Surprises & Paradoxes II

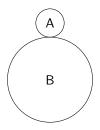
Coin Rotation Paradox

2 Braess Paradox

3 The Social Cost of Traffic at Equilibrium

Coin Rotation Paradox

The 1982 SAT Question Everyone Got Wrong



The radius of circle A is $\frac{1}{3}$ of the radius of circle B. Circle A rolls around circle B one trip back to its starting point. How many times will circle A revolve in total?

- (a) $\frac{3}{2}$ (b) 3 (c) 6 (d) $\frac{9}{2}$ (e) 9

Animation of Coin Rotation Paradox: Python Code I

```
import math
import matplotlib.pyplot as plt
from matplotlib.patches import Circle, Wedge
from matplotlib.transforms import Affine2D
from matplotlib.backends.backend_pdf import PdfPages
import os; os.chdir(os.path.dirname(__file__))
# Set the radii
R = 1.5 # Radius of the outer (stationary) coin
r = 0.5 # Radius of the inner (rolling) coin
def create_logo(ax, x, y, radius, angle):
    # Create a simple smiley face logo
    face = Circle((0, 0), radius * 0.7, fill=True, color="cyan")
    left eye = Circle((-radius * 0.3, radius * 0.2), radius *
    → 0.1, fill=True, color="black")
    right eye = Circle((radius * 0.3, radius * 0.2), radius *
    → 0.1, fill=True, color="black")
```

Animation of Coin Rotation Paradox: Python Code II

```
mouth = Wedge((0, 0), radius * 0.5, 225, 315, width=radius *
    → 0.1, color="black")
    # Create a transformation
   t = Affine2D().rotate(-angle).translate(x, y) + ax.transData
    # Apply the transformation to all elements
   face.set_transform(t)
    left eye.set transform(t)
   right eye.set transform(t)
   mouth.set transform(t)
    ax.add artist(face)
    ax.add artist(left eye)
    ax.add_artist(right_eye)
    ax.add_artist(mouth)
def create_frame(angle):
```

Animation of Coin Rotation Paradox: Python Code III

```
# Create figure and axis
fig, ax = plt.subplots(figsize=(8, 8))
# Set limits and aspect ratio
ax.set xlim(-(R + r + 1), R + r + 1)
ax.set ylim(-(R + r + 1), R + r + 1)
ax.set aspect("equal")
# Remove axis ticks
ax.set xticks([])
ax.set yticks([])
# Calculate position of the rolling coin's center
x = (R + r) * math.cos(angle)
y = (R + r) * math.sin(angle)
# Calculate position of the point on the edge of the moving
```

Animation of Coin Rotation Paradox: Python Code IV

```
point_x = x + r * math.cos((R + r) / r * angle)
point_y = y + r * math.sin((R + r) / r * angle)
# Draw stationary coin
stationary coin = Circle((0, 0), R, fill=False, color="blue")
ax.add artist(stationary coin)
# Draw rotating coin
rotating_coin = Circle((x, y), r, fill=False, color="red")
ax.add artist(rotating coin)
# Add logo to the rotating coin
create_logo(ax, x, y, r, (R + r) / r * angle)
# Draw point on the edge of the moving coin
ax.plot(point_x, point_y, "ro")
# Draw cardioid path
```

Animation of Coin Rotation Paradox: Python Code V

```
theta = [i * 2 * math.pi / 100 for i in range(int(100 * angle
    \rightarrow / (2 * math.pi)) + 1)]
    cardioid_x = [(R + r) * math.cos(t) + r * math.cos((R + r) / r)]

→ r * t) for t in thetal

    cardioid y = [(R + r) * math.sin(t) + r * math.sin((R + r))]
    \rightarrow r * t) for t in thetal
    ax.plot(cardioid x, cardioid y, "g-")
    # Add title
    ax.set title(f"Rotation: {angle/(2*math.pi)*360:.1f}")
    return fig
# Generate 512 frames and save them into a single PDF
num_frames = 512
pdf_filename = 'coin_rotation_paradox.pdf'
tex = \Pi
```

Animation of Coin Rotation Paradox: Python Code VI

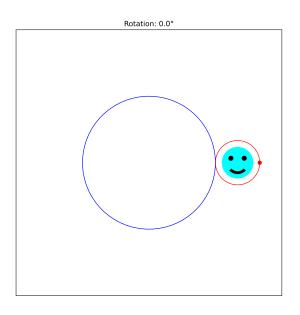
```
with PdfPages(pdf filename) as pdf:
    for i in range(num frames + 1):
        angle = i * 2 * math.pi / num frames
        fig = create frame(angle)
        pdf.savefig(fig, bbox inches='tight', pad inches=0.1,
        \rightarrow dpi=300)
        plt.close(fig)
        # Generate LaTeX commands referencing pages (1-based
        \rightarrow indexing)
        tex.append(f'\\only<{i +
        → 1}>{{\\includegraphics[width=.6\\textwidth, page={i}
        + 1}]{{fig/note02/coin_rotation_paradox.pdf}}\\noind_

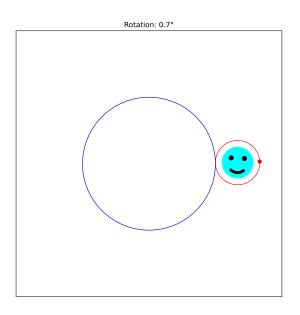
    ent}}')

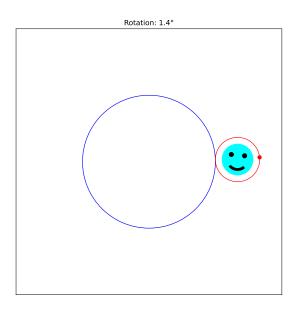
# Write the LaTeX overlayarea content to a .tex file
open('../../coin.tex',
→ 'w').write(f"\\hspace*{{\\dimexpr(\\textwidth - .6\\textwidt}
\rightarrow h)/2\relax}\n\begin{{overlayarea}}{{\\textwidth}}{{0.7\\t_|}}

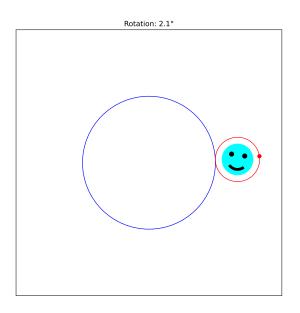
    extheight}\n{'\n'.join(tex)}\n\\end{{overlayarea}}")

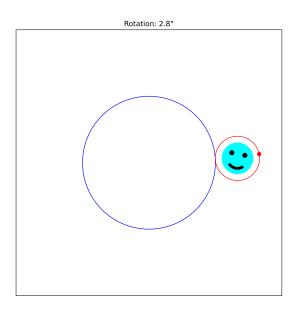
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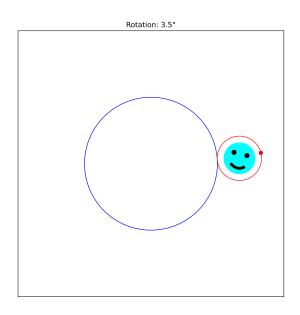


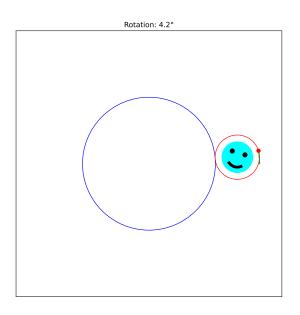


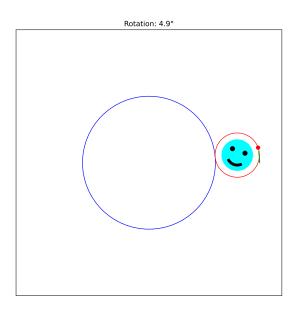


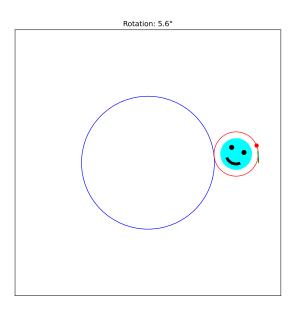


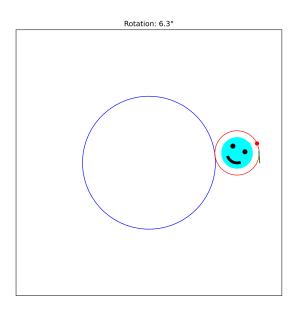


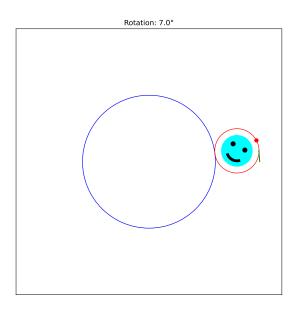


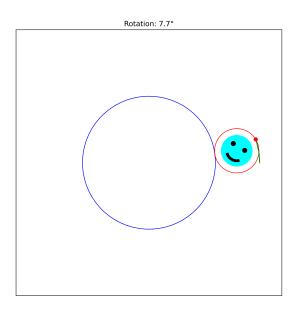


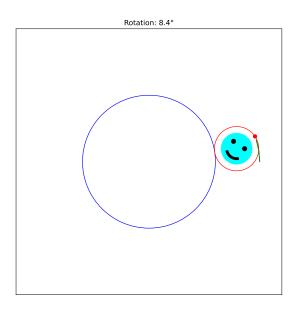


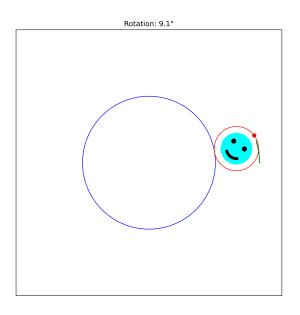


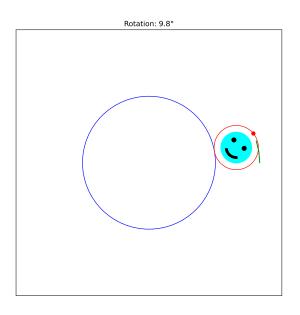


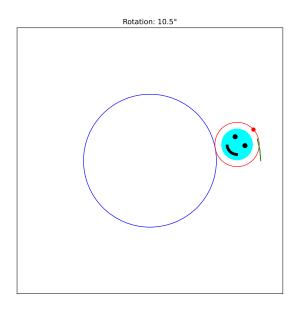


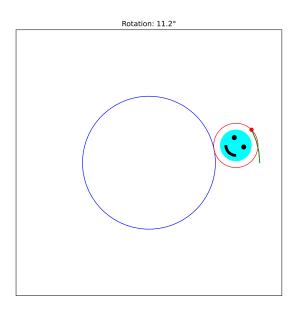


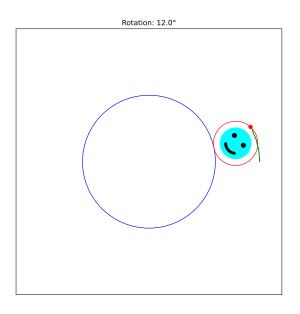


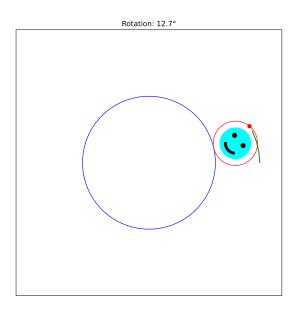


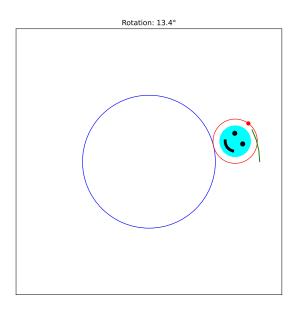


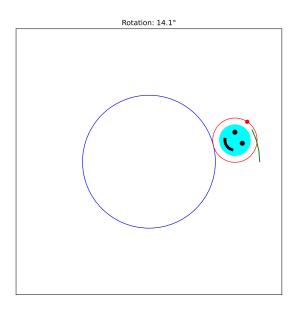


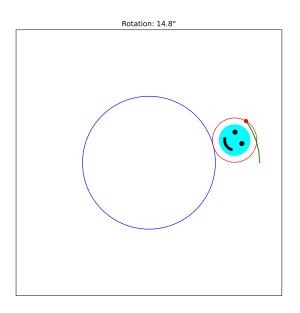


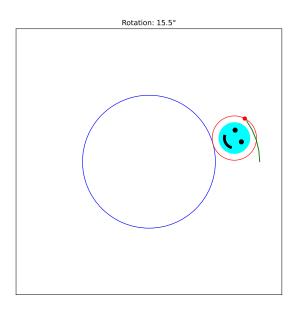


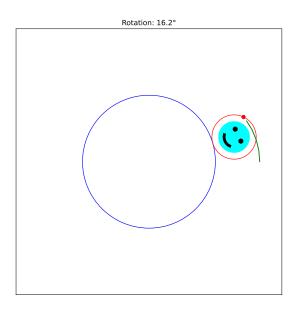


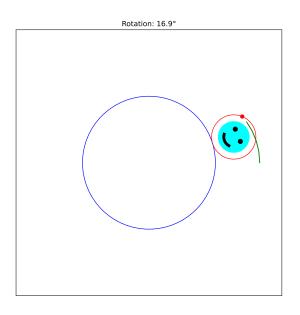


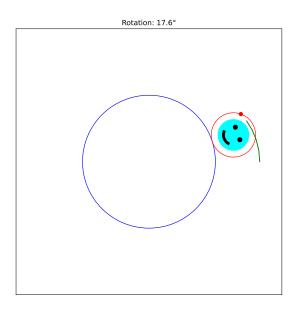


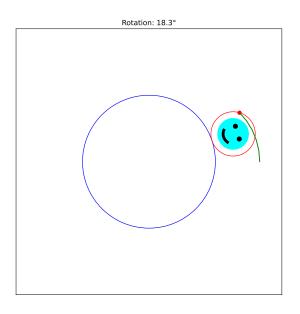


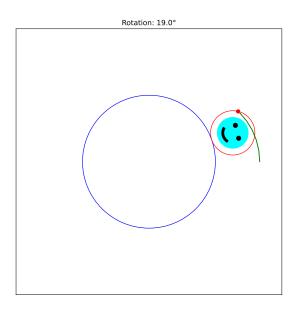


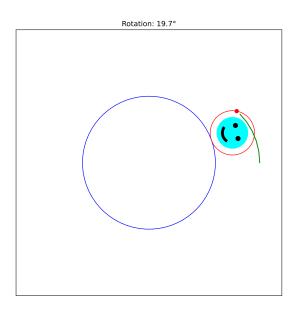


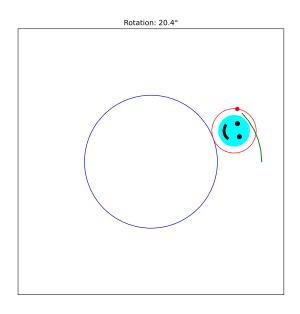


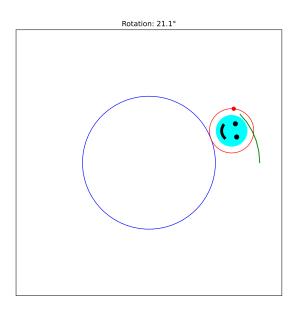


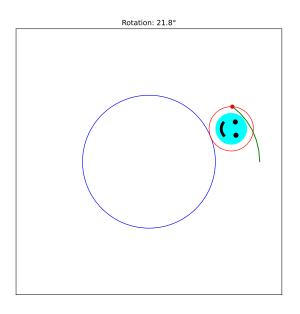


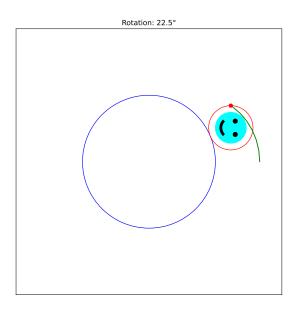


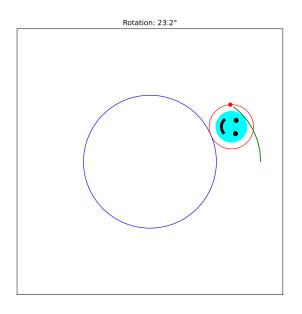


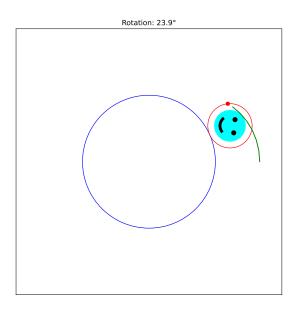


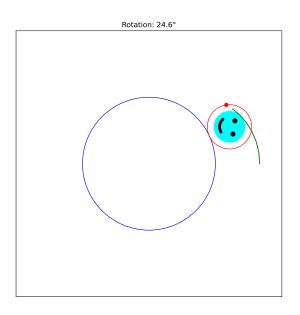


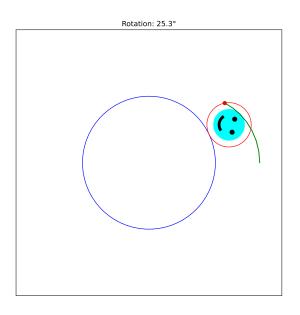


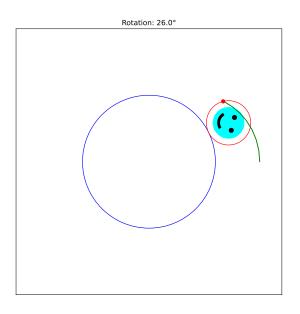


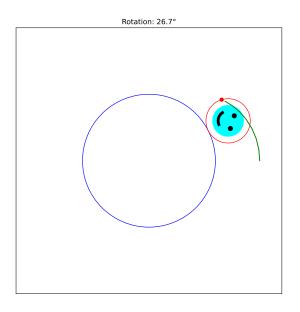


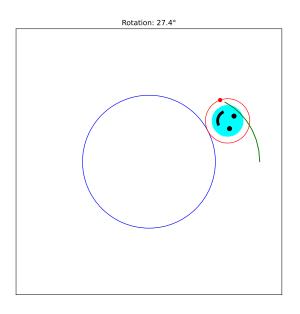


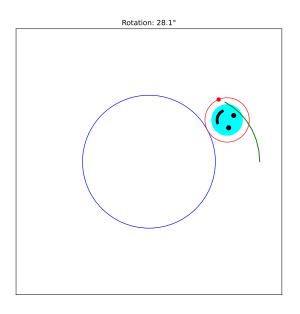


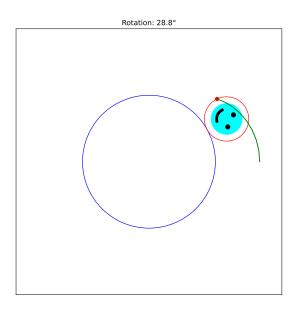


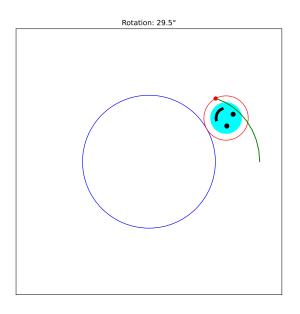


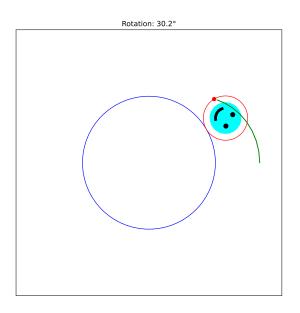


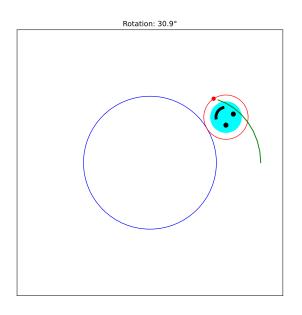


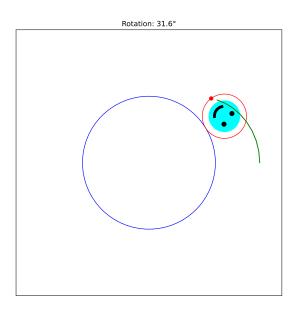


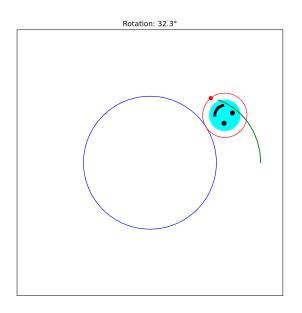


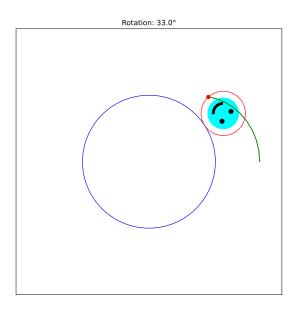


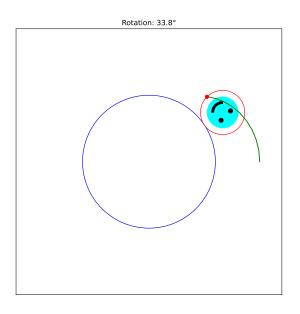


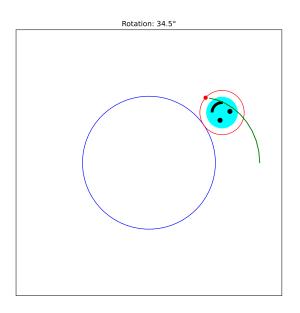


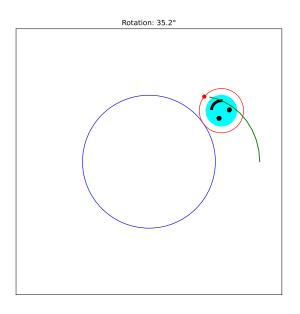


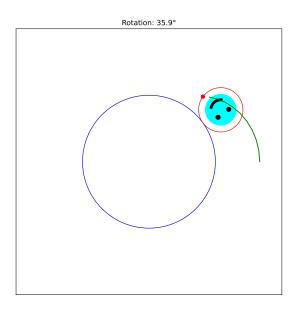


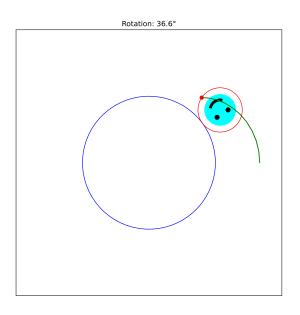


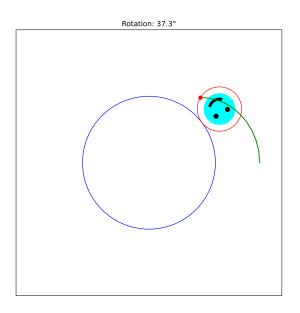


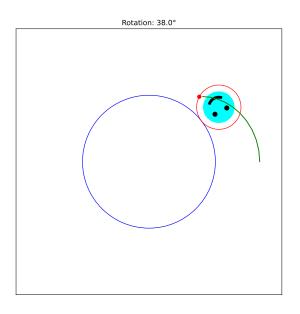


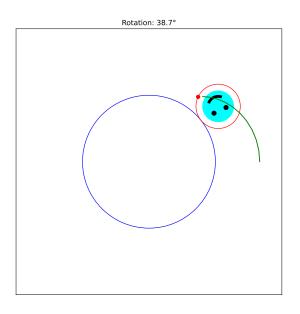


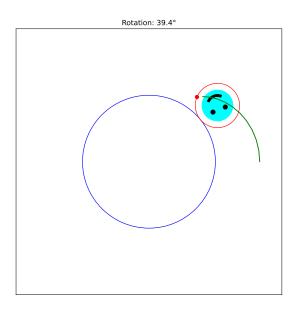


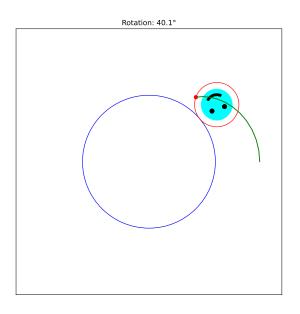


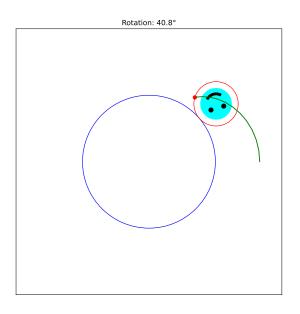


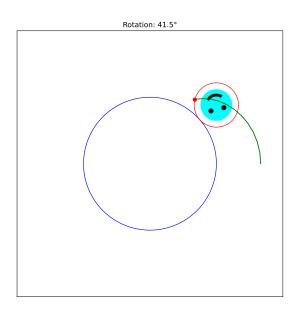


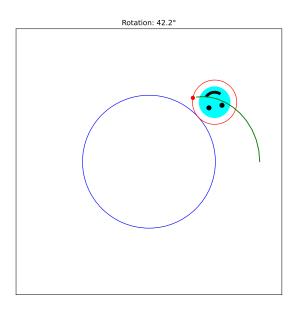


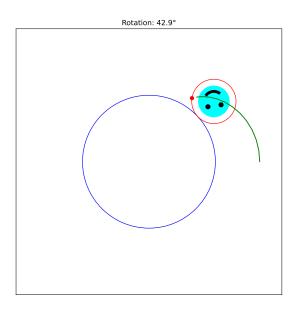


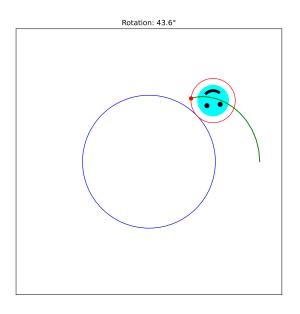


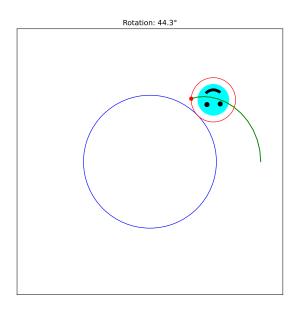


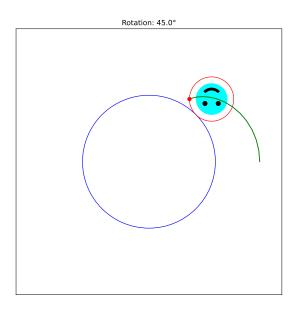


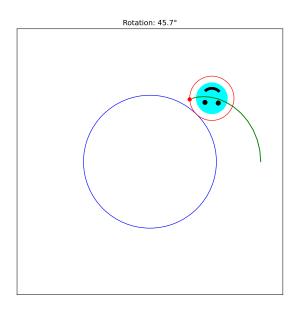


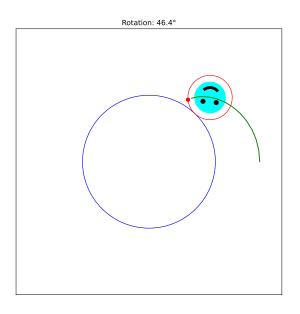


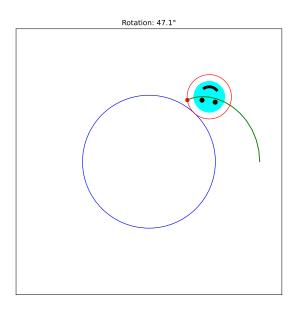


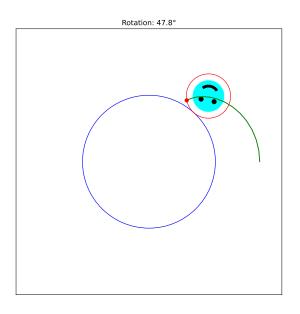


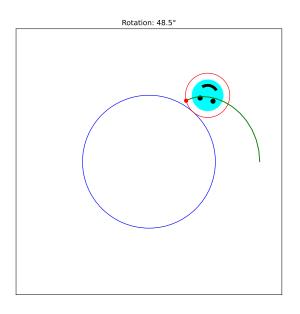


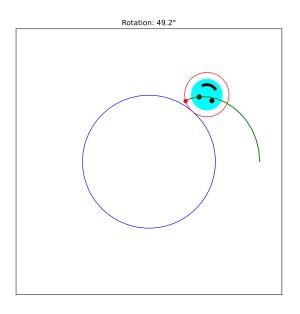


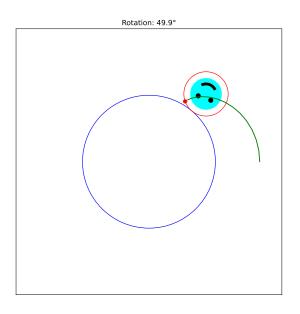


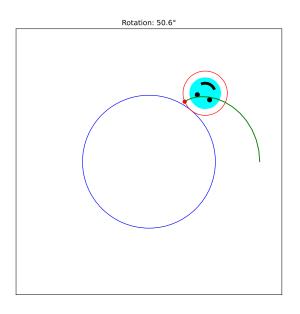


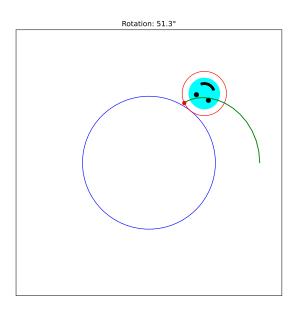


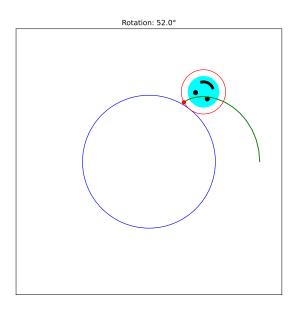


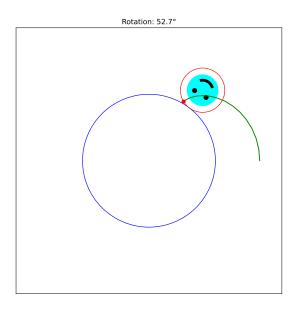


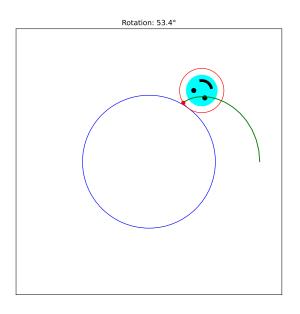


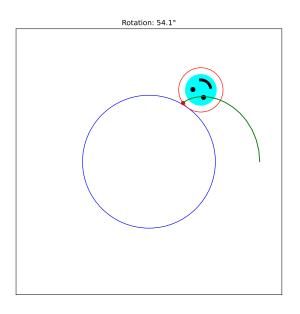


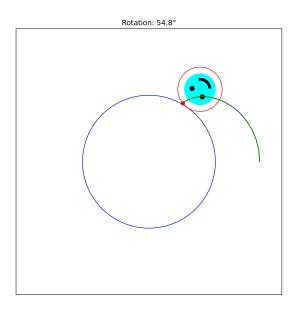


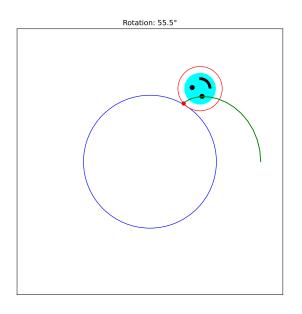


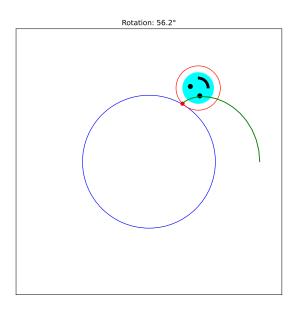


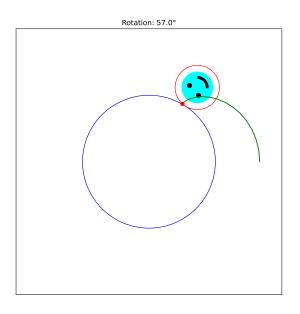


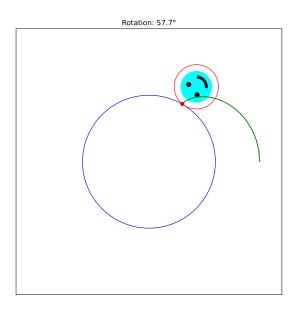


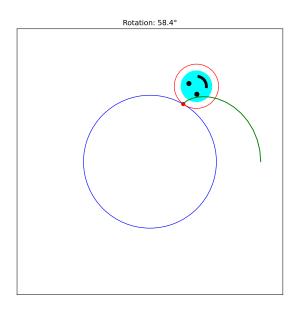


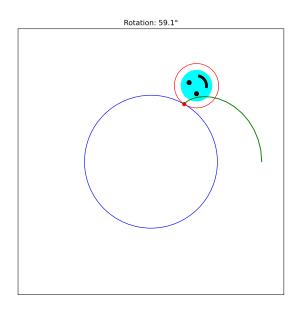


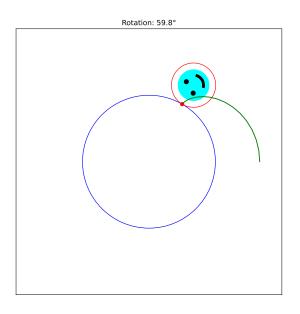


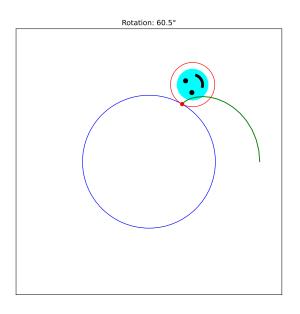


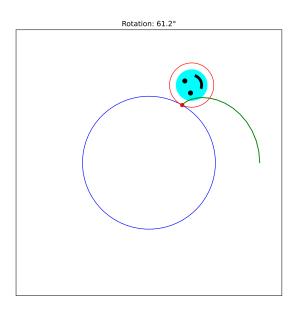


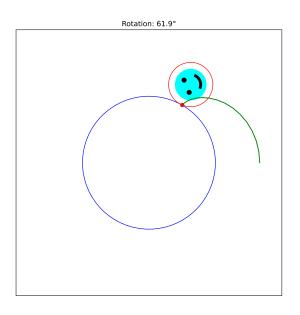


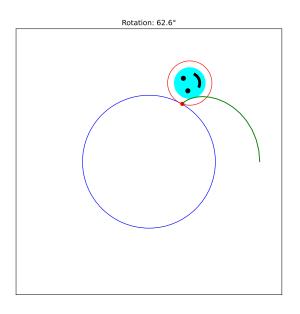


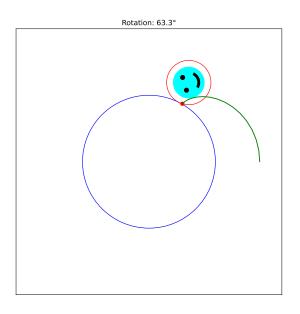


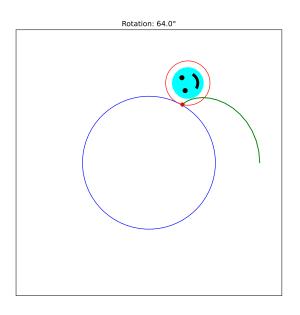


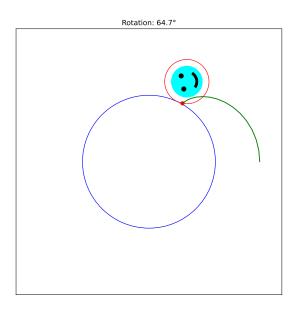


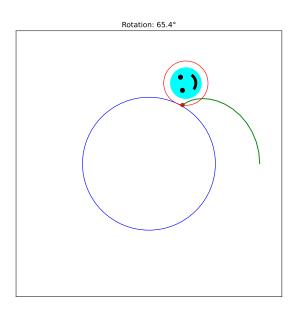


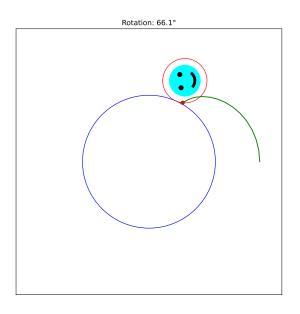


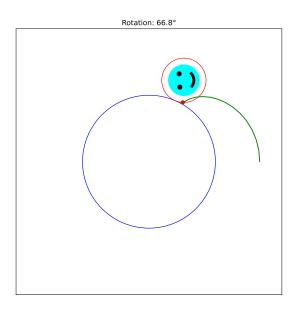


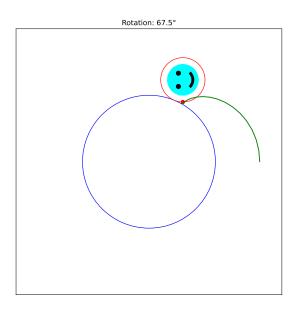


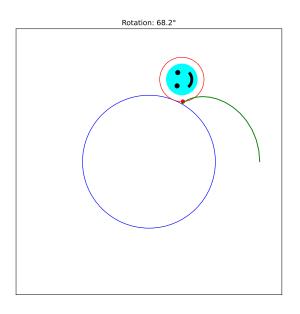


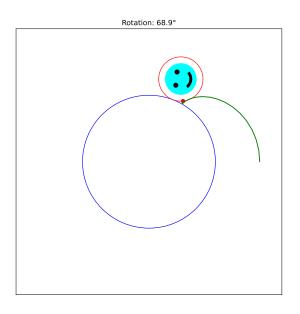


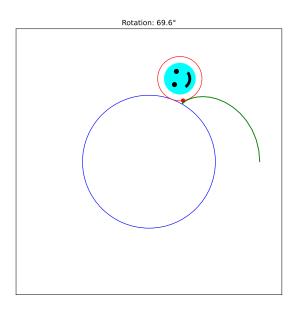


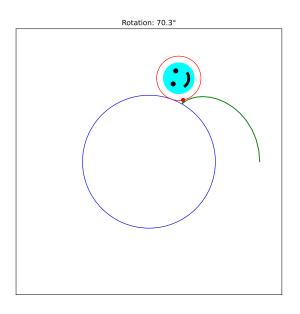


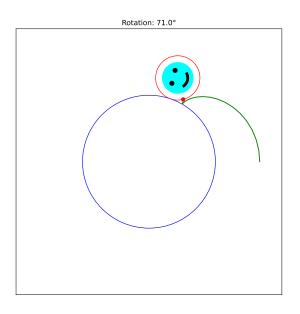


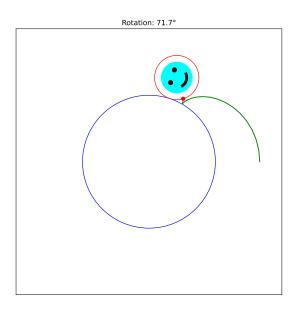


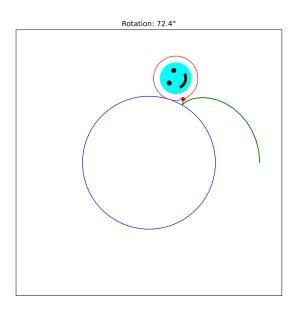


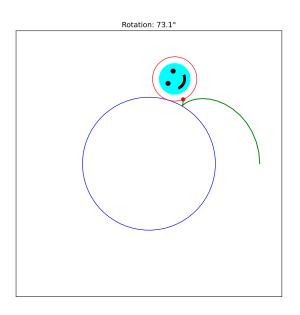


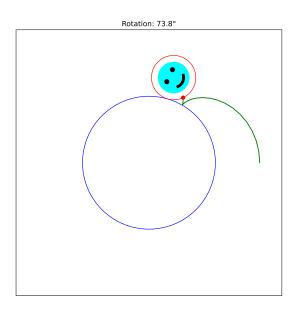


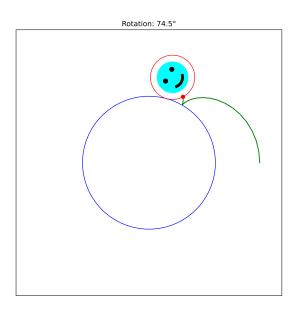


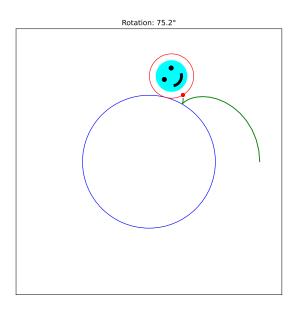


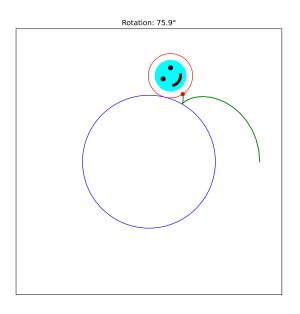


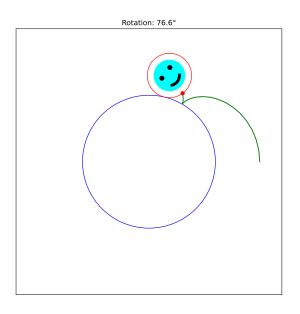


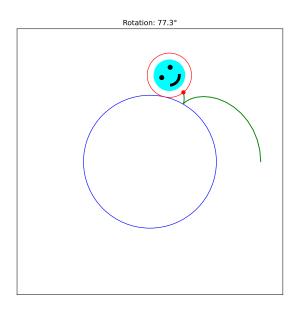


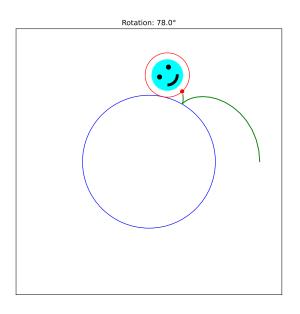


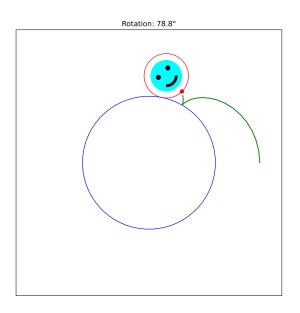


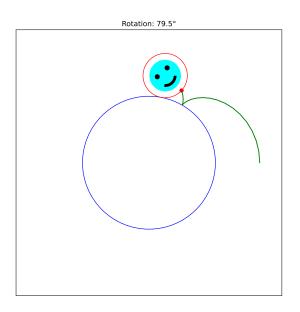


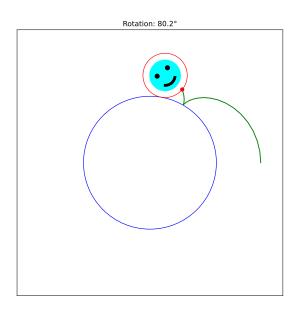


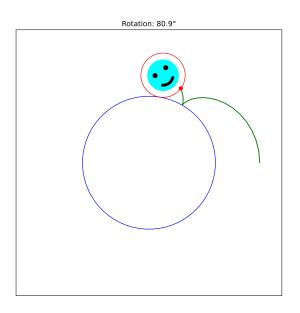


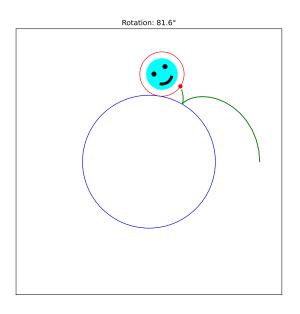


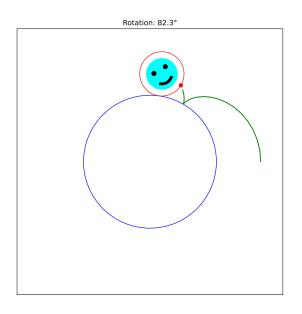


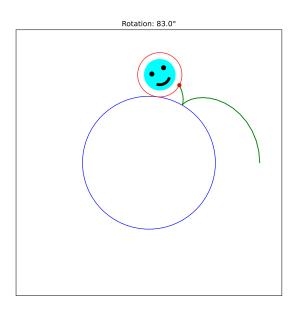


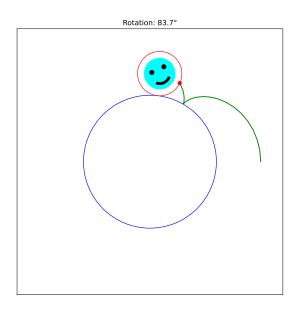


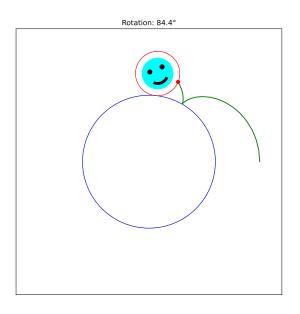


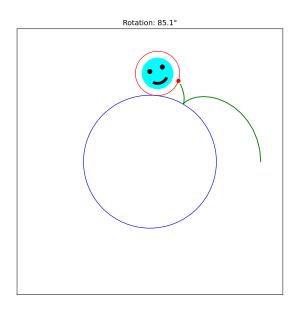


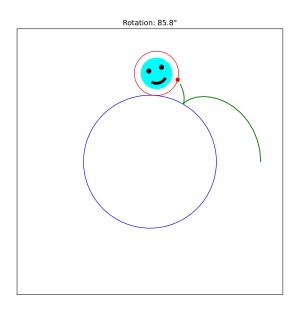


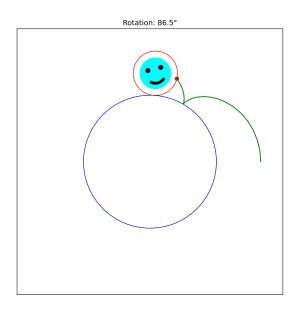


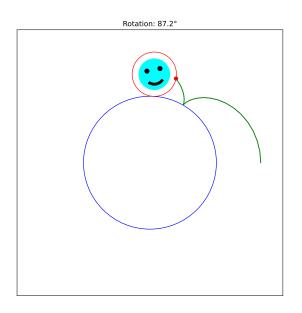


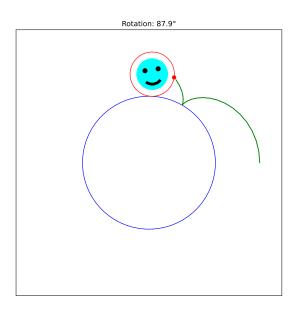


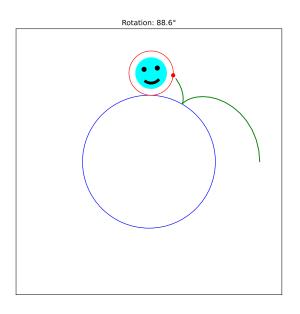


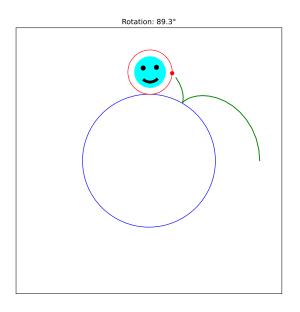


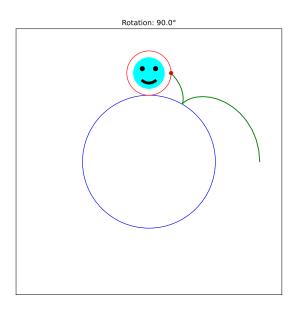


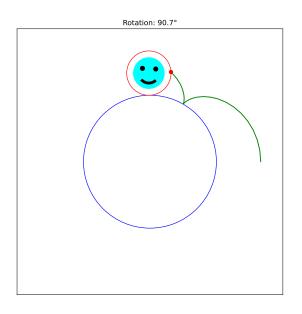


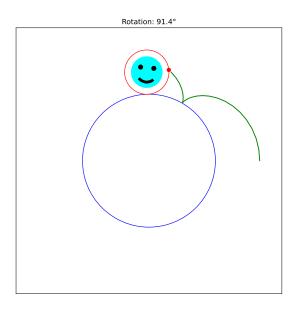


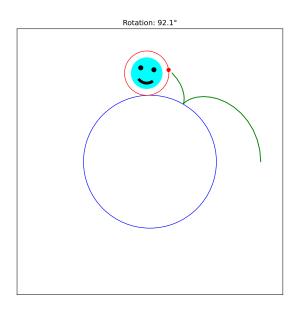


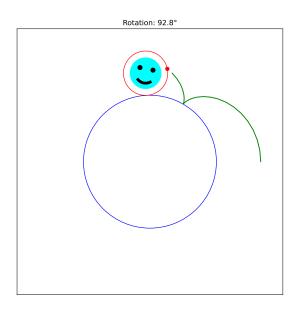


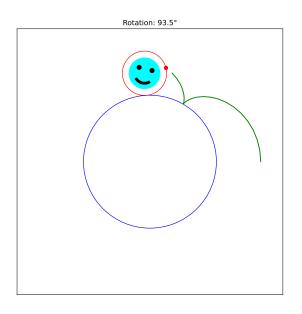


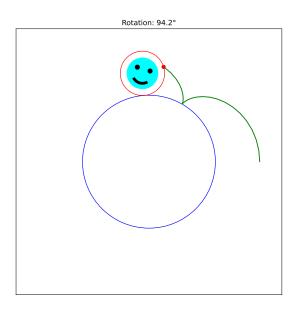


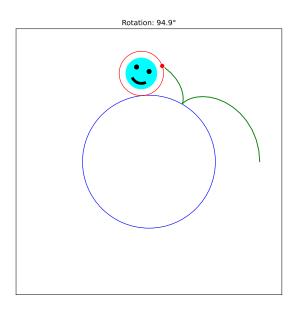


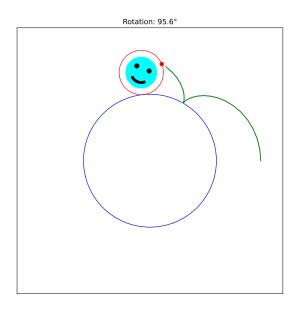


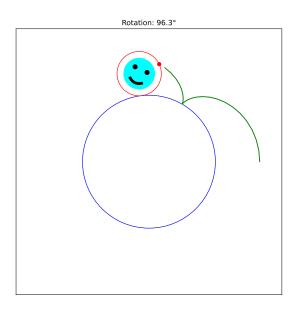


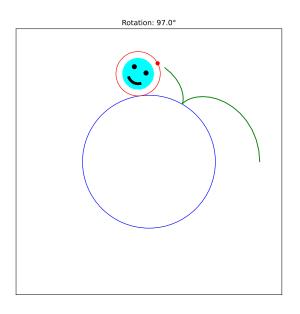


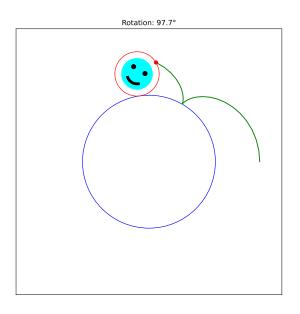


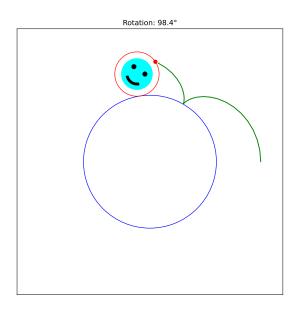


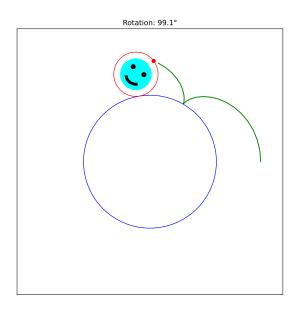


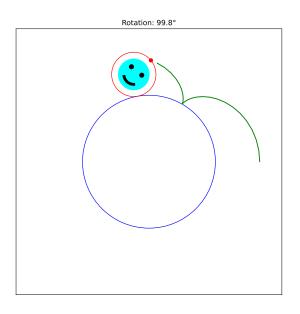


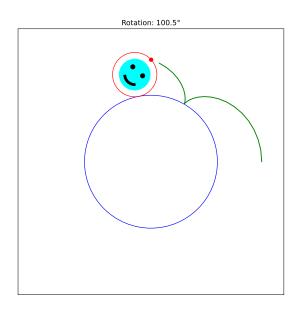


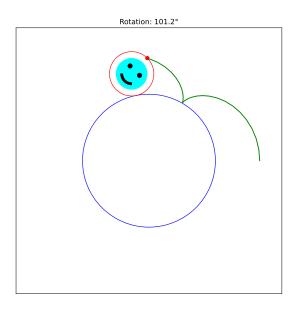


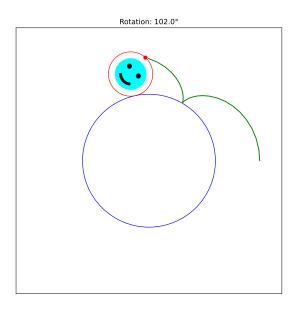


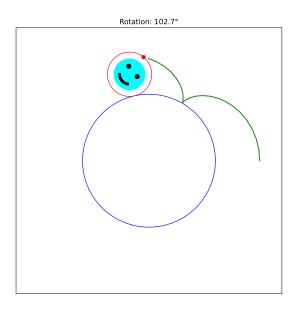


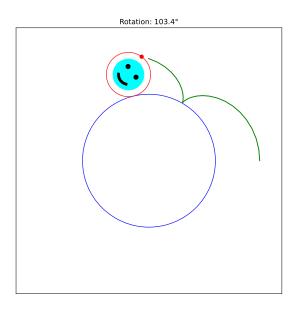


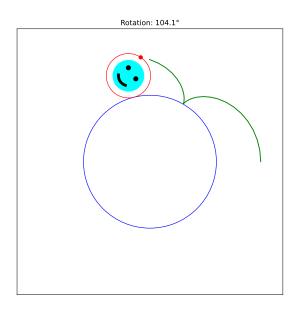


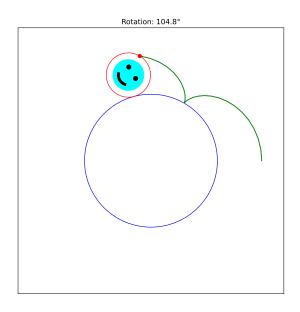


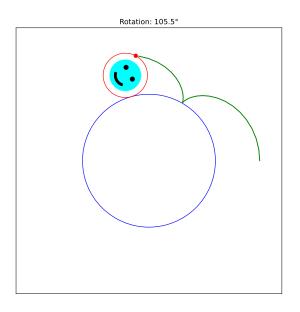


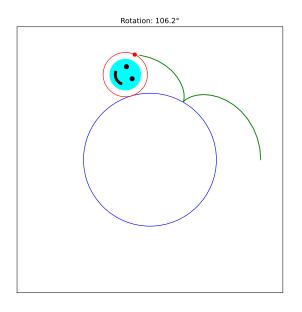


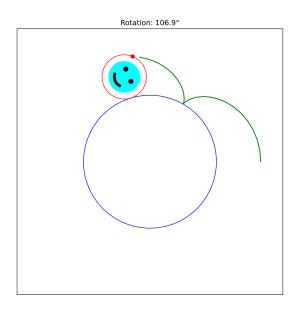


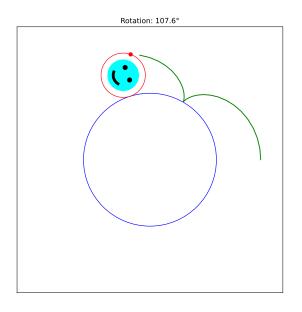


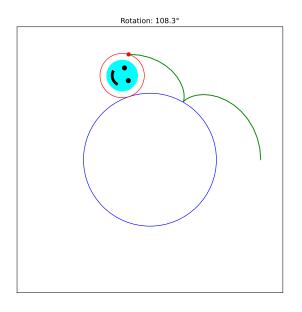


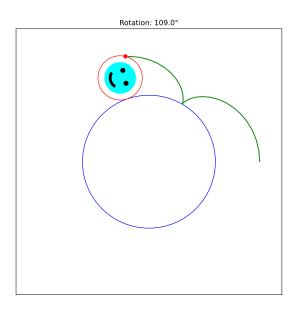


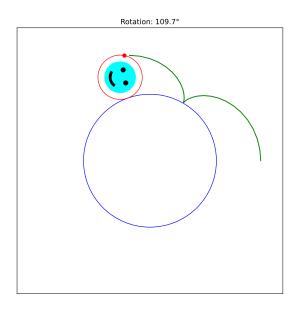


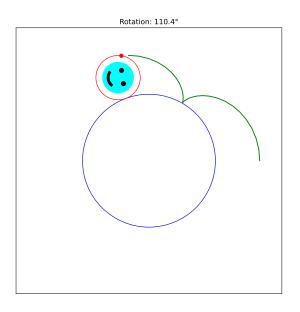


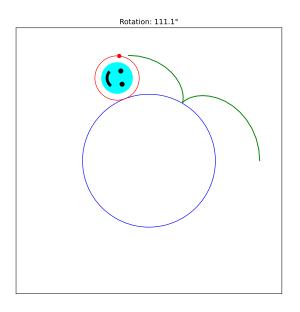


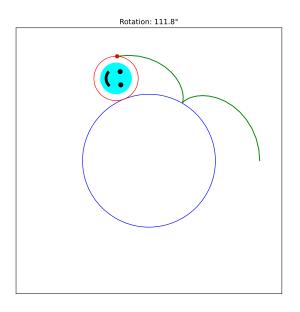


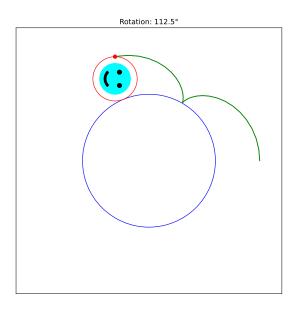


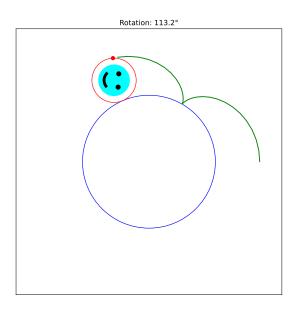


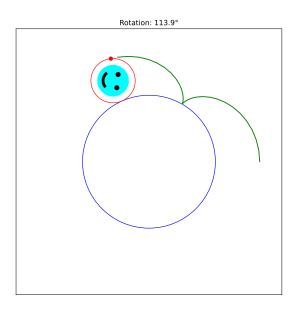


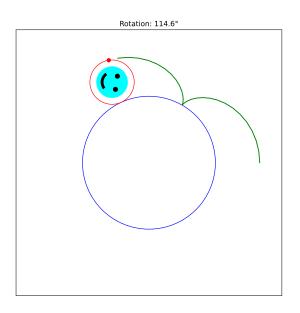


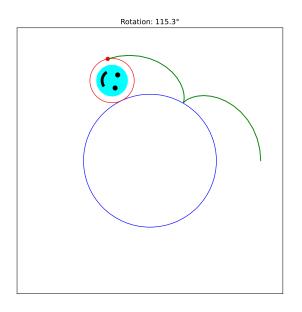


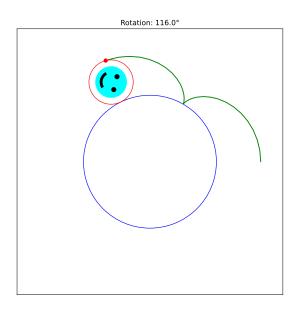


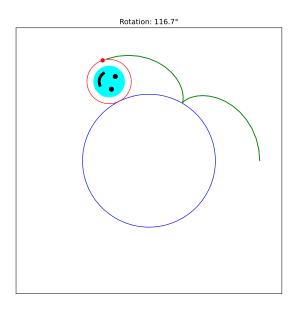


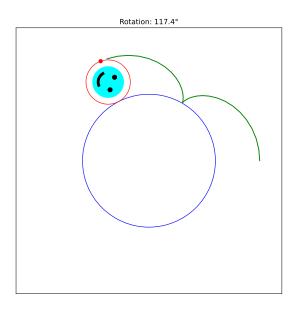


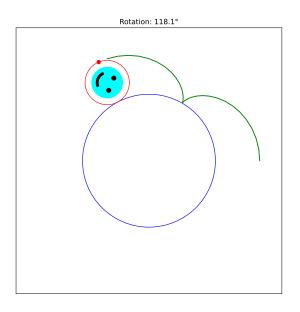


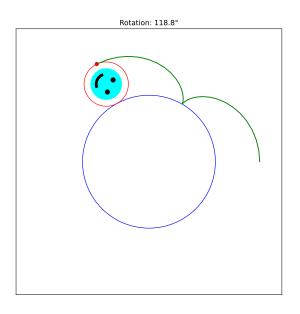


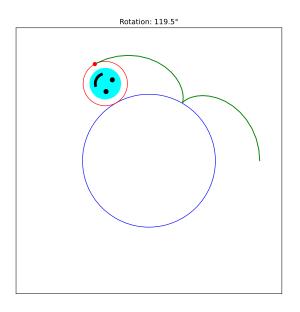


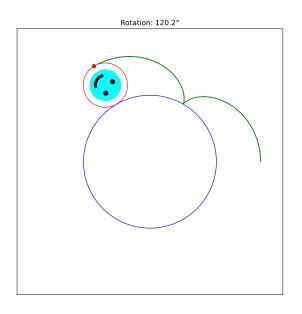


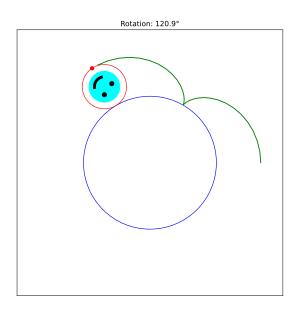


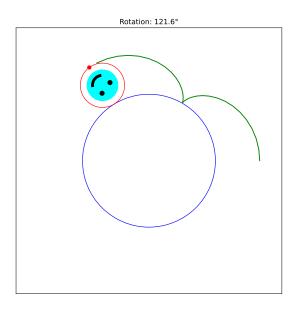


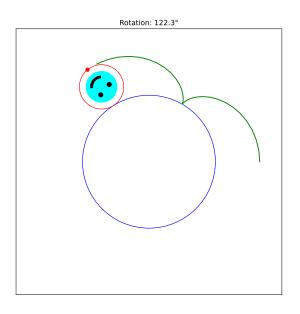


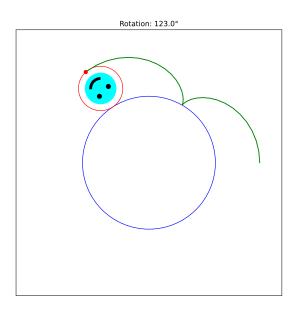


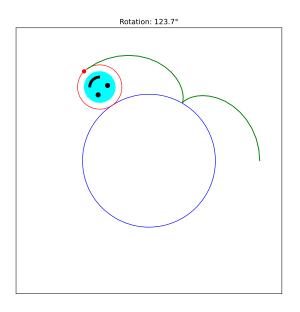


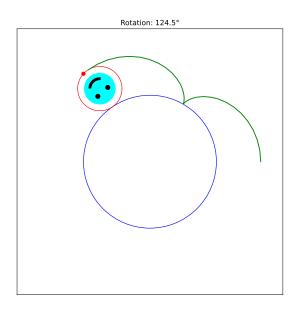


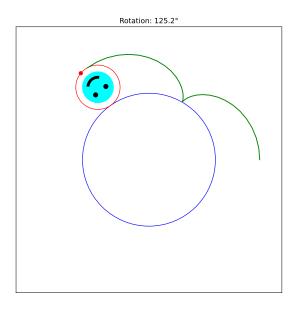


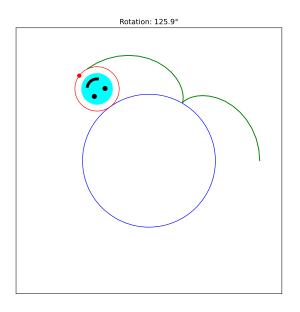


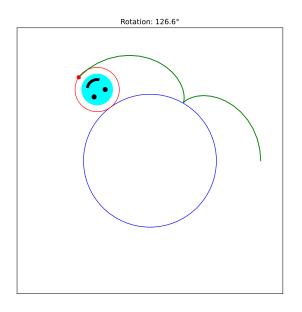


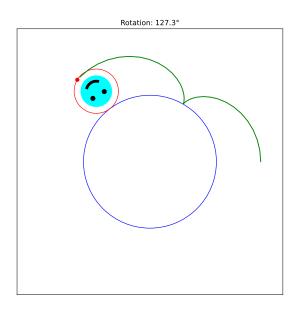


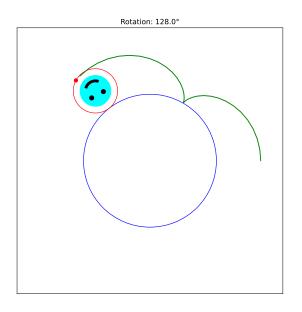


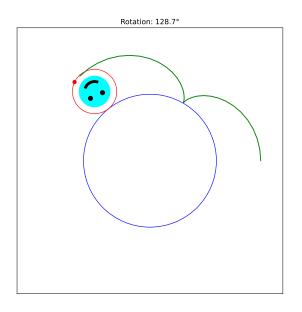


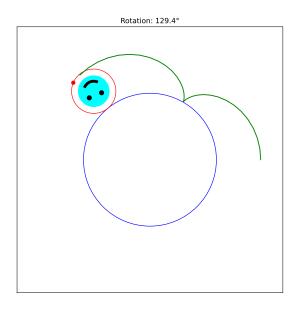


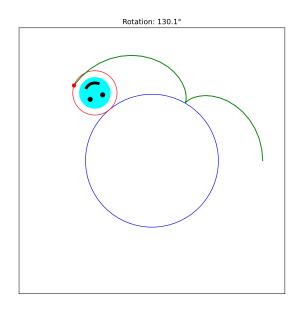


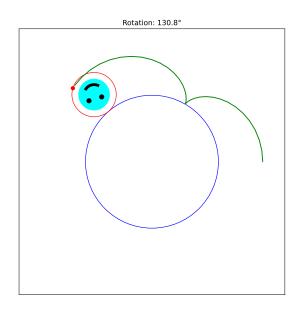


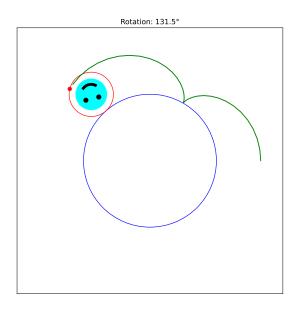


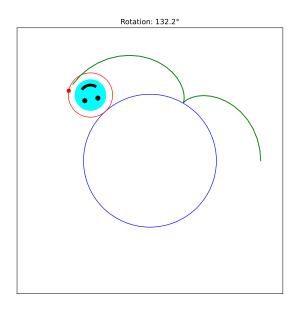


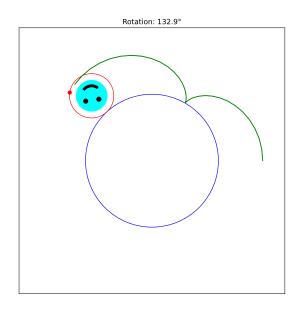


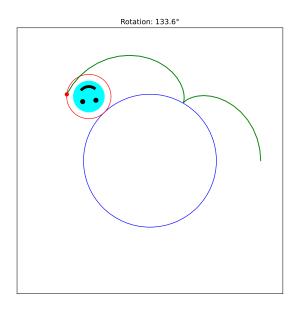


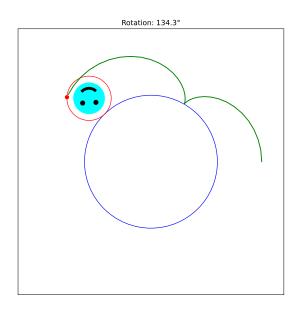


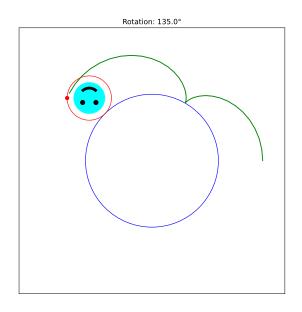


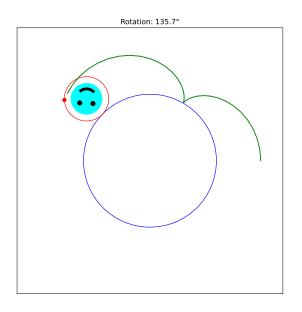


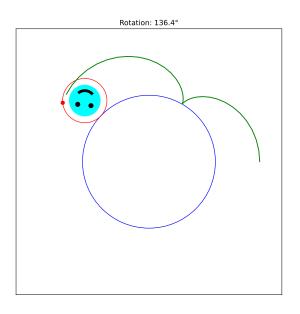


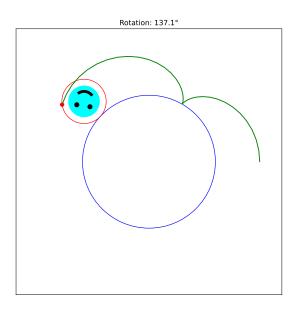


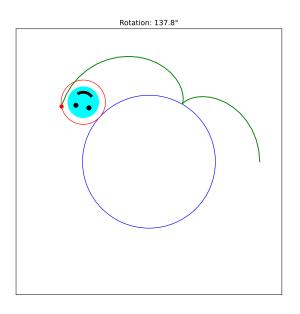


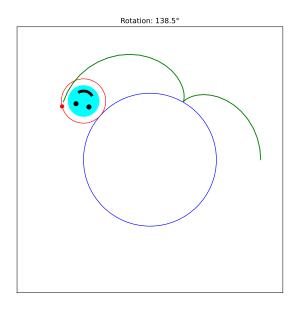


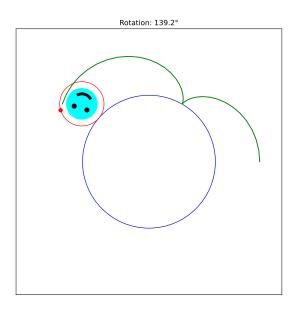


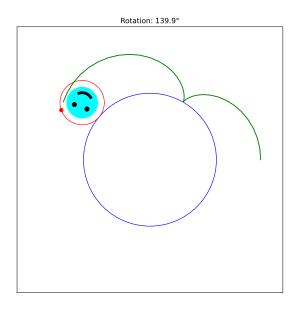


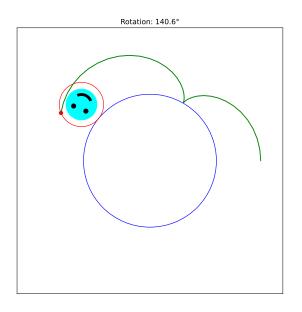


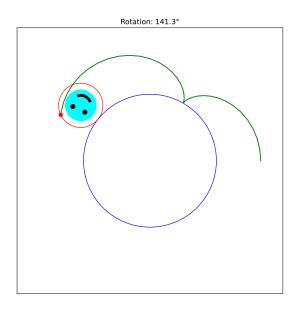


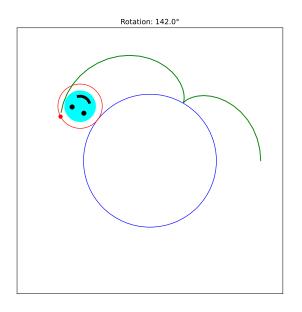


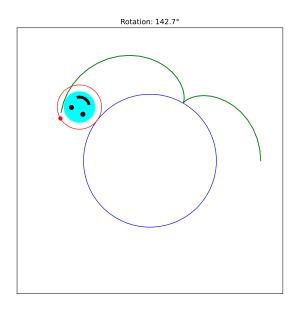


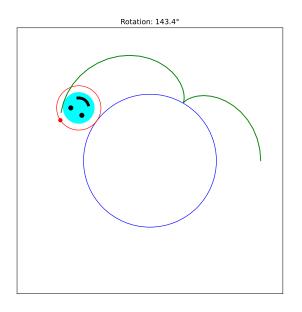


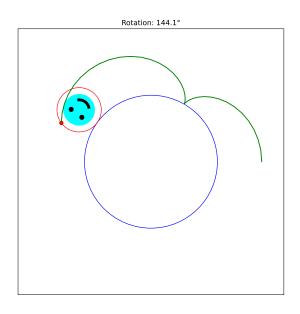


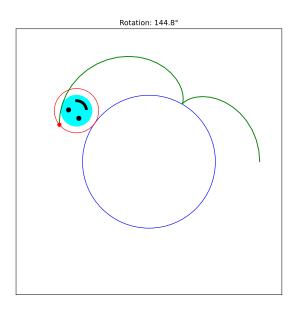


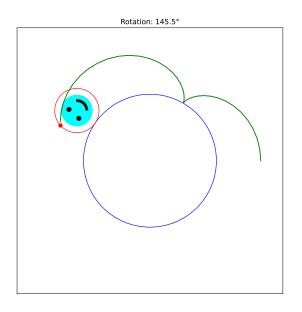


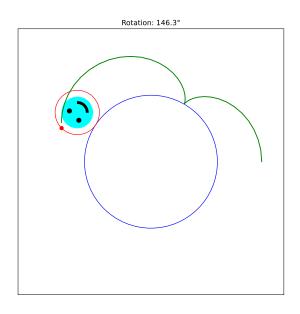


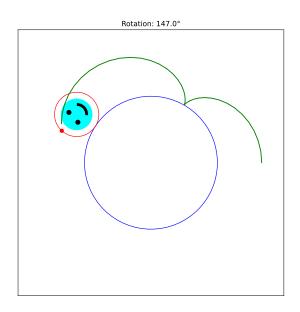


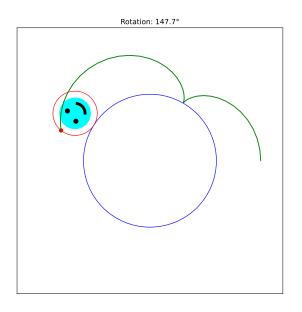


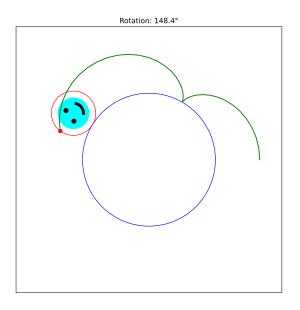


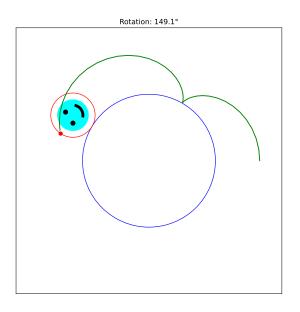


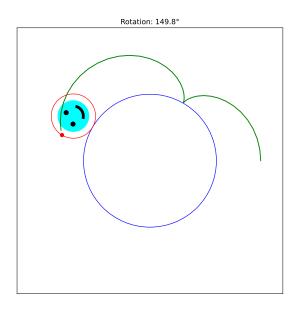


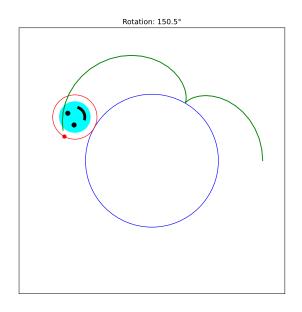


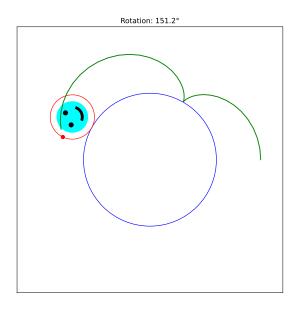


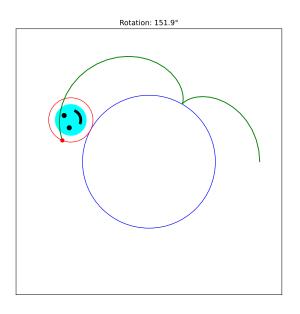


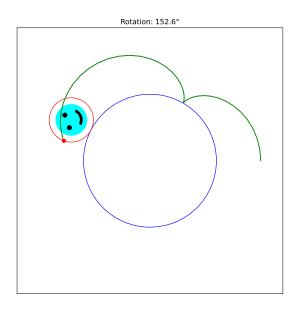


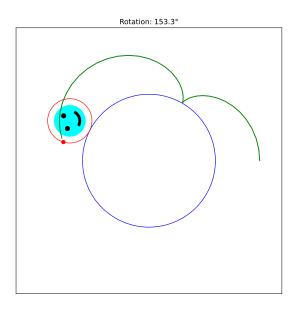


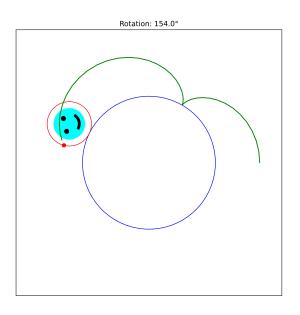


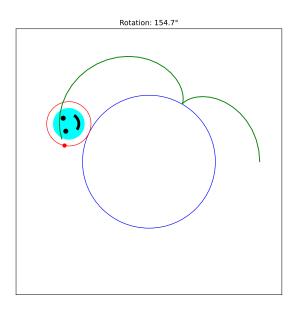


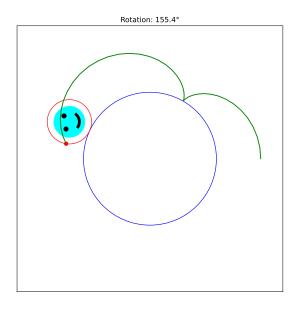


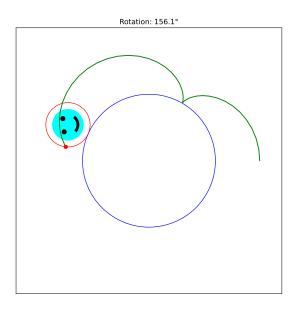


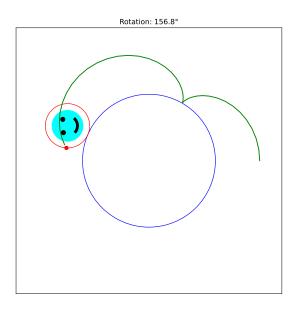


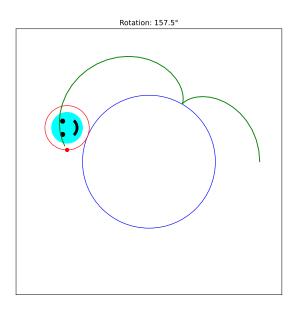


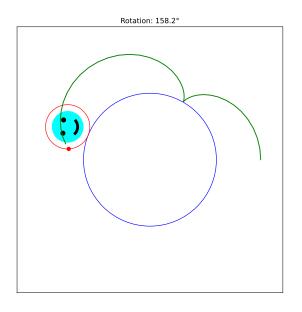


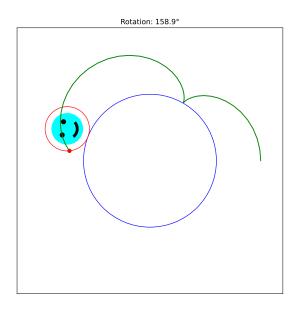


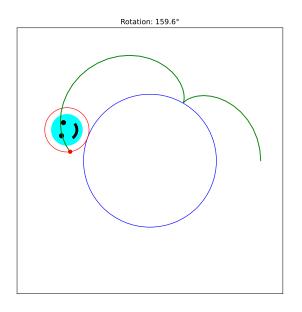


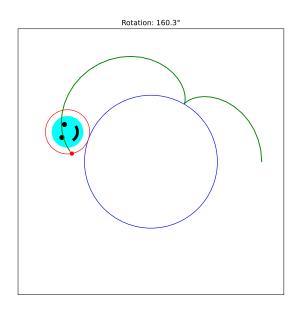


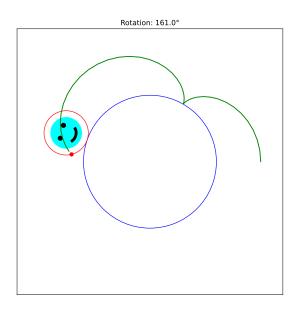


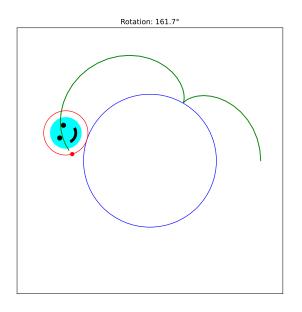


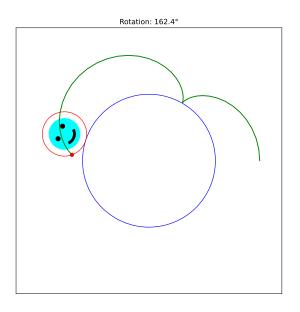


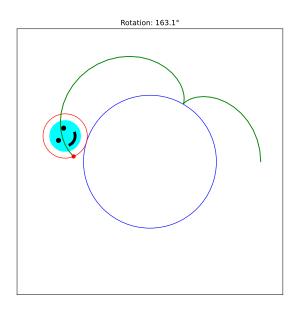


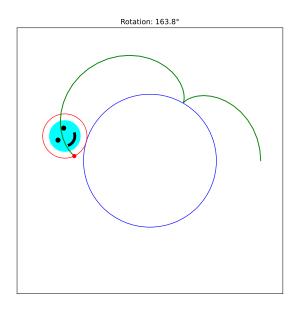


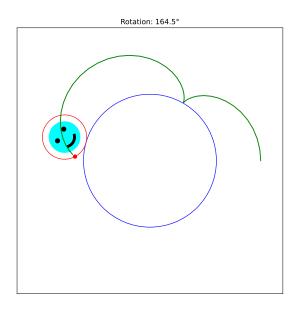


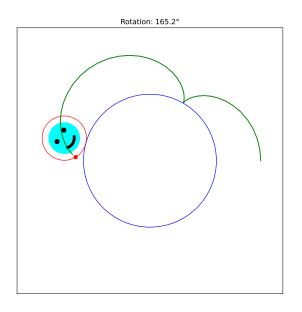


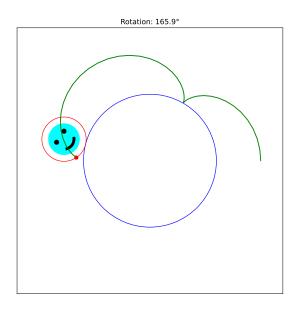


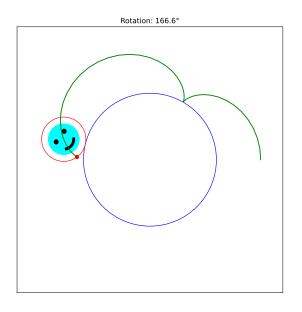


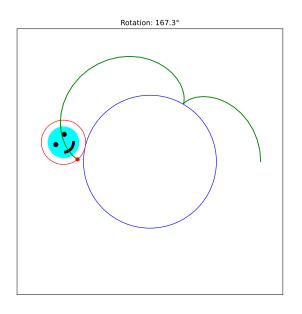


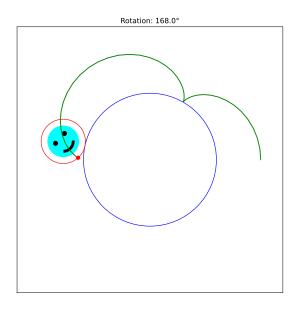


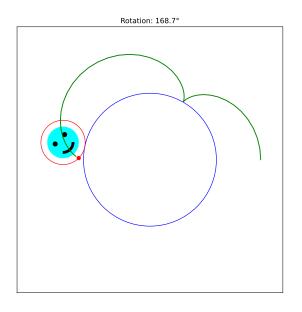


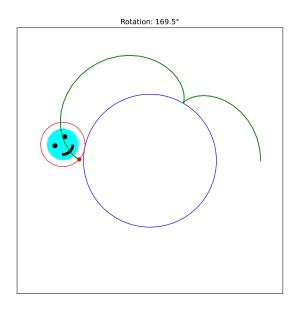


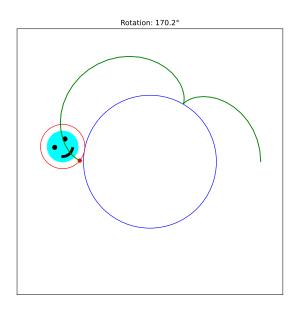


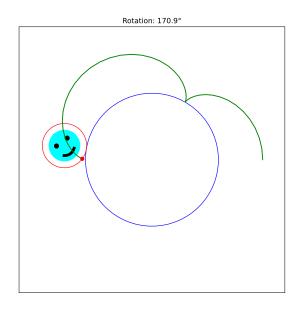


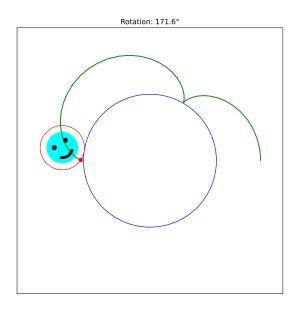


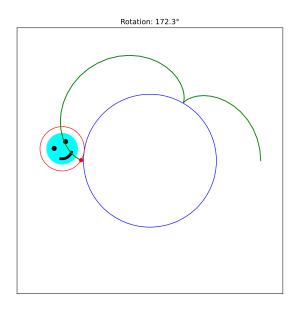


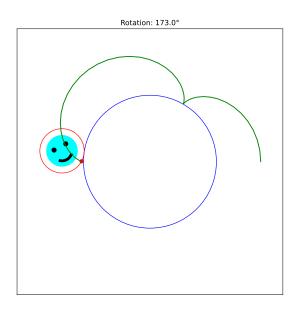


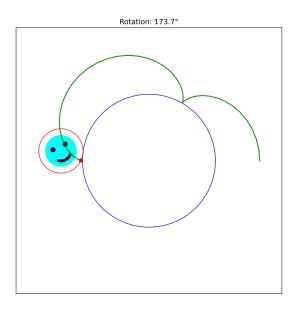


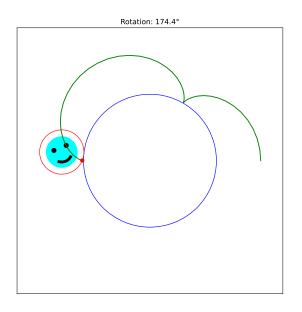


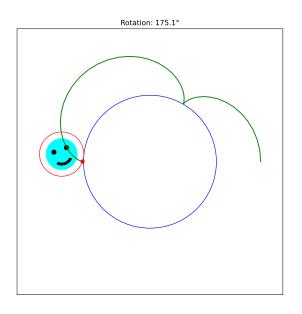


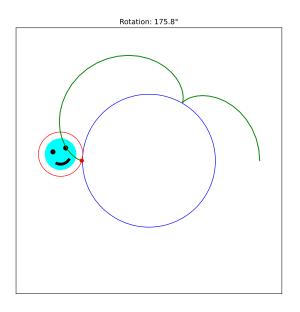


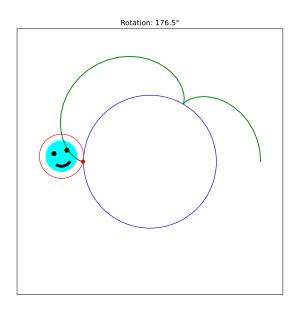


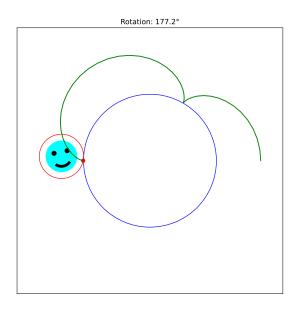


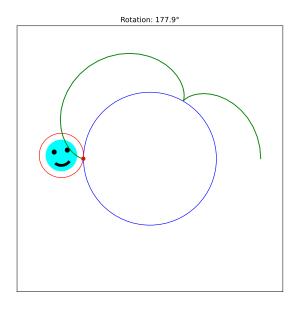


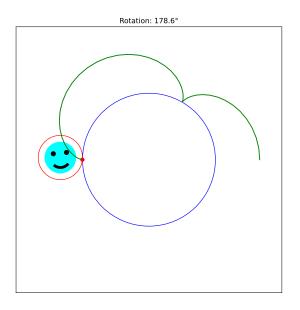


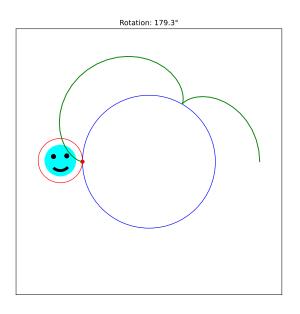


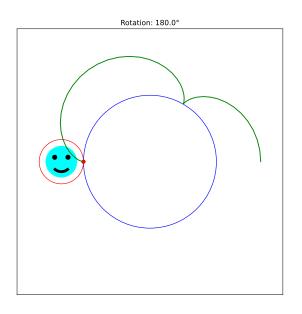


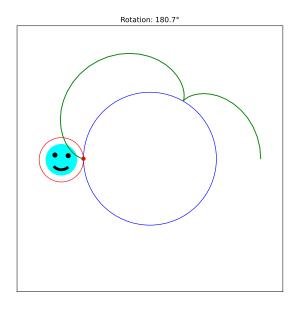


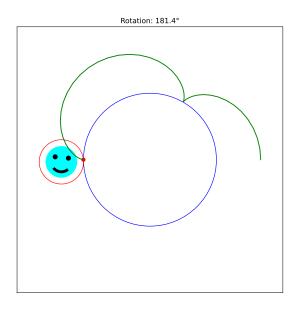


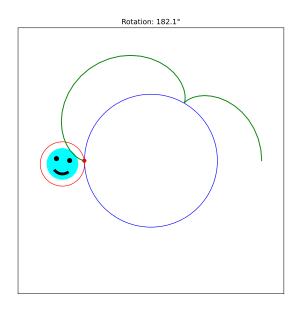


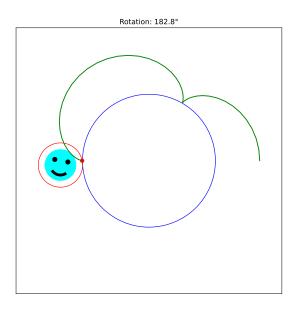


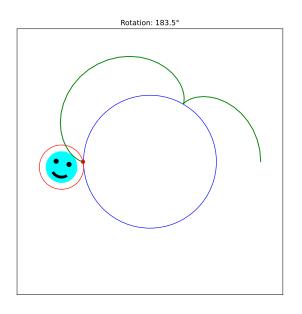


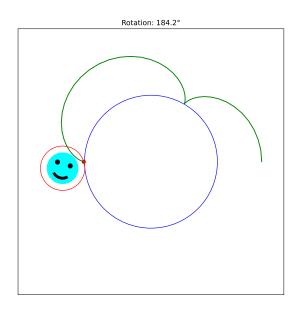


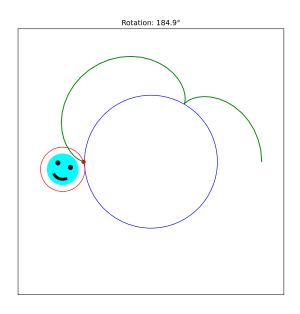


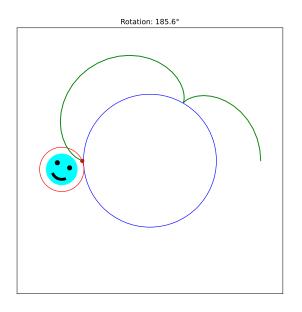


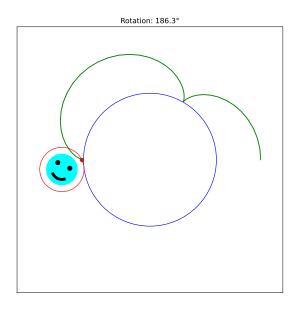


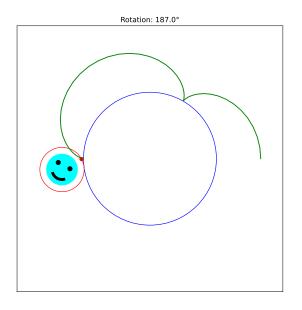


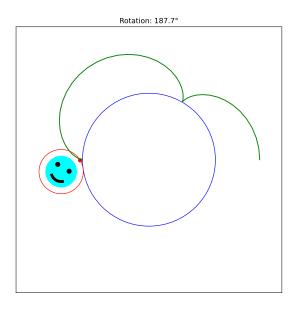


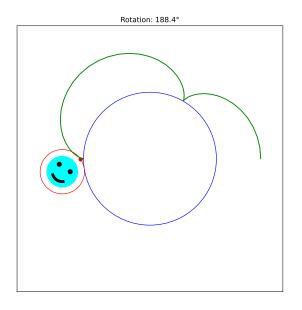


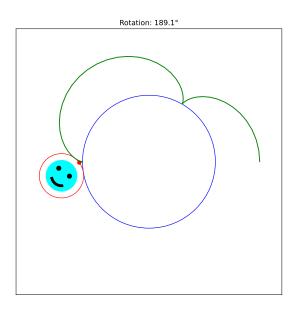


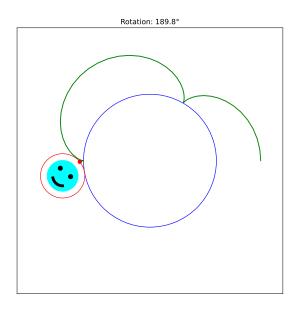


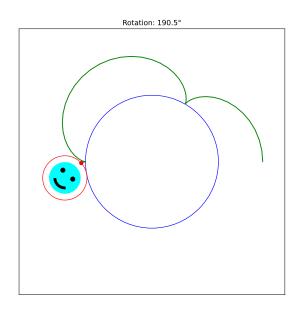


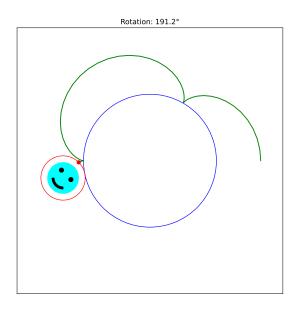


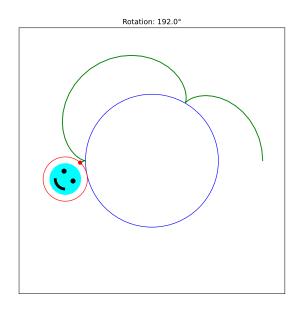


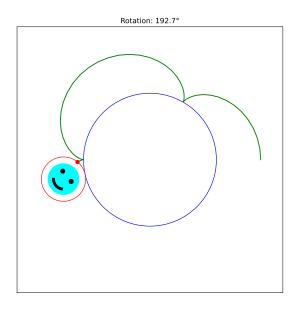


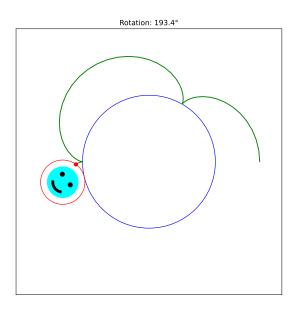


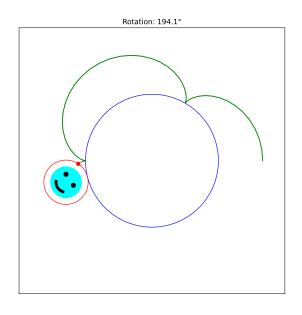


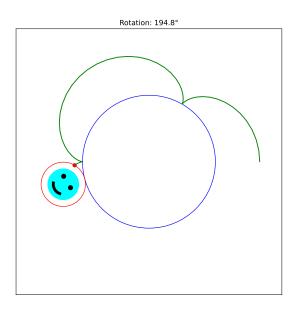


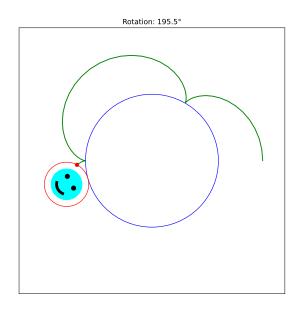


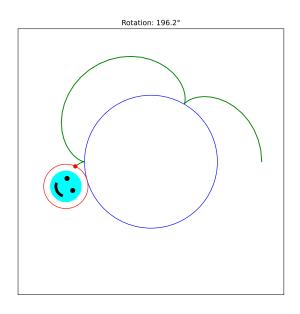


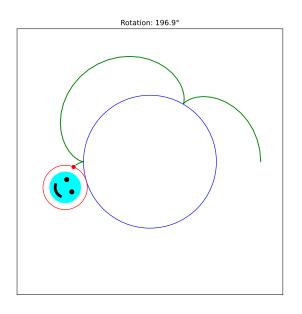


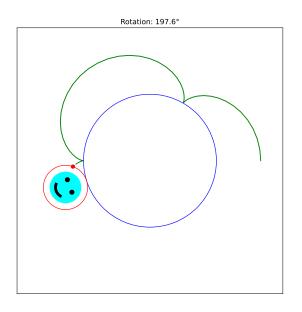


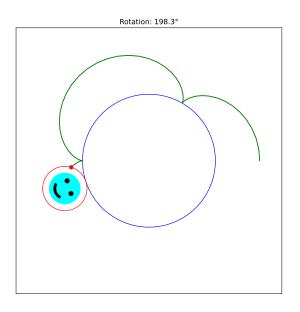


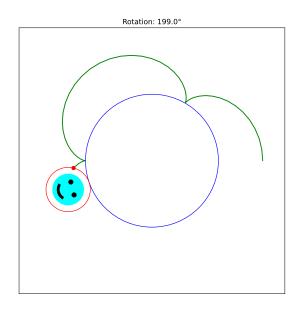


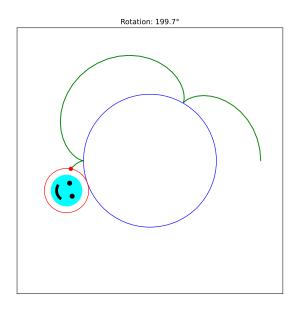


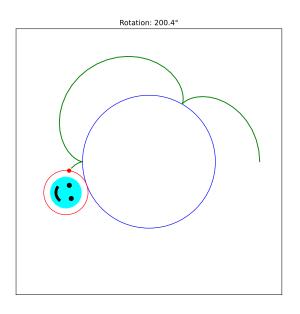


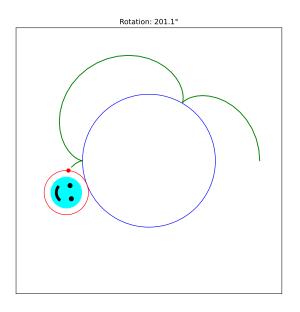


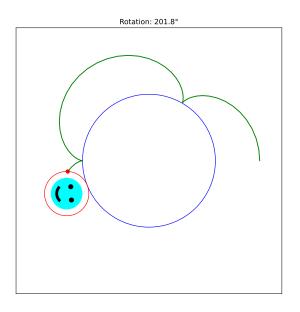


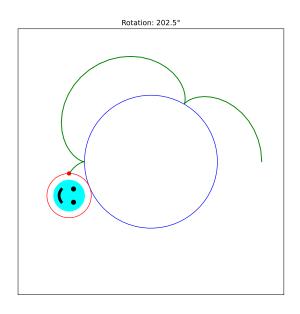


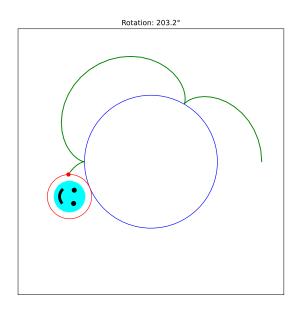


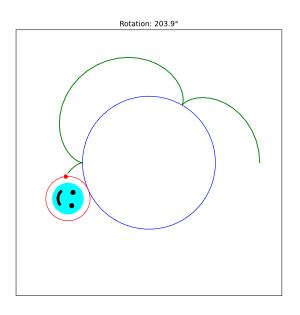


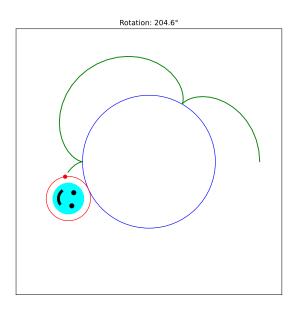


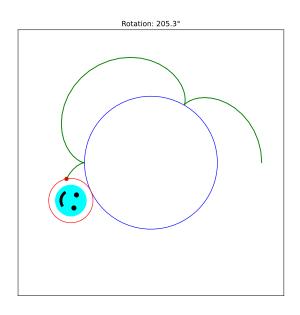


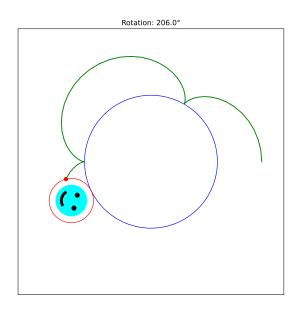


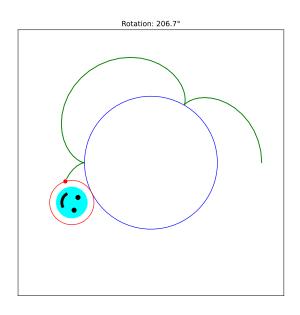


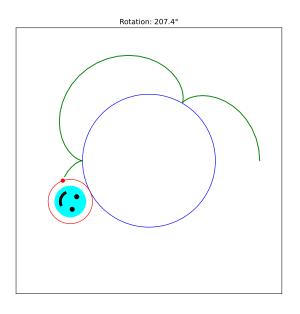


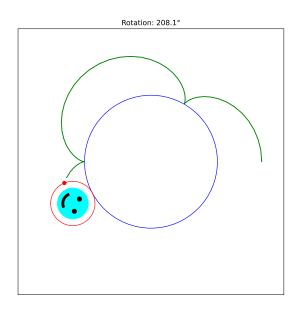


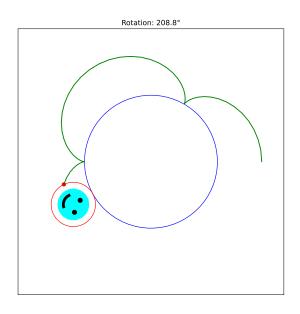


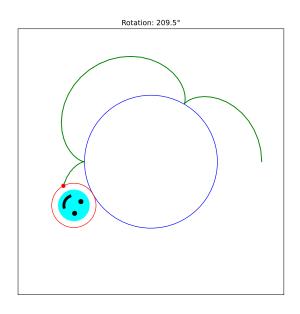


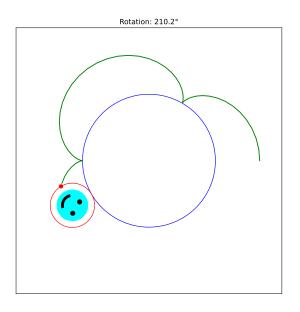


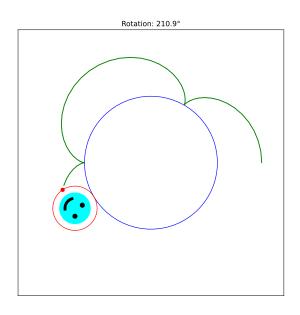


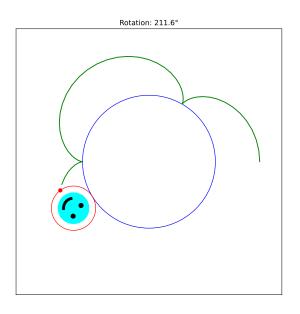


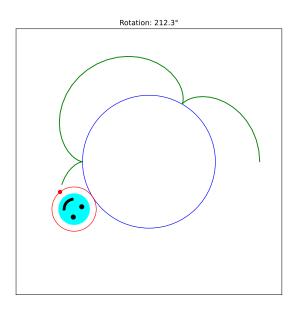


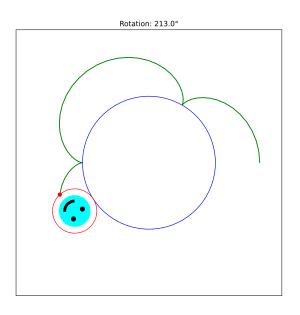


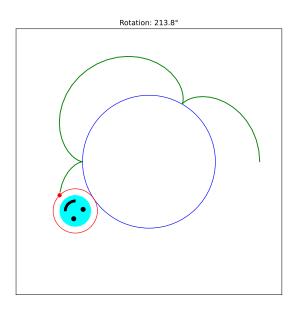


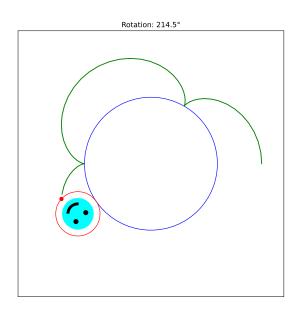


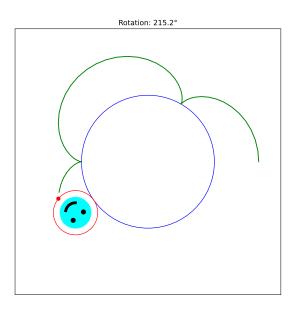


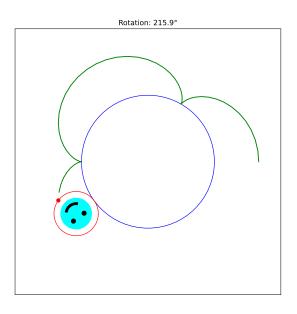


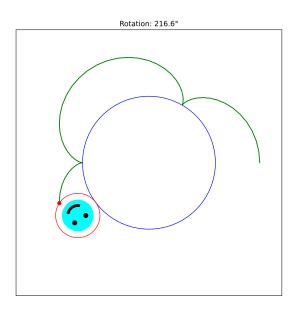


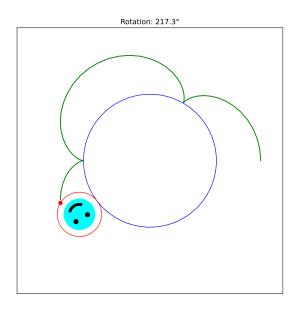


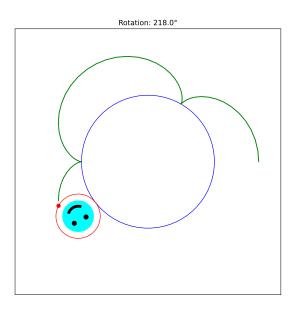


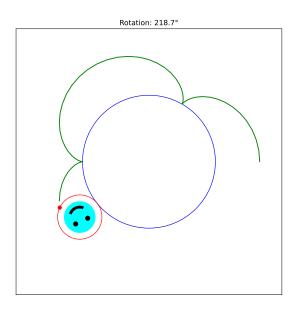


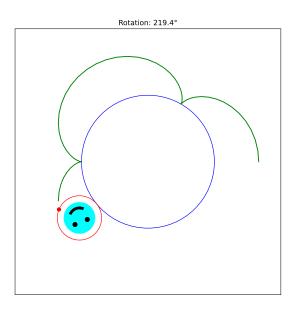


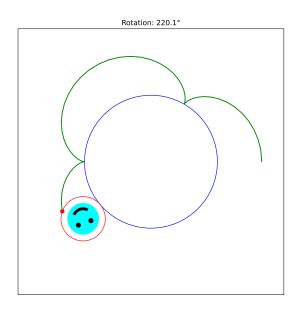


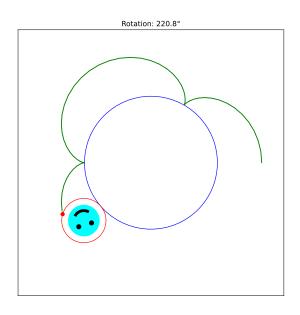


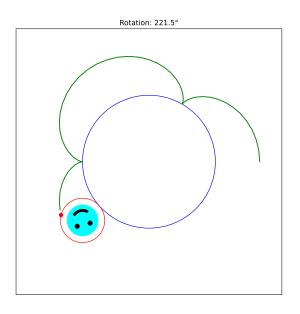


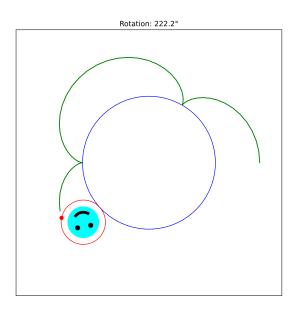


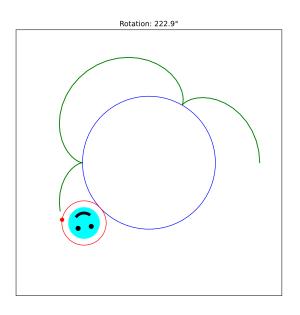


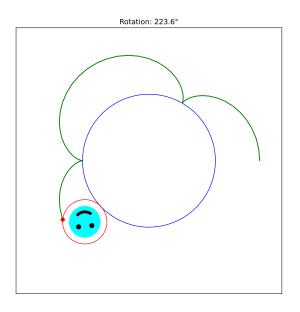


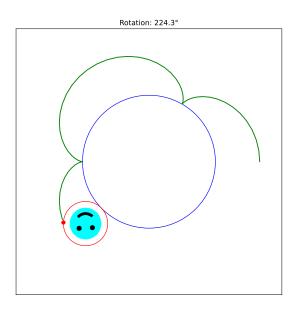


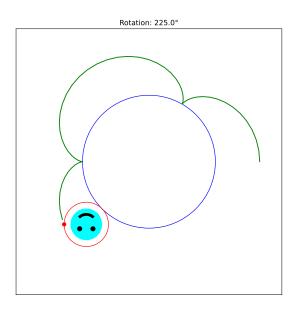


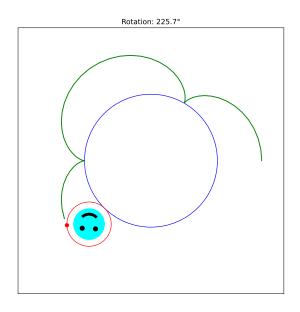


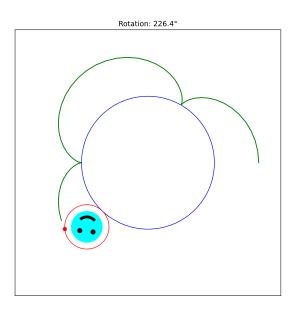


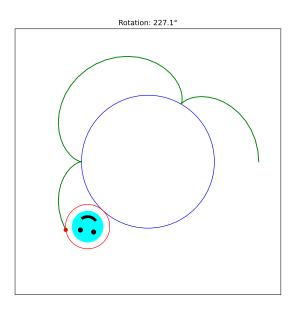


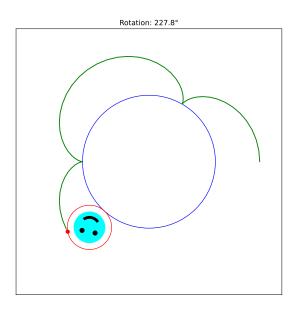


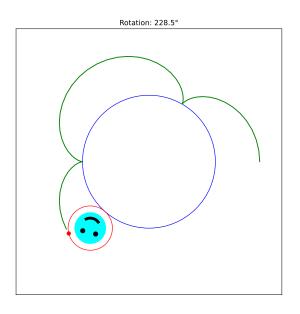


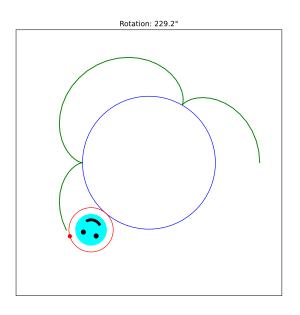


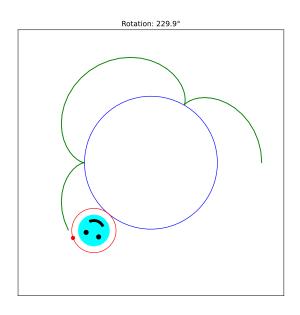


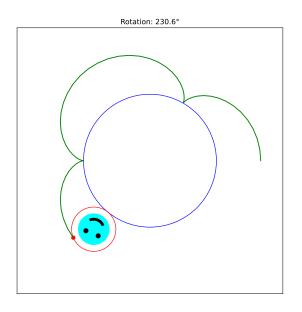


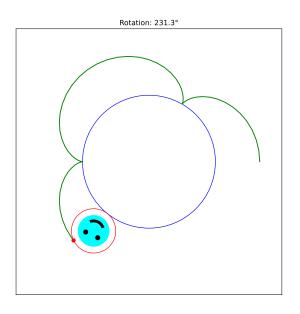


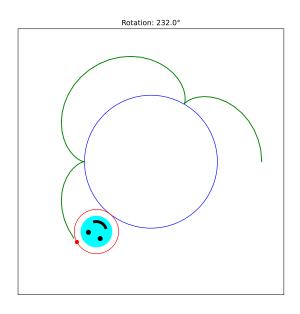


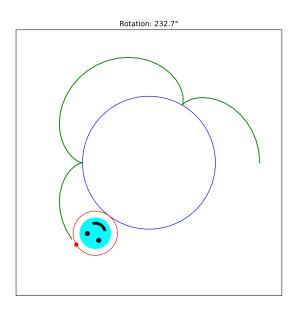


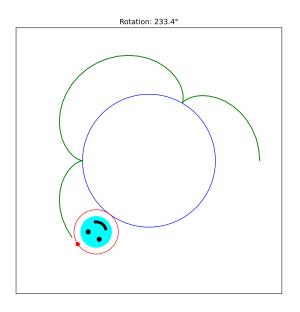


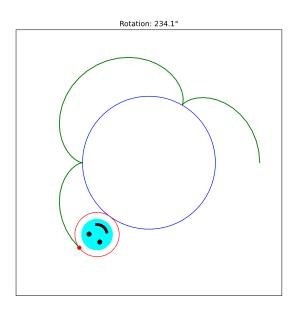


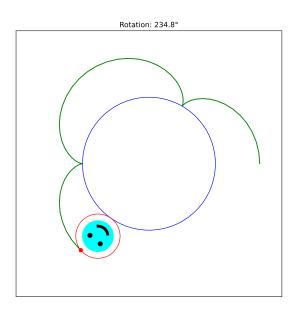


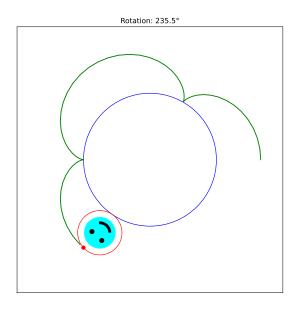


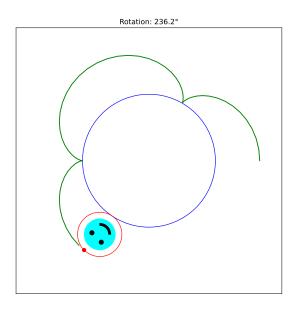


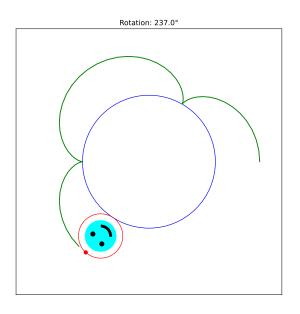


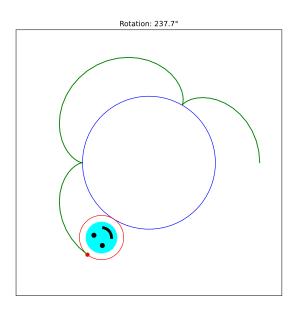


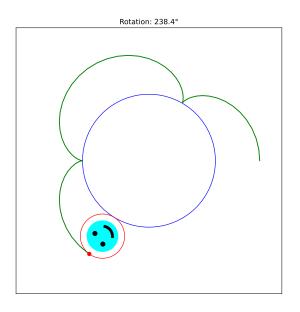


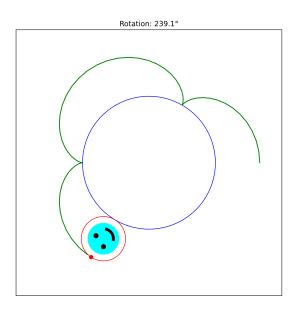


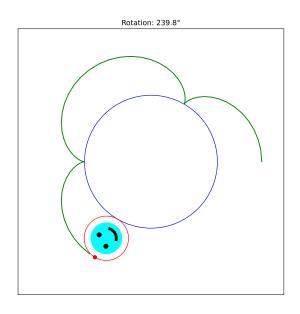


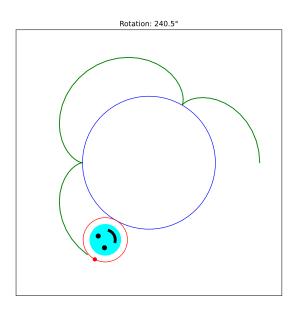


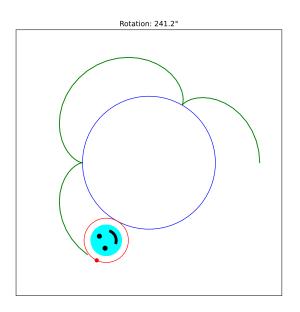


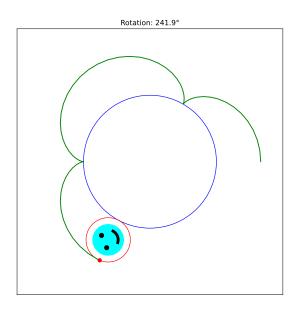


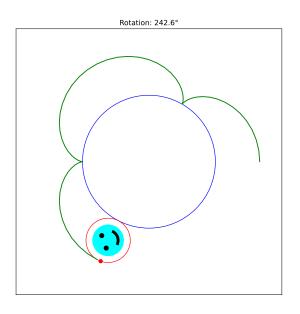


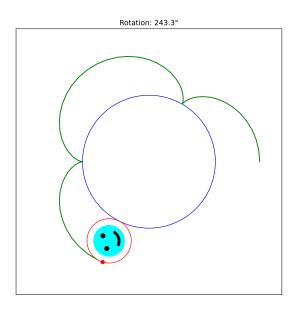


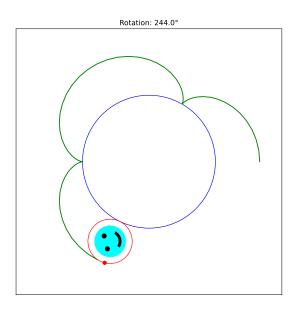


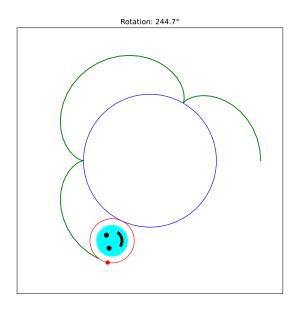


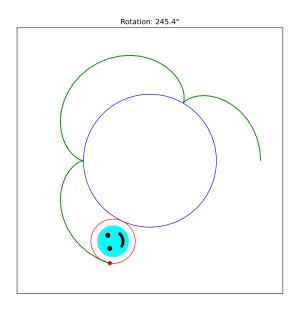


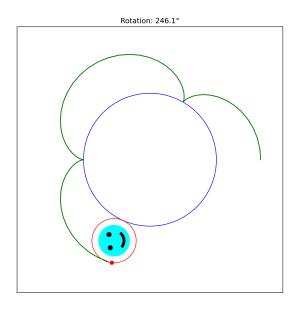


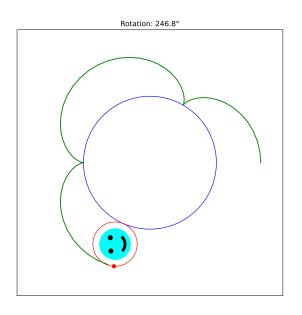


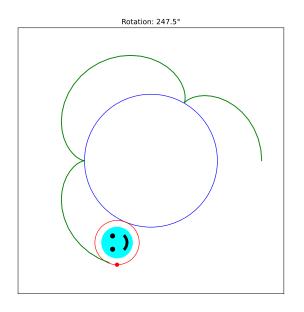


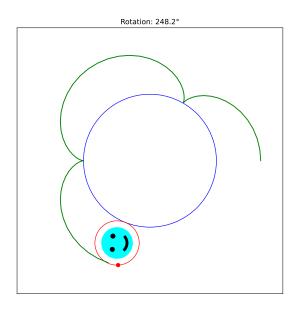


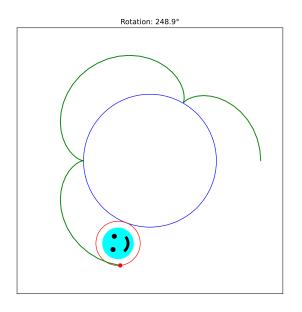


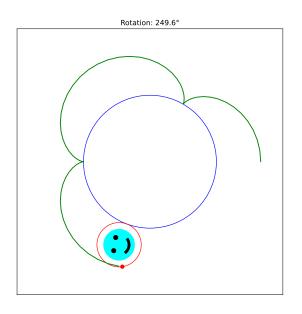


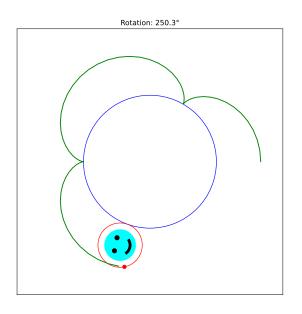


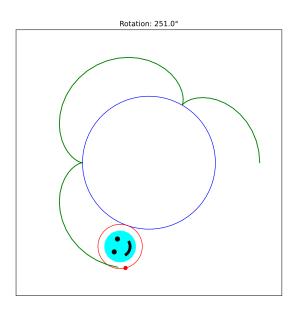


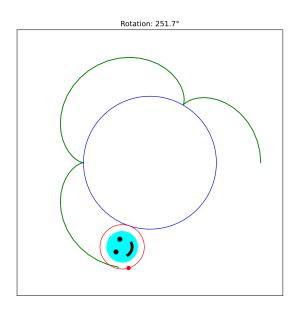


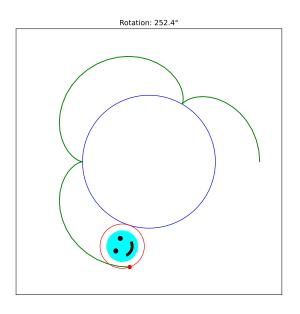


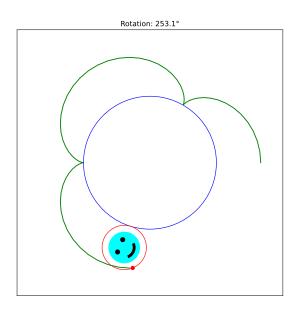


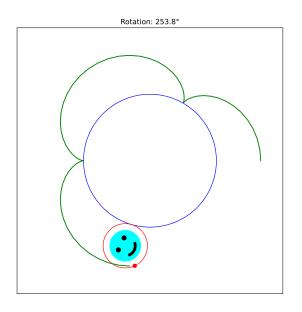


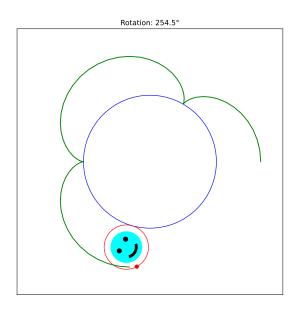


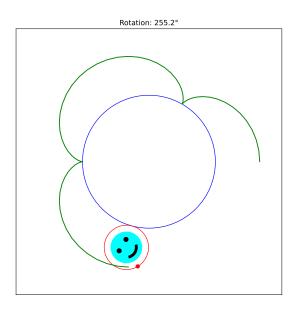


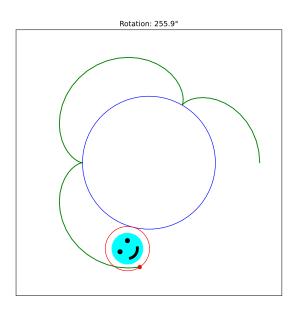


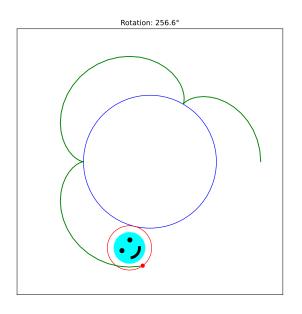


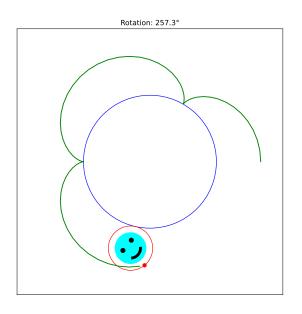


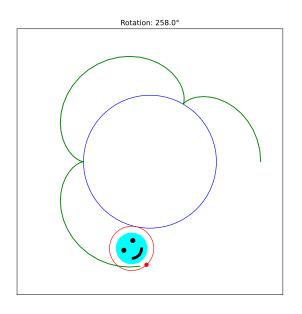


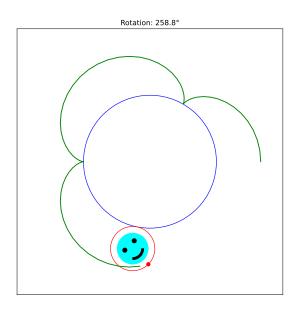


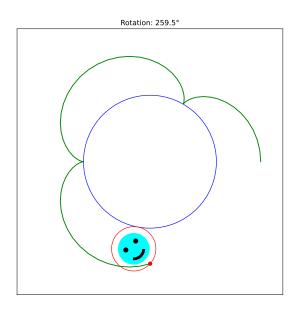


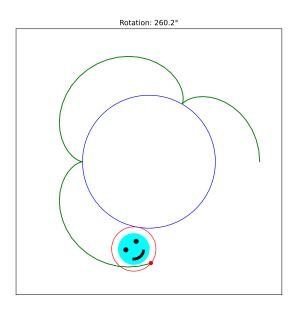


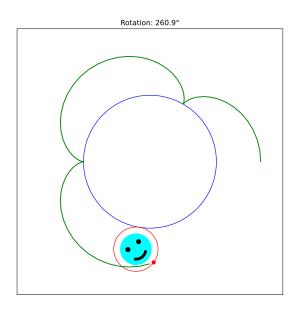


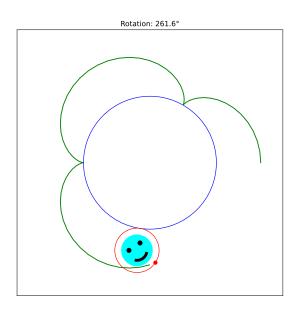


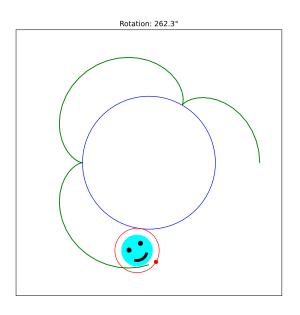


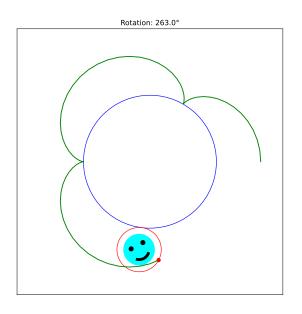


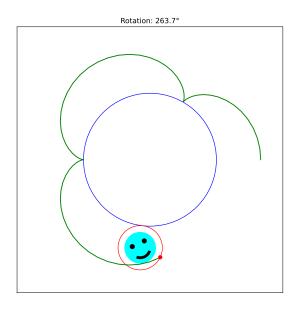


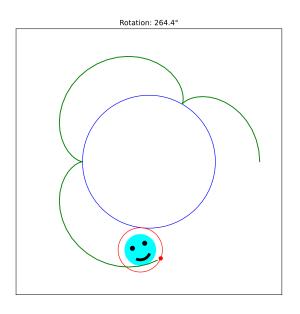


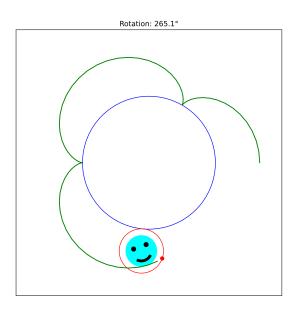


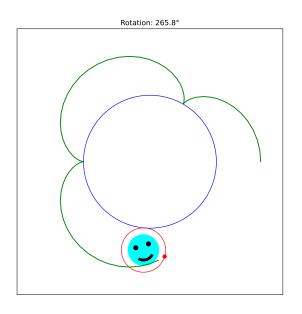


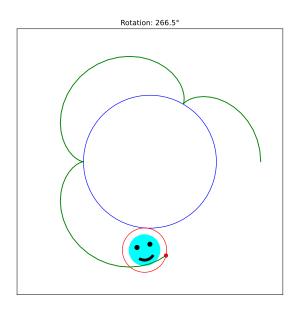


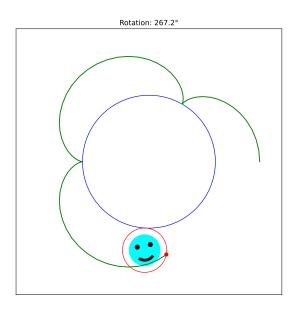


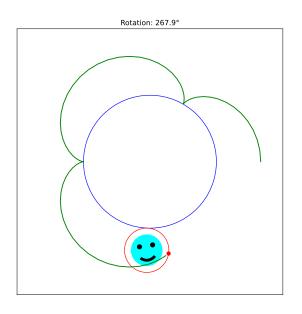


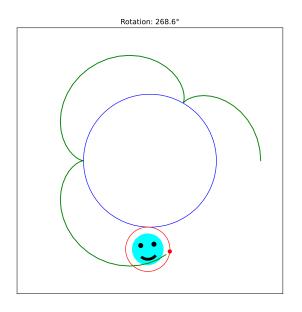


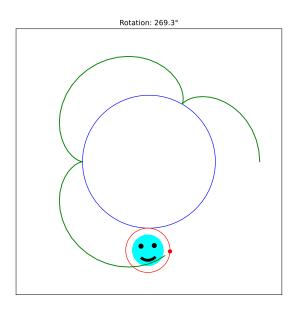


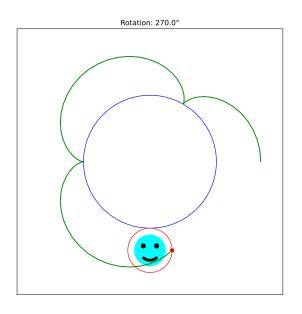


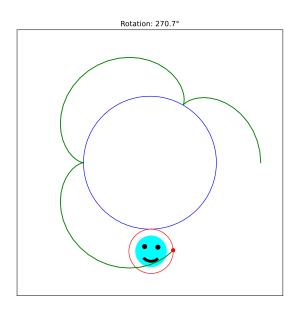


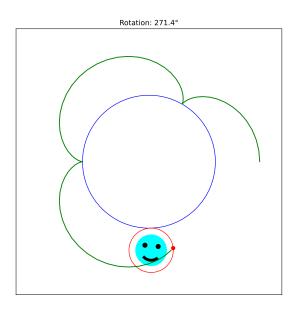


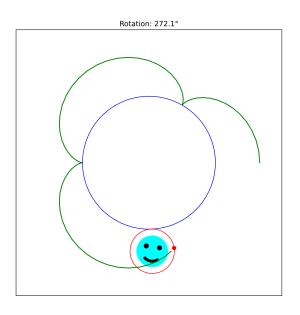


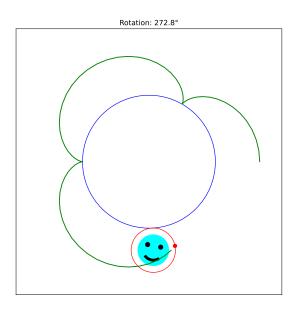


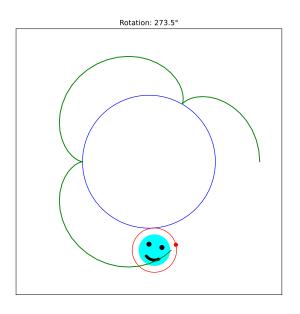


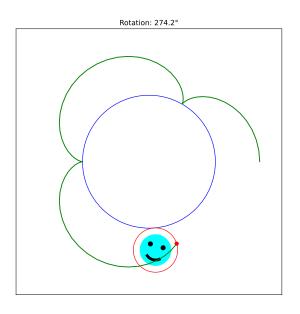


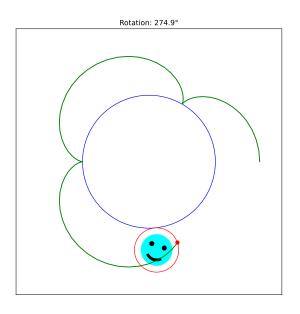


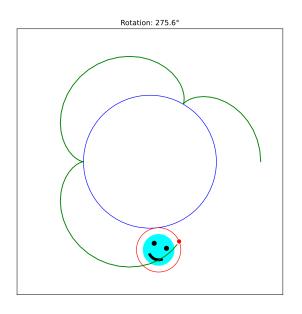


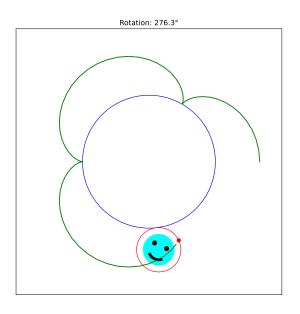


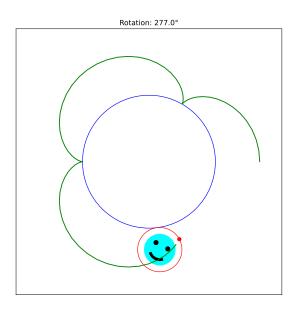


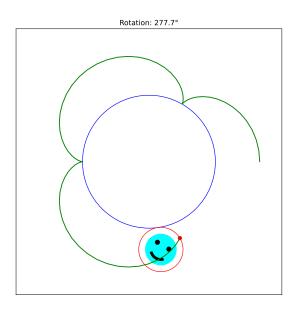


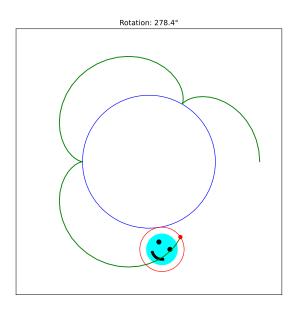


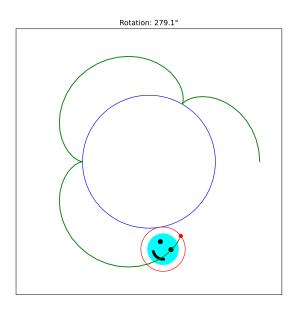


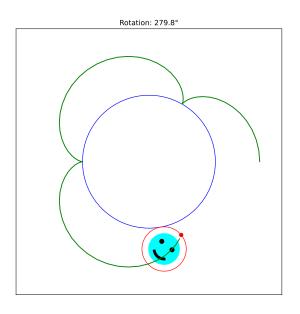


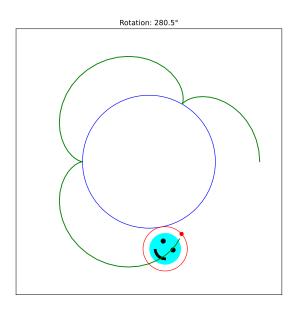


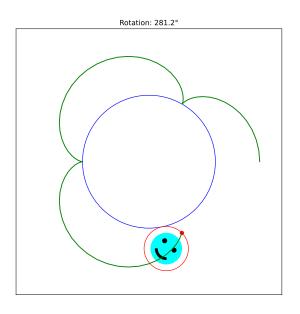


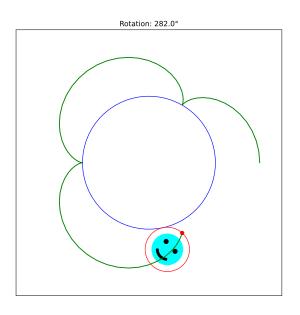


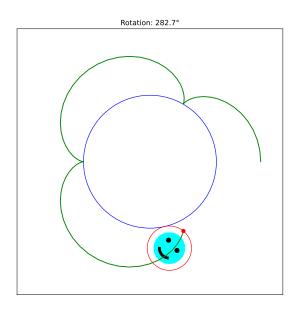


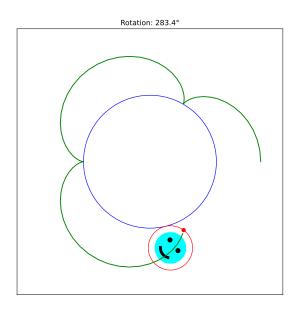


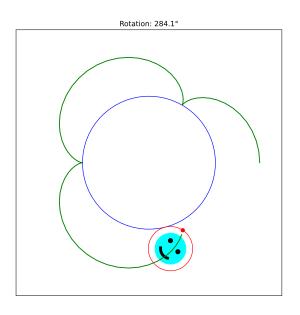


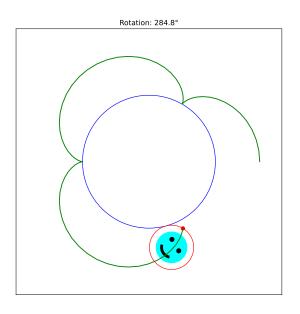


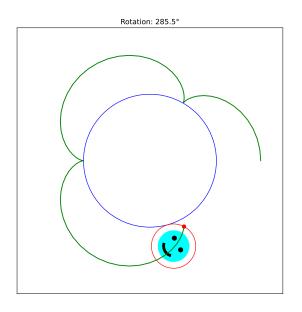


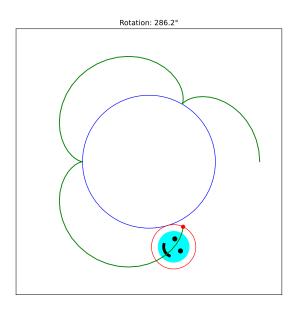


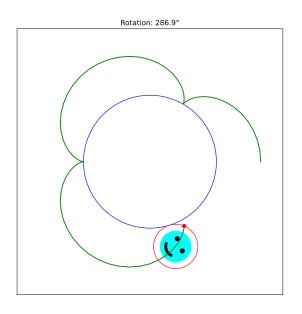


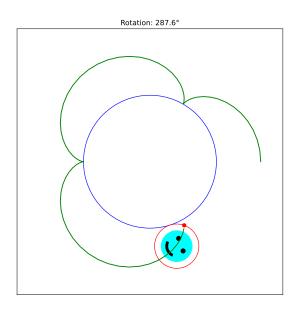


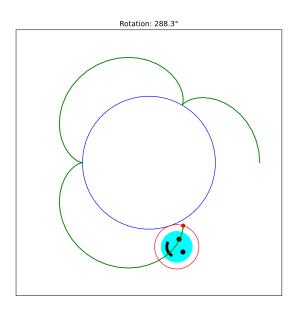


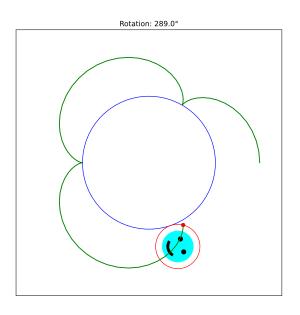


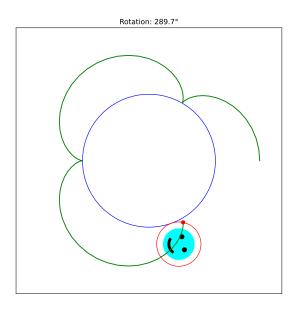


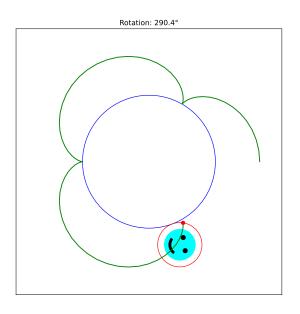


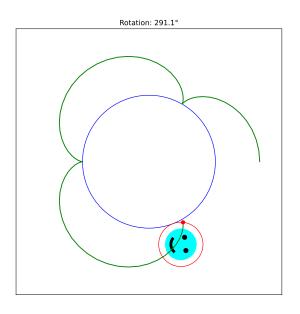


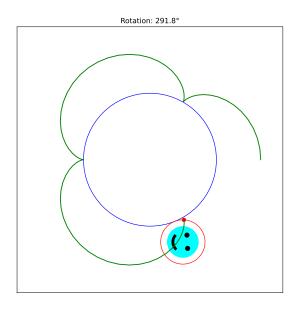


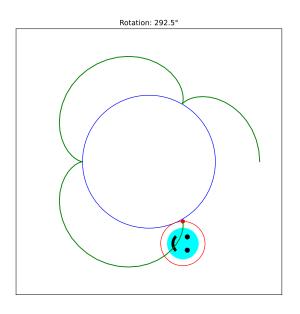


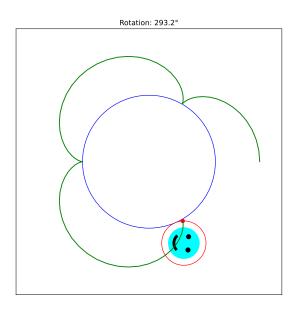


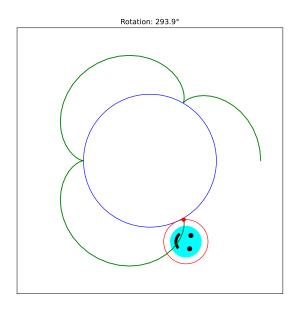


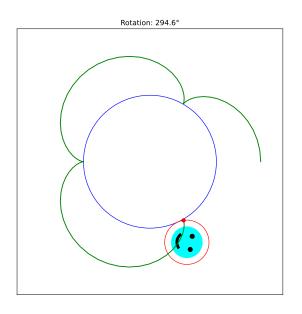


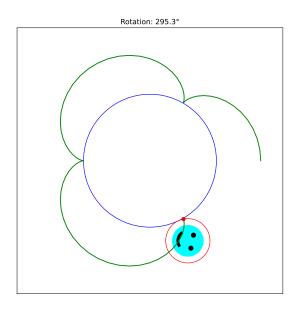


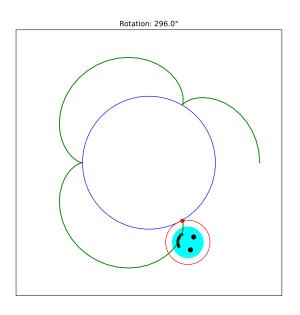


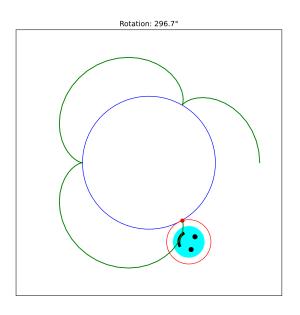


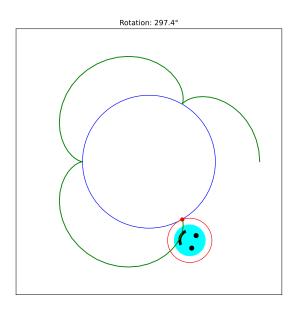


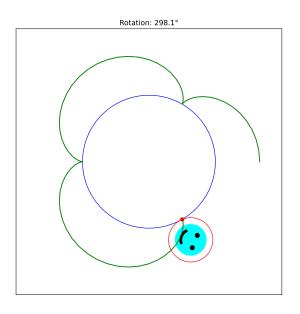


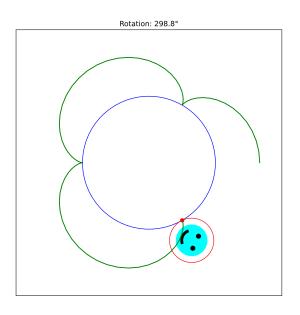


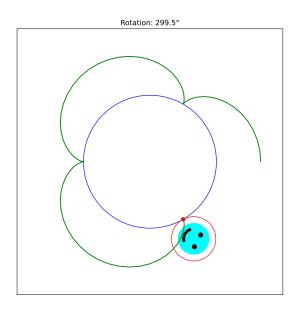


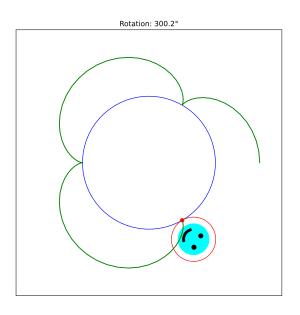


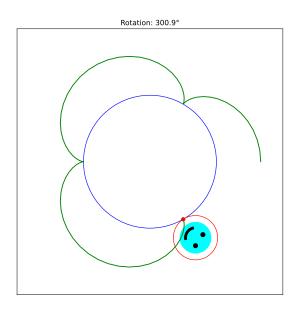


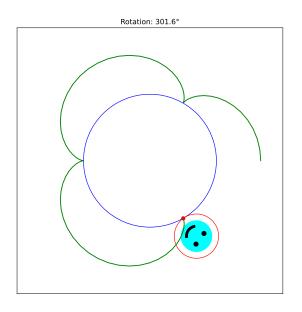


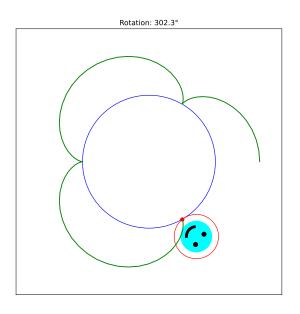


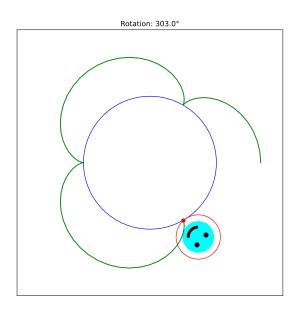


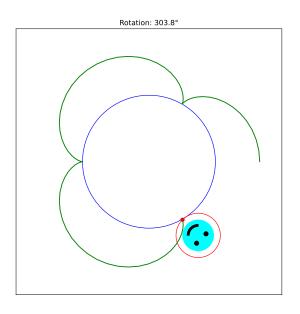


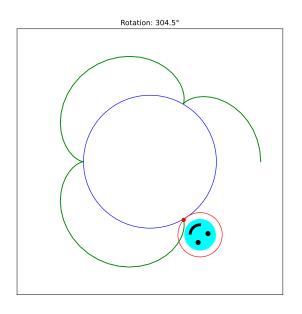


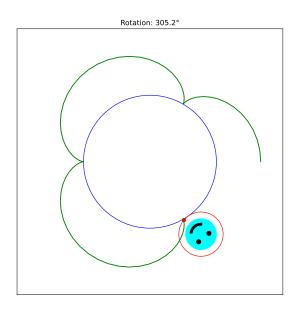


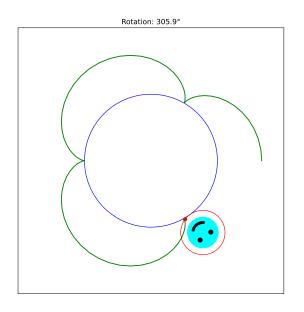


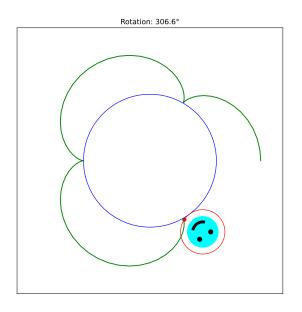


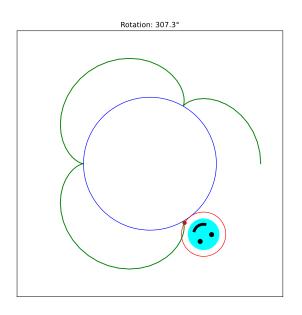


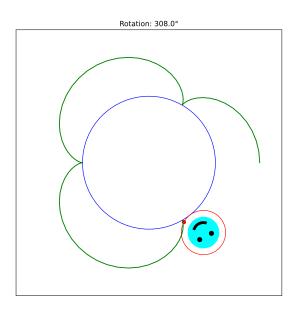


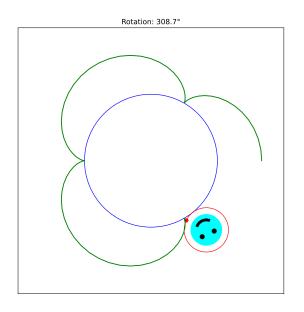


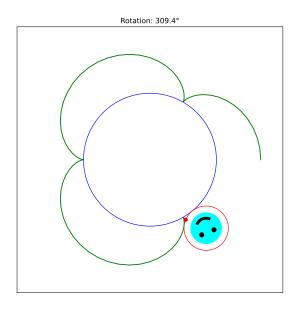


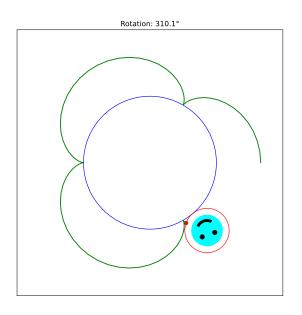


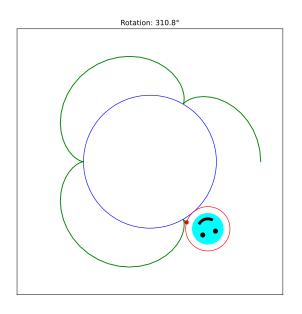


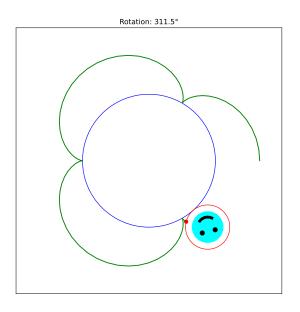


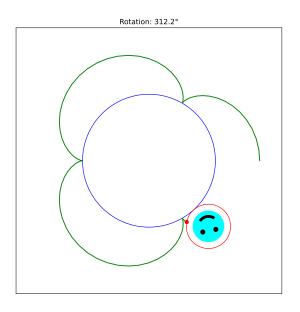


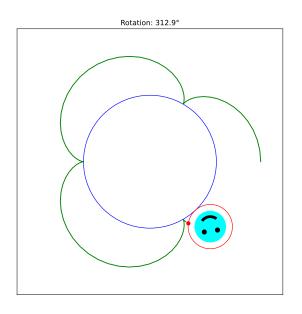


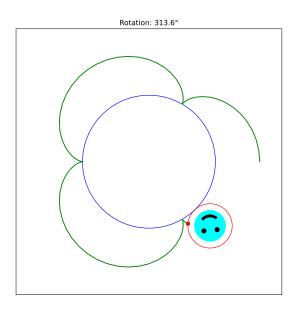


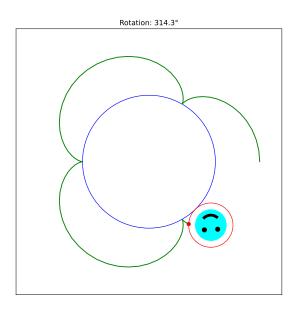


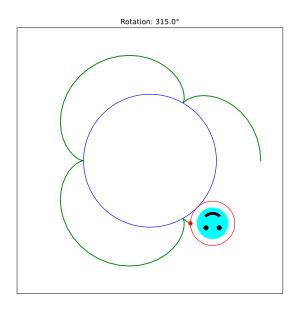


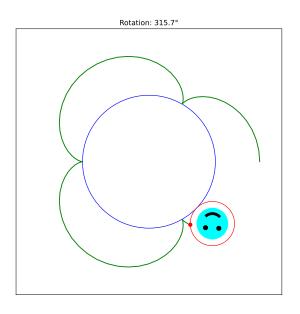


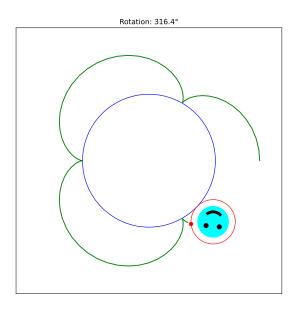


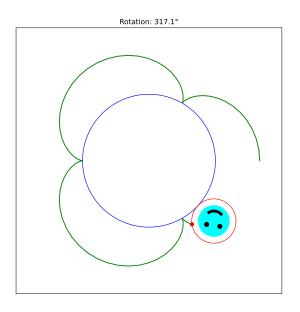


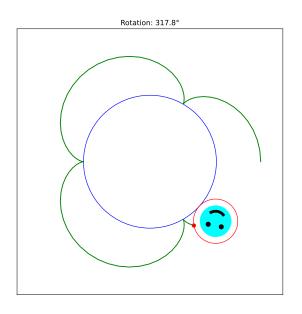


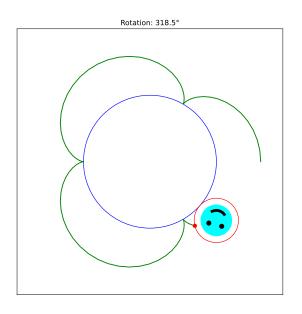


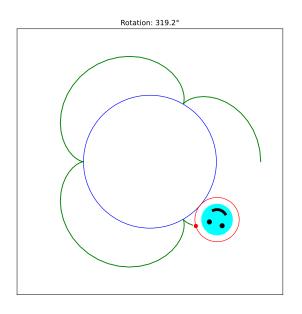


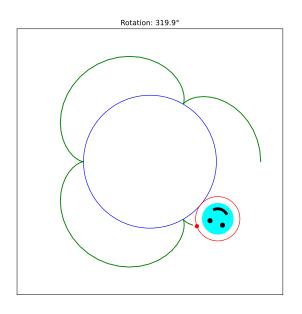


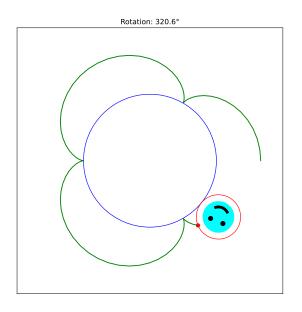


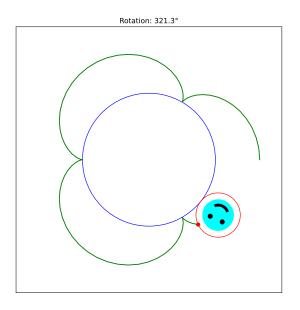


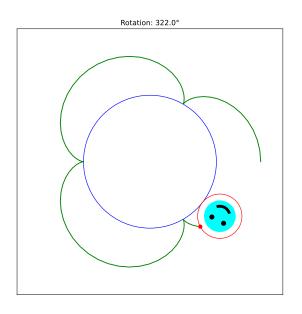


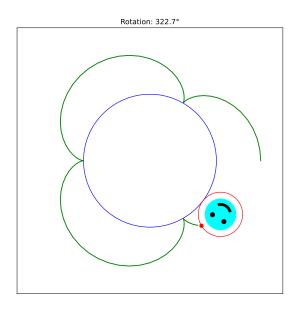


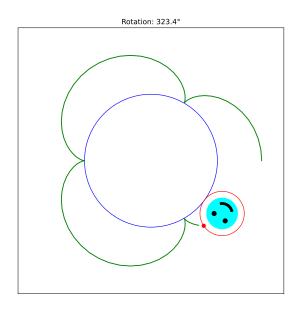


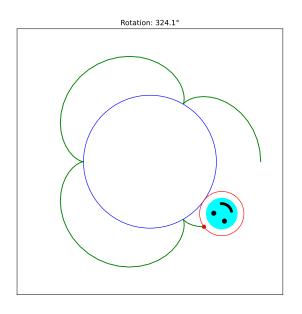


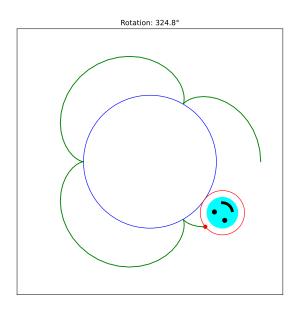


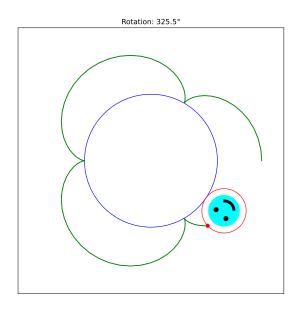


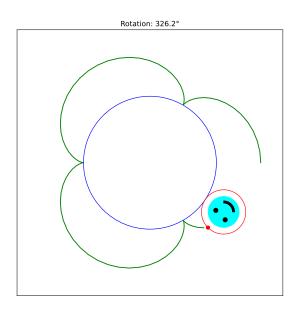


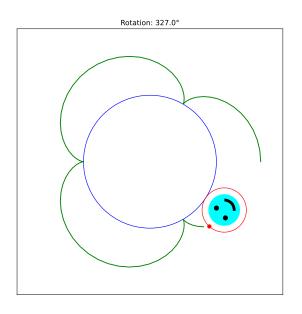


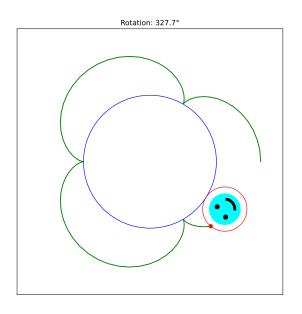


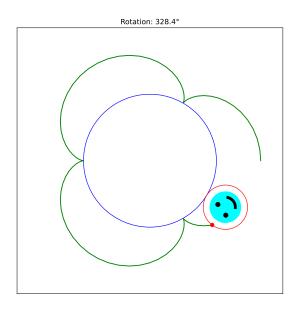


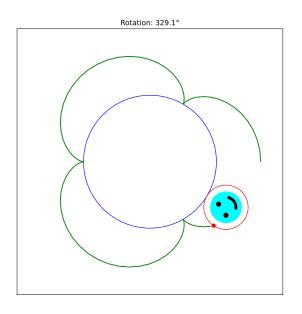


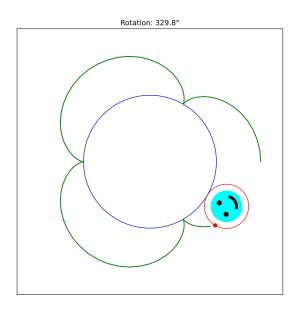


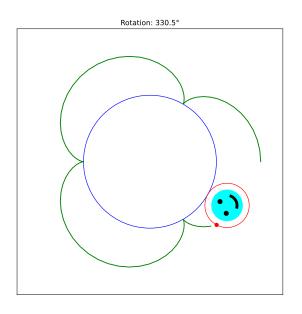


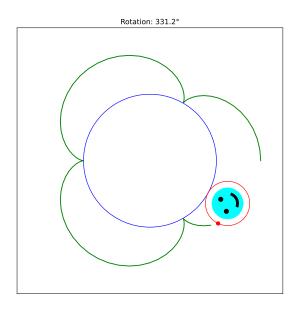


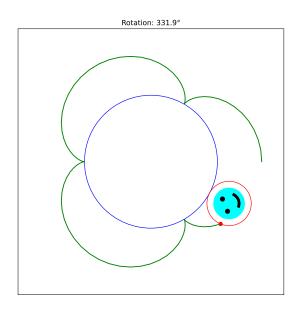


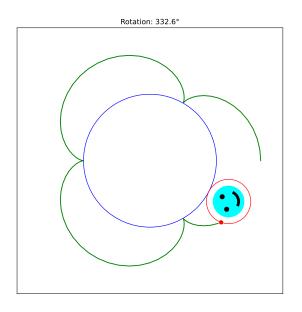


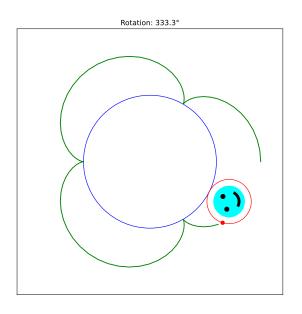


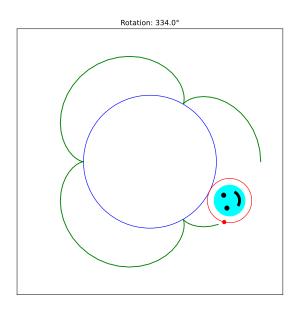


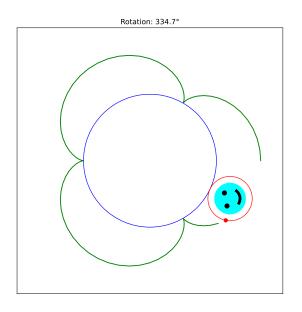


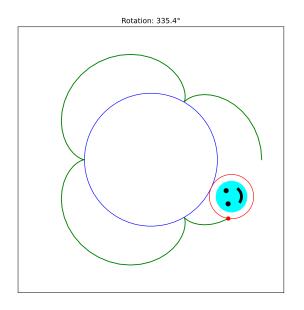


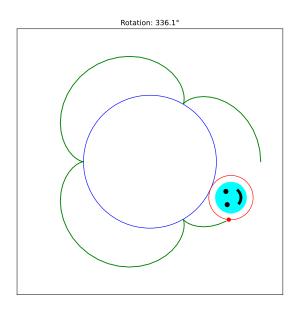


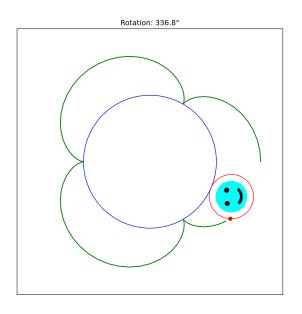


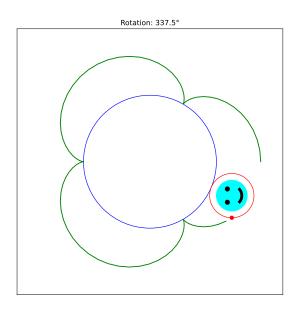


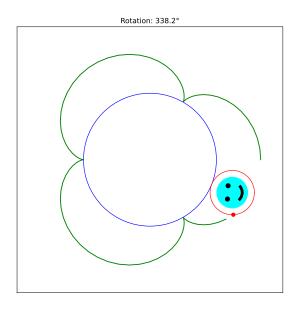


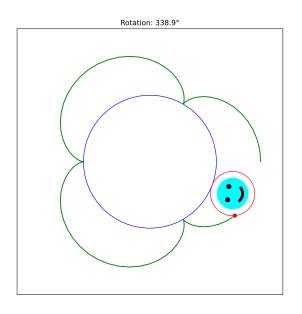


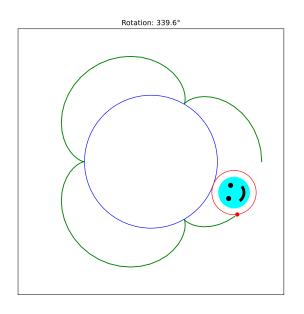


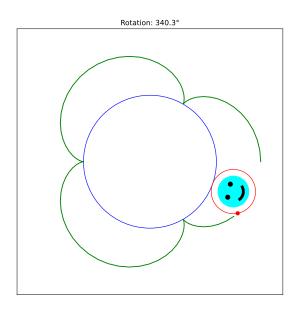


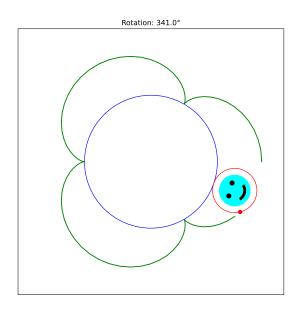


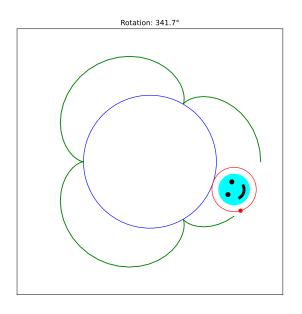


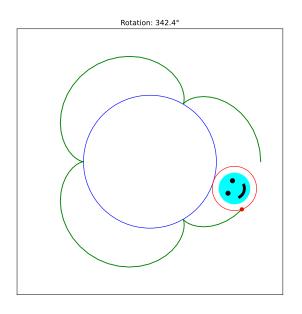


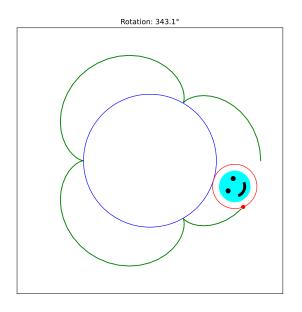


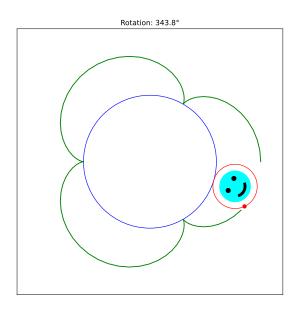


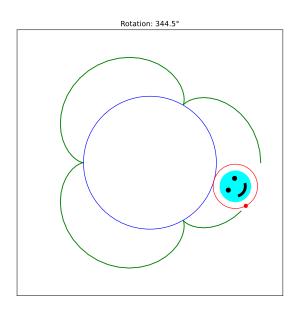


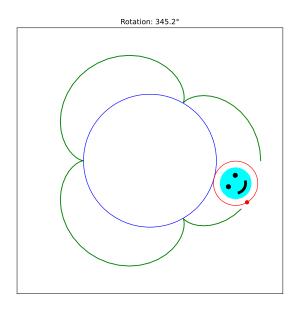


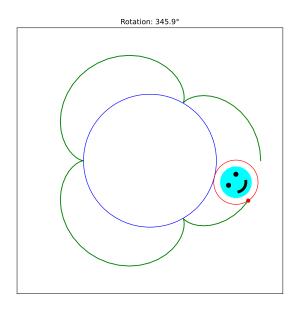


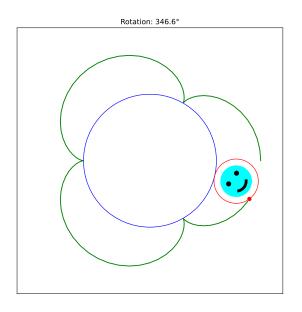


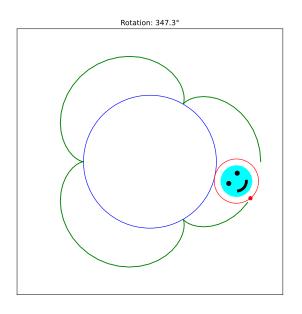


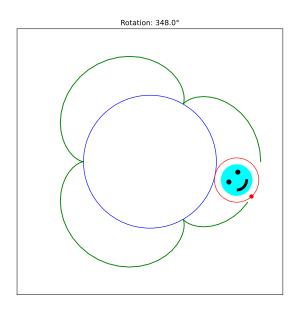


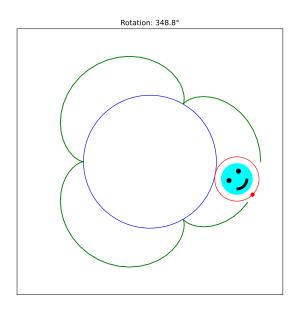


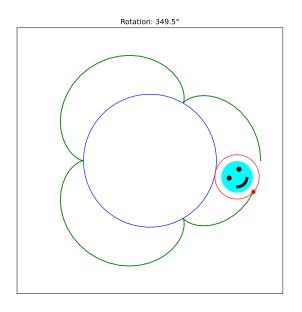


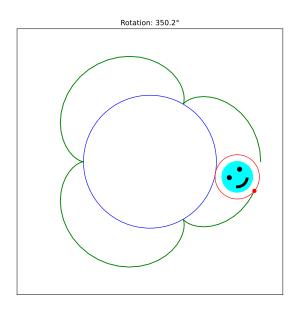


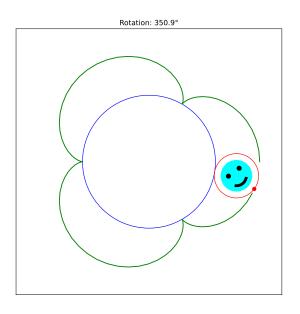


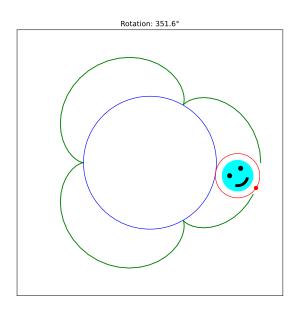


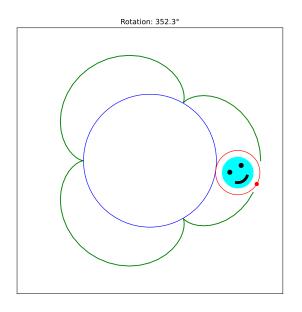


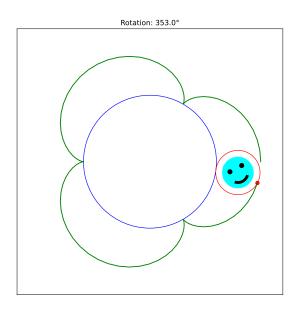


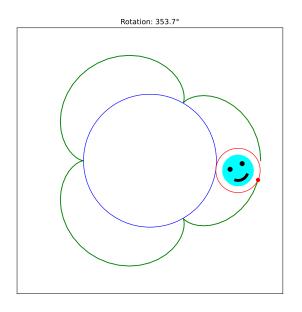


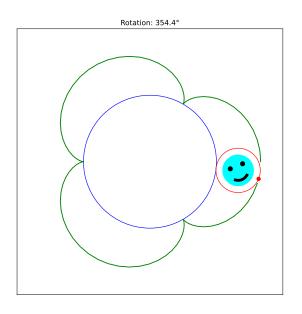


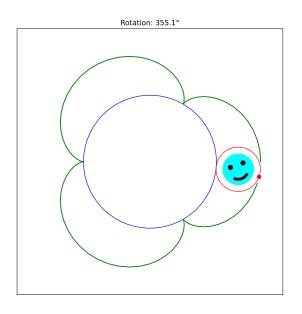


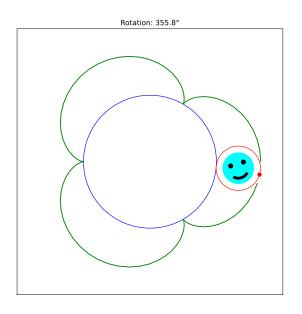


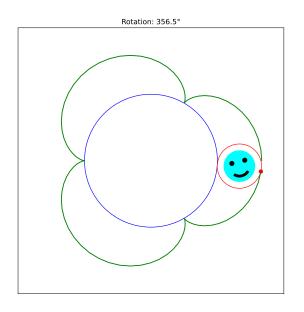


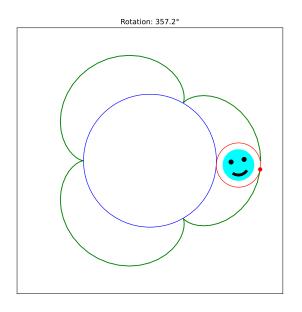


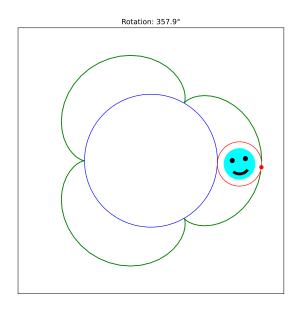


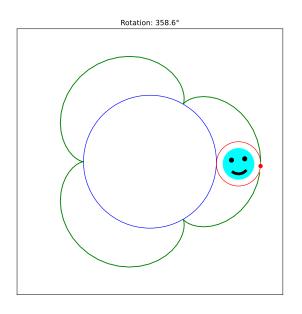


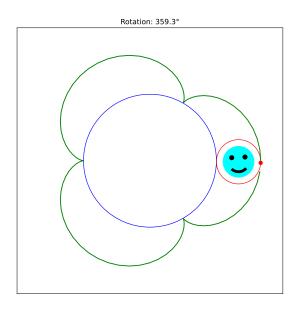


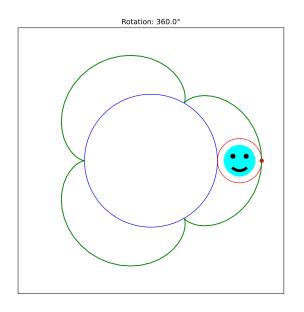




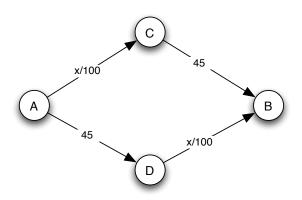


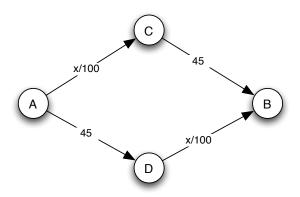




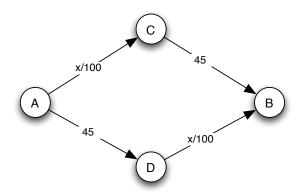


Braess Paradox

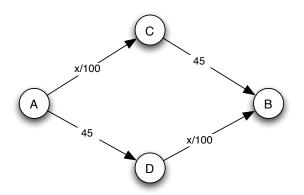




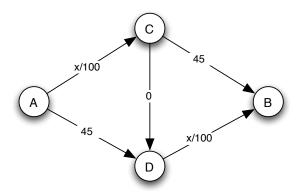
ullet A highway network, with each edge labeled by its travel time (in minutes) when there are x cars using it

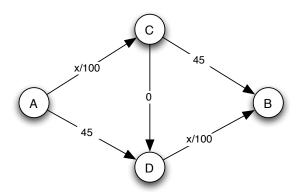


- A highway network, with each edge labeled by its travel time (in minutes) when there are x cars using it
- Suppose there are 4000 cars need to get from A to B

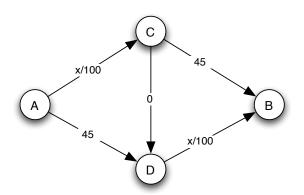


- A highway network, with each edge labeled by its travel time (in minutes) when there are x cars using it
- Suppose there are 4000 cars need to get from A to B
- They divide evenly over the two routes at equilibrium; the travel time is 45 + 2000/100 = 65 mins

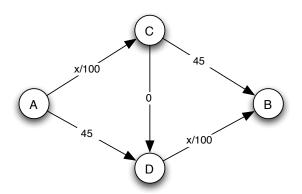




• Now a very fast edge is added from C to D to the previous highway network

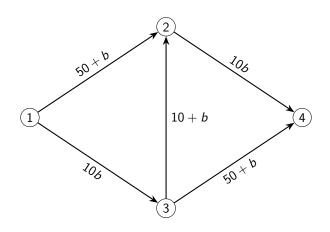


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- At equilibrium, every user uses the route through C and D

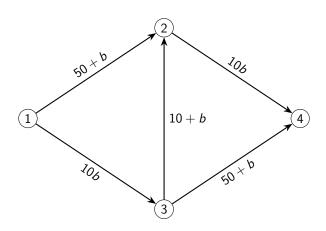


- Now a very fast edge is added from C to D to the previous highway network
- At equilibrium, every user uses the route through C and D
- As a result, the travel time is 4000/10 + 0 + 4000/100 = 80 mins!

[Braess et al., 2005] Example

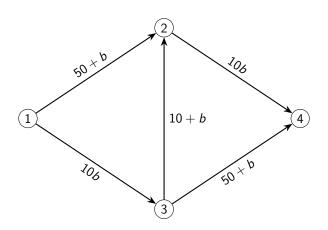


[Braess et al., 2005] Example



• Assume a flow of 6 units (e.g., 6000 vehicles) must travel from 1 to 4

[Braess et al., 2005] Example



- Assume a flow of 6 units (e.g., 6000 vehicles) must travel from 1 to 4
- Three paths exist: $B_1 = 124$, $B_2 = 1324$, $B_3 = 134$

• Case 1: All Paths Open. Split the flow equally (2 units each):

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$$b_{12} = 2$$
, $b_{24} = 4$, $b_{34} = 2$, $b_{13} = 4$, $b_{32} = 2$
 $d_{12} = 52$, $d_{24} = 40$, $d_{34} = 52$, $d_{13} = 40$, $d_{32} = 12$
 $L(B_1) = L(B_2) = L(B_3) = 92$

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All paths are critical (maximal length). This distribution is stable: Switching paths increases load and travel time beyond 92.

• Case 1: All Paths Open. Split the flow equally (2 units each):

$$b_{12} = 2$$
, $b_{24} = 4$, $b_{34} = 2$, $b_{13} = 4$, $b_{32} = 2$
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 $d_{12} = 53, \ d_{24} = 30, \ d_{13} = 30, \ d_{34} = 53$
 $L(B_1) = L(B_3) = 83$

[Braess et al., 2005] Example: Cont'd

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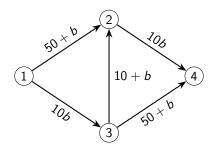
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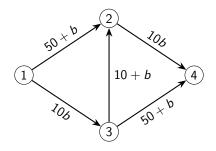
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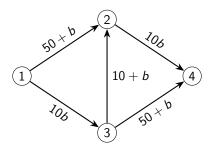
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Despite losing a connection, all paths are shorter!

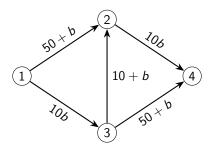




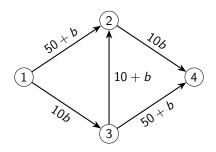
• Recall that $B_1 = 124$, $B_2 = 1324$, $B_3 = 134$



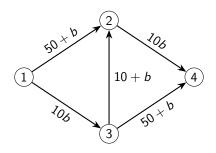
- Recall that $B_1 = 124$, $B_2 = 1324$, $B_3 = 134$
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- B_2 users clog k_{13} and k_{24} (factor 10 in load-time relation), worsening times for B_1 and B_3 beyond 83; Yet B_2 remains shortest, attracting more traffic and degrading the system for all.

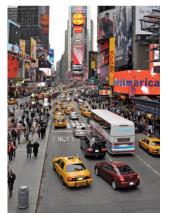
• The mechanical analogy: the spring paradox

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- Road closure in NY: Times and Herald Squares pedestrian plaza (2009 —)





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- Cheonggyecheon restoration project (2003 —)
 - Replaced a six lane highway with a five mile long park, traffic flow improved

The Social Cost of Traffic at Equilibrium

A traffic network is a directed graph with

• Nodes: Start and destination points for drivers.

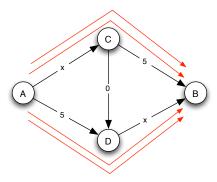
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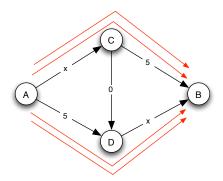
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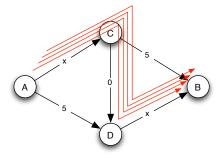
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- Nash Equilibrium: No driver can reduce their travel time by switching paths, given others' choices.



(a) The social optimum.







(b) The Nash equilibrium.

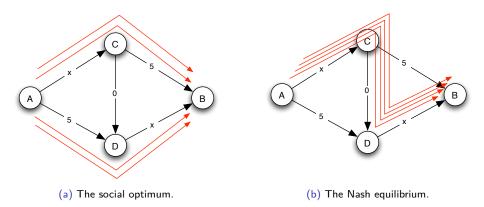


Figure: A version of Braess's Paradox: In the socially optimal traffic pattern, the social cost is 28, while in the unique Nash equilibrium, the social cost is 32.

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$$\mathsf{Energy}(e) = \mathit{T}_e(1) + \mathit{T}_e(2) + \cdots + \mathit{T}_e(x)$$

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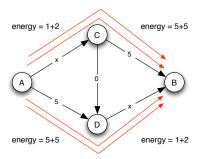
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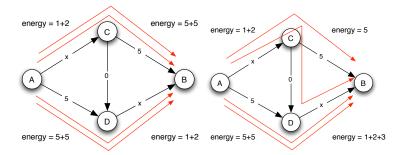
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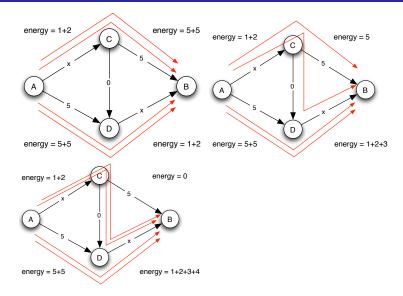
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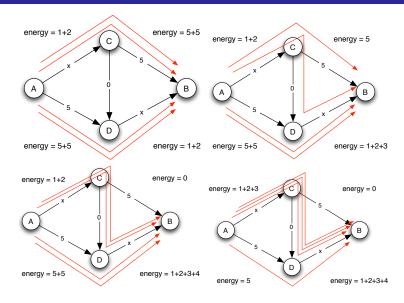
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- Social cost can increase or decrease with best-response steps (e.g., from 28 to 32 in the Braess example), but potential energy strictly decreases, serving as a progress measure.









Analyzing Dynamics

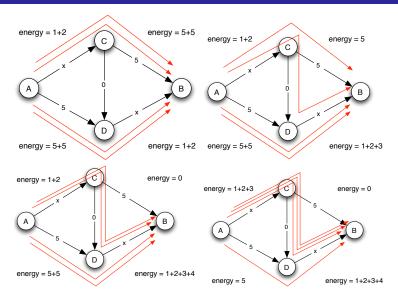


Figure: Steps of best-response dynamics with potential energy changes.

For any edge e,

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- With finite patterns, dynamics must stop at an equilibrium.

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Energy(e) =
$$a_e(1 + 2 + \dots + x) + b_e x = \frac{a_e x(x+1)}{2} + b_e x$$

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Proof:

$$\frac{a_e x(x+1)}{2} + b_e x \geqslant \frac{1}{2}(a_e x^2 + b_e x) \quad \text{and} \quad \leqslant a_e x^2 + b_e x$$

• For a pattern Z,

$$\frac{1}{2}$$
 · Social-Cost(Z) \leq Energy(Z) \leq Social-Cost(Z)

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$$\frac{1}{2} \cdot \mathsf{Social\text{-}Cost}(\mathit{Z}) \leqslant \mathsf{Energy}(\mathit{Z}) \leqslant \mathsf{Social\text{-}Cost}(\mathit{Z})$$

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- Implications: Network design and tolls can mitigate inefficiencies.

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