# Advanced Microeconomics II Strictly Competitive Games

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# Strictly Competitive Games

#### **Definition**

A strategic game  $\{\{1,2\},(A_i),(\succeq_i)\}$  is strictly competitive if for any  $a \in A$  and  $b \in A$  we have  $a \succeq_1 b$  if and only if  $b \succeq_2 a$ .

Example:

Player 2 
$$\frac{L}{R}$$
 Player 1  $\frac{U}{D}$   $\frac{3,4}{5,2}$   $\frac{6,1}{4,3}$ 

# Maxminimizing Strategies

## **Definition**

Let  $\{\{1,2\},(A_i),(\succeq_i)\}$  be a strictly competitive strategic game. The action  $x^* \in A_1$  is a maxminimizer for player 1 if

$$\min_{y \in A_2} u_1(x^*, y) \ge \min_{y \in A_2} u_1(x, y) \text{ for all } x \in A_1.$$

The action  $y^* \in A_2$  is a maxminimizer for player 2 if

$$\min_{x \in A_1} u_2(x, y^*) \ge \min_{x \in A_1} u_2(x, y) \text{ for all } y \in A_2.$$

## Example:

## Zero-sum Games

#### Definition

A strategic game  $\{\{1,2\},(A_i),(u_i)\}$  is zero-sum if for any  $a\in A$  we have  $u_2(a)=-u_1(a)$ .

Example:

Player 2 
$$L R$$

Player 1  $U 3,-3 6,-6 5,-5 4,-4$ 

# Zero-sum Games and Strictly Competitive Games

#### Lemma

Any strategic game can be translated into a zero-sum game that preserves player preference ordering over outcomes.

#### Lemma

Let  $\{\{1,2\},(A_i),(u_i)\}$  be a zero-sum strategic game. Then  $\max_{y \in A_2} \min_{x \in A_1} u_2(x, y) = -\min_{y \in A_2} \max_{x \in A_1} u_1(x, y)$ . Further,  $y \in A_2$ solves the problem  $\max_{y \in A_2} \min_{x \in A_1} u_2(x, y)$  if and only if it solves the problem  $\min_{y \in A_2} \max_{x \in A_1} u_1(x, y)$ .

- For any function  $f_1 \min_z f(z) = \max_z -f(z)$  and  $arg min_z f(z) = arg max_z - f(z)$ .
- $-\min_{x \in A_1} u_2(x, y) = \max_{x \in A_1} -u_2(x, y) = \max_{x \in A_1} u_1(x, y).$
- $\max_{v \in A_2} \min_{x \in A_1} u_2(x, y) = -\min_{v \in A_2} \max_{x \in A_1} u_1(x, y).$

## Proposition

Let  $G = \{\{1,2\}, (A_i), (u_i)\}$  be a zero-sum strategic game.

- If  $(x^*, y^*)$  is a Nash equilibrium of G then  $x^*$  is a maxminimizer for player 1 and  $y^*$  is a maxminimizer for player 2.
- ② If  $(x^*, y^*)$  is a Nash equilibrium of G then  $\max_x \min_y u_1(x, y) = \min_y \max_x u_1(x, y)$ , and thus all Nash equilibria of G yield the same payoffs.
- If  $\max_x \min_y u_1(x,y) = \min_y \max_x u_1(x,y)$ ,  $x^*$  is a maxminimizer for player 1, and  $y^*$  is a maxminimizer for player 2, then  $(x^*,y^*)$  is a Nash equilibrium of G.

- $(x^*, y^*)$  is a NE  $\Rightarrow u_2(x^*, y^*) \ge u_2(x^*, y)$  for all  $y \in A_2$   $\Rightarrow u_1(x^*, y^*) \le u_1(x^*, y)$  for all  $y \in A_2$  $\Rightarrow u_1(x^*, y^*) = \min_y u_1(x^*, y) \le \max_x \min_y u_1(x, y)$ .
- $(x^*, y^*)$  is a NE  $\Rightarrow u_1(x^*, y^*) \ge u_1(x, y^*)$  for all  $x \in A_1$   $\Rightarrow u_1(x^*, y^*) \ge \min_y u_1(x, y)$  for all  $x \in A_1$  $\Rightarrow u_1(x^*, y^*) \ge \max_x \min_y u_1(x, y)$ .
- Repeat for player 2.
- What have we proven?

- Let  $v^* = \max_x \min_y u_1(x, y) = \min_y \max_x u_1(x, y)$ .
- Since  $x^*$  is a maxminimizer for player 1,  $u_1(x^*, y) \ge v^*$  for all  $y \in A_2$ . In particular,  $u_1(x^*, y^*) \ge v^*$ .
- From the lemma,  $\max_y \min_x u_2(x,y) = -v^*$ .  $y^*$  is a maxminimizer for player  $2 \Rightarrow u_2(x,y^*) \ge -v^*$  for all  $x \in A_1$   $\Rightarrow u_1(x,y^*) \le v^*$  for all  $x \in A_1$ . In particular,  $u_2(x^*,y^*) \ge -v^* \Rightarrow u_1(x^*,y^*) \le v^*$ .
- $v^* = u_1(x^*, y^*) \ge u_1(x, y^*)$  for all  $x \in A_1$ .
- Repeat for player 2.

### Notes:

- A way to find Nash equilibria in strictly competitive games.
- Equilibria are interchangeable.
- $u_1(x^*, y^*)$  is the value of the game for player 1.