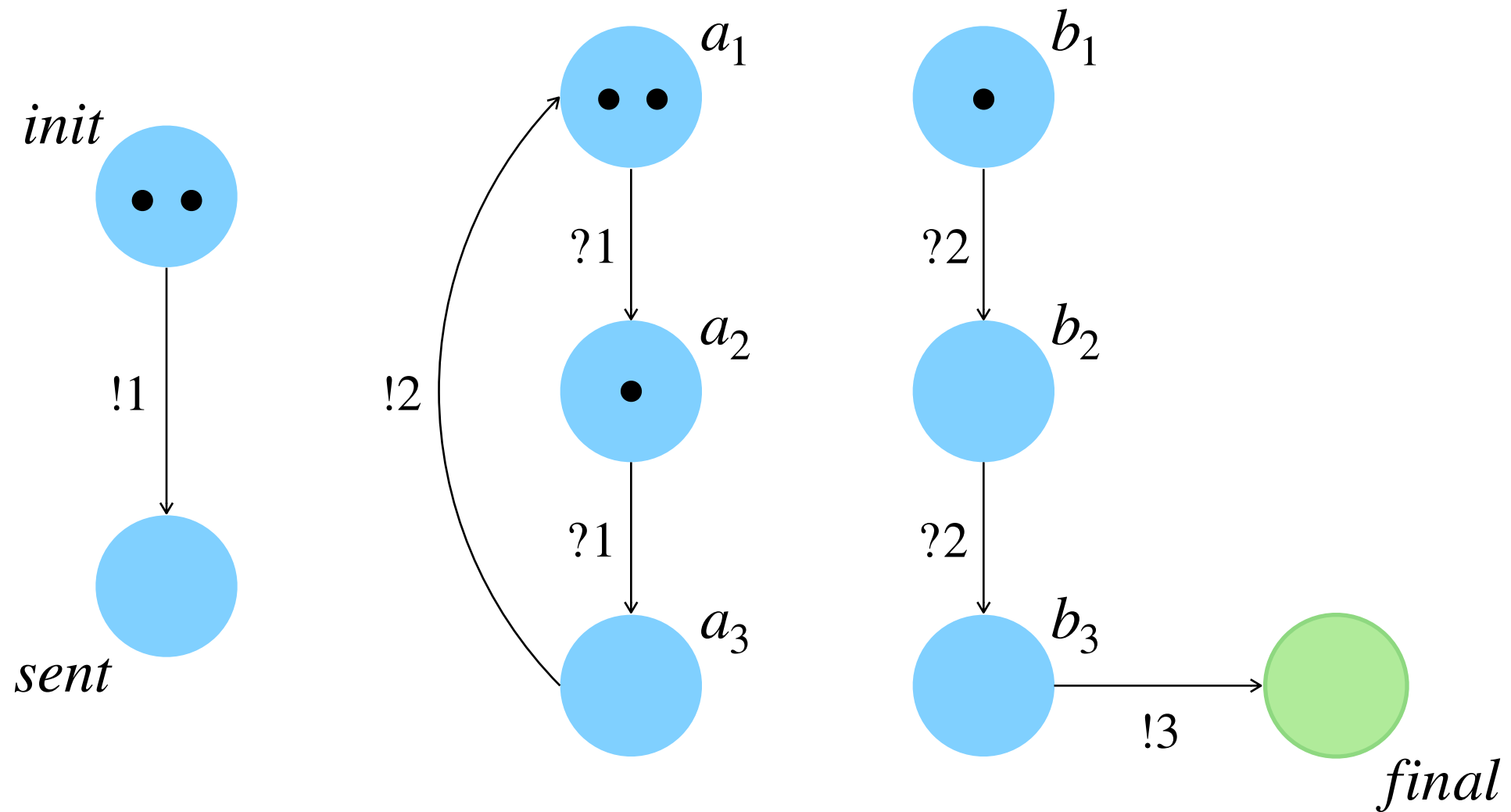


# RBN



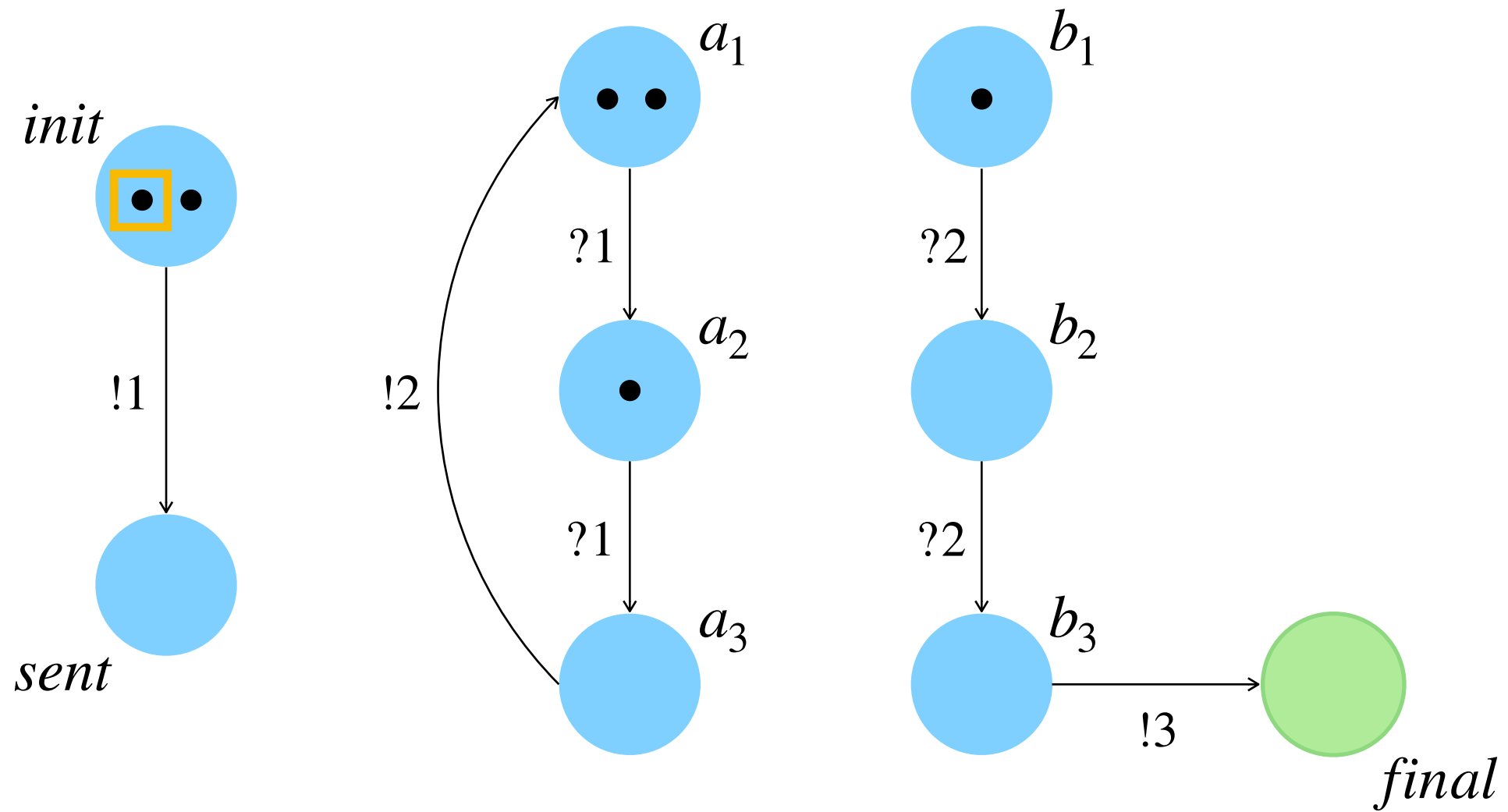
# RBN

Goal: put a process in *final*



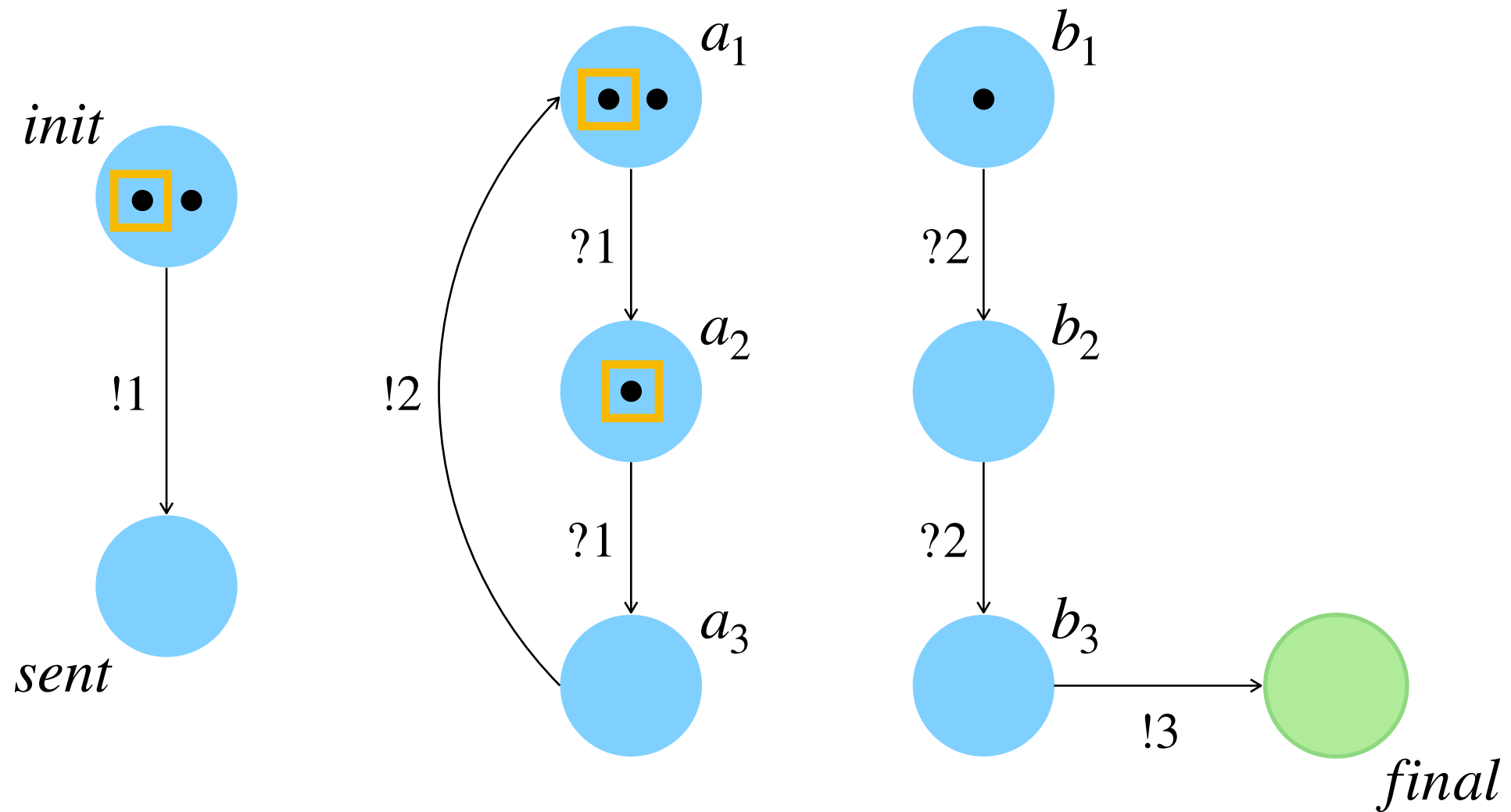
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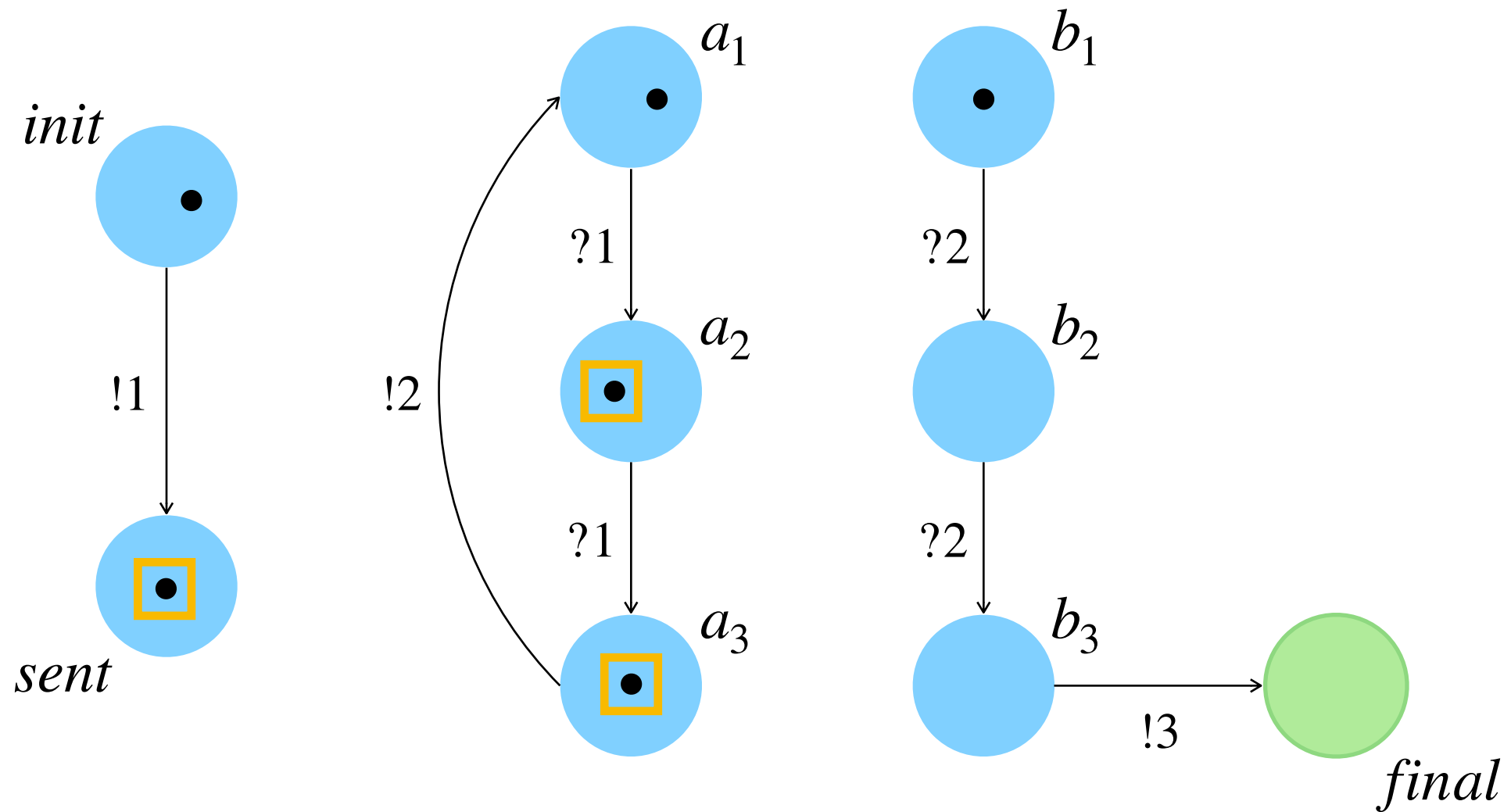
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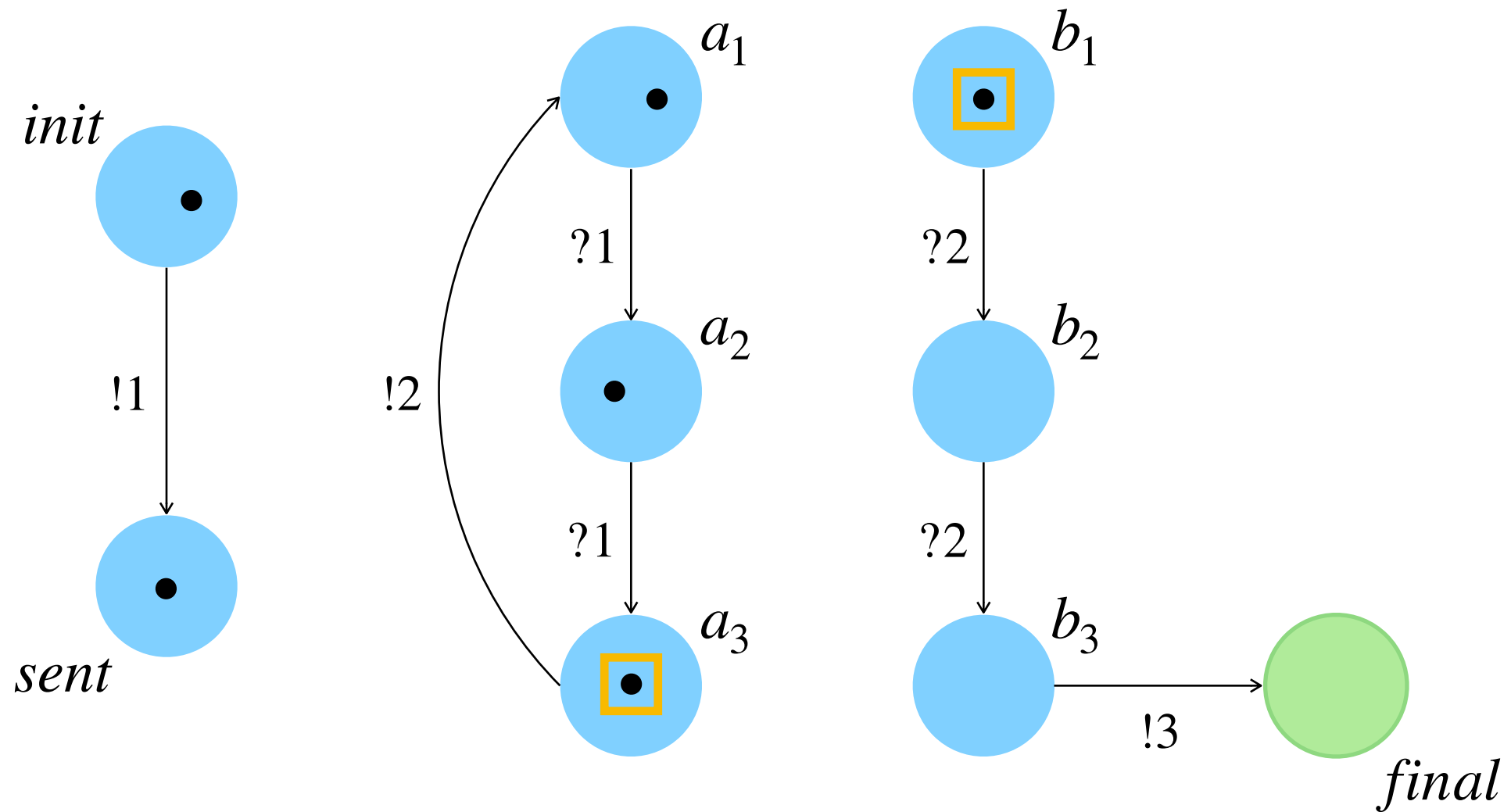
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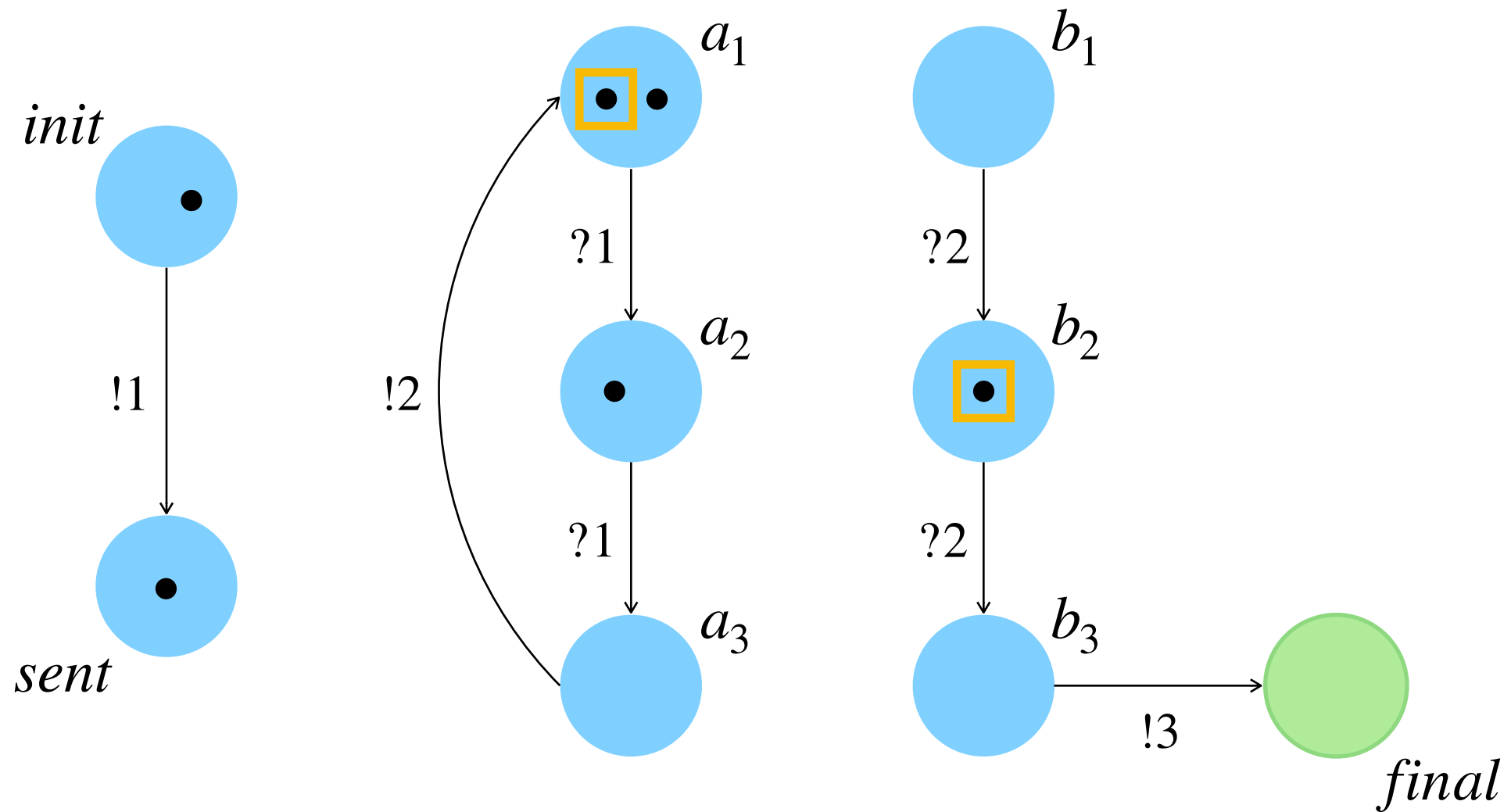
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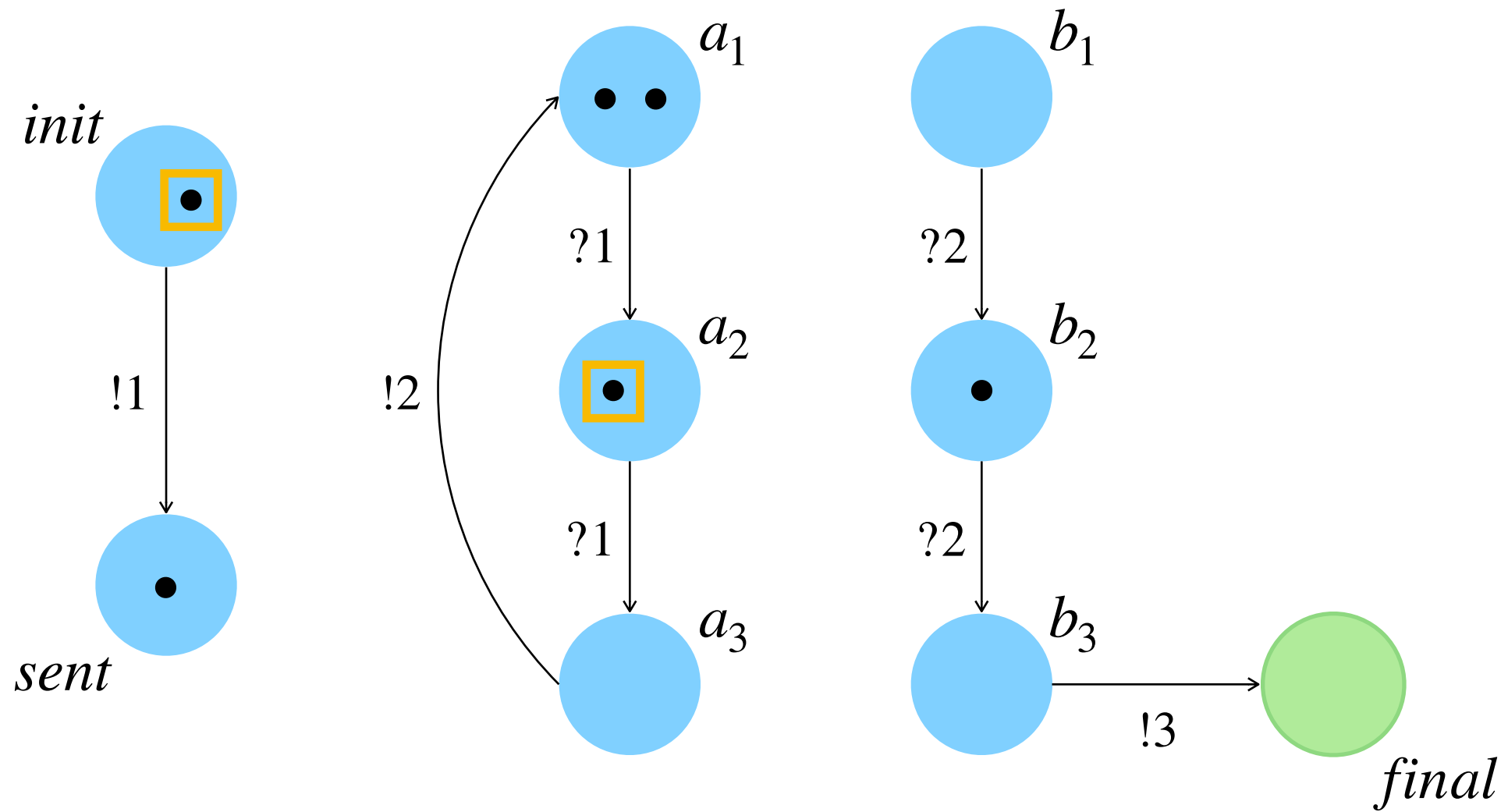
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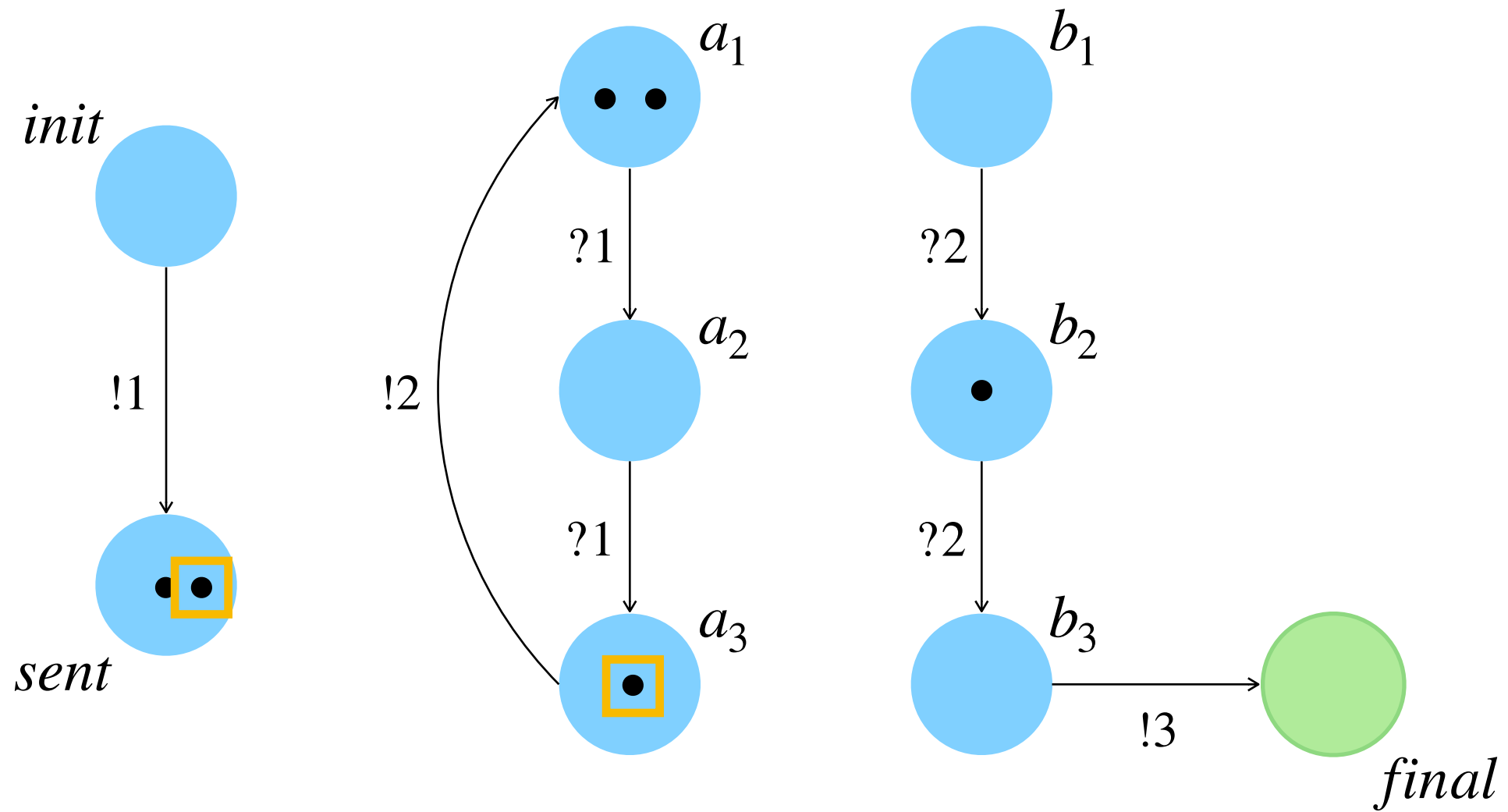
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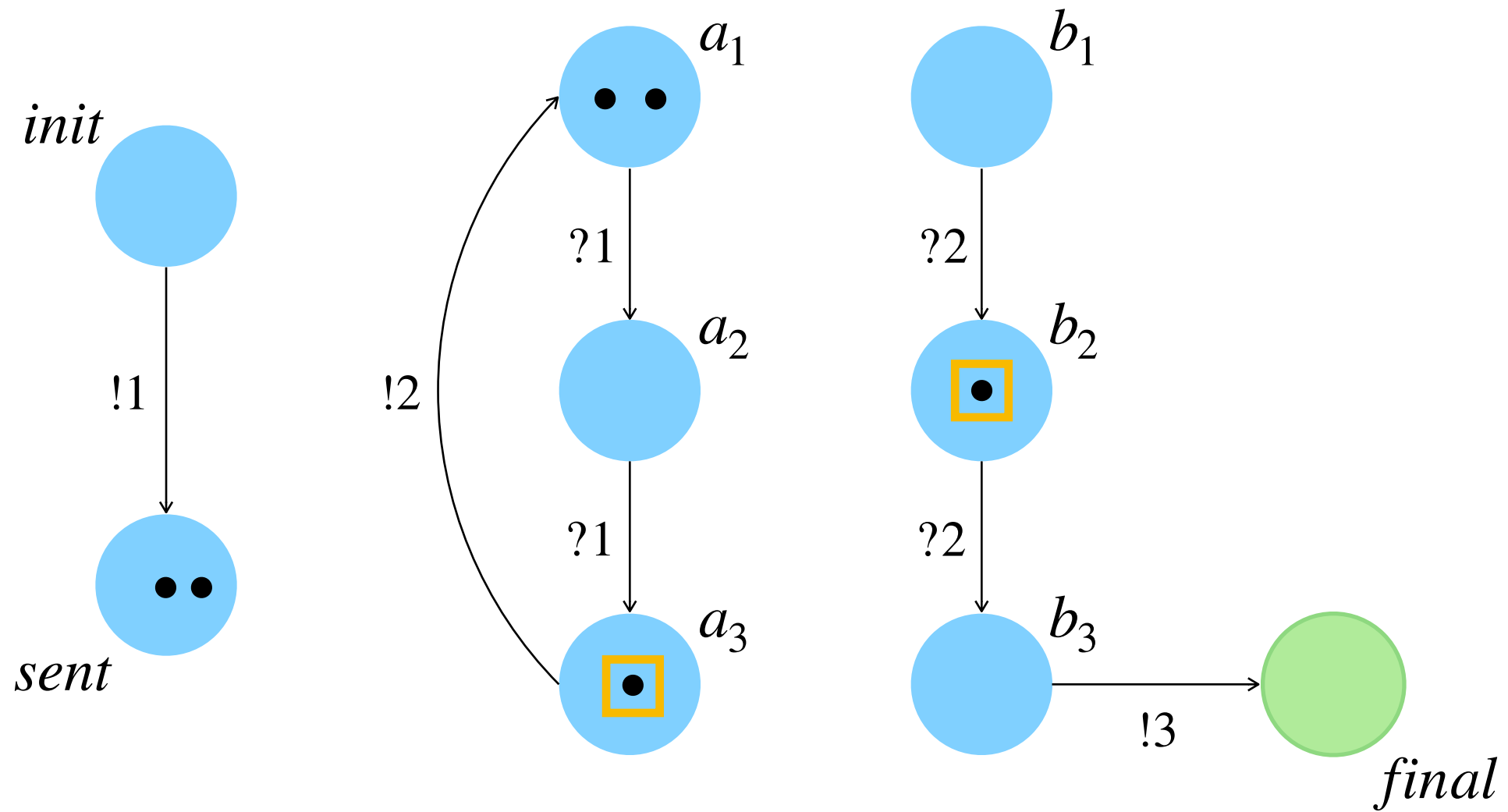
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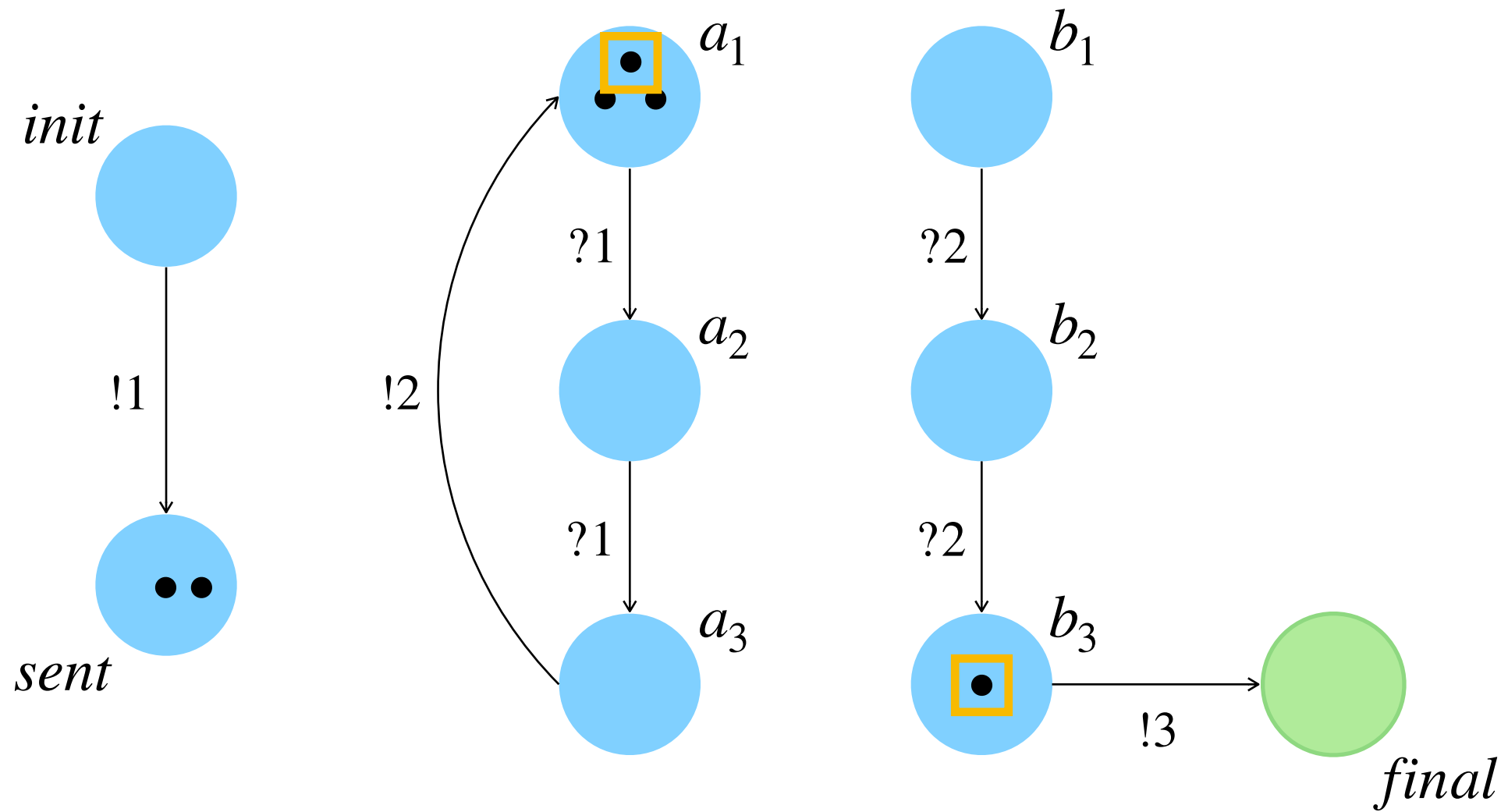
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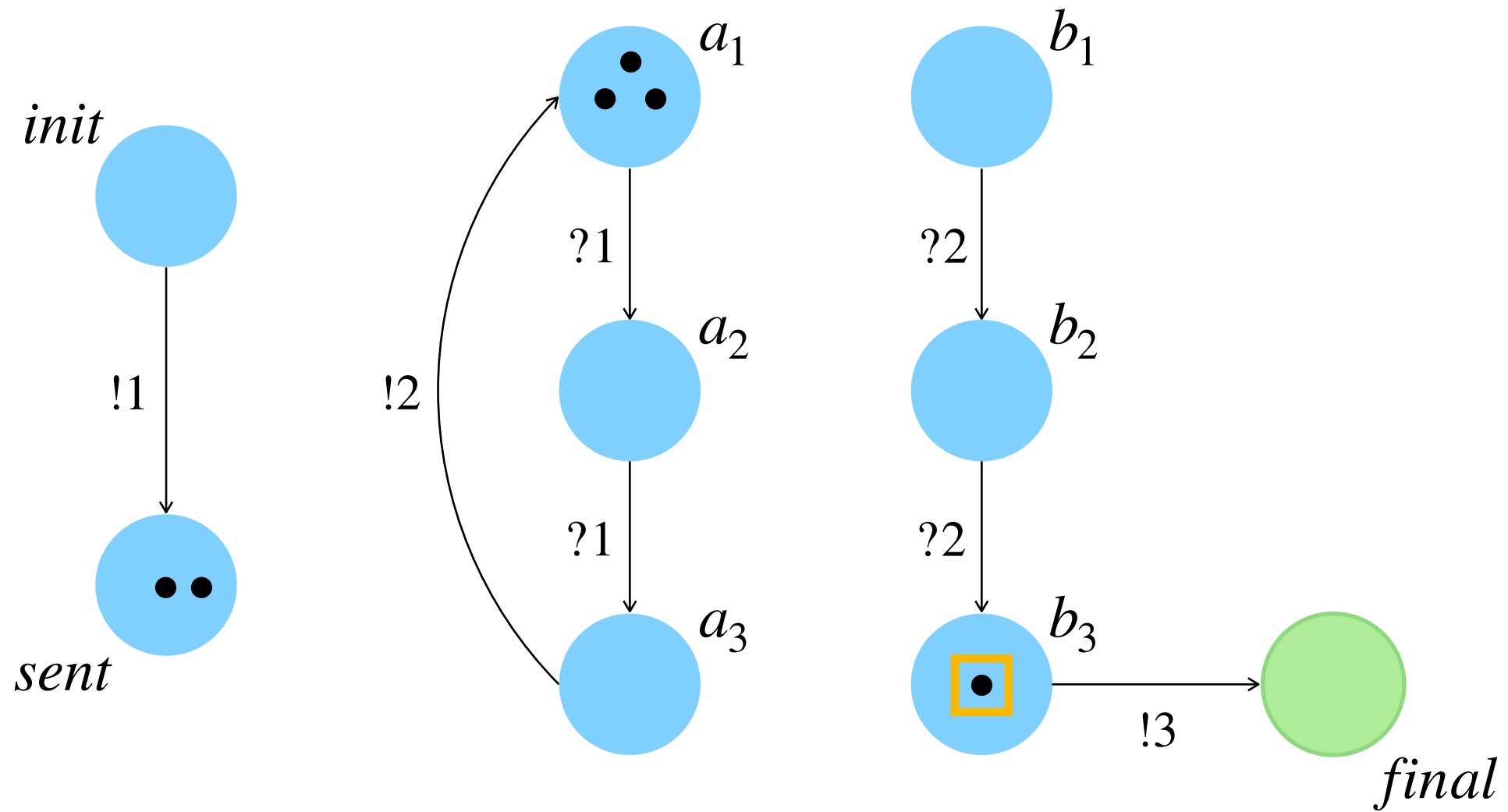
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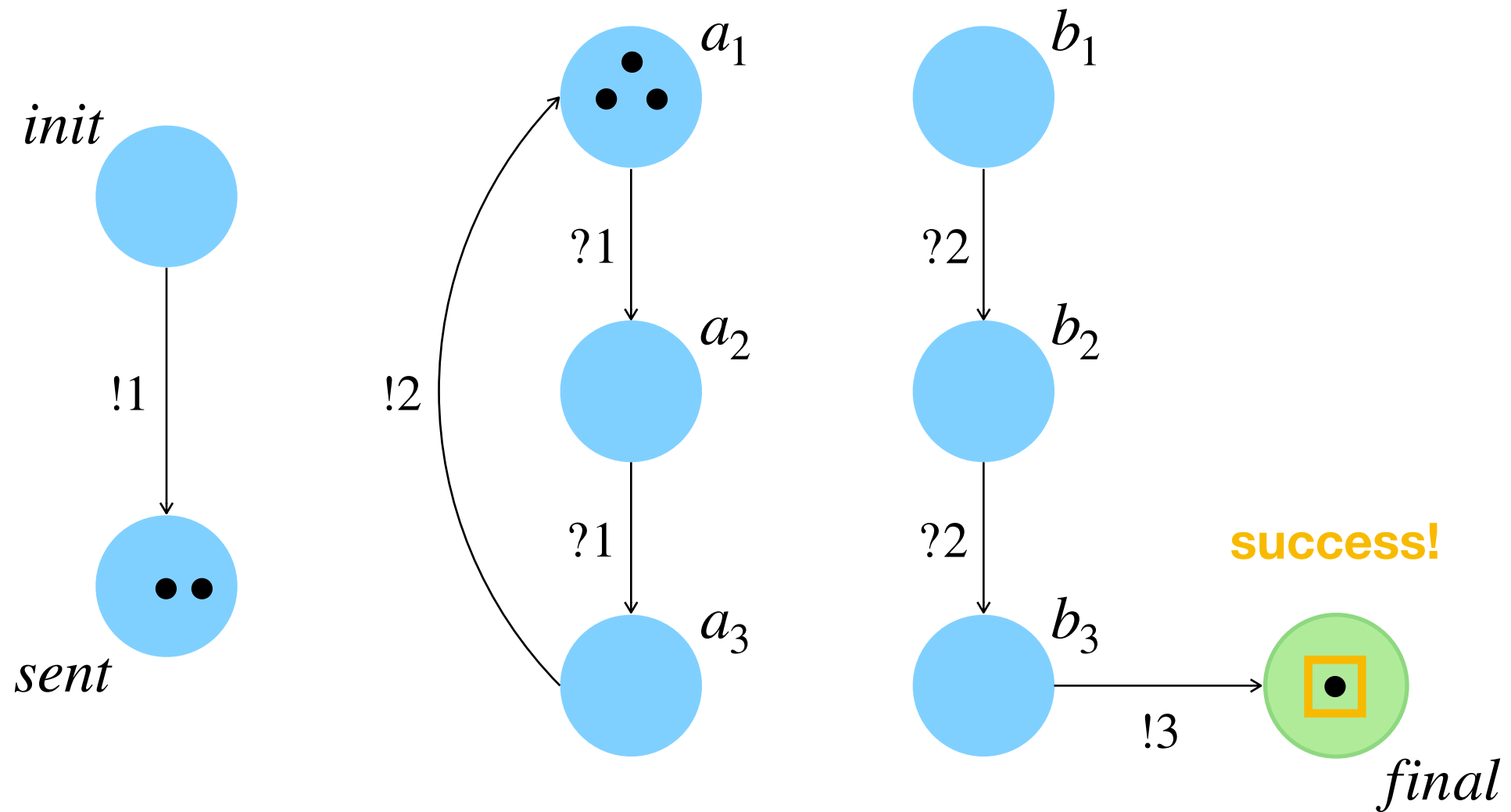
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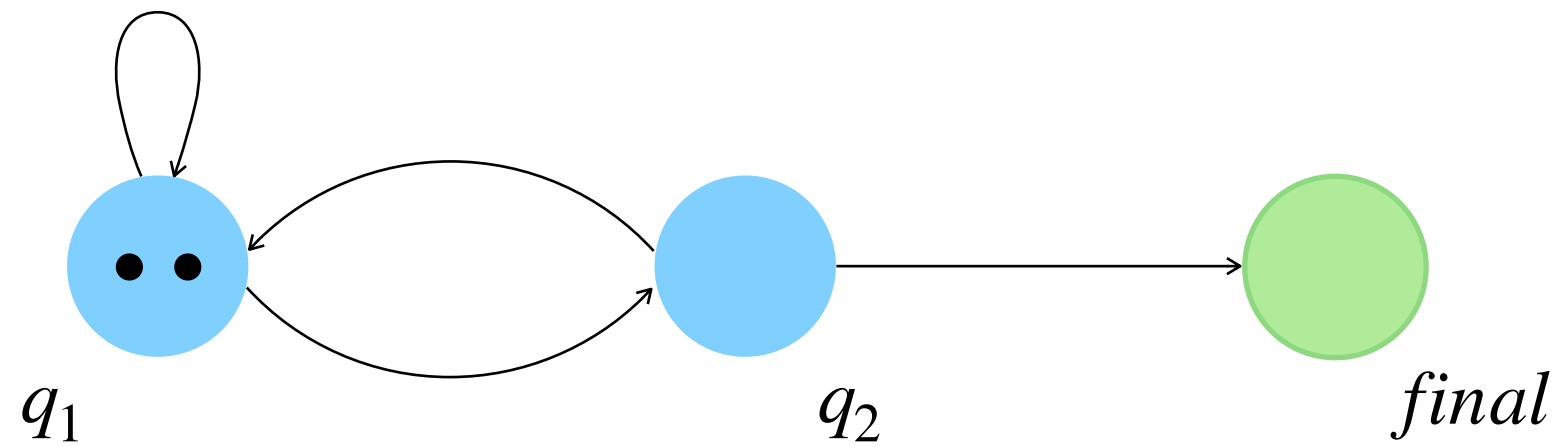
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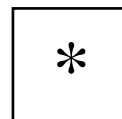
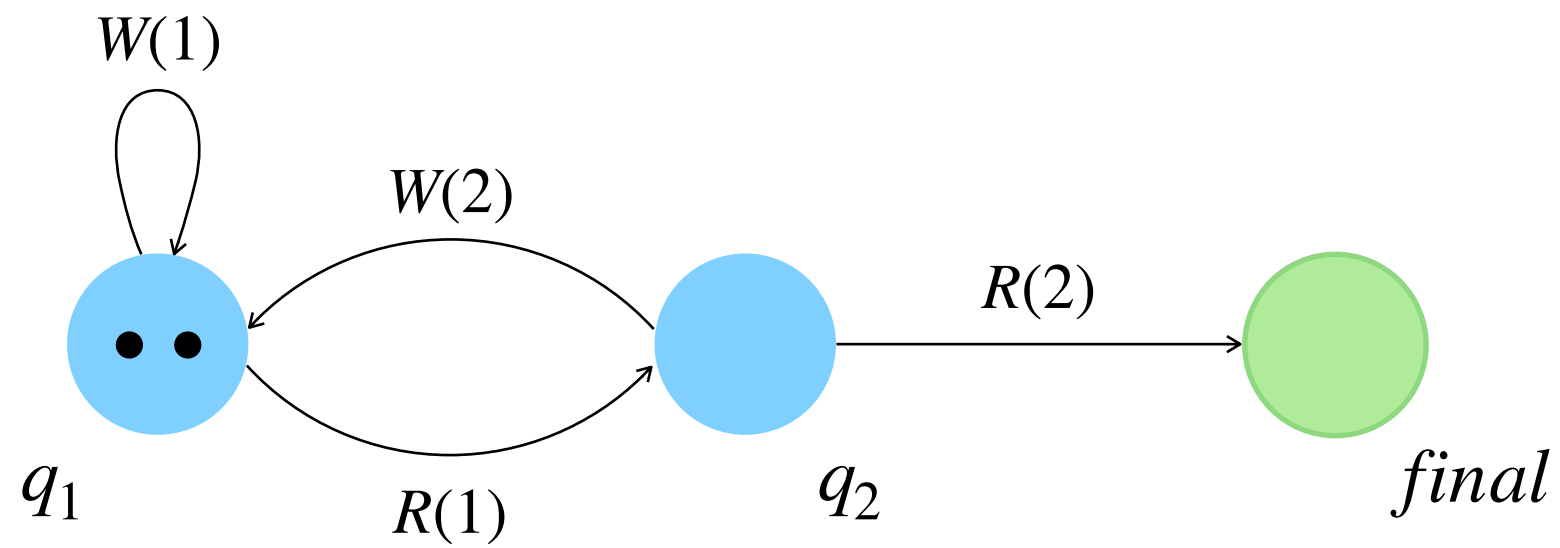




# ASMS

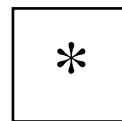
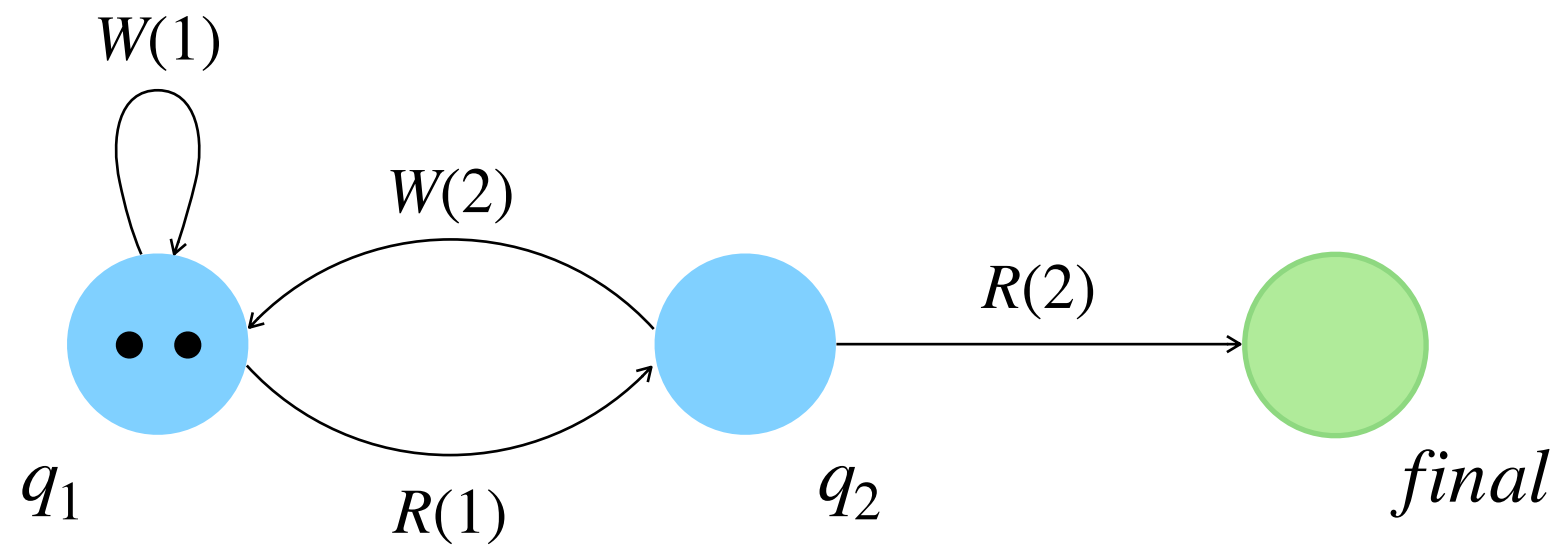


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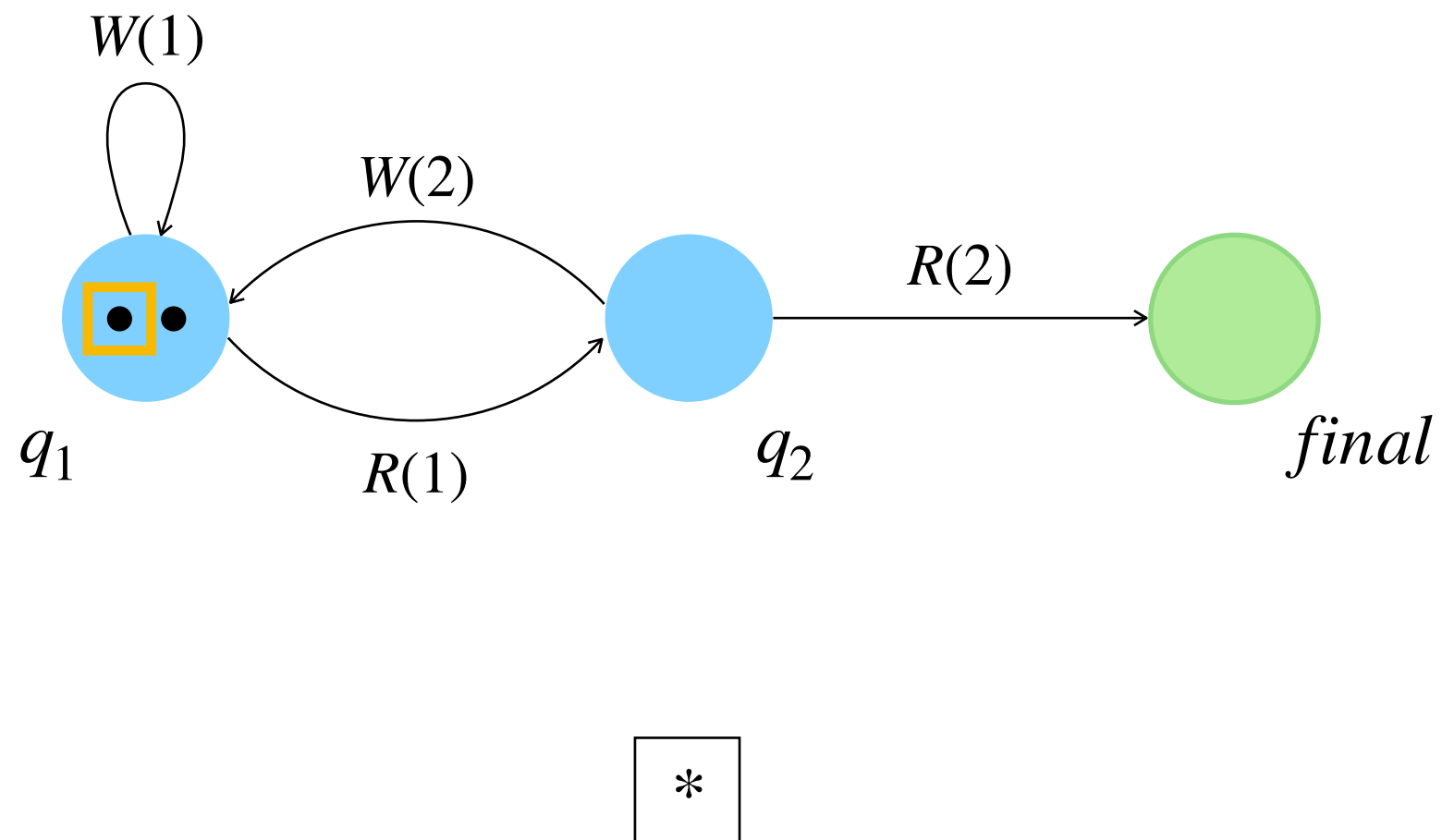
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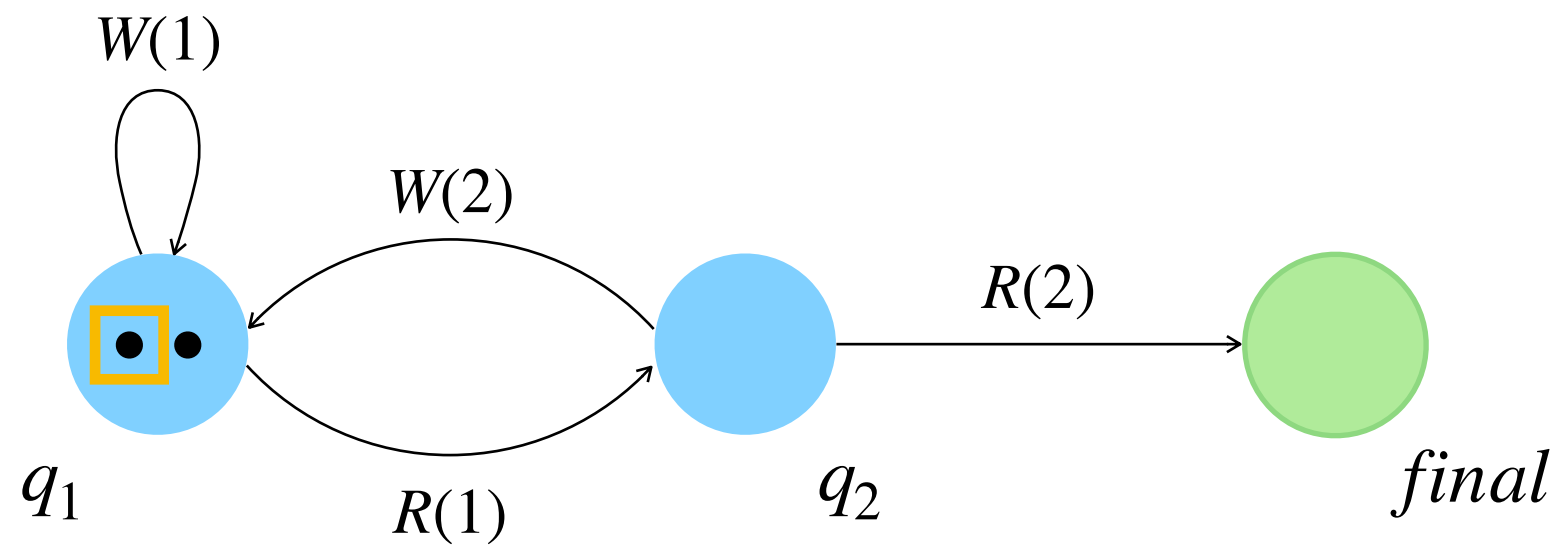
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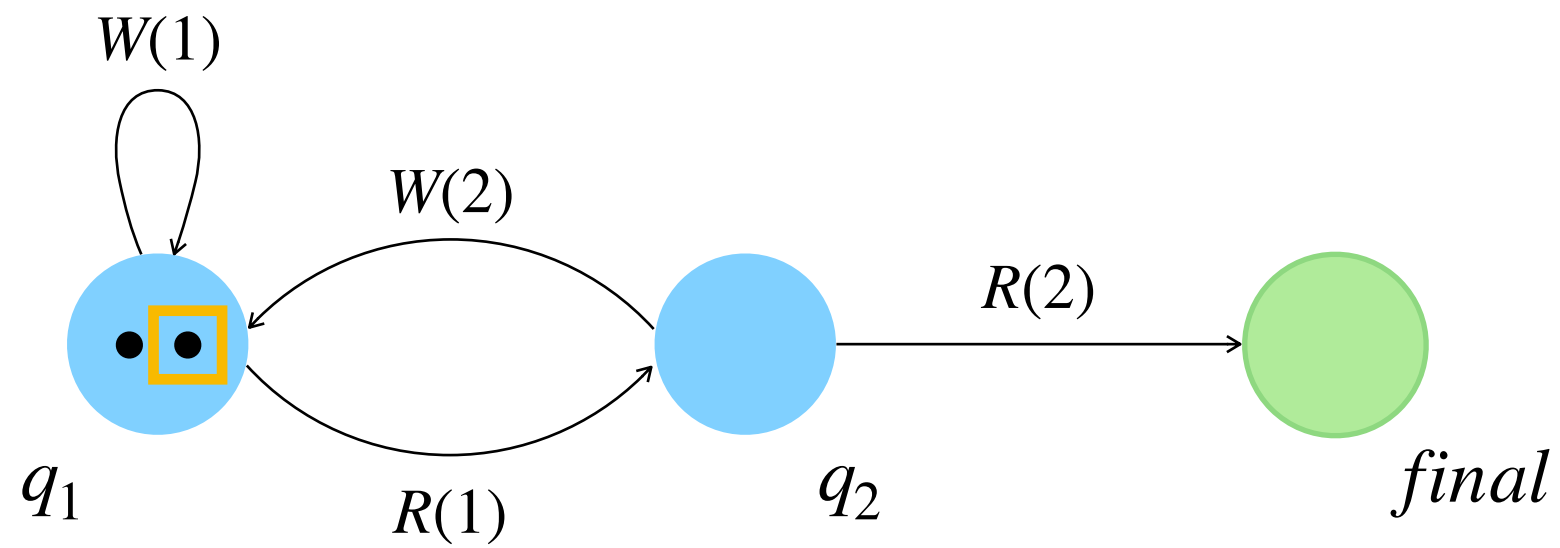
Goal: put a process in *final*



1

# ASMS

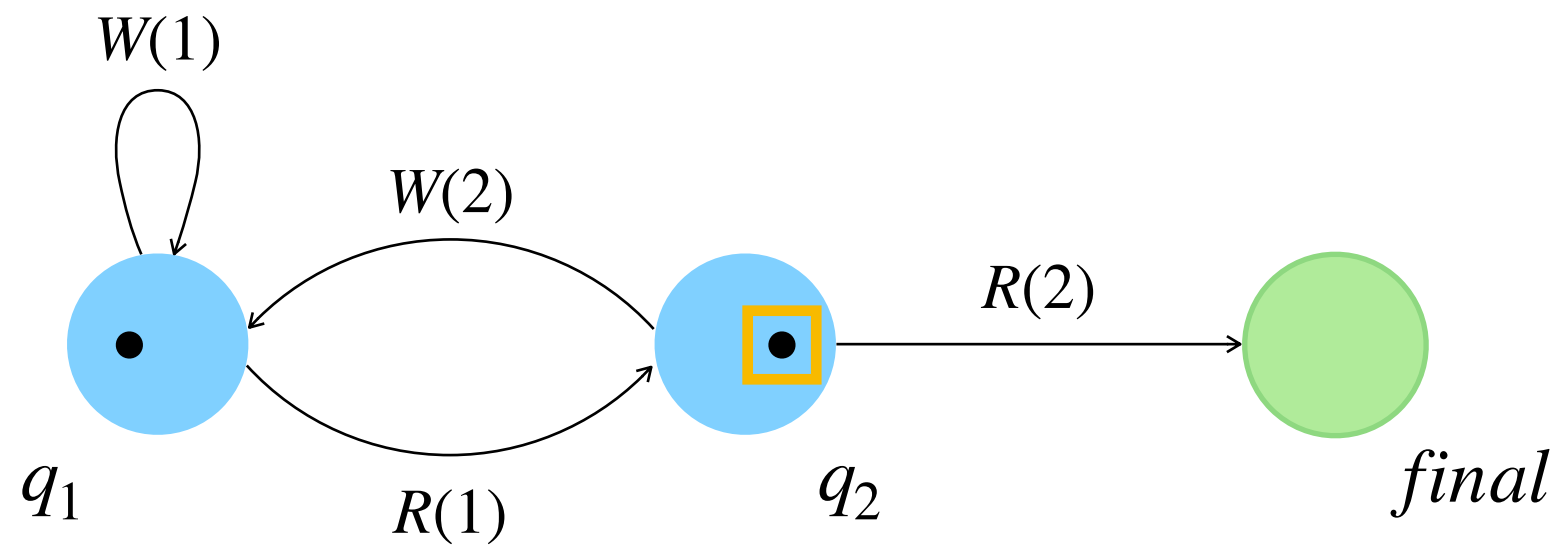
Goal: put a process in *final*



1

# ASMS

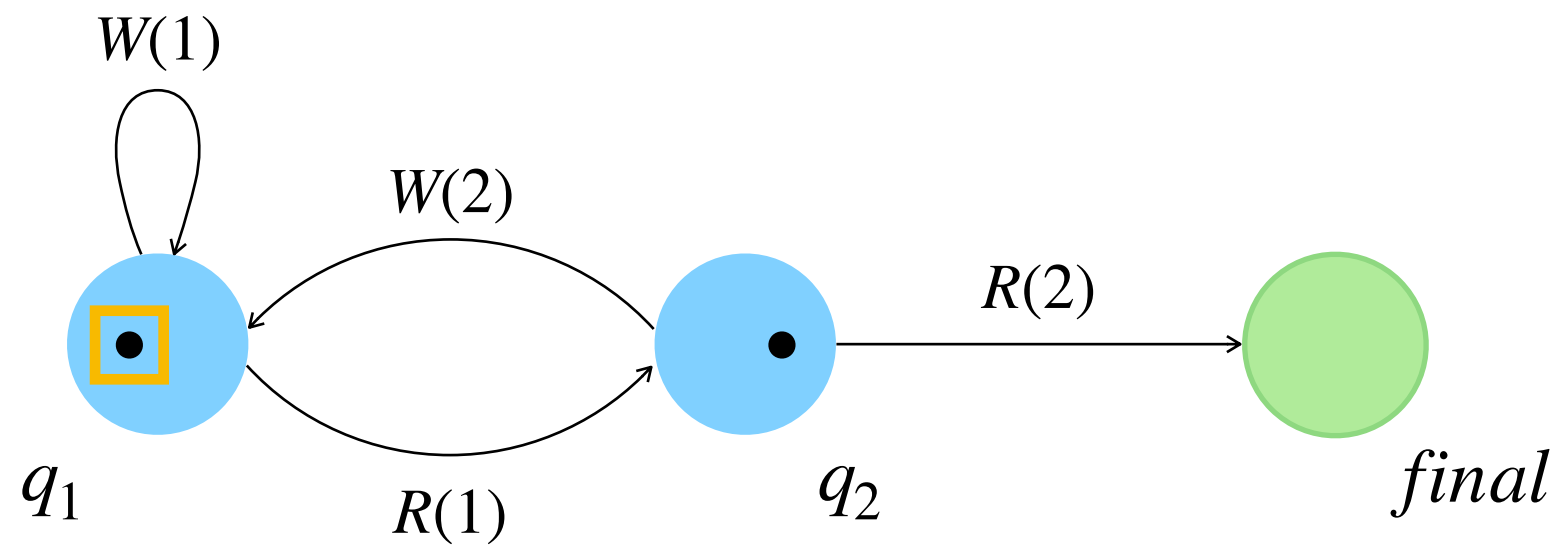
Goal: put a process in *final*



1

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Goal: put a process in *final*

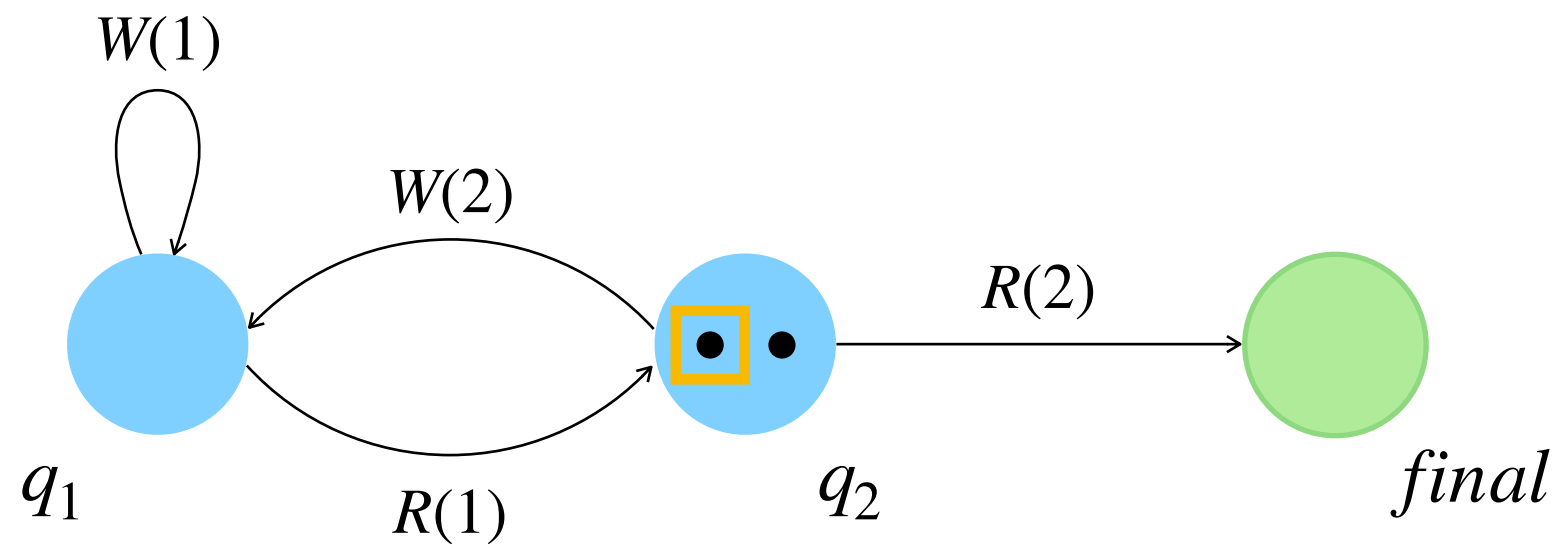


1



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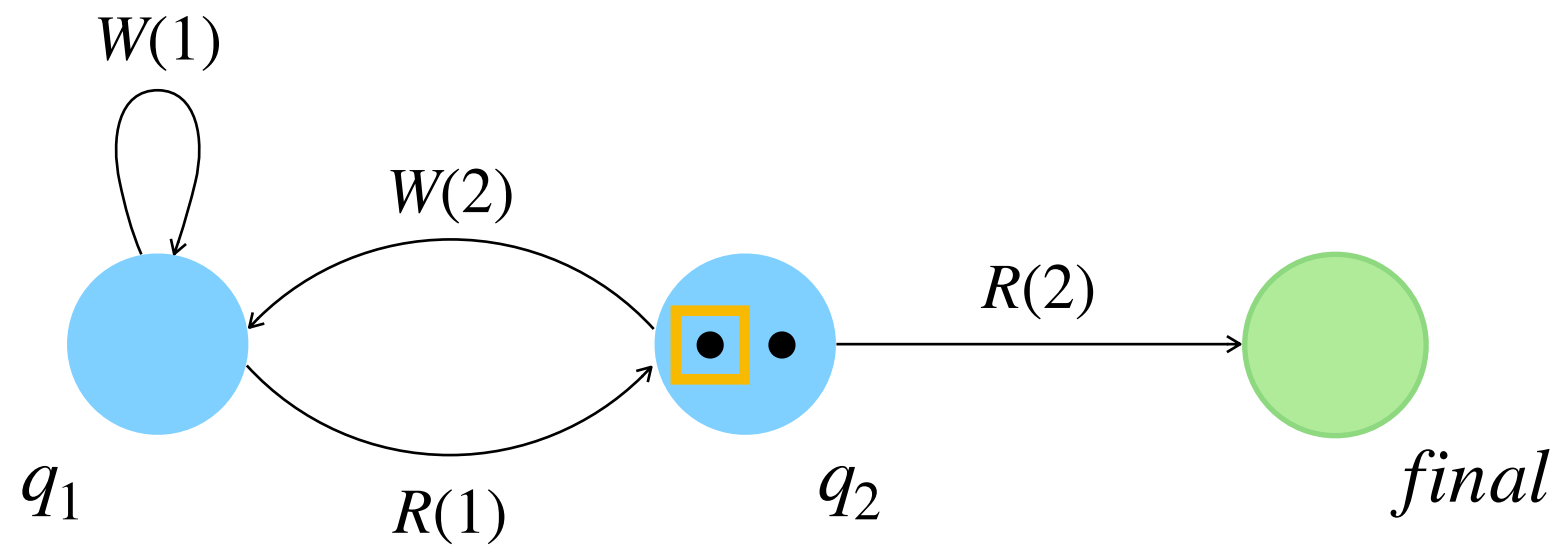
Goal: put a process in *final*



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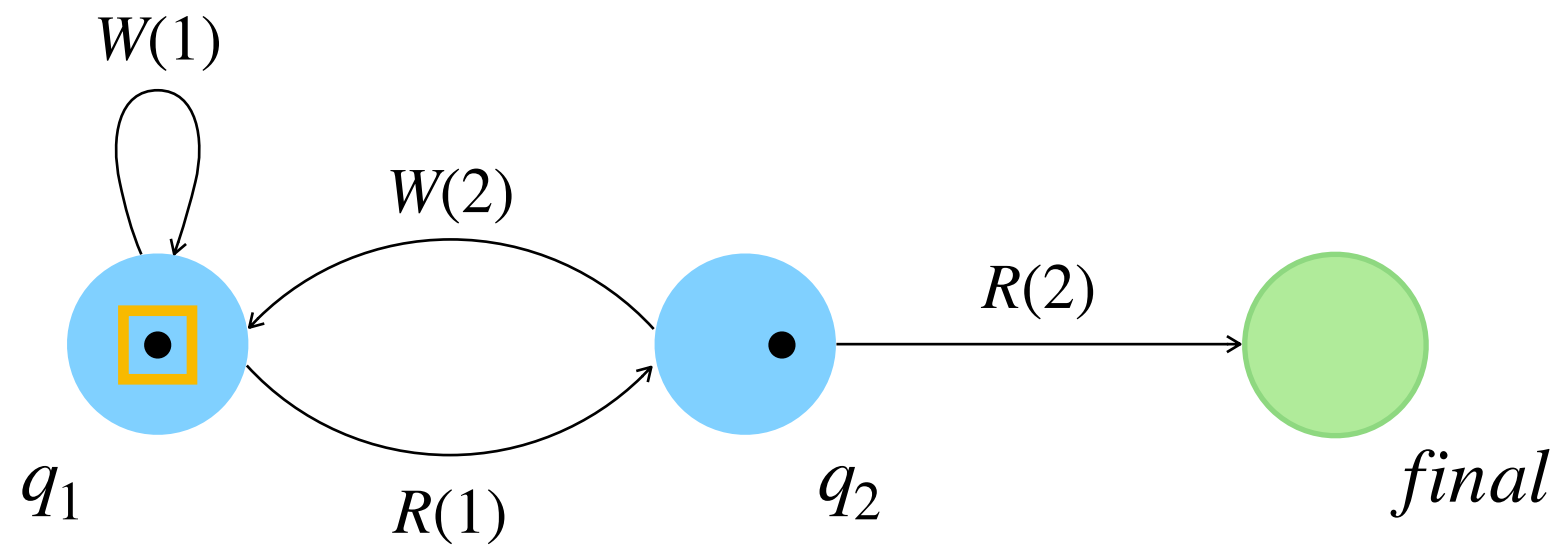
Goal: put a process in *final*



1

# ASMS

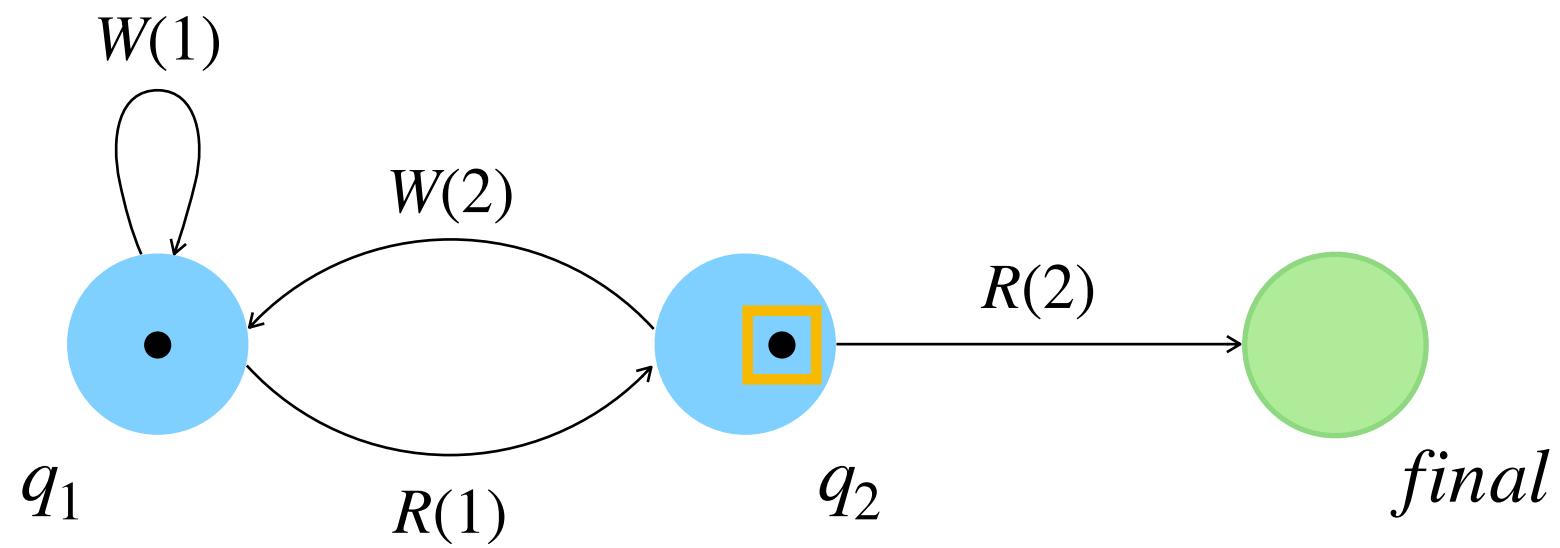
Goal: put a process in *final*



2

# ASMS

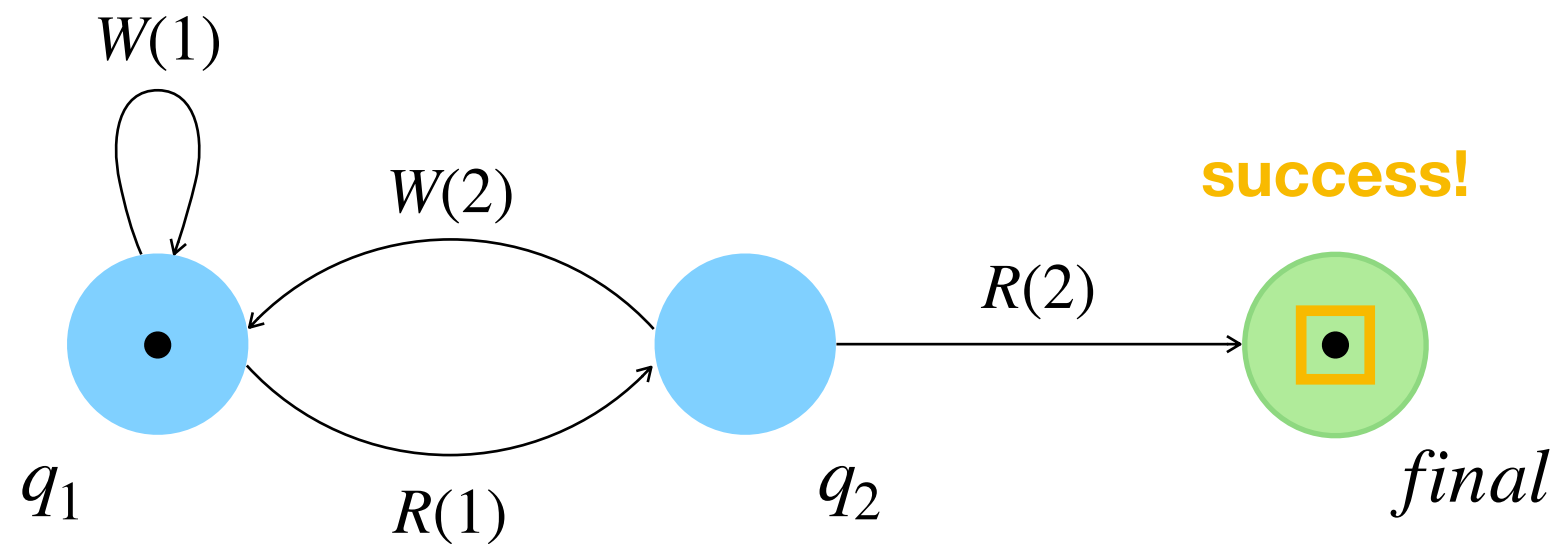
Goal: put a process in *final*



2

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Goal: put a process in *final*



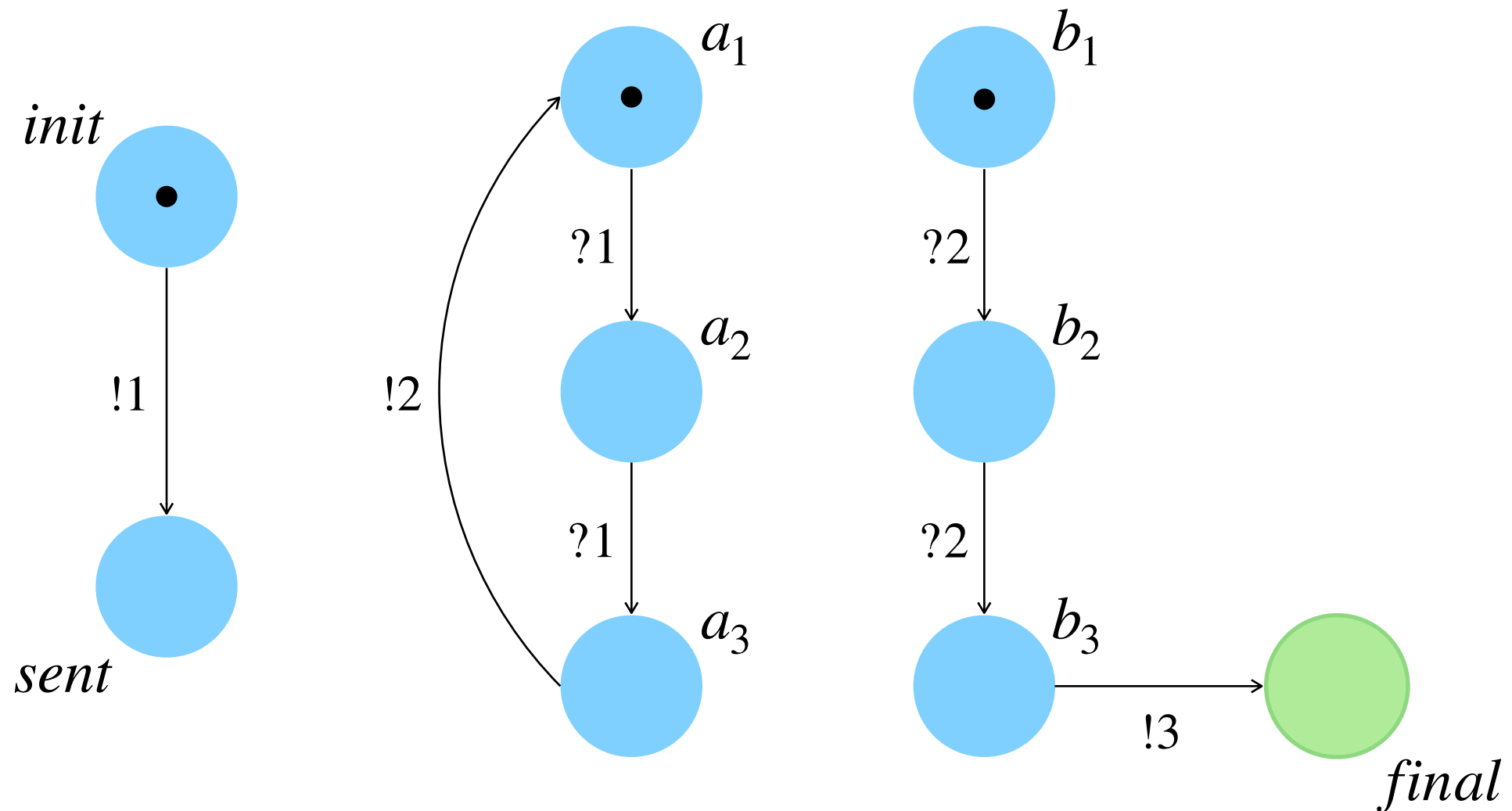
2

# ASMS simulates RBN

Simulation of polynomial-size with bijection between configurations

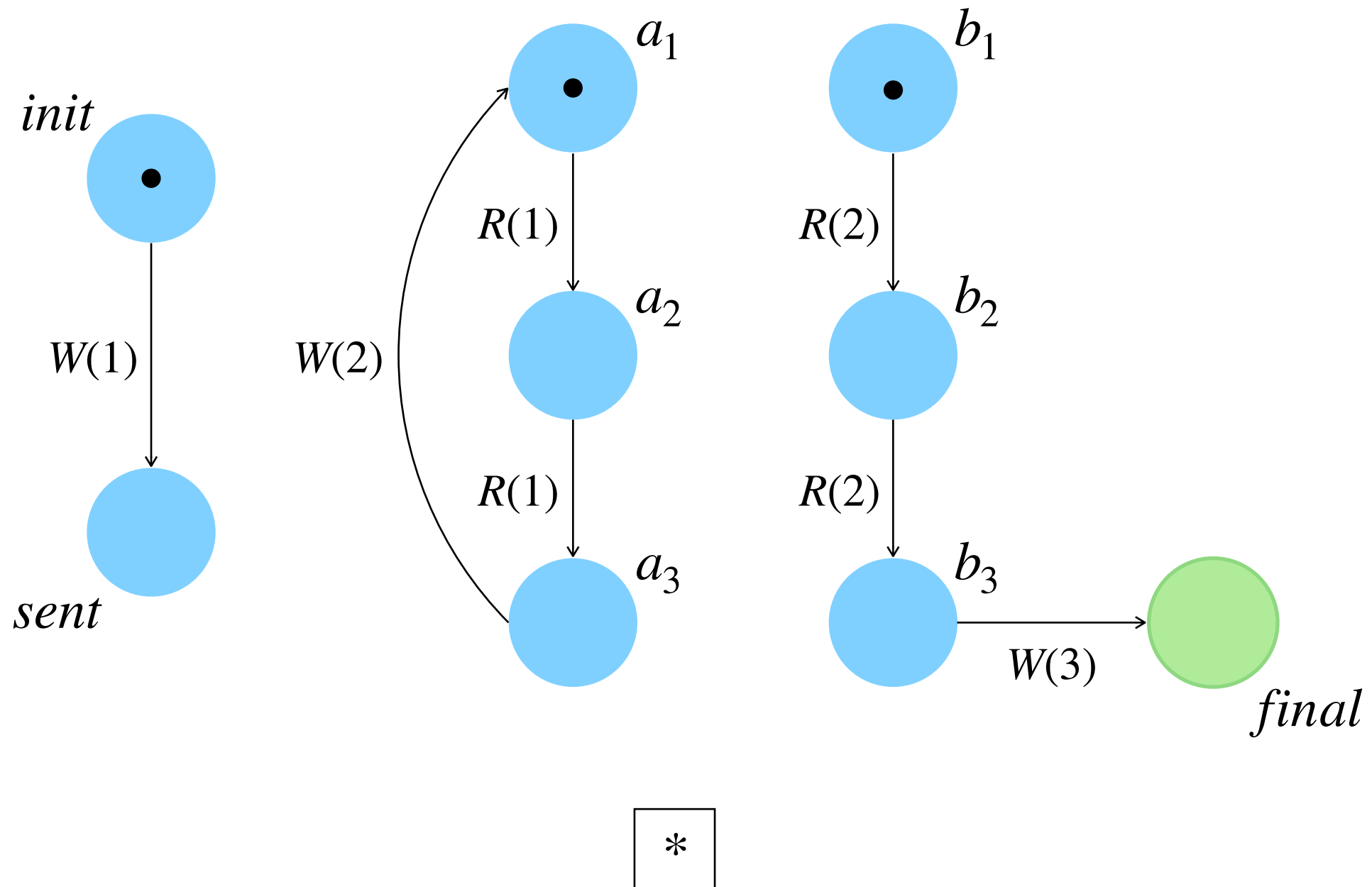
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Simulation of polynomial-size with bijection between configurations



# ASMS simulates RBN

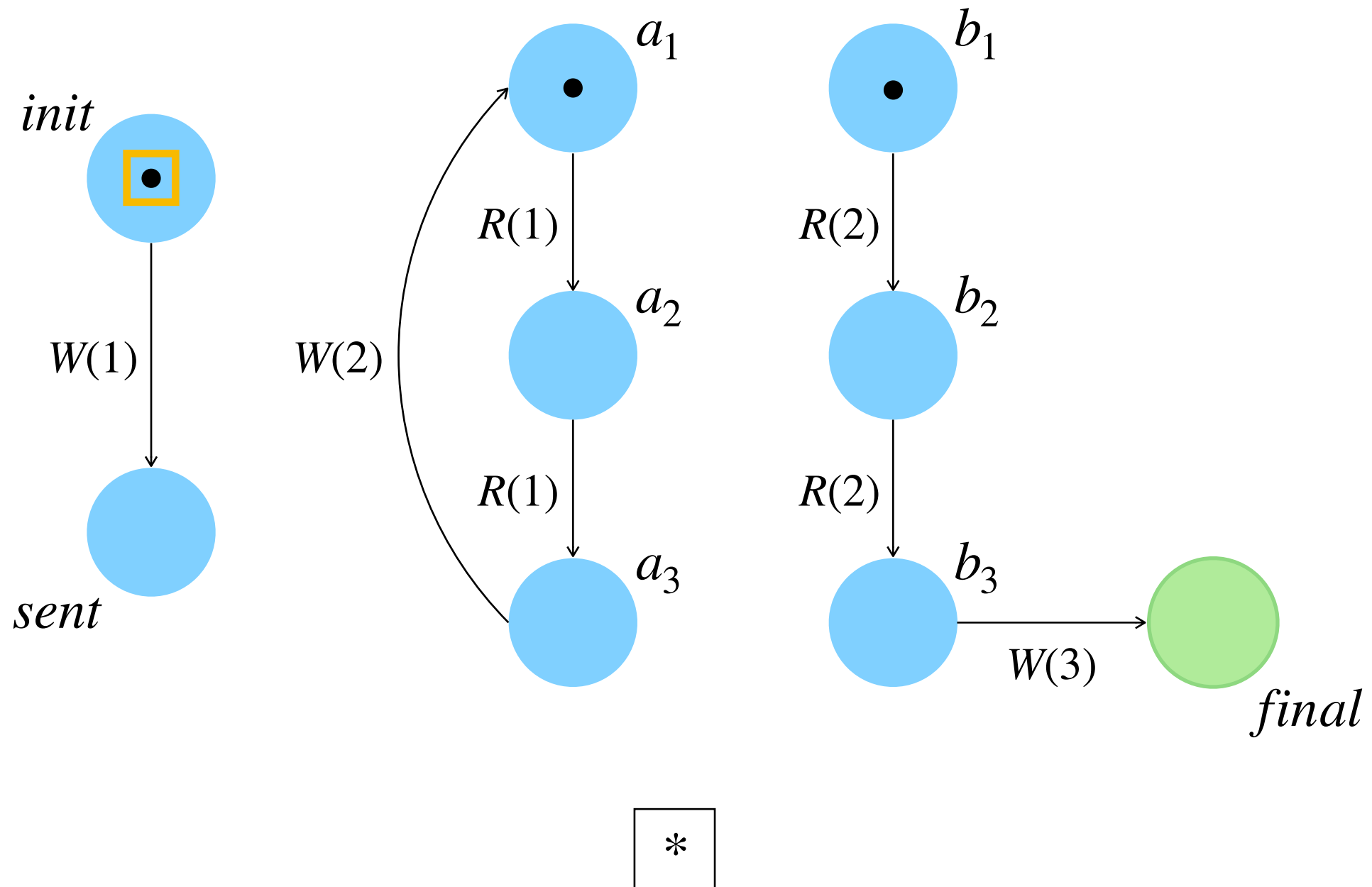
## Simulation idea #1





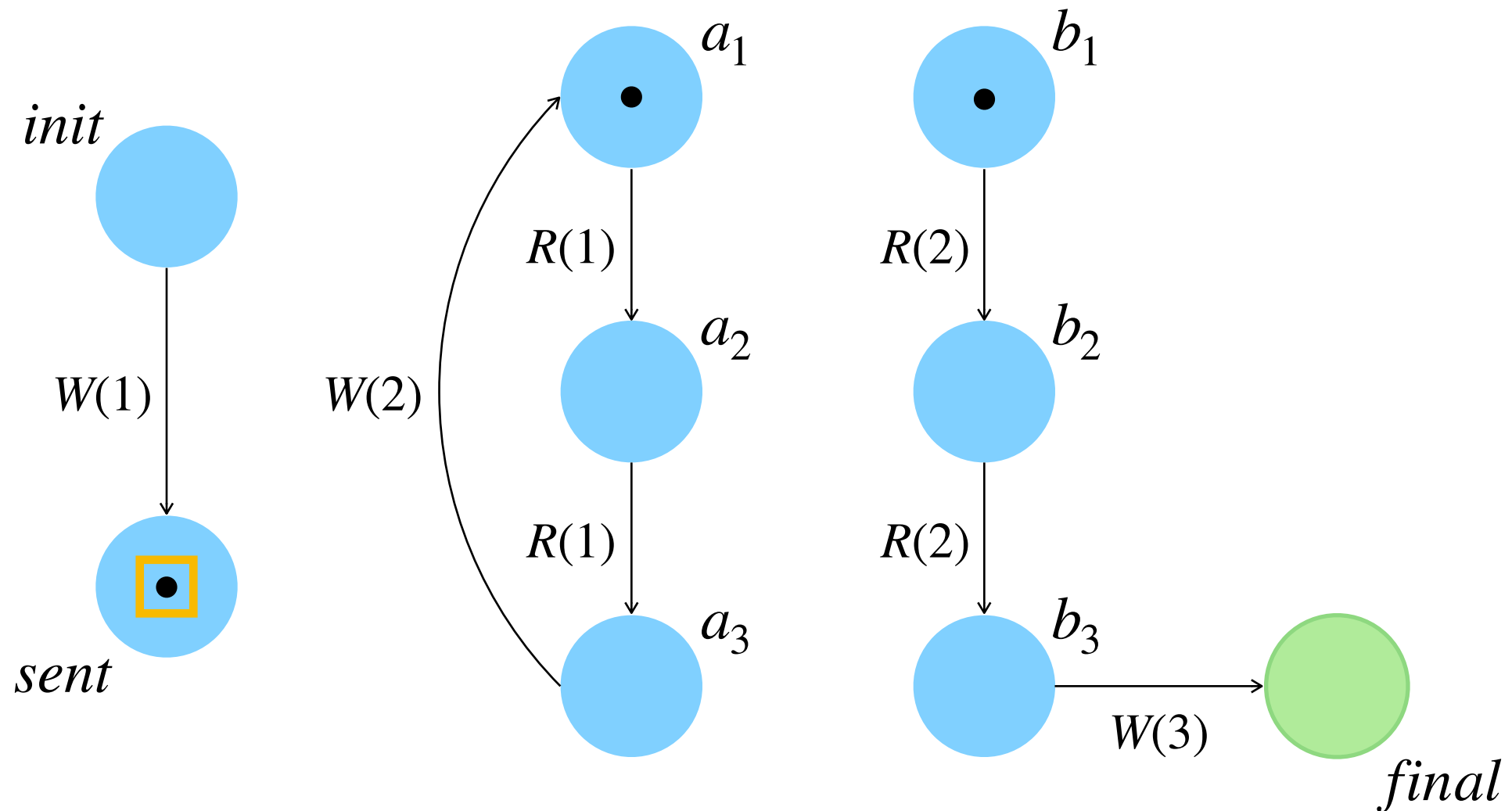
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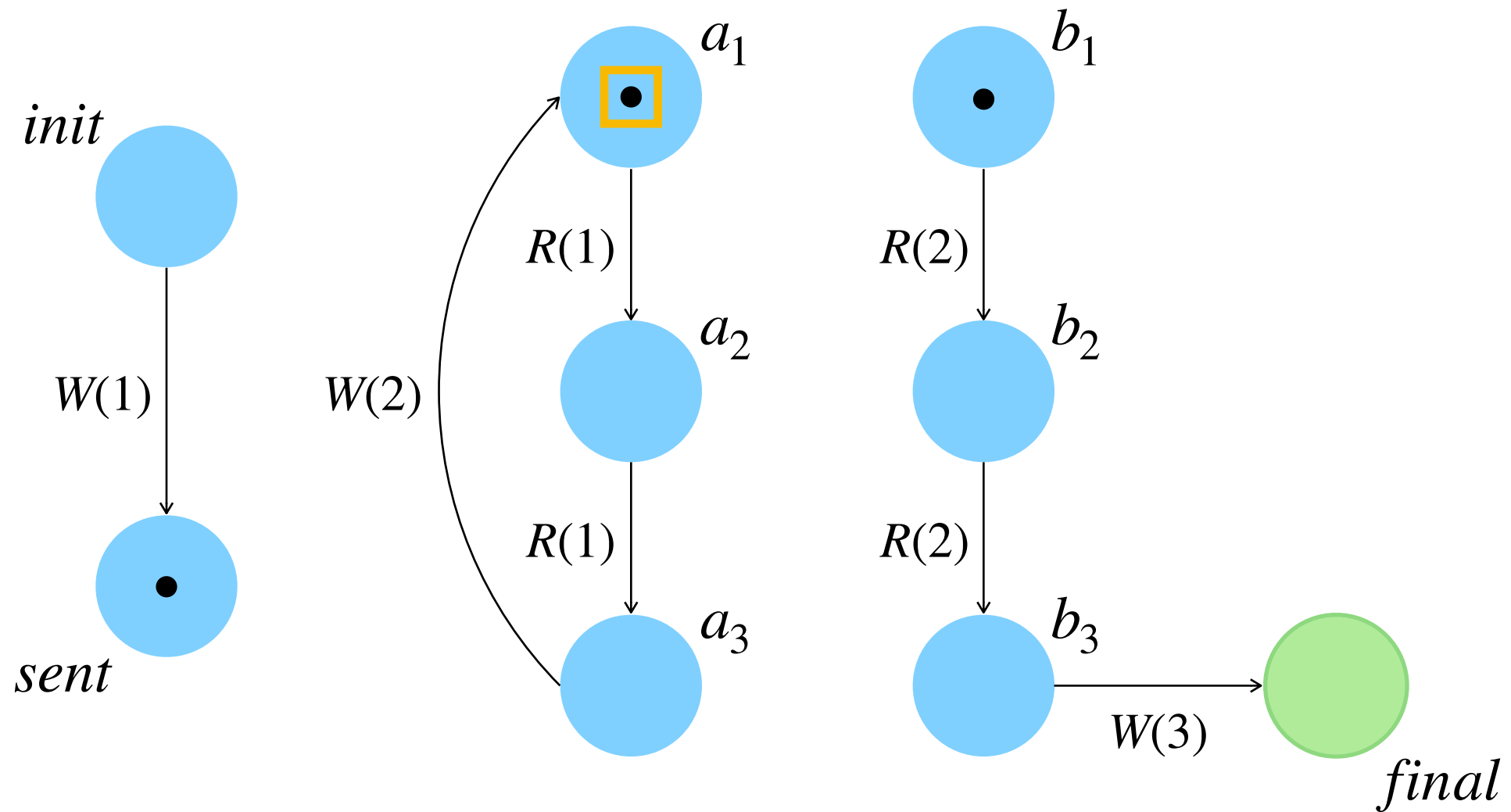
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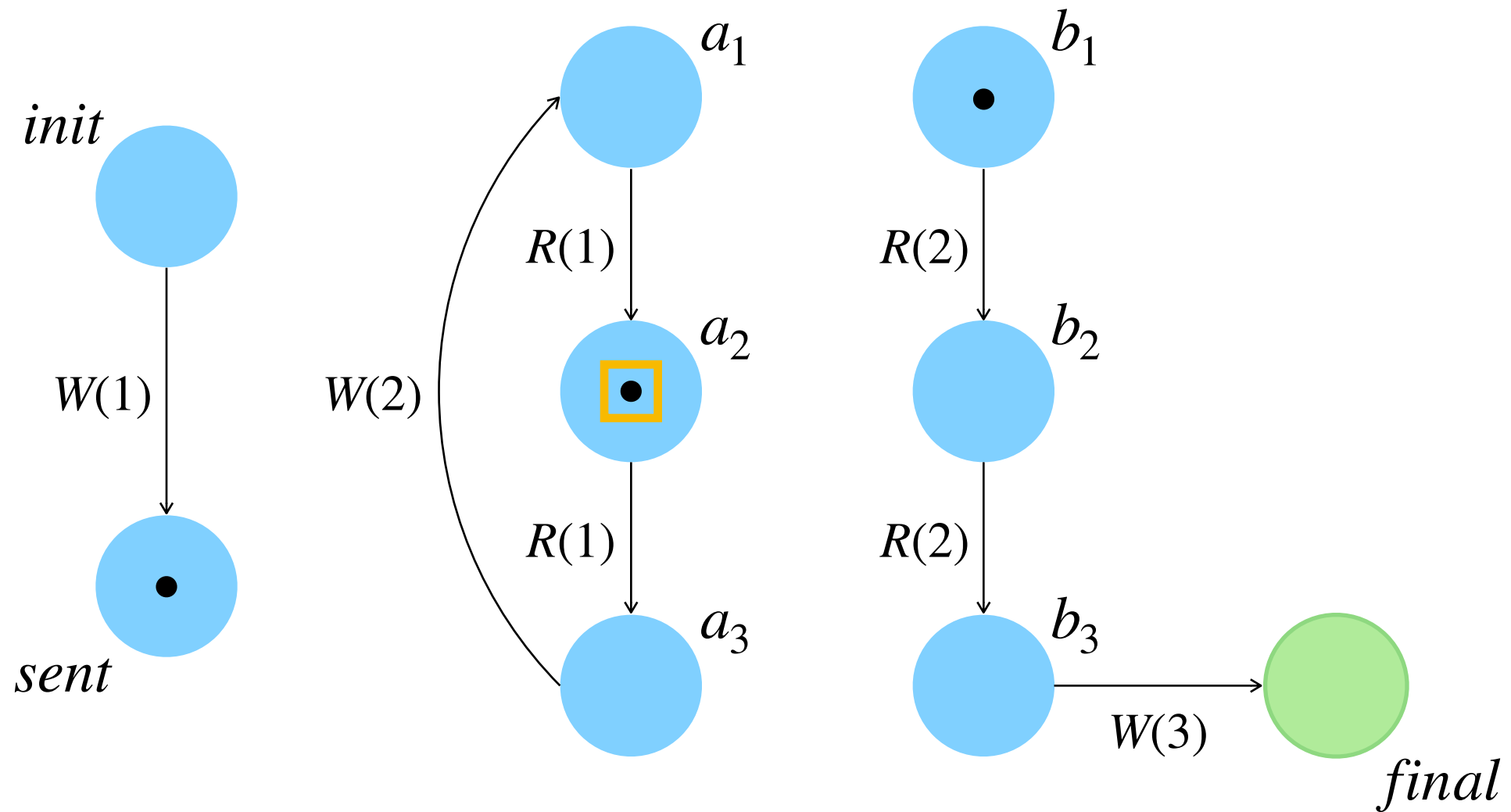
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1

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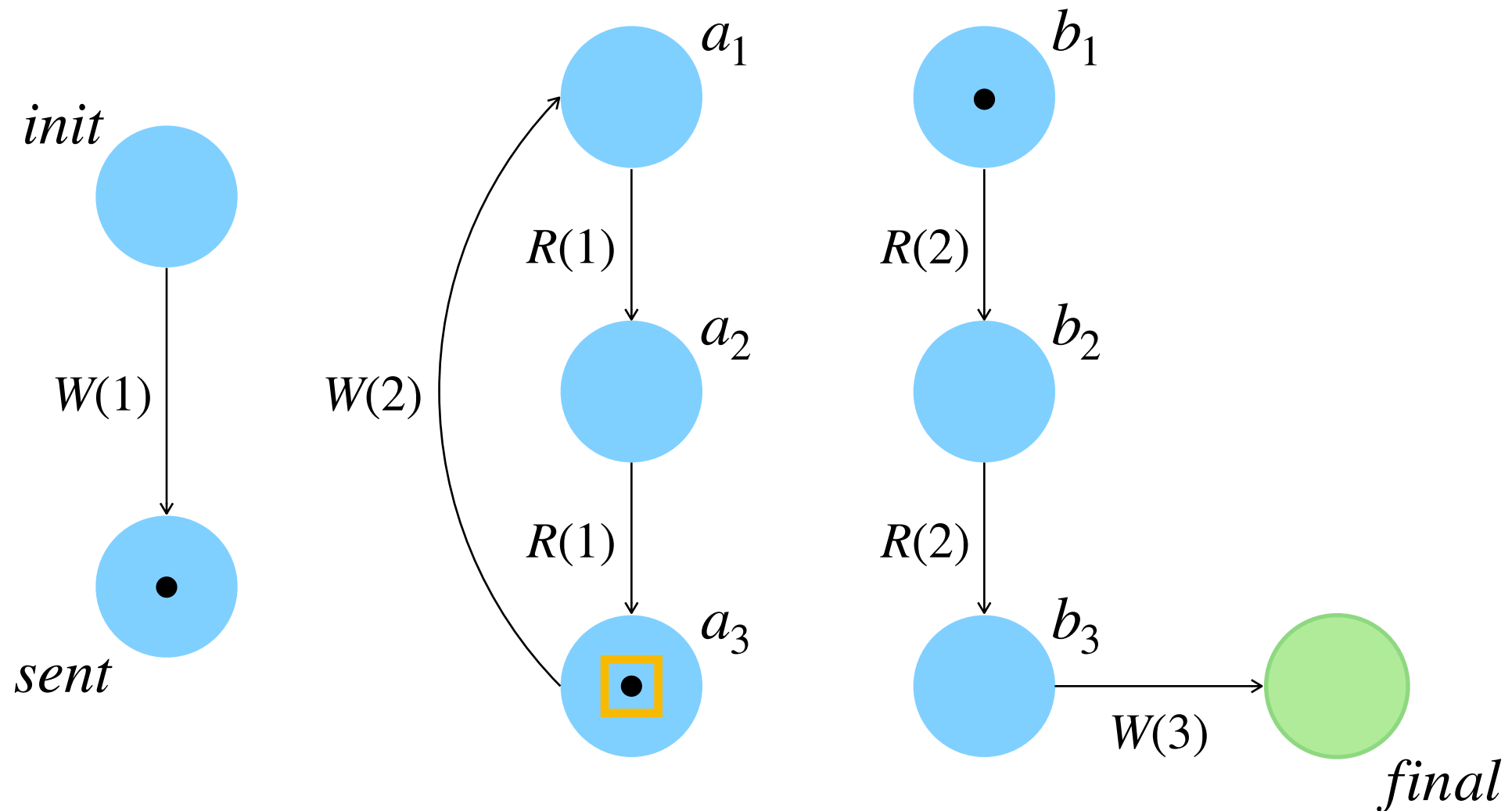
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1

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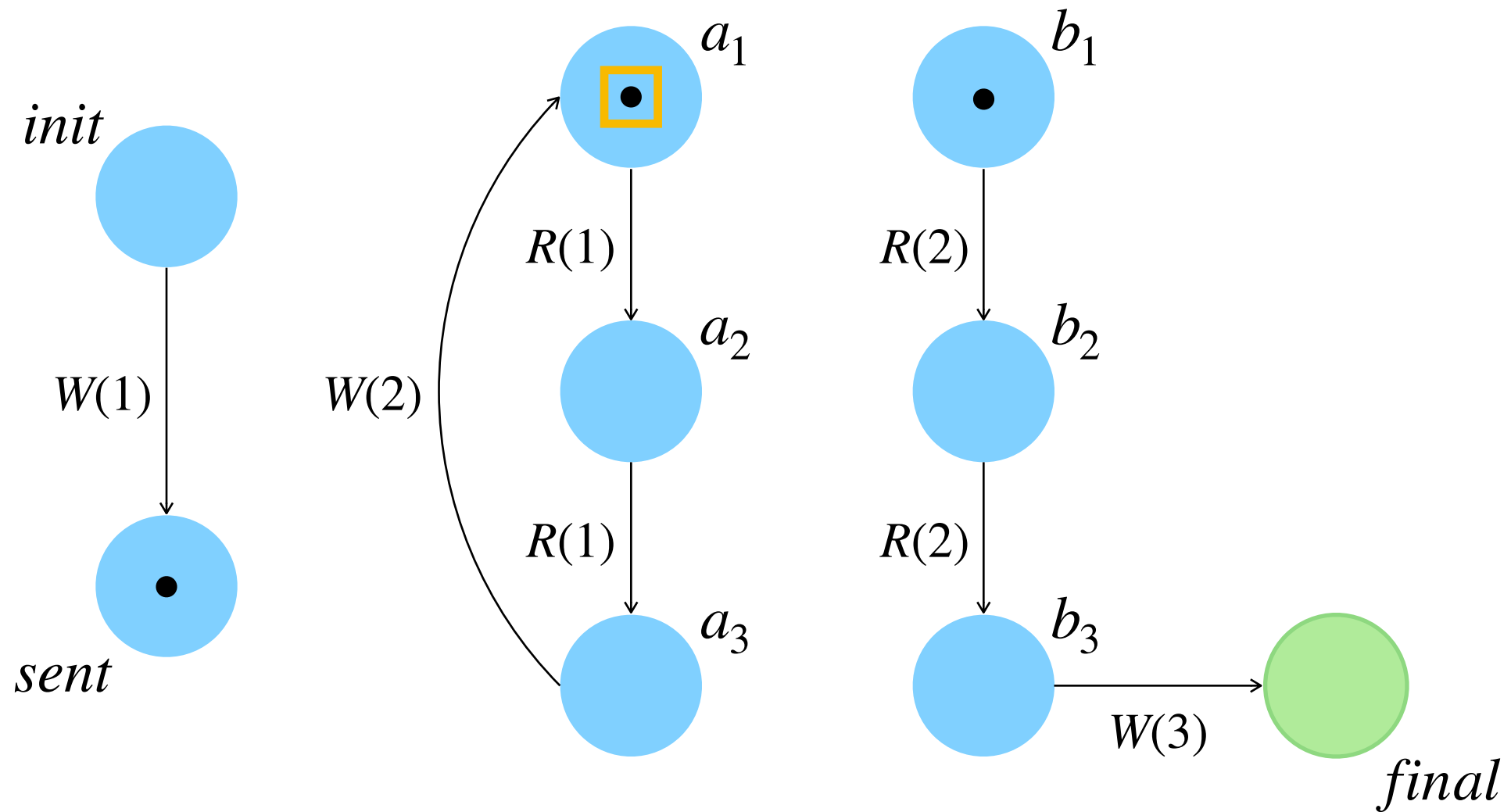
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1

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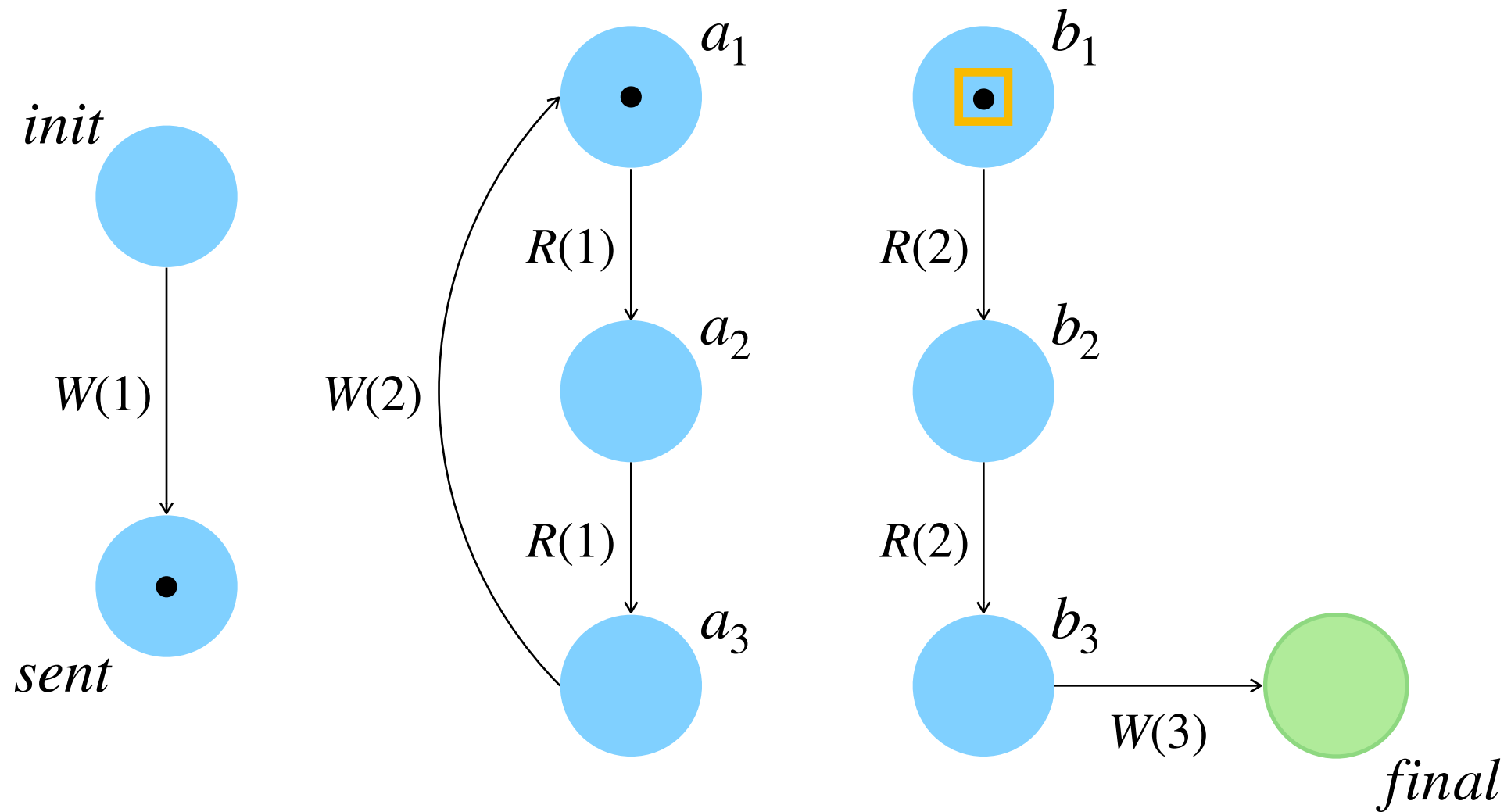
## Simulation idea #1



2

# ASMS simulates RBN

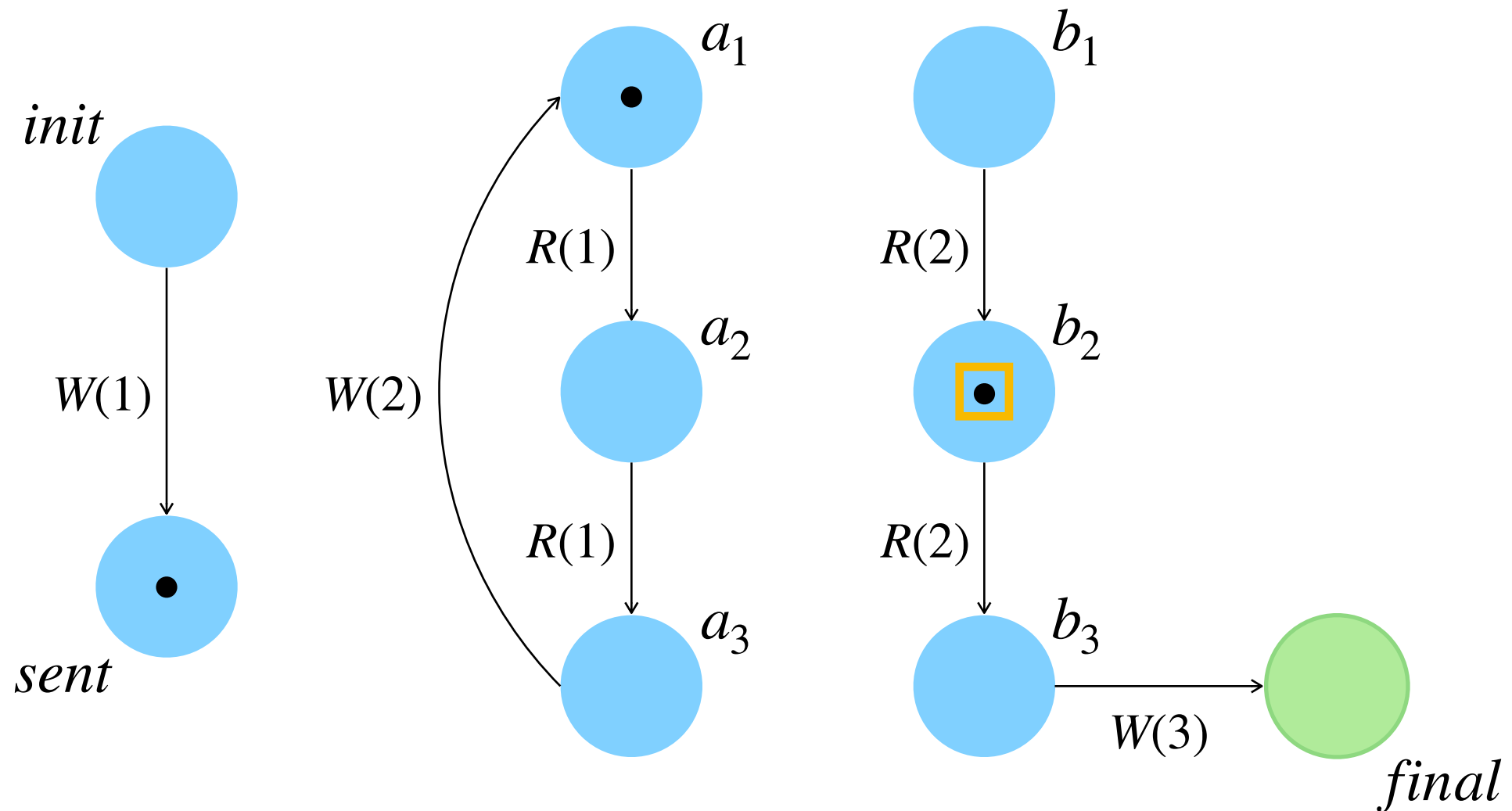
## Simulation idea #1



2

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## Simulation idea #1

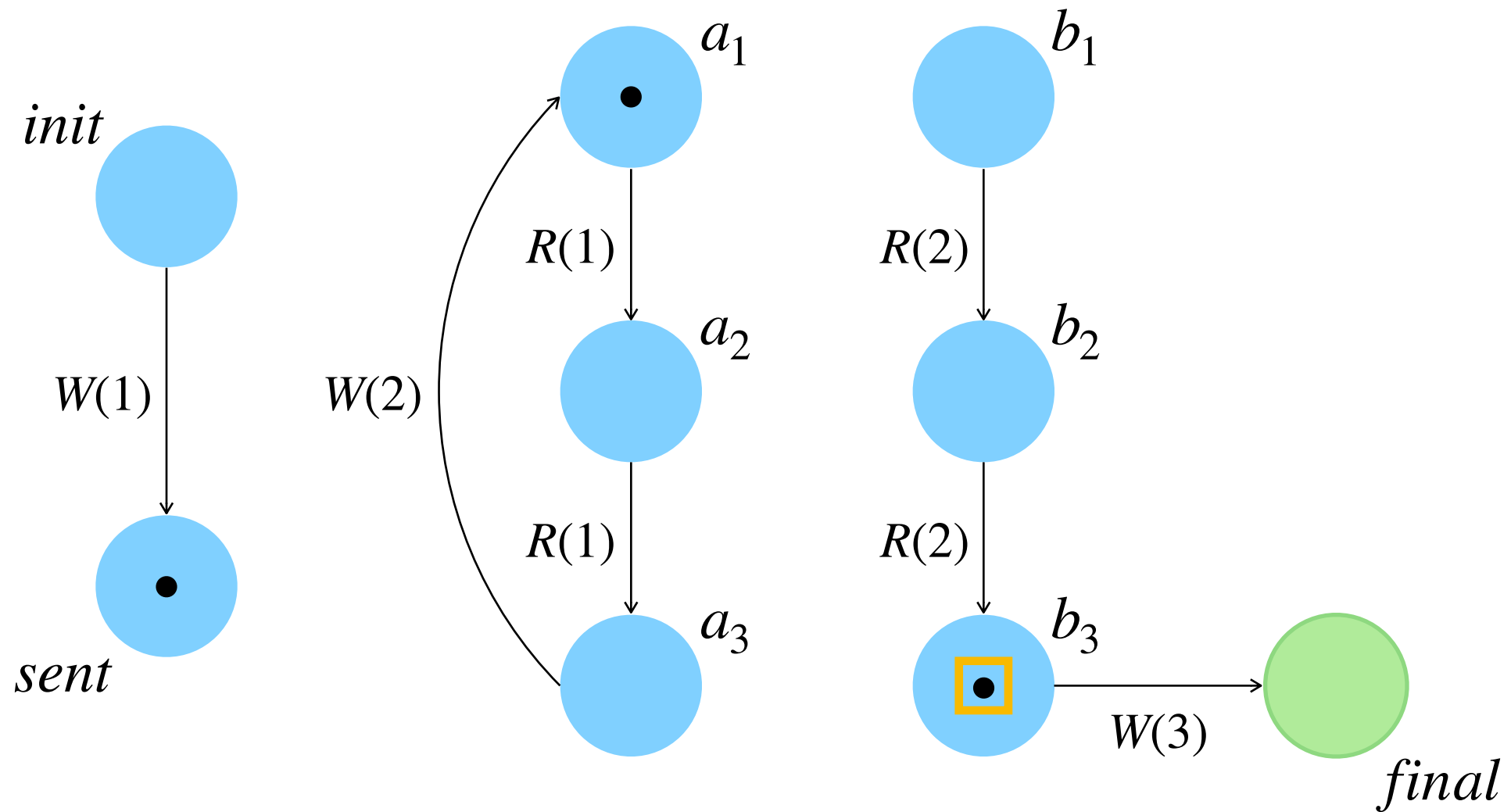


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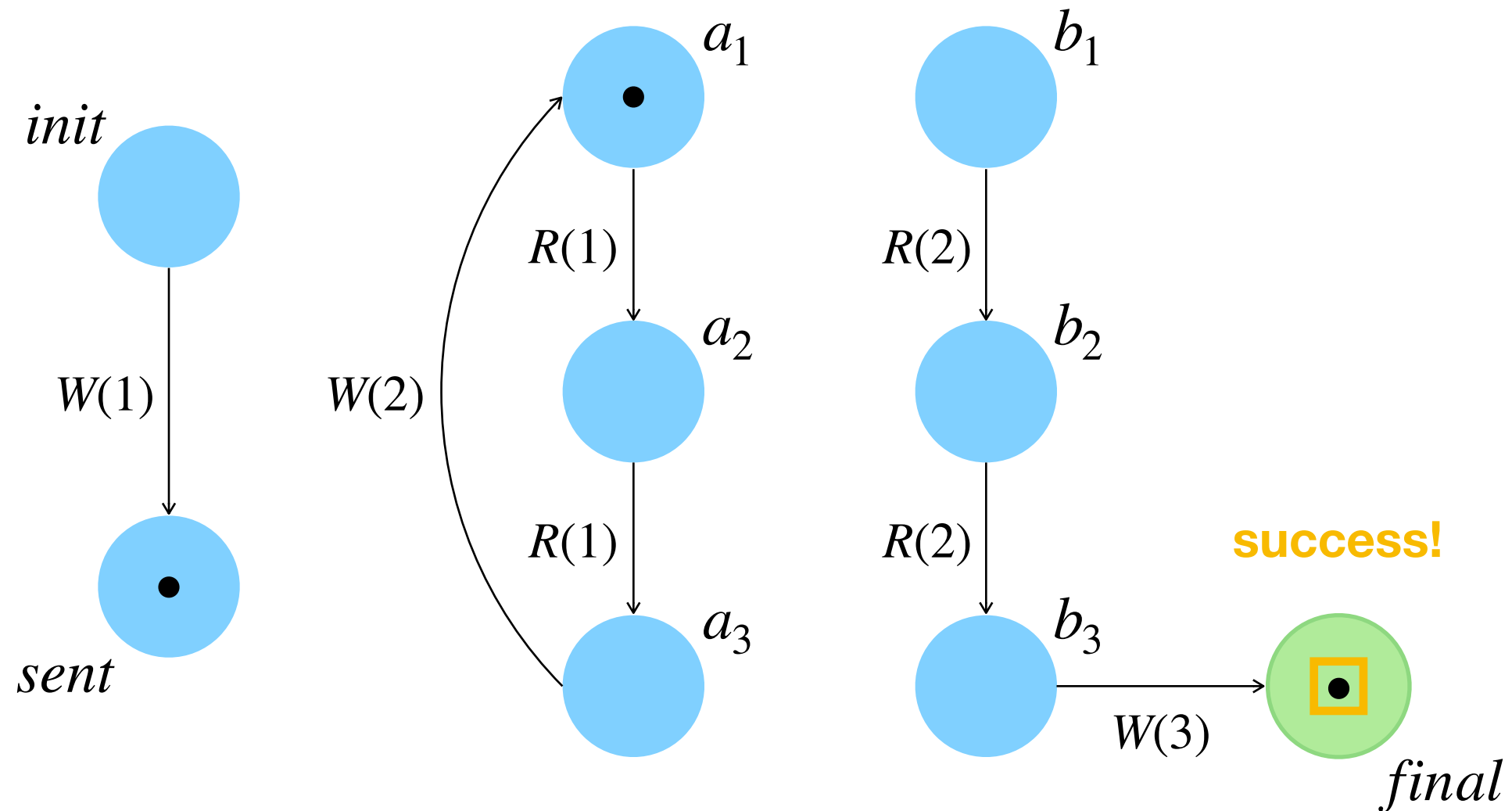
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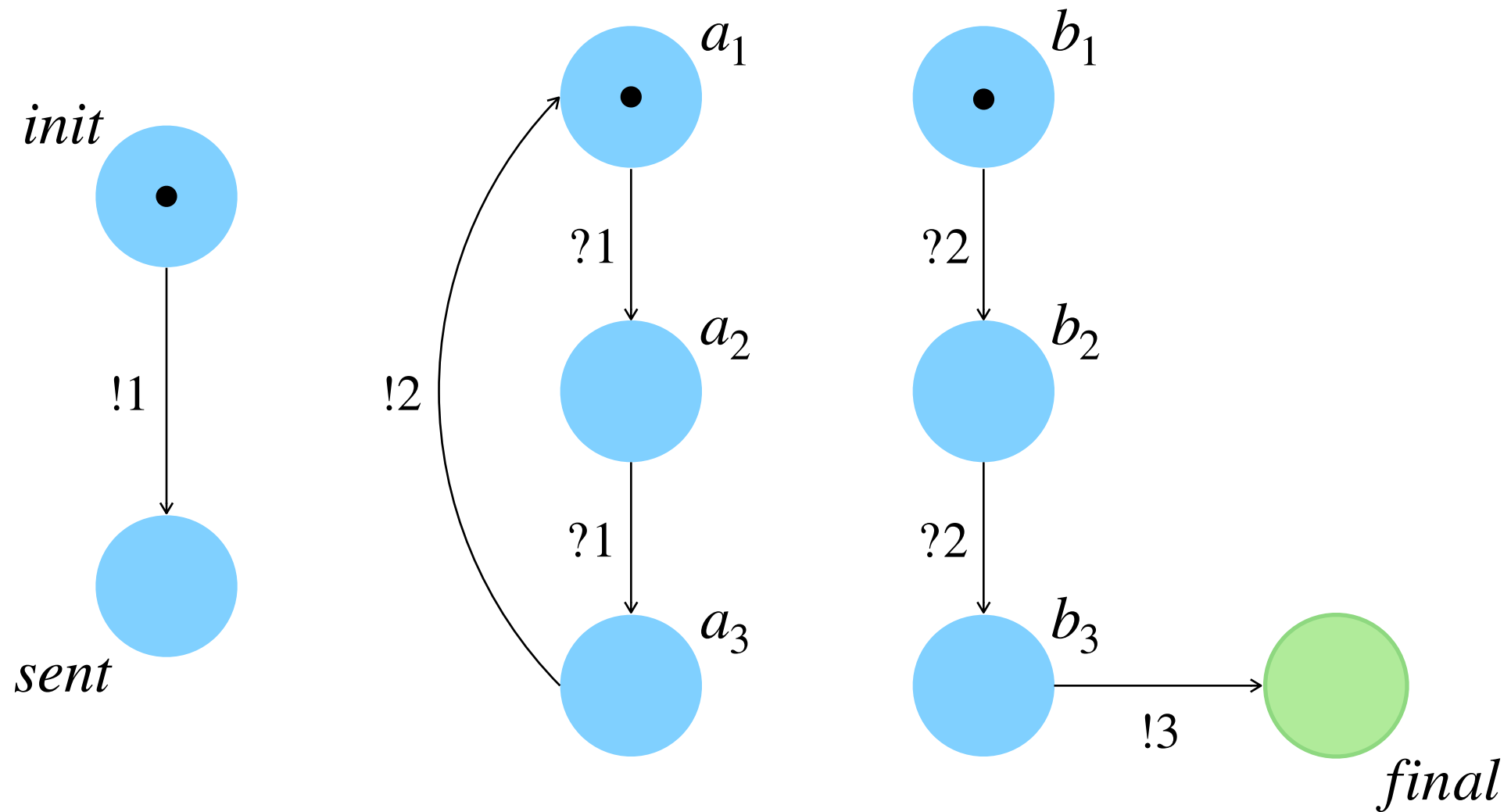
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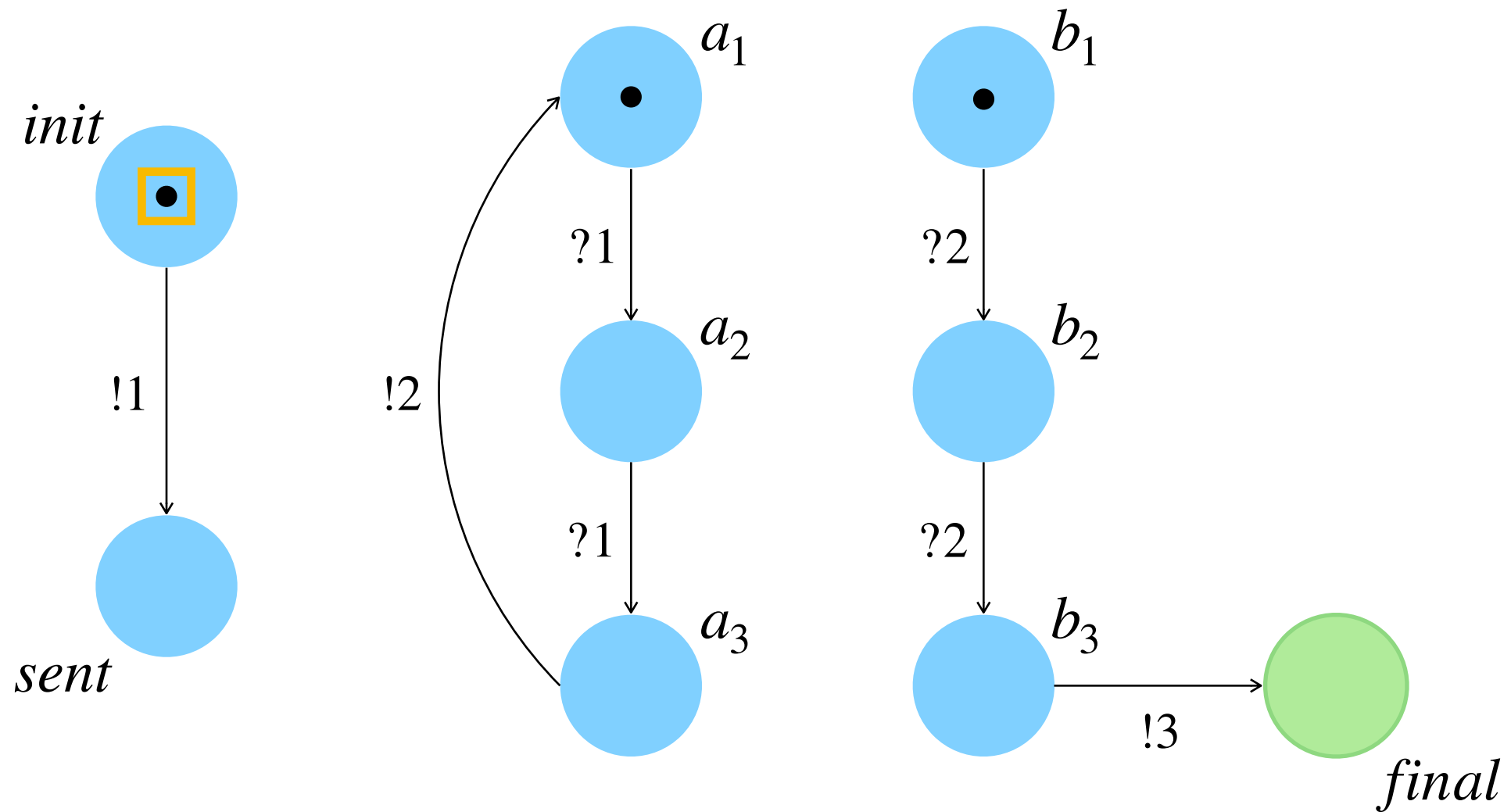


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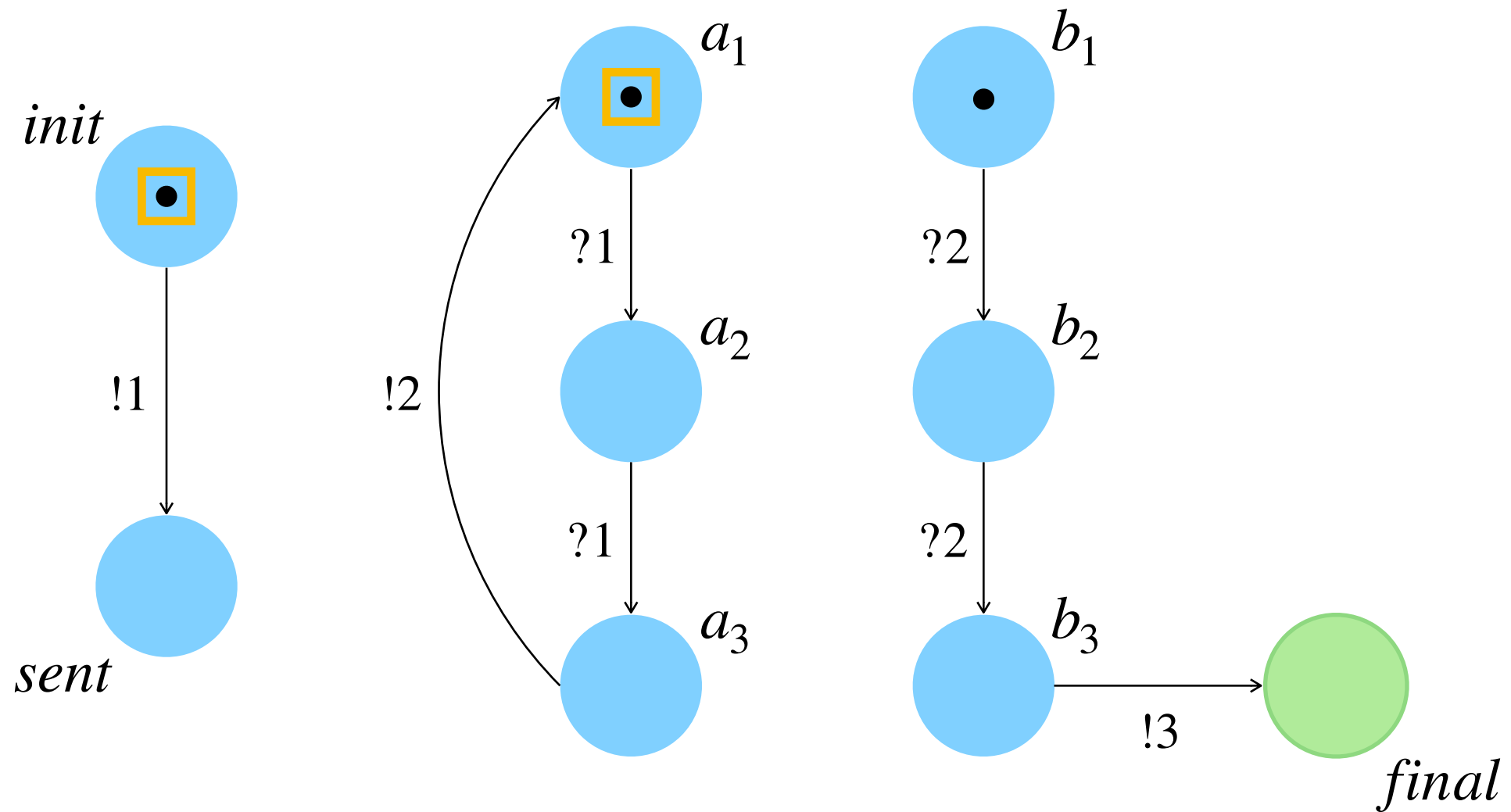
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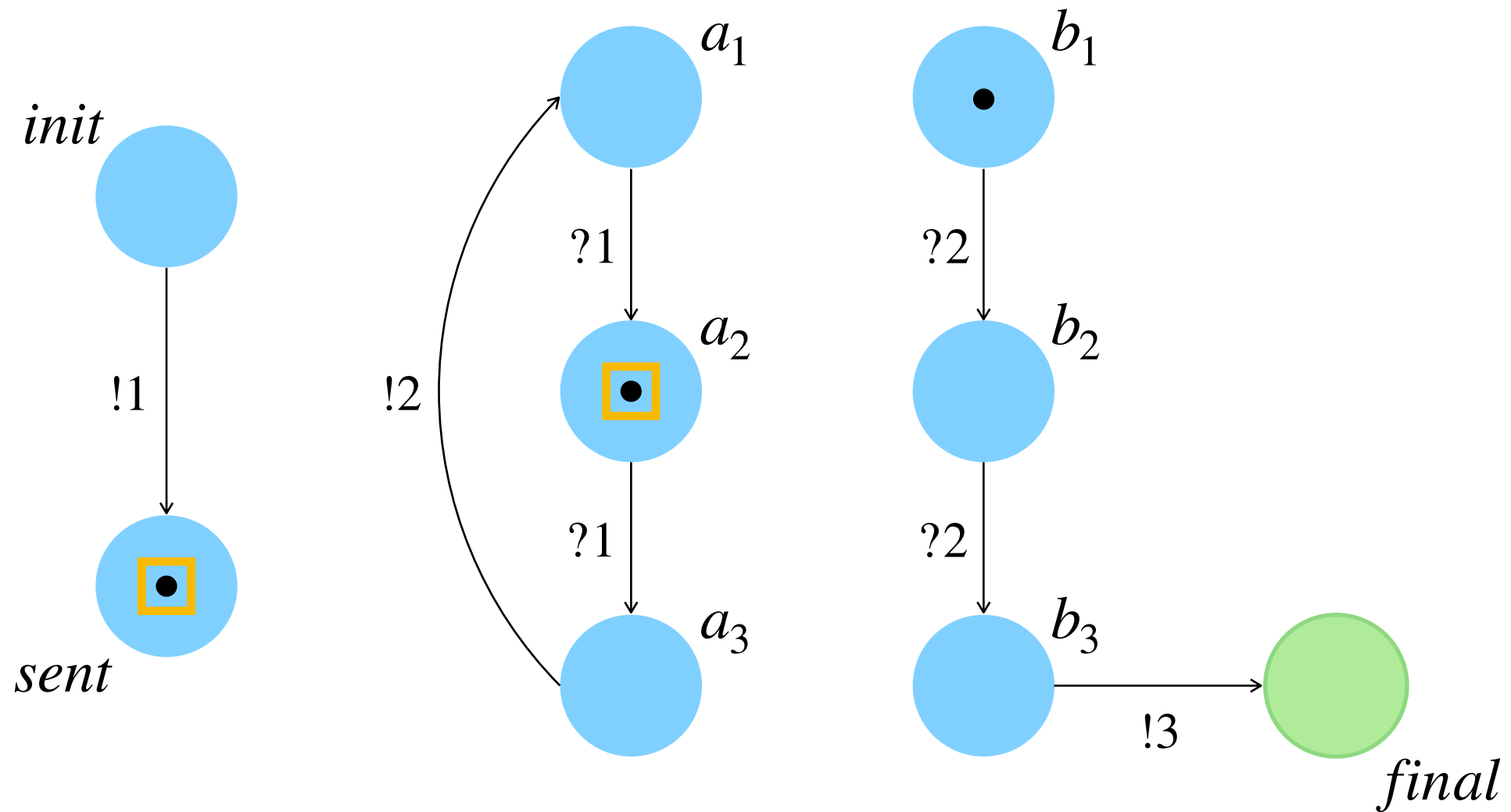
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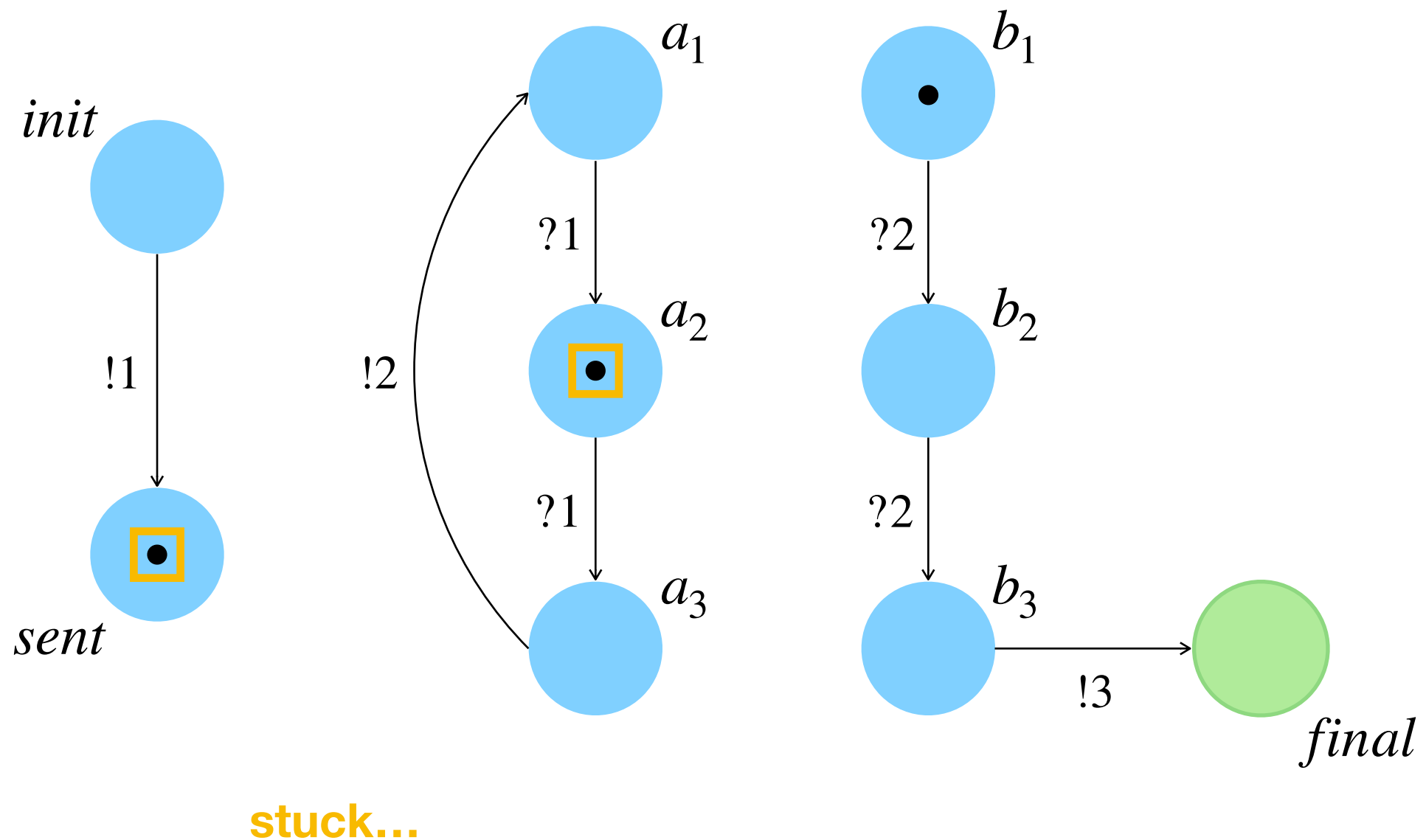
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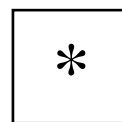
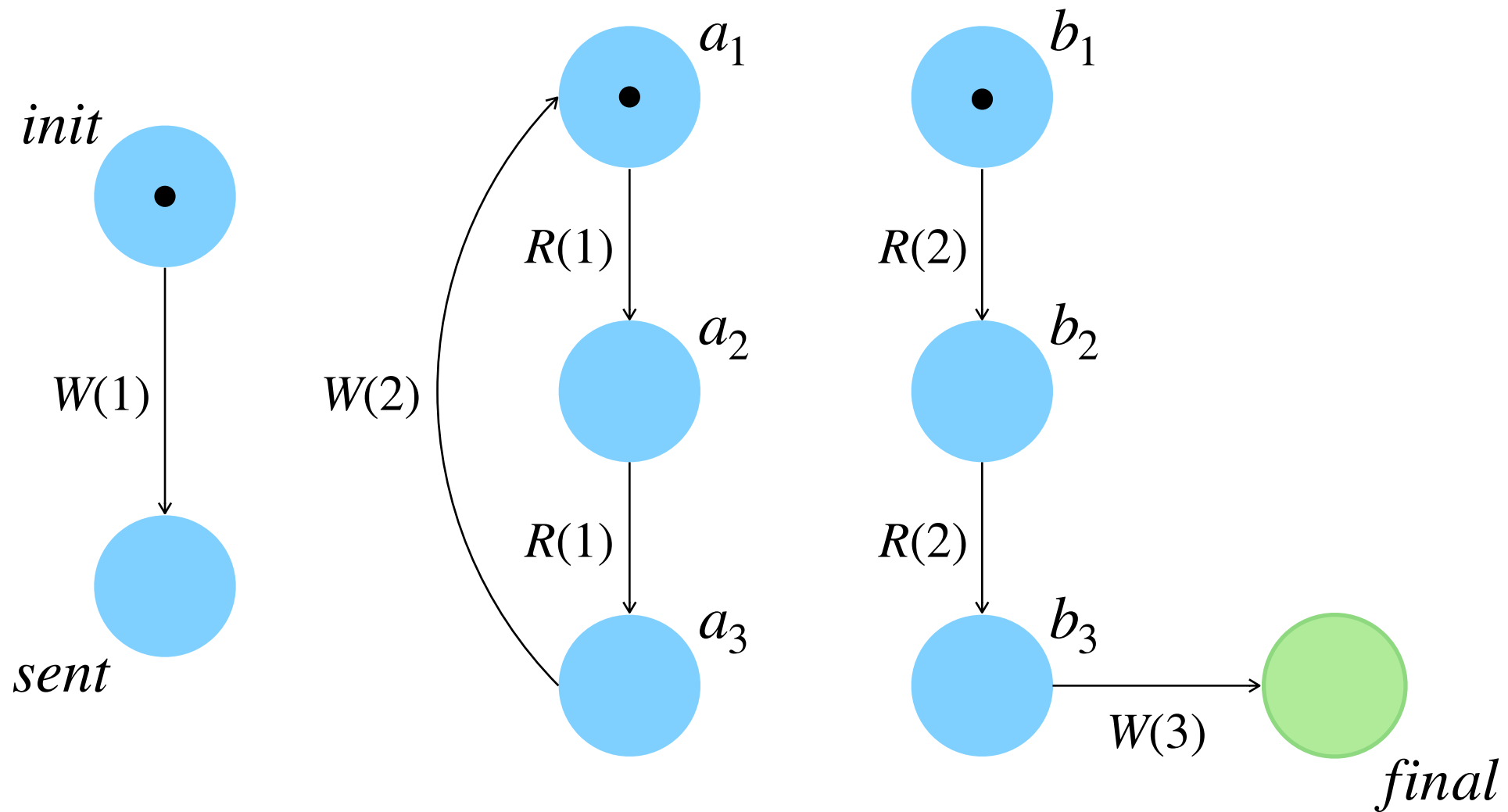


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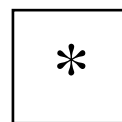
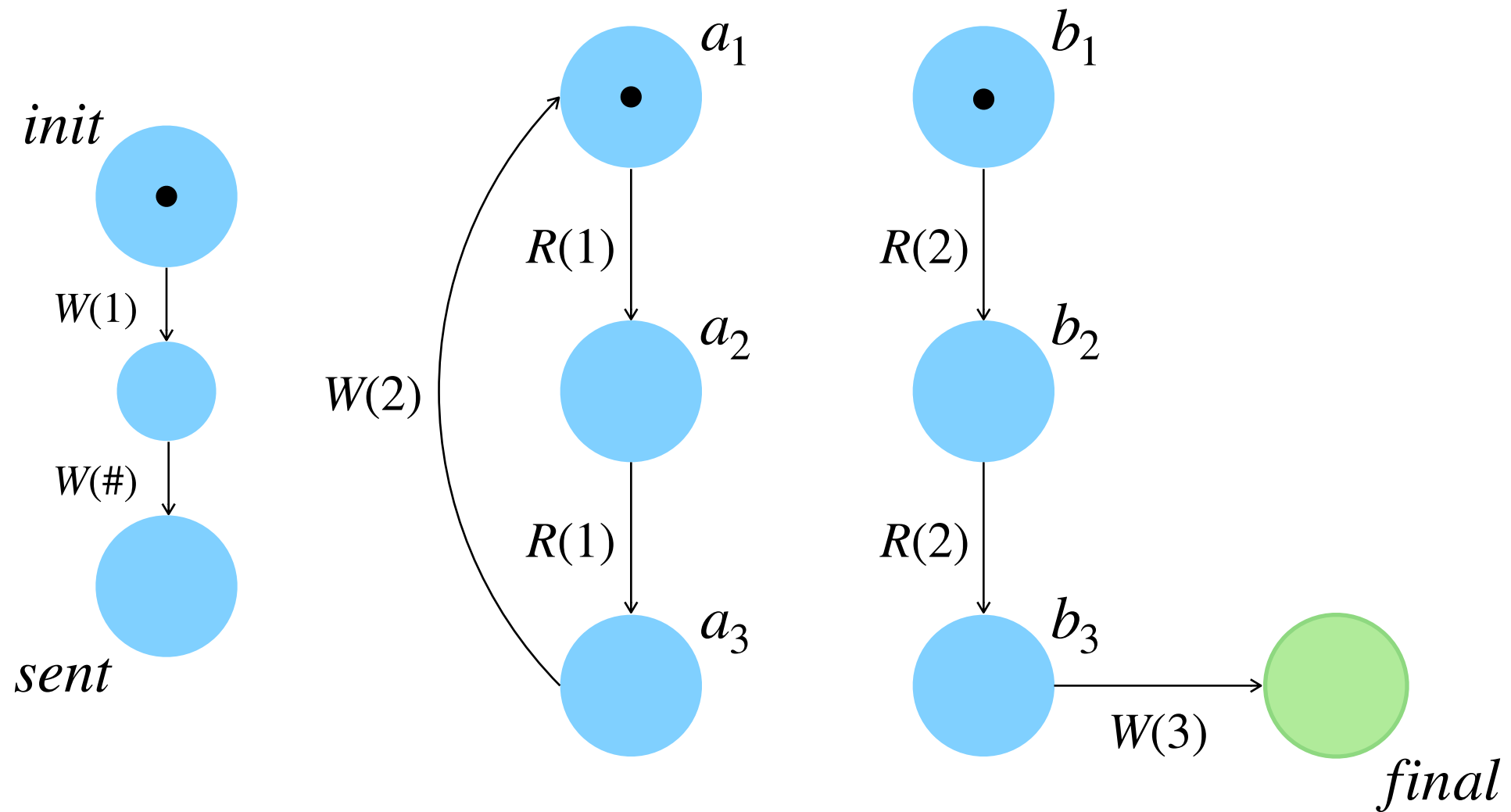
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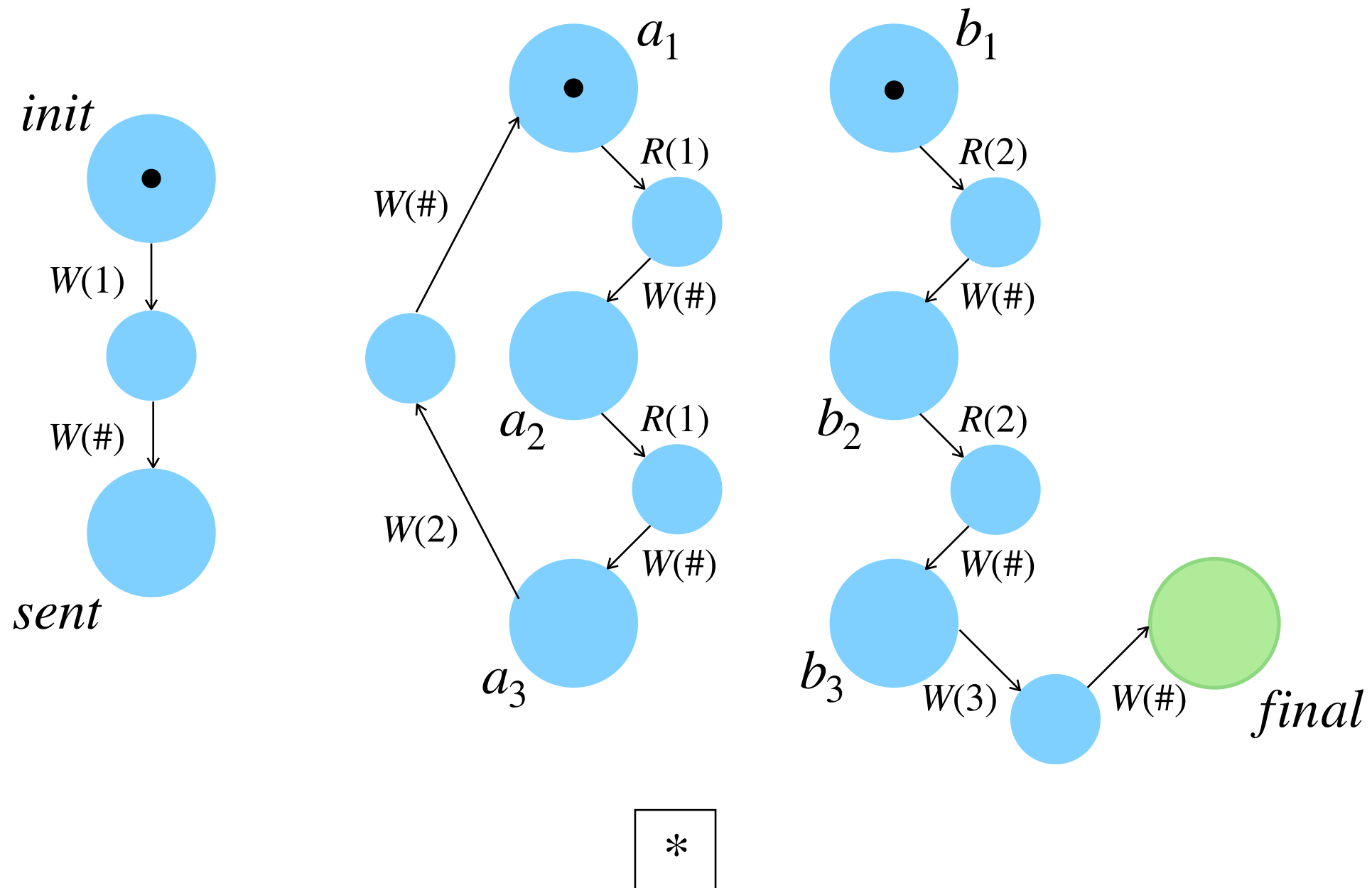
# ASMS simulates RBN

## Simulation idea #2



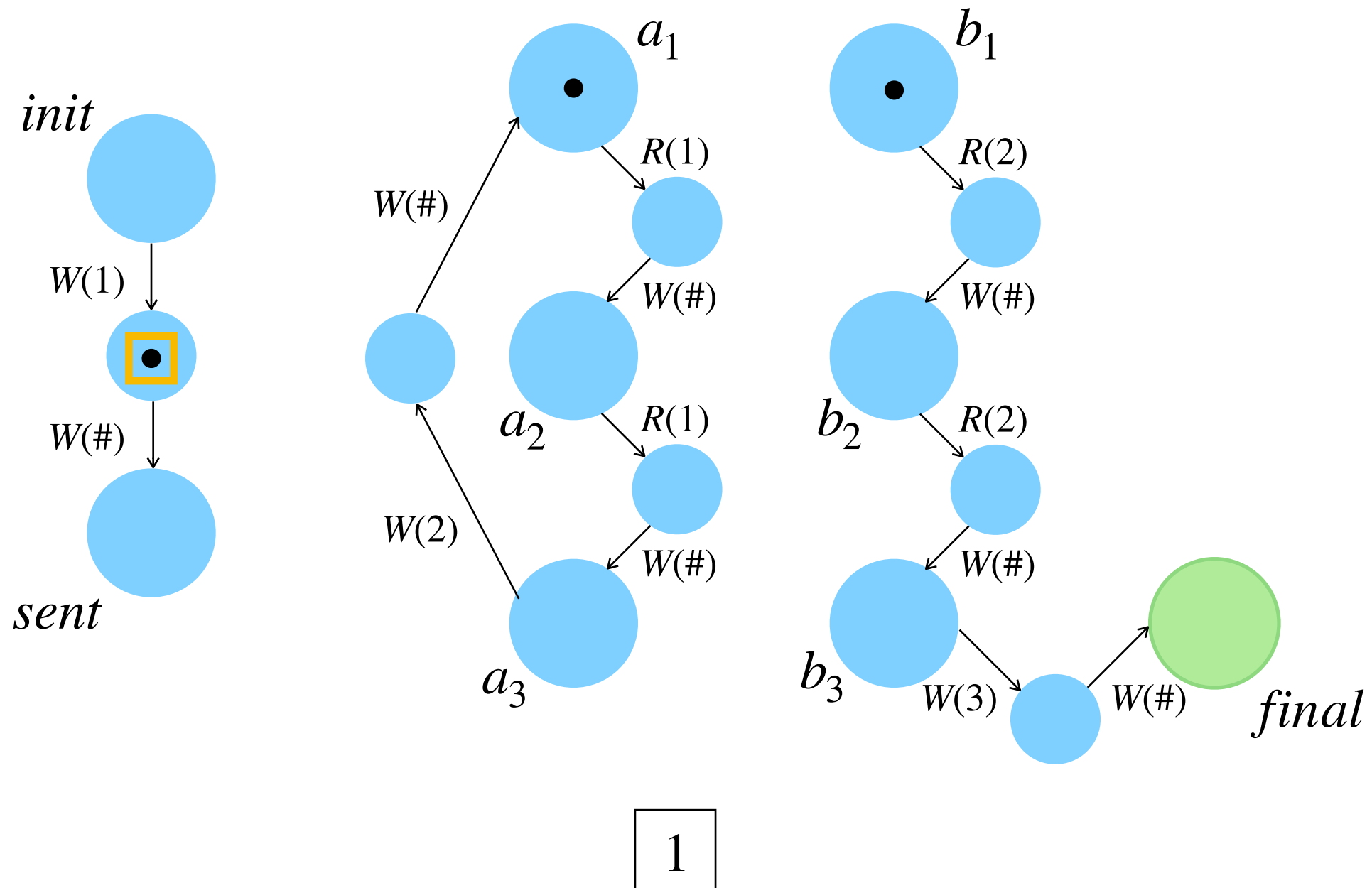
# ASMS simulates RBN

## Simulation idea #2



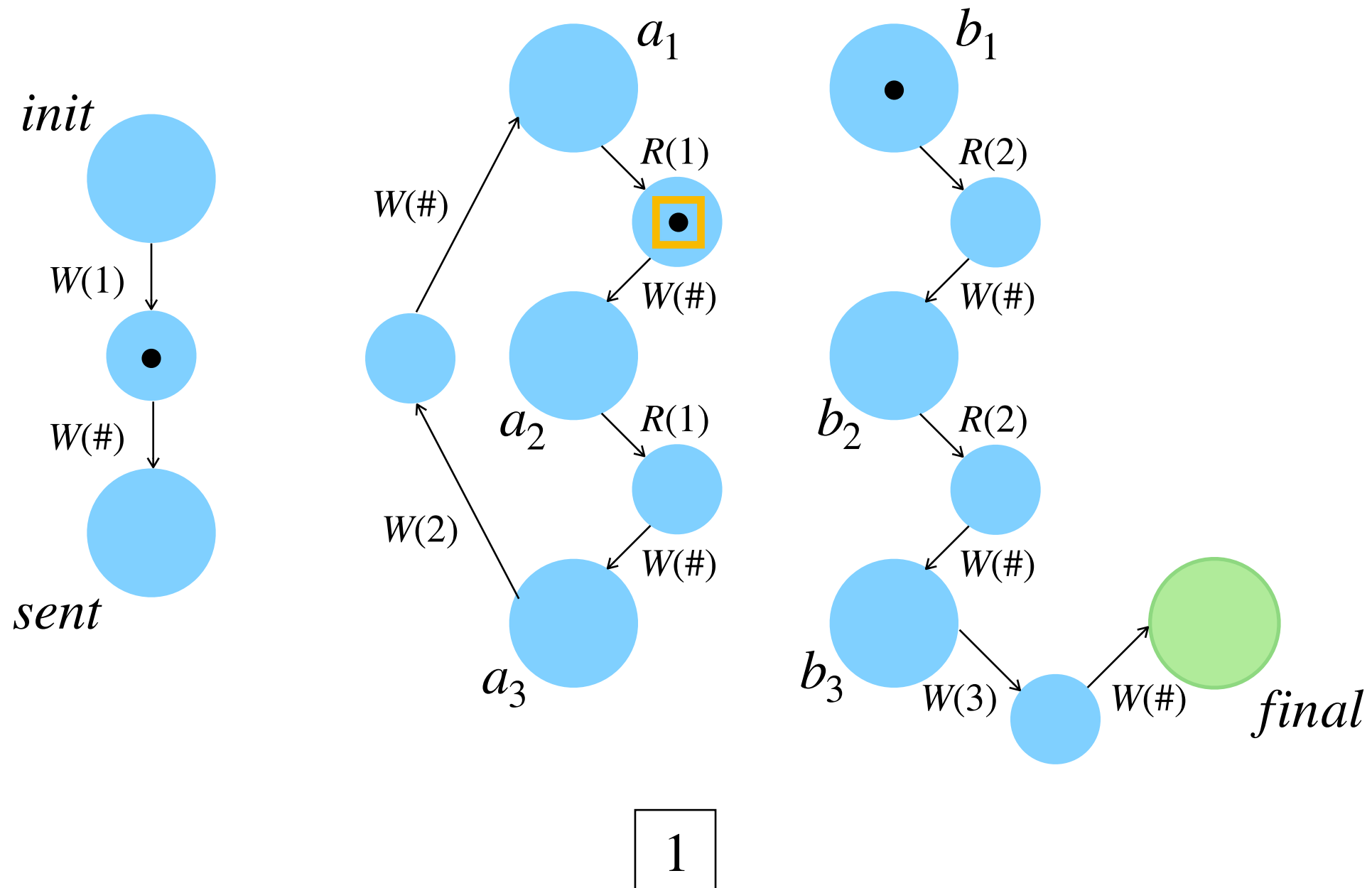
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## Simulation idea #2



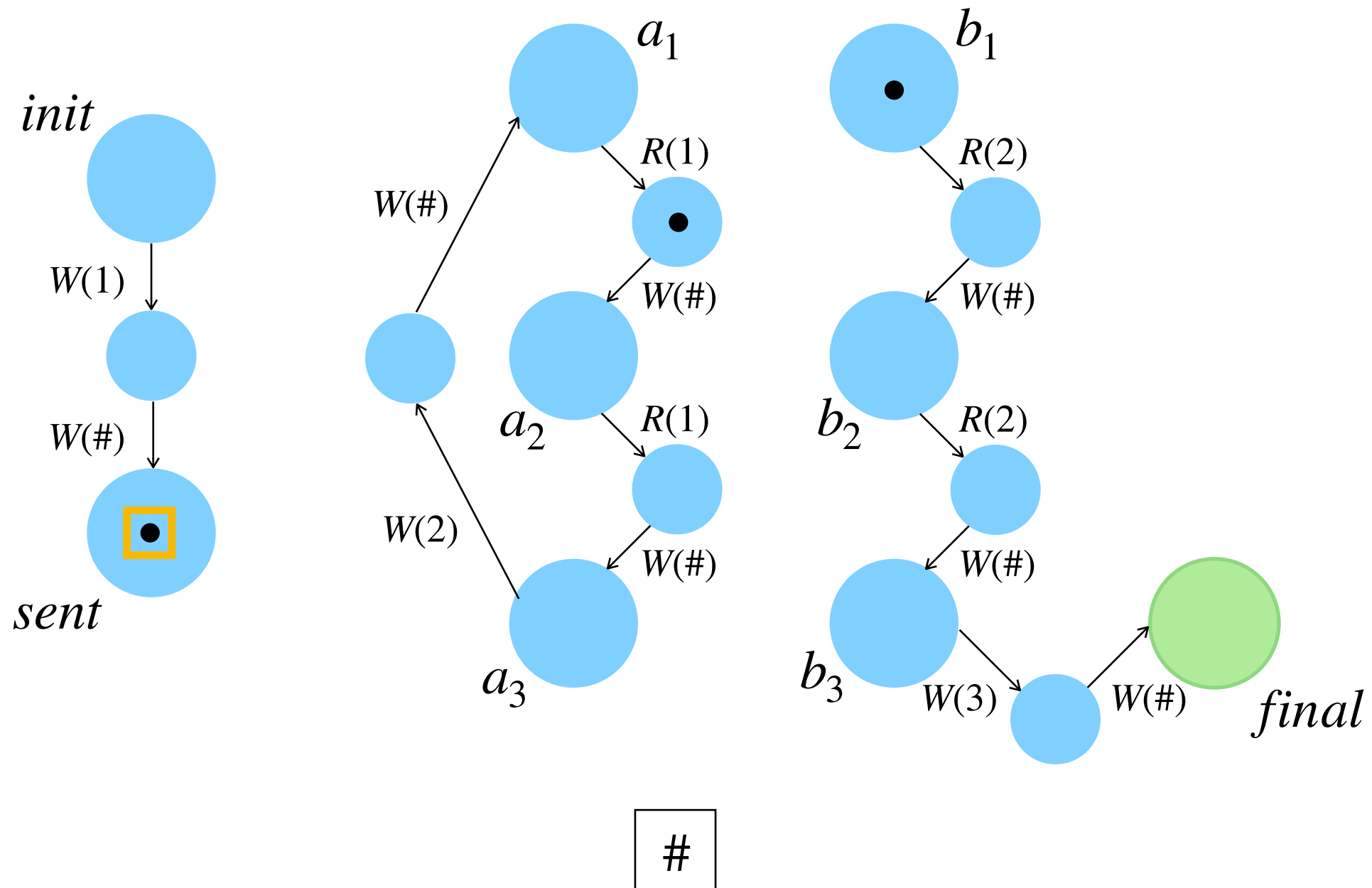
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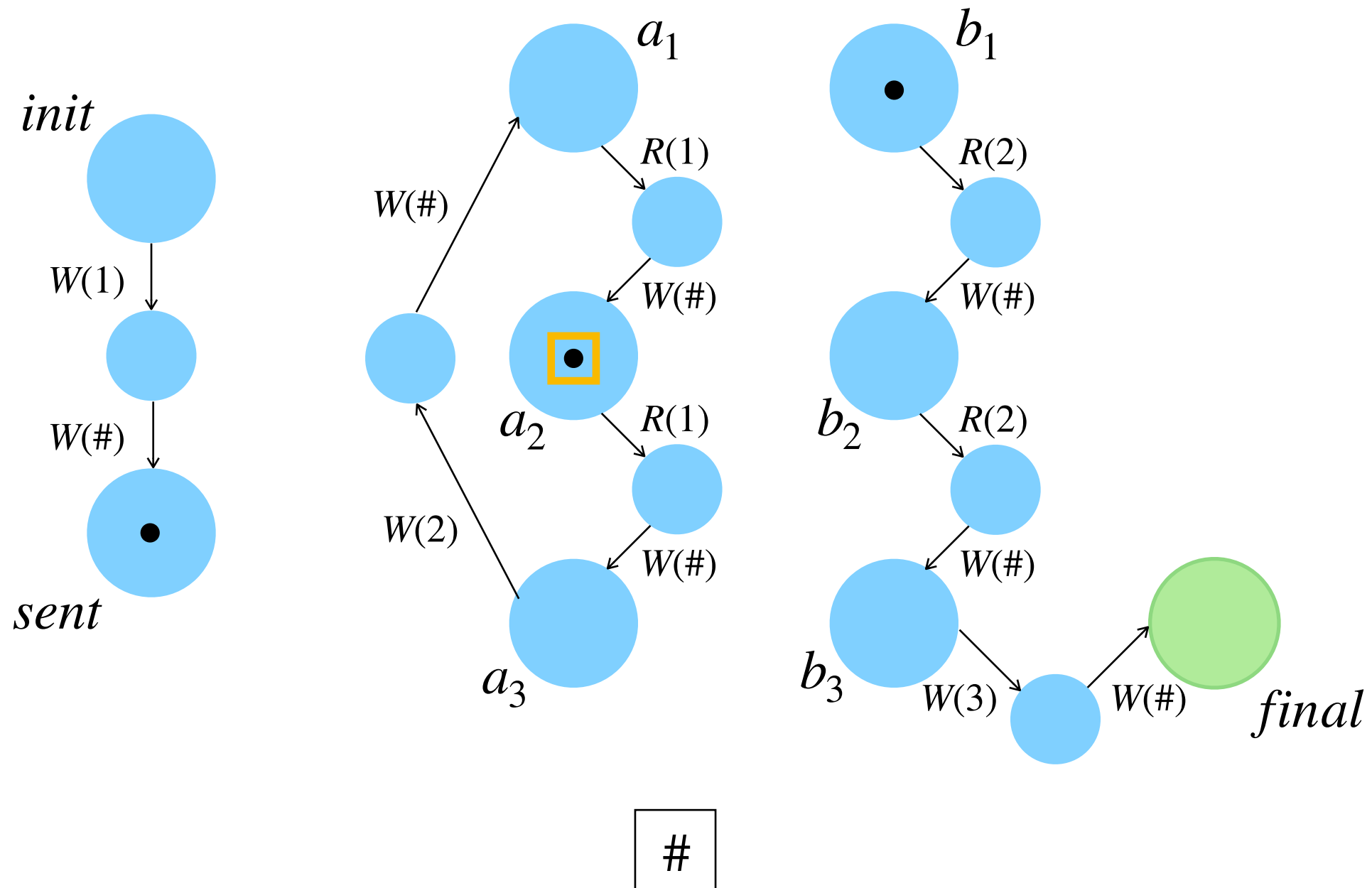
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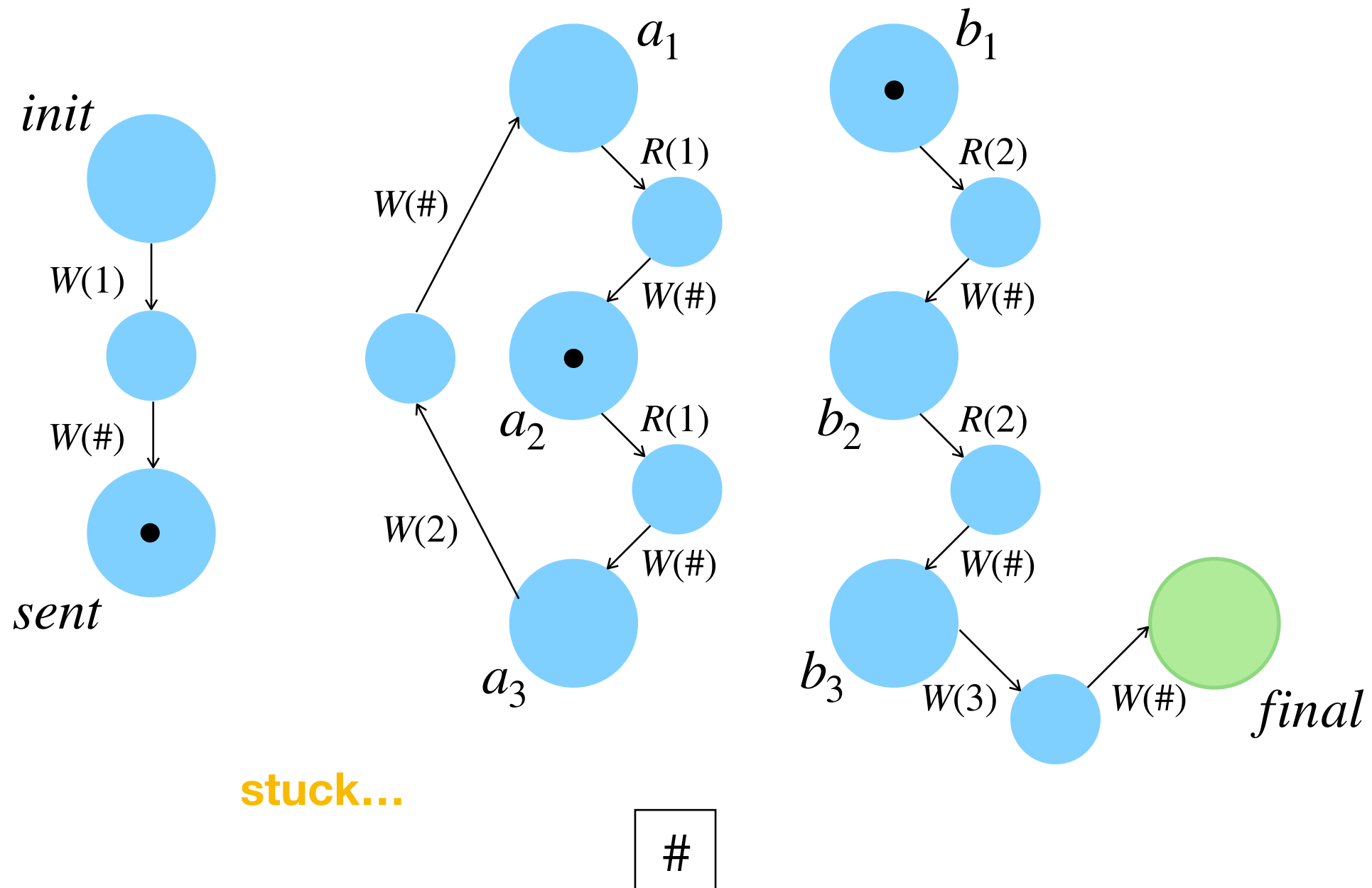
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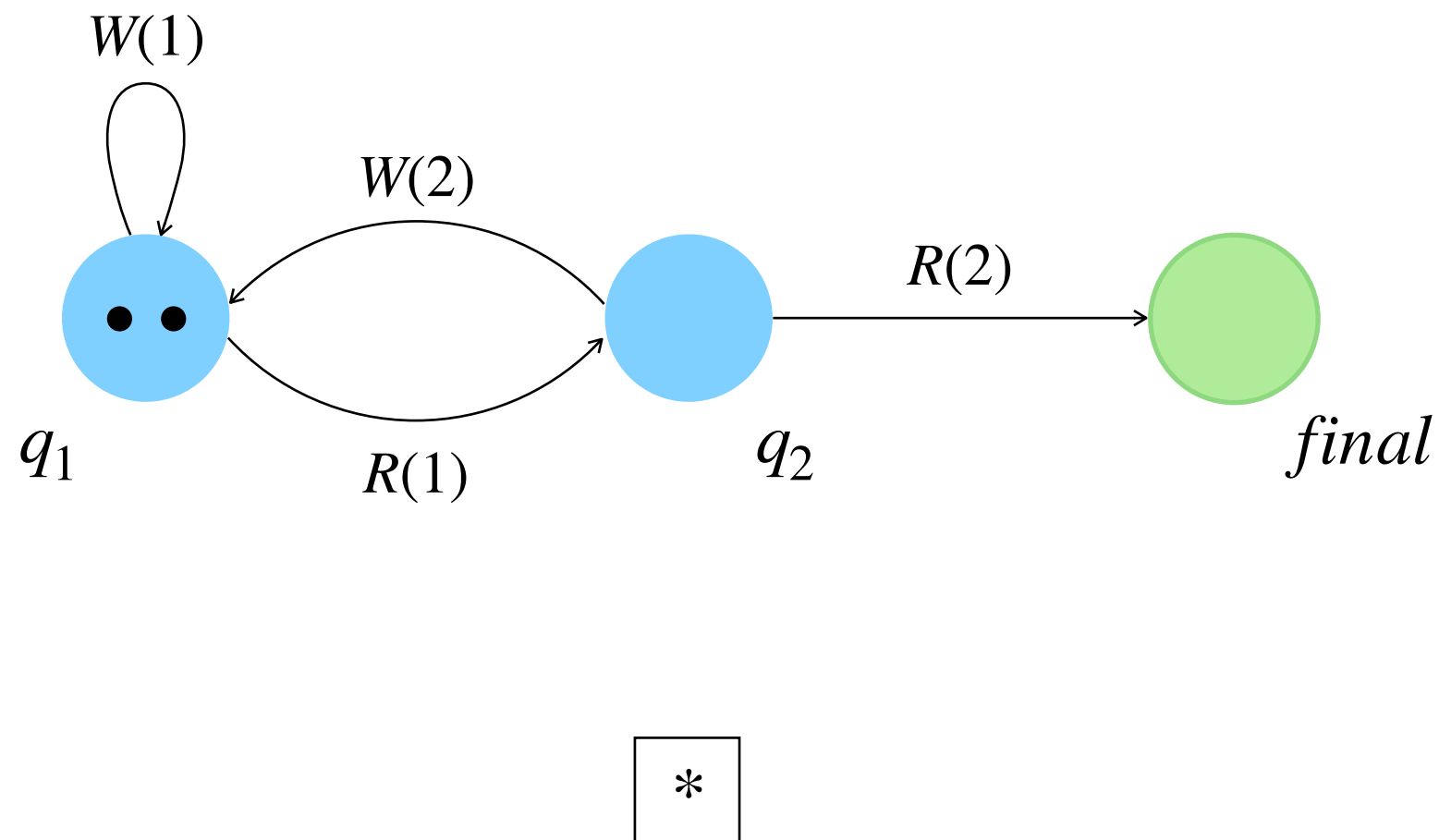
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# RBN simulates ASMS

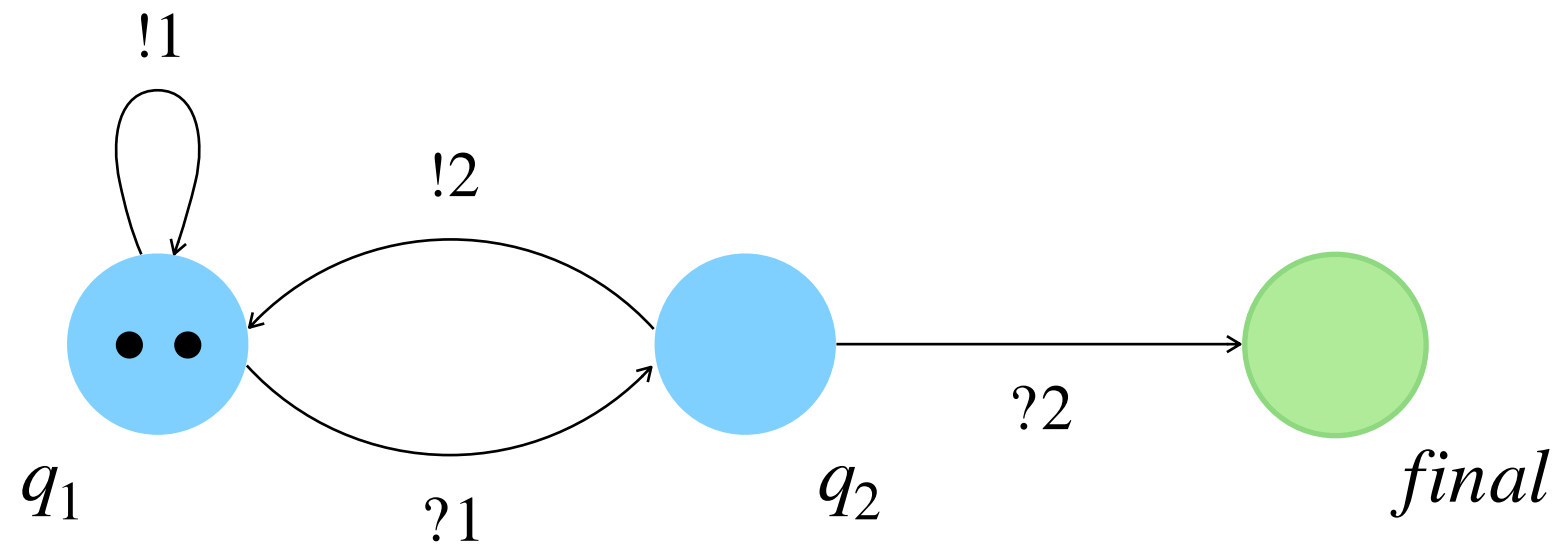
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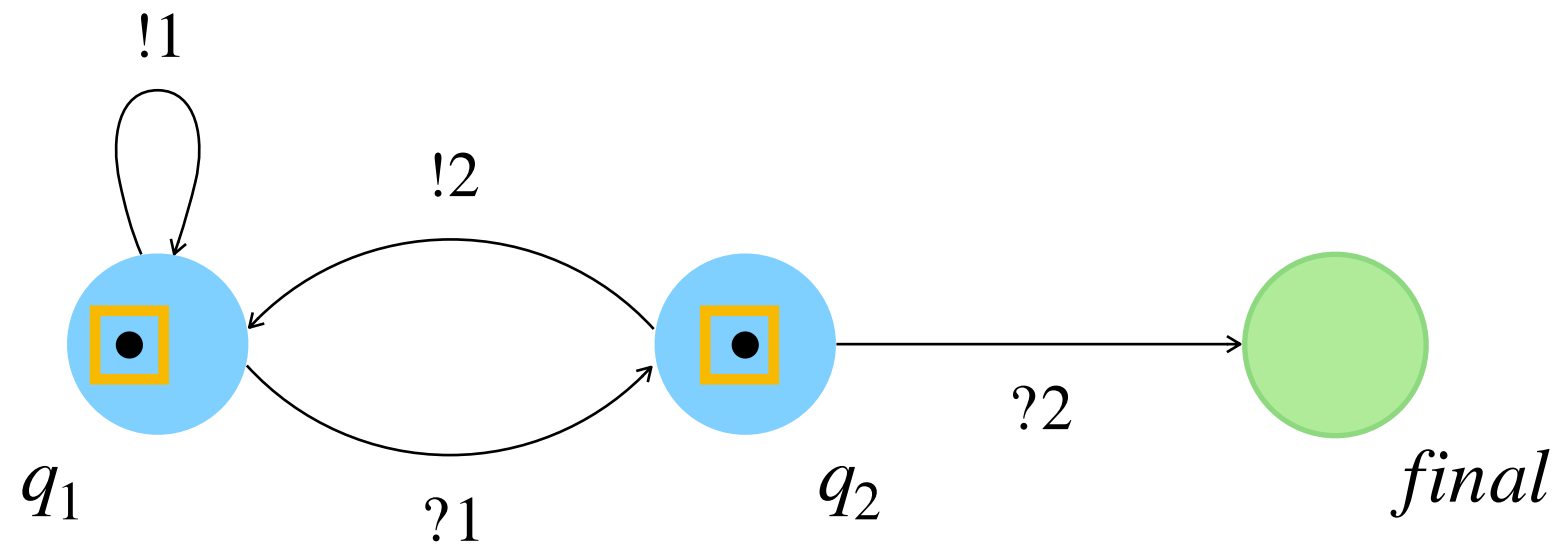
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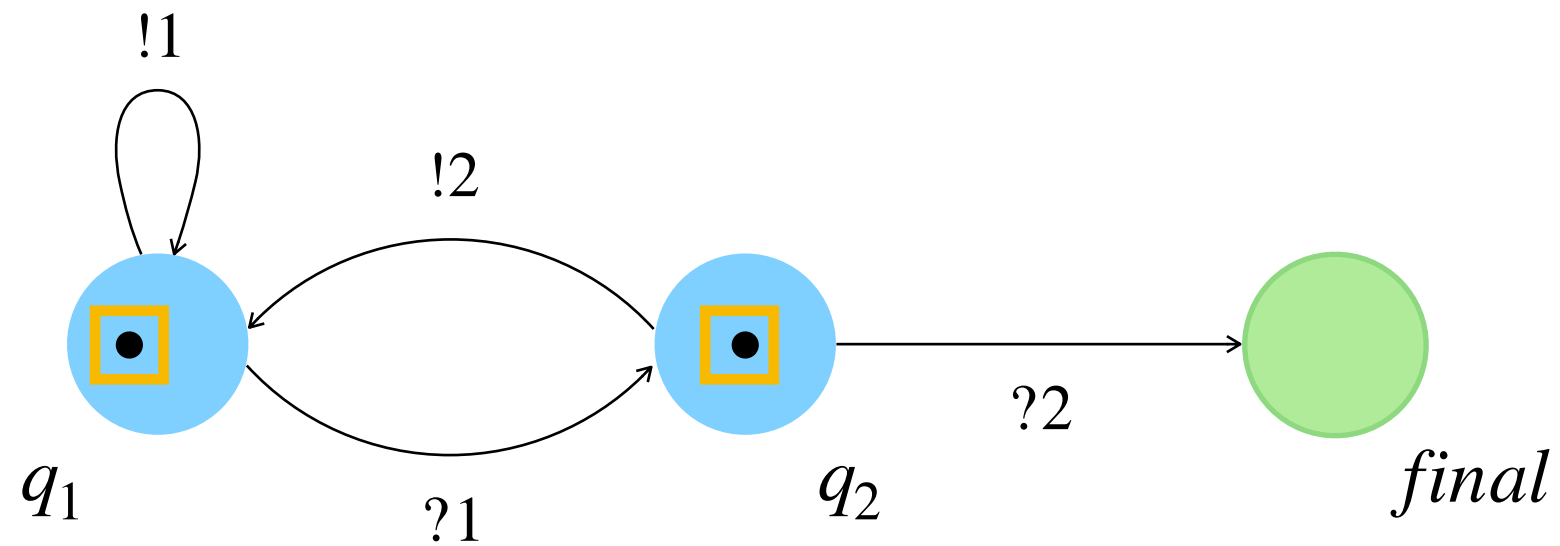
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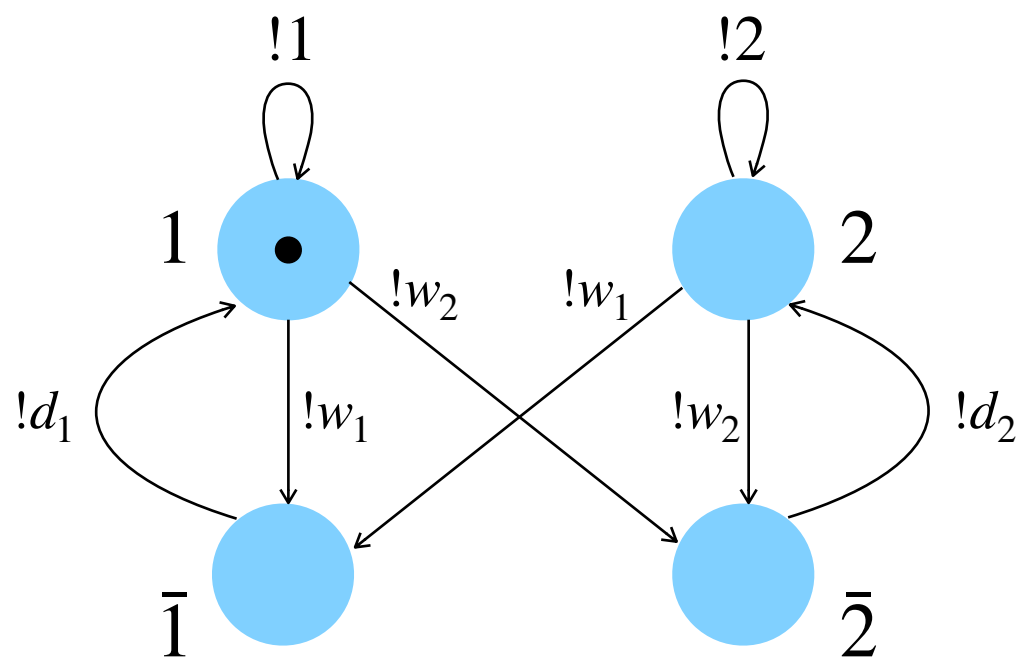
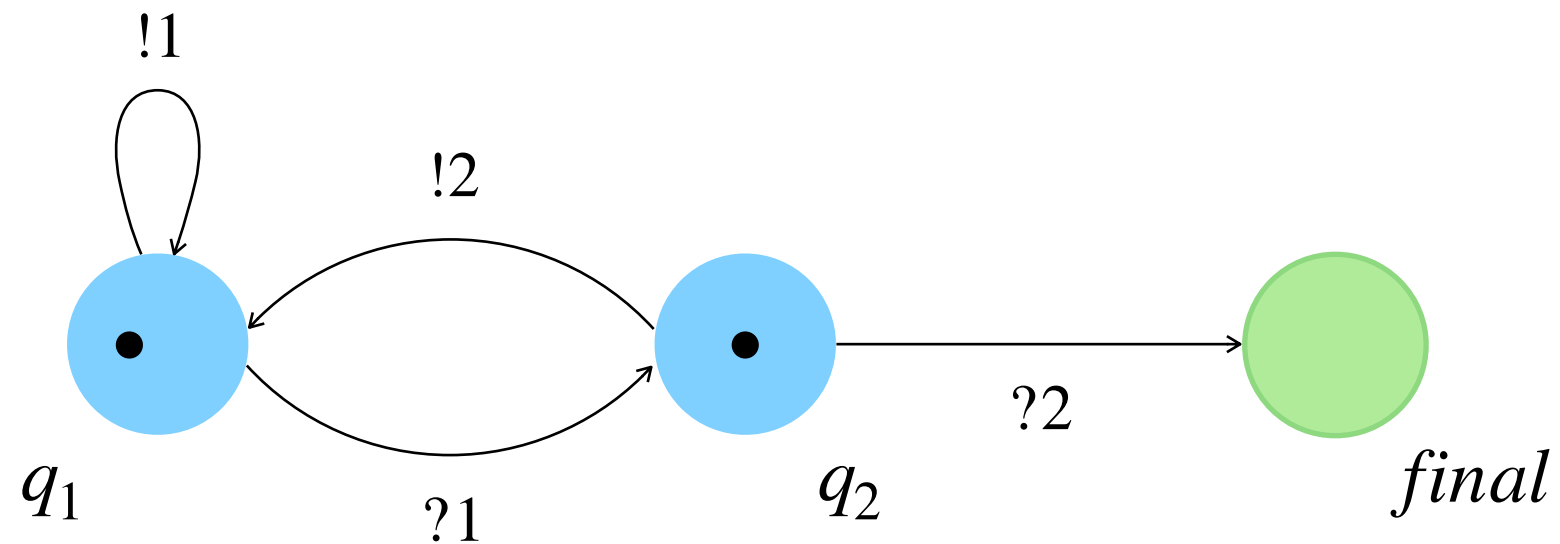
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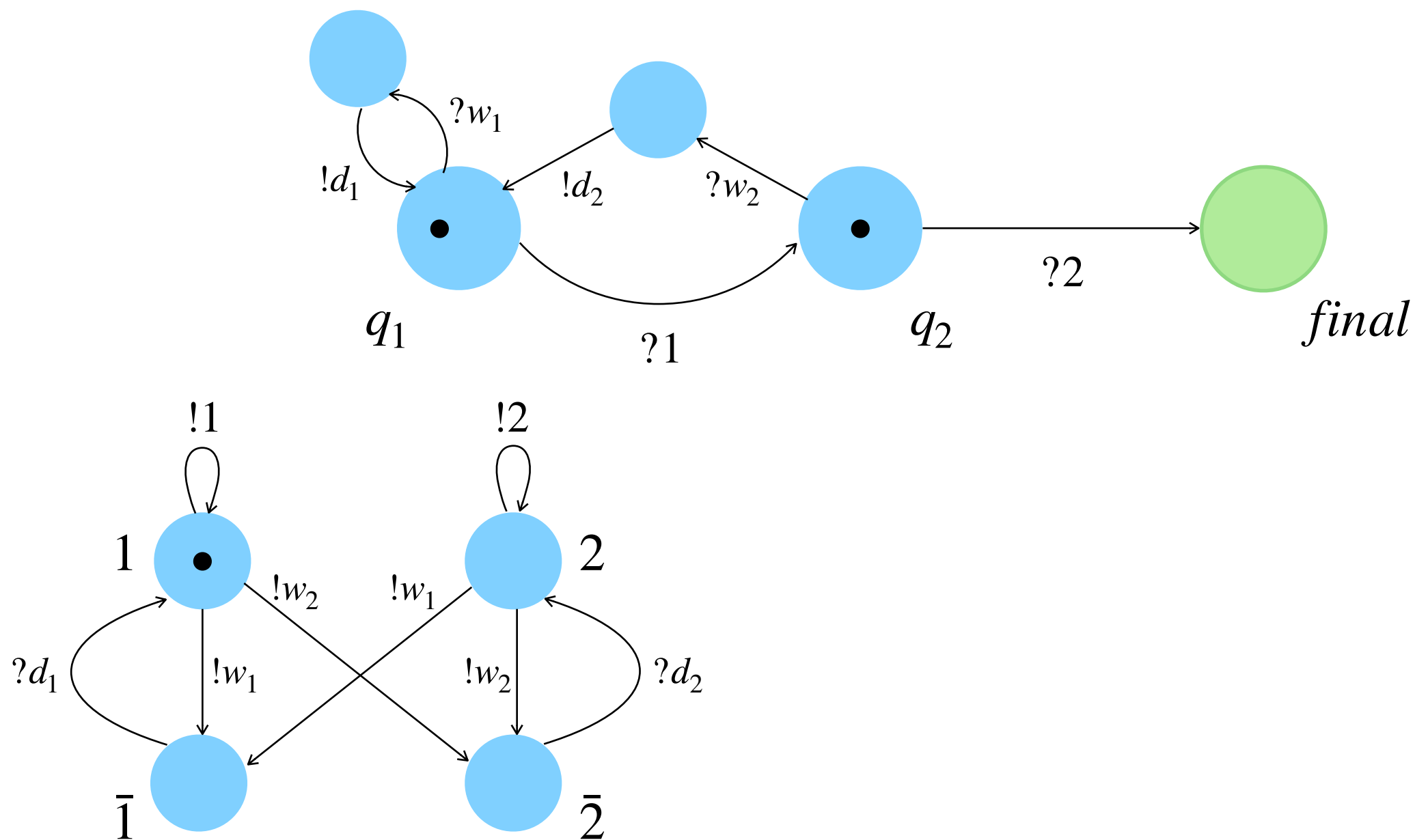
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## Simulation idea #2



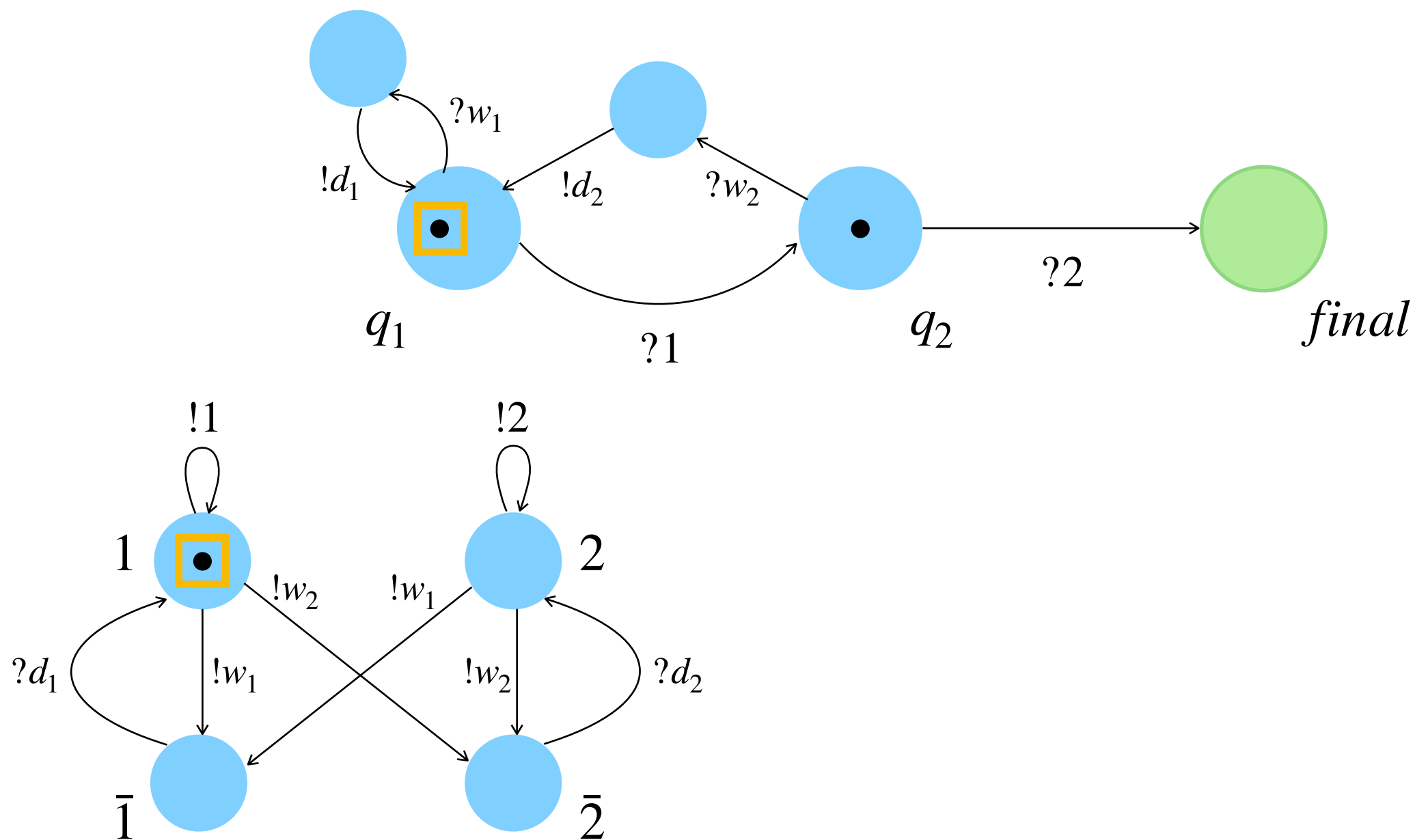
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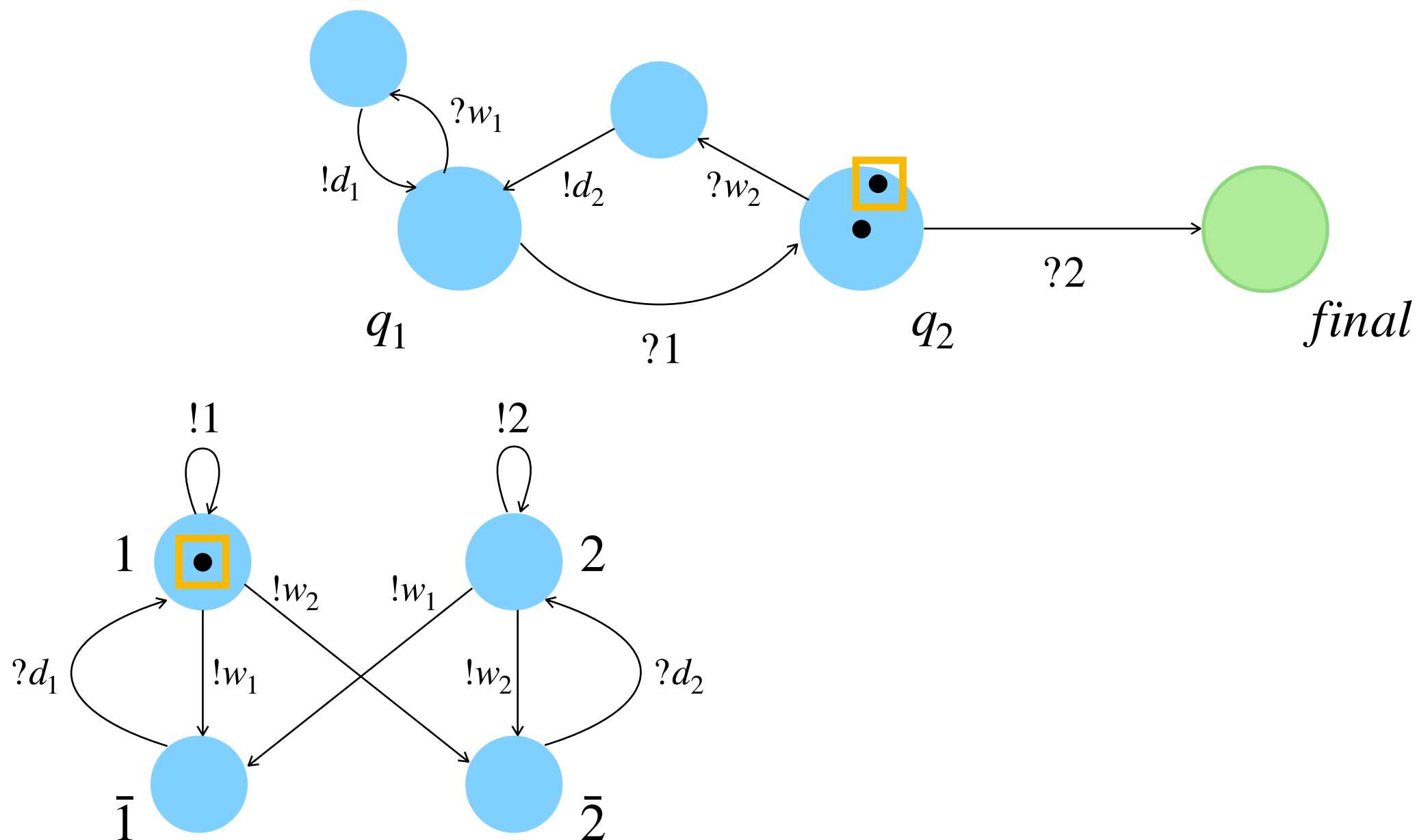
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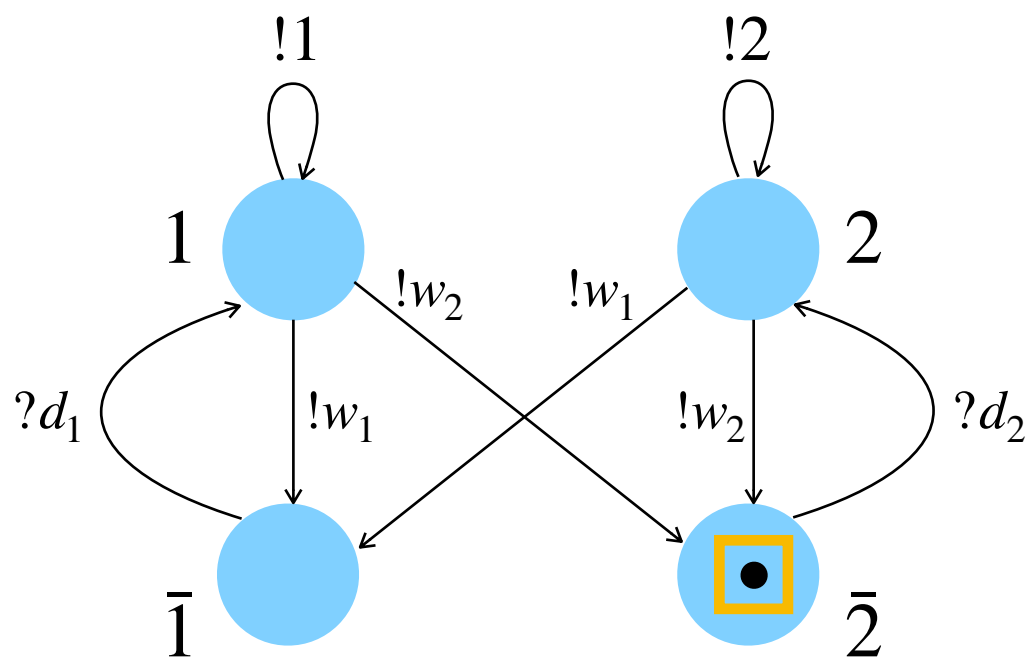
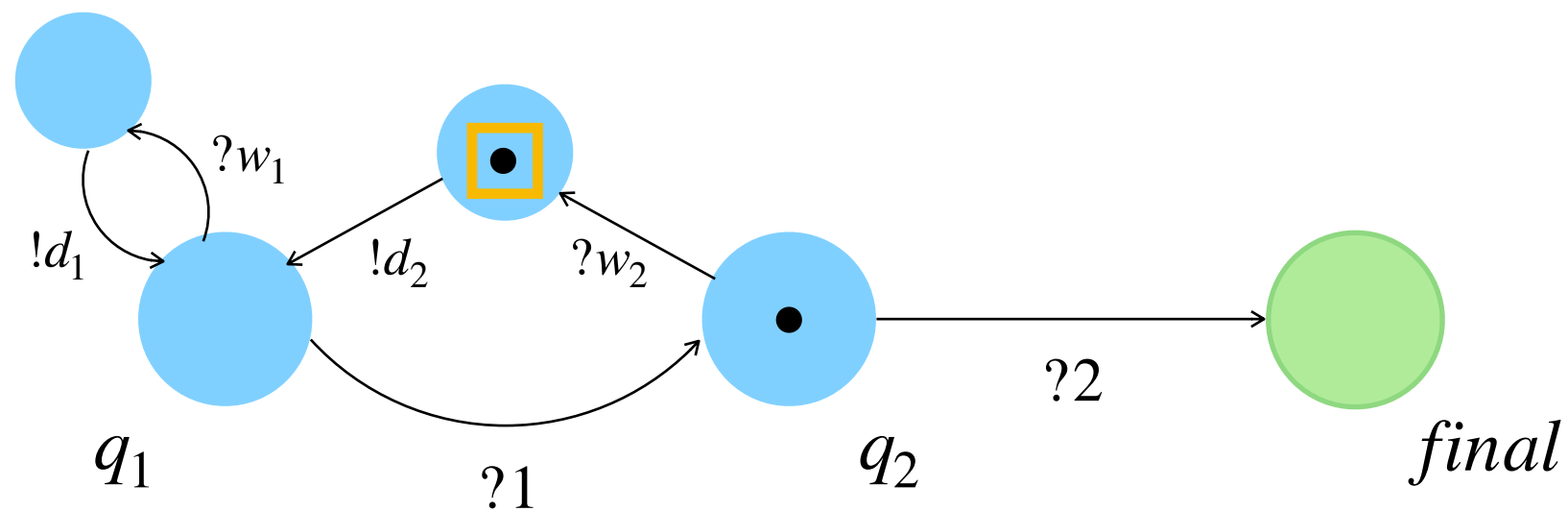
# RBN simulates ASMS

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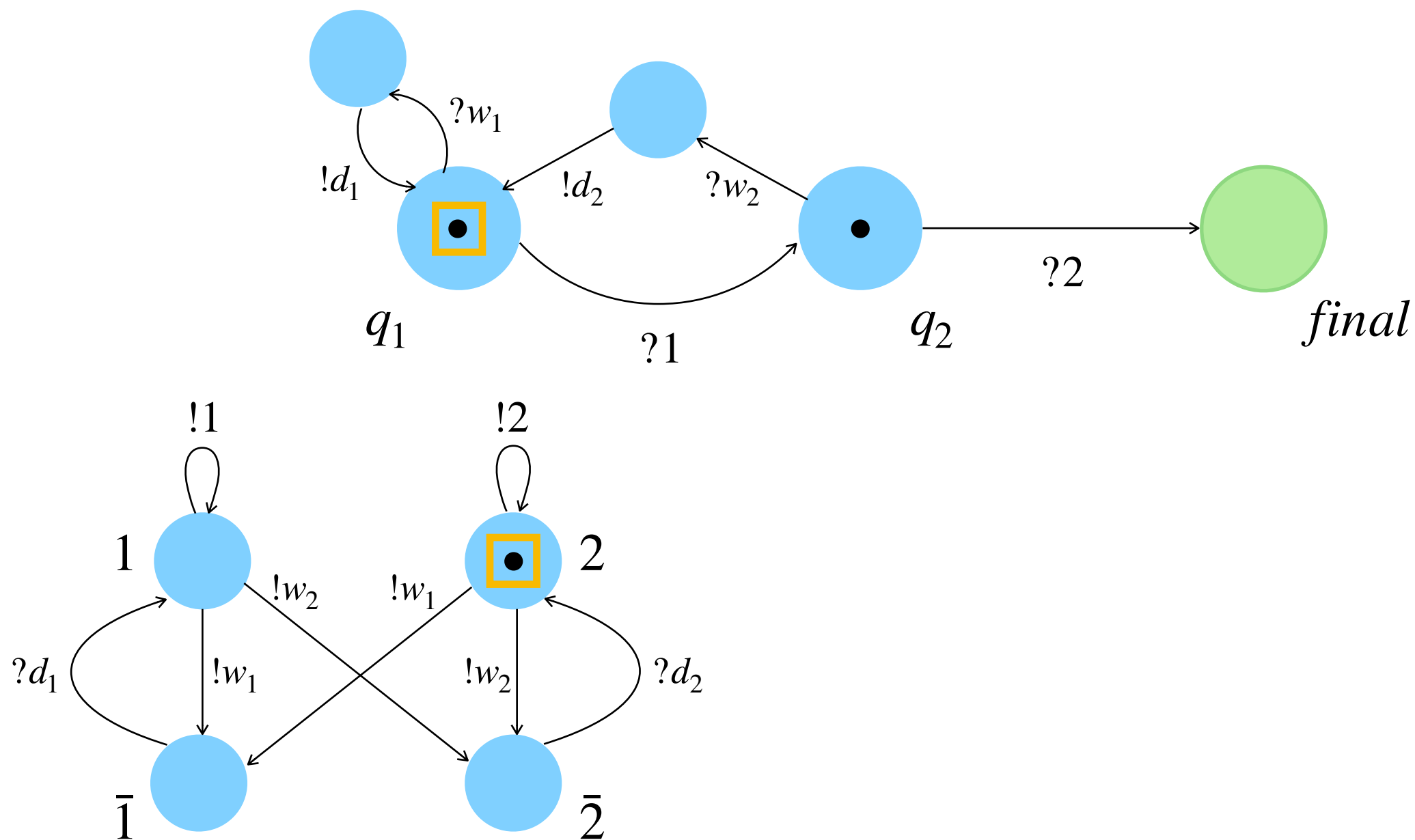
## Simulation idea #2





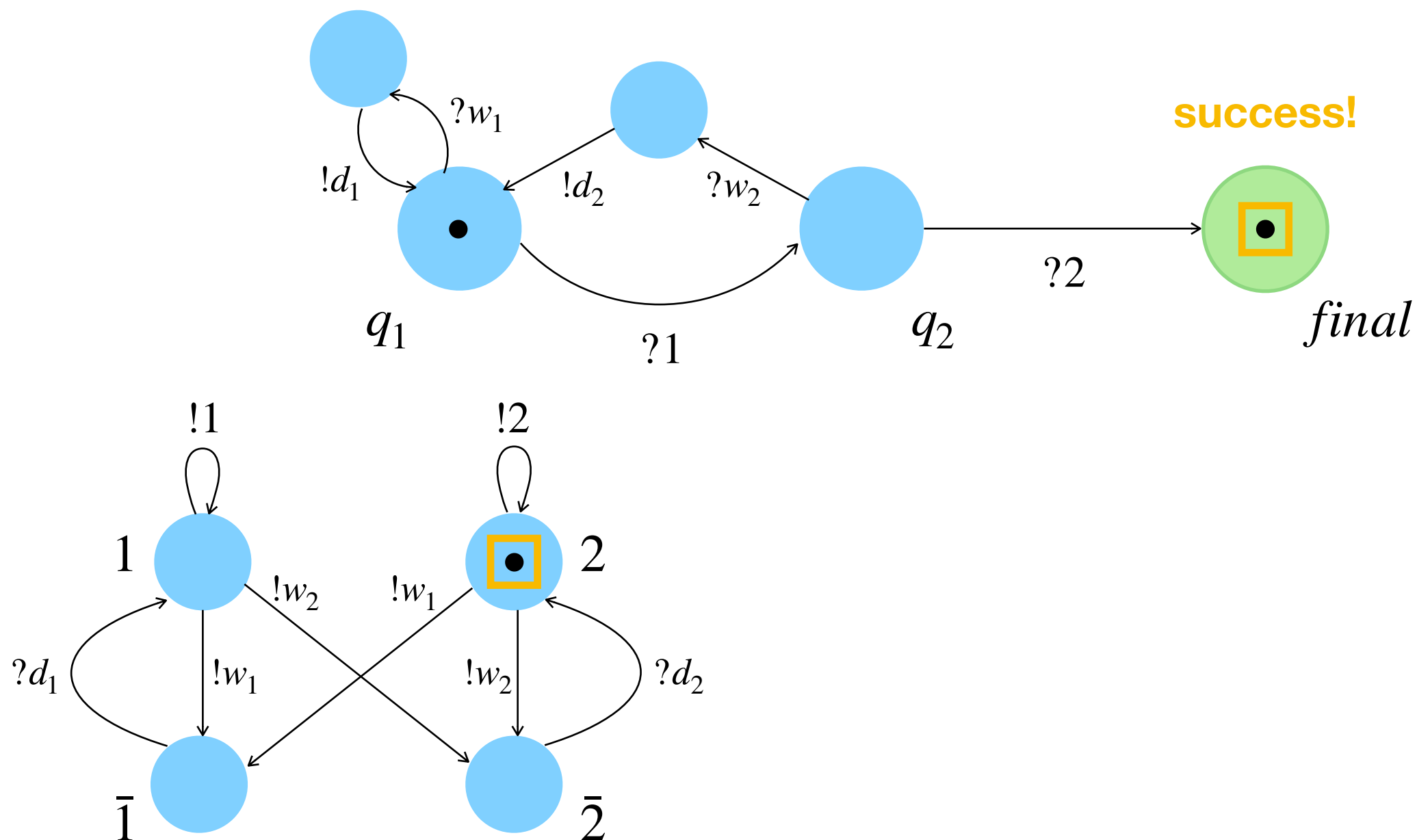
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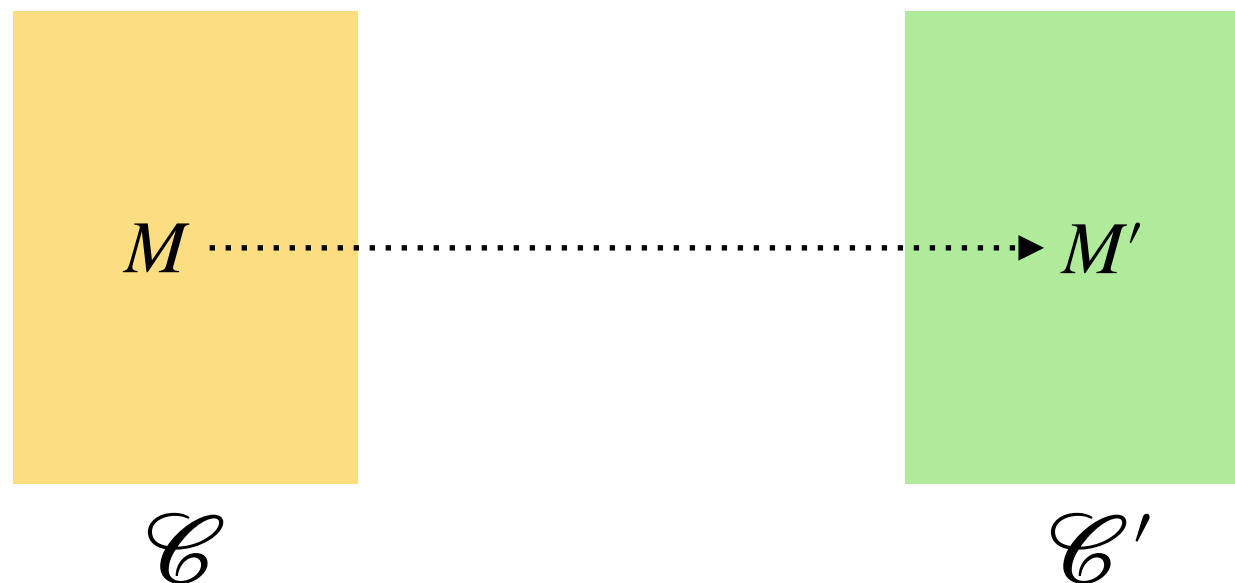
# Simulation

A **cube** is a boolean combination of constraints

number of processes in  $q$

$$a \leq \#q \leq b$$

$\in \mathbb{N}$   $\in \mathbb{N} \cup \infty$



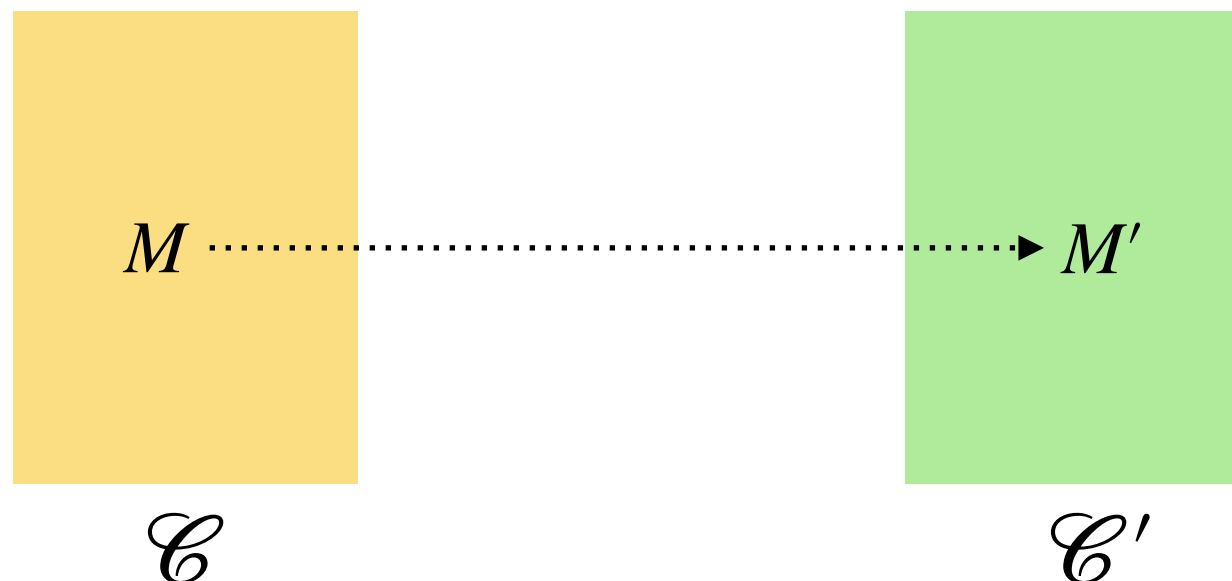
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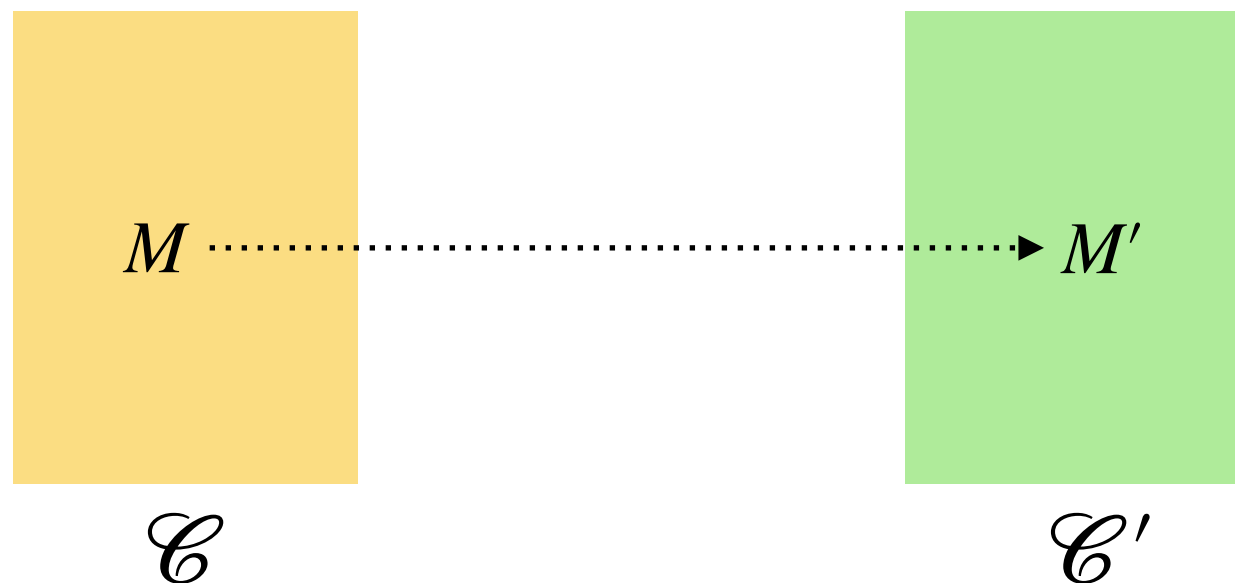
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**cube-reachability:** given cubes  $\mathcal{C}$  and  $\mathcal{C}'$ , does there exist  $M \in \mathcal{C}$  and  $M' \in \mathcal{C}'$  such that  $M$  reaches  $M'$ ?

# Simulation

RBN and ASMS are polynomial-time equivalent w.r.t. to  
cube-reachability

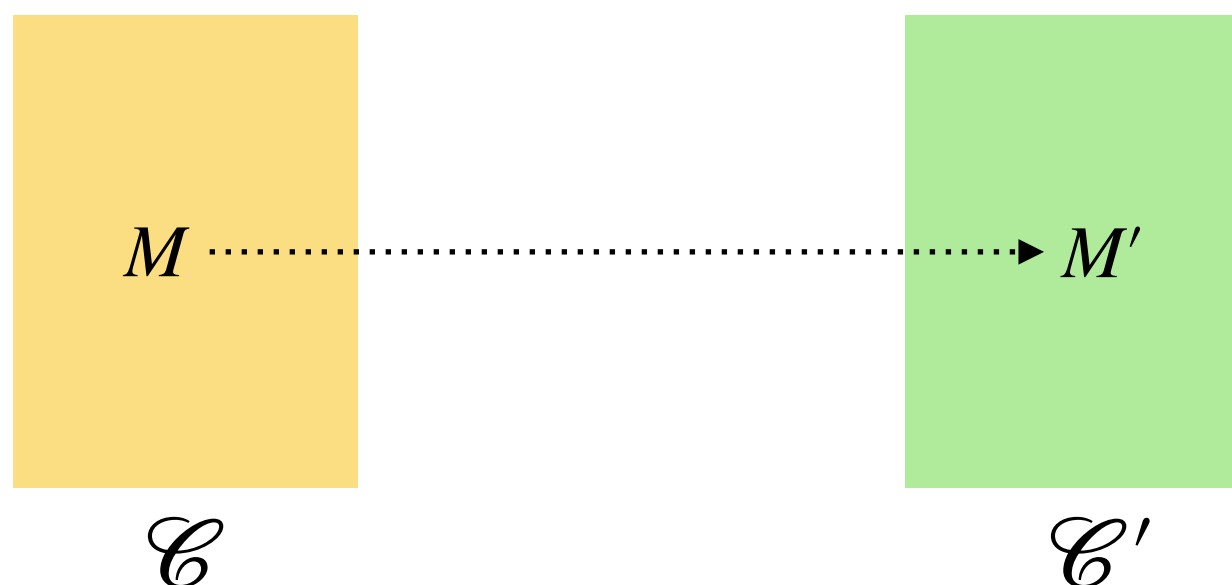


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# Transfer of results

[Delzanno et al., FSTTCS '12]

**Unbounded initial cube reachability** is PSPACE-complete for RBN



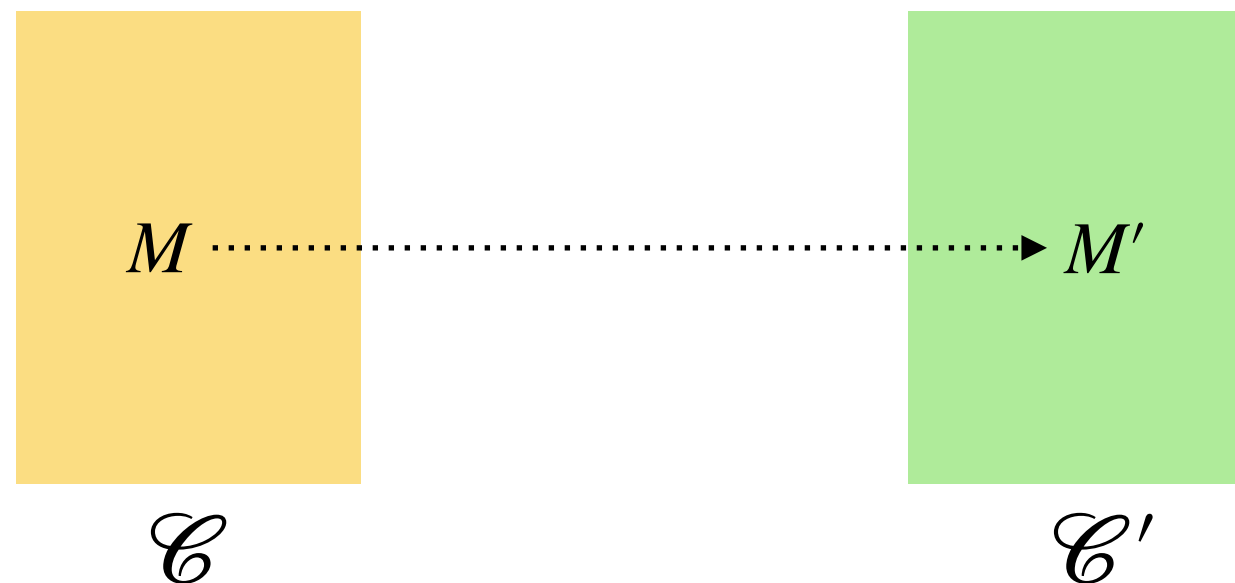
is a combination of constraints  
of the form

$$0 \leq \#q \leq 0 \text{ and } 0 \leq \#q < \infty$$

# Transfer of results

[Delzanno et al., FSTTCS '12]

**Unbounded initial cube reachability** is PSPACE-complete for RBN



PSPACE-complete for ASMS

# Transfer of results

[Bouyer et al., ICALP '16]

In ASMS,  $(\{k \cdot \textit{init}\}, r)$  **almost-surely covers** *final*  
if it covers *final* with probability 1 under a uniform stochastic scheduler



# Transfer of results

[Bouyer et al., ICALP '16]

In ASMS,  $(\{k \cdot \text{init}\}, r)$  **almost-surely covers** *final*  
if it covers *final* with probability 1 under a uniform stochastic scheduler

- $k \geq 1$  is a **positive cut-off** if  $(\{h \cdot \text{init}\}, r)$  almost-surely covers *final* for all  $h \geq k$
- $k \geq 1$  is a **negative cut-off** if  $(\{h \cdot \text{init}\}, r)$  does not almost-surely cover *final* for all  $h \geq k$

# Transfer of results

[Bouyer et al., ICALP '16]

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# Transfer of results

[Bouyer et al., ICALP '16]

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# Conclusion

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We also compare RBN / ASMS to **immediate observation nets (IO)**:

- anonymous, identical processes which can communicate by observation.
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**Thank you!**

# Parameterized problems

A **cube** is a boolean combination of constraints  $a \leq \#q \leq b$   
 $\in \mathbb{N}$   $\in \mathbb{N} \cup \infty$

**parameterized problems:** verifying predicates using boolean operators and reachability operators  $pre^*$  and  $post^*$  over cubes

PSPACE

$pre^*(\mathcal{C})$  is the set of markings that can reach  $\mathcal{C}$

$post^*(\mathcal{C})$  is the set of markings that  $\mathcal{C}$  can reach

e.g. reachability from cube  $\mathcal{C}$  to cube  $\mathcal{C}'$ :  $post^*(\mathcal{C}) \cap \mathcal{C}' \neq \emptyset$

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$$post^*(\mathcal{C}_{init}) \subseteq pre^*(\mathcal{C}_{final})$$