Winter 09 **Instructor:** Mihir Bellare
February 9, 2009

## Problem Set 4

Due: Wednesday February 18, 2009, in class.

See course information section (on course web page) for instructions and rules on working on problem sets and turning them in.

**Problem 1.** [30 points] Define the family of functions  $H: \{0,1\}^{64} \times \{0,1\}^{192} \to \{0,1\}^{128}$  as follows:

Show that H is not collision-resistant by presenting a practical adversary A such that  $\mathbf{Adv}_{H}^{\mathrm{cr2-kk}}(A)$  is close to one. (The better the attack, the more points you get.)

**Problem 2.** [40 points] Let  $h: \mathcal{K} \times \{0,1\}^{2b} \to \{0,1\}^{b}$  be a compression function. Define  $H: \mathcal{K} \times \{0,1\}^{4b} \to \{0,1\}^{b}$  as follows:

```
function H(K, M)

Break M into 2b-bit blocks, M = M_1 \parallel M_2

V_1 \leftarrow h(K, M_1); V_2 \leftarrow h(K, M_2)

V \leftarrow h(K, V_1 \parallel V_2)

return V
```

Show that if h is collision-resistant then so is H. Do this by stating and proving an analogue of Theorem 6.8 in the course notes.

**Problem 3.** [50 points] Let sha1:  $\{0,1\}^{672} \to \{0,1\}^{160}$  be the compression function underlying the SHA1 hash function. We define a message authentication scheme  $\Pi = (\mathcal{K}, \text{MAC}, \text{VF})$  as follows. The key generation algorithm returns a random 160 bit string as the key K, and the tagging and verifying algorithms are:

```
Algorithm \mathrm{MAC}_K(M)
Divide M into 512 bit blocks, M=M[1]\dots M[n]
C[0]\leftarrow K
For i=1,\dots,n do
C[i]\leftarrow \mathrm{shal}(C[i-1]\parallel M[i])
EndFor
Return C[n] Algorithm \mathrm{VF}_K(M,\sigma)
If \sigma=\mathrm{MAC}_K(M) then return 1
Else return 0
```

The message space is the set of all strings whose length is a positive multiple of 512.

Present a practical chosen-message attack that succeeds in forgery using one query to the tagging oracle.