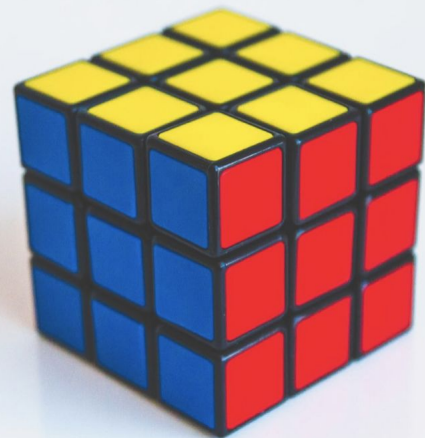


# Rubik's ParaCube: A Collection of Parallel Implementations for Optimal Rubik's Cube Solver

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# Project Background

- Rubik's Cube
  - Representations: F, B, L, R, U, D, F1, ..., U3, D3
  - Solver: input a series of scrambles, output a series of solution steps
- Why Iterative Deepening A\*?
  - Thistlethwaite's 52-move algorithm (not optimal)
  - Optimality, BFS Search?
  - High branching factor ( $K \approx 12$ ), Memory...
  - DFS? May never terminate!



# IDA\* and Korf's Algorithm

## IDA Iterations

0 ...

...

i: visit all blue paths

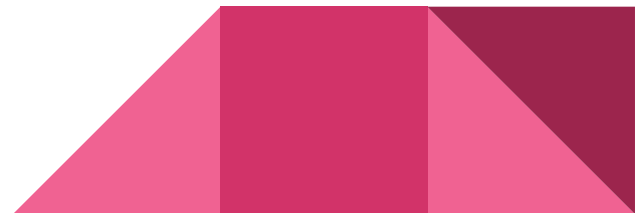
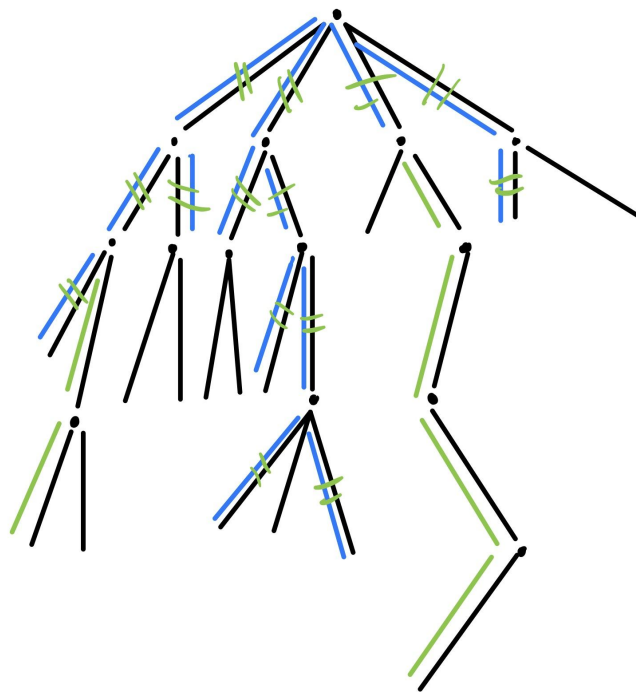
i+1: visit blue + green paths

...

## Korf Pattern Database

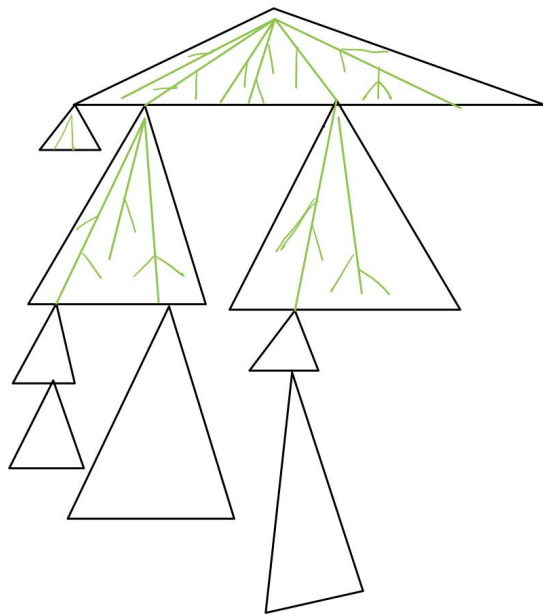
Map cube condition to least number of moves needed to achieve goal state.

To find the correct bound, (which nodes are leaf nodes for iteration i)



# OpenMP Implementation

- Spawn enough tasks
  - To provide load balance
- Allow freedom of task management and multi level task spawn through task steal
  - Exposing worker private problem to peers



# MPI Implementation

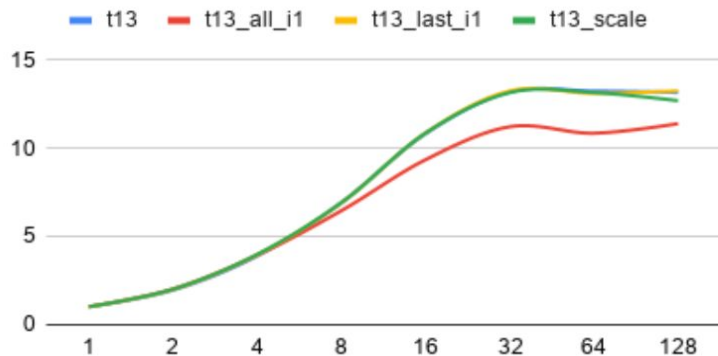
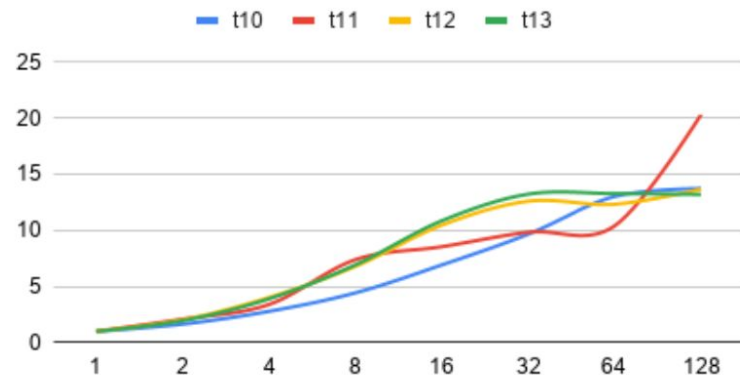
- Frontier Nodes Generation: "Best-First" Expansion
  - Nodes with similar f-values will have sub-trees of comparable sizes
  - Expand nodes with smaller f-value until we have enough frontier nodes
- Dynamic Work Assignment
  - Master generate enough frontier nodes ( $18 \times 18 \times 18$  in our case)
  - Master node send node to worker nodes where they independently perform IDA\* searching
  - Worker nodes report back searching results, if optimal solution detected

master node call `MPI_Abort()`, otherwise, master node update its bound

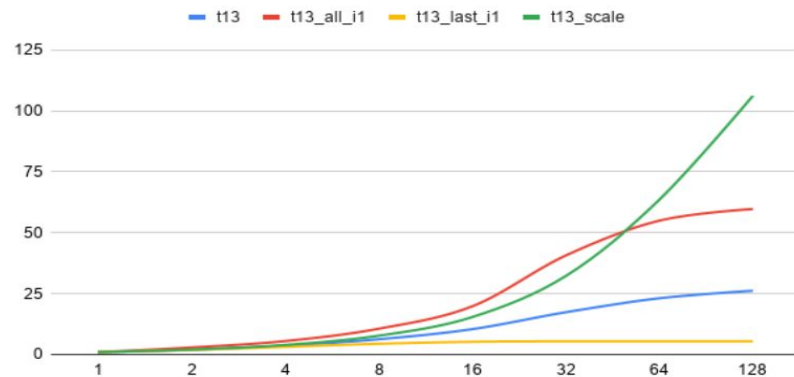
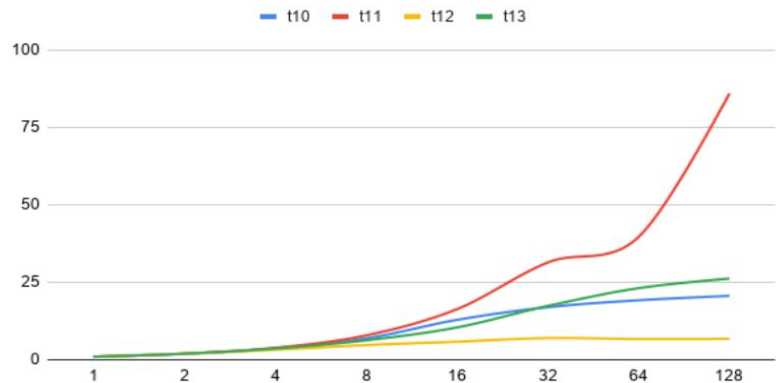


# Results

Speedup Factors for OpenMP on Latedays



Speedup Factors for OpenMPI on Bridges



# Analysis

- Property of IDA
  - For a given iteration, we have a graph to search,  $G_i$
  - Not the entire  $G_i$  is searched in the last iteration for sequential algorithm
  - For the previous iterations, fully searched  $G_{\{i-x\}}$  has  $1/K^x$  as much work
- Efficiency Factor
  - Traversing in different orders will yield different speedup results
  - Parallelization, incur unnecessary work (when not the entire graph  $G_i$  is searched)



# Summary: Efficiency Analysis

Why parallelizing IDA is hard!

And why we are proud of our results!

