



Chapel's Multiresolution Programming Model

Mixing High-level Parallel Abstractions with Lower-level Control

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Northwest C++ Users Group

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What is Chapel?



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What is Chapel?



Chapel: A productive parallel programming language

- portable
- open-source
- a collaborative effort

Goals:

- Support general parallel programming
 - “any parallel algorithm on any parallel hardware”
- Make parallel programming at scale far more productive



Scalable Parallel Programming Concerns



Q: What do HPC programmers need from a language?

A: *Serial Code*: Software engineering and performance

Parallelism: What should execute simultaneously?

Locality: Where should those tasks execute?

Mapping: How to map the program to the system?

Separation of Concerns: Decouple these issues

Chapel is a language designed to address these needs from first principles



Chapel and Other Languages



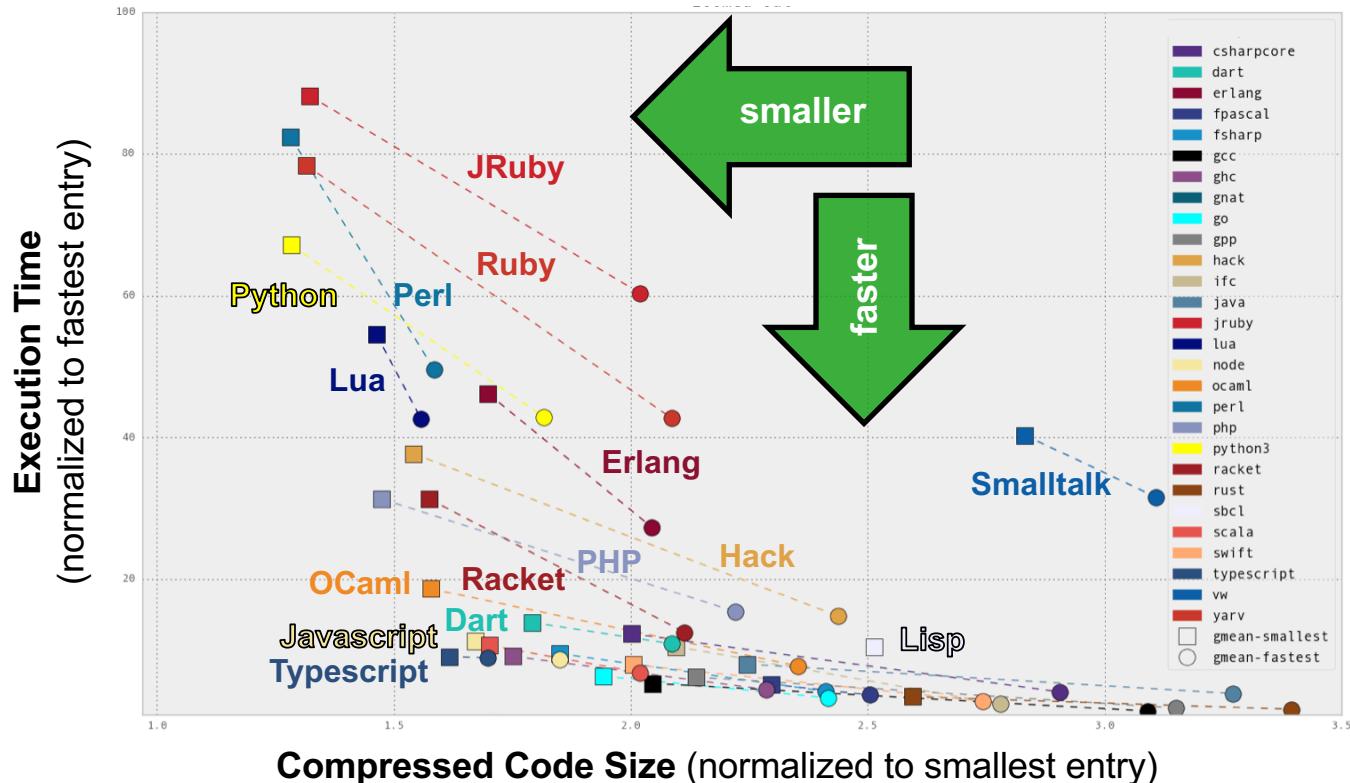
Chapel strives to be as...

- ...**programmable** as Python
- ...**fast** as Fortran
- ...**scalable** as MPI, SHMEM, or UPC
- ...**portable** as C
- ...**flexible** as C++
- ...**fun** as [your favorite programming language]



CLBG Cross-Language Summary

(Oct 2017 standings)



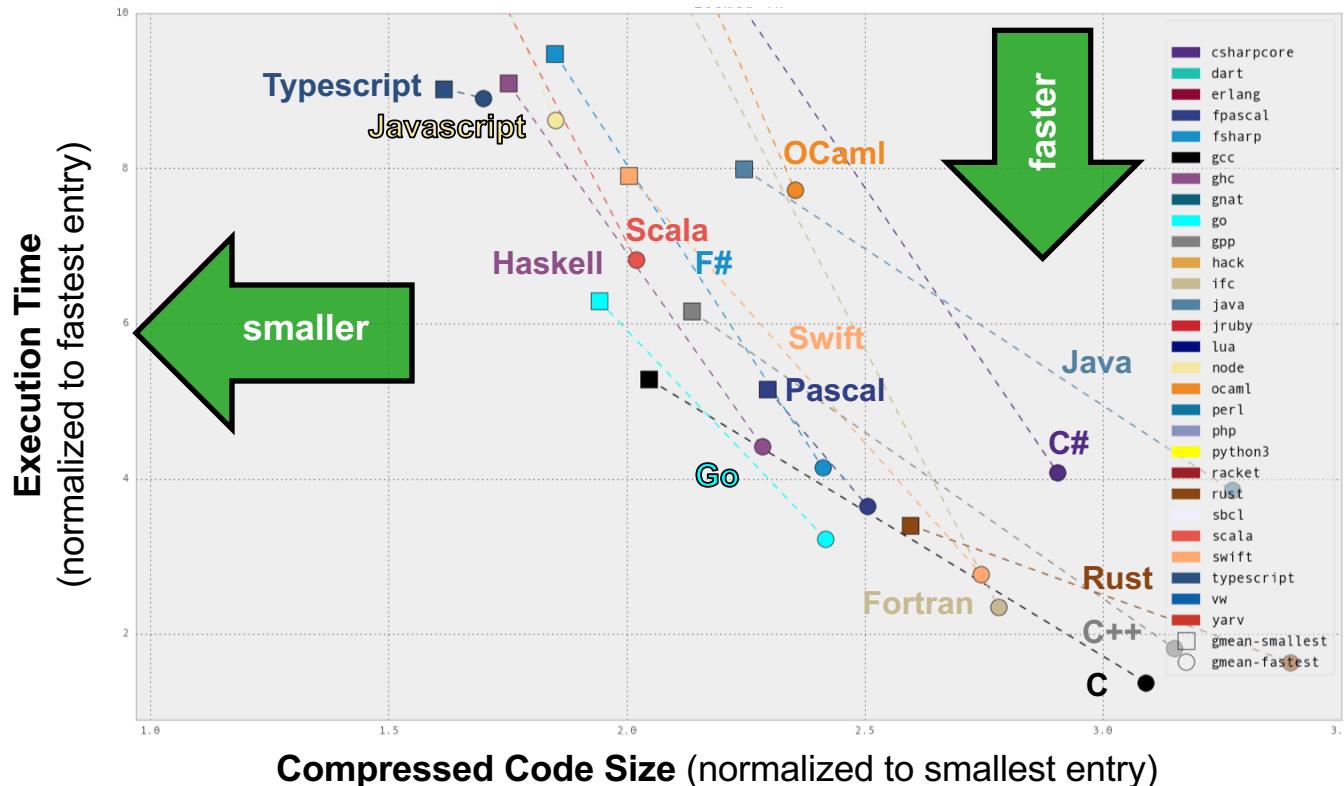
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CLBG Cross-Language Summary

(Oct 2017 standings, zoomed in)



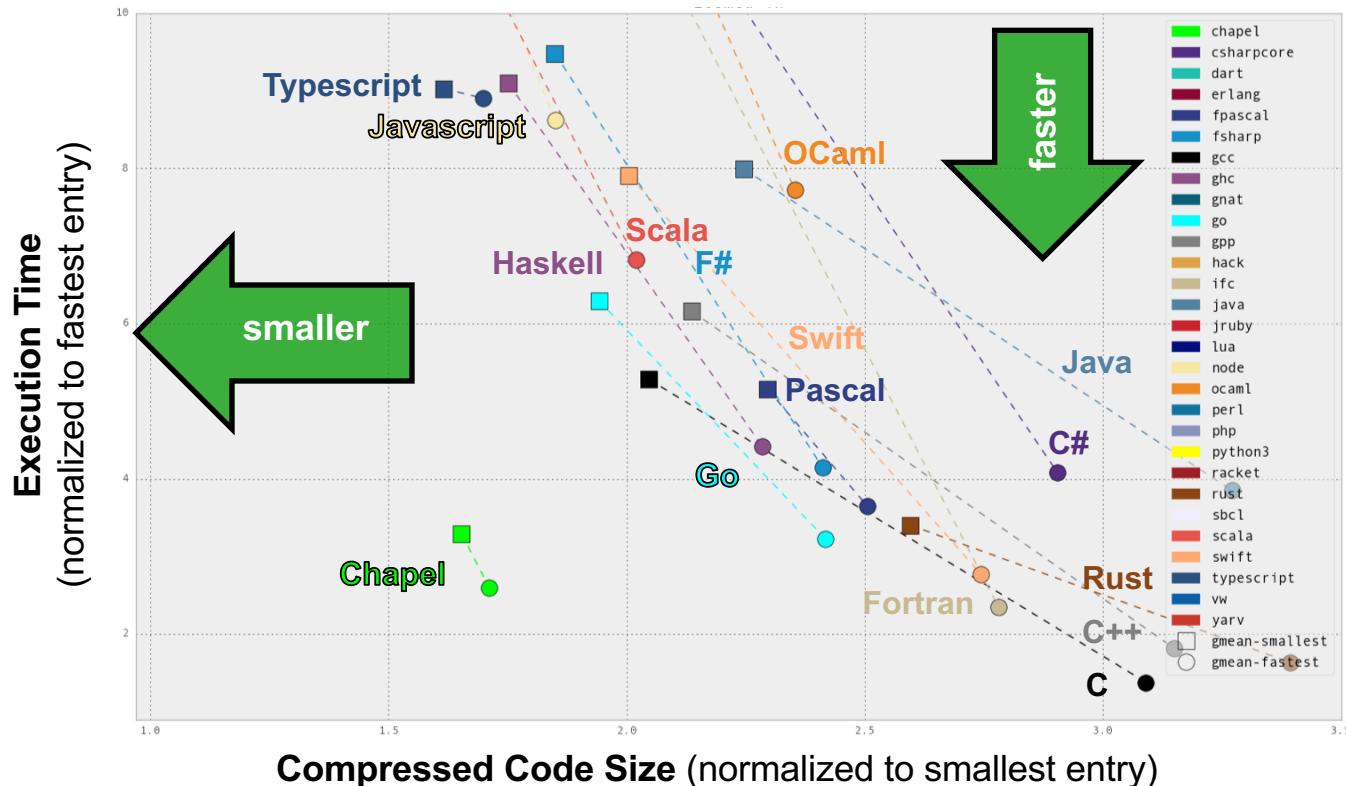
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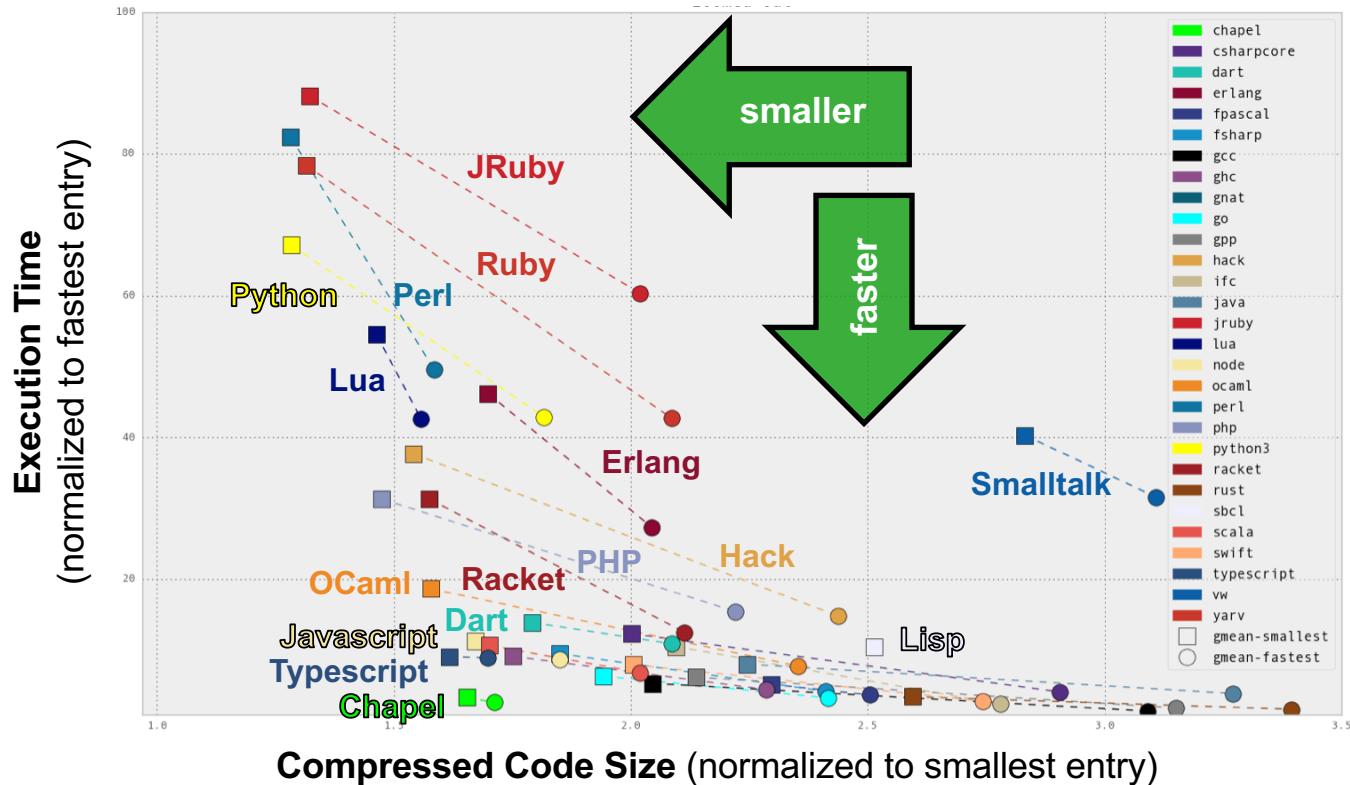
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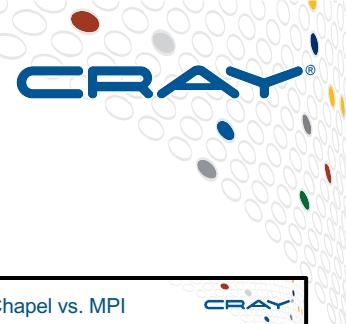


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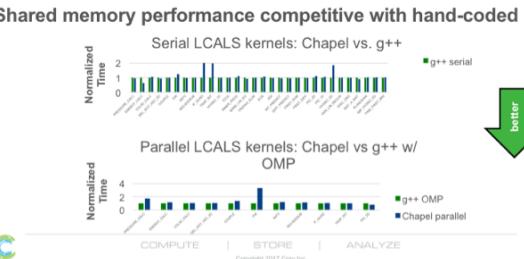
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Chapel Performance: HPC Benchmarks



LCALS: Chapel vs. C + OpenMP



LCALS

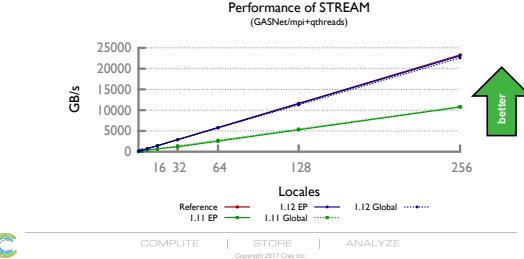
HPCC RA

STREAM
Triad

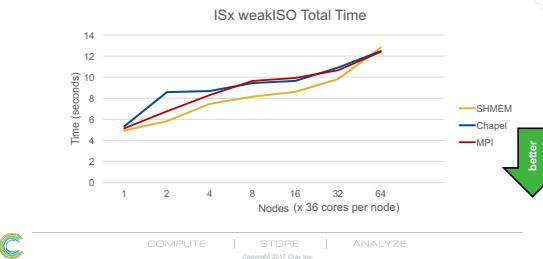
ISx

PRK
Stencil

HPCC Stream Triad: Chapel vs. MPI+OpenMP



ISx Performance: Chapel vs. MPI, SHMEM



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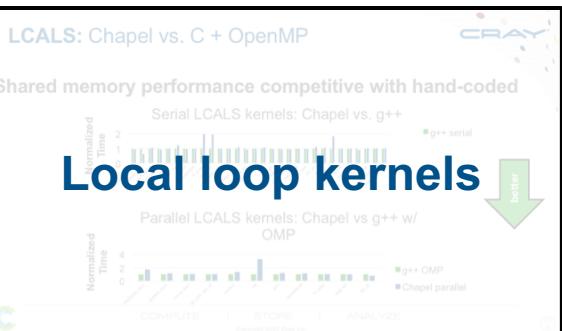
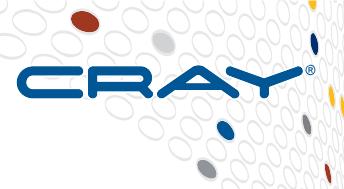
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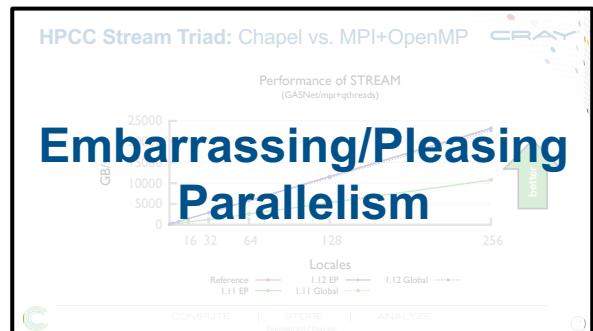
Nightly performance graphs online
at: <https://chapel-lang.org/perf>

Chapel Performance: HPC Benchmarks



LCALS

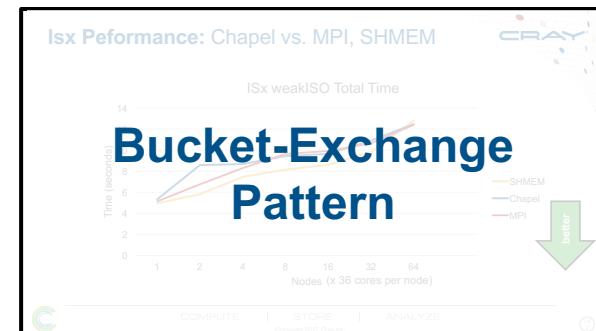
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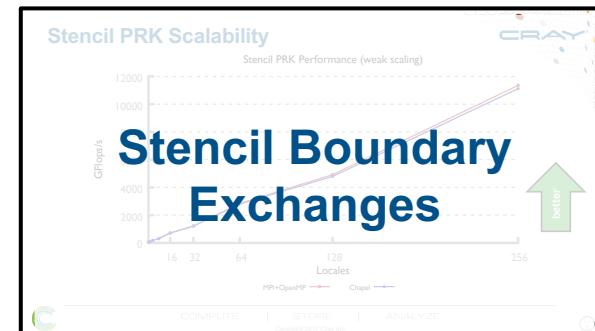
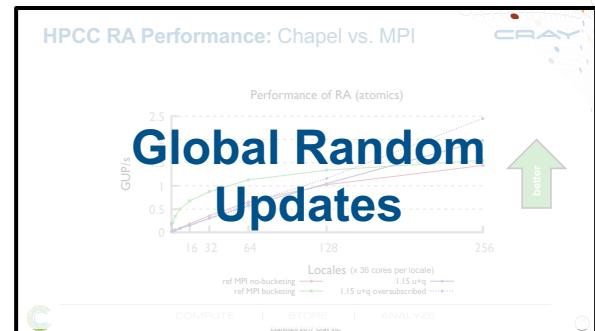


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Nightly performance graphs online
at: <https://chapel-lang.org/perf>

The Chapel Team at Cray (May 2017)



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Chapel Community Partners



Lawrence Berkeley
National Laboratory



Yale

(and several others...)

<https://chapel-lang.org/collaborations.html>



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Tonight's Plan



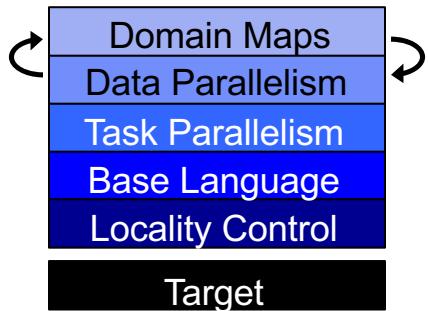
- **Cover features that we haven't in this forum before**
 - base language features of potential interest to C++ users
 - multiresolution features for user control over parallel abstractions
 - parallel iterators
 - domain maps
 - locale models
- **Review core features along the way**
 - goal: quicker than in previous talks
 - help refresh memories / bring new attendees up-to-speed
- **Please ask questions as we go**



Chapel language feature areas



Chapel language concepts



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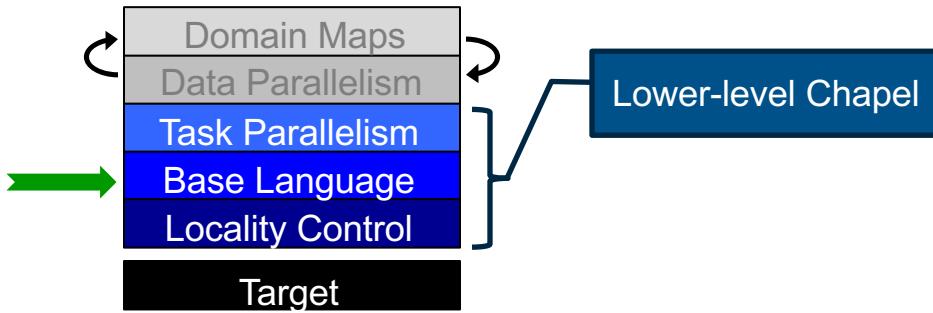
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Base Language



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Base Language Features, by example



```
iter fib(n) {
    var current = 0,
        next = 1;

    for i in 1..n {
        yield current;
        current += next;
        current <=gt; next;
    }
}
```

```
config const n = 10;

for f in fib(n) do
    writeln(f);
```

```
0  
1  
1  
2  
3  
5  
8  
...
```



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Base Language Features, by example



Configuration declarations
(to avoid command-line argument parsing)
./a.out --n=1000000

```
iter fib(n) {
    var current = 0,
        next = 1;

    for i in 1..n {
        yield current;
        current += next;
        current <=> next;
    }
}
```

```
config const n = 10;

for f in fib(n) do
    writeln(f);
```

```
0
1
1
2
3
5
8
...
...
```



Base Language Features, by example



Modern iterators

```
iter fib(n) {
    var current = 0,
        next = 1;

    for i in 1..n {
        yield current;
        current += next;
        current <= n;
    }
}
```

```
config const n = 10;

for f in fib(n) do
    writeln(f);
```

```
0  
1  
1  
2  
3  
5  
8  
...
```

Base Language Features, by example



Static type inference for:

- arguments
- return types
- variables

```
iter fib(n) {
    var current = 0,
        next = 1;

    for i in 1..n {
        yield current;
        current += next;
        current <= next;
    }
}
```

```
config const n = 10;

for f in fib(n) do
    writeln(f);
```

```
0  
1  
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8  
...
```

Base Language Features, by example



```
iter fib(n) {
    var current = 0,
        next = 1;

    for i in 1..n {
        yield current;
        current += next;
        current <= next;
    }
}
```

```
config const n = 10;

for (i,f) in zip(0..#n, fib(n)) do
    writeln("fib #", i, " is ", f);
```

Zippered iteration

```
fib #0 is 0
fib #1 is 1
fib #2 is 1
fib #3 is 2
fib #4 is 3
fib #5 is 5
fib #6 is 8
...
```



Base Language Features, by example



Range types and operators

```
iter fib(n) {
    var current = 0,
        next = 1;

    for i in 1..n {
        yield current;
        current += next;
        current <= next;
    }
}
```

```
config const n = 10;

for (i,f) in zip(0..#n, fib(n)) do
    writeln("fib #", i, " is ", f);
```

```
fib #0 is 0
fib #1 is 1
fib #2 is 1
fib #3 is 2
fib #4 is 3
fib #5 is 5
fib #6 is 8
...
```



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Base Language Features, by example



tuples

```
iter fib(n) {
    var current = 0,
        next = 1;

    for i in 1..n {
        yield current;
        current += next;
        current <= next;
    }
}
```

```
config const n = 10;

for (i,f) in zip(0..#n, fib(n)) do
    writeln("fib #", i, " is ", f);
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```
fib #0 is 0
fib #1 is 1
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Base Language Features, by example



```
iter fib(n) {
    var current = 0,
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    for i in 1..n {
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        current += next;
        current <= next;
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config const n = 10;

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```
fib #0 is 0
fib #1 is 1
fib #2 is 1
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fib #4 is 3
fib #5 is 5
fib #6 is 8
```

...



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Other Base Language Features of Potential Interest to C++ Users



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Other Base Language Features: OOP



Two flavors of object-oriented types:

- Value-based:

```
record Circle {  
    var radius: real;  
    proc area() { ... }  
}
```

```
var myCircle = new Circle(radius = 1.0),  
    myCircle2 = myCircle; // copy for record, alias for class  
myCircle.radius = 2.0;  
writeln(myCircle2.area()); // 1.0 for record, 4.0 for class
```

- Reference-based:

```
class Circle {  
    var radius: real;  
    proc area() { ... }  
}
```



Other Base Language Features: Generics

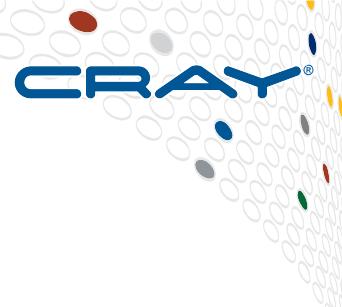


- Support for generic types and functions
 - w.r.t. types and statically known values (`param's)

```
class Arr {  
    param numDims: int;           // number of dimensions  
    type eltType;                // element type  
    var size: numDims*int;        // tuple storing per-dimension size  
}  
  
var myArr = new Arr(2, string, (100, 200)),  
    myArr2 = new Arr(3, real, (500, 500, 500));
```



Other Base Language Features: Generics



- Support for generic types and functions

- w.r.t. types and statically known values (`param's)

```
proc mypow(type t, x: t, param exponent: int) {  
    var result = 1:t;  
    for param i in 1..exponent do  
        result *= x;  
    return result;  
}
```

note: this is an utterly artificial and over-engineered way to write this function in Chapel, done merely to demonstrate type/param args in ~6 lines...

```
var twoSquared = mypow(int, 2, 2);  
var piCubed = mypow(real, 3.14159265, 3);
```



Other Base Language Features: Meta-Programming



- **Compile-time procedures to compute types / params**

```
proc computePacketSize(type t1, type t2) param {
    return numBits(t1) + numBits(t2);
}

proc c_intToChapelInt() type {
    return int(numBits(c_int));
}
```

- Also, support for config types / params

```
config param bitsPerInt = 16;
config type eltType = int(bitsPerInt);
```

```
chpl -sbitsPerInt=64 -seltType=real(32) myProg.chpl
```



Other Base Language Features: Meta-Programming



- Ability to unroll loops / fold conditionals or ‘void’ exprs

```
for param i in (1, 2.3, "hello", (5,7)) do  
    writeln("i: ", i, " has type: ", i.type:string);
```

- Reflection module:

- “Can this function / method be resolved”
- “Iterate over all fields in this record giving me their names / types”
- ...



Other Base Language Features



- Error-handling
- Modules (namespaces)
- Overloading, filtering
- Default args, arg intents, keyword-based arg passing
- Argument type queries / pattern-matching
- ...



Base Language Features: What's Missing?



- **better initializer (constructor) features**
 - currently being implemented and refined
- **delete-free programming / borrow-checking**
 - currently being designed and implemented
- **first-class functions**
 - prototyped, need strengthening
- **constrained generics / interfaces / concepts**
 - proposal drafted but not implemented
- **anti-function hijacking features**
 - currently under consideration



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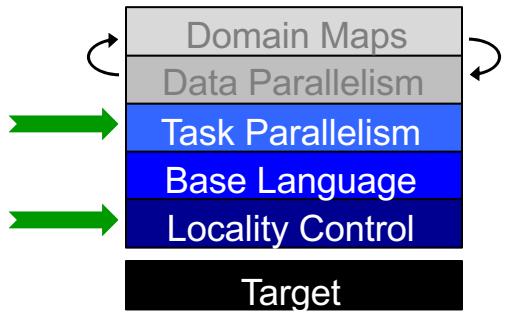
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Task Parallelism and Locality Control



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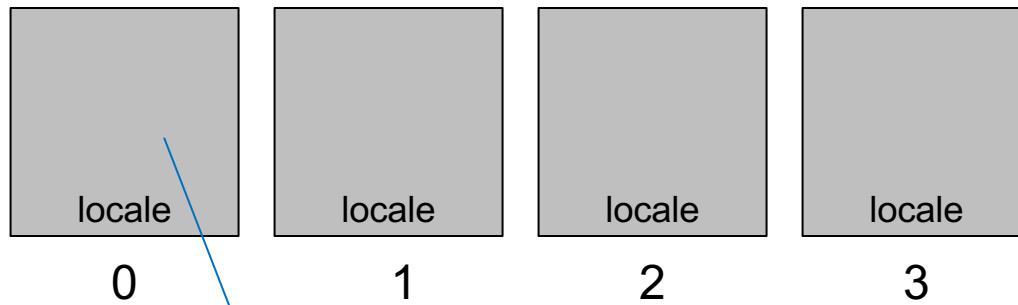
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Locales



- Unit of the target system useful for reasoning about locality
 - Each locale can run tasks and store variables
 - Has processors and memory (or can defer to something that does)
 - For most HPC systems, locale == compute node

Locales :



User's main() executes on locale #0



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Task Parallelism and Locality, by example



taskParallel.chpl

```
coforall loc in Locales do
    on loc {
        const numTasks = here.numPUs();
        coforall tid in 1..numTasks do
            writef("Hello from task %n of %n "+
                "running on %s\n",
                tid, numTasks, here.name);
    }
```

```
prompt> chpl taskParallel.chpl -o taskParallel
prompt> ./taskParallel --numLocales=2
Hello from task 1 of 2 running on n1033
Hello from task 2 of 2 running on n1032
Hello from task 2 of 2 running on n1033
Hello from task 1 of 2 running on n1032
```

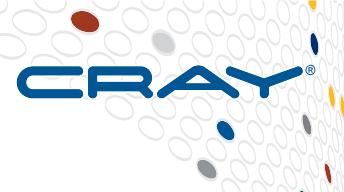


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Task Parallelism and Locality, by example



Abstraction of
System Resources

taskParallel.chpl

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Task Parallelism and Locality, by example



High-Level
Task Parallelism

taskParallel.chpl

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Task Parallelism and Locality, by example



Control of Locality/Affinity

taskParallel.chpl

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Task Parallelism and Locality, by example



Abstraction of
System Resources

taskParallel.chpl

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Task Parallelism and Locality, by example



High-Level
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Task Parallelism and Locality, by example



taskParallel.chpl

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        coforall tid in 1..numTasks do
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                "running on %s\n",
                tid, numTasks, here.name);
    }
```

Not seen here:
Data-centric task coordination
via atomic and full/empty vars

```
prompt> chpl taskParallel.chpl -o taskParallel
prompt> ./taskParallel --numLocales=2
Hello from task 1 of 2 running on n1033
Hello from task 2 of 2 running on n1032
Hello from task 2 of 2 running on n1033
Hello from task 1 of 2 running on n1032
```



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Task Parallelism and Locality, by example



taskParallel.chpl

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```

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prompt> chpl taskParallel.chpl -o taskParallel
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Hello from task 1 of 2 running on n1033
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```



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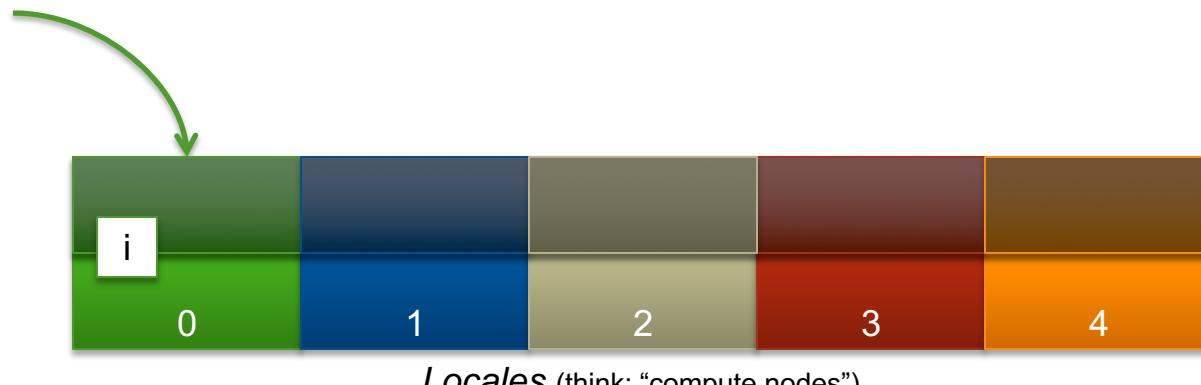
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Chapel: Scoping and Locality



```
var i: int;
```



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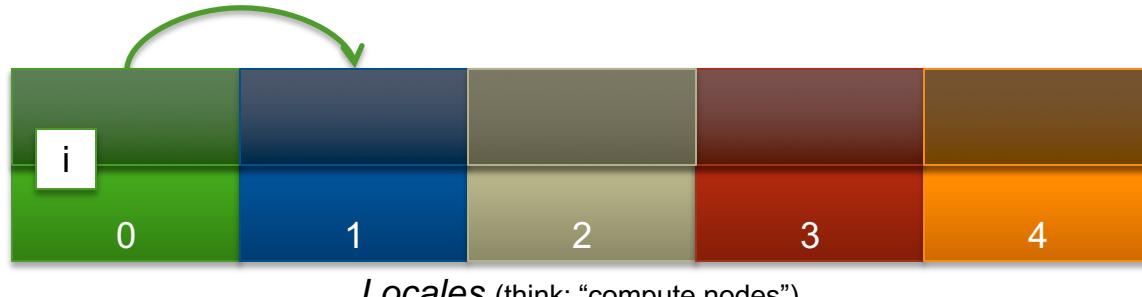
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Chapel: Scoping and Locality



```
var i: int;  
on Locales[1] {
```



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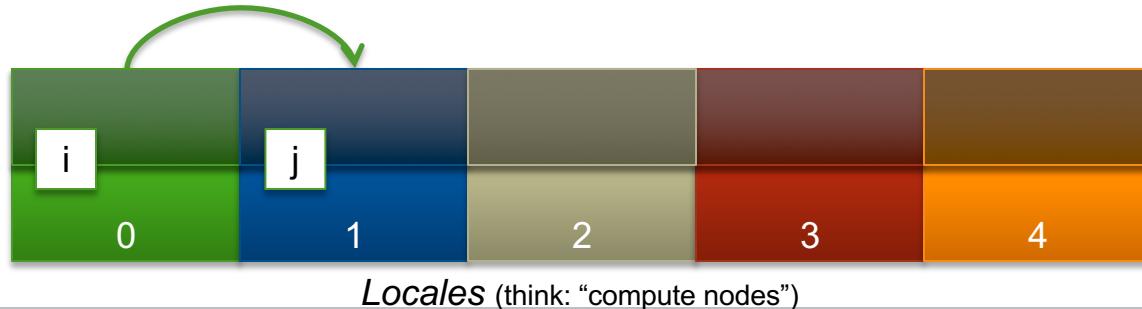
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Chapel: Scoping and Locality



```
var i: int;  
on Locales[1] {  
    var j: int;
```



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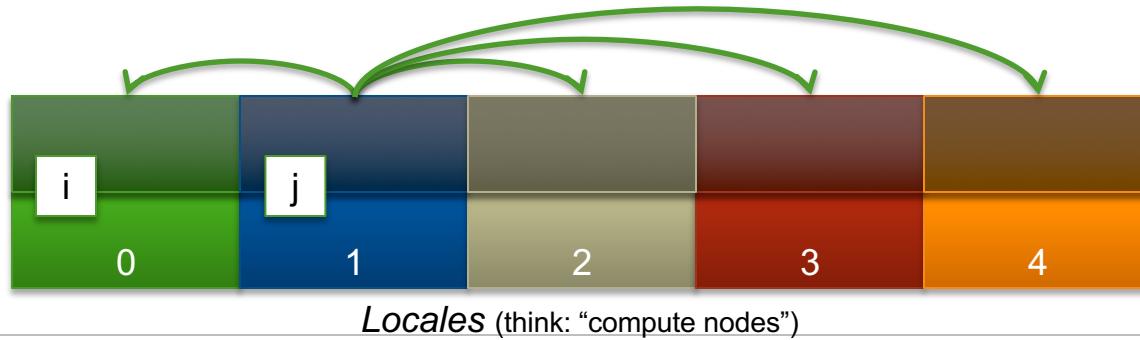
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Chapel: Scoping and Locality



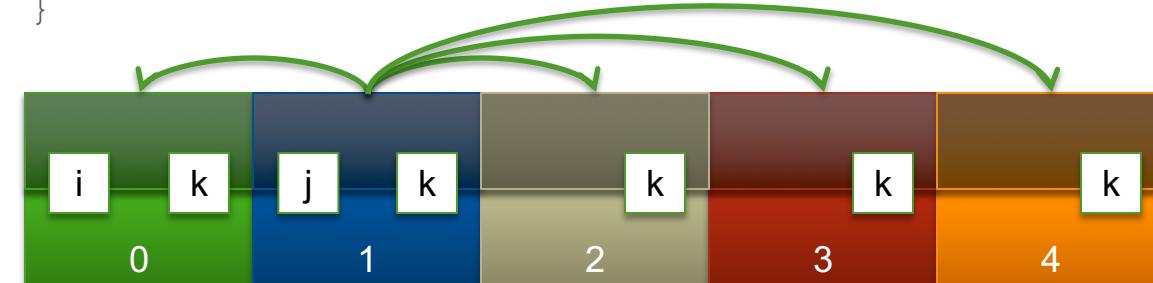
```
var i: int;  
on Locales[1] {  
    var j: int;  
    coforall loc in Locales {  
        on loc {
```



Chapel: Scoping and Locality



```
var i: int;  
on Locales[1] {  
    var j: int;  
    coforall loc in Locales {  
        on loc {  
            var k: int;  
            ...  
        }  
    }  
}
```



Locales (think: “compute nodes”)

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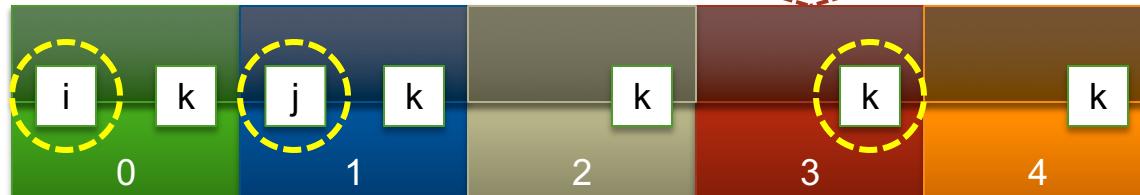
Chapel: Scoping and Locality



```
var i: int;  
on Locales[1] {  
    var j: int;  
    coforall loc in Locales {  
        on loc {  
            var k: int;  
            k = 2*i + j;  
        }  
    }  
}
```

OK to access i , j , and k
wherever they live

$k = 2*i + j;$



Locales (think: “compute nodes”)

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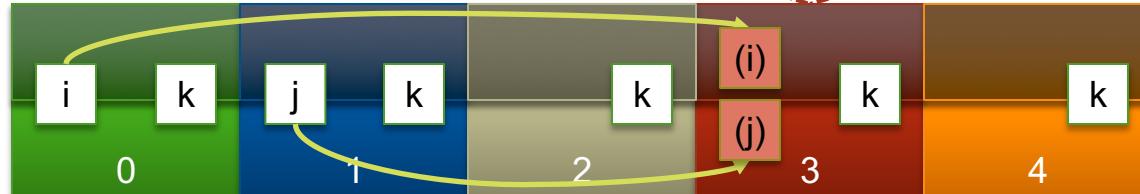


Chapel: Scoping and Locality



```
var i: int;  
on Locales[1] {  
    var j: int;  
    coforall loc in Locales {  
        on loc {  
            var k: int;  
            k = 2*i + j;  
        }  
    }  
}
```

here, *i* and *j* are remote, so
the compiler + runtime will
transfer their values



Locales (think: “compute nodes”)

COMPUTE

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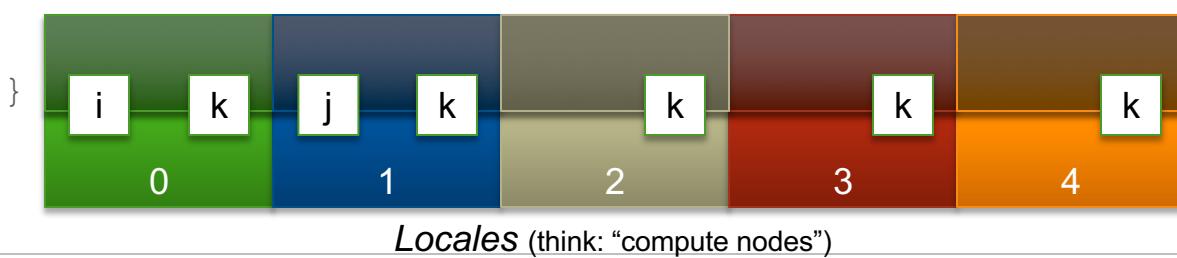
ANALYZE



Chapel: Locality queries



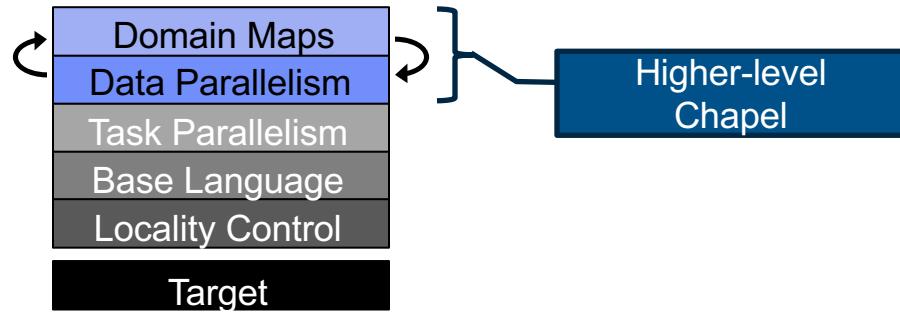
```
var i: int;  
on Locales[1] {  
    var j: int;  
    coforall loc in Locales {  
        on loc {  
            ...here...           // query the locale on which this task is running  
            ...j.locale...      // query the locale on which j is stored  
            ...here.physicalMemory (...) ... // query system characteristics  
            ...here.runningTasks () ...    // query runtime characteristics  
        }  
    }  
}
```



Data Parallelism in Chapel



Chapel language concepts



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Data Parallelism, by example



dataParallel.chpl

```
config const n = 1000;
var D = {1..n, 1..n};

var A: [D] real;
forall (i,j) in D do
    A[i,j] = i + (j - 0.5)/n;
writeln(A);
```

```
prompt> chpl dataParallel.chpl -o dataParallel
prompt> ./dataParallel --n=5
1.1 1.3 1.5 1.7 1.9
2.1 2.3 2.5 2.7 2.9
3.1 3.3 3.5 3.7 3.9
4.1 4.3 4.5 4.7 4.9
5.1 5.3 5.5 5.7 5.9
```



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Data Parallelism, by example



Domains (Index Sets)

dataParallel.chpl

```
config const n = 1000;
var D = {1..n, 1..n};

var A: [D] real;
forall (i,j) in D do
    A[i,j] = i + (j - 0.5)/n;
writeln(A);
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3.1 3.3 3.5 3.7 3.9
4.1 4.3 4.5 4.7 4.9
5.1 5.3 5.5 5.7 5.9
```



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Data Parallelism, by example



Arrays

dataParallel.chpl

```
config const n = 1000;
var D = {1..n, 1..n};

var A: [D] real;
forall (i,j) in D do
    A[i,j] = i + (j - 0.5)/n;
writeln(A);
```

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4.1 4.3 4.5 4.7 4.9
5.1 5.3 5.5 5.7 5.9
```



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Data Parallelism, by example



Data-Parallel Forall Loops

dataParallel.chpl

```
config const n = 1000;
var D = {1..n, 1..n};

var A: [D] real;
forall (i,j) in D do
    A[i,j] = i + (j - 0.5)/n;
writeln(A);
```

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3.1 3.3 3.5 3.7 3.9
4.1 4.3 4.5 4.7 4.9
5.1 5.3 5.5 5.7 5.9
```



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Data Parallelism, by example



This is a shared memory program
Nothing has referred to remote
locales, explicitly or implicitly

dataParallel.chpl

```
config const n = 1000;
var D = {1..n, 1..n};

var A: [D] real;
forall (i,j) in D do
    A[i,j] = i + (j - 0.5)/n;
writeln(A);
```

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3.1 3.3 3.5 3.7 3.9
4.1 4.3 4.5 4.7 4.9
5.1 5.3 5.5 5.7 5.9
```



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Data Parallelism, by example



This is a shared memory program
Nothing has referred to remote
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dataParallel.chpl

```
config const n = 1000;
var D = {1..n, 1..n};

var A: [D] real;
forall (i,j) in D do
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3.1 3.3 3.5 3.7 3.9
4.1 4.3 4.5 4.7 4.9
5.1 5.3 5.5 5.7 5.9
```



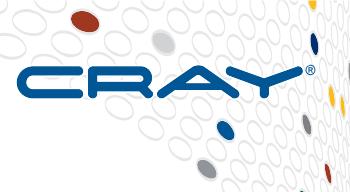
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Distributed Data Parallelism, by example



Domain Maps
(Map Data Parallelism to the System)

dataParallel.chpl

```
use CyclicDist;
config const n = 1000;
var D = {1..n, 1..n}
        dmapped Cyclic(startIdx = (1,1));
var A: [D] real;
forall (i,j) in D do
    A[i,j] = i + (j - 0.5)/n;
writeln(A);
```

```
prompt> chpl dataParallel.chpl -o dataParallel
prompt> ./dataParallel --n=5 --numLocales=4
1.1 1.3 1.5 1.7 1.9
2.1 2.3 2.5 2.7 2.9
3.1 3.3 3.5 3.7 3.9
4.1 4.3 4.5 4.7 4.9
5.1 5.3 5.5 5.7 5.9
```



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Distributed Data Parallelism, by example



magic!
HPF-like!
descriptive!

Not at all...

- Lowering of code is well-defined
- User can control details
- Part of Chapel's *multiresolution philosophy*...

dataParallel.chpl

```
use CyclicDist;
config const n = 1000;
var D = {1..n, 1..n}
    dmapped Cyclic(startIdx = (1,1));
var A: [D] real;
forall (i,j) in D do
    A[i,j] = i + (j - 0.5)/n;
writeln(A);
```

```
prompt> chpl dataParallel.chpl -o dataParallel
prompt> ./dataParallel --n=5 --numLocales=4
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3.1 3.3 3.5 3.7 3.9
4.1 4.3 4.5 4.7 4.9
5.1 5.3 5.5 5.7 5.9
```

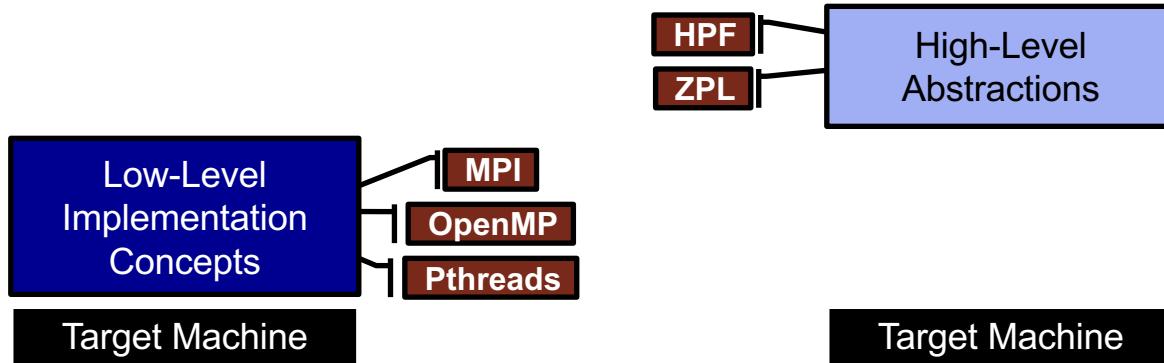


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Chapel's Multiresolution Design: Motivation



“Why is everything so tedious/difficult?”

*“Why don’t my programs trivially port
to new systems?”*

“Why don’t I have more control?”

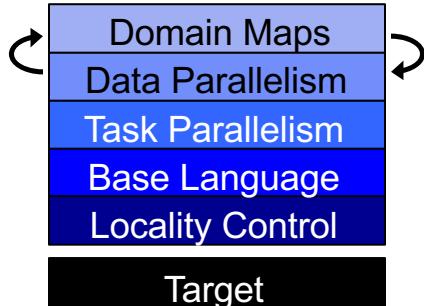


Chapel's Multiresolution Philosophy



Multiresolution Design: Support multiple tiers of features

- higher levels for programmability, productivity
- lower levels for greater degrees of control



- build the higher-level concepts in terms of the lower
- permit users to intermix layers arbitrarily



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Distributed Data Parallelism, by example



Chapel's prescriptive approach:

```
forall (i,j) in D do...
```

- ⇒ invoke and inline D's default parallel iterator
- defined by D's type / domain map

default domain map

- create a task per local core
- block indices across tasks

dataParallel.chpl

```
config const n = 1000;
var D = {1..n, 1..n};

var A: [D] real;
forall (i,j) in D do
    A[i,j] = i + (j - 0.5)/n;
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```

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prompt> chpl dataParallel.chpl -o dataParallel
prompt> ./dataParallel --n=5 --numLocales=4
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3.1 3.3 3.5 3.7 3.9
4.1 4.3 4.5 4.7 4.9
5.1 5.3 5.5 5.7 5.9
```



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Distributed Data Parallelism, by example



Chapel's prescriptive approach:

```
forall (i,j) in D do...
```

- ⇒ invoke and inline D's default parallel iterator
 - defined by D's type / domain map

default domain map cyclic domain map

- on each target locale...
 - create a task per core
 - block local indices across tasks

dataParallel.chpl

```
use CyclicDist;
config const n = 1000;
var D = {1..n, 1..n}
        dmapped Cyclic(startIdx = (1,1));
var A: [D] real;
forall (i,j) in D do
    A[i,j] = i + (j - 0.5)/n;
writeln(A);
```

```
prompt> chpl dataParallel.chpl -o dataParallel
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3.1 3.3 3.5 3.7 3.9
4.1 4.3 4.5 4.7 4.9
5.1 5.3 5.5 5.7 5.9
```



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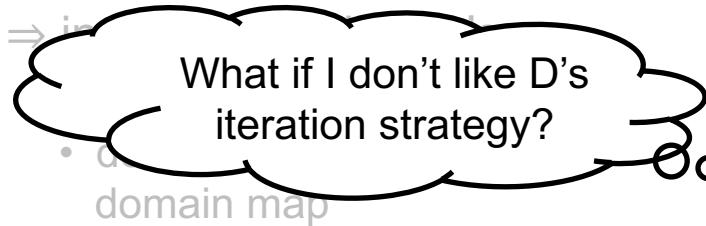
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Distributed Data Parallelism, by example



Chapel's prescriptive approach:

```
forall (i,j) in D do...
```



dataParallel.chpl

```
use CyclicDist;
config const n = 1000;
var D = {1..n, 1..n}
        dmapped Cyclic(startIdx = (1,1));
var A: [D] real;
forall (i,j) in D do
    A[i,j] = i + (j - 0.5)/n;
```

- Write and call your own parallel iterator:

```
forall (i,j) in myParIter(D) do...
```

- Or, use a different domain map:

```
var D = {1..n, 1..n} dmapped Block(...);
```

- Or, write and use your own domain map:

```
var D = {1..n, 1..n} dmapped MyDomMap(...);
```

```
rallel.chpl -o dataParallel
lel --n=5 --numLocales=4
```



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Write and Call Your Own Parallel Iterator



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Authoring Parallel Iterators



- Similar to serial iterators, but invoked by `forall` loops
 - Unlike serial iterators, these can contain parallel constructs

```
forall i in myParIter(D) { ... }
```

invokes:

```
iter myParIter(dom: domain, ... /* tag as a parallel iterator */) {  
    coforall tid in 0..#numTasks {  
        const myChunk = computeChunk(dom, tid, numTasks);  
        for i in myChunk do  
            yield i;  
    }  
}
```



Authoring Zippered Parallel Iterators



- Parallel iterators can also support zippered iteration

```
forall (i,j) in zip(myParIter(D), A) { ... }
```

- defined in terms of leader...

```
iter myParIter(dom: domain, ...) {
    coforall tid in 0..#numTasks do
        yield computeChunk(dom, tid, numTasks);
}
```

...and follower iterators:

```
iter myParIter(dom: domain, followThis, ...) {
    for i in followThis do yield i;
}
```



Use a Different Domain Map



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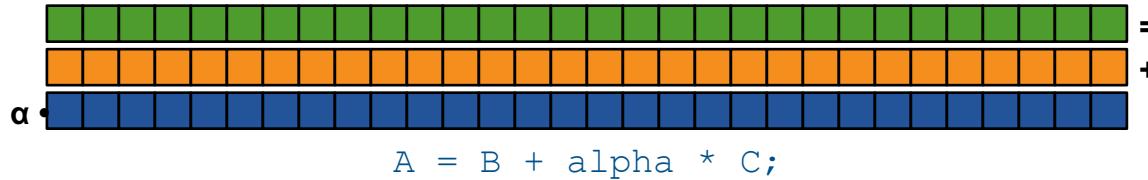
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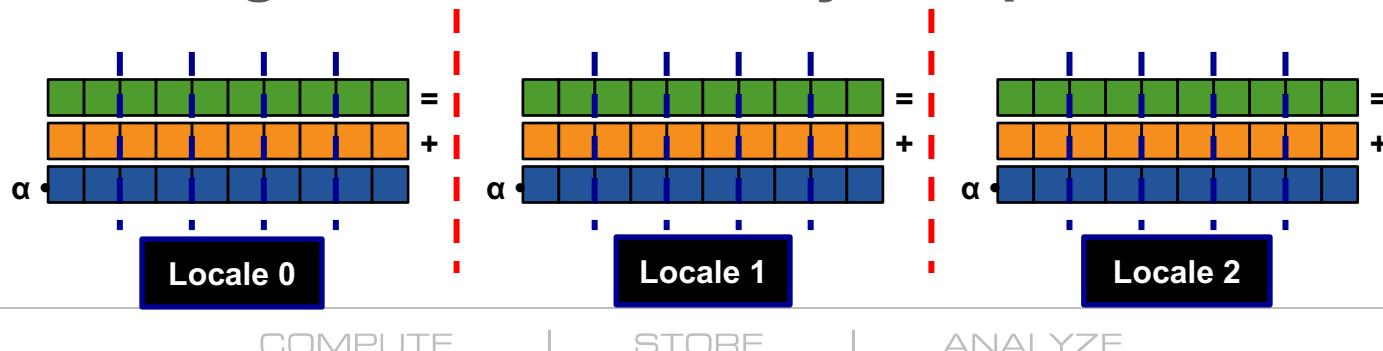
Domain Maps: A Multiresolution Feature



Domain maps are “recipes” that instruct the compiler how to map the global view of a computation...



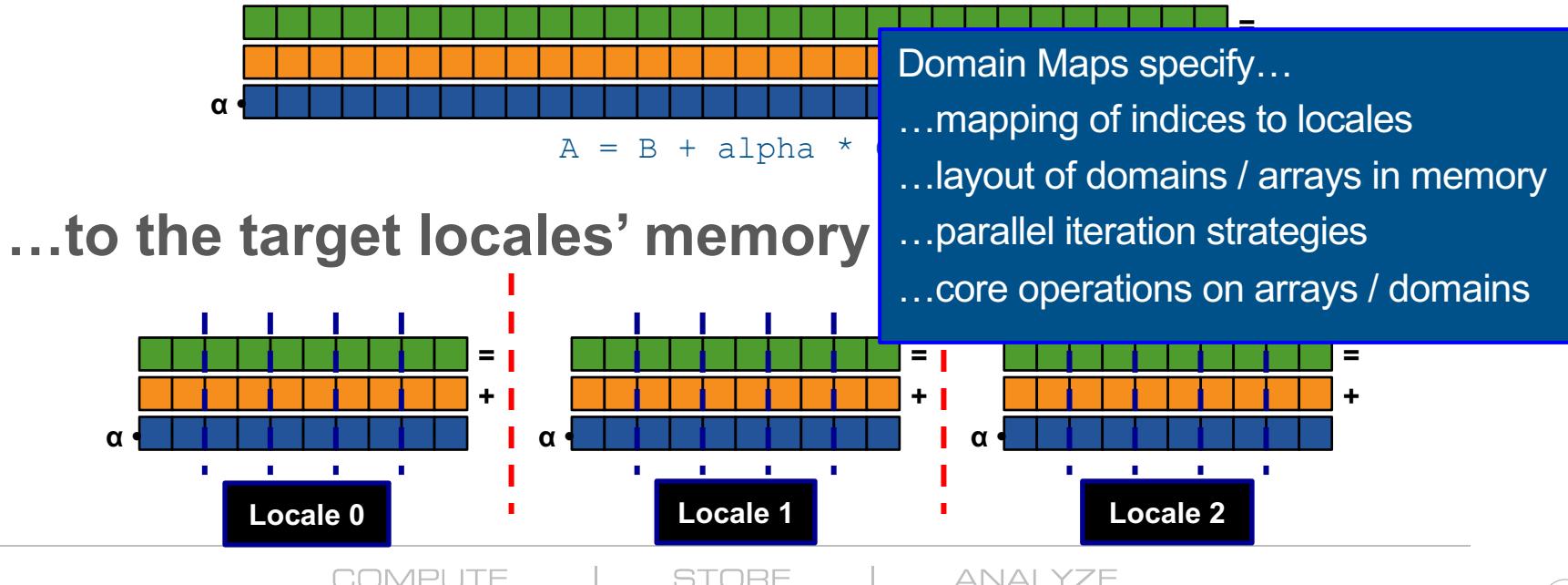
...to the target locales' memory and processors:



Domain Maps: A Multiresolution Feature



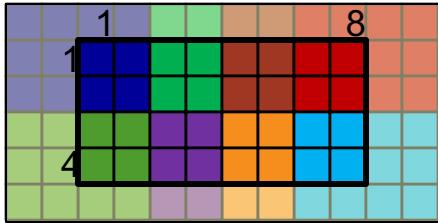
Domain maps are “recipes” that instruct the compiler how to map the global view of a computation...



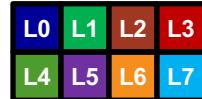
Sample Domain Maps: Block and Cyclic



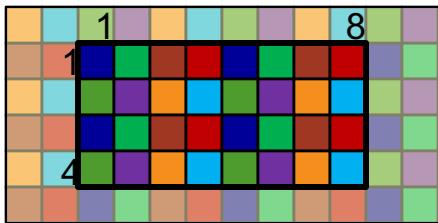
```
var Dom = {1..4, 1..8} dmapped Block( {1..4, 1..8} );
```



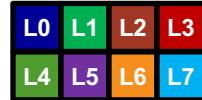
distributed to



```
var Dom = {1..4, 1..8} dmapped Cyclic( startIdx=(1,1) );
```



distributed to



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Write and Use Your Own Domain Map



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Chapel's Domain Map Philosophy

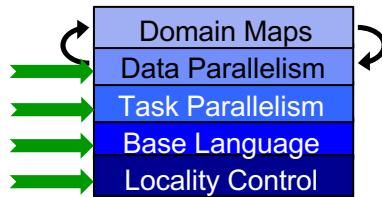


1. Chapel provides a library of standard domain maps

- to support common array implementations effortlessly

2. Expert users can write their own domain maps in Chapel

- to cope with any shortcomings in our standard library



3. Chapel's standard domain maps are written using the end-user framework

- to avoid a performance cliff between “built-in” and user-defined cases
- in fact every Chapel array is implemented using this framework

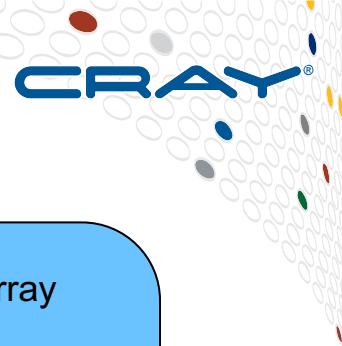


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Domain Map Descriptors



Domain Map

Represents: a domain map value
Generic w.r.t.: index type
State: the domain map's representation
Typical Size: $\Theta(1) \rightarrow \Theta(\text{numLocales})$

Required Interface:

- create new domains
- which locale owns index i ?

Domain

Represents: a domain
Generic w.r.t.: index type
State: representation of index set
Typical Size: $\Theta(1) \rightarrow \Theta(\text{numIndices})$

Required Interface:

- create new arrays
- queries: size, members
- iterators: serial, parallel
- domain assignment
- index set operations

Array

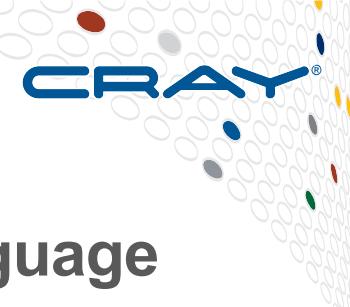
Represents: an array
Generic w.r.t.: index type, element type
State: array elements
Typical Size: $\Theta(\text{numIndices})$

Required Interface:

- (re-)allocation of elements
- random access
- iterators: serial, parallel
- get/set of sparse “zero” values
- ...



Chapel and Performance Portability



- **Avoid locking key policy decisions into the language**
 - Array memory layout?
 - Sparse storage format?
 - Parallel loop policies?



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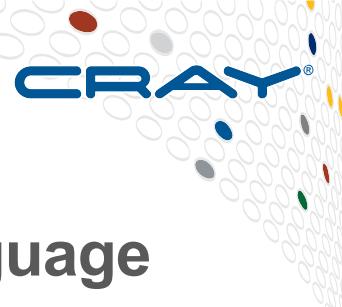
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Chapel and Performance Portability



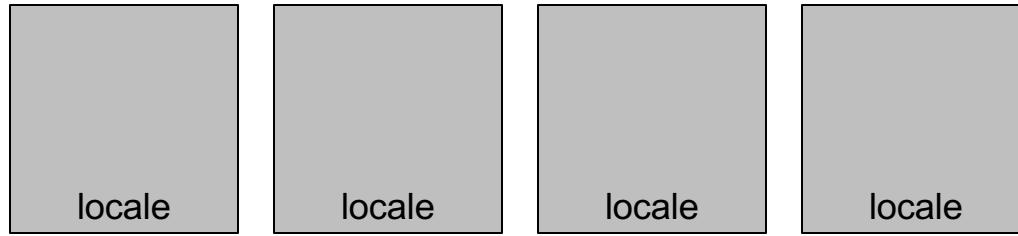
- **Avoid locking key policy decisions into the language**
 - Array memory layout? **not defined by Chapel**
 - Sparse storage format? **not defined by Chapel**
 - Parallel loop policies? **not defined by Chapel**
 - Abstract node architecture? **not defined by Chapel**
- **Instead, permit users to specify these in Chapel itself**
 - support performance portability through...
 - ...a separation of concerns
 - ...abstractions—known to the compiler, and therefore optimizable
 - **goal:** make Chapel a future-proof language



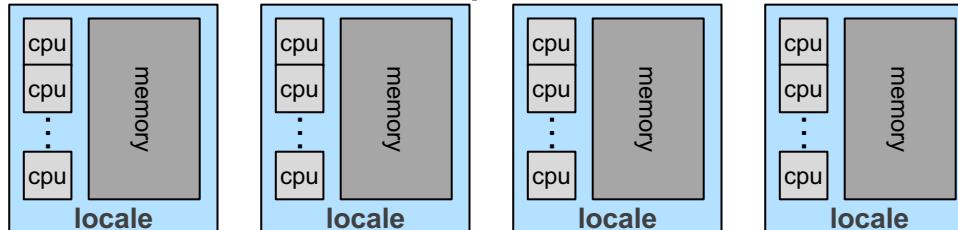
Classic Locales



- Historically, Chapel's locales were black boxes
 - Intra-node concerns handled by compiler, runtime, OS



- This was sufficient when compute nodes were simple



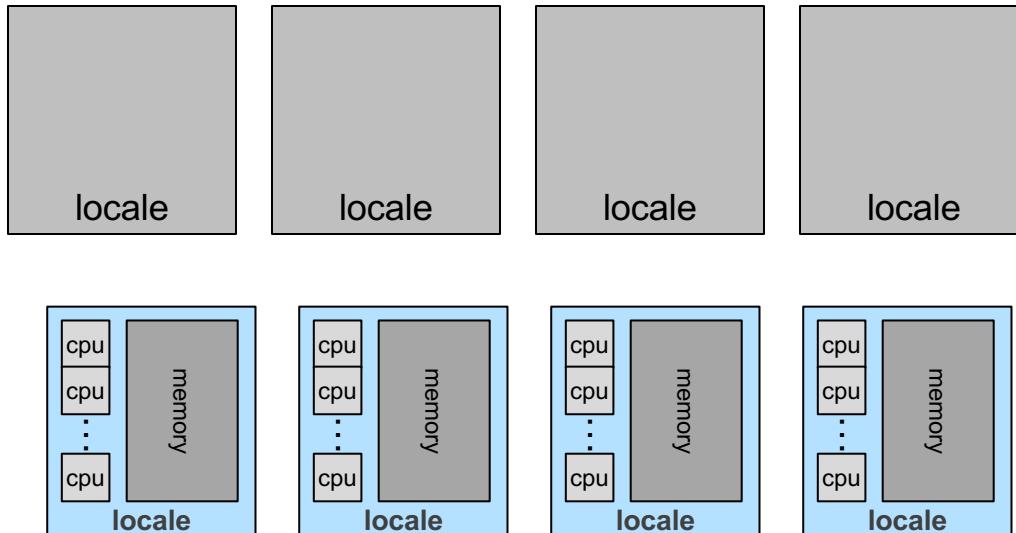
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Classic Locales



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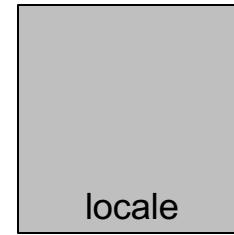
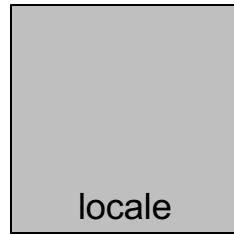
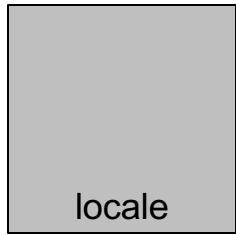
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Classic Locales



- Classic model breaks down for more complex cases
 - E.g. multiple flavors of memory or processors

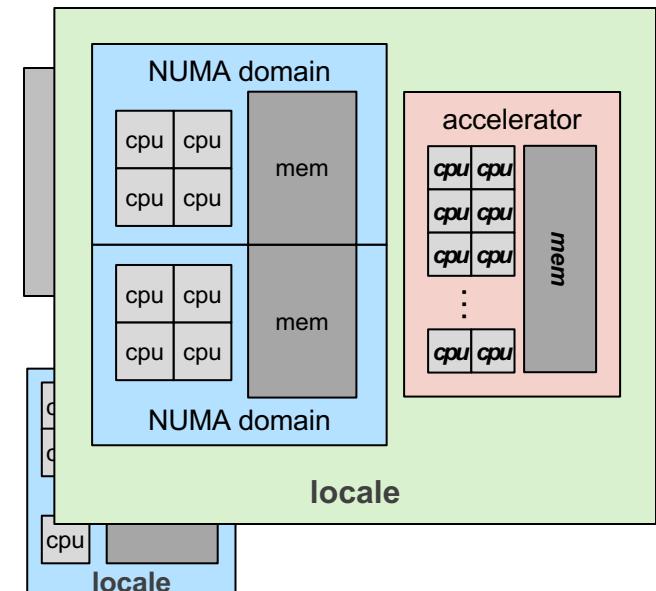


Could hope compilers will “simply get smart enough”

...but seems naïve and doesn't match Chapel's philosophy



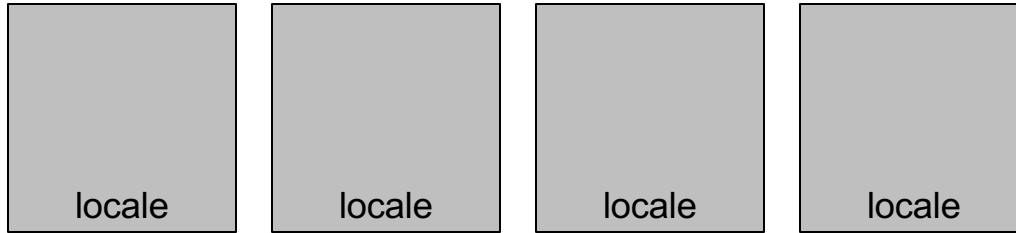
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Hierarchical Locales



- So, we made locales hierarchical



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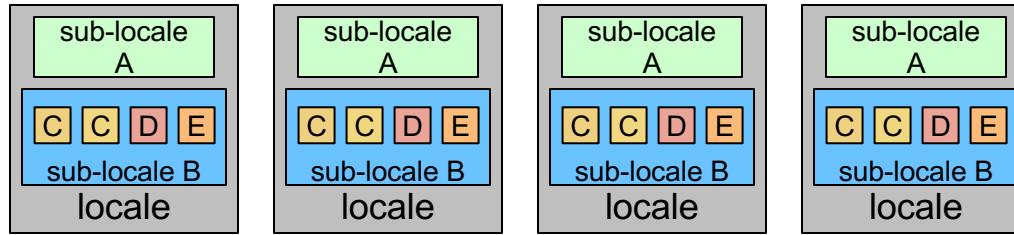
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Hierarchical Locales

- So, we made locales hierarchical
 - Locales can now themselves contain locales
 - E.g., an accelerator sub-locale, a scratchpad memory sub-locale



- Target sub-locales with on-clauses, as before
 - `on Locales[0].GPU do computationThatLikesGPUs();`
 - Ideally, hide such logic in abstractions: domain maps, parallel iterators
- Introduced a new multiresolution type: *locale models*

Chapel's Locale Models



- User-specified type representing locales
- Similar goals to domain maps:
 - Support user implementation of key high-level abstractions
 - Make language future-proof (w.r.t. emerging architectures)



Authoring a Locale Model

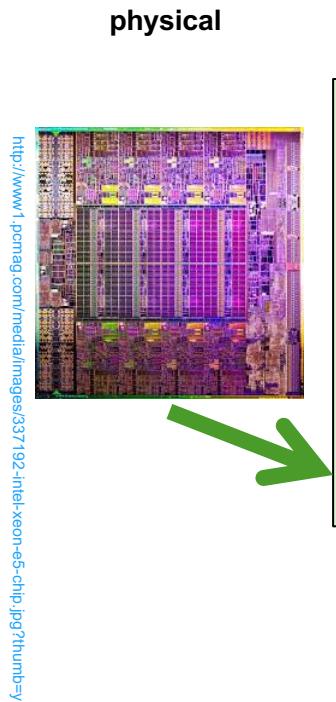
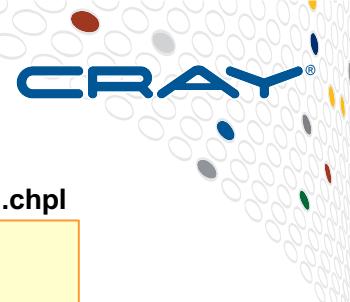


- **Creating a locale model:**

- Create a top-level locale object type
 - In turn, it can contain fields representing sub-locales
- Each locale / sub-locale type must meet a required interface:
 - **Memory:** How is it managed? (malloc, realloc, free)
 - **Tasking:** How do I launch and synchronize tasks?
 - **Communication:** How are data & control transferred between locales?
 - gets, puts, active messages
 - widening of pointers



An Example: The numa Locale Model



\$CHPL_HOME/modules/.../numa/LocaleModel.chpl

```
class NumaDomain : AbstractLocaleModel {
    const sid: chpl_sublocID_t;
}

// The node model
class LocaleModel : AbstractLocaleModel {
    const numSublocales: int;
    var childSpace: domain(1);
    var childLocales: [childSpace] NumaDomain;
}

// support for memory management
proc chpl_here_alloc(size:int, md:int(16)) { ... }

// support for "on" statements
proc chpl_executeOn
    (loc: chpl_localeID_t, // target locale
     fn: int,             // on-body func idx
     args: c_void_ptr,   // func args
     args_size: int(32)   // args size
    ) { ... }

// support for tasking stmts: begin, cobegin, coforall
proc chpl_taskListAddCoStmt
    (subloc_id: int,           // target subloc
     fn: int,                 // body func idx
     args: c_void_ptr,        // func args
     ref tlist: _task_list,   // task list
     tlist_node_id: int       // task list owner
    ) { ... }
```



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Locale Models: Status



- All Chapel compilations use a locale model
 - Set via environment variable or compiler flag
- Current locale models:
 - flat: the default, has no sublocales (as in the classic model)
 - numa: supports a sub-locale per NUMA domain within the node
 - knl: for Intel® Xeon Phi™: numa w/ sublocale for HBM/MCDRAM



Wrapping Up



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Summary

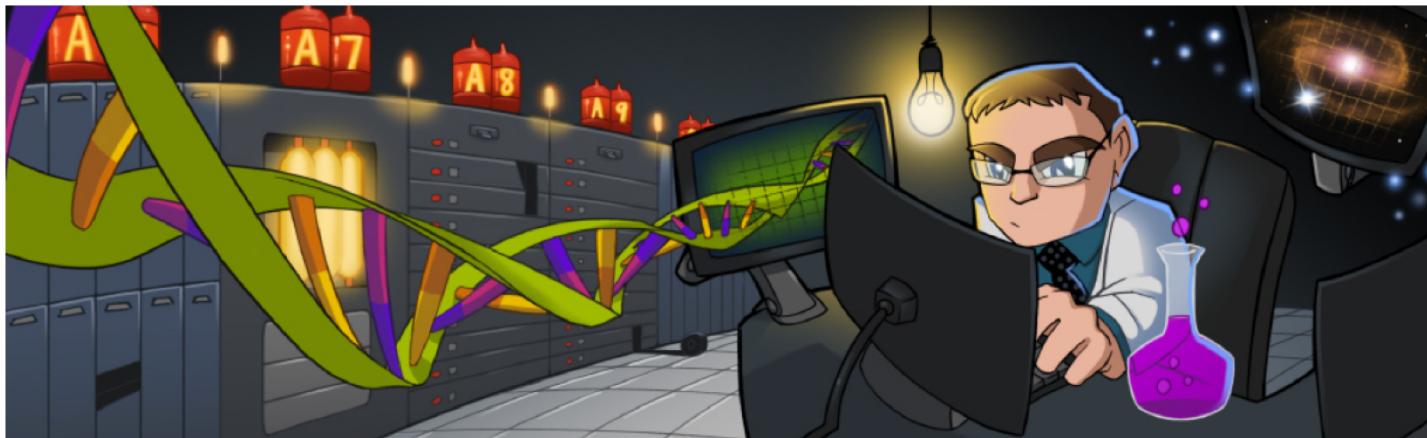


- Chapel's design uses a multiresolution philosophy
 - High-level for productivity
 - Low-level for control
 - User-extensible for flexibility, future-proof design
- Three key examples of multiresolution features:
 - **Parallel iterators:** specify the implementation of forall loops
 - **Domain maps:** specify the implementation of domains and arrays
 - **Locale models:** specify the capabilities of the target architecture



Chapel's Home in the Landscape of New Scientific Computing Languages (and what it can learn from the neighbours)

Jonathan Dursi, *The Hospital for Sick Children, Toronto*



Quote from CHIUW 2017 keynote



“My opinion as an outsider...is that Chapel is important, Chapel is mature, and Chapel is just getting started.

“If the scientific community is going to have frameworks for solving scientific problems that are actually designed for our problems, they’re going to come from a project like Chapel.

“And the thing about Chapel is that the set of all things that are ‘projects like Chapel’ is ‘Chapel.’”

—Jonathan Dursi

Chapel’s Home in the New Landscape of Scientific Frameworks

(and what it can learn from the neighbours)

CHIUW 2017 keynote

<https://ljdursi.github.io/CHIUW2017> / <https://www.youtube.com/watch?v=xj0rwdLOR4U>



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Chapel Resources



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Chapel Central: <https://chapel-lang.org/>





The Chapel Parallel Programming Language

What is Chapel?

Chapel is a modern programming language that is...

- **parallel:** contains first-class concepts for concurrent and parallel computation
- **productive:** designed with programmability and performance in mind
- **portable:** runs on laptops, clusters, the cloud, and HPC systems
- **scalable:** supports locality-oriented features for distributed memory systems
- **open-source:** hosted on [GitHub](#), permissively [licensed](#)

New to Chapel?

As an introduction to Chapel, you may want to...

- read a [blog article](#) or [book chapter](#)
- watch an [overview talk](#) or browse its [slides](#)
- [download](#) the release
- browse [sample programs](#)
- view [other resources](#) to learn how to trivially write distributed programs like this:

```
use CyclicDist;           // use the Cyclic distribution Library
config const n = 100;      // use ./a.out --n=<val> to override this default
forall i in {1..n} mapped Cyclic(startIdx=1) do
    writeln("Hello from iteration ", i, " of ", n, " running on node ", here.id);
```

What's Hot?

- Chapel 1.16 is now available—[download](#) a copy today!
- The [CHI UW 2018 call for participation](#) is now available!
- A recent [Cray blog post](#) reports on highlights from CHI UW 2017.
- Chapel is now one of the supported languages on [Try It Online!](#)
- Watch talks from [ACCU 2017](#), [CHI UW 2017](#), and [ATPESC 2016](#) on [YouTube](#).
- [Browse slides](#) from **PADAL**, **EAGE**, **EMBRACE**, **ACCU**, and other recent talks.
- See also: [What's New?](#)



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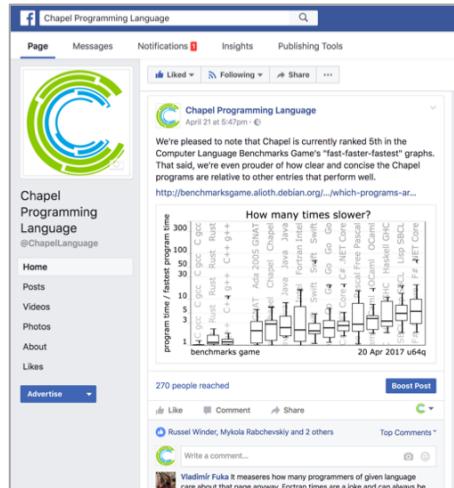
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How to Track Chapel

<http://facebook.com/ChapelLanguage>

<http://twitter.com/ChapelLanguage>

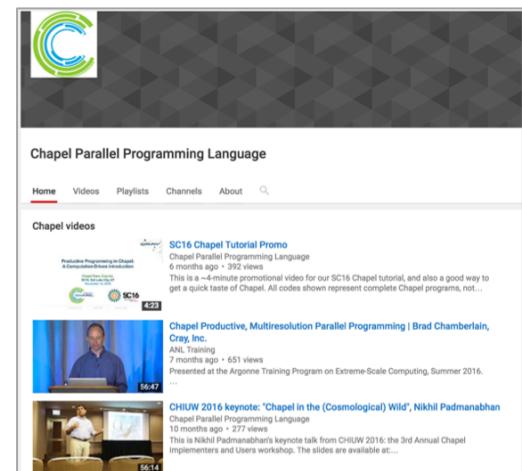
<https://www.youtube.com/channel/UCHmm27bYjhknK5mU7ZzPGsQ/>
chapel-announce@lists.sourceforge.net



The Facebook page for Chapel Programming Language has 129 followers. It features a post from April 21, 2017, at 8:47pm, which reads: "We're pleased to note that Chapel is currently ranked 5th in the Computer Language Benchmarks Game's "fast-faster-fastest" graphs. That said, we're even prouder of how clear and concise the Chapel programs are relative to other entries that perform well." Below the post is a chart titled "Program time / fastest program time" comparing various languages. The chart shows Chapel performing well, often faster than others like C and C++.



The Twitter profile for Chapel Language (@ChapelLanguage) has 222 tweets, 12 following, 129 followers, and 32 likes. A recent tweet from April 20, 2017, at 9:47pm, encourages users to submit interesting applications to the PAW 2017 workshop at SC17. The profile also includes a link to their GitHub repository: [sourcey/institute.github.io/PAW/](https://github.com/sourcey/institute). The bio states: "Chapel is a productive parallel programming language designed for large-scale computing whose development is being led by @cray_inc".



The YouTube channel for Chapel Parallel Programming Language has 129 subscribers. It features several video uploads, including a promotional video for SC16 Chapel Tutorial, a presentation on Chapel Productive, Multiresolution Parallel Programming by Brad Chamberlain, and a keynote talk by Nikhil Padmanabhan from CHI16. The channel also includes a video from the 2nd Annual PGAS Applications Workshop held in November 2017.



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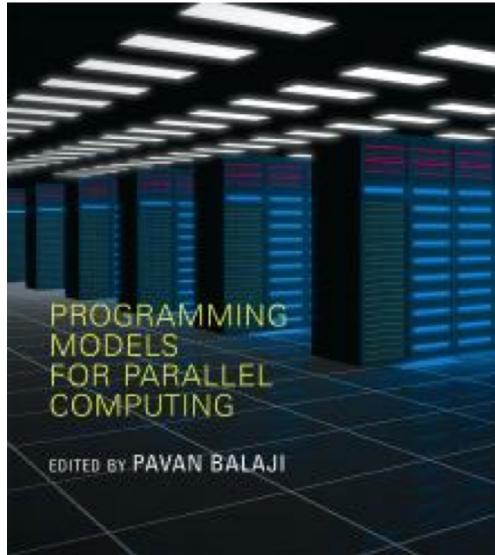
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Suggested Reading (healthy attention spans)



Chapel chapter from [Programming Models for Parallel Computing](#)

- a detailed overview of Chapel's history, motivating themes, features
- published by MIT Press, November 2015
- edited by Pavan Balaji (Argonne)
- chapter is now also available [online](#)



Other Chapel papers/publications available at <https://chapel-lang.org/papers.html>

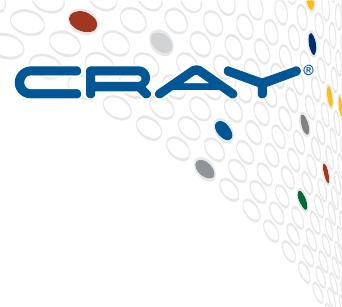


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Suggested Reading (short attention spans)



[CHIUW 2017: Surveying the Chapel Landscape](#), Cray Blog, July 2017.

- *a run-down of recent events*

[Chapel: Productive Parallel Programming](#), Cray Blog, May 2013.

- *a short-and-sweet introduction to Chapel*

[Six Ways to Say “Hello” in Chapel](#) (parts [1](#), [2](#), [3](#)), Cray Blog, Sep-Oct 2015.

- *a series of articles illustrating the basics of parallelism and locality in Chapel*

[Why Chapel?](#) (parts [1](#), [2](#), [3](#)), Cray Blog, Jun-Oct 2014.

- *a series of articles answering common questions about why we are pursuing Chapel in spite of the inherent challenges*

[Ten] Myths About Scalable Programming Languages, [IEEE TCSC Blog](#)

(index available on chapel-lang.org “blog posts” page), Apr-Nov 2012.

- *a series of technical opinion pieces designed to argue against standard reasons given for not developing high-level parallel languages*



Chapel StackOverflow and GitHub Issues



A screenshot showing two side-by-side web pages. On the left is the Stack Overflow 'chapel' tag page, displaying several questions about the Chapel language. On the right is the GitHub repository 'chapel-lang/chapel' issues page, listing 292 open pull requests. Both pages have a dark header and footer.



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Where to..



Submit bug reports:

[GitHub issues for chapel-lang/chapel](#): public bug forum
chapel_bugs@cray.com: for reporting non-public bugs

Ask User-Oriented Questions:

[StackOverflow](#): when appropriate / other users might care
[#chapel-users \(irc.freenode.net\)](#): user-oriented IRC channel
chapel-users@lists.sourceforge.net: user discussions

Discuss Chapel development

chapel-developers@lists.sourceforge.net: developer discussions
[#chapel-developers \(irc.freenode.net\)](#): developer-oriented IRC channel

Discuss Chapel's use in education

chapel-education@lists.sourceforge.net: educator discussions

Directly contact Chapel team at Cray: chapel_info@cray.com



Questions?



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