

Review of Prerequisite Mathematics

Introduction to Elementary Optimization

Definition (The Jacobian)

Let V be open in \mathbb{R}^n , $\mathbf{x} \in V$, and $g_i : V \rightarrow \mathbb{R}$, $i = 1, 2, \dots, m$ be C^1 on V . The Jacobian of $\mathbf{g}(\mathbf{x}) : \mathbb{R}^n \rightarrow \mathbb{R}^m$ is defined as

$$D\mathbf{g}(\mathbf{x}) = \begin{pmatrix} \frac{\partial g_1}{\partial x_1} & \frac{\partial g_1}{\partial x_2} & \cdots & \frac{\partial g_1}{\partial x_n} \\ \frac{\partial g_2}{\partial x_1} & \frac{\partial g_2}{\partial x_2} & \cdots & \frac{\partial g_2}{\partial x_n} \\ \vdots & \vdots & \ddots & \vdots \\ \frac{\partial g_m}{\partial x_1} & \frac{\partial g_m}{\partial x_2} & \cdots & \frac{\partial g_m}{\partial x_n} \end{pmatrix} (\mathbf{x})$$

Theorem (The Chain Rule)

Suppose that \mathbf{f} and \mathbf{g} are vector functions. If \mathbf{g} is differentiable at \mathbf{a} and \mathbf{f} is differentiable at $\mathbf{g}(\mathbf{a})$, then $\mathbf{f} \circ \mathbf{g}$ is differentiable at \mathbf{a} and

$$D(\mathbf{f} \circ \mathbf{g})(\mathbf{a}) = D\mathbf{f}(\mathbf{g}(\mathbf{a})) D\mathbf{g}(\mathbf{a})$$

More explicitly, if f is a differentiable function of x_1, x_2, \dots, x_n , and each x_j is a differentiable function of t_1, t_2, \dots, t_m , $n, m \geq 1$; then f is a differentiable function of t_1, t_2, \dots, t_m with

$$\frac{\partial f}{\partial t_i} = \frac{\partial f}{\partial x_1} \frac{\partial x_1}{\partial t_i} + \frac{\partial f}{\partial x_2} \frac{\partial x_2}{\partial t_i} + \cdots + \frac{\partial f}{\partial x_n} \frac{\partial x_n}{\partial t_i}$$

Example

Let $w = f(xz, yz)$, where f is a differentiable function. Prove that

$$x \frac{\partial w}{\partial x} + y \frac{\partial w}{\partial y} = z \frac{\partial w}{\partial z}.$$

Solution I

Write $u(x, y, z) = xz$ and $v(x, y, z) = yz$ so that $w(x, y, z) = f(u(x, y, z), v(x, y, z))$. By the chain rule,

$$\begin{aligned}\frac{\partial w}{\partial x}(x, y, z) &= \frac{\partial}{\partial x}[f(u(x, y, z), v(x, y, z))] \\&= \frac{\partial f}{\partial u}(u(x, y, z), v(x, y, z)) \frac{\partial u}{\partial x}(x, y, z) + \frac{\partial f}{\partial v}(u(x, y, z), v(x, y, z)) \frac{\partial v}{\partial x}(x, y, z) \\&= z \frac{\partial f}{\partial u}(xz, yz)\end{aligned}$$

$$\begin{aligned}\frac{\partial w}{\partial y}(x, y, z) &= \frac{\partial}{\partial y}[f(u(x, y, z), v(x, y, z))] \\&= \frac{\partial f}{\partial u}(u(x, y, z), v(x, y, z)) \frac{\partial u}{\partial y}(x, y, z) + \frac{\partial f}{\partial v}(u(x, y, z), v(x, y, z)) \frac{\partial v}{\partial y}(x, y, z) \\&= z \frac{\partial f}{\partial v}(xz, yz)\end{aligned}$$

Solution II

$$\begin{aligned}\frac{\partial w}{\partial z}(x, y, z) &= \frac{\partial}{\partial z}[f(u(x, y, z), v(x, y, z))] \\&= \frac{\partial f}{\partial u}(u(x, y, z), v(x, y, z)) \frac{\partial u}{\partial z}(x, y, z) + \frac{\partial f}{\partial v}(u(x, y, z), v(x, y, z)) \frac{\partial v}{\partial z}(x, y, z) \\&= x \frac{\partial f}{\partial u}(xz, yz) + y \frac{\partial f}{\partial v}(xz, yz)\end{aligned}$$

So

$$\begin{aligned}\frac{\partial w}{\partial x} + y \frac{\partial w}{\partial y} &= xz \frac{\partial f}{\partial u}(xz, yz) + yz \frac{\partial f}{\partial v}(xz, yz) \\&= z \left[x \frac{\partial f}{\partial u}(xz, yz) + y \frac{\partial f}{\partial v}(xz, yz) \right] = z \frac{\partial w}{\partial z}\end{aligned}$$

Unconstrained Optimization Problems

Theorem

Given $S \subseteq \mathbb{R}^n$ and continuous $f : S \rightarrow \mathbb{R}$; if S is compact, then

$$M = \sup \{f(\mathbf{x}) : \mathbf{x} \in S\} \quad \text{and} \quad m = \inf \{f(\mathbf{x}) : \mathbf{x} \in S\}$$

are finite real numbers. Moreover, there exists points $\mathbf{x}_M, \mathbf{x}_m \in S$ such that $M = f(\mathbf{x}_M)$ and $m = f(\mathbf{x}_m)$.

Definition

Given $S \subseteq \mathbb{R}^n$, $f : S \rightarrow \mathbb{R}$ and $B(\mathbf{x}, h) \equiv \{\mathbf{y} \in \mathbb{R}^n \mid |\mathbf{y} - \mathbf{x}| < h\}$, f achieves

- global maximum $f(\mathbf{x}_M)$ at $\mathbf{x}_M \in S$: $f(\mathbf{x}_M) \geq f(\mathbf{x})$, $\forall \mathbf{x} \in S$.
- global minimum $f(\mathbf{x}_m)$ at $\mathbf{x}_m \in S$: $f(\mathbf{x}_m) \leq f(\mathbf{x})$, $\forall \mathbf{x} \in S$.
- local maximum $f(\mathbf{x}_0)$ at $\mathbf{x}_0 \in S$: $\exists h_0 > 0$ s.t. $f(\mathbf{x}_0) \geq f(\mathbf{x})$, $\forall \mathbf{x} \in B(\mathbf{x}_0, h_0) \cap S$.
- local minimum $f(\mathbf{x}_1)$ at $\mathbf{x}_1 \in S$: $\exists h_1 > 0$ s.t. $f(\mathbf{x}_1) \leq f(\mathbf{x})$, $\forall \mathbf{x} \in B(\mathbf{x}_1, h_1) \cap S$.

Theorem (necessary conditions for extremum)

Given $S \subseteq \mathbb{R}^n$ and differentiable $f : S \rightarrow \mathbb{R}$, if f achieves extremum at an interior $\mathbf{c} \in S$, then $\nabla f(\mathbf{c}) = \mathbf{0}$.

Proof

If $\mathbf{c} = (c_1, c_2, \dots, c_n)$, let

$$g_j(t) \equiv f(c_1, c_2, \dots, c_{j-1}, t, c_{j+1}, \dots, c_n), \quad j = 1, 2, \dots, n$$

For f achieves extremum at \mathbf{c} , $f(\mathbf{c}) = g_j(c_j)$, g_j achieves extremum at $c_j \implies g'_j(t)|_{t=c_j} = 0 \implies D_j f(\mathbf{c}) = 0 \forall j$, so $\nabla f(\mathbf{c}) = \mathbf{0}$.

Theorem

Given $S \subseteq \mathbb{R}^n$, if $f : S \rightarrow \mathbb{R}$ achieves extremum at $\mathbf{c} \in S$, then \mathbf{c} can possibly be a

- critical point: $\nabla f(\mathbf{c}) = \mathbf{0}$.
- singular point: f is non-differentiable at \mathbf{c} .
- boundary point of S .

Definition (Hessian Matrix)

Given $S \subseteq \mathbb{R}^n$, an interior point \mathbf{c} of S , and a differentiable function $f : S \rightarrow \mathbb{R}$,

$$\mathbf{H}(f, \mathbf{c}) = \begin{pmatrix} f_{11} & f_{12} & \cdots & f_{1n} \\ f_{21} & f_{22} & \cdots & f_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ f_{n1} & f_{n2} & \cdots & f_{nn} \end{pmatrix}, \quad f_{ij} = \frac{\partial^2 f}{\partial x_j \partial x_i}(\mathbf{c}), \quad i, j = 1, 2, \dots, n.$$

Definition (Matrix Positive/Negative Definiteness)

Given an $n \times n$ real symmetric matrix \mathbf{A} . For any $\mathbf{v} \in \mathbb{R}^n \neq \mathbf{0}$, \mathbf{A} is

- positive-definite: $\mathbf{v}\mathbf{A}\mathbf{v}^\top > 0$
- positive-semidefinite: $\mathbf{v}\mathbf{A}\mathbf{v}^\top \geq 0$
- negative-definite: $\mathbf{v}\mathbf{A}\mathbf{v}^\top < 0$
- negative-semidefinite: $\mathbf{v}\mathbf{A}\mathbf{v}^\top \leq 0$

Theorem

Given a real symmetric matrix \mathbf{A} .

- \mathbf{A} is positive-definite \iff all eigenvalues of \mathbf{A} are positive.
- \mathbf{A} is positive-semidefinite \iff all eigenvalues of \mathbf{A} are nonnegative.

Definition (Minor)

Given an $n \times n$ matrix $\mathbf{A} = \begin{pmatrix} a_{11} & a_{12} & \cdots & a_{1n} \\ a_{21} & a_{22} & \cdots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & \cdots & a_{nn} \end{pmatrix}$ and minor

$$\mathbf{A} \begin{pmatrix} i_1, i_2, \dots, i_k \\ j_1, j_2, \dots, j_k \end{pmatrix} = \begin{vmatrix} a_{i_1 j_1} & a_{i_1 j_2} & \cdots & a_{i_1 j_k} \\ a_{i_2 j_1} & a_{i_2 j_2} & \cdots & a_{i_2 j_k} \\ \vdots & \vdots & \ddots & \vdots \\ a_{i_k j_1} & a_{i_k j_2} & \cdots & a_{i_k j_k} \end{vmatrix}, \quad 1 \leq k \leq n,$$

$$1 \leq i_1 < i_2 < \cdots < i_k \leq n, \quad 1 \leq j_1 < j_2 < \cdots < j_k \leq n.$$

- $\Delta_k \equiv \mathbf{A} \begin{pmatrix} i_1, i_2, \dots, i_k \\ i_1, i_2, \dots, i_k \end{pmatrix}$ is the k -th order principal minor of A .
- $M_k \equiv \mathbf{A} \begin{pmatrix} 1, 2, \dots, k \\ 1, 2, \dots, k \end{pmatrix}$ is the k -th order leading principal minor of A .

Theorem (Criteria for Matrix Positive/Negative Definiteness)

Given an $n \times n$ real symmetric matrix \mathbf{A} , then $\forall k \leq n$, \mathbf{A} is

- positive-definite $\iff M_k > 0$
- negative-definite $\iff (-1)^k M_k > 0$
- positive-semidefinite $\iff \Delta_k \geq 0$
- negative-semidefinite $\iff (-1)^k \Delta_k \geq 0$

Equalities Constrained Optimization: The Lagrange Multipliers Method

Theorem

Given an open set $S \subseteq \mathbb{R}^n$, differentiable functions $f : S \rightarrow \mathbb{R}$ and $g_j : S \rightarrow \mathbb{R}$, $j = 1, 2, \dots, m$, $m < n$, and $X_0 = \{\mathbf{x} \in S \mid g_j(\mathbf{x}) = 0, j = 1, 2, \dots, m\}$. If f has an extremum at $\mathbf{x}_0 \in S \cap X_0$ and $\det(D_i g_j(\mathbf{x}_0)) \neq 0$, then

$$\exists \lambda_1, \lambda_2, \dots, \lambda_m \quad \text{s.t.} \quad D_i f(\mathbf{x}_0) + \sum_{j=1}^m \lambda_j D_i g_j(\mathbf{x}_0) = 0, \quad i = 1, 2, \dots, n$$

Remark

Let $\mathcal{L} \equiv f + \sum_{j=1}^m \lambda_j g_j$, the sufficient condition can be rewritten as

$$\begin{aligned} D_i \mathcal{L}(\mathbf{x}_0) &= 0, & i &= 1, 2, \dots, n \\ g_j(\mathbf{x}_0) &= 0, & j &= 1, 2, \dots, m \end{aligned}$$

Example

Find the maximum and minimum values of $x^2 - 10x - y^2$ on $x^2 + 4y^2 = 16$.

Solution

Let $\mathcal{L} = x^2 - 10x - y^2 + \lambda(x^2 + 4y^2 - 16)$, then

$$\frac{\partial \mathcal{L}}{\partial x} = 2x - 10 + 2\lambda x = 0 \implies x - 5 + \lambda x = 0 \quad (1)$$

$$\frac{\partial \mathcal{L}}{\partial y} = -2y + 8\lambda y = 0 \implies -y + 4\lambda y = 0 \quad (2)$$

$$x^2 + 4y^2 - 16 = 0 \quad (3)$$

From (2) $(1 - 4\lambda)y = 0$, so $y = 0 \vee \lambda = \frac{1}{4}$. If $y = 0$, from (3) $x = \pm 4$; if $\lambda = \frac{1}{4}$, from (1) $(1 + \lambda)x = 5 \implies x = 4$, substituting into (3) gives $y = 0$. Therefore, the extremum points are $(x, y) = (4, 0), (-4, 0)$; $x^2 - 10x - y^2$ has a maximum value of 56 (at $(x, y) = (-4, 0)$), and a minimum value of -24 (at $(x, y) = (4, 0)$).

Example

Find the maximum and minimum values of $f(x, y, z) = (x + z) e^y$ on $x^2 + y^2 + z^2 = 6$.

Solution

Let $\mathcal{L} = (x + z) e^y + \lambda (x^2 + y^2 + z^2 - 6)$, then

$$\frac{\partial \mathcal{L}}{\partial x} = e^y + 2\lambda x = 0 \quad (4)$$

$$\frac{\partial \mathcal{L}}{\partial y} = (x + z) e^y + 2\lambda y = 0 \quad (5)$$

$$\frac{\partial \mathcal{L}}{\partial z} = e^y + 2\lambda z = 0 \quad (6)$$

$$x^2 + y^2 + z^2 - 6 = 0 \quad (7)$$

From (4), (6) $2\lambda(x - z) = 0$, so $\lambda = 0 \vee x = z$. If $\lambda = 0$, then from (4) $e^y = 0$ which is impossible, so $x = z$. From (4) $e^y = -2\lambda x$, substituting into (5) $2x(-2\lambda x) + 2\lambda y = 0 \implies y = 2x^2$, substituting into (7) gives $x^2 + 4x^4 + x^2 = 6 \implies (4x^2 + 6)(x^2 - 1) = 0 \implies x = \pm 1$. Therefore, the extremum points are $(x, y, z) = (1, 2, 1), (-1, 2, -1)$; $(x + z) e^y$ has a maximum value of $2e^2$ (at $(x, y, z) = (1, 2, 1)$), and a minimum value of $-2e^2$ (at $(x, y, z) = (-1, 2, -1)$).

Example

If L is the curve of intersection of $z^2 = x^2 + y^2$ and $x - 2z = 3$, find the point on L that is closest to the origin and the shortest distance.

Solution

The objective is $x^2 + y^2 + z^2$ with constraints $x^2 + y^2 - z^2 = 0$ and $x - 2z - 3 = 0$. Let $\mathcal{L} = x^2 + y^2 + z^2 + \lambda_1 (x^2 + y^2 - z^2) + \lambda_2 (x - 2z - 3)$, then

$$\frac{\partial \mathcal{L}}{\partial x} = 2x + 2\lambda_1 x + \lambda_2 = 0 \implies 2(1 + \lambda_1)x + \lambda_2 = 0 \quad (8)$$

$$\frac{\partial \mathcal{L}}{\partial y} = 2y + 2\lambda_1 y = 0 \implies (1 + \lambda_1)y = 0 \quad (9)$$

$$\frac{\partial \mathcal{L}}{\partial z} = 2z - 2\lambda_1 z - 2\lambda_2 = 0 \implies (1 - \lambda_1)z - \lambda_2 = 0 \quad (10)$$

$$x^2 + y^2 - z^2 = 0 \quad (11)$$

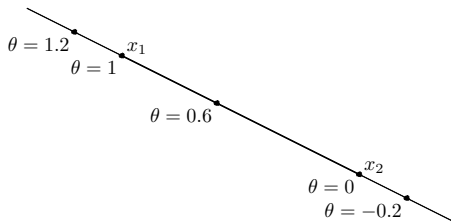
$$x - 2z - 3 = 0 \quad (12)$$

From (9) $(1 + \lambda_1)y = 0$, so $y = 0 \vee \lambda_1 = -1$. If $y = 0$, from (11) $x^2 = z^2 \implies x = \pm z$. If $x = z$, from (12) $x = z = -3$. If $x = -z$, from (12) $x = 1, z = -1$; if $\lambda_1 = -1$, from (8) $\lambda_2 = 0$, from (10) $z = 0$, substituting into (11) gives $x = y = 0$, which contradicts (12). Therefore, the extremum points are $(x, y, z) = (-3, 0, -3), (1, 0, -1)$; optimizer: $(1, 0, -1)$.

Introduction to Convex Programming

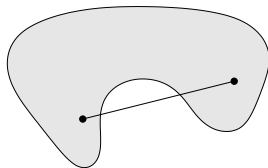
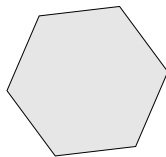
Affine Set

- **line** through x_1, x_2 : all points of the form $x = \theta x_1 + (1 - \theta)x_2, \theta \in \mathbb{R}$
- **affine set** contains the line through any two distinct points in the set
- e.g. solution set of linear equations $\{x \mid Ax = b\}$; every affine set can be expressed as solution set of system of linear equations



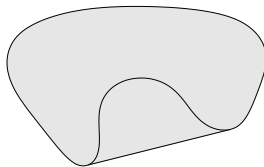
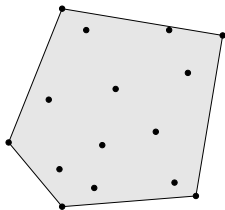
Convex Set

- **line segment** through x_1, x_2 : all points of the form $x = \theta x_1 + (1 - \theta)x_2, 0 \leq \theta \leq 1$
- **convex set** contains the line segment between any two distinct points in the set:
$$x_1, x_2 \in S \implies \forall 0 \leq \theta \leq 1, \theta x_1 + (1 - \theta)x_2 \in S$$



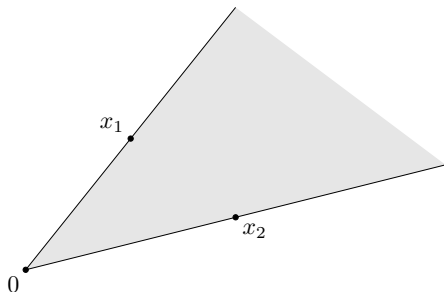
Convex Combination, Convex Hull

- **convex combination** of x_1, x_2, \dots, x_k : any point x of the form $x = \theta_1 x_1 + \theta_2 x_2 + \dots + \theta_k x_k$ with $\theta_1 + \theta_2 + \dots + \theta_k = 1, \theta_i \geq 0$
- **convex hull** $\text{conv } S$: sets of all convex combinations of points in S



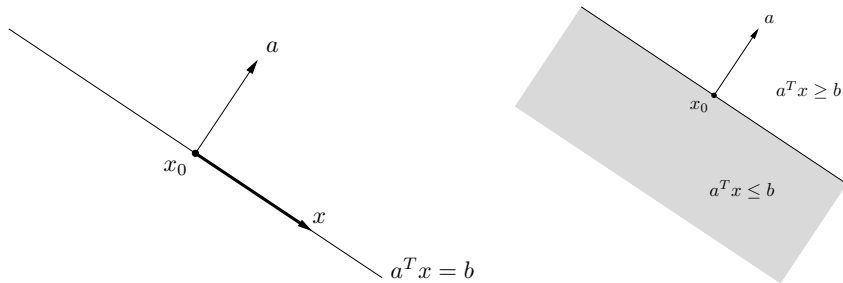
Convex Cone

- **conic (nonnegative) combination** of x_1 and x_2 : any point x of the form $x = \theta_1 x_1 + \theta_2 x_2$ with $\theta_i \geq 0$
- **convex cone** set that contains all conic combinations of points in the set



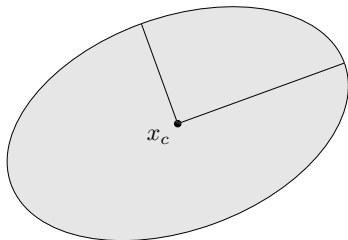
Hyperplane, Halfspace

- **hyperplane**: set of the form $\{x \mid a^\top x = b\}$ with $a \neq 0$
 - **halfspace**: set of the form $\{x \mid a^\top x \leq b\}$ with $a \neq 0$
 - a : normal vector
- hyperplanes are affine and convex, halfspaces are convex



Euclidean Ball, Ellipsoid

- **(Euclidean) ball** with center x_c and radius r :
 $B(x_c, r) = \{x \mid \|x - x_c\|_2 \leq r\} = \{x_c + r u \mid \|u\|_2 \leq 1\}$
- **ellipsoid**: set of the form $\{x \mid (x - x_c)^\top P^{-1}(x - x_c) \leq 1\}$ with $P \in \mathbf{S}_{++}^n$ (P symmetric positive definite), or $\{x_c + A u \mid \|u\|_2 \leq 1\}$ with nonsingular A



Norm Ball, Norm Cone

- **norm**: a function $\|\cdot\|$ that satisfies
 - $\|x\| \geq 0$; $\|x\| = 0 \iff x = 0$
 - $\|tx\| = |t|\|x\|$, $\forall t \in \mathbb{R}$
 - $\|x + y\| \leq \|x\| + \|y\|$
- **norm ball** with center x_c and radius r : $\{x \mid \|x - x_c\| \leq r\}$
- **norm cone**: $\{(x, t) \mid \|x\| \leq t\}$
- norm balls and norm cones are convex
- notation for different norms: $\|\cdot\|_2$, $\|\cdot\|_{\text{symb}}$

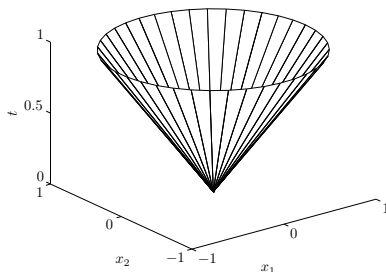
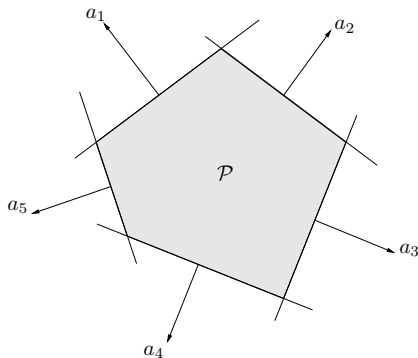


Figure: Boundary of second-order cone in \mathbb{R}^3 , $\{(x_1, x_2, t) \mid (x_1^2 + x_2^2)^{\frac{1}{2}} \leq t\}$.

Polyhedra

- **polyhedron**: solution set of finitely many linear equalities and inequalities $\{x \mid Ax \preceq b, Cx = d\}$, where $A \in \mathbb{R}^{m \times n}$, $C \in \mathbb{R}^{p \times n}$, \preceq is componentwise inequality
- intersection of finite number of halfspaces and hyperplanes



Positive Semidefinite Cone

- S^n : set of symmetric $n \times n$ matrices
- $S_+^n = \{X \in S^n \mid X \succcurlyeq 0\}$: set of positive semidefinite (symmetric) $n \times n$ matrices; $X \in S_+^n \iff z^\top X z \geq 0 \forall z$; a convex cone, the **positive semidefinite cone**; Below: $\begin{pmatrix} x & y \\ y & z \end{pmatrix} \in S_+^2$
- $S_{++}^n = \{X \in S^n \mid X \succ 0\}$: set of positive definite (symmetric) $n \times n$ matrices

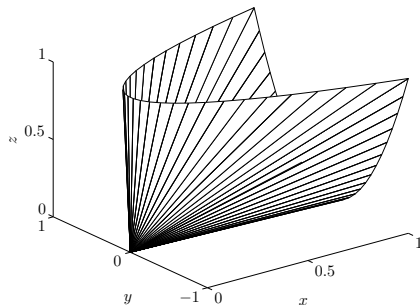


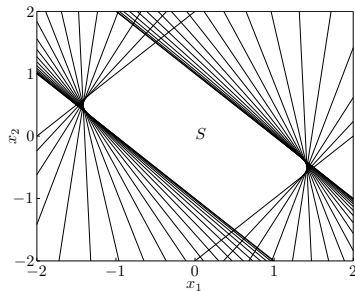
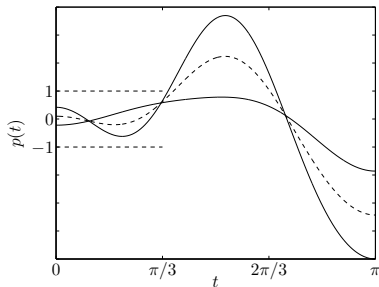
Figure: $x \geq 0, z \geq 0, xz \geq y^2$.

Showing a Set is Convex

- apply definition: $x_1, x_2 \in S \implies \theta x_1 + (1 - \theta)x_2 \in S, \forall 0 \leq \theta \leq 1$
recommended only for simple sets
- use convex functions (later)
- show that the set is obtained from other simple convex sets (e.g. hyperplanes, halfspaces, norm balls) by operations that preserve convexity:
 - intersection
 - affine mapping
 - perspective mapping
 - linear-fractional mapping
- mostly using last two

Intersection

- intersection of (any number of) convex sets is convex
- e.g. $S = \left\{ x \in \mathbb{R}^m \mid |p(t)| \leq 1 \ \forall |t| \leq \frac{\pi}{3} \right\}$, $p(t) = \sum_{k=1}^m x_k \cos kt$
is convex by $S = \bigcap_{|t| \leq \frac{\pi}{3}} \{x \mid |p(t)| \leq 1\}$; intersection of convex slabs. Below:
 $m = 2$.



Affine Mappings

- suppose $f : \mathbb{R}^n \rightarrow \mathbb{R}^m$ is **affine**, i.e.

$$f(x) = Ax + b \quad \text{with } A \in \mathbb{R}^{m \times n}, b \in \mathbb{R}^m$$

- the **image** of a convex set under f is convex:

$$S \subseteq \mathbb{R}^n \text{ convex} \implies f(S) = \{f(x) \mid x \in S\} \text{ convex}$$

- the **inverse image** of a convex set under f is convex:

$$C \subseteq \mathbb{R}^m \text{ convex} \implies f^{-1}(C) = \{x \in \mathbb{R}^n \mid f(x) \in C\} \text{ convex}$$

- e.g. scaling $aS + b = \{ax + b \mid x \in S\}$, $a, b \in \mathbb{R}$ is convex
- e.g. projection $\text{proj}_x(S) = \{x \mid (x, y) \in S\}$ is convex

Perspective and Linear-Fractional Function

- **perspective function** $p : \mathbb{R}^{n+1} \rightarrow \mathbb{R}^n$:

$$p(x, t) = \frac{x}{t} \quad \text{dom } p = \{(x, t) \mid t > 0\}$$

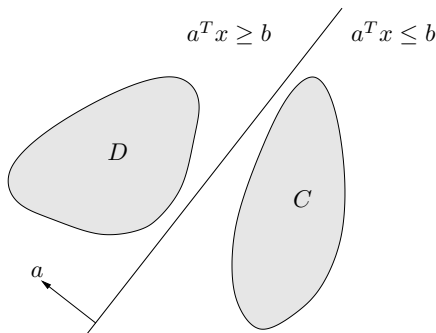
- **linear-fractional function** $f : \mathbb{R}^n \rightarrow \mathbb{R}^m$:

$$f(x) = \frac{Ax + b}{c^\top x + d} \quad \text{dom } f = \{x \mid c^\top x + d > 0\}$$

- images and inverse images of convex sets under perspective and linear-fractional functions are all convex

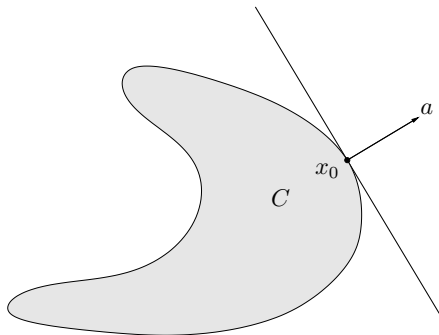
Separating Hyperplane Theorem

- if C, D are nonempty disjoint ($C \cap D = \emptyset$) convex sets, $\exists a \neq 0, b$ s.t.
 $a^\top x \leq b$ for $x \in C$, $a^\top x \geq b$ for $x \in D$
- the hyperplane $\{x \mid a^\top x = b\}$ **separates** C and D
- strict separating requires additional assumptions (e.g. C is closed; D is a singleton)



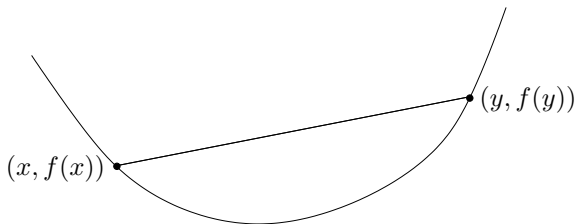
Supporting Hyperplane Theorem

- suppose x_0 is a boundary point of $C \subseteq \mathbb{R}^n$
- **supporting hyperplane** to C at x_0 : $\{x \mid a^\top x = a^\top x_0\}$, where $a \neq 0$ and $a^\top x \leq a^\top x_0 \ \forall x \in C$.
- if C is convex, then there exists a supporting hyperplane at every boundary point of C



Convex Function

- $f : \mathbb{R}^n \rightarrow \mathbb{R}$ is **convex** if $\text{dom } f$ is convex and $\forall x, y \in \text{dom } f, 0 \leq \theta \leq 1$,
 $f(\theta x + (1 - \theta)y) \leq \theta f(x) + (1 - \theta)f(y)$
- $f : \mathbb{R}^n \rightarrow \mathbb{R}$ is **strictly convex** if $\text{dom } f$ is convex and $\forall x, y \in \text{dom } f, x \neq y$,
 $0 < \theta < 1, f(\theta x + (1 - \theta)y) < \theta f(x) + (1 - \theta)f(y)$
- f is **concave** if $-f$ is convex



Example Functions on \mathbb{R}

- convex functions

- affine: $ax + b$, $\forall a, b \in \mathbb{R}$
- exponential: e^{ax} , $\forall a \in \mathbb{R}$
- power: x^α on $x > 0$, $\forall \alpha \geq 1 \vee \alpha \leq 0$
- power of absolute value: $|x|^\alpha$, $\forall \alpha \geq 1$
- positive part (relu): $\max\{x, 0\}$

- concave functions

- affine: $ax + b$, $\forall a, b \in \mathbb{R}$
- power: x^α on $x > 0$, $\forall 0 \leq \alpha \leq 1$
- logarithm: $\log x$ on $x > 0$
- entropy: $-x \log x$ on $x > 0$
- negative part: $\min\{x, 0\}$

Example Convex Functions on \mathbb{R}^n

- affine: $a^\top x + b$
- any norm
 - $\|x\|_p = (|x_1|^p + |x_2|^p + \dots + |x_n|^p)^{\frac{1}{p}}, \forall p > 1$
 - $\|x\|_\infty = \max\{|x_1|, |x_2|, \dots, |x_n|\}$
- sum of squares: $\|x\|_2^2 = x_1^2 + x_2^2 + \dots + x_n^2$
- max function: $\max(x) = \max\{x_1, x_2, \dots, x_n\}$
- softmax / log-sum-exp: $\log(e^{x_1} + e^{x_2} + \dots + e^{x_n})$

Example Functions on $\mathbb{R}^{m \times n}$

- Let $X \in \mathbb{R}^{m \times n}$ be the variable
- general affine function

$$f(X) = \text{tr}(A^\top X) + b = \sum_{i=1}^m \sum_{j=1}^n A_{ij} X_{ij} + b, \quad A \in \mathbb{R}^{m \times n}, \quad b \in \mathbb{R}$$

- spectral norm (maximum singular value) is convex:

$$f(X) = \|X\|_2 = \sigma_{\max}(X) = \sqrt{\lambda_{\max}(X^\top X)}$$

- log determinant is concave:

$$f(X) = \log \det X, \quad X \in \mathbf{S}_{++}^n$$

Extended-Value Extension

- suppose f is convex on \mathbb{R}^n
- its extended-value extension $\tilde{f} : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{\infty\}$ is defined as

$$\tilde{f}(x) = \begin{cases} f(x) & x \in \text{dom } f \\ \infty & x \notin \text{dom } f \end{cases}$$

- this often simplifies notation; e.g. the condition

$$0 \leq \theta \leq 1 \implies \tilde{f}(\theta x + (1 - \theta)y) \leq \theta \tilde{f}(x) + (1 - \theta) \tilde{f}(y)$$

(as an inequality in $\mathbb{R} \cup \{\infty\}$), means the same as the two conditions combine

- $\text{dom } f$ is convex
- $x, y \in \text{dom } f, 0 \leq \theta \leq 1 \implies f(\theta x + (1 - \theta)y) \leq \theta f(x) + (1 - \theta)f(y)$

Restriction of a Convex Function to a Line

- $f : \mathbb{R}^n \rightarrow \mathbb{R}$ is convex (concave) $\iff g : \mathbb{R} \rightarrow \mathbb{R}$,

$$g(t) = f(x + tv), \quad \text{dom } g = \{t \mid x + tv \in \text{dom } f\}$$

is convex (concave) in t for all $x \in \text{dom } f$ and $v \in \mathbb{R}^n$

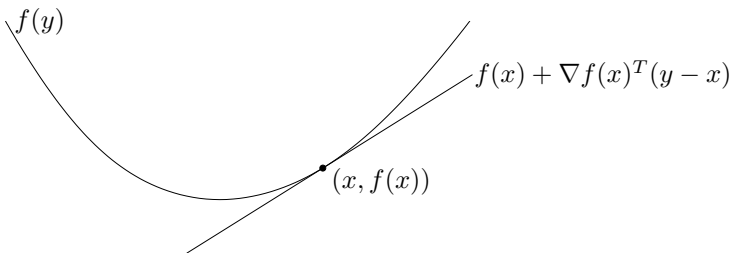
- useful for checking convexity / concavity of multivariate f ; e.g. to check the concavity of log determinant: Let $X \in S_{++}^n$, $V \in S^n$,

$$\begin{aligned} g(t) &= f(X + tV) = \log \det(X + tV) \\ &= \log \det \left(X^{\frac{1}{2}} \left(I + t X^{-\frac{1}{2}} V X^{-\frac{1}{2}} \right) X^{\frac{1}{2}} \right) \\ &= \log \det X + \log \det \left(I + t X^{-\frac{1}{2}} V X^{-\frac{1}{2}} \right) \\ &= \log \det X + \sum_{i=1}^n \log(1 + t\lambda_i) \end{aligned}$$

where λ_i are the eigenvalues of $X^{-\frac{1}{2}} V X^{-\frac{1}{2}}$; g is concave in t

First-Order Condition

- $f : \mathbb{R}^n \rightarrow \mathbb{R}$ is twice differentiable if $\text{dom } f$ is open and the gradient ∇f exists at each $x \in \text{dom } f$.
- **first-order condition** differentiable f with convex domain is convex $\iff f(y) \geq f(x) + \nabla f(x)^\top (y - x), \forall x, y \in \text{dom } f$
- first order Taylor approximation of convex f is a **global underestimator** of f



Second-Order Condition

- $f : \mathbb{R}^n \rightarrow \mathbb{R}$ is **differentiable** if $\text{dom } f$ is open and the Hessian matrix $\nabla^2 f \in \mathbb{S}^n$ exists at each $x \in \text{dom } f$:

$$\{\nabla^2 f(x)\}_{ij} = \frac{\partial^2 f}{\partial x_i \partial x_j}(x)$$

- **second-order condition** for twice differentiable f with convex domain is convex:
 - f is convex $\iff \nabla^2 f \succcurlyeq 0, \forall x \in \text{dom } f$
 - $\nabla^2 f \succ 0, \forall x \in \text{dom } f \implies f$ is strictly convex

Examples I

- **quadratic function:** $f(x) = \frac{1}{2} x^\top P x + q^\top x + r$ with $P \in \mathbb{S}^n$

$$\nabla f(x) = P x + q, \quad \nabla^2 f(x) = P$$

convex if $P \succcurlyeq 0$ (concave if $P \preccurlyeq 0$)

- **least-squares objective:** $f(x) = \|A x - b\|^2$

$$\nabla f(x) = 2A^\top (A x - b), \quad \nabla^2 f(x) = 2A^\top A$$

convex for any A

- **quadratic-over-linear function:** $f(x, y) = \frac{x^2}{y}, y > 0$

$$\nabla f(x, y) = \left(\frac{2x}{y} \quad -\frac{x^2}{y^2} \right), \quad \nabla^2 f(x, y) = \frac{2}{y^3} \begin{pmatrix} y^2 & -xy \\ -xy & x^2 \end{pmatrix}$$

convex for $y > 0$

Examples II

- **log-sum-exp function:** $f(x) = \log \left(\sum_{k=1}^n e^{x_k} \right)$ is convex:

$$\nabla^2 f(x) = \frac{1}{\mathbf{1}^\top z} \text{diag}(z) - \frac{1}{(\mathbf{1}^\top z)^2} z z^\top, \quad z_k = e^{x_k}$$

- to show that $\nabla^2 f(x) \succcurlyeq 0$, one must verify $v^\top \nabla^2 f(x) v \geq 0 \forall v$:

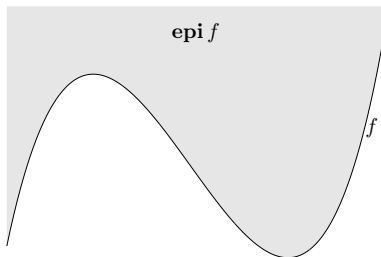
$$v^\top \nabla^2 f(x) v = \frac{(\sum_k z_k v_k^2)(\sum_k z_k) - (\sum_k v_k z_k)^2}{(\sum_k z_k)^2} \geq 0$$

by Cauchy-Schwarz inequality $\left(\sum_k z_k v_k^2 \right) \left(\sum_k z_k \right) \geq \left(\sum_k v_k z_k \right)^2$

- **geometric-mean function:** $f(x) = \left(\prod_{k=1}^n x_k \right)^{\frac{1}{n}}$ on $x \succ 0$ is concave

Epigraph, Sublevel Set

- **α -sublevel set** of $f : \mathbb{R}^n \rightarrow \mathbb{R}$: $C_\alpha = \{x \in \text{dom } f \mid f(x) \leq \alpha\}$
- sublevel sets of convex functions are convex sets
- **epigraph** of $f : \mathbb{R}^n \rightarrow \mathbb{R}$: $\text{epi } f = \{(x, t) \in \mathbb{R}^{n+1} \mid x \in \text{dom } f, f(x) \leq t\}$
- f is convex $\iff \text{epi } f$ is a convex set



Jensen's Inequality

- **basic form:** if f is convex, then for $x, y \in \text{dom } f$, $0 \leq \theta \leq 1$

$$f(\theta x + (1 - \theta)y) \leq \theta f(x) + (1 - \theta)f(y)$$

- **extension:** if f is convex and z is a random variable on $\text{dom } f$,

$$f(\mathbb{E} z) \leq \mathbb{E} f(z)$$

- basic form is special case with discrete distribution

$$\mathbb{P}\{z = x\} = \theta, \quad \mathbb{P}\{z = y\} = 1 - \theta$$

- e.g. for $z \sim \mathcal{N}(\mu, \sigma^2)$, let $f(x) = e^x$, then

$$f(\mathbb{E} z) = f(\mu) = e^\mu \leq e^{\mu + \frac{\sigma^2}{2}} = \mathbb{E} f(z)$$

Showing Convexity of a Function

- apply definition (often simplified by restricting to a line)
- for twice differentiable functions, show $\nabla^2 f(x) \succcurlyeq 0$
- show that f is obtained from simple convex functions by operations that preserve convexity
 - nonnegative multiple, sum, integral
 - composition with affine function
 - pointwise maximum and supremum
 - partial minimization
 - composition
 - perspective

Nonnegative Multiple, Sum, Integral

- **nonnegative multiple:** αf is convex if f is convex and $\alpha \geq 0$
- **sum:** $f_1 + f_2$ is convex if f_1, f_2 is convex
- **infinite sum:** if each of f_i is convex, then $\sum_{i=1}^{\infty} f_i$ is convex
- **integral:** if $f(x, \alpha)$ is convex in x for each $\alpha \in \mathcal{A}$, then

$$\int_{\alpha \in \mathcal{A}} f(x, \alpha) d\alpha$$

is convex

- analogous rules for concave functions

Composition with Affine Function

- $f(Ax + b)$ is convex if f is convex
- e.g.
 - log barrier for linear inequalities

$$f(x) = - \sum_{i=1}^m \log(b_i - a_i^\top x)$$

$$\text{dom } f = \{x \mid a_i^\top x < b_i, \ i = 1, 2, \dots, m\}$$

- norm approximation error (any norm)

$$f(x) = \|Ax - b\|$$

Pointwise Maximum

- $f(x) = \max \{f_1(x), f_2(x), \dots, f_m(x)\}$ is convex if each f_i is convex
- e.g.
 - piecewise linear function

$$f(x) = \max_i (a_i^\top x + b_i)$$

- sum of r largest components of $x \in \mathbb{R}^n$

$$f(x) = x_{[1]} + x_{[2]} + \dots + x_{[r]}$$

where $x_{[i]}$ is i -th largest component of x . Note that

$$f(x) = \max \{x_{i_1} + x_{i_2} + \dots + x_{i_r} \mid 1 \leq i_1 < i_2 < \dots < i_r \leq n\}$$

Pointwise Supremum

- $g(x) = \sup_{y \in \mathcal{A}} f(x, y)$ is convex if $f(x, y)$ is convex in x for each $y \in \mathcal{A}$
- e.g.
 - distance to farthest point in a set C

$$f(x) = \sup_{y \in C} \|x - y\|$$

- maximum eigenvalue of symmetric matrix

$$\lambda_{\max}(X) = \sup_{\|y\|_2=1} y^\top X y, \quad X \in \mathbf{S}^n$$

- support function of a set C

$$S_C(x) = \sup_{y \in C} y^\top x$$

Partial Minimization

- the function $g(x) = \inf_{y \in C} f(x, y)$ is called the **partial minimization** of f w.r.t. y
- if $f(x, y)$ is convex in (x, y) and C is a convex set, then partial minimization g is convex
- e.g.
 - let $f(x, y) = x^\top A x + 2x^\top B y + y^\top C y$ with $\begin{pmatrix} A & B \\ B^\top & C \end{pmatrix} \succcurlyeq 0, C \succ 0$;
minimizing over y gives

$$g(x) = \inf_{y \in C} f(x, y) = x^\top (A - BC^{-1}B^\top) x$$

g is convex, hence Schur complement $A - BC^{-1}B^\top \succcurlyeq 0$

- distance to a convex set S

$$\text{dist}(x, S) = \inf_{y \in S} \|x - y\|$$

Composition with Scalar Functions

- composition of $g : \mathbb{R}^n \rightarrow \mathbb{R}$ and $h : \mathbb{R} \rightarrow \mathbb{R}$ is $f(x) = h(g(x))$ ($f = h \circ g$)
- composition f is convex if
 - g convex, h convex, \tilde{h} nondecreasing; or
 - g concave, h convex, \tilde{h} nonincreasing
- proof for $n = 1$, differentiable g, h

$$f''(x) = h''(g(x)) g'(x)^2 + h'(g(x)) g''(x)$$

- e.g.
 - $f(x) = e^{g(x)}$ is convex if g is convex
 - $f(x) = \frac{1}{g(x)}$ is convex if g is concave and positive

Composition: General

- composition of $g : \mathbb{R}^n \rightarrow \mathbb{R}^k$ and $h : \mathbb{R}^k \rightarrow \mathbb{R}$ is
$$f(x) = h(g_1(x), g_2(x), \dots, g_k(x))$$
- composition f is convex if h is convex and for each i , one of the following holds:
 - g_i convex, \tilde{h} nondecreasing in its i -th argument
 - g_i concave, \tilde{h} nonincreasing in its i -th argument
 - g_i affine
- e.g.
 - $\log \left(\sum_{i=1}^m e^{g_i(x)} \right)$ is convex if each g_i is convex
 - $\frac{p(x)^2}{q(x)}$ is convex if p is nonnegative and convex and q is positive and concave

- perspective of $f : \mathbb{R}^n \rightarrow \mathbb{R}$ is the function $g(x, t) : \mathbb{R}^n \times \mathbb{R} \rightarrow \mathbb{R}$ defined as

$$g(x, t) = t f\left(\frac{x}{t}\right), \quad \text{dom } g = \left\{ (x, t) \mid \frac{x}{t} \in \text{dom } f, t > 0 \right\}$$

- g is convex if f is convex

- e.g.

- $f(x) = x^\top x$ is convex, so $g(x, t) = \frac{x^\top x}{t}$ is convex if $t > 0$
- $f(x) = -\log x$ is convex, so the **relative entropy**

$$g(x, t) = t \log t - t \log x$$

is convex on $x > 0, t > 0$

Convexity Verification: An Example

- test the convexity of $f(x, y) = \frac{(x - y)^2}{1 - \max(x, y)}$, $x < 1$, $y < 1$
- x , y , and 1 are affine
- $\max(x, y)$ is convex; $x - y$ is affine
- $1 - \max(x, y)$ is concave
- $\frac{u^2}{v}$ is convex, monotone decreasing in v for $v > 0$
- f is composition of $\frac{u^2}{v}$ with $u = x - y$, $v = 1 - \max(x, y)$, hence convex

Convexity Verification: A Caveat

- test the convexity of $f(x) = \sqrt{1 + x^2}$
- $\sqrt{\cdot}$ is concave
- $1, x^2$ are convex
- $\sqrt{1 + x^2}$ is ... indefinite ?
- but, note that $\|\cdot\|_2$ is convex
- $\sqrt{1 + x^2}$ can be represented as the 2-norm of vector $(1, x)$ — $\|(1, x)\|_2$, hence is convex
- The general composition rules are only sufficient, not necessary

Standard Form of General Optimization

$$\begin{array}{ll}\text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i = 1, 2, \dots, m \\ & h_i(x) = 0, \quad i = 1, 2, \dots, p\end{array}$$

- $x \in \mathbb{R}^n$ is the optimization variable
- $f_0(x) : \mathbb{R}^n \rightarrow \mathbb{R}$ is the objective / cost
- $f_i(x) : \mathbb{R}^n \rightarrow \mathbb{R}, i = 1, 2, \dots, m$ are the inequality constraints
- $h_i(x) : \mathbb{R}^n \rightarrow \mathbb{R}, i = 1, 2, \dots, p$ are the equality constraints

Standard Form of Convex Optimization

$$\begin{array}{ll}\text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq 0, \quad i = 1, 2, \dots, m \\ & a_i^\top x = b_i, \quad i = 1, 2, \dots, p\end{array}$$

- objective and inequality constraints f_0, f_1, \dots, f_m are convex
- equality constraints are affine, often written as $Ax = b$
- feasible and optimal sets of a convex optimization problem are convex

Local and Global Optima

Theorem

Locally optimal point of a convex optimization problem is (globally) optimal.

Proof

- suppose x is locally optimal, but $\exists y$ with $f_0(y) < f_0(x)$
- x locally optimal means $\exists R > 0$ such that if x' is feasible and $\|x' - x\| \leq R$, then $f_0(x') \geq f_0(x)$
- set $z = \theta y + (1 - \theta)x$ with $\theta = \frac{R}{2\|y - x\|_2}$
- $\|y - x\|_2 > R$, so $0 < \theta < \frac{1}{2}$
- z is a convex combination of two feasible points, hence also feasible
- $\|z - x\|_2 = \frac{R}{2}$ and $f_0(z) \leq \theta f_0(y) + (1 - \theta)f_0(x) < f_0(x)$, which contradicts that x is locally optimal