#### **HIGHLIGHTS**

September 18th, 2020

# Branching Immediate Observation Petri Nets

Chana Weil-Kennedy joint work with Javier Esparza and Mikhail Raskin



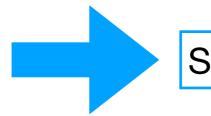


**Reachability problem:** Given a Petri net  $\mathcal{N}$ , and markings  $M_0$  and M

can marking  $M_0$  reach marking M in  $\mathcal N$  ?

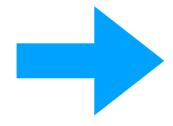
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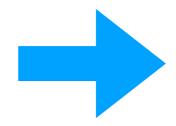


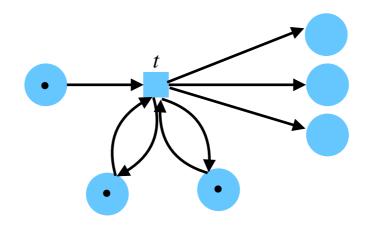
Study subclasses of Petri nets

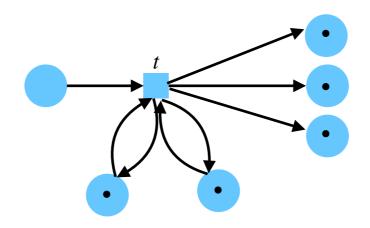
**Reachability problem:** Given a Petri net  $\mathcal{N}$ , and markings  $M_0$  and M can marking  $M_0$  reach marking M in  $\mathcal{N}$  2 complexity non-elementary complexity. Lasota, Lazic, Leroux, Mazowiecki, 19]

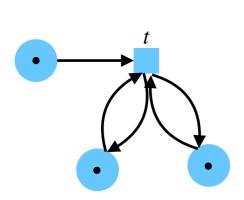


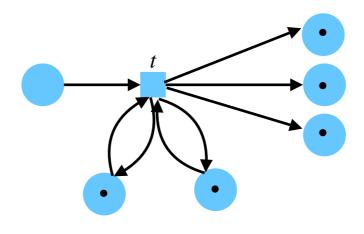
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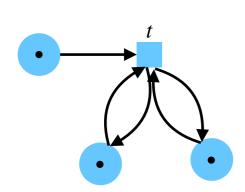


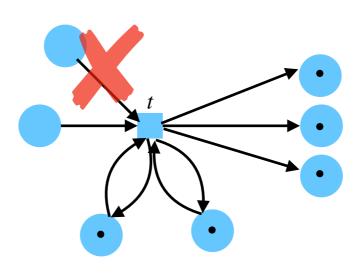


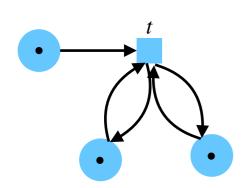


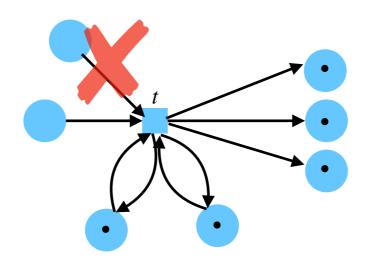


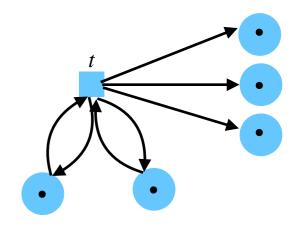


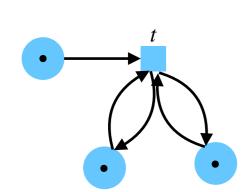


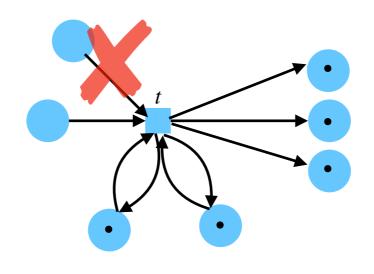


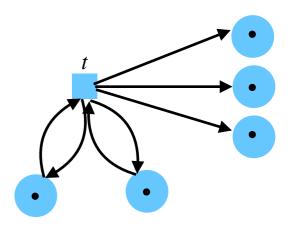






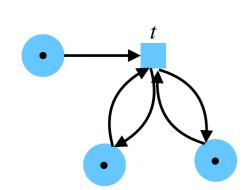


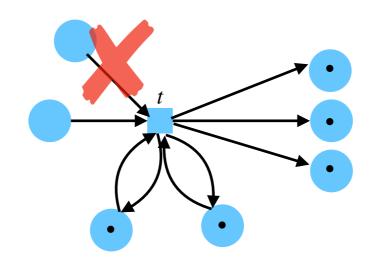


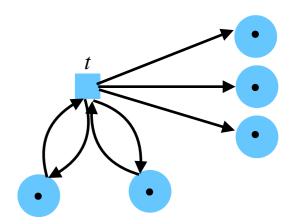


$$Card(^{\bullet}t - t^{\bullet}) \leq 1$$

≡ at most one "pure input" place





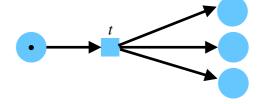


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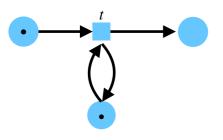
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Extend

• BPP nets:



• IO nets:



[Leroux, Sutre, '05]

**flat**  $\exists$  sequence  $t_1^*t_2^*...t_l^*$  such that

$$M_0 \xrightarrow{*} M \text{ iff } M_0 \xrightarrow{t_1^{k_1} t_2^{k_2} \dots t_l^{k_l}} M$$

[Leroux, Sutre, '05]

flat 
$$\exists$$
 sequence  $t_1^*t_2^*...t_l^*$  such that  $M_0 \stackrel{*}{\to} M$  iff  $M_0 \stackrel{t_1^{k_1}t_2^{k_2}...t_l^{k_l}}{\to} M$  BPP, IO nets

pre\*-flat 
$$\forall M, \exists$$
 sequence  $t_1^*t_2^*...t_l^*$  such that  $M_0 \xrightarrow{*} M$  iff  $M_0 \xrightarrow{t_1^{k_1}t_2^{k_2}...t_l^{k_l}} M$ 

BIO nets

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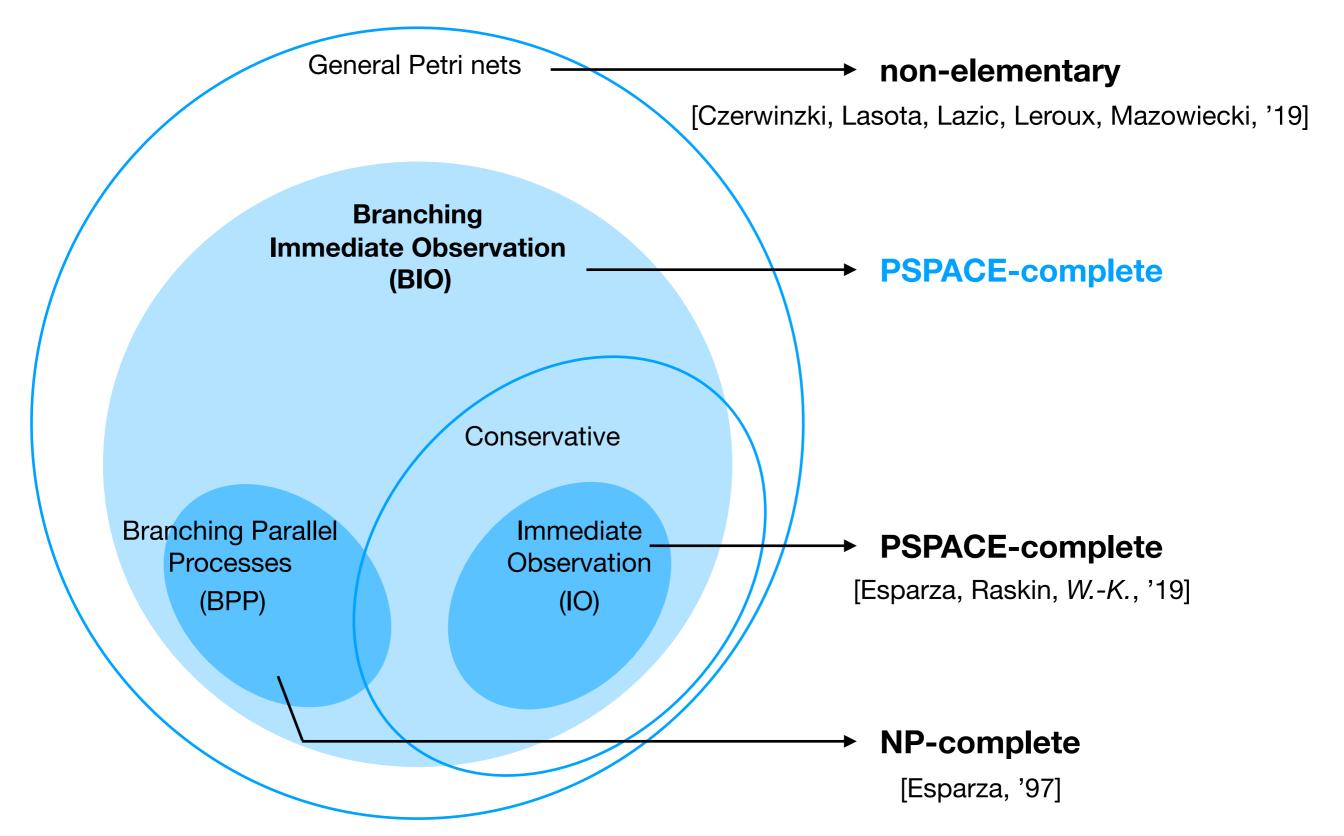
$$M_0 \stackrel{*}{\rightarrow} M \text{ iff } M_0 \stackrel{t_1^{k_1}t_2^{k_2}...t_l^{k_l}}{\longrightarrow} M$$
 BPP, IO nets

pre\*-flat  $\forall M, \exists$  sequence  $t_1^*t_2^*...t_l^*$  such that  $M_0 \xrightarrow{*} M$  iff  $M_0 \xrightarrow{t_1^{k_1}t_2^{k_2}...t_l^{k_l}} M$  ↓

**BIO** nets

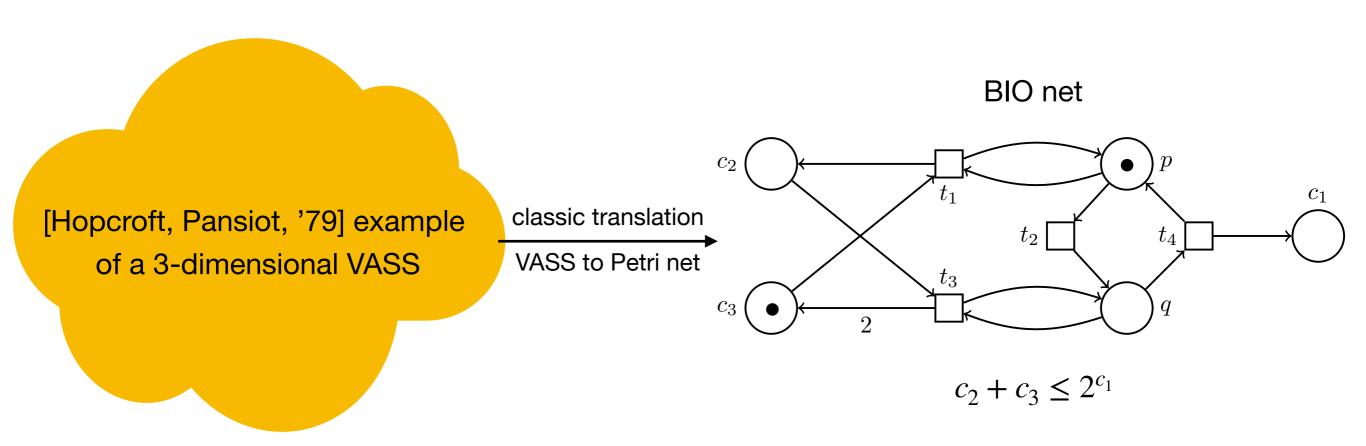
check **reachability properties** with model checking **tools** that use acceleration techniques **e.g. FAST** [Bardin, Finkel, Leroux, Petrucci, '03]

### A strong class with simple reachability

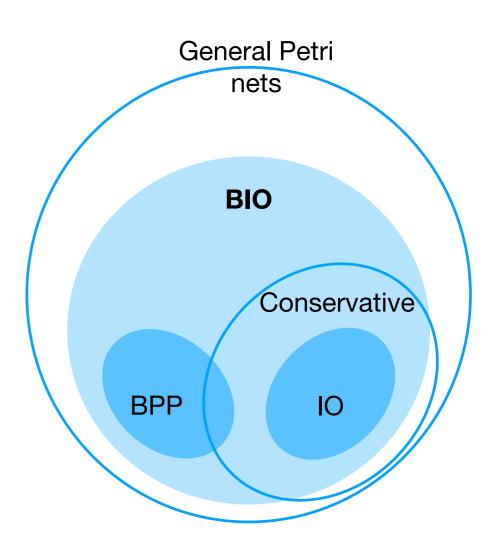


### A strong class with simple reachability

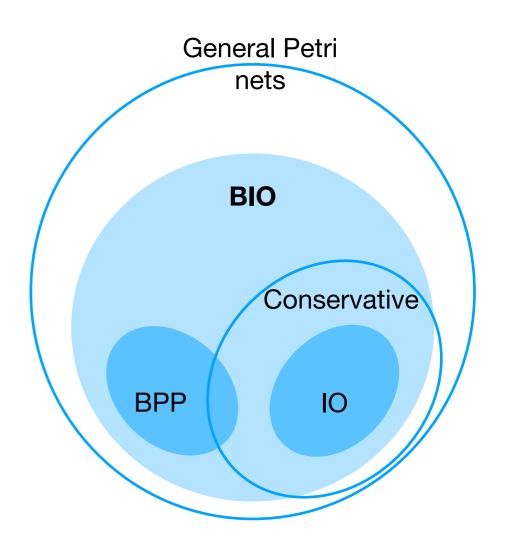
BIO nets can have non-semilinear reachability set



- unbounded (token creation and destruction)
- pre\*-flat reachability relation → use of model-checking tools like FAST
- PSPACE-complete reachability problem
- non-semilinear reachability



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- **Open Problems:** Applications for BIO nets (e.g. chemical reaction networks)
  - Consequences of this result in other domains (data nets, process calculi, formal languages...)

Article: https://drops.dagstuhl.de/opus/volltexte/2020/12857/

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