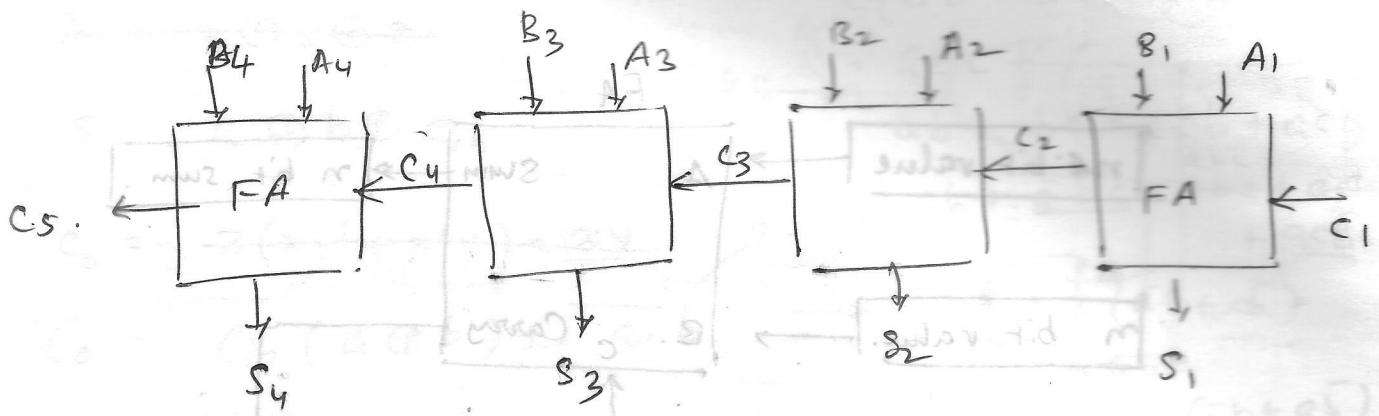


$$(x_1 + x_2)x_3 + x_4 = x \quad \therefore$$

say true at input 2

## 4-bit Binary Parallel Adder (Ripple carry Adder)



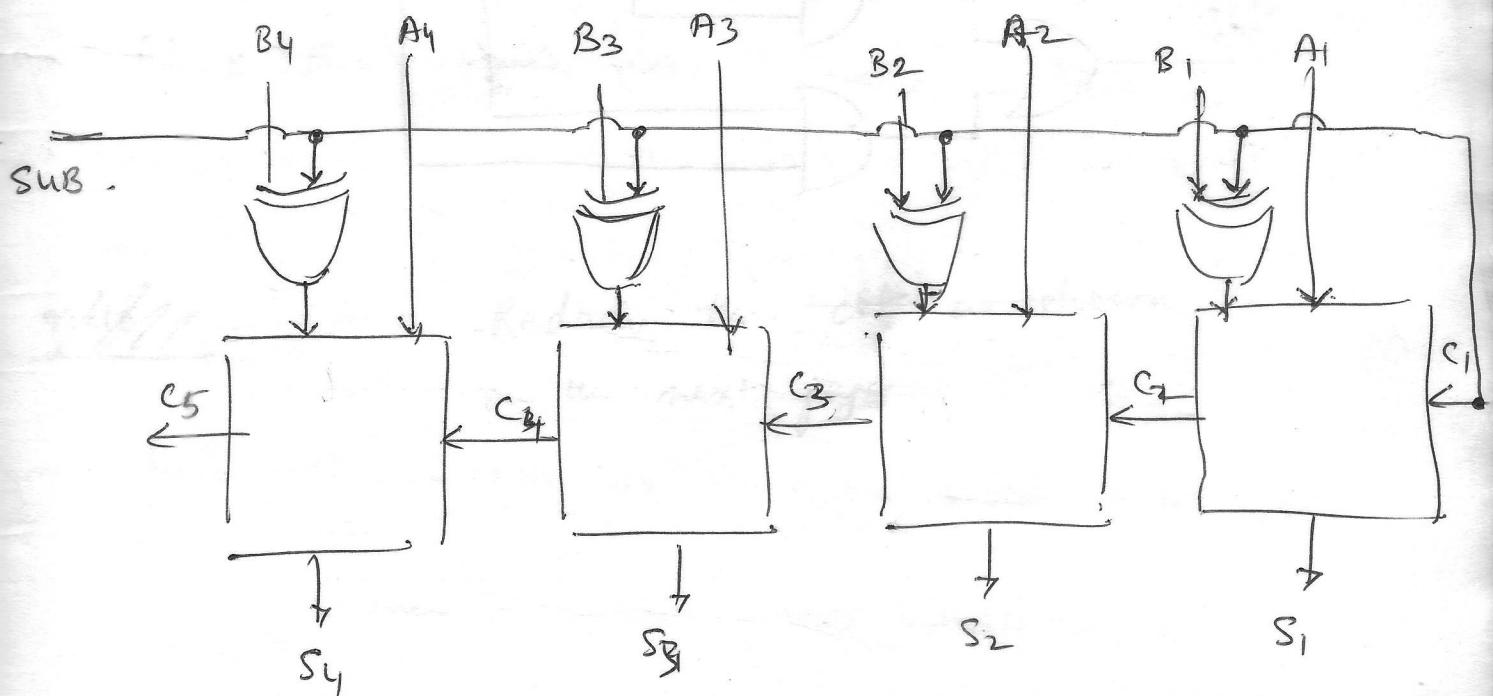
In general,

An  $n$ -bit adder requires  $n$  full adders.

Propagation delay =  $n \times (\text{FA}_{\text{delay}})$ .

∴ 32 bit adder  $\Rightarrow$  32 cycles!

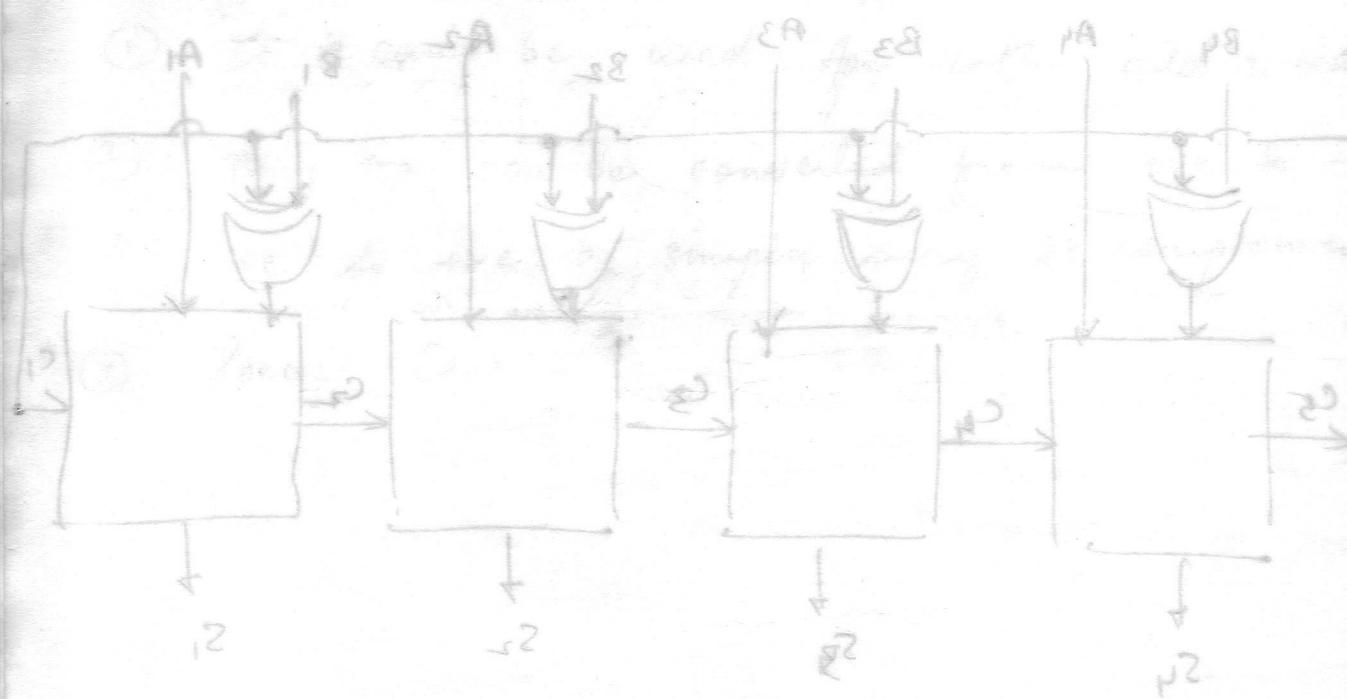
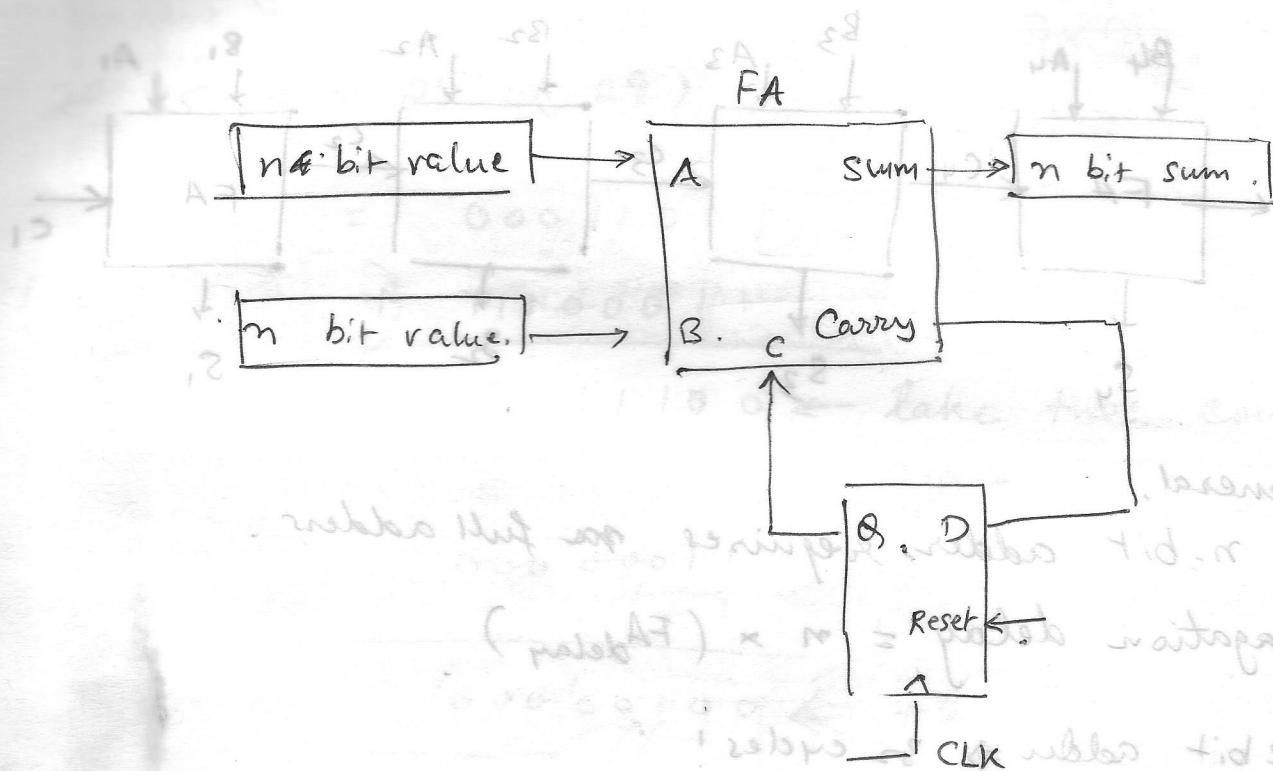
## 4-bit Binary Adder - Subtractor



$\text{SUB} = 0 \Rightarrow \text{Adder}$ .

$\text{SUB} = 1 \Rightarrow \text{Subtractor}$ .

# Serial Adder



$$nba \Leftarrow 0 = \alpha N^2$$

$$\rightarrow \text{robust?} \Leftarrow 1 - \alpha w^2$$

## Review of Full Adder

$$ab + bc + ac(\bar{b} + \bar{b})$$

$$S = x \oplus y \oplus z$$

$$S = a \oplus b \oplus c_{in}$$

$$C_0 = \cancel{z(x'y' + xy)} + \cancel{xy}$$

$$C_0 = C_{in}(a \oplus b) + ab$$

$$C_{out} = C_0 + C_{in}$$

x  
y → ADD

$$c(ab + \bar{a}\bar{b}) + ab$$

$$= abc + \cancel{ab} + \cancel{ac} + \cancel{bc}$$

$$= ab + bc + ac + abc$$

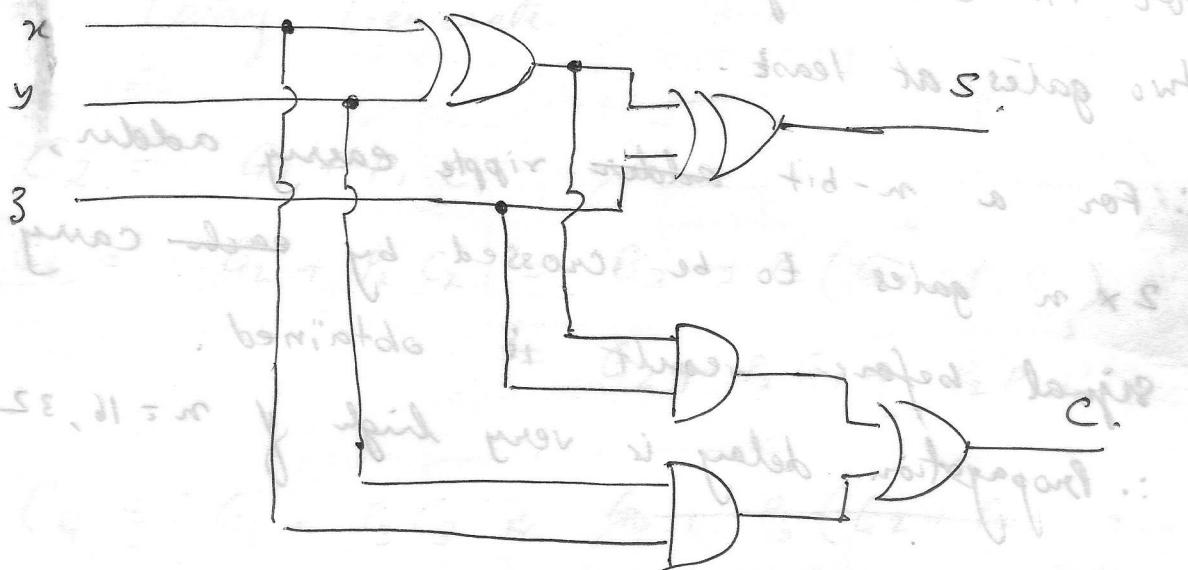
$$= ab + \cancel{\bar{a}bc} + \cancel{abc} + \cancel{acb}$$

$$+ \cancel{acb}$$

$$= ab + c(\bar{a}b + ab)$$

$$= ab + c(\bar{c}b + c\bar{b})$$

: then :

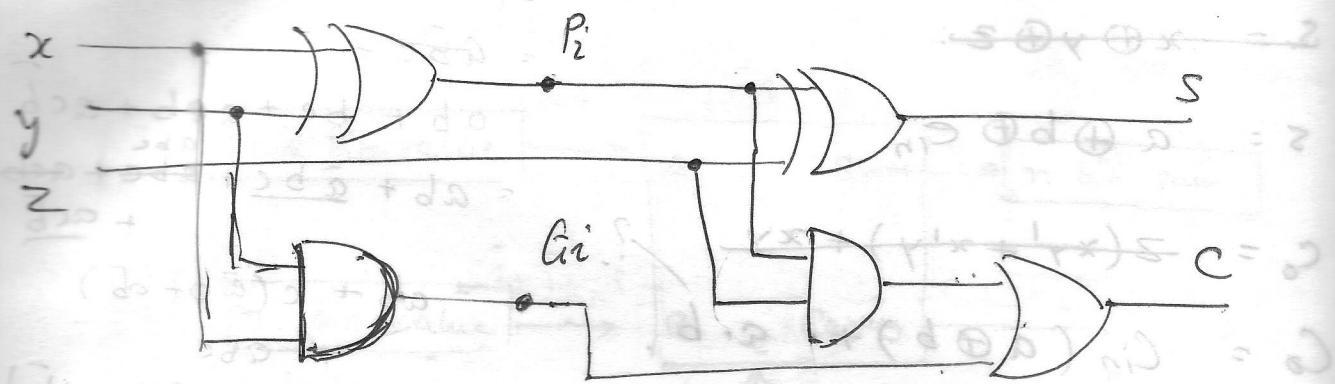


Notes:

Redraw this ckt as shown

on the next page

new w/ web w/ new books - look



$\therefore$  note :

- 1) For FA each signal has to go through two gates at least.
- 2)  $\therefore$  For a  $n$ -bit adder ripple carry adder,  $2 \times n$  gates to be crossed by each carry signal before result is obtained.  
 $\therefore$  Propagation delay is very high if  $n = 16, 32$  etc.

16.

### Ripple Carry Adder

- 1)  $2 \times n$  prop delay
- 2) Each o/p depends on previous sum and carry  
 $\therefore$  Look-ahead carry adder is used.

Let us define

$P_i = A_i \oplus B_i$ , where  $A_i$  and  $B_i$  are the  $i$ th bits  
in an  $n$ -bit number to be added.

Simil

$$G_i = A_i \cdot B_i$$

$$\therefore S_i = A_i \oplus B_i \oplus C_i = P_i \oplus C_i$$

$$C_{i+1} = G_i + P_i C_i$$

$P_i$  = carry propagate

$G_i$  = Carry Generate

$$C_2 = G_1 + P_1 C_1$$

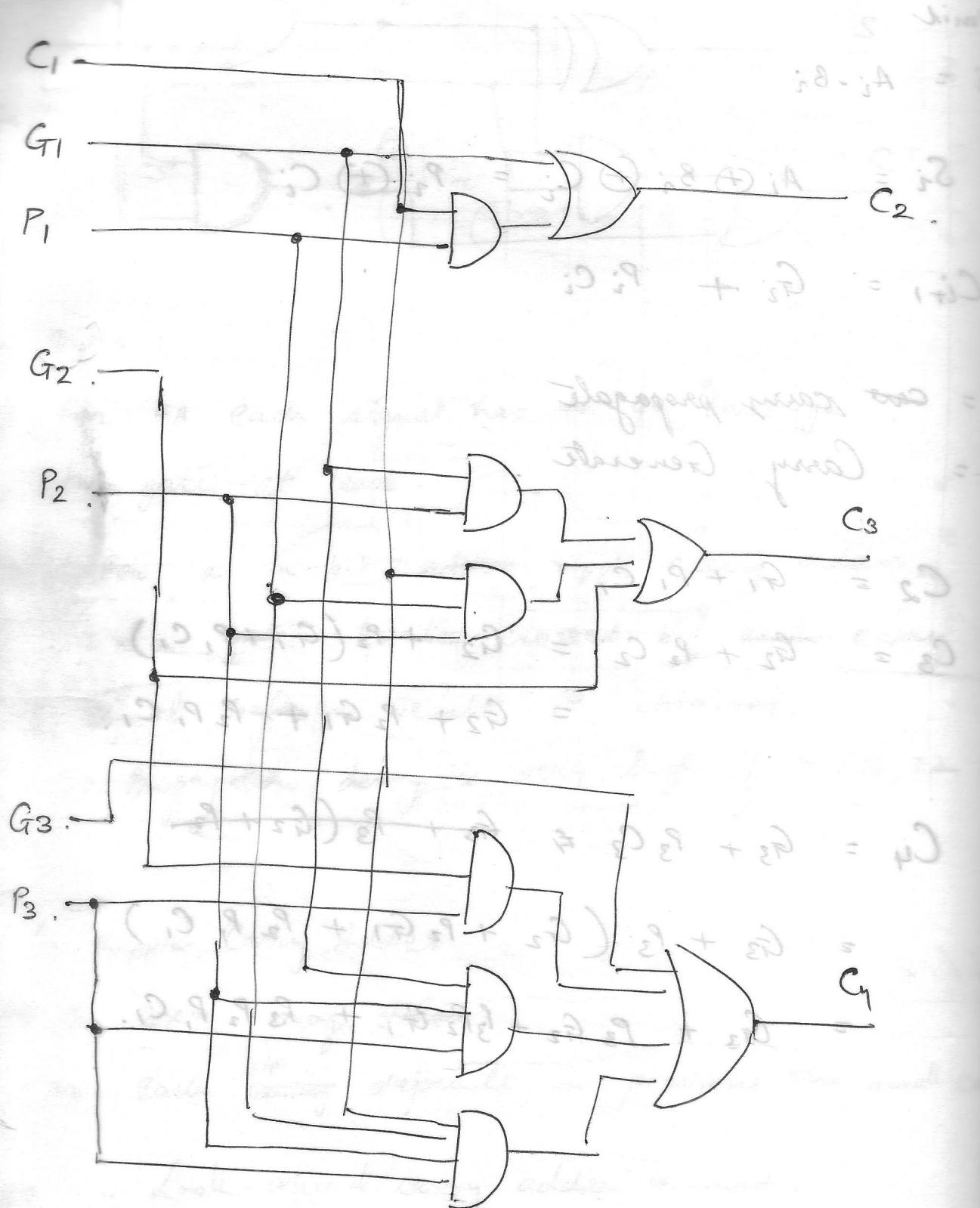
$$\begin{aligned} C_3 &= G_2 + P_2 C_2 = G_2 + P_2 (G_1 + P_1 C_1) \\ &= G_2 + P_2 G_1 + P_2 P_1 C_1. \end{aligned}$$

$$C_4 = G_3 + P_3 C_3 \Leftarrow G_3 + P_3 (G_2 + P_2)$$

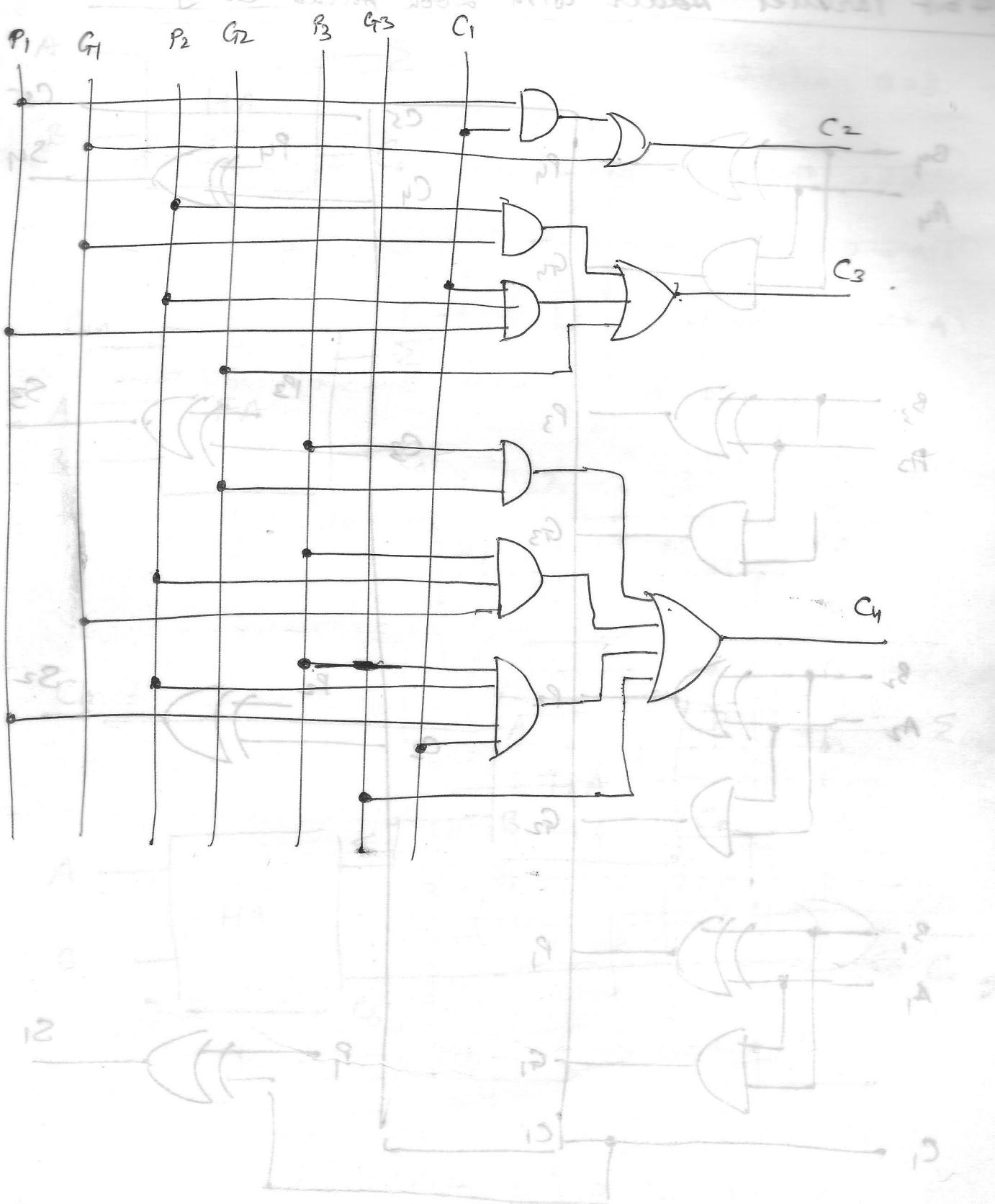
$$= G_3 + P_3 (G_2 + P_2 G_1 + P_2 P_1 C_1)$$

$$= G_3 + P_3 G_2 + P_3 P_2 G_1 + P_3 P_2 P_1 C_1.$$

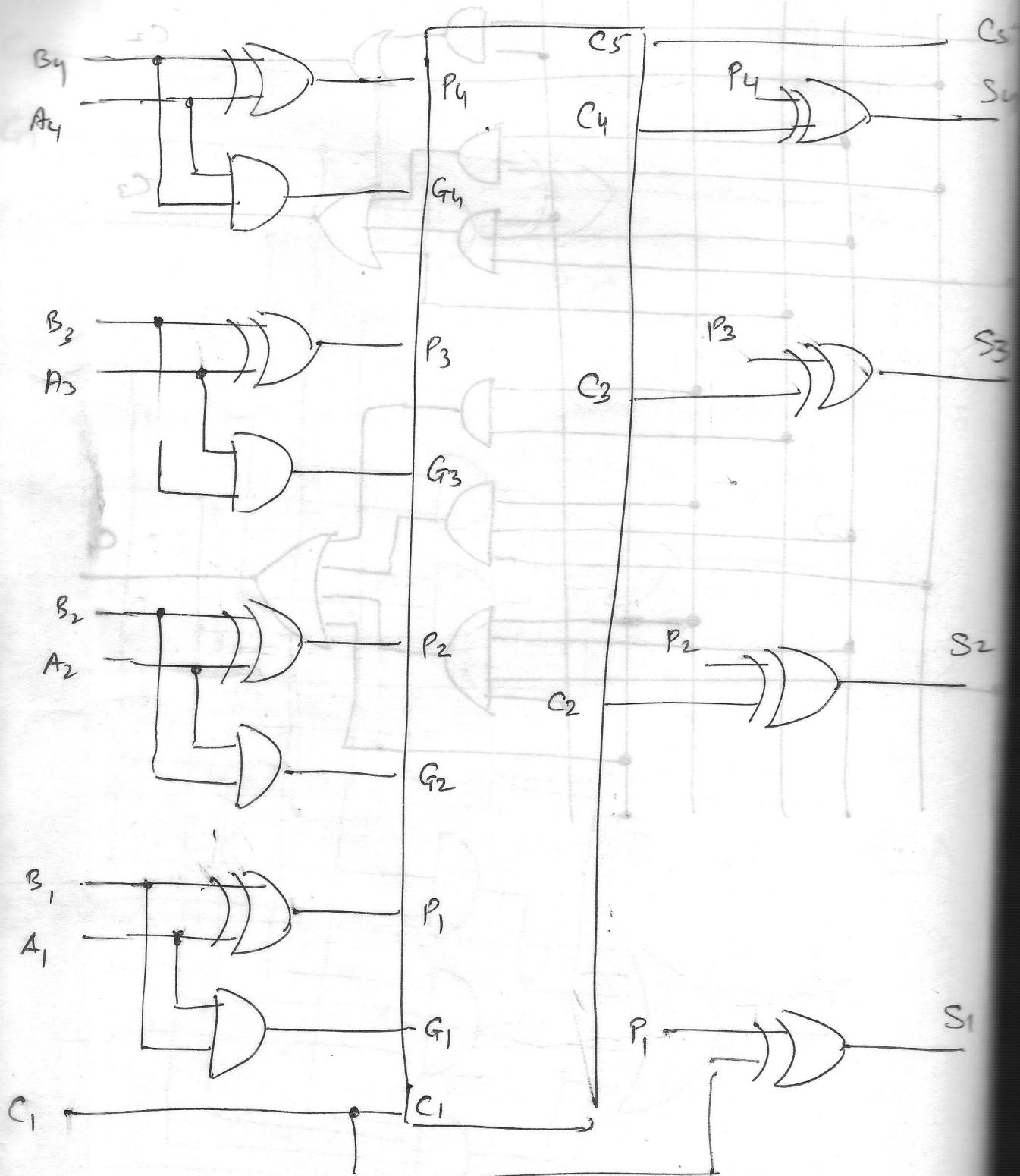
# 4-bit Parallel adder with look-ahead carry



redraw.



# 4 bit Parallel Adder with Look Ahead Carry



# BCD Adder

- BCD addition : if sum  $> 9$  we add 6.

Take two 4 bit nos.  $A = a_3 a_2 a_1 a_0$  and  $B = b_3 b_2 b_1 b_0$

$A = 0$  to 9,  $B = 0$  to 9. ( $\because$  BCD)

$A + B = 0$  to 18 (10010).

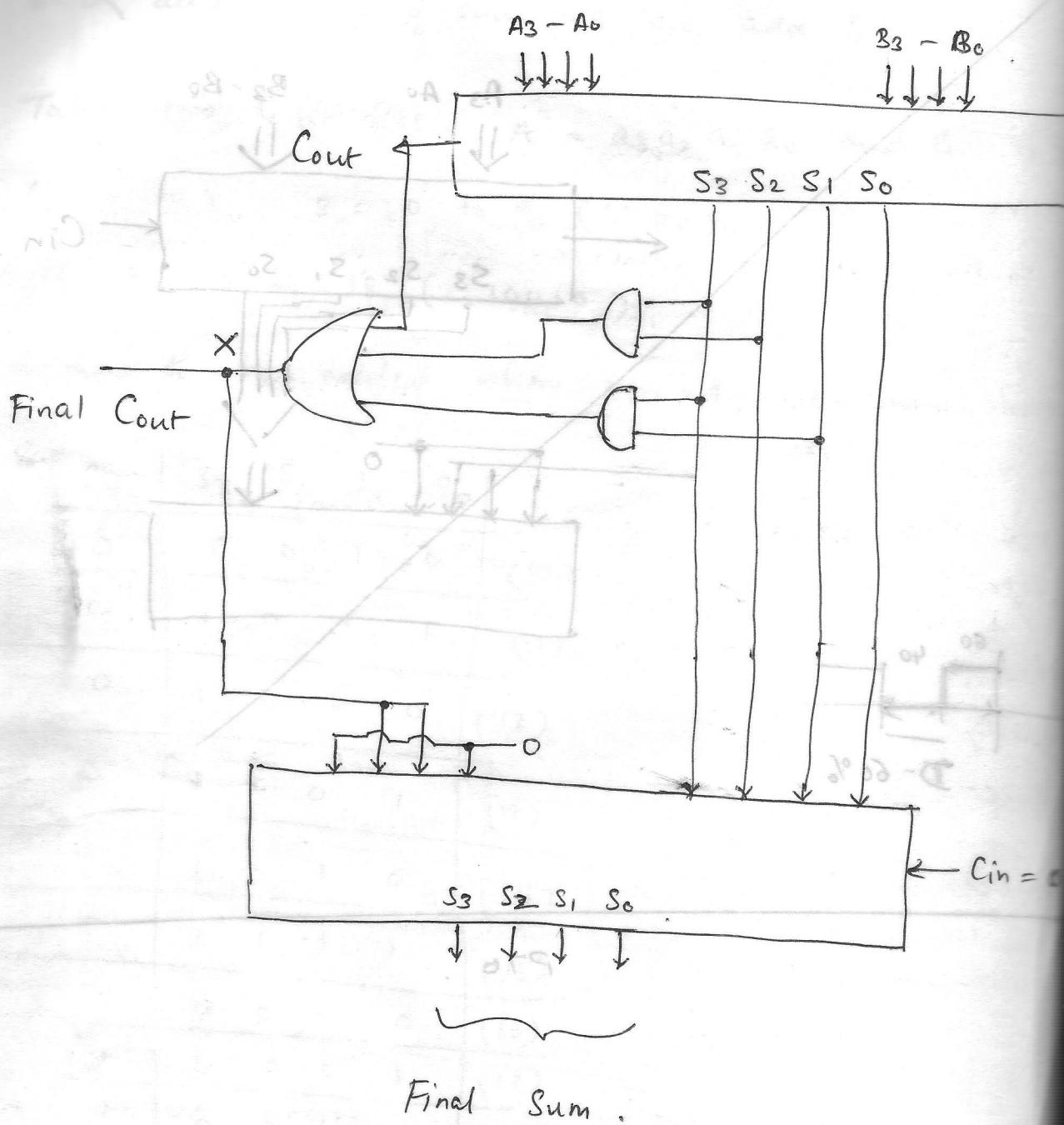
6 has to be added when sum  $> 9$ , i.e. sum = 10 -

$C_4$	$S_3$	$S_2$	$S_1$	$S_0$	
0	1	0	1	0	(10)
0	1	0	1	1	(11)
0	1	1	0	0	(12)
0	1	1	0	1	(13)
0	1	1	1	0	(14)
0	1	1	1	1	(15)
1	0	0	0	0	(16)
1	0	0	0	1	(17)
1	0	0	1	0	(18)

$$\therefore X = C_4 + S_3(S_2 + S_1)$$

See diagram on next page

# Block Diagram of BCD Adder.



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