

# Euler Graph & Euler Tour

Joy Mukherjee

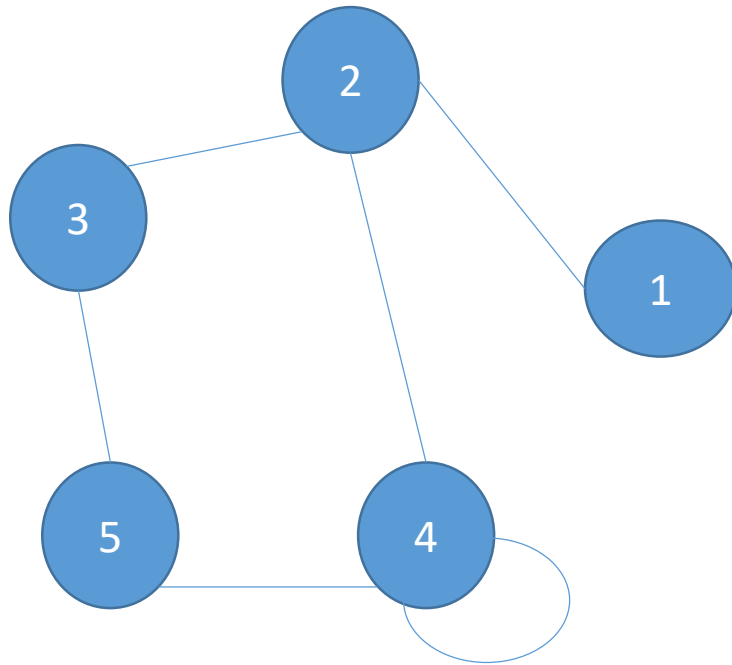
# Walk and Trail

- A **walk** of length  $k$  is a sequence,  $v_0, e_1, v_1, e_2, \dots, e_k, v_k$  of vertices and edges such that  $e_i = v_{i-1}v_i$  for all  $i$
- A **trail** is a walk with no repeated edge
- A **path** is a walk with no repeated vertex
- A walk is **closed** if it has length at least one and its endpoints are equal
- A cycle is a **closed** trail in which “first = last” is the only vertex repetition
- A loop is a cycle of length one

# Path & Cycle

- A **path** in a graph is a single vertex or an ordered list of distinct vertices  $v_1, \dots, v_k$  such that  $v_{i-1}v_i$  is an edge for all  $2 \leq i \leq k$ .
- The ordered list is a **cycle** if  $v_kv_1$  is also an edge.
- A path is an **u, v-path** if  $u$  and  $v$  are respectively the first and last vertices on the path.
- A path of  $n$  vertices is denoted by  $P_n$ , and a cycle of  $n$  vertices is denoted by  $C_n$ .
- A graph  $G$  is **connected** if it has a  $u, v$ -path for each pair  $u, v \in V(G)$ .

# Walk, Trail, Path, Cycle



**Walk:** 1 2 3 5 3 2 4 5

**Closed Walk:** 2 3 2

**Trail:** 1 2 3 5 4 4 2 (No repeated edge)

**Closed Trail:** 4 2 3 5 4 4

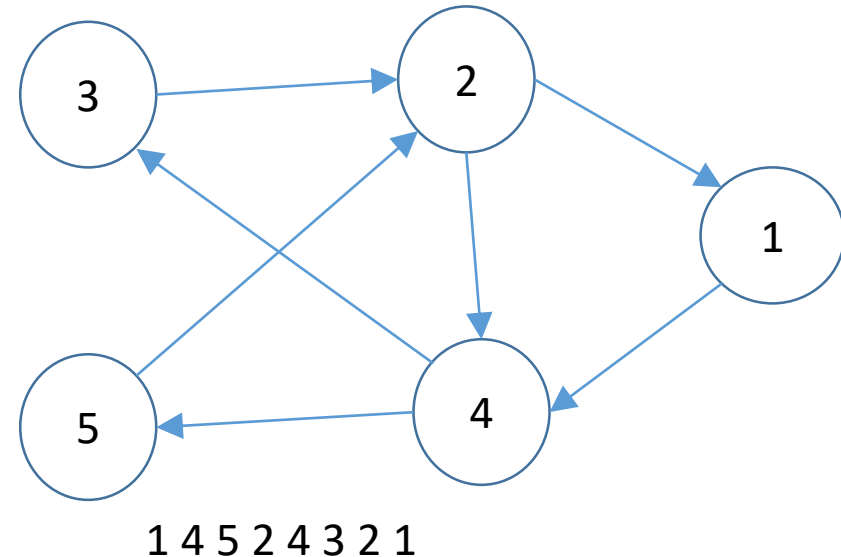
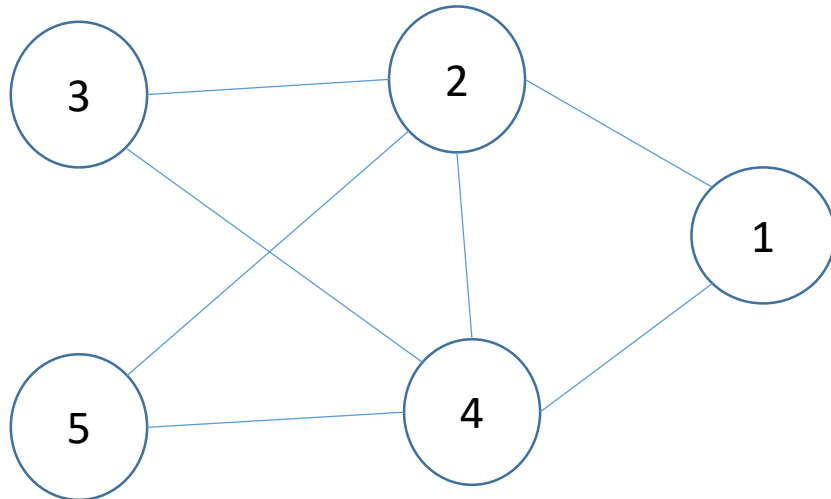
**Path:** 1 2 3 5 4 (No repeated vertex)

**Cycle:** 4 4 (self loop: cycle of length 1)

2 3 5 4 2 (Closed Path)

# Euler Graph

- A trail is a walk with no repeated edges.
- A graph is **Eulerian** if it has a **closed trail containing all edges**.
- An **even graph** is a graph with vertex degrees all even.
- A vertex is odd (or even) when its degree is odd (or even).



# Euler Trails in Directed Graphs

**Input:** A connected digraph  $G$  with  $d^+(u) = d^-(u)$  for all  $u \in V(G)$ .

**Output:** A directed Euler trail

**Algorithm:**

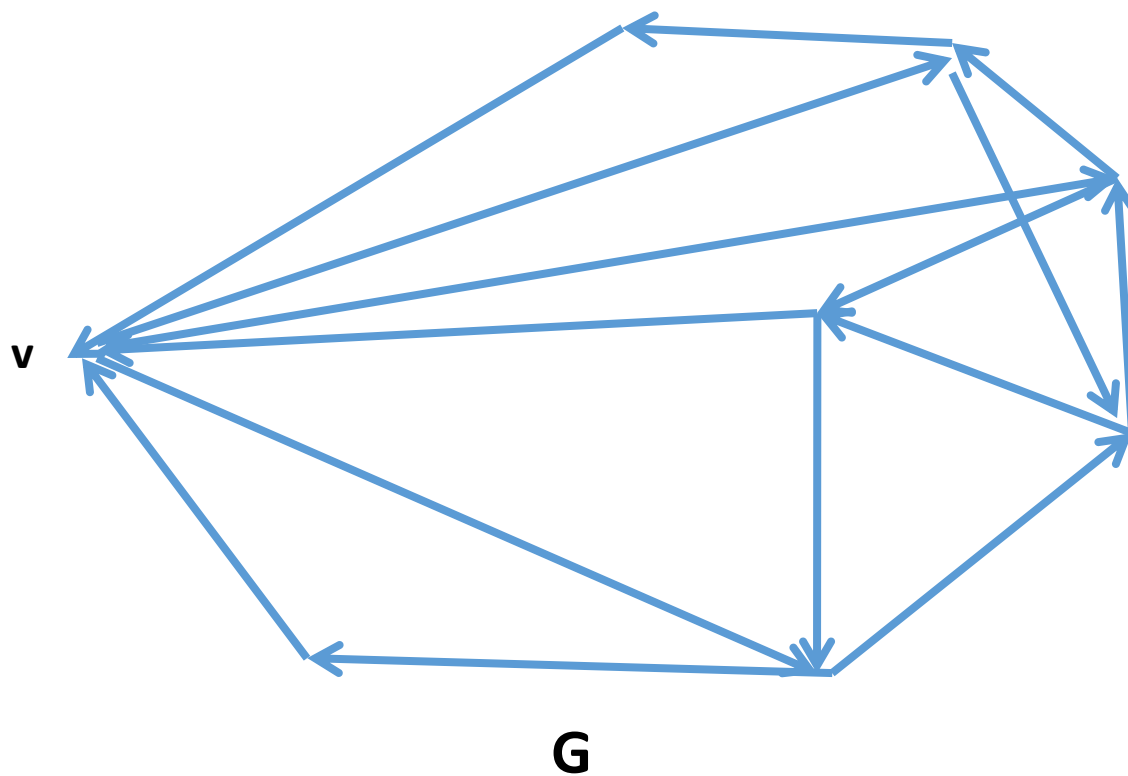
**Step 1:**  $G' =$  Digraph obtained from  $G$  by reversing direction of each edge.

**Step 2:** DFS on  $G'$  to construct  $T'$  consisting of paths from  $s$  to  $V(G')-s$ .

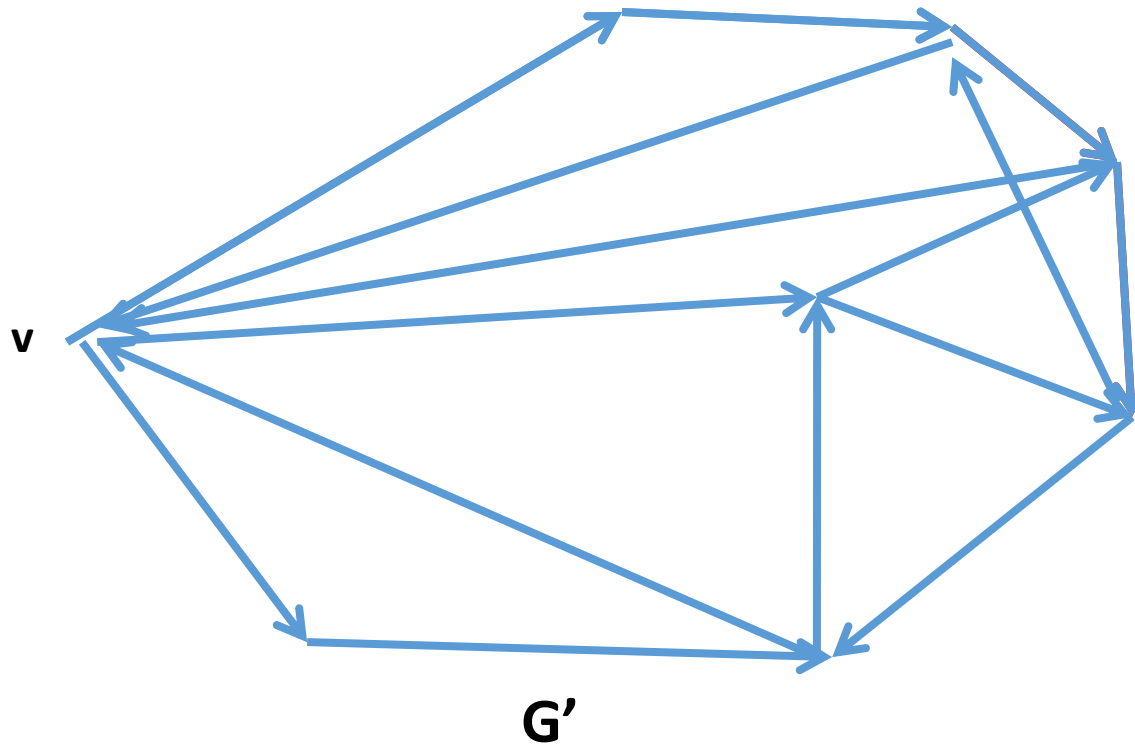
**Step 3:** Let  $T$  be the reversal of  $T'$ .  $T$  contains a  $u, s$ -path in  $G$  for each  $u \in V(G)$ .

**Step 4:** Construct an Eulerian circuit from  $s$  as follows: Whenever  $u$  is the current vertex, exit along the next unused edge of  $G - T$  as long as possible. If there is no other edge from  $u$ , then use edges in  $T$ .

# Euler Trails in Directed Graphs

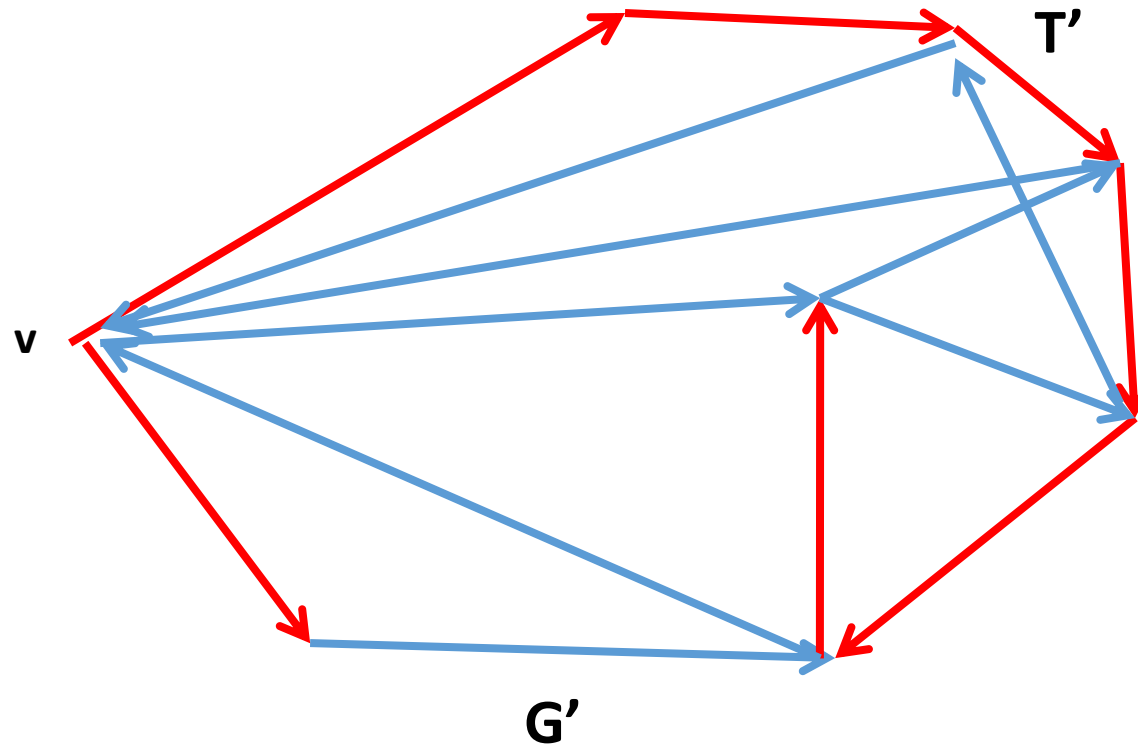


Reverse the edges



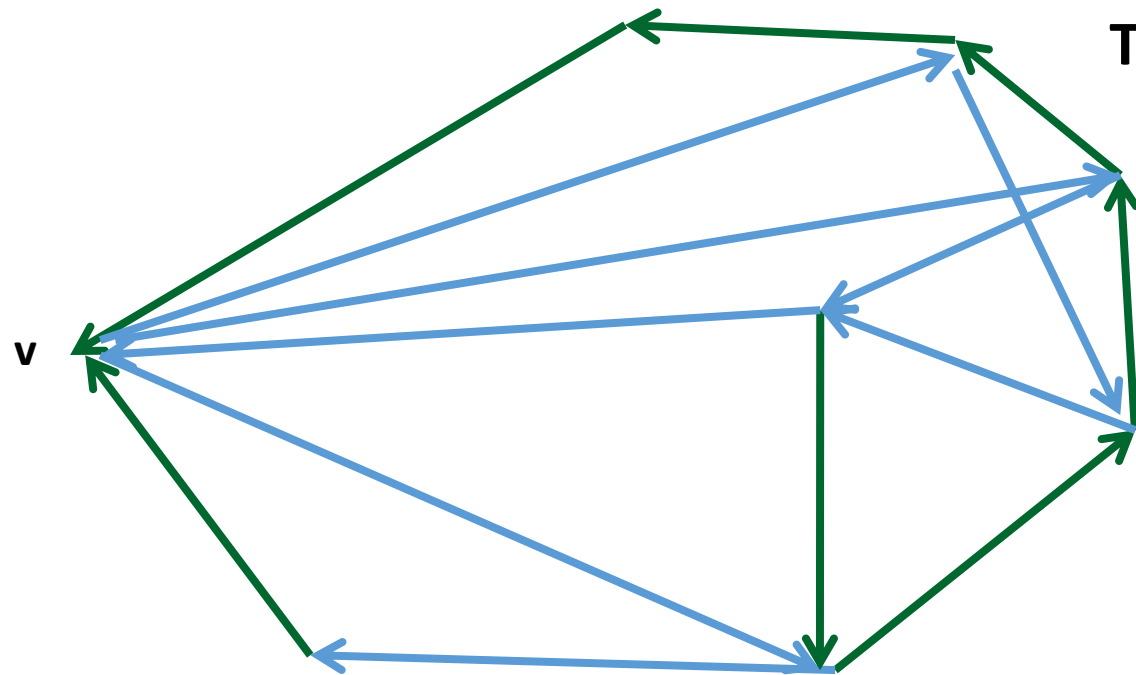


# DFS



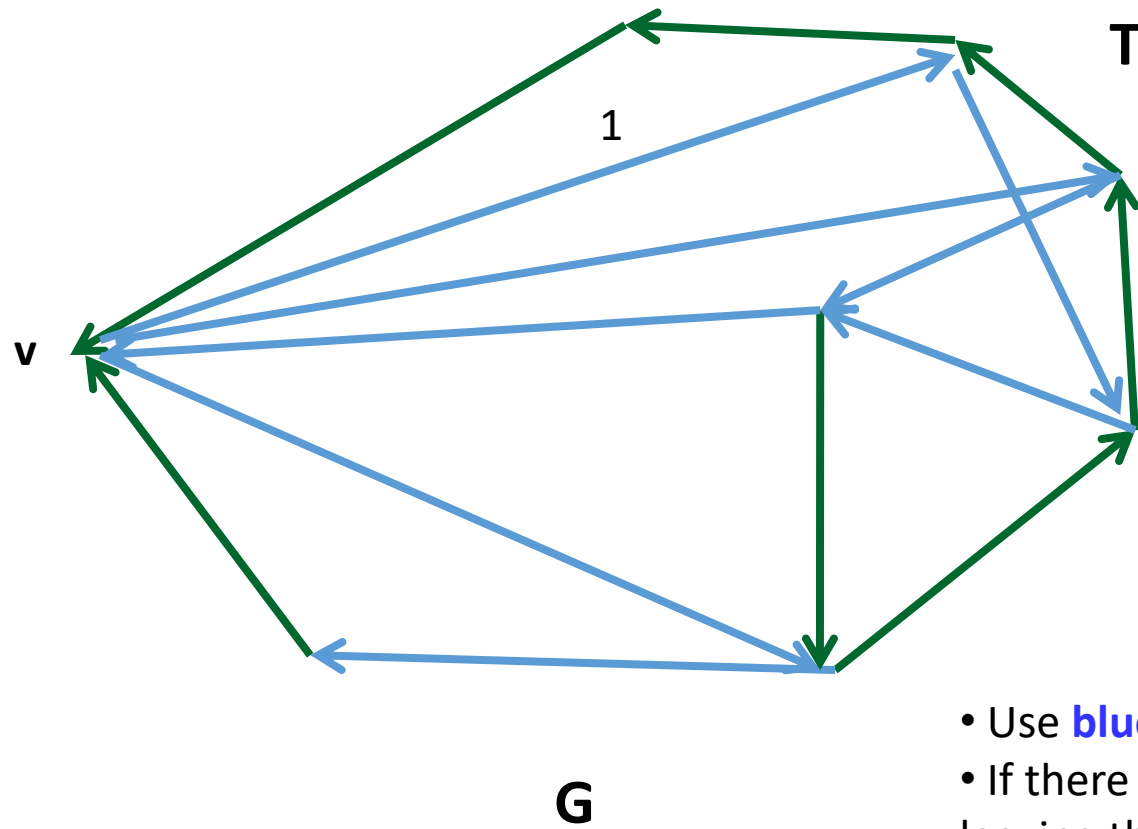
- DFS Tree  $T'$  consists of **red** edges

# Eulerian Circuit Construction



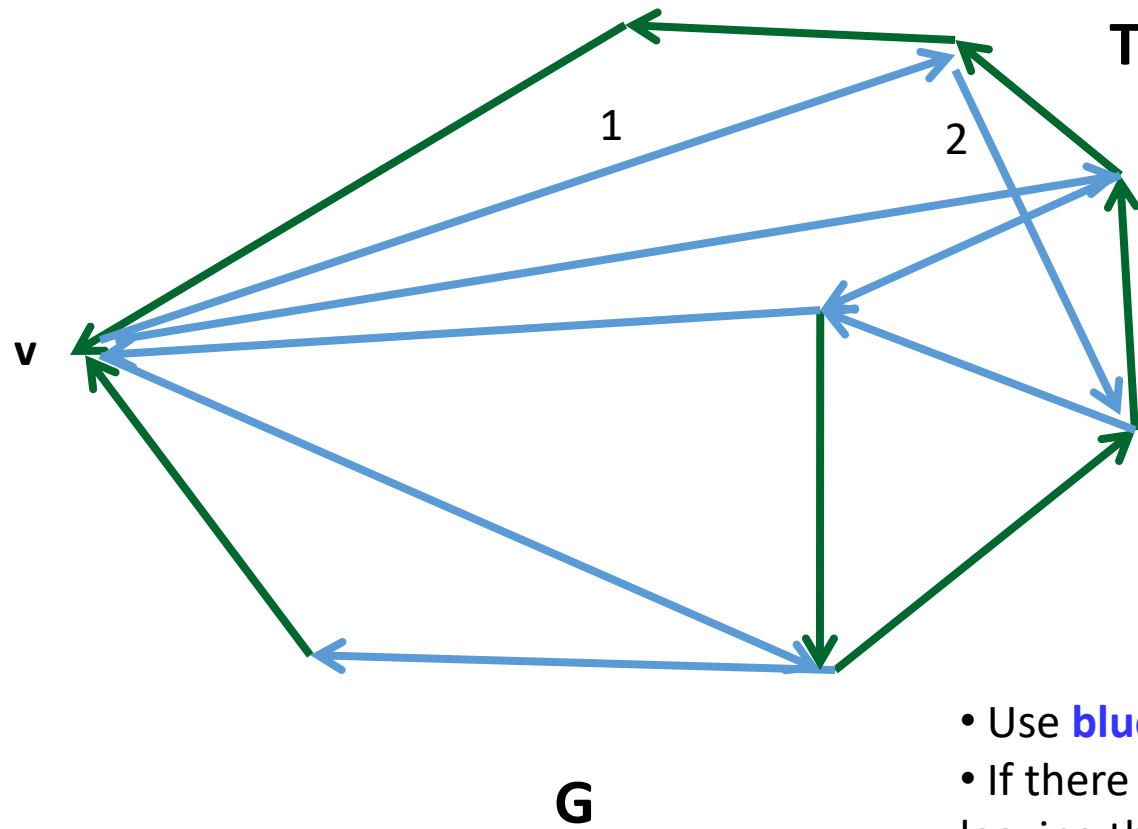
- Use **blue** edges for leaving a vertex.
- If there is no **blue** edges left for leaving the vertex, use **green** edges

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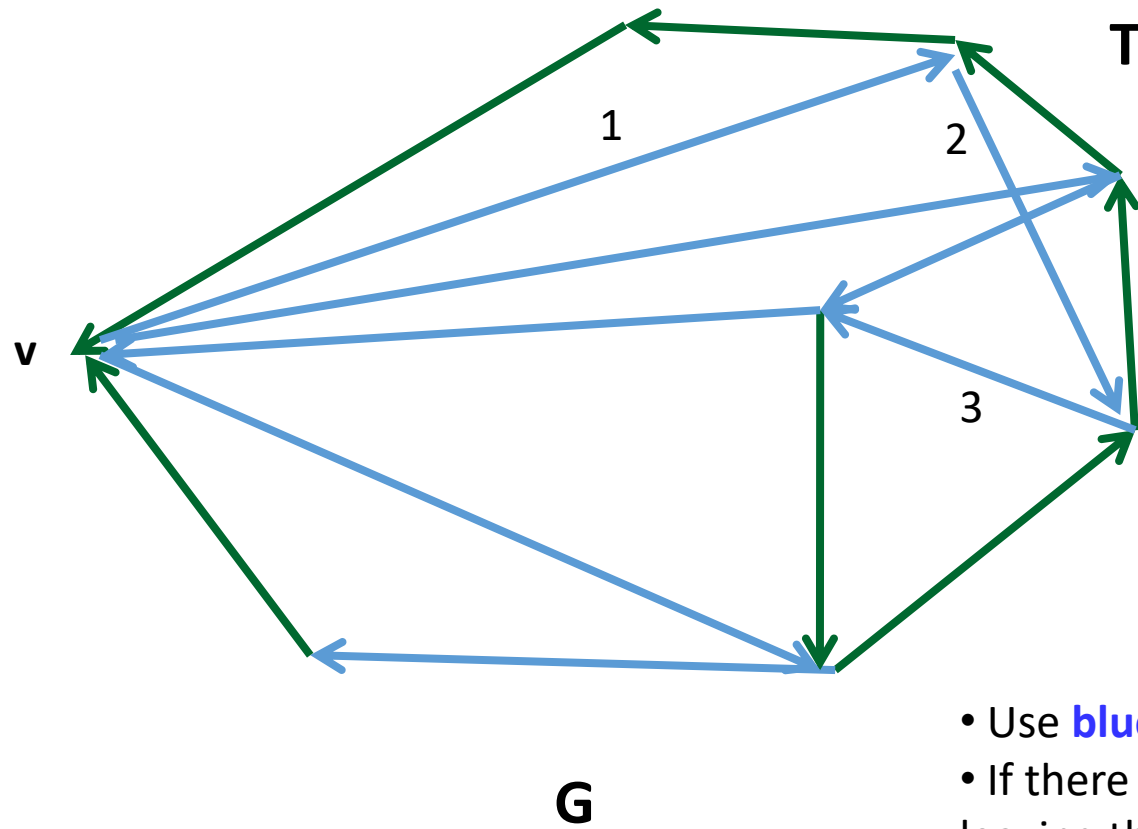
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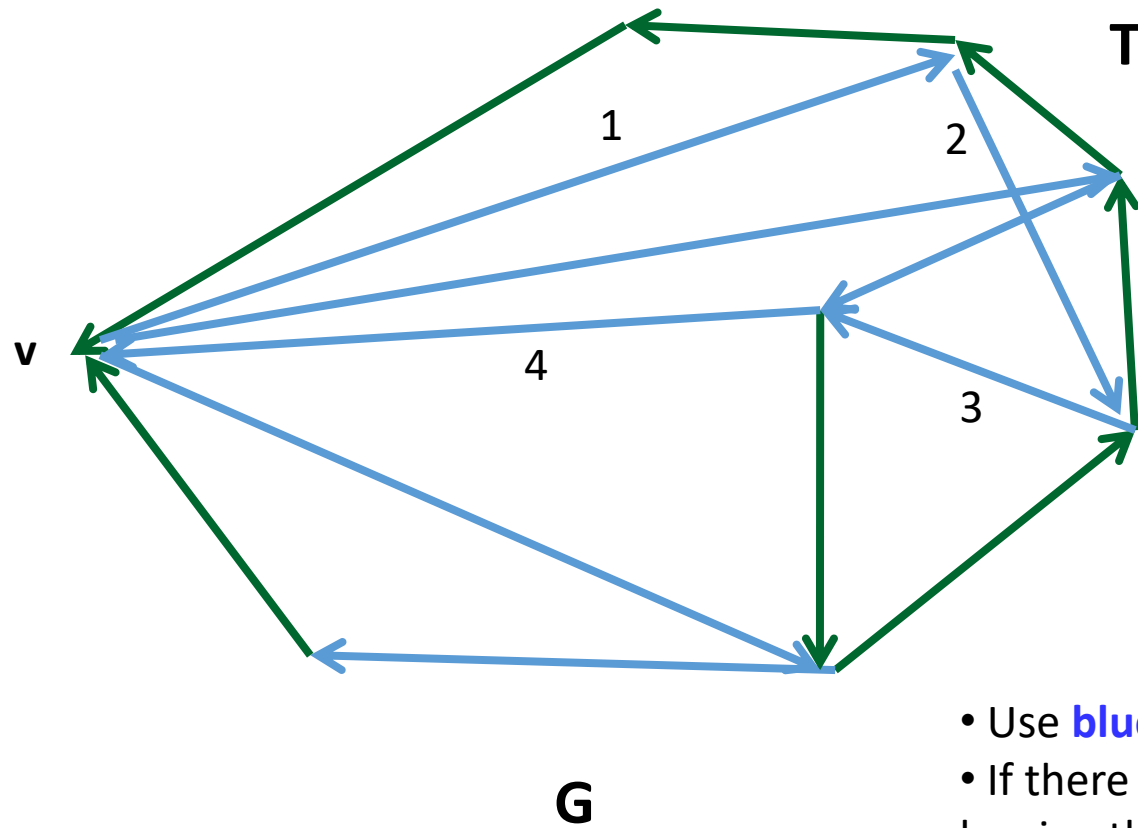
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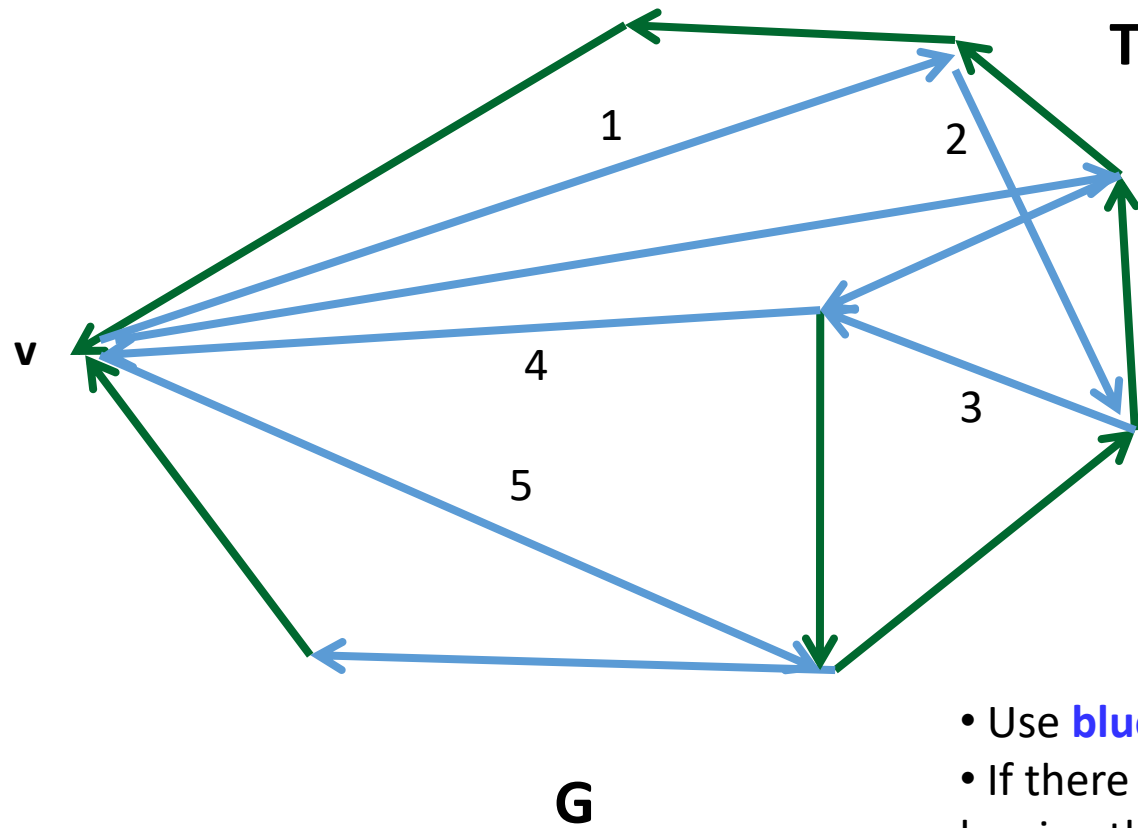
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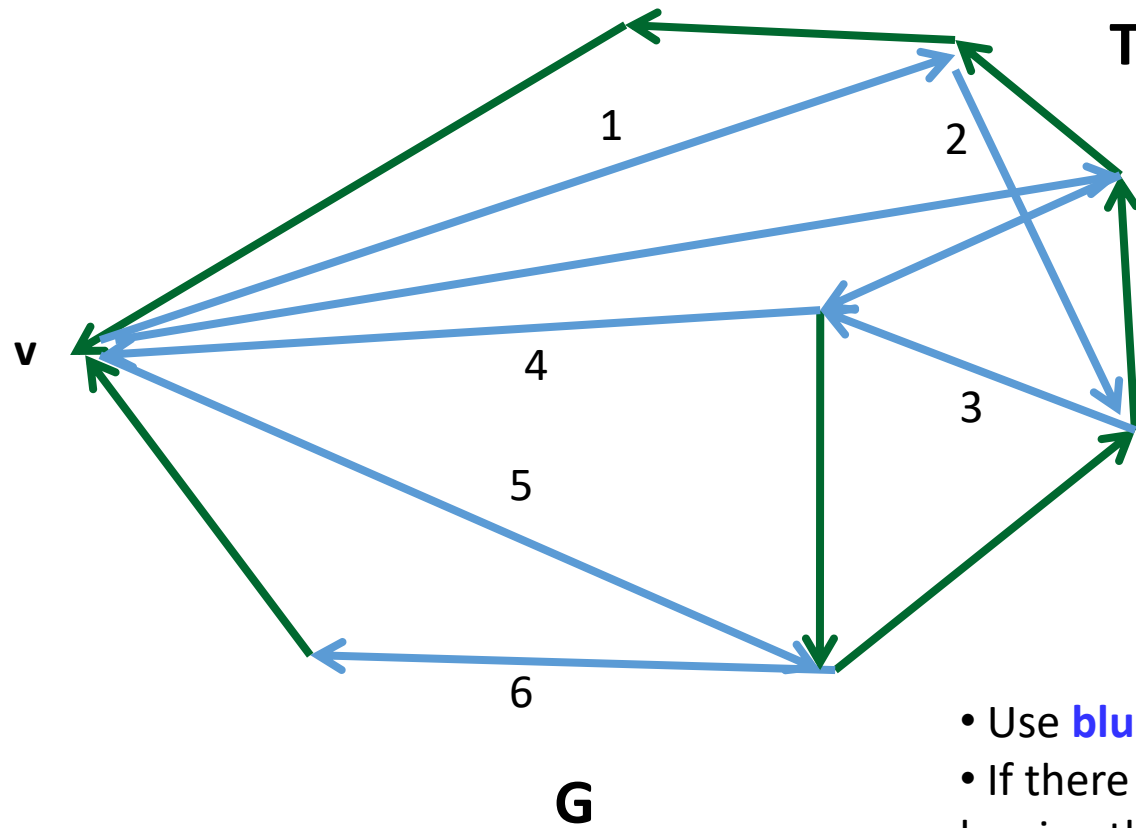
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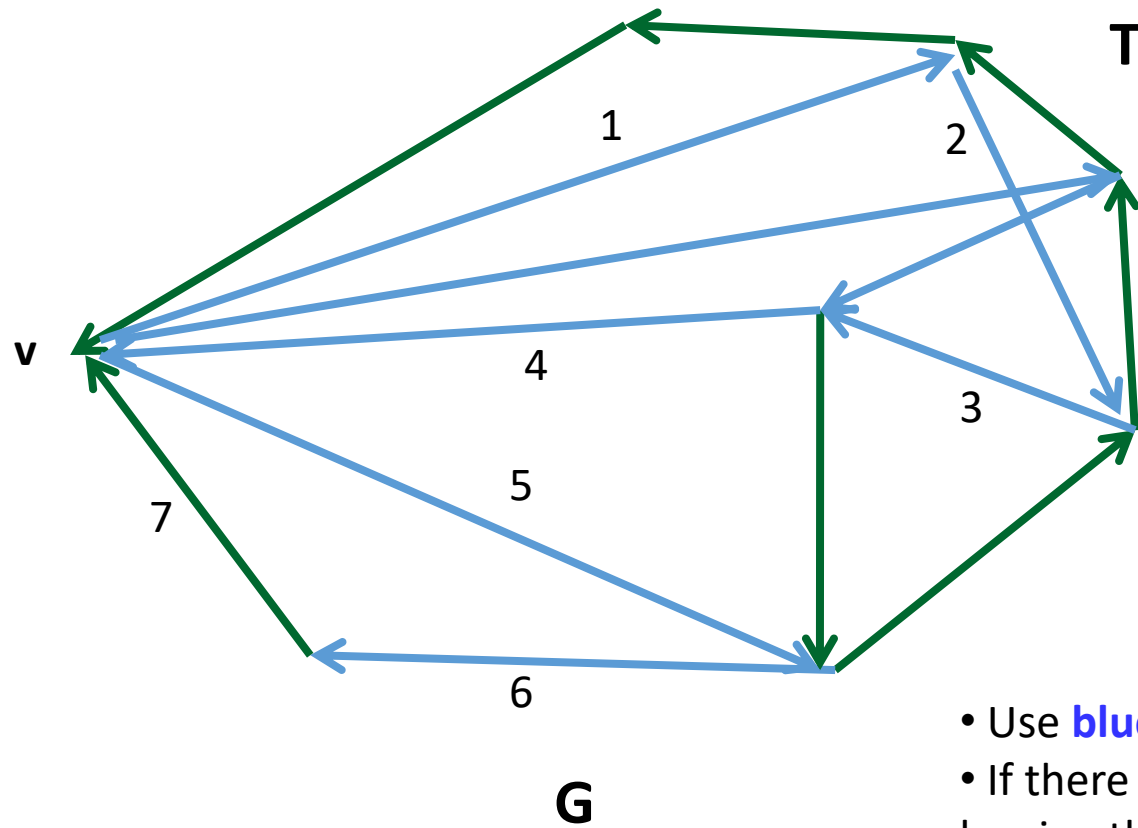
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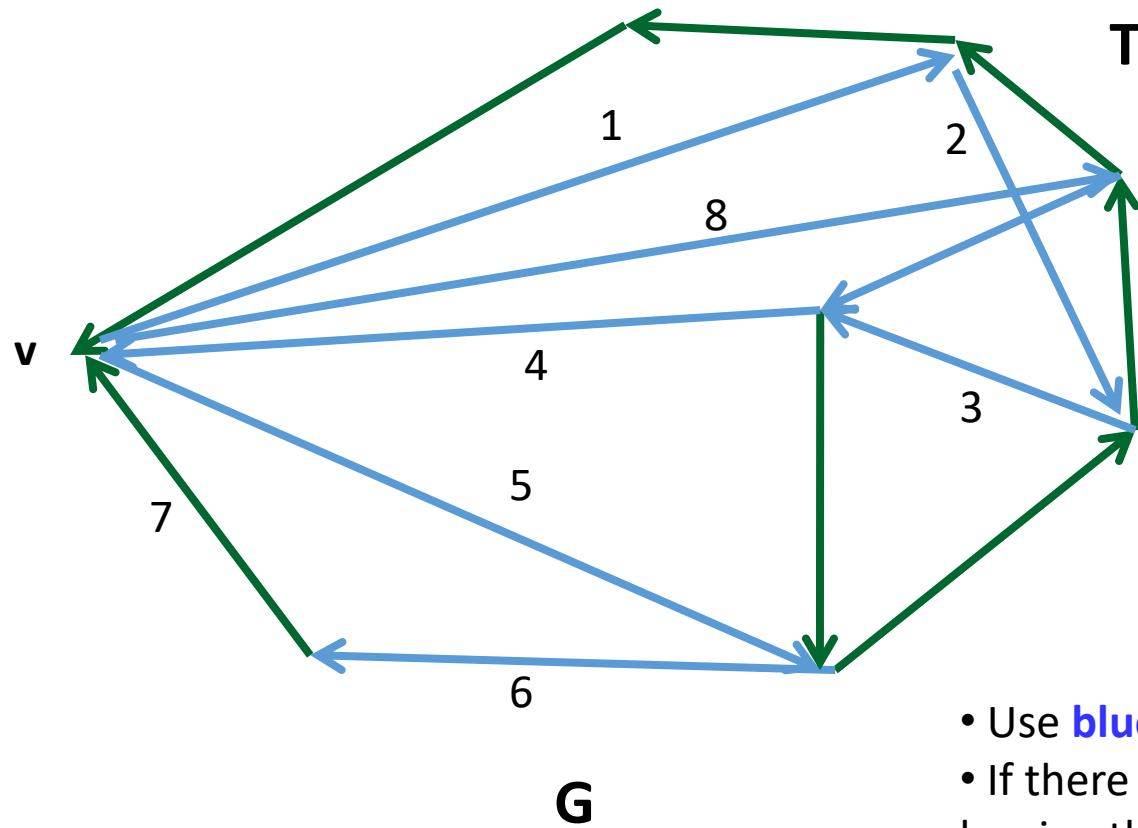


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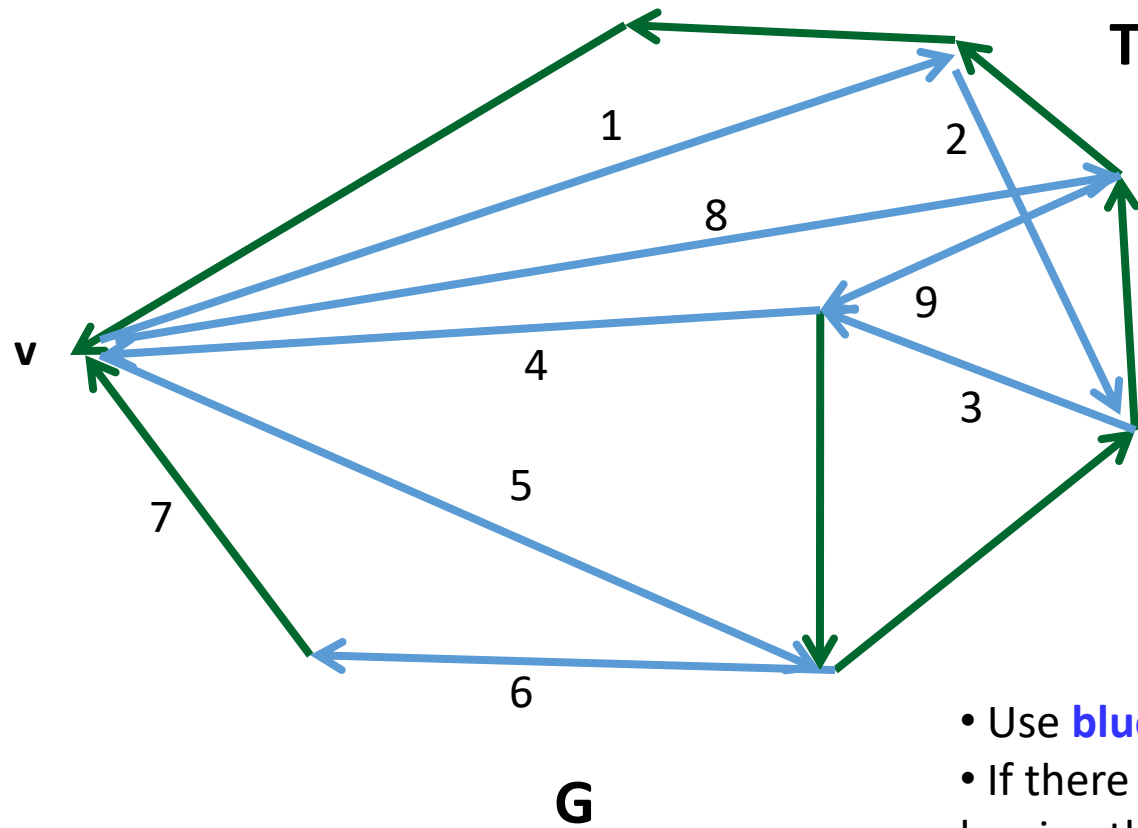
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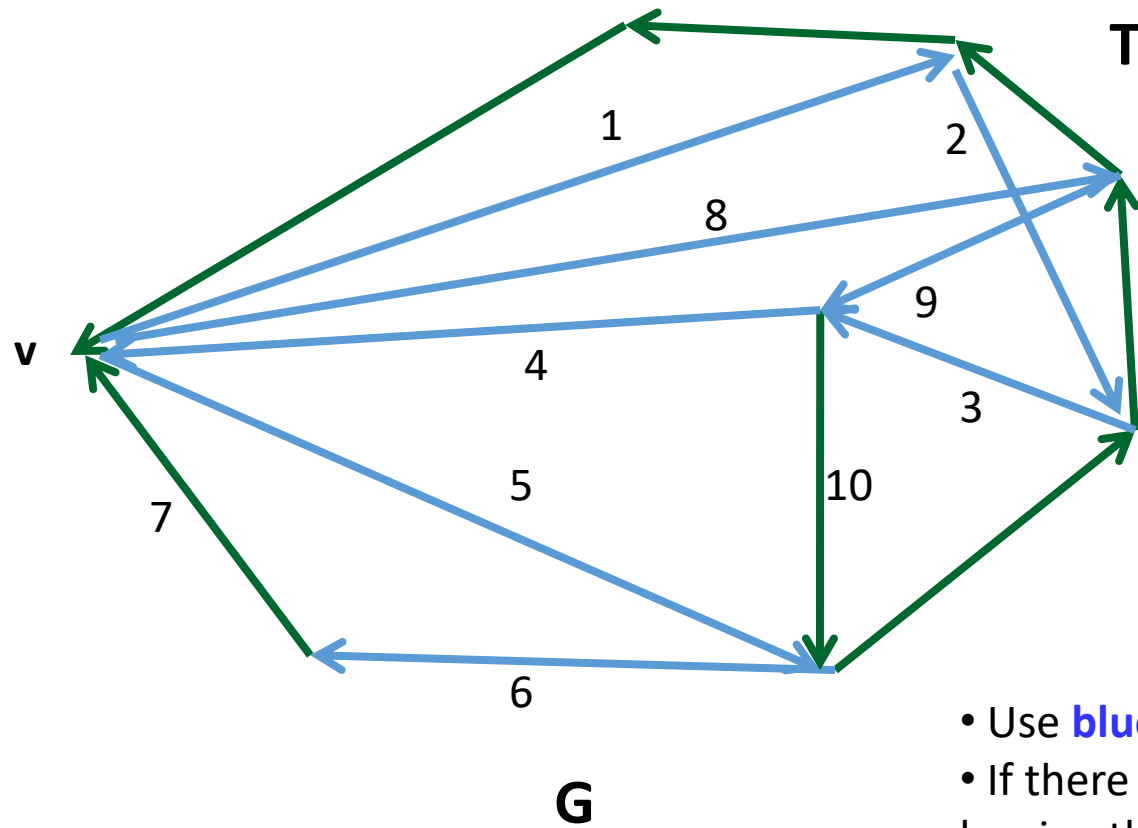
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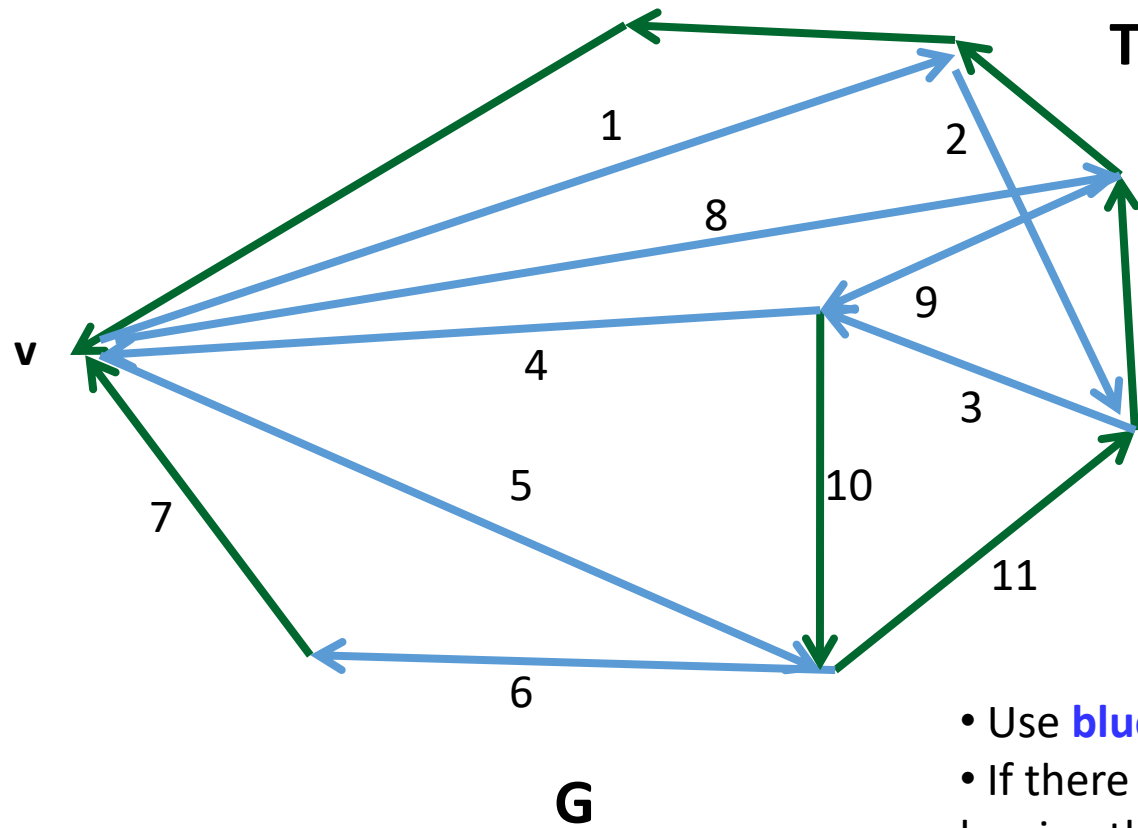
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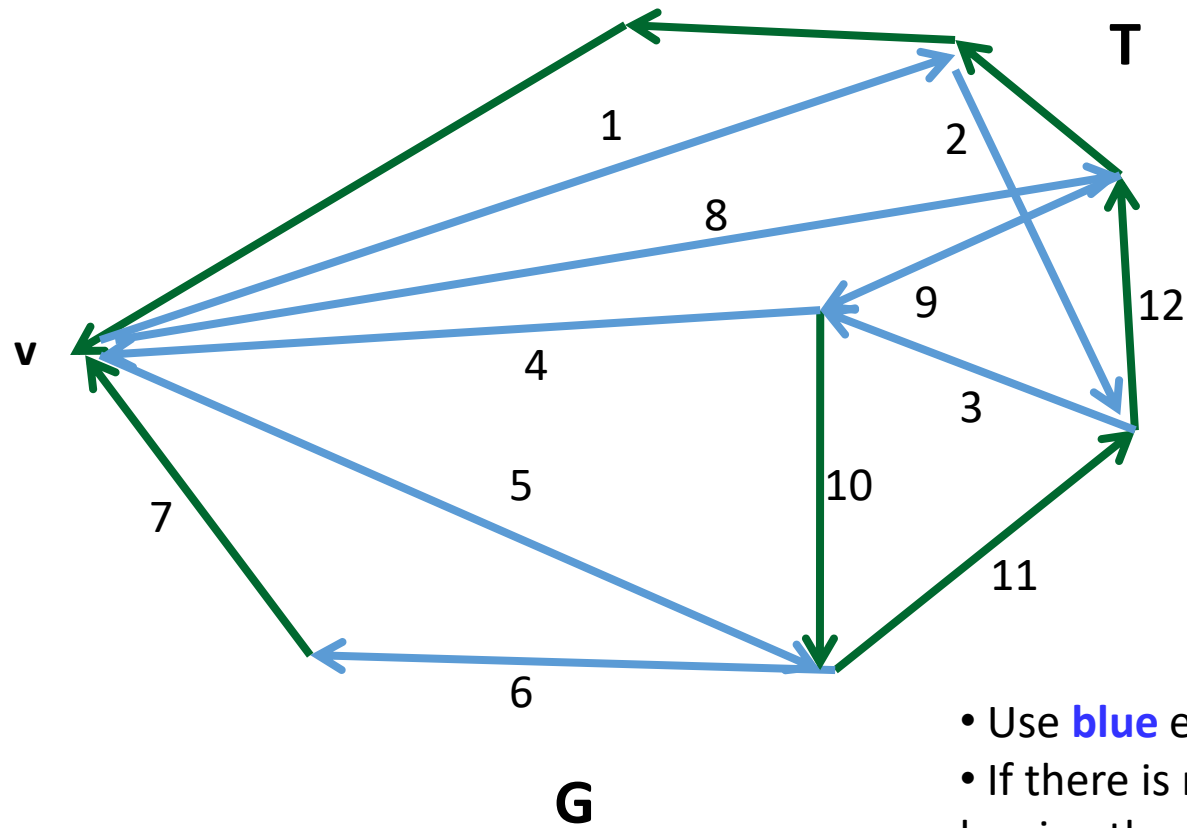
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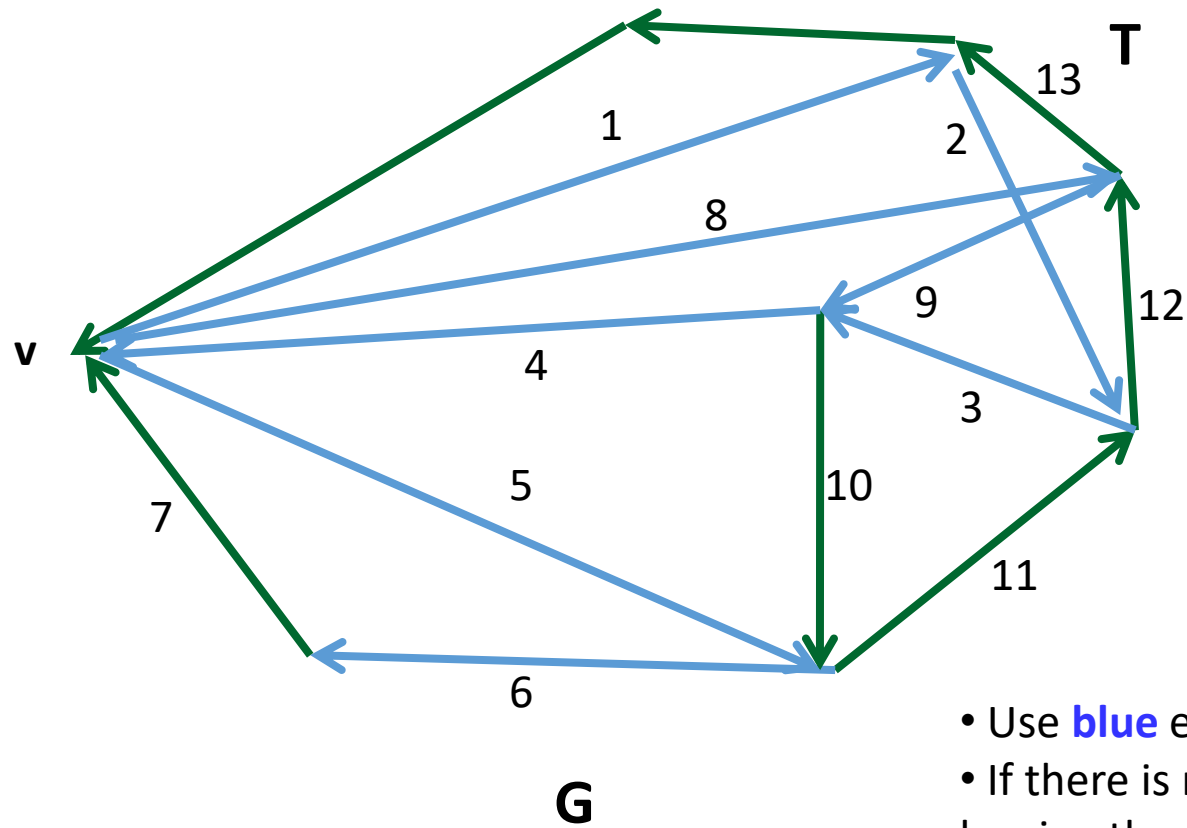
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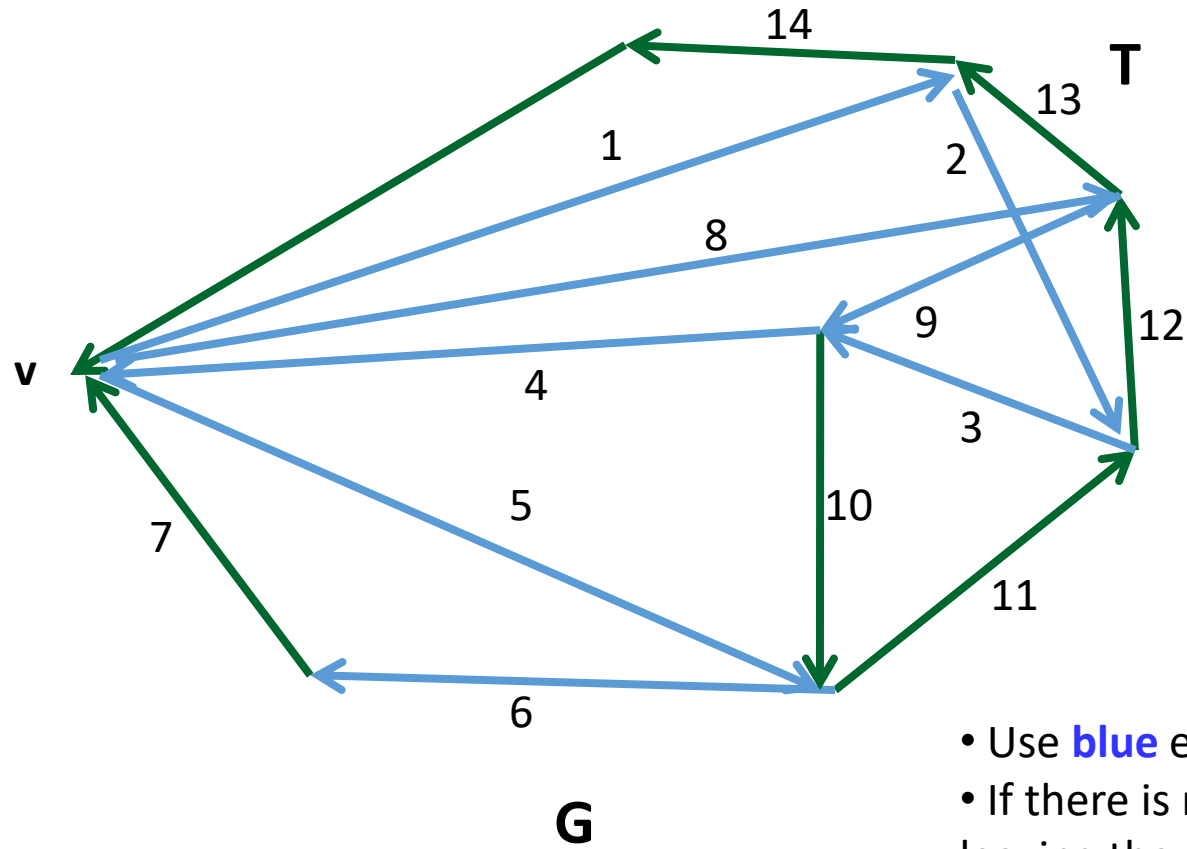
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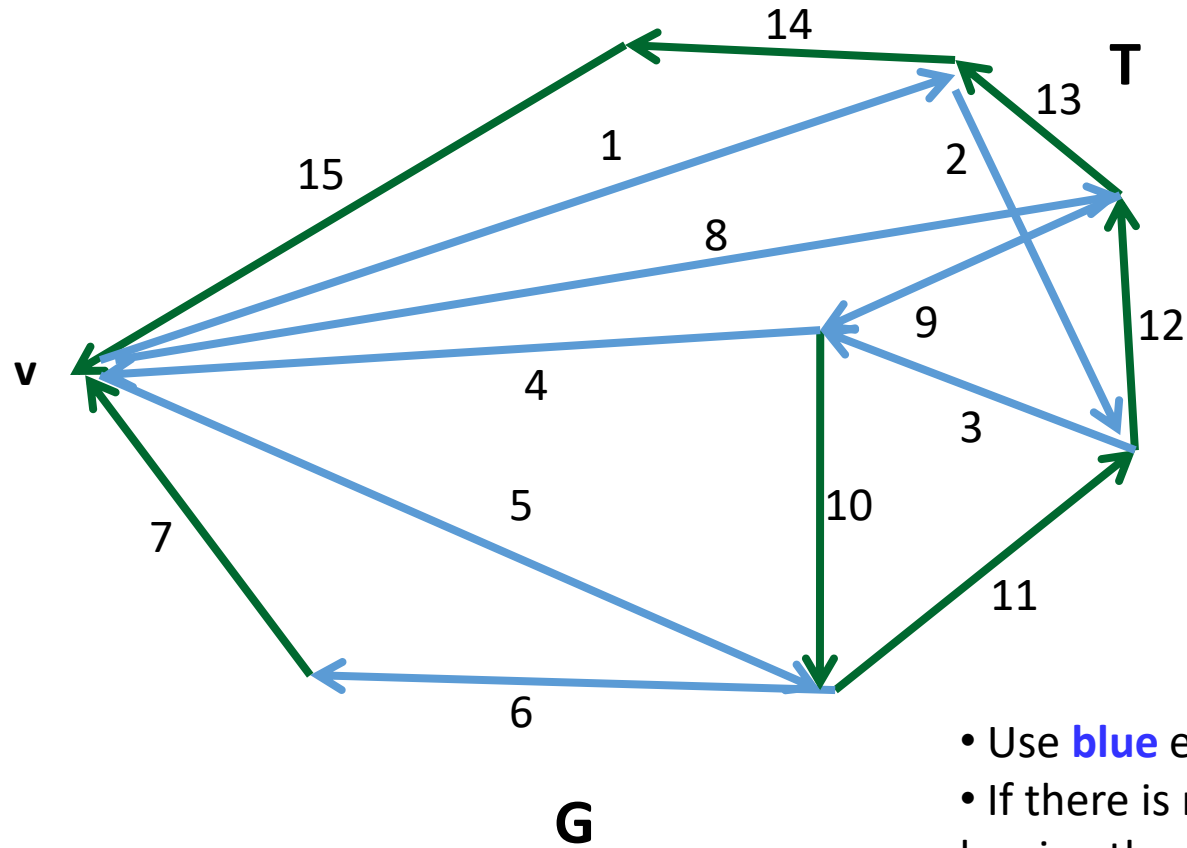
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