

# 1 Results

**Definition 1. (Interference)** Given two sets  $U, W$ , we say  $U$   $k$ -interferes with  $W$  if

$$|U \cap W| \geq \frac{|W|}{k} \quad (1)$$

for some  $k \in (0, |W|]$

**Corollary 2.** If  $|U| = |W|$ , then  $U$   $k$ -interferes with  $W$  if and only if  $W$   $k$ -interferes with  $U$ .

We restrict the upper range of  $k$  to  $|W|$  for convenience, as beyond that all values of  $\frac{|W|}{k}$  will be less than 1.

**Definition 3. ((r,T,k)-Subgraph capacity)** Given a vertex set  $V = \{v_1, \dots, v_n\}$ , the  $(r, T, k)$ -subgraph capacity of  $V$  is the maximum number of subgraphs of size  $r$  subject to the constraint that for any randomly generated subgraph  $U$  of size  $r$ ,

$$\mathbb{E}[X] \leq T \quad (2)$$

where  $X$  is the number of interferences caused by adding  $U$  to the existing collection of subgraphs.

**Lemma 4.** Given a vertex set  $V$  with  $n$  vertices and two subsets  $U, W$  of size  $r$ , the probability that  $U$   $k$ -interferes with  $W$  is

$$\sum_{y=\lceil \frac{r}{k} \rceil}^r \frac{\binom{r}{y} \binom{n-r}{r-y}}{\binom{n}{r}}$$

*Proof.* if  $V = \{v_1, \dots, v_n\}$ , we can represent a subset  $U$  as a vector of length  $n$ ,  $u$  defined by

$$u_i = \begin{cases} 1 & \text{if } v_i \in U \\ 0 & \text{if } v_i \notin U \end{cases}$$

With this representation,  $U, W$  intersect at the indices where both vectors  $u, w$  have a 1. Let  $Y$  be a random variable denoting the number of indices where both  $u, w$  have a 1. Then

$$P(Y = y) = \frac{\binom{r}{y} \binom{n-r}{r-y}}{\binom{n}{r}} \quad (3)$$

Then the probability that  $U$   $k$ -interferes with  $W$  is

$$\sum_{y=\lceil \frac{r}{k} \rceil}^r \frac{\binom{r}{y} \binom{n-r}{r-y}}{\binom{n}{r}} \quad (4)$$

□

**Theorem 5.** Given a vertex set  $V$  with  $n$  vertices, the  $(r, T, k)$ -subgraph capacity of  $V$  is

$$\left\lfloor \frac{T}{\sum_{y=\lceil \frac{r}{k} \rceil}^r \frac{\binom{r}{y} \binom{n-r}{r-y}}{\binom{n}{r}}} + 1 \right\rfloor$$

*Proof.* Let  $M - 1$  be the number of subgraphs of size  $r$  currently generated. Add an arbitrary subgraph  $U$  of size  $r$ . From lemma 4., we know that the probability of  $U$   $k$ -interfering with another subgraph is  $\sum_{y=\lceil \frac{r}{k} \rceil}^r \frac{\binom{r}{y} \binom{n-r}{r-y}}{\binom{n}{r}}$ . This can be interpreted as the expected number of additional  $k$ -interferences  $U$  will cause with one subgraph. Since there are  $M - 1$  other subgraphs, the expected number of  $k$ -interferences by adding  $U$  is

$$\sum_{y=\lceil \frac{r}{k} \rceil}^r \frac{\binom{r}{y} \binom{n-r}{r-y}}{\binom{n}{r}} (M - 1) \quad (5)$$

From equation (2), we have

$$\sum_{y=\lceil \frac{r}{k} \rceil}^r \frac{\binom{r}{y} \binom{n-r}{r-y}}{\binom{n}{r}} (M - 1) \leq T \implies M \leq \frac{T}{\sum_{y=\lceil \frac{r}{k} \rceil}^r \frac{\binom{r}{y} \binom{n-r}{r-y}}{\binom{n}{r}}} + 1 \quad (6)$$

The  $(r, T, k)$ -subgraph capacity of  $V$  is the largest integer  $M$  that satisfies equation (6). □

*Alternate proof.* Suppose we have  $M$  subgraphs in the collection. Pick two subgraphs  $U, W$ . From lemma 4., we know that the probability of  $U$   $k$ -interfering with  $W$  is  $\sum_{y=\lceil \frac{r}{k} \rceil}^r \frac{\binom{r}{y} \binom{n-r}{r-y}}{\binom{n}{r}}$ . Since all subgraphs have the same size, this is the probability that  $U, W$  pair will cause 2  $k$ -interferences. So the expected number of interferences caused by one pair is

$$2 \sum_{y=\lceil \frac{r}{k} \rceil}^r \frac{\binom{r}{y} \binom{n-r}{r-y}}{\binom{n}{r}}$$

We know that there are  $\binom{M}{2} = M(M-1)/2$  such pairings so the expected number of total interferences are

$$2 \frac{M(M-1)}{2} \sum_{y=\lceil \frac{r}{k} \rceil}^r \frac{\binom{r}{y} \binom{n-r}{r-y}}{\binom{n}{r}} = M(M-1) \sum_{y=\lceil \frac{r}{k} \rceil}^r \frac{\binom{r}{y} \binom{n-r}{r-y}}{\binom{n}{r}}$$

Since there are  $M$  subgraphs, the expected number of interferences due to one subgraph is

$$\frac{M(M-1)}{M} \sum_{y=\lceil \frac{r}{k} \rceil}^r \frac{\binom{r}{y} \binom{n-r}{r-y}}{\binom{n}{r}} = (M-1) \sum_{y=\lceil \frac{r}{k} \rceil}^r \frac{\binom{r}{y} \binom{n-r}{r-y}}{\binom{n}{r}}$$

From equation (2), we have

$$(M-1) \sum_{y=\lceil \frac{r}{k} \rceil}^r \frac{\binom{r}{y} \binom{n-r}{r-y}}{\binom{n}{r}} \leq T \implies M \leq \frac{T}{\sum_{y=\lceil \frac{r}{k} \rceil}^r \frac{\binom{r}{y} \binom{n-r}{r-y}}{\binom{n}{r}}} + 1 \quad (7)$$

The  $(r, T, k)$ -subgraph capacity of  $V$  is the largest integer  $M$  that satisfies equation (6). □