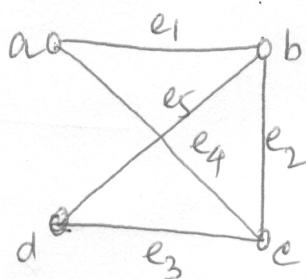


# GRAPH THEORY

①

GRAPH → A mathematical model

Definition → A Graph  $G = (V, E)$  consists of a non-empty set  $V$  of vertices or nodes and a set of edges. Each edge has either one or two vertices associated with it called endpoints. An edge is said to connect its endpoints.



$$V = \{a, b, c, d\}$$

$$|V| = 4$$

$$E = \{e_1, e_2, e_3, e_4, e_5\}$$

$$|E| = 5$$

Fig 1.

\* The connections made by edges matters, not the geometry (length of the edges, whether edges cross, how vertices are depicted and so on do not matter)

Infinite Graph: A Graph with infinite vertex set

Finite Graph: A Graph with finite vertex set

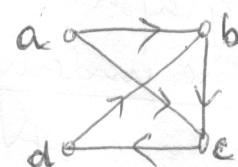
Notations:

\*  $V(G)$  — Set of Vertices (Domain)      || Defining  $G$  in terms of logic notation  
 Represents

$$E(G) — Relation : E(G) \subseteq V(G) \times V(G)$$

Type of Graph:

① Directed Graph (Di-Graph)  $G = (V, E)$  with  $V$  a set of vertices ( $v$ ) and set  $E$  of directed edges. Each edge is an ordered pair of vertices  $(u, v)$  where  $u, v \in V$



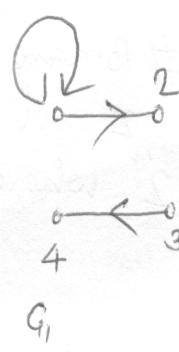
$$V = \{a, b, c, d\}$$

$$E = \{(a, b), (b, c), (a, c), (d, b), (c, d)\}$$

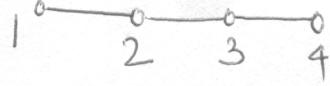
$E(G) = V(G) \times V(G)$   
 represents relation

$$\text{Given, } V(G) = \{1, 2, 3, 4\}$$

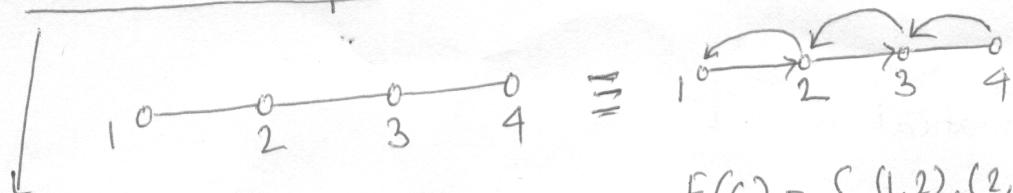
$$E(G) = \{(1, 2), (1, 3), (3, 4)\}$$



if  $E(G) = \{(1, 2), (2, 3), (3, 4)\}$



(2) Undirected Graph  $\rightarrow$  Does not have directed edges



A Special case of  
directed graph

$$E(G) = \{(1,2), (2,1), (2,3), (3,2), (3,4), (4,3)\}$$

Set notation

$$\equiv \{\{1,2\}, \{2,3\}, \{3,4\}\}$$

Note:

For a Directed Graph

$$E(G) = V(G) \times V(G)$$

A binary relation = Directed Graph

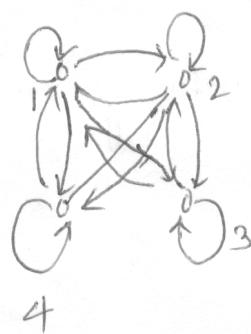
Case 1: If  $V(G) = \{1, 2, 3, 4\}$  || empty/null graph

$$E(G) = \emptyset$$

$$\begin{matrix} 0 & 0 & 0 & 0 \\ 1 & 2 & 3 & 4 \end{matrix}$$

Case 2: If  $V(G) = \{1, 2, 3, 4\}$

$$E(G) = V(G) \times V(G) = \{(1,1), (2,1), (3,1), (4,1), (1,2), (2,2), (3,2), (4,2), (1,3), (2,3), (3,3), (4,3), (1,4), (2,4), (3,4), (4,4)\}$$



If  $V(G) = \{1, 2, \dots, n\}$

Q: How many different directed graphs are possible?

Ans: Any binary set<sup>n</sup> corresponds to a directed graph.

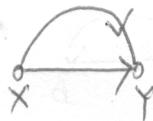
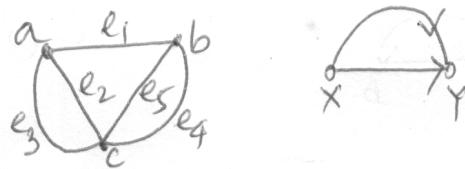
Similarly, given a  $\neq$  DAC, one can derive the underlying relation

$$\# \text{Directed graph} = \# \text{Binary relations}$$
$$= 2^{n^2} \text{ where } n \text{ is no. of vertices}$$

Given, a set of  $n$  items,  
~~no~~ Size of cross product  
i.e. if  $|V| = n$  then  $|V \times V| = n^2$   
 $\therefore \# \text{Binary relations}$  $= 2^{n^2}$

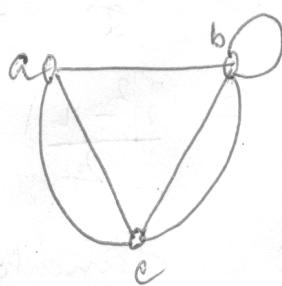
## Terminologies

- (i) Simple graphs → each edge connects two different vertices and no two edges connect the same pair of vertices
- (ii) Multigraphs → graphs that may have multiple edges connecting the same two vertices

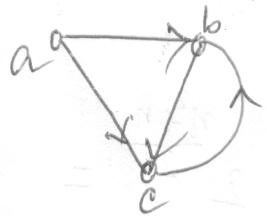


when 'm' different edges connect vertices  $u$  &  $v$  then we say  $\{u, v\}$  is an edge of multiplicity  $m$   
 For ex: the two graphs above,  $\{a, c\}$  &  $\{b, c\}$  have multiplicity 2

- (iii) Self loop → An edge that connects a vertex to itself is called self loop
- (iv) Pseudograph → may include loops, as well as multiple edges connecting the same pair of vertices.



\* Simple directed graph has no loops & no multiple edges.



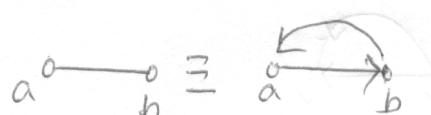
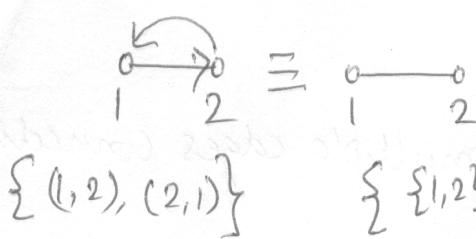
$$\begin{aligned} \# \text{Directed Simple graphs} &= \text{no. of irreflexive binary relations of } n \text{ vertices} \\ &= \frac{n^2-n}{2} \end{aligned}$$

$$\begin{aligned} \# \text{Undirected Simple graphs} &= \text{no. of irreflexive binary relations of } n \text{ vertices} \\ &= \frac{n^2-n}{2} \end{aligned}$$

## Undirected Graphs

$$V(G) = \{1, 2, \dots, n\}$$

$$E(G) \subseteq V(G) \times V(G)$$



Q: How many possible undirected graphs are possible on  $n$  nodes/vertices

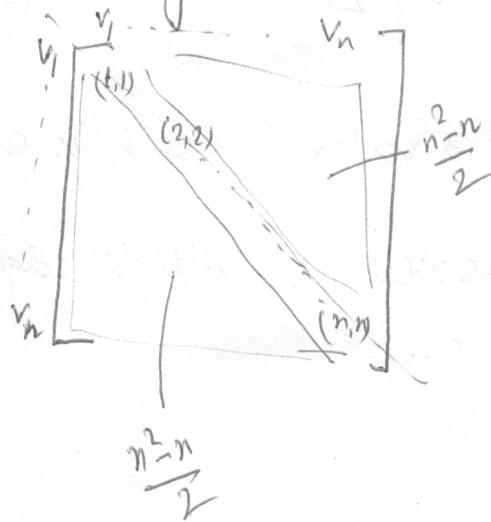
Ans: Any undirected graph corresponds to symmetric binary relation

$$\begin{matrix} & 1 & 2 & 3 \\ 1 & 0 & 1 & 0 \\ 2 & 1 & 0 & 1 \\ 3 & 0 & 1 & 0 \end{matrix} \quad \left\{ (1, 2), (2, 1) \right\} \quad \left\{ (2, 3), (3, 2) \right\} = \left\{ \{1, 2\}, \{2, 3\} \right\}$$

The underlying relation is always symmetric

# undirected graphs = # symmetric binary rel's on  $n$  nodes

To Count # Symmetric rel's on  $n$  nodes using matrix representation



The diagonal elements =  $\{(1,1), (2,2), \dots, (n,n)\}$

$$|D| = n$$

# elements in the lower triangle

$$= \frac{n^2 - n}{2} \quad \text{Total - Diagonals}$$

$$= \frac{n^2 - n}{2}$$

Lower triangle is symmetric to upper triangle  $\therefore$  elements  $(v_i, v_j)$  &  $(v_j, v_i)$  are included

$$\therefore \# \text{Binary rel}^n = \frac{\# \text{elements}}{2}$$

$$\# \text{elements in any one row} \times 2$$

$$\# \text{elements in diag}$$

$$= 2^{\frac{n^2 - n}{2}} \times 2^n = 2^{\frac{n^2 + n}{2}} =$$

all possible  
subsets of  
 ~~$(v_i, v_j)$~~  & its  
counterpart

## Undirected Graphs

Neighbourhood of a vertex  $v \in V(G)$

$$N_G(v) = \{u \mid \{v, u\} \in E(G)\}$$

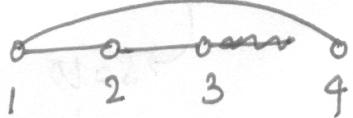
$$V(G) = \{1, 2, 3, 4\}$$

$$N_G(1) = \{2, 4\}$$

$$E(G) = \{\{1, 2\}, \{2, 3\}, \{1, 4\}\}$$

$$N_G(2) = \{1, 3\}$$

$$N_G(3) = \{2, 4\}$$



Degree of a vertex  $v \in V(G)$

$$d_G(v) = |N_G(v)|$$

$$d_G(1) = 2$$

$$d_G(3) = 1 \quad d_G(2) = 2 \quad d_G(4) = 1$$

Degree sequence of a graph G

$$D(G) = (d_{v_1}(G), d_{v_2}(G), \dots, d_{v_n}(G))$$

$$D(G) = (2, 1, 2, 1)$$

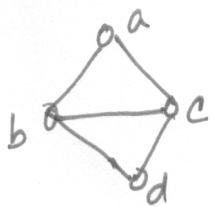
If a graph G is given,  $\Rightarrow D(G)$  exists

Q: If a sequence is given, are these sequences graphic?

$\hookrightarrow$  given an arbitrary vector, does there exist a graph whose degree sequence correspond to the vector?

Q: Is  $(1, 1, 3, 4, 2)$  a graphic sequence  $\Rightarrow$  trying to construct a graph on 5 vertices s.t.  $D(G) = \{1, 1, 3, 4, 2\}$

Observation:



$$\forall v \in V(G) \quad n = |V(G)|$$

$$0 \leq d_G(v) \leq n-1$$

degree of any vertex

ex: b, c  
already to all

Is  $(4, 3, 2, 1)$  graphic?

$\hookrightarrow$  graph on 4 vertices

$$D(G) = \{4, 3, 2, 1\}$$

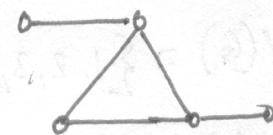
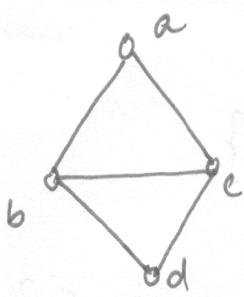
$D_G(v) = 4$  but any node in a graph having 4 vertices can at max have degree = 3

$\cancel{\text{No graph exists corresponding to it}}$

$(4, 3, 2, 1)$  is not graphic

R: Is  $(3, 3, 3, 3, 3)$  is graphic?

↪ not graphic - why?



$$D(G) = (1, 3, 2, 3, 1)$$

$$(2, 3, 3, 2), \text{ Degree sum } S = \sum_{v_i \in V(G)} d_G(v_i)$$

$$S = 2+3+3+2=10$$

$$\hookrightarrow S=10$$

$$\underline{2 \times |E(G)|}$$



Observation: For any  $G$ ,  $S = \sum d_G(v_i)$  is even

every edge contributes 1 to each of the vertices it connects

∴ in a graph, degree of a vertex is contributed by each edges it connects

each edge is counted once for left hand pt. & once for rt hand point

$$D(G) = (2, 3, 3, 2)$$
$$S = \sum d(v_i) = 1+1+1+1 = 4$$
$$\sum d_G(v_i) = 10 = 2 \times 5 = 2|E(G)|$$

⇒ sum of is counted once w.r.t  $d_G(u)$

again " "  $d_G(v)$

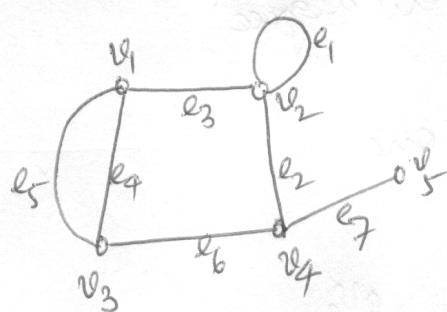
$$\therefore \sum d_G(v_i) = 2 \times |E(G)|$$

Sum of degrees is always even.

(3)

## Incidence & Degree

When a vertex  $v_i$  is an end vertex of some edge  $e_j$ ,  $v_i$  &  $e_j$  are said to be incident (on or to) each other



$e_2, e_6$  &  $e_7$  are incident with  $v_4$

Two non parallel edges are adjacent if they are incident on a common vertex. For ex:  $e_2$  &  $e_7$  are adjacent

Two vertices are adjacent if they are the end vertices of same edge  
 $v_4$  &  $v_5$  are adjacent while  $v_1$  &  $v_4$  are not

# The no. of edges incident on a vertex  $v_i$ , with self-loop counted twice is called degree,  $d(v_i)$  of vertex  $v_i$

$$d(v_1) = d(v_3) = d(v_4) = 3 \quad d(v_2) = 4 \quad d(v_5) = 1$$

Degree is also called Valency

degree sequence: (3, 3, 3, 4, 1)

# Let us consider a Graph  $G$  with  $e$  edges and  $n$  vertices

$v_1, v_2, \dots, v_n$

each edge contributes two degrees, the sum of degrees of all vertices in  $G$  is twice the no. of edges

i.e.

$$\sum_{i=1}^n d(v_i) = 2e$$

*degree sequence sum* → *the sum of degree sequence of vertices in a graph are even and equal to twice no. of edges*

Ex:  $d(v_1) + d(v_2) + d(v_3) + d(v_4) + d(v_5) = 3+4+3+3+1 = 14 = 2 \times \text{edges}$

Theorem: The no. of vertices of odd degree in a graph is even

Proof: Consider vertices with odd & even degree separately.

\* Given a sequence (3, 3, 3, 3, 5)  
Is it graphic?  
Solv. No., ∵ the sum of sequence  $3+3+3+3+5=17$  is not even

$$\sum_{i=1}^n d(v_i) = \sum_{\text{even}} d(v_j) + \sum_{\text{odd}} d(v_k)$$

↓  
2x#edges Even

∴ L.H.S is even & 1st term of R.H.S is even  
 $\therefore \sum_{\text{odd}} d(v_k)$  is also even

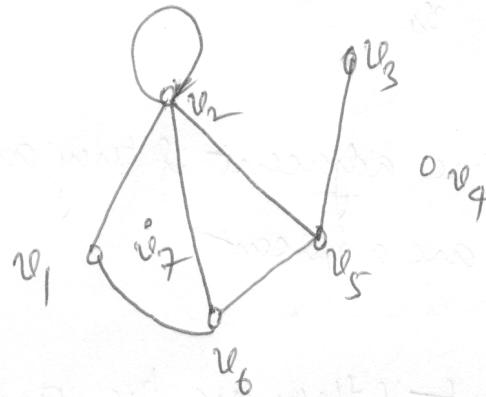
### Isolated Vertex, Pendant Vertex & Null Graph

Isolated Vertex  $\rightarrow$  Vertex with no incident edge  
i.e. Vertices with zero degree

Pendant Vertex  $\rightarrow$  Vertex with degree one

$v_4, v_7 \rightarrow$  isolated vertices

$v_3$  - pendant vertex



Two adjacent edges are said to be in series if their common vertex is of degree 2. Ex: Two edges incident on  $v_7$  are in series

Null Graph  $\rightarrow$  Graph without edges

$$\begin{matrix} 0 & 2 & 3 & 4 \\ & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 5 & 6 & 7 & 8 \end{matrix} \quad V = \{1, 2, 3, 4, 5, 6, 7, 8\} \quad E = \emptyset \quad |V| = 8$$

If  $V = \emptyset \rightarrow$  then there is no graph