



Global-View Abstractions for User-Defined Reductions and Scans

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Steve Deitz
Brad Chamberlain
David Callahan
Larry Snyder





A Bit of Context



- DARPA HPCS: High Productivity Computing Systems
 - HPLS: High Productivity Language Systems
- ◆ Increase productivity for HEC community by 2010

```
Productivity = Programmability +
Performance +
Portability +
Robustness
```

- Revolutionary results (not evolutionary)
- Marketable to people other than program sponsors
- Phase II competition: Cray, IBM, and Sun





What is Chapel?



- Chapel: Cascade High-Productivity Language
- Overall goal:
 - Simplify the creation of parallel programs
 - Support evolution to production-grade codes
 - Emphasize generality
- Motivating Language Technologies:
 - 1) Global-view multithreaded parallel programming
 - 2) Locality-aware programming
 - 3) Object-oriented programming
 - 4) Generic programming and type inference





A Bit of Background (Reduce)



Definition

The **reduce operator** takes a binary operation \oplus and a sequence of values (/a₁, a₂, ..., a_n/) and returns the value

$$a_1 \oplus a_2 \oplus ... \oplus a_n$$
.

Example

+ reduce (/1, 2, 3, 4, 5, 6, 7, 8, 9, 10/)

 \rightarrow 55





A Bit of Background (Scan)



Definition

The **scan operator** takes a binary operation \oplus and a sequence of values ($/a_1, a_2, ..., a_n$ /) and returns the sequence of values

$$(/a_1, a_1 \oplus a_2, ..., a_1 \oplus a_2 \oplus ... \oplus a_n/).$$

Example

+ scan (/1, 2, 3, 4, 5, 6, 7, 8, 9, 10/)

 \rightarrow (/1, 3, 6, 10, 15, 21, 28, 36, 45, 55/)





Parallel Performance



- Performance of reduce and scan
 - Highly efficient implementation
 - Log-tree for reduce
 - Parallel-prefix for scan
 - Associative operator > parallel performance
 - Commutative operator → more efficient
- Application of reduce and scan
 - Common in scientific computing
 - Used to test for convergence in iterative algorithms
 - Used in matrix multiplication and sorting kernels
 - → ~9% of MPI calls in NPB are reductions (static count)
 - Supported on today's systems





Outline



- ◆ Reduce and Scan Overview
- Local-View vs. Global-View Overview
- MPI and Chapel Abstraction for Reduce and Scan
- Application to MPI and Quantitative Results





Local vs. Global View



- ◆ Local-View Programming:
 - Programmer codes on a per-processor basis
 - Breaks distributed data structures into per-process chunks
 - Breaks work into per-process iterations/control flow
- ◆ Global-View Programming:
 - Programmer codes independent of processors (mostly)
 - Relies on compiler to manage decomposition details
 - Provides guide to decomposition at a high level

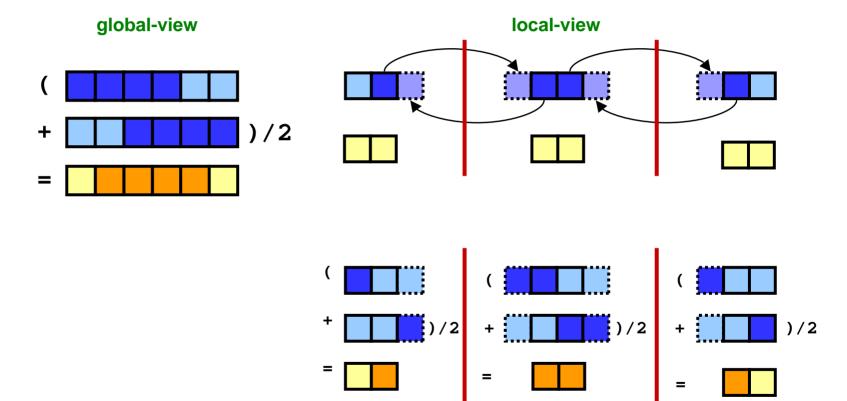




Local vs. Global Example



Example: "Apply a stencil to a vector"







Local vs. Global Example



Example: "Apply a stencil to a vector"

global-view

```
var n: int = 1000;
var a, b: [1..n] float;
forall i in 2..n-1 do
  b(i) = (a(i-1) + a(i+1))/2;
```

local-view

```
var n: int = 1000;
var locN: int = n/numProcs;
var a, b: [0..locN+1] float;
var innerLo: int = 1;
var innerHi: int = locN;
if (iHaveRightNeighbor) {
  send(right, a(locN));
  recv(right, a(locN+1));
} else
  innerHi = locN-1;
if (iHaveLeftNeighbor) {
  send(left, a(1));
  recv(left, a(0));
} else
  innerLo = 2i
forall i in innerLo..innerHi do
  b(i) = (a(i-1) + a(i+1))/2;
```





3D Stencil in NAS MG



```
submoutane: comma(n/n1.n2/n3.kk)
    amcsfeintrinaidain(3) = [-1..1, -1..1, -1..1];
  implicit pone
 iRaram wicasepsth.float
iRaram wicasepsth.float
include *giobars.fl.25, 0.125, 0.0625/);
param w3d: [(i,j,k) in Stencil] float
  integer n1; w2(in30)kk (j!=0) + (k!=0));
    pohše specisionma(n1,n2,n3)
        (_npoh_ dead() | huntbeduce

do axis in recently

if (npcos_ne-1) then

call (synd_all()) it her

call (synd_all()) it call ( yet all ()) it call () it ca
                          call give3( axis, -1, u, n1, n2, n3, kk
                          coll cume oll()
                          call take3( axis, -1, u, n1, n2, n3 )
                          call take3( axis, +1, u, n1, n2, n3 )
                else call commlp( axis, u, n1, n2, n3, kk )
                endif
        enddo
else
        do axis = 1, 3
                 call sync_all(
                 call sync all()
                                                                                                                                                              endif
        call zero3(u,n1,n2,n3)
andif
subroutine give3( axis, dir. u. n1, n2, n3, k )
 use caf intrinsics
implicit none
include 'cafnpb.h'
integer axis, dir, n1, n2, n3, k, ierr
double precision u( n1, n2, n3 )
integer i3, i2, i1, buff len,buff id
buff id = 2 + dir
if( axis .eq. 1 )then
        if( dir .eq. -1 )then
                 do i3=2 n3=1
                          do i2=2,n2-1
                                  buff len = buff len + 1
                                   buff(buff_len,buff_id ) = u( 2,
                          enddo
                                                                                                                                                               endif
               buff(1:buff len.buff id+1)[nbr(axis.dir.k)]
                buff(1-buff len buff id)
        else if( dir .eq. +1 ) then
                do i3=2.n3-1
                                  i2=2,n2-1
buff_len = buff_len + 1
                                  buff(buff len, buff id ) = u( n1-1,
               i2,i3)
                enddo
               buff(1:buff len.buff id+1)[nbr(axis.dir.k)]
                 buff(1:buff_len,buff_id)
endif
if( axis .eq. 2 )then
        if( dir .eg. -1 )then
```

```
do i3=2 n3=1
         do i1=1,n1
            buff_len = buff_len + 1
            buff(buff len, buff id ) = u( i1,
     buff(1:buff_len,buff_id+1)[nbr(axis,dir,k)]
      buff(1:buff len.buff id)
   else if( dir eg +1 ) then
         do i1=1.n1
            buff_len = buff_len + 1
buff(buff len, buff id )= u( i1,n2-
      enddo
enddo
     buff(1:buff len.buff id+1)[nbr(axis.dir.k)]
      buff(1:buff_len,buff_id)
if( axis .eq. 3 )then
   if( dir .eg. -1 )then
      do i2=1.n2
         do i1=1.n1
            buff_len = buff_len + 1
            buff(buff len, buff id ) = u(
     41 42 21
      anddo
     buff(1-buff len buff id+1)[nbr(avis dir k)]
      buff(1:buff_len,buff_id)
   else if( dir .eq. +1 ) then
      do i2=1 n2
         do i1=1,n1
            buff len = buff len + 1
             buff(buff_len, buff_id ) = u(
     i1,i2,n3-1)
     buff(1:buff_len,buff_id+1)[nbr(axis,dir,k)]
      buff(1-buff len buff id)
   endif
subroutine take3( axis, dir. u. n1, n2, n3 )
implicit none
include 'cafnpb.h'
integer axis, dir, n1, n2, n3 double precision u( n1, n2, n3 )
integer buff id. indx
integer i3, i2, i1
buff id = 3 + dir
if( axis .eq. 1 )then
   if( dir .eq. -1 )then
      do i3=2.n3-1
         do i2=2,n2-1
```

indx = indx + 1

```
u(n1 i 2 i 3) = buff(indx buff id)
   else if( dir .eq. +1 ) then
      do i3=2 n3=1
        do i2=2.n2-1
             u(1,i2,i3) = buff(indx, buff_id )
   andif
if( axis .eq. 2 )then if( dir .eq. -1 )then
      do i3=2 n3=1
            indx = indx + 1
            u(i1,n2,i3) = buff(indx, buff_id )
   else if( dir .eq. +1 ) then
         do i1=1.n1
            indy = indy + 1
            u(i1,1,i3) = buff(indx, buff_id )
      enddo
   endif
endif
if( axis .eq. 3 )then if( dir .eq. -1 )then
      do i2=1,n2
            indy = indy + 1
             u(i1,i2,n3) = buff(indx, buff_id)
      enddo
   else if( dir .eq. +1 ) then
      do i2=1,n2
        do i1=1,n1
indx = i
            u(i1,i2,1) = buff(indx, buff_id)
subroutine comm1p( axis, u, n1, n2, n3, kk )
use caf_intrinsics
implicit none
include !cafnph h
integer axis, dir, n1, n2, n3
integer i3, i2, i1, buff len,buff id
integer i, kk, indx
dir = -1
buff_id = 3 + dir
```

```
do i=1,nm2
   buff(i,buff_id) = 0.0D0
enddo
buff_id = 3 + dir
buff len = nm2
   buff(i,buff id) = 0.000
buff id = 2 + div
if( axis .eq. 1 )then
do i3=2,n3-1
      do i2=2 n2=1
          buff_len = buff_len + 1
buff(buff_len, buff_id) = u( n1-1,
andif
if( axis .eg. 2 )then
   do i3=2,n3-1
do i1=1,n1
         buff len = buff len + 1
          buff(buff_len, buff_id )= u( i1,n2-
      enddo
endif
if( axis .eq. 3 )then
   do i2=1,n2
      do i1=1,n1
         buff len = buff len + 1
          buff(buff_len, buff_id ) = u( i1,i2,n3-
      enddo
buff id = 2 + dir
if( axis .eq. 1 )then
do i3=2,n3-1
      do i2=2,n2-1
          buff_len = buff_len + 1
buff(buff_len,buff_id ) = u( 2, i2,i3)
endif
if( axis .eq. 2 )then
do i3=2,n3-1
      do i1=1.n1
          buff_len = buff_len + 1
          buff(buff len. buff id ) = u( i1.
   enddo
if( axis .eq. 3 )then
  do i2=1,n2
      do i1=1.n1
          buff_len = buff_len + 1
          buff(buff_len, buff_id ) = u( i1,i2,2)
      enddo
endif
   buff(i.4) = buff(i.3)
   buff(i,2) = buff(i,1)
dir = -1
```

```
buff_id = 3 + dir
if( axis .eq. 1 )then
   do i3=2 n3=1
         u(n1 i2 i3) = buff(indy buff id )
endif
if( axis .eq. 2 )then
   do i3=2,n3-1
     do i1=1.n1
         u(i1,n2,i3) = buff(indx, buff_id )
      enddo
   andda
if( axis .eq. 3 )then
do i2=1,n2
     do i1=1 n1
         indx = indx + 1
u(i1,i2,n3) = buff(indx, buff_id)
      enddo
enddo
endif
dir - 41
buff_id = 3 + dir
if( axis .eq. 1 )then
   do i3=2,n3-1
     do i2=2,n2-1
         u(1,i2,i3) = buff(indx, buff id)
if( axis .eq. 2 )then
do i3=2,n3-1
     do i1=1 n1
         u(i1,1,i3) = buff(indx, buff_id)
      enddo
if( axis .eq. 3 )then
   do i2=1,n2
      do i1=1,n1
         u(i1,i2,1) = buff(indx, buff id )
      enddo
endif
return
```





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- ◆ Local-View vs. Global-View Overview
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User Defined Reduction MinK



Compute the k minimum values in a sequence of values

$$\rightarrow$$
 (1, 2)

E.g. mink(2) reduce (/Cray, Gonzaga, IBM, Sun, Duke/)

→ (Duke, Gonzaga)



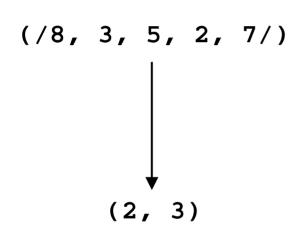


Local vs. Global View

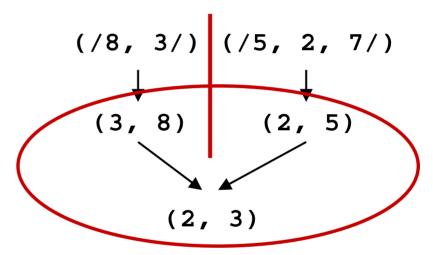


◆ Example: "Reduce a sequence with Min2"

global-view



local-view



Local-view abstraction applies only to this part of the reduction





MPI Reductions



- Has two varieties
 - Reduce returns reduction on root process
 - Allreduce returns result back to all processes

```
int MPI_Reduce(
    void* sendbuf,
    void* recvbuf,
    int count,
    MPI_Datatype datatype,
    MPI_Op op,
    int root,
    MPI Comm comm)
int MPI_Allreduce(
    void* sendbuf,
    void* recvbuf,
    int count,
    MPI_Datatype datatype,
    MPI_Op op,
    MPI_Op op,
    MPI_Comm comm)
```

- Supports nine built-in reduction operators
- Accepts arbitrary combine function
- MPI_Op assumes associativity, can be commutative

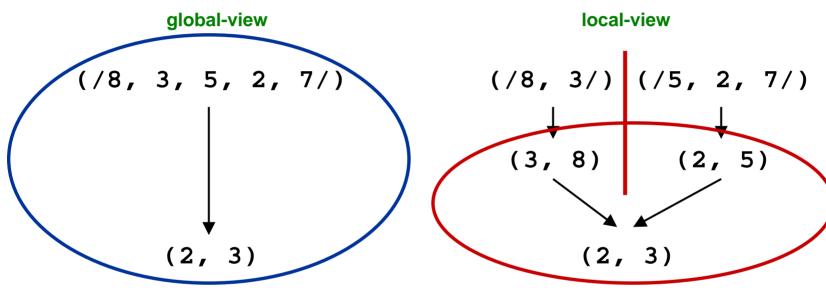




Local vs. Global View



◆ Example: "Reduce a sequence with Min2"



Global-view abstraction needs to hide the details of communication

Local-view abstraction applies only to this part of the reduction





Brief Intro to Chapel Classes



- Similar to classes in other OO languages
 - Fields are accessed via "." syntax
 - Methods are invoked "." syntax
 - Support multiple inheritance, dynamic dispatch
- Automatic default constructor

```
class Point {
  var x : int, y : int;
}
class ColorfulPoint : Point {
  var c : color;
}
var cp = ColorfulPoint(x=2,y=3,c=red);
```

Default constructor provides named argument passing





Brief Intro to Chapel Classes



Classes can be generic

```
class Triple {
  type elt_type;
  var x, y, z : elt_type;
}
def Triple.add(e : elt_type) {
  x += e; y += e; z += e;
}
var t = Triple(float);
t.add(1.0);
```

Class can be instantiated on different types





Chapel Reduction Class



```
class reduction {
                                                  All reductions are generic to the
                                                  sequence element type and two
  type elt type;
                                                 parameters that can be overridden
  param commutative = true;
  param associative = true;
  def accumulate(x : elt type);
  def combine(s : reduction);
  def gen();
  def scan_gen(x : elt_type) do
    return gen();
                                                     Default scan and reduce
  def red gen() do
                                                     generate functions call a
    return gen();
                                                    common generate function
def reduce(r, s) {
  for e in s do
                                            Serial reduce is implemented
    r.accumulate(e);
                                            as is any overloaded operator
  return r.red gen();
```





MinK in Chapel



```
class mink : reduction {
  var k : int;
  var v : [1..k] elt_type = max(elt_type);
  def accumulate(x : elt type) do
    if x < v(1) {
      v(1) = x_i
      for i in 2..k do
        if v(i-1) < v(i) then
          (v(i-1),v(i)) = (v(i),(v(i-1));
  def combine(s: mink(elt_type)) do
    for i in 1..k do
      accumulate(s.v(i));
  def gen() do
    return v;
... mink(2) reduce (/8, 3, 5, 2, 7/) ...
```

Default constructor stores maxes in k-element array.

Operator class stores state.

Combine function takes other thread's locally accumulated state.





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Global-View in MPI



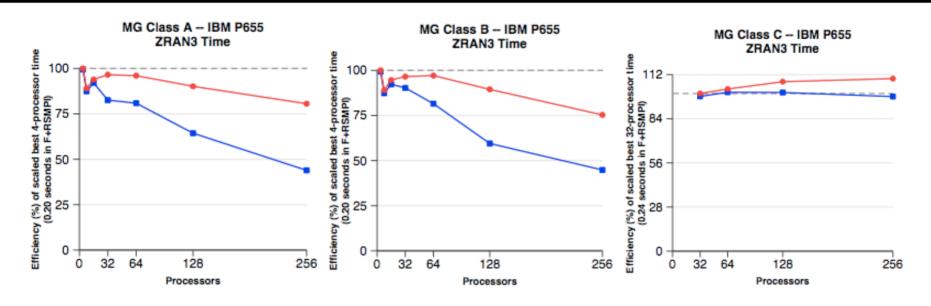
- RSMPI (Reduce and Scan in MPI)
 - Proof of concept of global view
 - Potential use for MPI programmers
- Minimal changes to MPI to support global view
 - Adds a construct to iterate over the local part of an ordered set
 - Uses a preprocessor to generate the accumulate loop based on the iterator construct.
 - Adds an operator as a collection of functions
 - Identity function
 - Accumulate function
 - Combine function
 - Generate function
 - Requires a single RSMPI_Reduce call

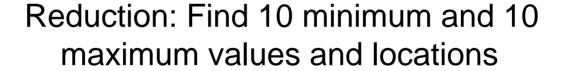


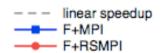


MG Initialization Results









F+MPI: Uses 10 min and 10 max reductions

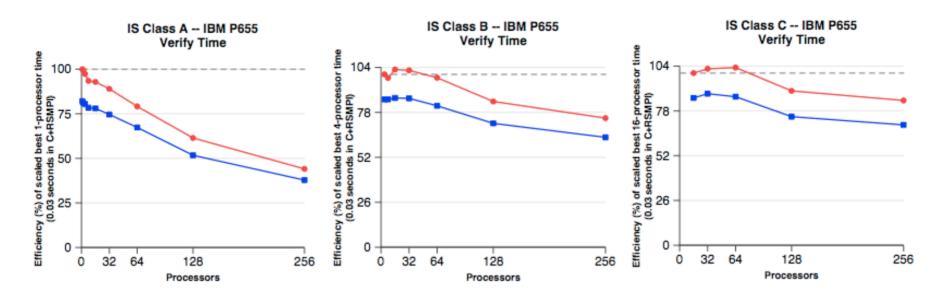
F+RSMPI: Uses one big reduction



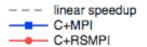


IS Verification Results





Reduction: Determine if array is sorted



C+MPI: Compare boundary elements, determine if local portion is sorted, and 'sum' reduce local determinations
C+RSMPI: Use one big reduction





Summary



- High-level Chapel abstraction
 - Takes advantage of generic OOP capabilities
 - Contains both accumulate and combine function
 - Uses default constructor to produce identity
 - Supports generate function to return result that is different from the combining state
 - Supports a scan generate function that can return different result from a reduction
- RSMPI
 - Shows performance does not suffer at global view
 - Suggests path for incremental improvements
- Global-View
 - Orthogonal to changes in the distribution/layout





Questions?



For more information:

http://chapel.cs.washington.edu

email: chapel_info@cray.com

