# A constrained optimization problem

The general form of a constrained optimization problem is given as:

(P1) Min 
$$f(x)$$
  
s/t  $g_i(x) \le 0$ ;  $i = 1, 2, ..., m$ 

where  $f: \mathbb{R}^n \longrightarrow \mathbb{R}, \ g_i: \mathbb{R}^n \longrightarrow \mathbb{R}; \ i = 1, 2, \dots, m.$ 

## Classification:

- If f and  $g_i(\forall i)$  are linear, then the problem (P1) is called a **linear** programming problem(LPP).
- If the problem (P1) is not linear, then we call it a non-linear programming problem(NLPP).
- If in (P1), f is quadratic and g<sub>i</sub>(∀i) are linear, then the problem is called quadratic programming problem(QPP). It is a special class of a non-linear programming problem.

# Convex Programming problem

If f and  $g_i$  (i = 1, 2, ..., m) in (P1) are convex functions, then we (P1) is called a **convex programming problem(CPP).** 

# Different formats of CPP

Optimization problem	Conditions for (CPP)
	$f$ and $g_i$ ( $\forall i$ ) are convex.
	$f$ is concave and $g_i$ ( $\forall i$ ) are convex.

$f$ is convex and $g_i$ ( $\forall i$ ) are concave.
$f$ and $g_i$ ( $\forall i$ ) are concave.

## **Examples:**

(P1) Min 
$$x_1 + x_2$$
 (P2) Max  $2x_1 - x_2$  subject to:  $x_1^2 + x_2^2 \le 4$ , subject to:  $x_1 + x_2 \le 3$ ,  $x_1^2 \le x_2$ ,  $x_1 x_2 \ge 0$ .  $x_1, x_2 \ge 0$ .

Here, (P1) is convex while (P2) is not a convex programming problem.

#### **Theorem**

Let  $g_i$  for each i = 1, 2, ...., m be a convex function. Then,

$$S = \{x \in \mathbb{R}^n : g_i(x) \le 0; i = 1, 2, ..., m\}$$

is a convex set.

(PI) Min 
$$f(x) = x_1 + x_2$$
  
SIt  $x_1^2 + x_2^2 \le y$   $g_1 = x_1^2 + x_2^2 - y \le 0$   
 $x_1^2 \le x_2$   $g_2 = x_1^2 - x_2 \le 0$   
 $x_1^2 \le x_2$   $g_3 = -x_1 \le 0$   
 $x_1^2 \le x_2$   $y_2^2 = -x_2 \le 0$   
 $y_3 = -x_1 \le 0$   
 $y_4 = -x_2 \le 0$ 

$$\nabla^2 g_1 = \begin{pmatrix} 2 & 0 \\ 0 & 2 \end{pmatrix}$$

Minor of order  $1 \times 1$ : 2, 2

Minor of order  $2 \times 2 = |\nabla^2 g|$ 
 $\Rightarrow \nabla^2 g_1$  is positive definite

 $\Rightarrow g_1$  is convex.

$$g_2 = \chi_1^2 - \chi_2$$

$$\nabla^2 g_2 = \begin{pmatrix} 2 & 0 \\ 0 & 0 \end{pmatrix}$$

Minors of order 1×1 c 2,0 Miners of order 2×2 = 0

=> 92 is a convex function

# Quadratic Programming Problem

The general form of a quadratic programming problem is given as:

(QPP) Min 
$$f(x) = c^T x + \frac{1}{2} x^T Q x$$
  
 $s/t$   $Ax \le b$ ,  $x \ge 0$ 

where  $c, x \in \mathbb{R}^n$ ,  $b \in \mathbb{R}^m$ , A is a matrix of order  $m \times n$  and Q is a symmetric matrix of order n.

#### Lemma

Let M be a symmetric positive semi-definite matrix of order n. Then, for any  $x, y \in \mathbb{R}^n$ ,

$$x^T M y \leq \frac{1}{2} [x^T M x + y^T M y]$$

#### **Theorem**

A quadratic programming problem is a **convex** programming problem when *Q* in the objective function is a symmetric **positive semi-definite matrix**.

## Examples:

(QP1) Min 
$$f(x) = x_1^2 + 2x_2^2 + 3x_3^2 + 4x_1x_3 + 2x_2x_3 + x_1 - x_2 + 2x_3$$
  
subject to  $x_1 + x_2 + x_3 \le 10$ ,  $x_1, x_2, x_3 \ge 0$ .

(QP2) Min 
$$f(x) = x_1^2 + 2x_2^2 - 2x_1x_2 - 6x_1 - 8x_2$$
  
subject to  $2x_1 - x_2 \le 13$ ,  $x_1, x_2 \ge 0$ .

Here, (QP2) is a convex QPP while (QP1) is not a convex QPP.

(PI) 
$$S = \nabla^2 f = \begin{pmatrix} 2 & 0 & 4 \\ 0 & 4 & 2 \\ 4 & 2 & 6 \end{pmatrix}$$

Minon of order 
$$1 \times 1$$
: 2, 4, 6

Minon of order  $2 \times 2$ :  $\begin{vmatrix} 2 & 0 \\ 0 & 4 \end{vmatrix} = 8$ 

$$\begin{vmatrix} 4 & 2 \\ 2 & 6 \end{vmatrix} = 24 - 4 = 20$$

$$\begin{vmatrix} 2 & 4 \\ 4 & 6 \end{vmatrix} = 12 - 16 = -\frac{4}{2}$$

# **Constrained Optimization Problem**

The general form of a constrained optimization is given as:

(P2) Min 
$$f(x)$$
  
s/t  $g_i(x) \le 0$ ;  $i = 1, 2, ..., m$ .

where f and  $g_i(\forall i) : \mathbb{R}^n \to \mathbb{R}$  are defined and continuously differentiable functions.

87

#### **Necessary Condition**

Let  $\bar{x}$  be a local min point of the problem at which basic constraint qualification holds. Then there exist multipliers (called KKT-multipliers)  $\bar{\lambda}_i$ , i = 1, 2, ..., m such that the following conditions hold:

- $g_i(\bar{x}) \leq 0, i = 1, 2, ..., m,$
- $\bar{\lambda}_i g_i(\bar{x}) = 0, i = 1, 2, ..., m,$
- $\bar{\lambda}_i \geq 0$  for all i.

These conditions are called KKT-conditions.

## **Sufficient Condition**

Let  $(\bar{x}, \bar{\lambda}_1, \bar{\lambda}_2, ..., \bar{\lambda}_m)$  satisfy the KKT-conditions (1) - (4). Let f and  $g_i(\forall i)$  be differentiable convex functions. Then  $\bar{x}$  is a global min point of the problem (P2).

#### Remark:

Without the convexity assumptions on f and  $g_i$ , the KKT conditions are not sufficient for a point  $\bar{x}$  to be a local min/global min point.

## For example:

Min 
$$-x_2$$
  
subject to:  $x_1^2 + x_2^2 \le 4$   
 $-x_1^2 + x_2 \le 0$ .

The point (0,0) satisfy KKT-conditions but it is not a local/global min point.

Min 
$$- \frac{1}{2} \frac{1}{2}$$

SIT  $\frac{1}{3} = \frac{1}{1} + \frac{1}{12} - \frac{1}{2} \le 0$ 
 $\frac{1}{3} = -\frac{1}{12} + \frac{1}{12} \le 0$ 

Thus,  $\frac{1}{3} = \frac{1}{12} + \frac{1}{12} = 0$ 
 $\frac{1}{3} = -\frac{1}{12} + \frac{1}{12} = 0$ 

#### **Problems**

Show that (3/2, 9/4)<sup>T</sup> is a unique global optimal solution for the following problem:

(P1) Min 
$$f(x) = \left(x_1 - \frac{9}{4}\right)^2 + (x_2 - 2)^2$$
  
subject to  $x_1^2 \le x_2$ ,  
 $x_1 + x_2 \le 6$ ,  
 $x_1, x_2 \ge 0$ .

Solve the following problem:

(P2) Min 
$$f(x) = x_1^2 + x_2^2 - 6x_1 - 4x_2 + 13$$
  
subject to  $x_1^2 + x_2^2 \le 52$ ,  
 $x_1, x_2 \ge 0$ .

Min 
$$f = (n_1 - 9/4)^2 + (n_2 - 2)^2$$
 $5/t \quad g_1 = n_1^2 - n_2 \leq 0$ 
 $g_2 = n_1 + n_2 - 6 \leq 0$ 
 $g_3 = -n_1 \leq 0$ 
 $g_4 = -n_2 \leq 0$ 
 $7/f = \begin{pmatrix} 2 & 0 \\ 0 & 2 \end{pmatrix} \longrightarrow positive definite \\ \Rightarrow f is convex$ 
 $\sqrt{2}g_1 = \begin{pmatrix} 2 & 0 \\ 0 & 0 \end{pmatrix} \longrightarrow positive sense-definite$ 
 $\Rightarrow g_1$  is convex.

Lagrange multiplier.  $3 \times \text{min} Z = f(x, y). \quad F_{x} = f_{x} - \lambda g_{x} - 0 = ) \quad f_{x} = \lambda g_{x}$   $S \cdot t \quad F_{y} = f_{y} - \lambda g_{y} = 0 = ) \quad f_{y} = \lambda g_{y}.$   $g(x, y) = c \quad F_{x} = - \left(g(x, y) - c\right) = 0 = \left(g(x, y) - c\right) = 0 = \left(g(x, y) - c\right) = 0 = 0$ Let à be the Lagrange multiplier let F (x,y,h)= f(x,y)-h[g(x,y)-c]

# Unconstrained optimization problems

Consider the following unconstrained minimization problem:

$$(P) \qquad \min_{x \in R^n} f(x).$$

The question arises how to find a point  $\bar{x} \in R^n$  which solves (or atleast approximately solves) (P). Because in general, our analytical approach may not work for all types of optimization problems. So, we move to search techniques.

## Basic scheme

A common basic scheme is of the form:

$$x_{k+1} = x_k + \alpha_k d_k$$

where  $x_k$  is the current solution,  $d_k$  is the direction of movement from  $x_k$  and  $\alpha_k > 0$  is the step size (distance upto which we move from  $x_k$  in the direction  $d_k$ ). How to find  $\alpha_k$  and  $d_k$  to find next iteration  $x_{k+1}$  such that we move to the solution of (P) in an efficient manner?

#### Descent property

An algorithm for solving (P) is said to have a descent property if  $f(x_{k+1}) < f(x_k)$  for all k. That is, as we proceed, the value of objective function should decrease.

## Order of convergence

Let a sequence  $\{x_k\}$  converge to a point  $\bar{x}$  and let  $x_k \neq \bar{x}$  for sufficiently large k. The quantity  $||x_k - \bar{x}||$  is called the error of the  $k^{th}$  iteration. Suppose there exist p and  $0 < \alpha < \infty$  such that

$$\lim_{k \to \infty} \frac{||x_{k+1} - \bar{x}||}{||x_k - \bar{x}||^p} = \alpha,$$

then p is called the order of convergence of the sequence  $\{x_k\}$ .

# Continued...

- For p = 1, we say that the sequence  $\{x_k\}$  is linearly convergent.
- For p = 2, the sequence  $\{x_k\}$  is called quadratic convergent.
- Larger the value of *p*, faster the algorithm will converge.

## Unimodal function

The function  $f:[a,b] \longrightarrow R$  is said to be a unimodal function if it has only one peak in the given interval [a,b].

Consider, a unimodal min function  $f : [a, b] \longrightarrow R$ . Then there exists  $a \le x \le b$ , such that

- $\bigcirc$  f is strictly decreasing in [a, x).
- ② f is strictly increasing in [x, b].

Similarly, we can define for a unimodal max function.

## Continued...

Let f(x) be the unimodal min function on the interval of uncertainty [a, b]. Take two distinct points (called experiments)  $x_1$  and  $x_2$  such that  $x_1 < x_2$ , then the following cases may arise

- $f(x_1) < f(x_2) \implies x_{min} \in [a, x_2]$
- $f(x_1) > f(x_2) \implies x_{min} \in [x_1, b]$
- $f(x_1) = f(x_2) \implies x_{min} \in [x_1, x_2].$

## Measure of effectiveness

The measure of effectiveness of any search technique,  $\alpha$  is defined as

$$\alpha = \frac{L_n}{L_0}$$

where,  $L_n$  is the width of interval of uncertainty after n-experiments and  $L_0$  is the initial width of uncetainty.

## Steepest Descent method

Consider the following unconstrained minimization problem:

$$\min_{x \in R^n} f(x)$$

where f has continuous first order partial derivatives in  $\mathbb{R}^n$ .

Choose the starting point as  $X_1$  and move toward the optimal point according to the following rule:

$$X_{k+1} = X_k + \lambda_k d_k$$

where  $d_k = -\nabla f(X_k)$  and  $\lambda_k$  is the optimal step size which can be obtained by  $\min\{f(X_k + \lambda_k d_k)\}.$ 

Stopping rule:  $||\nabla f(X_k)|| < \epsilon$  or  $||f(X_{k+1}) - f(X_k)|| < \epsilon'$ .

## Steepest Descent algorithm

- is globally convergent.
- has order of convergence unity.
- has descent property.

## Example

Use the steepest descent method to minimize  $f(x_1, x_2) = x_1^2 - x_1x_2 + x_2^2$  such that

$$|f(X_{k+1}) - f(X_k)| < 0.05$$
. Take  $X_1 = \left(1, \frac{1}{2}\right)^T$ .

Min 
$$f = \chi_1^2 - \chi_1 \eta_2 + \eta_2^2$$
  
 $X_1 = (1, \frac{1}{2})^T$   $\nabla f = (2\chi_1 - \chi_2 - \chi_1 + 2\chi_2)^T$   
 $\nabla f(X_1) = (\frac{3}{2} \circ)^T$   
 $d_1 = -\nabla f(X_1) = (-\frac{3}{2}, \circ)^T$   
 $X_2 = X_1 + d_1 d_1 = X_1 + d_1 (-\frac{3}{2}, \circ)$   
 $= (1\frac{1}{2})^T + d_1 (-\frac{3}{2}, \circ)^T$   
 $X_2 = (1 - \frac{3}{2}\chi_1)$ 

$$f(x_{2}) = (1 - \frac{3}{2} \alpha_{1})^{2} - (1 - \frac{3}{2} \alpha_{1})(\frac{1}{2}) + \frac{1}{9}$$

$$\frac{df}{d\alpha_{1}} = 0 \implies 2(1 - \frac{3}{2} \alpha_{1})(-\frac{3}{2}) + \frac{3}{9} = 0$$

$$\frac{d^{1}f}{d\alpha_{2}} > 0 \implies \text{minima}$$

$$-3 + \frac{9}{2} \alpha_{1} + \frac{3}{4} = 0$$

$$\implies \frac{9}{2} \alpha_{1} = 3 - \frac{3}{9} = \frac{9}{9}$$

$$(\alpha_{1} = \frac{1}{2})$$

$$x_{2} = (\frac{1}{9} \alpha_{1})^{2}$$

$$||f(x_{2}) - f(x_{1})||$$

#### Solution

$$f(X_1) = f\left(1, \frac{1}{2}\right) = \frac{3}{4}, \quad \nabla f(x_1, x_2) = (2x_1 - x_2, -x_1 + 2x_2)^T$$
and  $\nabla f(X_1) = \left(\frac{3}{2}, 0\right)^T = -d_1$ 

$$X_2 = X_1 + \lambda_1 d_1$$

$$= \left(1, \frac{1}{2}\right)^T + \lambda_1 \left(-\frac{3}{2}, 0\right)^T = \left(1 - \frac{3}{2}\lambda_1, \frac{1}{2}\right)^T.$$
Now, to determine,  $\lambda_1$ ,
$$f(X_2) = f\left(1 - \frac{3}{2}\lambda_1, \frac{1}{2}\right) = \left(\frac{2 - 3\lambda_1}{2}\right)^2 - \left(\frac{2 - 3\lambda_1}{4}\right) + \frac{1}{4}$$

$$\frac{df(X_2)}{d\lambda_1} = 0 \implies \lambda_1 = \frac{1}{2}.$$
Therefore,  $X_2 = \left(\frac{1}{4}, \frac{1}{2}\right)^T$ . Since  $|f(X_2) - f(X_1)| = 0.75 \neq 0.05$ 

#### Continued...

So find the next iteration,

Now,

$$X_3 = X_2 + \lambda_2 d_2$$

$$= \left(\frac{1}{4}, \frac{1}{2}\right)^T + \lambda_2 \left(0, -\frac{3}{4}\right)^T, d_2 = -\nabla f(X_2)$$

$$= \left(\frac{1}{4}, \frac{1}{2} - \frac{3}{4}\lambda_2\right)^T$$

$$f(X_3) = \frac{1}{16} - \left(\frac{1}{2} - \frac{3}{4}\lambda_2\right) \left(\frac{1}{4}\right) + \left(\frac{1}{2} - \frac{3}{4}\lambda_2\right)^2$$

$$\frac{df}{d\lambda_2} = 0 \implies \lambda_2 = \frac{1}{2}.$$

Hence, 
$$X_3 = \left(\frac{1}{4}, \frac{1}{8}\right)^T$$
. Also,  $|f(X_3) - f(X_2)| = \frac{9}{64} < 0.05$ .

# Newton's method

#### **Basic Scheme**

Newton's method is an iterative method used for finding real roots of the equation  $g(y) = 0, y \in \mathbb{R}$ . The iterative formula for finding roots is given as:

$$y_{k+1} = y_k - \frac{g(y_k)}{g'(y_k)}$$

where  $y_k$  is the current iterate or the current approximation.

## For unconstrained optimization

Consider the following unconstrained minimization problem:

$$(P) \min_{\mathbf{x} \in \mathbf{R}^n} f(\mathbf{x})$$

where  $f: \mathbb{R}^n \longrightarrow \mathbb{R}$  is a differentiable function. For solving (P), we have to find  $\bar{x} \in \mathbb{R}^n$  such that  $\nabla f(\bar{x}) = 0$ . So, by the Newton scheme (in numerical methods), we have

$$x_{k+1} = x_k - (H_f(x_k))^{-1} \nabla f(x_k).$$
 (1)

#### **Proof**

The quadratic approximation the function f in (P), in a neighbourhood of  $x_k$  by the Taylor series is given as:

$$f(x) \approx f(x_k) + (x - x_k)^T \nabla f(x_k) + \frac{1}{2} (x - x_k)^T H_f(x_k) (x - x_k).$$

For minimization,  $\nabla f(x) = 0$ . This implies,

$$\nabla f(x_k) + H_f(x_k)(x - x_k) = 0$$

$$\Longrightarrow H_f(x_k)(x - x_k) = -\nabla f(x_k)$$

$$\Longrightarrow x - x_k = -(H_f(x_k))^{-1} \nabla f(x_k)$$
or  $x_{k+1} = x_k - (H_f(x_k))^{-1} \nabla f(x_k)$ .

This method has order of convergence, p = 2 and it has descent property. For solving quadratic functions (involving positive definite quadratic form), it will take exactly one iteration to find the optimal solution.

## Example

Use Newton's method to minimize

$$f(x_1,x_2)=x_1^2-x_1x_2+3x_2^2,\ (x_1,x_2)\in R^2.$$

Take initial approximation  $x_1 = (1, 2)^T$ .

#### Solution

$$x_{k+1} = x_k - (H_f(x_k))^{-1} \nabla f(x_k).$$

$$H_f(x) = \begin{bmatrix} 2 & -1 \\ -1 & 6 \end{bmatrix}, \nabla f(x) = \begin{bmatrix} 2x_1 - x_2 \\ -x_1 + 6x_2 \end{bmatrix}$$

$$(H_f(x))^{-1} = \frac{1}{11} \begin{bmatrix} 6 & 1 \\ 1 & 2 \end{bmatrix}, \nabla f(x_1) = (0, 11)^T$$