

Operations Research  
11. The Simplex Method. Dual of Linear Program

# THE SIMPLEX METHOD: A DEMONSTRATION

An iterative process starting with suboptimal solution and stopping at a feasible solution which cannot be improved

$$\begin{array}{ll}\text{maximize} & 5x_1 + 4x_2 + 3x_3 \\ \text{subject to} & 2x_1 + 3x_2 + x_3 \leq 5 \\ & 4x_1 + x_2 + 2x_3 \leq 11 \\ & 3x_1 + 4x_2 + 2x_3 \leq 8 \\ & x_1 \geq 0, x_2 \geq 0, x_3 \geq 0\end{array}$$

maximize	$\zeta =$	$+ 5 x_1 + 4 x_2 + 3 x_3$
subject to	$w_1 =$	$5 - 2 x_1 - 3 x_2 - 1 x_3$
	$w_2 =$	$11 - 4 x_1 - 1 x_2 - 2 x_3$
	$w_3 =$	$8 - 3 x_1 - 4 x_2 - 2 x_3$
$x_1, x_2, x_3, w_1, w_2, w_3 \geq 0$		

Turn inequalities to equalities: add nonnegative slack variables  $w_1, w_2, w_3$

# THE SIMPLEX METHOD: A DEMONSTRATION (CONT'D)

$$\begin{array}{lcl}
 \text{maximize} & \zeta = & + 5 x_1 + 4 x_2 + 3 x_3 \\
 \text{subject to} & w_1 = & 5 - 2 x_1 - 3 x_2 - 1 x_3 \\
 & w_2 = & 11 - 4 x_1 - 1 x_2 - 2 x_3 \\
 & w_3 = & 8 - 3 x_1 - 4 x_2 - 2 x_3
 \end{array}$$

$$x_1, x_2, x_3, w_1, w_2, w_3 \geq 0$$

$$\begin{array}{lcl}
 \text{maximize} & \zeta = & \frac{25}{2} - \frac{5}{2} w_1 - \frac{7}{2} x_2 + \frac{1}{2} x_3 \\
 \text{subject to} & x_1 = & \frac{5}{2} - \frac{1}{2} w_1 - \frac{3}{2} x_2 - \frac{1}{2} x_3 \\
 & w_2 = & 1 + 2 w_1 + 5 x_2 \\
 & w_3 = & \frac{1}{2} + \frac{3}{2} w_1 + \frac{1}{2} x_2 - \frac{1}{2} x_3
 \end{array}$$

$$x_1, x_2, x_3, w_1, w_2, w_3 \geq 0$$

1st step: take  $x_1 = 0$ ,  $x_2 = 0$ ,  $x_3 = 0$ , then  $w_1 = 5$ ,  $w_2 = 11$ ,  $w_3 = 8$

Find in  $\zeta$  the variable with the greatest positive coefficient: here  $+5x_1$

Increase  $x_1$  and keep  $x_2 = x_3 = 0$ ; Note the constraints  $w_1, w_2, w_3 \geq 0$ .

$$w_1 = 5 - 2x_1 \geq 0 \implies x_1 \leq \frac{5}{2}; \quad w_2 = 11 - 4x_1 \geq 0 \implies x_1 \leq \frac{11}{4}; \quad w_3 = 8 - 3x_1 \geq 0 \implies x_1 \leq \frac{8}{3}$$

2nd step: take  $x_1 = \frac{5}{2}$ ,  $x_2 = 0$ ,  $x_3 = 0$ , then  $w_1 = 0$ ,  $w_2 = 1$ ,  $w_3 = \frac{1}{2}$

Swap the roles of  $x_1$  and  $w_1$ , then the first constraint becomes  $x_1 = \frac{5}{2} - \frac{1}{2}w_1 - \frac{3}{2}x_2 - \frac{1}{2}x_3$

Rewrite all other constraints and  $\zeta$  with  $x_1 = \frac{5}{2} - \frac{1}{2}w_1 - \frac{3}{2}x_2 - \frac{1}{2}x_3$

# THE SIMPLEX METHOD: A DEMONSTRATION (CONT'D)

$$\begin{array}{rcl}
 \text{maximize} & \zeta = \frac{25}{2} - \frac{5}{2} w_1 - \frac{7}{2} x_2 & + \frac{1}{2} x_3 \\
 \text{subject to} & x_1 = \frac{5}{2} - \frac{1}{2} w_1 - \frac{3}{2} x_2 - \frac{1}{2} x_3 \\
 & w_2 = 1 + 2 w_1 + 5 x_2 \\
 & w_3 = \frac{1}{2} + \frac{3}{2} w_1 + \frac{1}{2} x_2 - \frac{1}{2} x_3
 \end{array}$$

$$x_1, x_2, x_3, w_1, w_2, w_3 \geq 0$$

$$\begin{array}{rcl}
 \text{maximize} & \zeta = 13 - 1 w_1 - 3 x_2 - 1 w_3 \\
 \text{subject to} & x_1 = 2 - 2 w_1 - 2 x_2 + 1 w_3 \\
 & w_2 = 1 + 2 w_1 + 5 x_2 \\
 & x_3 = 1 + 3 w_1 + 1 x_2 - 2 w_3
 \end{array}$$

$$x_1, x_2, x_3, w_1, w_2, w_3 \geq 0$$

Find in  $\zeta$  the variable with the greatest positive coefficient: here  $+\frac{1}{2}x_3$

Increase  $x_3$  and keep  $w_1 = x_2 = 0$ ; Note the constraints  $x_1, w_3 \geq 0$ .

$$x_1 = \frac{5}{2} - \frac{1}{2}x_3 \geq 0 \implies x_3 \leq 5; \quad w_3 = \frac{1}{2} - \frac{1}{2}x_3 \geq 0 \implies x_3 \leq 1$$

3rd step: take  $x_1 = 2, x_2 = 0, x_3 = 1$ , then  $w_1 = 0, w_2 = 1, w_3 = 0$

Swap the roles of  $x_3$  and  $w_3$ , then the last constraint becomes  $x_3 = 1 + 3w_1 + x_2 - 2w_3$

Rewrite all other constraints and  $\zeta$  with  $x_3 = 1 + 3w_1 + x_2 - 2w_3$

Now  $\zeta$  has no variable with positive coefficient; stop.  $\max \zeta = 13$  with  $x_1 = 2, x_2 = 0, x_3 = 1$ .

# THE SIMPLEX METHOD: ANOTHER EXAMPLE

$$\begin{array}{ll}\text{maximize} & 2x_1 + 3x_2 + 4x_3 \\ \text{subject to} & 2x_2 + 3x_3 \leq 5 \\ & x_1 + x_2 + 2x_3 \leq 4 \\ & x_1 + 2x_2 + 3x_3 \leq 7 \\ & x_1 \geq 0, x_2 \geq 0, x_3 \geq 0\end{array}$$

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$$\begin{array}{ll}\text{maximize} & \zeta = \quad + 2x_1 + 3x_2 + 4x_3 \\ \text{subject to} & w_1 = 5 \quad - 2x_2 - 3x_3 \\ & w_2 = 4 - 1x_1 - 1x_2 - 2x_3 \\ & w_3 = 7 - 1x_1 - 2x_2 - 3x_3\end{array}$$

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$$x_1, x_2, x_3, w_1, w_2, w_3 \geq 0$$

Turn inequalities to equalities: add nonnegative slack variables  $w_1, w_2, w_3$

# THE SIMPLEX METHOD: ANOTHER EXAMPLE (CONT'D)

$$\begin{array}{rcl}
 \text{maximize} & \zeta = & + 2 x_1 + 3 x_2 + 4 x_3 \\
 \text{subject to} & w_1 = 5 & - 2 x_2 - 3 x_3 \\
 & w_2 = 4 - 1 x_1 - 1 x_2 - 2 x_3 \\
 & w_3 = 7 - 1 x_1 - 2 x_2 - 3 x_3
 \end{array}$$

$$x_1, x_2, x_3, w_1, w_2, w_3 \geq 0$$

$$\begin{array}{rcl}
 \text{maximize} & \zeta = & \frac{20}{3} + 2 x_1 + \frac{1}{3} x_2 - \frac{4}{3} w_1 \\
 \text{subject to} & x_3 = & \frac{5}{3} - \frac{2}{3} x_2 - \frac{1}{3} w_1 \\
 & w_2 = & \frac{3}{3} - 1 x_1 + \frac{1}{3} x_2 + \frac{3}{3} w_1 \\
 & w_3 = & 2 - 1 x_1 + 1 w_1
 \end{array}$$

$$x_1, x_2, x_3, w_1, w_2, w_3 \geq 0$$

1st step: take  $x_1 = 0, x_2 = 0, x_3 = 0$ , then  $w_1 = 5, w_2 = 4, w_3 = 7$

Find in  $\zeta$  the variable with the greatest positive coefficient: here  $+4x_3$

Increase  $x_3$  and keep  $x_1 = x_2 = 0$ ; Note the constraints  $w_1, w_2, w_3 \geq 0$ .

$$w_1 = 5 - 3x_3 \geq 0 \implies x_3 \leq \frac{5}{3}; \quad w_2 = 4 - 2x_3 \geq 0 \implies x_3 \leq 2; \quad w_3 = 7 - 3x_3 \geq 0 \implies x_3 \leq \frac{7}{3}$$

2nd step: take  $x_1 = 0, x_2 = 0, x_3 = \frac{5}{3}$ , then  $w_1 = 0, w_2 = \frac{2}{3}, w_3 = 2$

Swap the roles of  $x_3$  and  $w_1$ , then the first constraint becomes  $x_3 = \frac{5}{3} - \frac{2}{3}x_2 - \frac{1}{3}w_1$

Rewrite all other constraints and  $\zeta$  with  $x_3 = \frac{5}{3} - \frac{2}{3}x_2 - \frac{1}{3}w_1$

# THE SIMPLEX METHOD: ANOTHER EXAMPLE (CONT'D)

$$\begin{array}{ll}
 \text{maximize} & \zeta = \frac{20}{3} + 2x_1 + \frac{1}{3}x_2 - \frac{4}{3}w_1 \\
 \text{subject to} & x_3 = \frac{5}{3} - \frac{2}{3}x_2 - \frac{1}{3}w_1 \\
 & w_2 = \frac{2}{3} - 1x_1 + \frac{1}{3}x_2 + \frac{1}{3}w_1 \\
 & w_3 = 2 - 1x_1 + 1w_1
 \end{array}$$

$$x_1, x_2, x_3, w_1, w_2, w_3 \geq 0$$

$$\begin{array}{ll}
 \text{maximize} & \zeta = 8 - 2w_2 + 1x_2 \\
 \text{subject to} & x_3 = \frac{5}{3} - \frac{2}{3}x_2 - \frac{1}{3}w_1 \\
 & x_1 = \frac{2}{3} - 1w_2 + \frac{1}{3}x_2 + \frac{1}{3}w_1 \\
 & w_3 = \frac{4}{3} + 1w_2 - \frac{1}{3}x_2 + \frac{1}{3}w_1
 \end{array}$$

$$x_1, x_2, x_3, w_1, w_2, w_3 \geq 0$$

Find in  $\zeta$  the variable with the greatest positive coefficient: here  $+2x_1$

Increase  $x_1$  and keep  $x_2 = w_1 = 0$ ; Note the constraints  $w_2, w_3 \geq 0$ .

$$w_2 = \frac{2}{3} - 1x_1 \geq 0 \implies x_1 \leq \frac{2}{3}; \quad w_3 = 2 - 1x_1 \geq 0 \implies x_1 \leq 2$$

3rd step: take  $x_1 = \frac{2}{3}$ ,  $x_2 = 0$ ,  $x_3 = \frac{5}{3}$ , then  $w_1 = 0$ ,  $w_2 = 0$ ,  $w_3 = \frac{4}{3}$

Swap the roles of  $x_1$  and  $w_2$ , then the second constraint becomes  $x_1 = \frac{2}{3} - w_2 + \frac{1}{3}x_2 + \frac{2}{3}w_1$

Rewrite all other constraints and  $\zeta$  with  $x_1 = \frac{2}{3} - w_2 + \frac{1}{3}x_2 + \frac{2}{3}w_1$

# THE SIMPLEX METHOD: ANOTHER EXAMPLE (CONT'D)

$$\begin{array}{rcl}
 \text{maximize} & \zeta = 8 - 2 w_2 & + 1 x_2 \\
 \text{subject to} & x_3 = \frac{5}{3} & - \frac{2}{3} x_2 - \frac{1}{3} w_1 \\
 & x_1 = \frac{4}{3} - 1 w_2 & + \frac{1}{3} x_2 + \frac{2}{3} w_1 \\
 & w_3 = \frac{4}{3} + 1 w_2 & - \frac{1}{3} x_2 + \frac{1}{3} w_1
 \end{array}$$

$$x_1, x_2, x_3, w_1, w_2, w_3 \geq 0$$

$$\begin{array}{rcl}
 \text{maximize} & \zeta = \frac{21}{2} - 2 w_2 - \frac{3}{2} x_3 & - \frac{1}{2} w_1 \\
 \text{subject to} & x_2 = \frac{5}{2} & - \frac{3}{2} x_3 - \frac{1}{2} w_1 \\
 & x_1 = \frac{3}{2} - 1 w_2 & - \frac{1}{2} x_3 + \frac{1}{2} w_1 \\
 & w_3 = \frac{1}{2} + 1 w_2 & + \frac{1}{2} x_3 + \frac{1}{2} w_1
 \end{array}$$

$$x_1, x_2, x_3, w_1, w_2, w_3 \geq 0$$

Find in  $\zeta$  the variable with the greatest positive coefficient: here  $+1x_2$

Increase  $x_2$  and keep  $w_1, w_2 = 0$ ; Note the constraints  $x_3, w_3 \geq 0$ .

$$x_3 = \frac{5}{3} - \frac{2}{3}x_2 \geq 0 \implies x_2 \leq \frac{5}{2}; \quad w_3 = \frac{4}{3} - \frac{1}{3}x_2 \geq 0 \implies x_2 \leq 4$$

3rd step: take  $x_1 = \frac{3}{2}$ ,  $x_2 = \frac{5}{2}$ ,  $x_3 = 0$ , and  $w_1 = 0$ ,  $w_2 = 0$ ,  $w_3 = \frac{1}{2}$

Swap the roles of  $x_2$  and  $x_3$ , then the first constraint becomes  $x_2 = \frac{5}{2} - \frac{3}{2}x_3 - \frac{1}{2}w_1$

Rewrite all other constraints and  $\zeta$  with  $x_2 = \frac{5}{2} - \frac{3}{2}x_3 - \frac{1}{2}w_1$

Now  $\zeta$  has no variable with positive coefficient; stop.  $\max \zeta = \frac{21}{2}$  with  $x_1 = \frac{3}{2}$ ,  $x_2 = \frac{5}{2}$ ,  $x_3 = 0$ .



# THE SIMPLEX METHOD: A FORMAL INTRODUCTION

LP in standard form:

$$\begin{aligned} & \text{maximize} && \sum_{j=1}^n c_j x_j \\ & \text{subject to} && \sum_{j=1}^n a_{ij} x_j \leq b_i, \quad i = 1, 2, \dots, m \\ & && x_j \geq 0, \quad j = 1, 2, \dots, n \end{aligned}$$

Introduce slack variables  $w_i$ ,  $i = 1, 2, \dots, m$   
and  $\zeta$  for the objective function value:

$$\begin{aligned} \zeta &= \sum_{j=1}^n c_j x_j \\ w_i &= b_i - \sum_{j=1}^n a_{ij} x_j, \quad i = 1, 2, \dots, m \end{aligned}$$

maximize	$\zeta = \frac{25}{2} - \frac{5}{2} w_1 - \frac{7}{2} x_2 + \frac{1}{2} x_3$
subject to	$x_1 = \frac{5}{2} - \frac{1}{2} w_1 - \frac{3}{2} x_2 - \frac{1}{2} x_3$ $w_2 = 1 + 2 w_1 + 5 x_2$ $w_3 = \frac{1}{2} + \frac{3}{2} w_1 + \frac{1}{2} x_2 - \frac{1}{2} x_3$

dictionary

basic variables (lhs):  $x_1, w_2, w_3$

nonbasic variables (rhs):  $w_1, x_2, x_3$

entering variable: the variable

nonbasic  $\Rightarrow$  basic

leaving variable: the variable

basic  $\Rightarrow$  nonbasic

rename  $(x_1, x_2, \dots, x_n, w_1, w_2, \dots, w_m) \Rightarrow$   
 $(x_1, x_2, \dots, x_n, x_{n+1}, x_{n+2}, \dots, x_{n+m})$ , so

$$x_{n+i} = b_i - \sum_{j=1}^n a_{ij} x_j, \quad i = 1, 2, \dots, m$$

The starting dictionary

$$\zeta = \sum_{j=1}^n c_j x_j$$

$$x_{n+i} = b_i - \sum_{j=1}^n a_{ij} x_j, \quad i = 1, 2, \dots, m$$

Each dictionary has  $m$  basic variables and  $n$  nonbasic variables

$\mathcal{B}$ : subset from  $\{1, 2, \dots, n+m\}$  that of the basic variables

$\mathcal{N}$ : subset from  $\{1, 2, \dots, n+m\}$  that of the nonbasic variables

Initially  $\mathcal{N} = \{1, 2, \dots, n\}$  and  $\mathcal{B} = \{n+1, n+2, \dots, n+m\}$

The current dictionary is of the form

$$\zeta = \bar{\zeta} + \sum_{j \in \mathcal{N}} \bar{c}_j x_j$$

$$x_i = \bar{b}_i - \sum_{j \in \mathcal{N}} \bar{a}_{ij} x_j, \quad i \in \mathcal{B}$$

The entering variable is chosen to increase  $\zeta$ : pick  $k \in \{j \in \mathcal{N} : \bar{c}_j > 0\}$ . If no such  $k$ , then the current solution is optimal.

The leaving variable is chosen to preserve nonnegativity of the current basic variables, so

$$x_i = \bar{b}_i - \bar{a}_{ik} x_k \implies \bar{b}_i - \bar{a}_{ik} x_k \geq 0, \quad i \in \mathcal{B}$$

The rule for selecting leaving variable:

pick  $l \in \{i \in \mathcal{B} : \bar{a}_{ik} > 0 \wedge \frac{\bar{b}_i}{\bar{a}_{ik}} \text{ is minimal}\}$

# THE SIMPLEX METHOD: CASE OF NEGATIVE RHS

Previously in standard form LP

$$\begin{aligned} &\text{maximize} && \sum_{j=1}^n c_j x_j \\ &\text{subject to} && \sum_{j=1}^n a_{ij} x_j \leq b_i, \quad i = 1, 2, \dots, m \\ &&& x_j \geq 0, \quad j = 1, 2, \dots, n \end{aligned}$$

We assume  $b_i \geq 0, \forall i = 1, 2, \dots, m$ .

If not the case, then consider the auxiliary problem

$$\begin{aligned} &\text{maximize} && -x_0 \\ &\text{subject to} && \sum_{j=1}^n a_{ij} x_j - x_0 \leq b_i, \quad i = 1, 2, \dots, m \\ &&& x_j \geq 0, \quad j = 0, 1, 2, \dots, n \end{aligned}$$

The original problem has a feasible solution iff the optimal solution of the auxiliary problem is zero.

First try to convert the initial infeasible dictionary with into a feasible one by one pivot of the “most infeasible variable”.

Proceed with the usual simplex steps.

Discard all  $x_0$  terms; reintroduce the original objective with the substituted nonbasic variables.

An example:

$$\begin{aligned} &\text{maximize} && -2x_1 - x_2 \\ &\text{subject to} && -x_1 + x_2 \leq -1 \\ &&& -x_1 - 2x_2 \leq -2 \\ &&& x_2 \leq 1 \\ &&& x_1 \geq 0, x_2 \geq 0 \end{aligned}$$

# THE SIMPLEX METHOD: CASE OF NEGATIVE RHS (CONT'D)

The original problem:

$$\begin{array}{ll}
 \text{maximize} & -2x_1 - x_2 \\
 \text{subject to} & -x_1 + x_2 \leq -1 \\
 & -x_1 - 2x_2 \leq -2 \\
 & x_2 \leq 1 \\
 & x_1 \geq 0, x_2 \geq 0
 \end{array}$$

The auxiliary problem:

$$\begin{array}{ll}
 \text{maximize} & -x_0 \\
 \text{subject to} & -x_1 + x_2 - x_0 \leq -1 \\
 & -x_1 - 2x_2 - x_0 \leq -2 \\
 & x_2 - x_0 \leq 1 \\
 & x_1 \geq 0, x_2 \geq 0, x_0 \geq 0
 \end{array}$$

Setup the initial infeasible dictionary:

maximize	$\zeta =$	$-1 x_0$
subject to	$w_1 =$	$-1 + 1 x_1 - 1 x_2 + 1 x_0$
	$w_2 =$	$-2 + 1 x_1 + 2 x_2 + 1 x_0$
	$w_3 =$	$1 - 1 x_2 + 1 x_0$
$x_1, x_2, x_0, w_1, w_2, w_3 \geq 0$		

The “most infeasible variable” is  $w_2$ ;  
 Substitute  $x_0$  with  $w_2 + 2 - x_1 - x_2$ :

maximize	$\zeta =$	$-2 + 1 x_1 + 2 x_2 - 1 w_2$
subject to	$w_1 =$	$1 - 3 x_2 + 1 w_2$
	$x_0 =$	$2 - 1 x_1 - 2 x_2 + 1 w_2$
	$w_3 =$	$3 - 1 x_1 - 3 x_2 + 1 w_2$
$x_1, x_2, x_0, w_1, w_2, w_3 \geq 0$		

# THE SIMPLEX METHOD: CASE OF NEGATIVE RHS (CONT'D)

Proceed with ordinary simplex method:

$$\begin{array}{l}
 \text{maximize} \quad \zeta = 0 - 1 x_0 \\
 \text{subject to} \quad x_2 = \frac{1}{3} - \frac{1}{3} w_1 + \frac{1}{3} w_2 \\
 \quad \quad \quad x_1 = \frac{2}{3} - 1 x_0 + \frac{2}{3} w_1 + \frac{1}{3} w_2 \\
 \quad \quad \quad w_3 = \frac{2}{3} + 1 x_0 + \frac{1}{3} w_1 - \frac{1}{3} w_2 \\
 \hline
 x_1, x_2, x_0, w_1, w_2, w_3 \geq 0
 \end{array}$$

Discard all  $x_0$  terms; Reintroduce the original objective

$$\text{maximize} \quad -2x_1 - x_2 \equiv -3 - w_1 - w_2$$

with the substituted nonbasic variables

$$\begin{aligned}
 x_1 &= \frac{4}{3} + \frac{2}{3} w_1 + \frac{1}{3} w_2 \\
 x_2 &= \frac{1}{3} - \frac{1}{3} w_1 + \frac{1}{3} w_2
 \end{aligned}$$

Combine the auxiliary dictionary with the substituted original object, we have the dictionary

$$\begin{array}{l}
 \text{maximize} \quad \zeta = -3 - 1 w_1 - 1 w_2 \\
 \text{subject to} \quad x_2 = \frac{1}{3} - \frac{1}{3} w_1 + \frac{1}{3} w_2 \\
 \quad \quad \quad x_1 = \frac{4}{3} + \frac{2}{3} w_1 + \frac{1}{3} w_2 \\
 \quad \quad \quad w_3 = \frac{2}{3} + \frac{1}{3} w_1 - \frac{1}{3} w_2 \\
 \hline
 x_1, x_2, w_1, w_2, w_3 \geq 0
 \end{array}$$

This is the final form, for the coefficients of all nonbasic variables are negative.

The final maximum of the original problem is  $-3$ .

# THE SIMPLEX METHOD: EXERCISES

Using [Simple Pivot Tool](#) to find the extrema of the following LP problems:

$$\begin{array}{ll}\text{maximize} & 6x_1 + 8x_2 + 5x_3 + 9x_4 \\ \text{subject to} & 2x_1 + x_2 + x_3 + 3x_4 \leq 5 \\ & x_1 + 3x_2 + x_3 + 2x_4 \leq 3 \\ & x_1, x_2, x_3, x_4 \geq 0\end{array}$$

$$\begin{array}{ll}\text{maximize} & 3x_1 + 2x_2 + 4x_3 \\ \text{subject to} & x_1 + x_2 + 2x_3 \leq 4 \\ & 2x_1 + 3x_3 \leq 5 \\ & 2x_1 + x_2 + 3x_3 \leq 7 \\ & x_1, x_2, x_3 \geq 0\end{array}$$

$$\begin{array}{ll}\text{maximize} & 5x_1 + 5x_2 + 3x_3 \\ \text{subject to} & x_1 + 3x_2 + x_3 \leq 3 \\ & -x_1 + 3x_3 \leq 2 \\ & 2x_1 - x_2 + 2x_3 \leq 4 \\ & 2x_1 + 3x_2 - x_3 \leq 2 \\ & x_1, x_2, x_3 \geq 0\end{array}$$

$$\begin{array}{ll}\text{maximize} & -x_1 - 3x_2 - x_3 \\ \text{subject to} & 2x_1 - 5x_2 + x_3 \leq -5 \\ & 2x_1 - x_2 + 2x_3 \leq 4 \\ & x_1, x_2, x_3 \geq 0\end{array}$$

$$\begin{array}{ll}\text{maximize} & 5x_1 + 6x_2 + 9x_3 + 8x_4 \\ \text{subject to} & x_1 + 2x_2 + 3x_3 + x_4 \leq 5 \\ & x_1 + x_2 + 2x_3 + 3x_4 \leq 3 \\ & x_1, x_2, x_3, x_4 \geq 0\end{array}$$

$$\begin{array}{ll}\text{minimize} & x_1 + 8x_2 + 9x_3 + 2x_4 + 7x_5 + 3x_6 \\ \text{subject to} & x_1 + x_2 + x_3 \geq 1 \\ & -x_1 + x_4 + x_5 = 0 \\ & -x_2 - x_4 + x_6 = 0 \\ & x_3 + x_5 + x_6 \leq 1 \\ & x_1, x_2, x_3, x_4, x_5, x_6 \geq 0\end{array}$$

LP in standard form:

$$\begin{array}{ll} \text{maximize} & x_1 + 2x_2 + x_3 + x_4 \\ \text{subject to} & x_1 + 2x_2 + x_3 \leq 2 \end{array} \quad (1)$$

$$x_2 + x_4 \leq 1 \quad (2)$$

$$x_1 + 2x_3 \leq 1 \quad (3)$$

$$x_j \geq 0, \quad j = 1, \dots, 4$$

Suppose the LP-solver find an “optimal”

solution  $x_1 = 1, x_2 = \frac{1}{2}, x_3 = 0, x_4 = \frac{1}{2}$

with maximum  $\frac{5}{2}$ , how can we check this?

Scale (1) by  $\frac{1}{2}$ , add (2), and add (3)

scaled by  $\frac{1}{2}$ , we get for every feasible  $x_j$ ,  
 $j = 1, \dots, 4$

$$x_1 + 2x_2 + \frac{3}{2}x_3 + x_4 \leq \frac{5}{2}$$

The objective

$$x_1 + 2x_2 + x_3 + x_4 \leq x_1 + 2x_2 + \frac{3}{2}x_3 + x_4 \leq \frac{5}{2}$$

so  $\frac{5}{2}$  is indeed optimal.

# THE DUAL OF LINEAR PROGRAM (CONT'D)

LP in standard form:

$$\begin{aligned} & \text{maximize} && \sum_{j=1}^n c_j x_j \\ & \text{subject to} && \sum_{j=1}^n a_{1j} x_j \leq b_1 \\ & && \sum_{j=1}^n a_{2j} x_j \leq b_2 \\ & && \dots \dots \dots \\ & && \sum_{j=1}^n a_{mj} x_j \leq b_m \\ & && x_j \geq 0, \quad j = 1, 2, \dots, n \end{aligned}$$

For every choice of the “scaling factors”  $y_i \geq 0$ ,  $i = 1, 2, \dots, m$ ,

$$y_1 \sum_{j=1}^n a_{1j} x_j + y_2 \sum_{j=1}^n a_{2j} x_j + \dots + y_m \sum_{j=1}^n a_{mj} x_j \leq y_1 b_1 + y_2 b_2 + \dots + y_m b_m$$

Rearrange the inequality as

$$x_1 \sum_{i=1}^m a_{i1} y_i + x_2 \sum_{i=1}^m a_{i2} y_i + \dots + x_n \sum_{i=1}^m a_{in} y_i \leq y_1 b_1 + y_2 b_2 + \dots + y_m b_m$$



# THE DUAL OF LINEAR PROGRAM (CONT'D)

Choose  $y_i$ ,  $i = 1, 2, \dots, m$  such that

$$c_1 \leq \sum_{i=1}^m a_{i1}y_i, \quad c_2 \leq \sum_{i=1}^m a_{i2}y_i, \quad \dots, \quad c_n \leq \sum_{i=1}^m a_{in}y_i$$

Then

$$c_1x_1 + c_2x_2 + \dots + c_nx_n \leq x_1 \sum_{i=1}^m a_{i1}y_i + x_2 \sum_{i=1}^m a_{i2}y_i + \dots + x_n \sum_{i=1}^m a_{in}y_i \leq y_1b_1 + y_2b_2 + \dots + y_mb_m$$

To make the upper bound tighter, we have the LP in dual form:

$$\begin{aligned} & \text{minimize} && \sum_{i=1}^n b_i y_i \\ & \text{subject to} && \sum_{i=1}^n a_{i1} y_i \geq c_1 \\ & && \sum_{i=1}^n a_{i2} y_i \geq c_2 \\ & && \dots \dots \dots \\ & && \sum_{i=1}^n a_{in} y_i \geq c_m \\ & && y_i \geq 0, \quad i = 1, 2, \dots, m \end{aligned}$$

The Primal

$$\begin{array}{ll}\text{maximize} & c^\top x \\ \text{subject to} & Ax \preceq b \\ & x \succeq 0\end{array}$$

The Primal

$$\begin{array}{ll}\text{minimize} & c^\top x \\ \text{subject to} & Ax \succeq b \\ & x \succeq 0\end{array}$$

The Dual

$$\begin{array}{ll}\text{minimize} & b^\top y \\ \text{subject to} & A^\top y \succeq c \\ & y \succeq 0\end{array}$$

The Dual

$$\begin{array}{ll}\text{maximize} & b^\top y \\ \text{subject to} & A^\top y \preceq c \\ & y \succeq 0\end{array}$$

Primal (max)	Dual (min)
$\geq$ constraint	$\leq 0$ variable
$\leq$ constraint	$\geq 0$ variable
$=$ constraint	free variable
$\geq 0$ variable	$\leq$ constraint
$\leq 0$ variable	$\geq$ constraint
free variable	$=$ constraint