

# **Benchmark Improvements**

Chapel Team, Cray Inc.
Chapel version 1.15
April 6, 2017



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### **Outline**

- LCALS Benchmark Improvements
- ISx Benchmark Improvements
- Computer Language Benchmarks Game (CLBG)





# **LCALS Benchmark Improvements**



## **LCALS Improvements: Background**



- Livermore Compiler Analysis Loop Suite
- Chapel port first released with version 1.13.0
- Version 1.14.0 saw significant performance improvements
  - Serial variant matching reference version
  - Parallel variant still had room for improvement



## **LCALS Improvements: This Effort**



## Kernel-by-kernel improvements

- Made nested loops into loops over 2D domains
- Improved style of array initializations
- Changed a 64-bit based range to 32-bit to match the reference version

## Changed some compiler/execution default options

- dropped --data-par-min-granularity=1000
- added –no-ieee-float
  - both option changes done for a more equal comparison to the reference

## Added SPMD variant of parallel kernels

- Start all tasks
- Block the kernels up across tasks
- Execute blocks serially within each task
- Barrier at the end



## **LCALS Improvements: Task Startup Times**



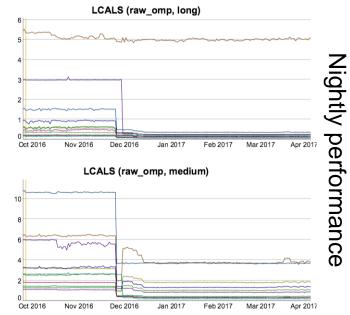
- The amount of work in each kernel is small
  - Split up across ~24 tasks the amount of work per task is very small
  - Kernels are repeated a large number of times
  - Leads to millions of tasks doing very little work each
- Task startup times are a large factor in total execution time
- Optimization added to the Chapel tasking interface
  - Enable threads looking for work to more quickly pick up new tasks
  - Dramatically improves task startup times



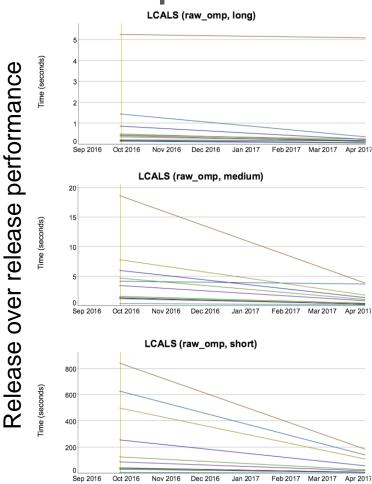
# **LCALS Improvements: Impact**



## Improved task startup time increased performance overall



Improvement for short loop size was masked in nightly testing by another change, but is clear in release-over-release testing.

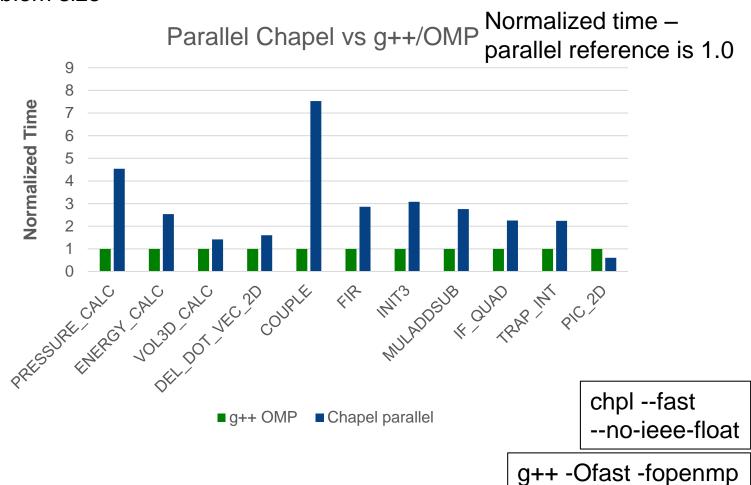




## LCALS Status: Impact: Parallel Perf. v1.14.0



Long problem size



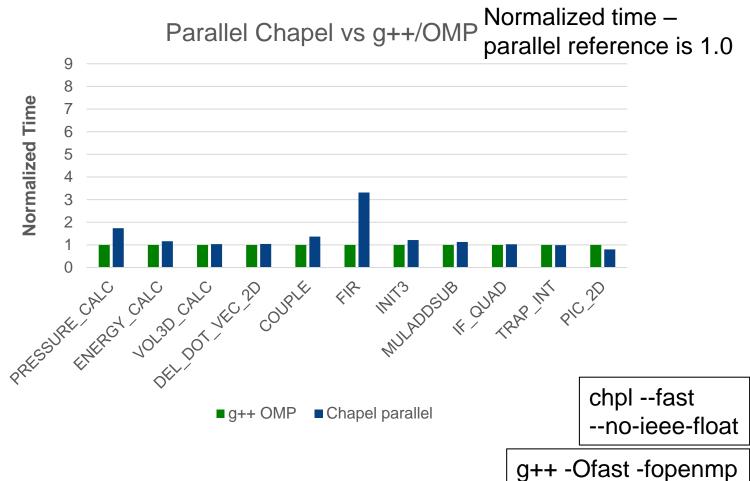


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## LCALS Status: Impact: Parallel Perf. v1.15.0



Long problem size





# **LCALS Improvements: Status and Next Steps**



#### Status:

- The parallel kernels have improved dramatically since version 1.14.0
- A few still lag slightly behind the C+OpenMP reference versions
- But most are effectively matching reference versions

## **Next Steps:**

- There is still some startup overhead to work through
  - Some ideas:
    - start multiple tasks simultaneously
    - improve task counting/sync mechanism





# **ISx Benchmark Improvements**



## **ISx: Background**



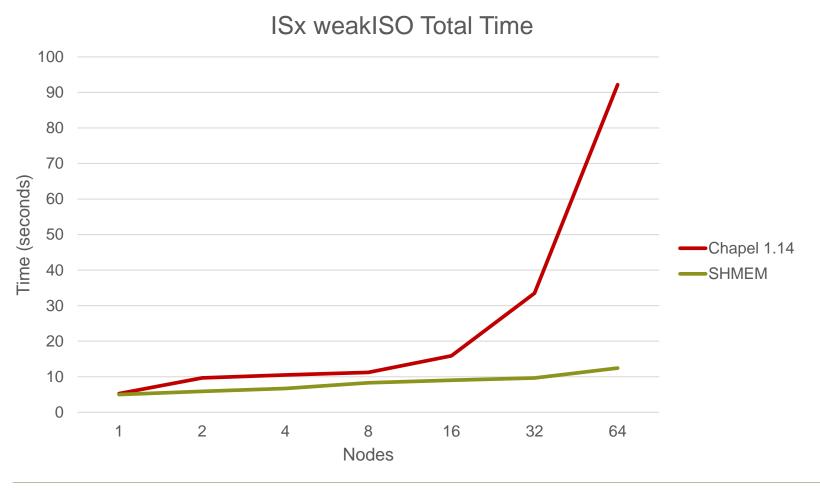
- Scalable Integer Sort benchmark
  - Developed at Intel, published at PGAS 2015
  - SPMD-style computation with barriers
  - Punctuated by all-to-all bucket exchange pattern
  - References implemented in SHMEM and MPI
- Chapel implementation introduced in 1.13 release
  - Motivation: a common distributed pattern
- Initial results showed roughly 10x worse performance



# **ISx: Background**



## 64 nodes on Cray XC





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## **ISx:** Background



## Developed an inelegant optimized version for study

- Uses language internals
- Also called the 'heroic' version, referring to user effort
- Competitive with SHMEM reference

## Found several areas for improvement:

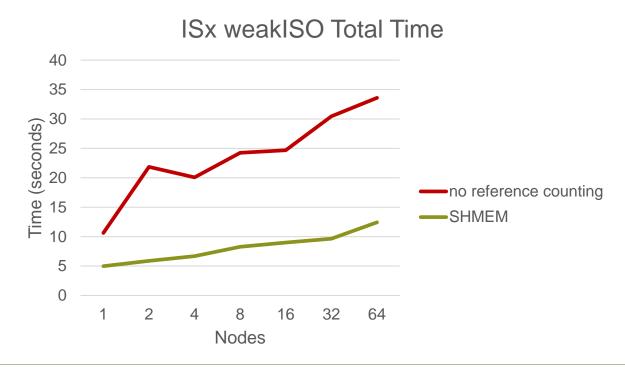
- Reference counted arrays
- Sliced assignment between arrays
- Wide-pointer overhead for serial loops



## **ISx:** This Effort - Reference Counting



- Reference counting arrays was costly at scale
  - Lots of on-statements and atomic operations
- Array memory management work resolved this problem
  - Eliminated need for reference counting entirely





## ISx: This Effort – Sliced Assignment Overhead



## The following pattern was quite slow at scale:

Core operation in all-to-all exchange step
 // 'Dest' is a remote 1D array, 'Src' is local
 Dest[1..10] = Src[1..10];

## Slicing the remote array involved an on-statement

- Used to initialize array metadata
- In all-to-all step, this might mean numCores<sup>2</sup> on-statements
- Problem worsens as we scale to more locales

## on-statement also used for slicing Src's domain

Src was declared over a remote const global domain



## **ISx: This Effort – Sliced Assignment Overhead**



## Fix #1: ArrayViews

- Allows for much cheaper creation of a remote array slice
- Avoids the on-statement entirely
- Still able to bulk-transfer the array data (single PUT/GET)

#### Fix #2: Declare Src over local domain

- Slicing 'Src' now an entirely local operation
- Before:

```
var Src : [NumTasksSpace] keyType;
```

After:

```
var Src : [0..#numTasks] keyType;
```

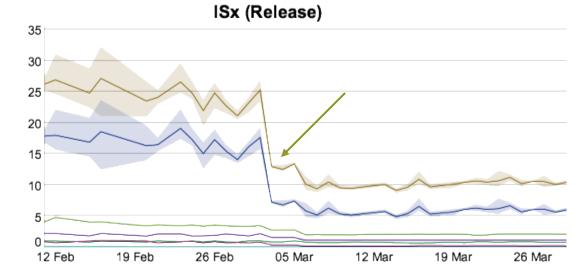


## **ISx: Impact – Sliced Assignment Overhead**



## • Result: 2x performance improvement

- Also less noise
- ugni-qthreads, 16 nodes on Cray XC



- total time (gnu+ugni-gthreads)
- exchange time (gnu+ugni-qthreads)
- bucketize time (gnu+ugni-qthreads)
- bucket count time (gnu+ugni-qthreads)
- input time (gnu+ugni-qthreads)
- count keys time (gnu+ugni-qthreads)
- bucket offset time (gnu+ugni-qthreads)



Time (seconds)

# **ISx:** This Effort – Wide-pointer overhead



- Compiler inserts wide-pointers for potentially-remote data
  - At its simplest, a struct containing a locale ID and pointer
  - Sometimes thwart backend compiler optimizations
- Can introduce overhead for array accesses
  - Problematic for subsections like this in ISx:

```
for key in myKeys do
  sizes[key/width] += 1;
```

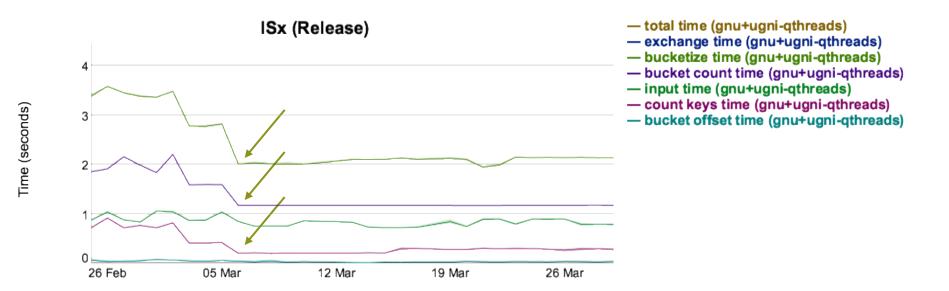
- 1.14 optimized some cases, regressed in 1.15 dev cycle
- For 1.15 a similar approach is used, but only for arrays
  - Disabled with "--[no-]infer-local-fields" compiler flag



## **ISx:** Impact – Wide-pointer overhead



Small, but useful, improvements in serial loops



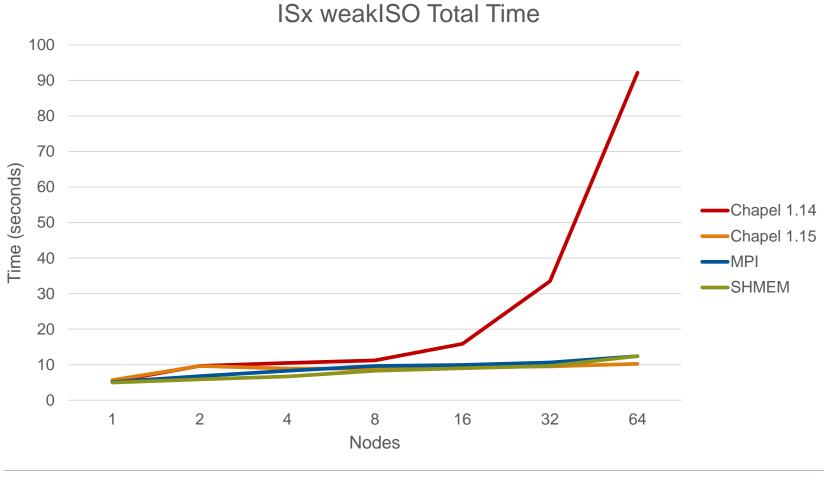
- Work remains to match timed subsections of reference
  - Roughly 20% worse in some cases
  - Note: exchange step, not impacted here, dominates at scale



# **ISx: Performance Summary**



## Gathered on Cray XC with default problem size

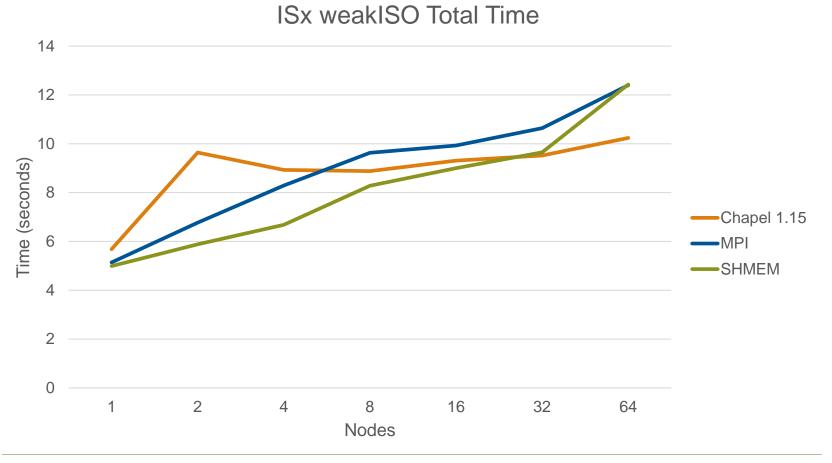




# **ISx: Performance Summary**



- Gathered on Cray XC with default problem size
  - Let's zoom in...





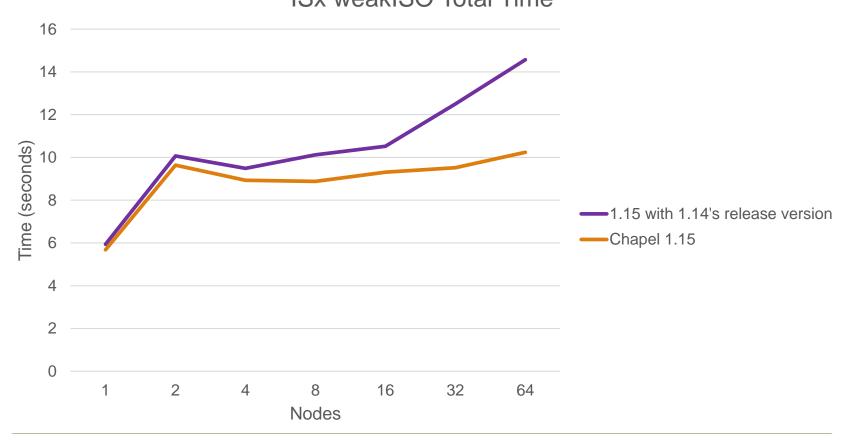
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# **ISx: Performance Summary**



- 1.15 with 1.14's release version of source code
  - 1.14 release version did not declare arrays over local domains ISx weakISO Total Time





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## **ISx: Status and Next Steps**



Status: Big improvements with minor user-code changes

## Next Steps:

- Address known performance issues
  - Under-counting tasks bug leading to oversubscription
    - Likely to impact other benchmarks as well
  - Eliminate additional wide-pointer overhead
  - Investigate performance at larger scales
  - Compiler should optimize references to remote const domains
- Continue to investigate performance results
  - Why does Chapel beat SHMEM and MPI at 64 nodes?
  - Why does Chapel scale poorly from 1-4 nodes, then recover?





# **Computer Language Benchmarks Game (CLBG)**



# Computer Language Benchmarks Game (CLBG)



#### 64-bit quad core data set

Will your <u>toy benchmark program</u> be faster if you write it in a different programming language? It depends how you write it!

#### Which programs are fast?

Which are succinct? Which are efficient?

Ada	<u>C</u>	Chape	1	Cloj	ure	<u>C#</u>	<u>C++</u>
Dart	Er]	lang	F#	For	tran	Go	Hack
Haskel	1	Java	Jav	aScri	ipt	Lisp	Lua
OCam	1	<u>Pascal</u>	P	erl	PHP	Pyt	hon
Rack	et	Ruby	JR	Ruby	Rust	t Sc	ala
<u>.</u>	Smal	ltalk	Swi	.ft	Types	Script	

## Website that supports crosslanguage game / comparisons

- 13 toy benchmark programs
- exercises key features like:
  - memory management
  - tasking and synchronization
  - vectorization
  - big integers
  - strings and regular expressions
- specific approach prescribed

## Take results w/ grain of salt

- other programs may be different
  - not to mention other programmers
- specific to this platform / OS / ...

That said, it's one of the only games in town...



# **Computer Language Benchmarks Game (CLBG)**



#### The Computer Language Benchmarks Game

#### 64-bit quad core data set

Will your toy benchmark program be faster if you write it in a different programming language? It depends how you write it!

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Ada	<u>C</u>	Chape	1	Cloj	ure	<u>C#</u>	<u>C++</u>
Dart	Er.	lang	F#	Fort	tran	Go	Hack
Haskel	1_	Java	Jav	aScri	.pt	Lisp	Lua
OCam	1	Pascal	P	erl	PHP	Pyt	hon
Rack	et	Ruby	JR	uby	Rust	t So	ala
9	Smal	ltalk	Swi	ft	Types	Script	

## Chapel's approach to CLBG:

- want to know how we compare
- strive for entries that are elegant rather than heroic
  - e.g., "Want to learn how program x works? Check out the Chapel version."



### **CLBG: Website**



### Can sort results by execution time, code size, memory or CPU use:

The Computer Language Benchmarks Game

chameneos-redux description

program source code, command-line and
measurements

×	source	secs	mem	gz	cpu	cpu load
1.0	<b>C</b> gcc #5	0.60	820	2863	2.37	100% 100% 98% 100%
1.2	<b>C++</b> g++ #5	0.70	3,356	1994	2.65	100% 100% 91% 92%
1.7	Lisp SBCL #3	1.01	55,604	2907	3.93	97% 96% 99% 99%
2.3	Chapel #2	1.39	76,564	1210	5.43	99% 99% 98% 99%
3.3	Rust #2	2.01	56,936	2882	7.81	97% 98% 98% 98%
5.6	C++ g++ #2	3.40	1,880	2016	11.88	100% 51% 100% 100%
6.8	Chapel	4.09	66,584	1199	16.25	100% 100% 100% 100%
8.0	Java #4	4.82	37,132	1607	16.73	98% 98% 54% 99%
8.5	Haskell GHC	5.15	8,596	989	9.26	79% 100% 2% 2%
10	Java	6.13	53,760	1770	8.78	42% 45% 41% 16%
10	Haskell GHC #4	6.34	6,908	989	12.67	99% 100% 2% 1%
11	C# .NET Core	6.59	86,076	1400	22.96	99% 82% 78% 91%
11	Go	6.90	832	1167	24.19	100% 96% 56% 100%
13	Go #2	7.59	1,384	1408	27.65	91% 99% 99% 78%
13	Java #3	7.94	53,232	1267	26.86	54% 96% 98% 94%

The Computer Language Benchmarks Game

chameneos-redux

description

program source code, command-line and
measurements

×	source	secs	mem	gz	cpu	cpu load
1.0	Erlang	58.90	28,668	734	131.19	62% 60% 51% 53%
1.0	Erlang HiPE	59.39	25,784	734	131.58	60% 56% 56% 54%
1.1	Perl #4	5 min	14,084	785	7 min	40% 40% 29% 28%
1.1	Racket	5 min	132,120	791	5 min	1% 0% 0% 100%
1.1	Racket #2	175.88	116,488	842	175.78	100% 1% 1% 0%
1.2	Python 3 #2	236.84	7,908	866	5 min	24% 48% 27% 45%
1.3	Ruby	90.52	9,396	920	137.53	35% 35% 35% 34%
1.3	Ruby JRuby	48.78	628,968	928	112.15	65% 60% 49% 58%
1.3	<b>Go</b> #5	11.05	832	957	32.48	75% 74% 75% 73%
1.3	Haskell GHC #4	6.34	6,908	989	12.67	99% 100% 2% 1%
1.3	Haskell GHC	5.15	8,596	989	9.26	79% 100% 2% 2%
1.6	OCaml #3					32% 38% 37% 39%
1.6	Go	gz == 0	code s	ize m	etric	100% 96% 56% 100%

strip comments and extra

whitespace, then gzip



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1.6 Chapel

1.6 Chapel #2

0% 100% 100% 100%

99% 99% 98% 99%

### **CLBG: Website**



# Can also compare languages pair-wise:

The Computer Language Benchmarks Game

Chapel programs versus Go all other Chapel programs & measurements

by benchmark task performance

#### regex-redux

source	secs	mem	gz	cpu	cpu load
Chapel	10.02	1,022,052	477	19.68	99% 72% 14% 12%
Go	29.51	352,804	798	61.51	77% 49% 43% 40%

#### binary-trees

source	secs	mem	gz	cpu	cpu load
Chapel	14.32	324,660	484	44.15	100% 58% 78% 75%
Go	34.77	269,068	654	132.04	95% 97% 95% 95%

#### fannkuch-redux

source	secs	mem	gz	cpu	cpu load
Chapel	11.38	46,056	728	45.18	100% 99% 99% 100%
Go	15.81	1,372	900	62.92	100% 100% 99% 99%



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#### **CLBG: Website**



#### Can also browse program source code (but this requires actual thought):

```
proc main() {
  printColorEquations();
  const group1 = [i in 1..popSize1] new Chameneos(i, ((i-1)%3):Color);
  const group2 = [i in 1..popSize2] new Chameneos(i, colors10[i]);
  cobegin {
   holdMeetings(group1, n);
    holdMeetings(group2, n);
  print(group1);
  print(group2);
  for c in group1 do delete c;
  for c in group2 do delete c;
// Print the results of getNewColor() for all color pairs.
proc printColorEquations() {
  for c1 in Color do
    for c2 in Color do
      writeln(c1, " + ", c2, " -> ", getNewColor(c1, c2));
 writeln();
// Hold meetings among the population by creating a shared meeting
// place, and then creating per-chameneos tasks to have meetings.
proc holdMeetings(population, numMeetings) {
  const place = new MeetingPlace(numMeetings);
  coforall c in population do
                                        // create a task per chameneos
    c.haveMeetings(place, population);
  delete place;
```

```
void get_affinity(int* is_smp, cpu_set_t* affinity1, cpu_set_t* affinity2)
                                active cpus;
   cpu set t
   FILE*
   char
                                buf [2048];
   char const*
                                pos;
                                cpu idx;
                                physical id;
   int
                                core_id;
   int
                                cpu cores;
   int
                                apic id;
   size t
                                cpu count;
   size t
   char const*
                                processor str
                                                     = "processor":
   size t
                                processor str len
                                                    = strlen(processor str);
   char const*
                                physical id str
                                                    = "physical id";
   size t
                                physical id str len = strlen(physical id str);
   char const*
                                core id str
                                                    = "core id";
   size t
                                core id str len
                                                    = strlen(core_id_str);
   char const*
                                cpu cores str
                                                    = "cpu cores";
                                                    = strlen(cpu cores str);
   size t
                                cpu cores str len
   CPU ZERO(&active cpus);
   sched getaffinity(0, sizeof(active cpus), &active cpus);
   cpu count = 0;
   for (i = 0; i != CPU SETSIZE; i += 1)
       if (CPU ISSET(i, &active cpus))
            cpu count += 1:
   if (cpu count == 1)
        is smp[0] = 0;
   is smp[0] = 1;
   CPU ZERO(affinity1);
```

#### excerpt from 1210 gz Chapel #2 entry

excerpt from 2863 gz C gcc #5 entry



### What's new with the CLBG?



## Two programs changed their official definitions:

#### binary-trees:

- improved checksum to avoid false positives at 1/2, 1/4, 1/8 the memory
- eliminated per-node data field
- changed what trees are allocated and freed, slightly
- increased the problem size

#### regex:

- changed the regular expression used
- renamed the test to regex-redux
- several versions are not currently passing due to these changes
  - our current standings may be due in part to this

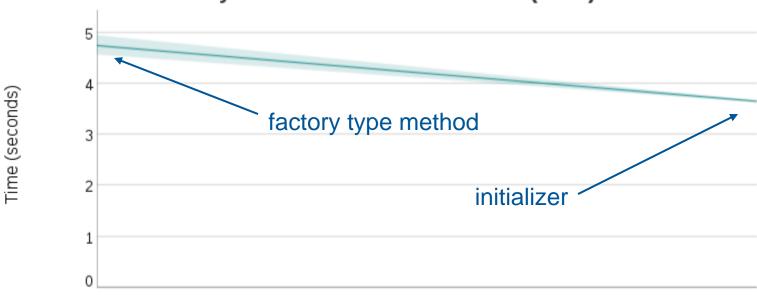




#### • We've submitted some new versions:

binary-trees: used an initializer rather than a factory type method

#### Binary Trees Shootout Benchmark (n=20)

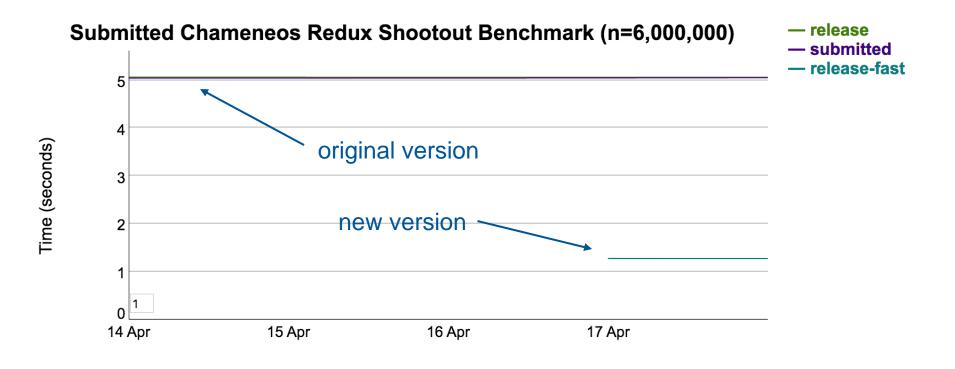






#### • We've submitted some new versions:

binary-trees: used an initializer rather than a factory type method chameneos-redux: increased parallelism and tuned a spin-wait



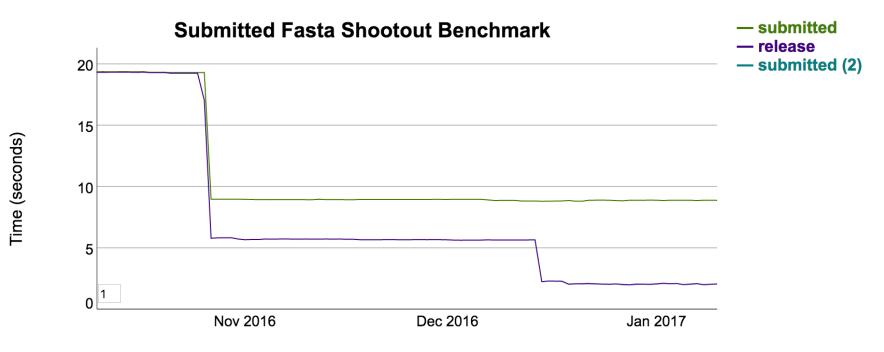




#### • We've submitted some new versions:

binary-trees: used an initializer rather than a factory type method chameneos-redux: increased parallelism and tuned a spin-wait fasta: implemented a parallel version and tuned for clarity and speed

• also, changed some 'var' declarations due to const-checking improvements



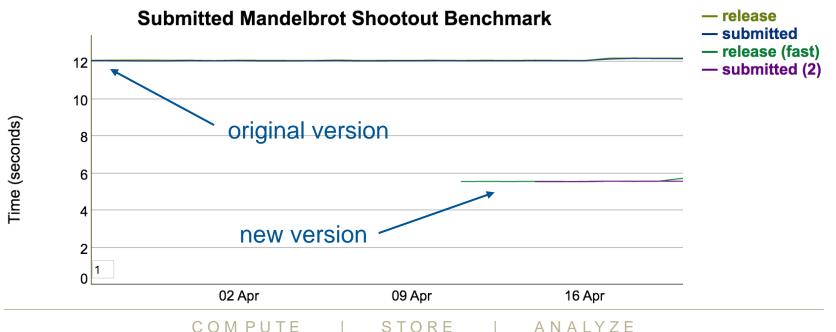




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• also, changed some 'var' declarations due to const-checking improvements mandelbrot: accelerated by hoisting values and using tuples of values





#### What's new with the Chapel CLBG entries?



#### • We've submitted some new versions:

binary-trees: used an initializer rather than a factory type method chameneos-redux: increased parallelism and tuned a spin-wait fasta: implemented a parallel version and tuned for clarity and speed

- also, changed some 'var' declarations due to const-checking improvements mandelbrot: accelerated by hoisting values and using tuples of values meteor-fast: fixed a race condition caused by array memory changes
- textbook example of an array being used by a 'begin' task
   pidigits: submitted a version that uses 'bigint's
  - currently the #1 fastest version, and also quite elegant

#### Note that some of these changes followed the 1.15 release

As such, not all are found in examples/benchmarks/shootout/ for 1.15



#### CLBG: Chapel Standings as of Oct 17th



- - one #1 fastest: pidigits
  - 2 others in the top-5 fastest: meteor-contest thread-ring
  - 2 others in the top-10 fastest: chameneos-redux fannkuch-redux
  - 3 others in the top-20 fastest: binary-trees n-body spectral-norm

- 8 / 13 programs in top-20 fastest: 8 / 13 programs in top-20 smallest:
  - two #1 smallest: n-body thread-ring
  - 2 others in the top-5 smallest: pidigits spectral-norm

4 others in the top-20 smallest: chameneos-redux mandelbrot meteor-contest regex-dna



#### CLBG: Chapel Standings as of Apr 20th



- - one #1 fastest: pidigits
  - 3 others in the top-5 fastest: chameneos-redux meteor-contest thread-ring
  - 3 others in the top-10 fastest: fannkuch-redux fasta mandelbrot
  - 5 others in the top-20 fastest: binary-trees k-nucleotide n-body regex-redux spectral-norm

- 12 /13 programs in top-20 fastest:
   8 / 13 programs in top-20 smallest:
  - two #1 smallest: n-body thread-ring
  - 2 others in the top-5 smallest: pidigits spectral-norm
  - 1 other in the top-10 smallest: regex-redux

3 others in the top-20 smallest: chameneos-redux mandelbrot meteor-contest



### CLBG: Website's fast-faster-fastest graph



#### Site summary: relative performance (sorted by geometric mean)

How many times slower? 300 program 100 50 program time / fastest 30 10 5 benchmarks game 20 Apr 2017 u64q



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#### **CLBG: Website**



site has voiced good philosophy about too-easy answers

**We want easy answers**, but easy answers are often incomplete or wrong. You and I know, there's more we should understand:

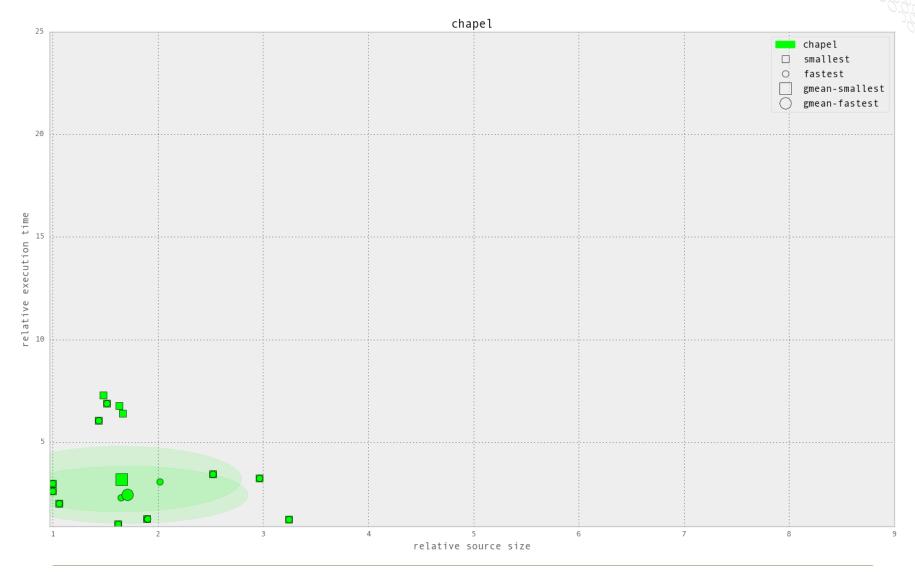
<u>stories</u> <u>details</u> <u>fast?</u> <u>conclusions</u>

- yet, most readers probably still jump to conclusions
  - execution time dominates default (or only) views of results
  - it's simply human nature
- we're interested in elegance as well as performance
  - elegance is obviously in the eye of the beholder
    - we compare source codes manually
    - but then use CLBG's code size metric as a quantitative stand-in
  - want to be able to compare both axes simultaneously
  - to that end, we used scatter plots to compare implementations



# **Chapel entries**





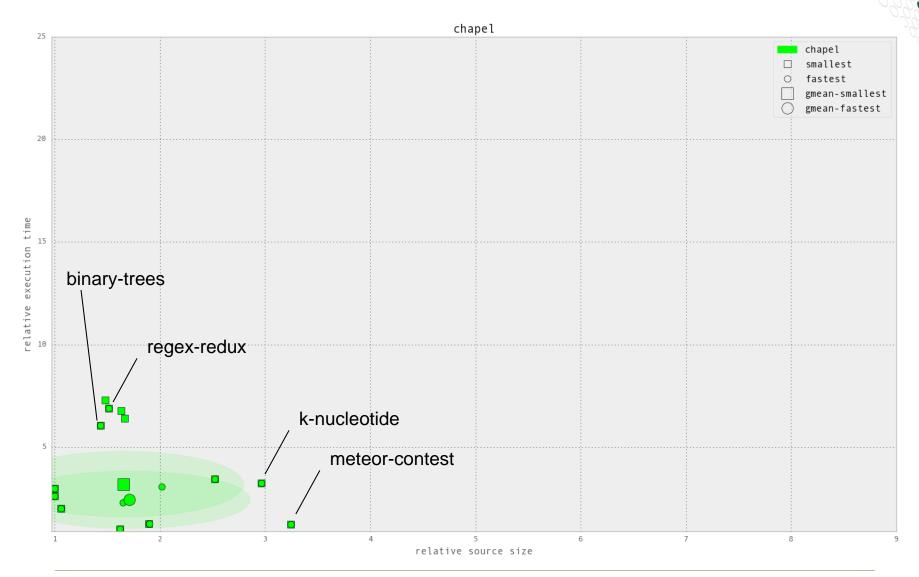


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# **Chapel entries (noting outliers)**

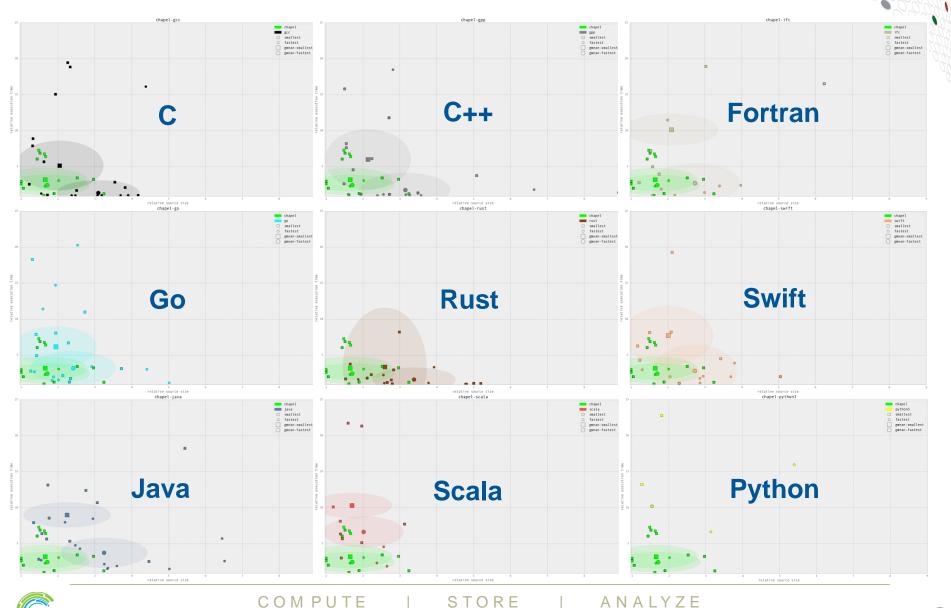






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# Chapel vs. 9 other languages





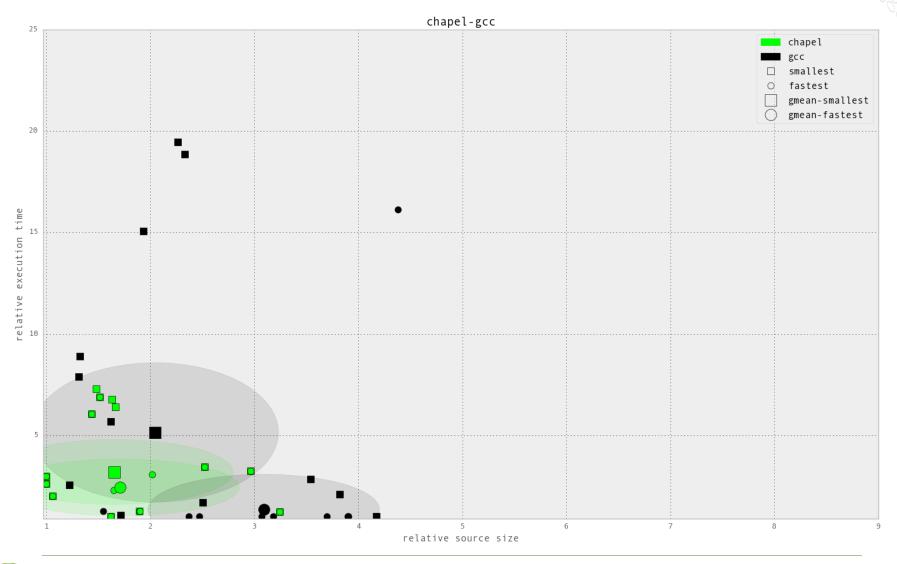
#### Chapel vs. 9 other languages (zoomed out)





# Chapel vs. C

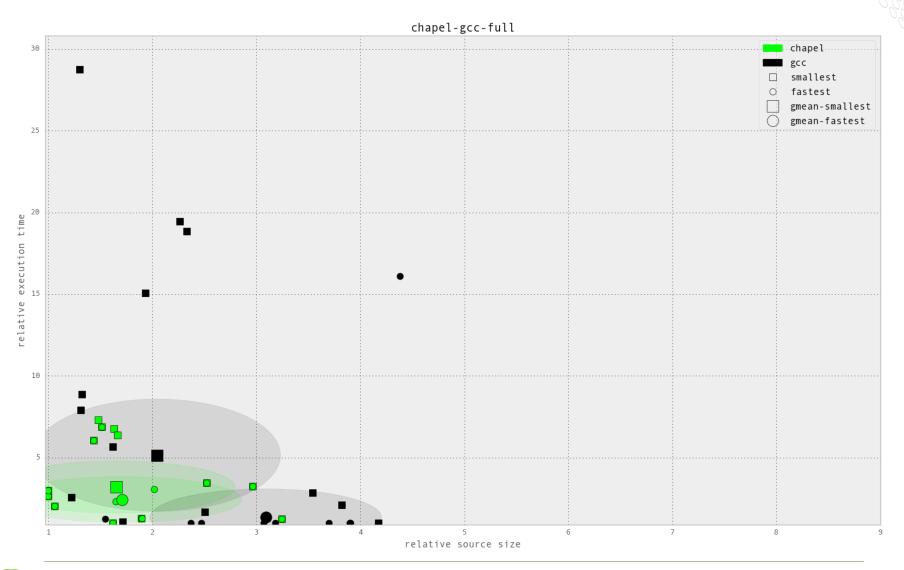






# Chapel vs. C (zoomed out)



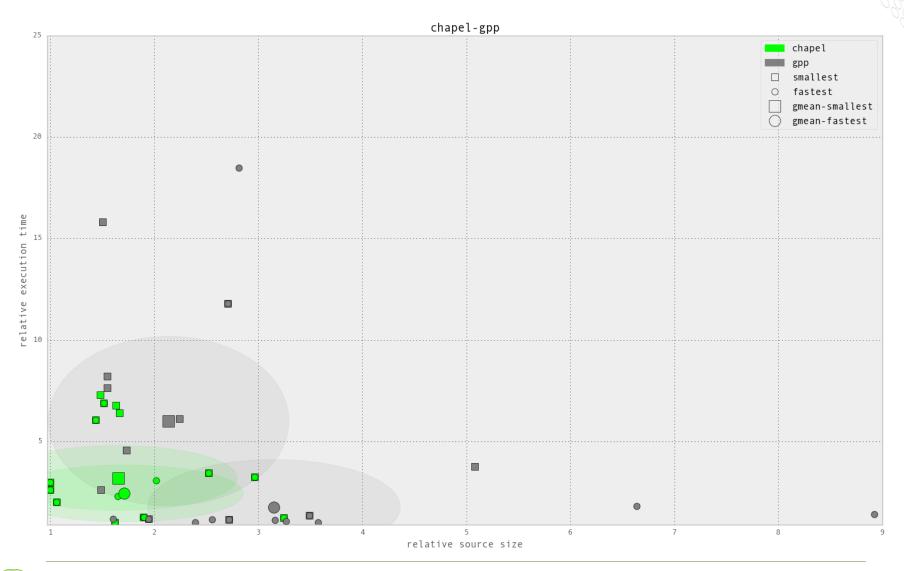




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# Chapel vs. C++

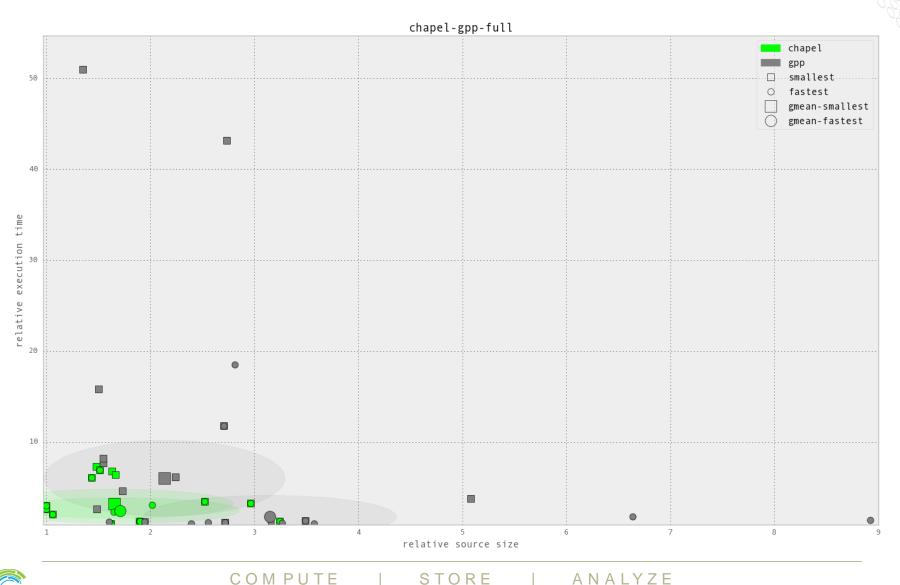






# Chapel vs. C++ (zoomed out)



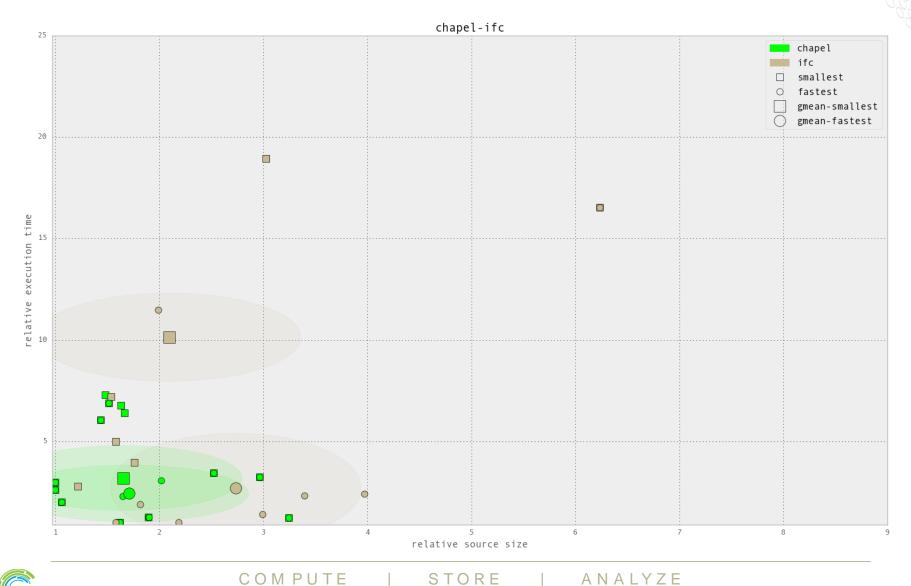




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## **Chapel vs. Fortran**



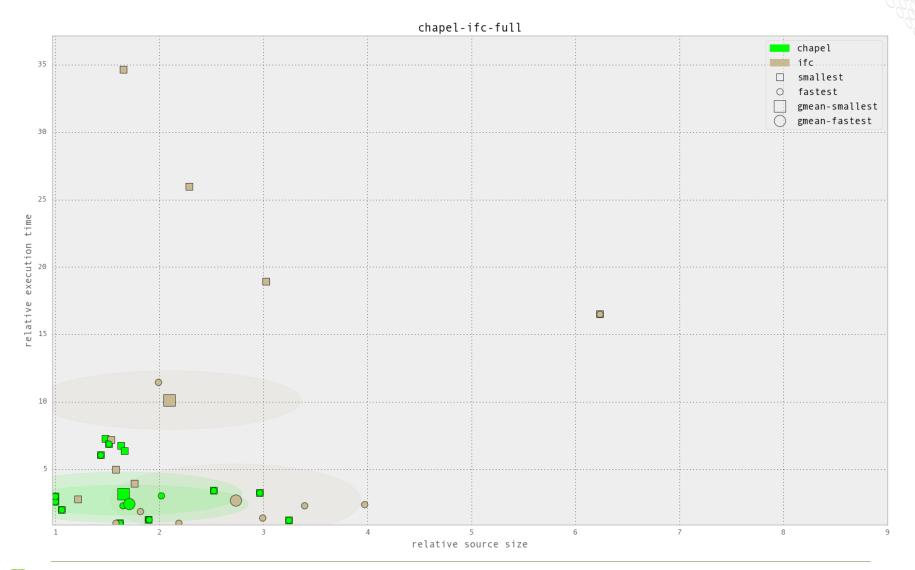




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# **Chapel vs. Fortran (zoomed out)**

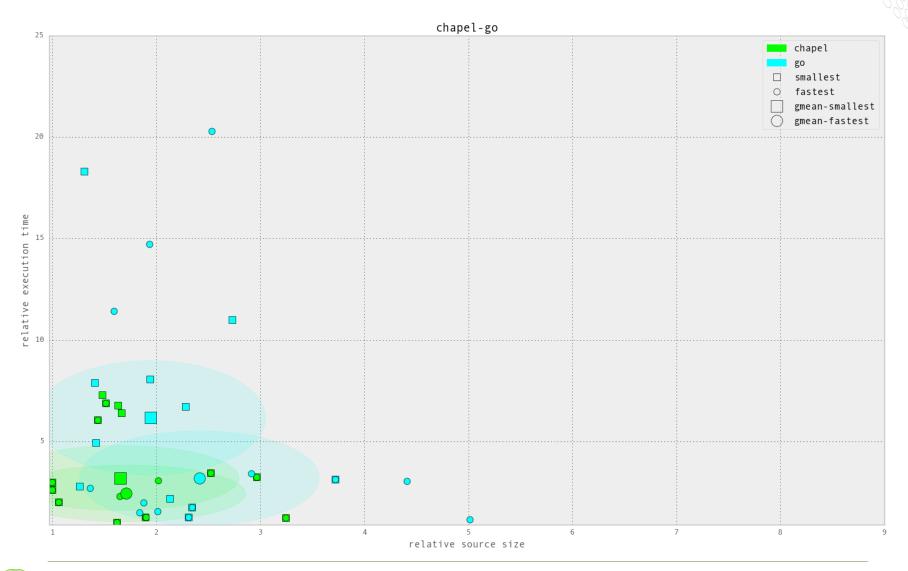






# Chapel vs. Go

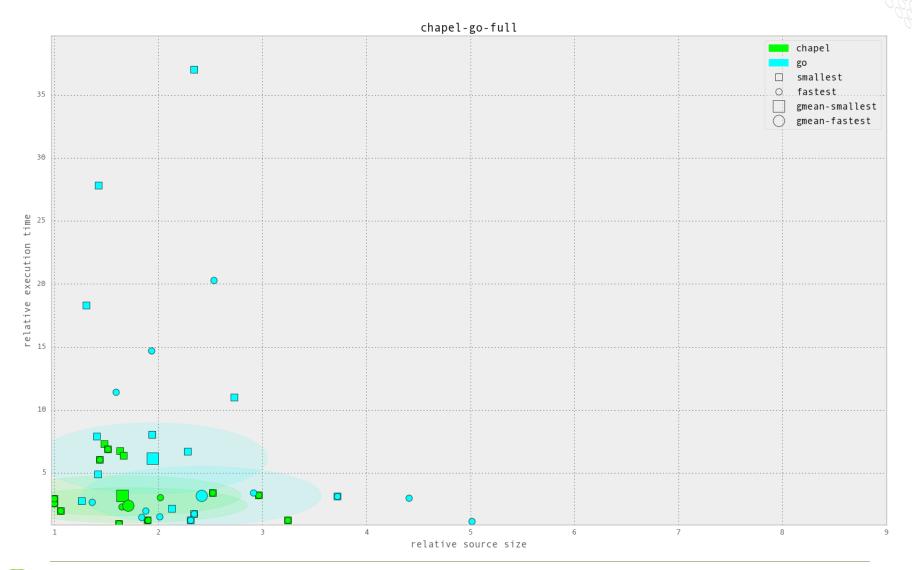






# Chapel vs. Go (zoomed out)

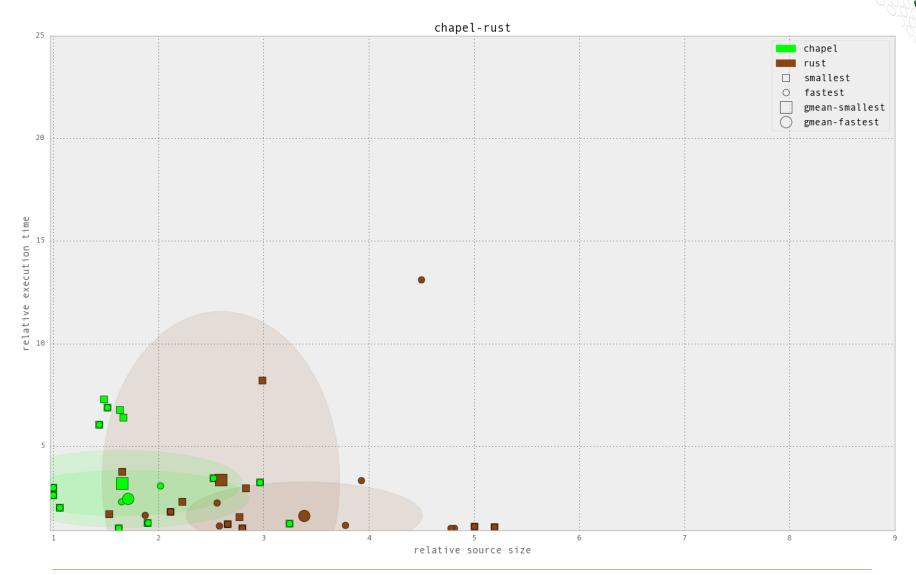






# **Chapel vs. Rust**

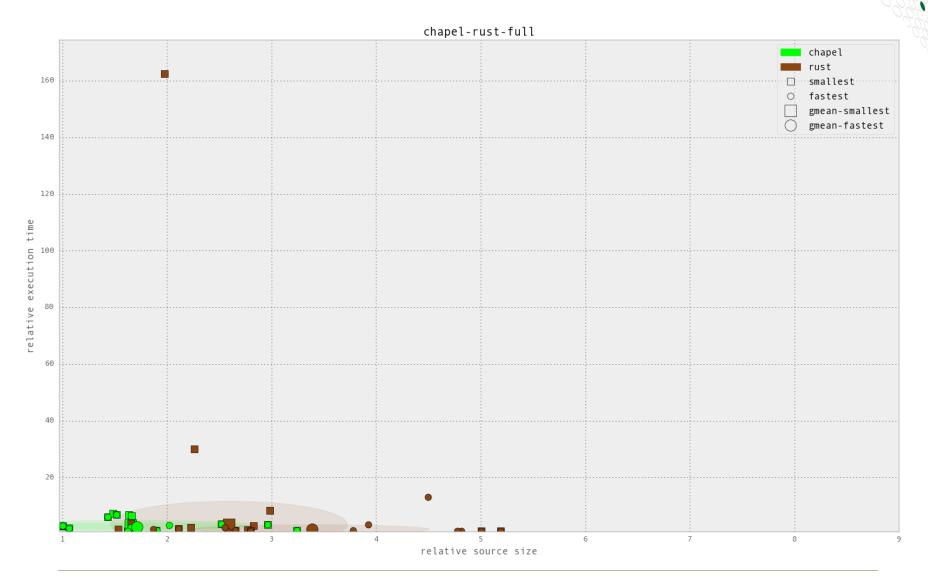






# Chapel vs. Rust (zoomed out)

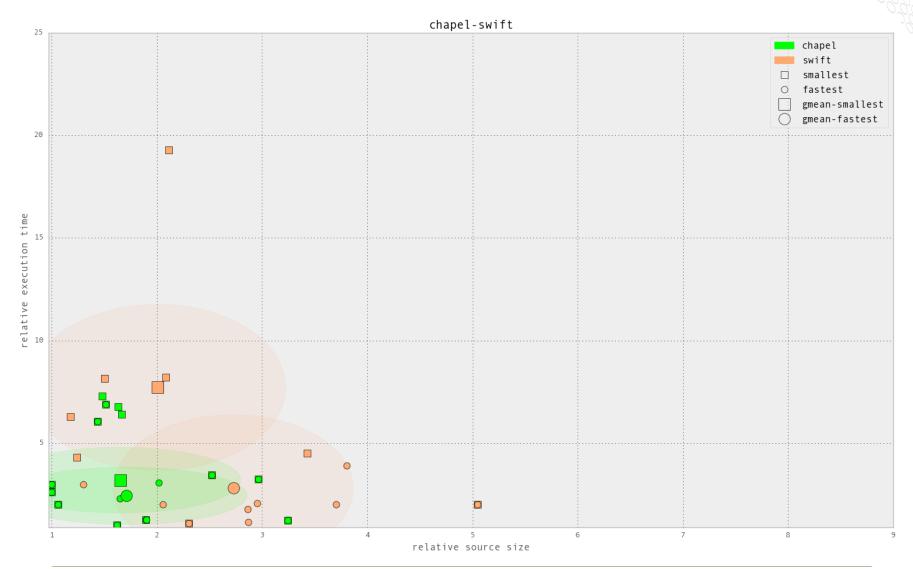






# **Chapel vs. Swift**





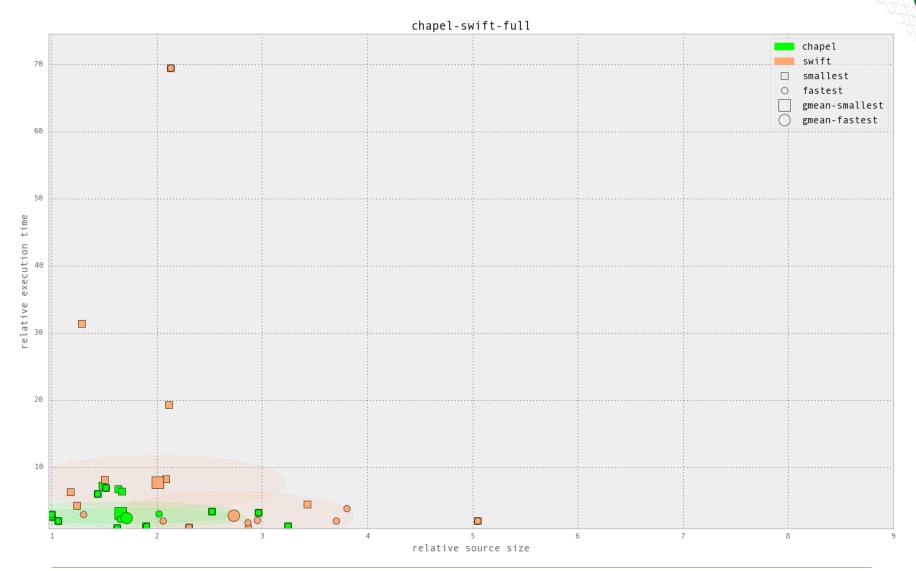


STORE

COMPUTE

# **Chapel vs. Swift (zoomed out)**



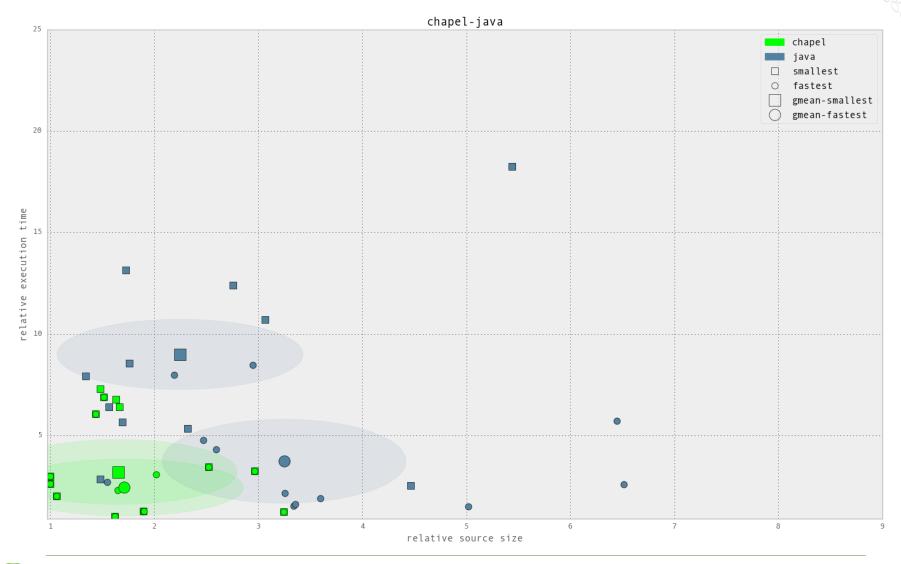




STO

# Chapel vs. Java

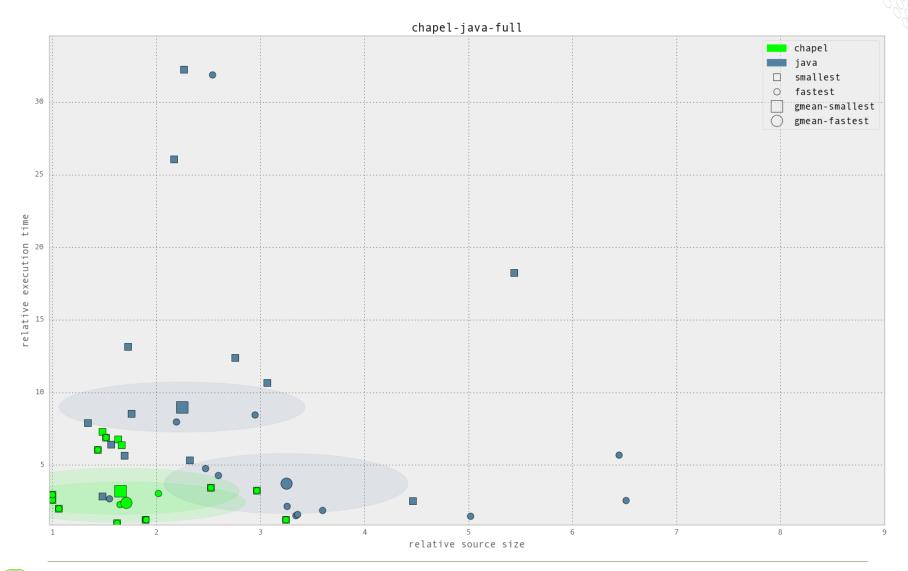






# Chapel vs. Java (zoomed out)

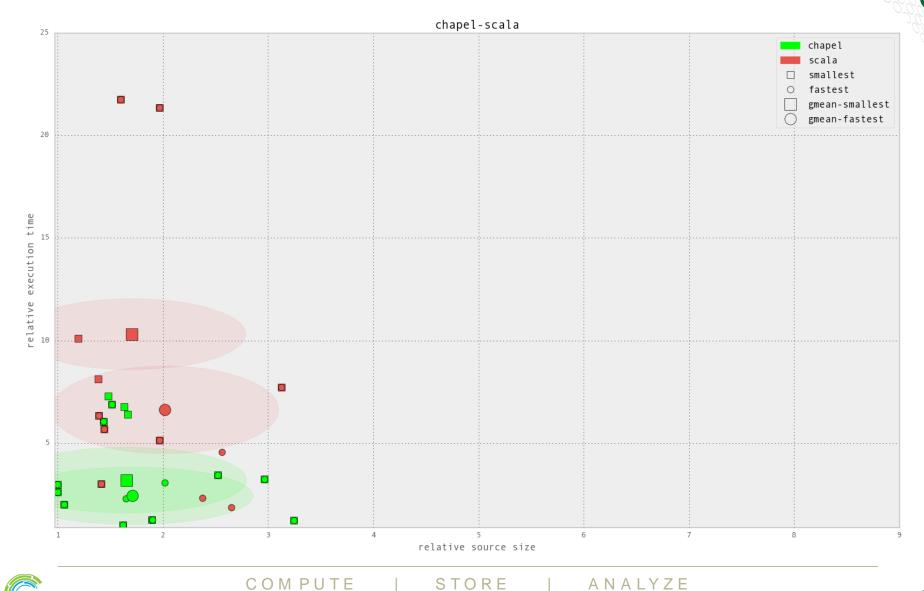






# Chapel vs. Scala



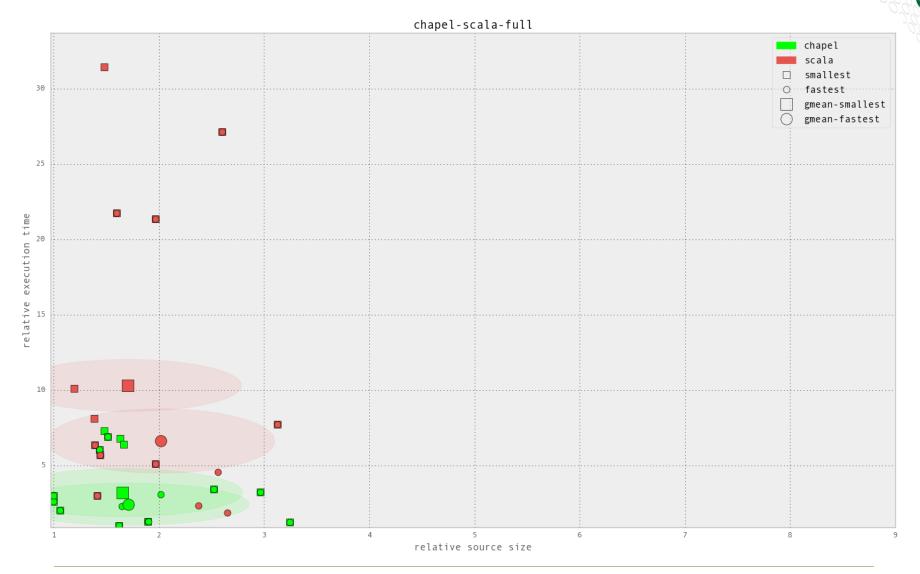




STORE

# Chapel vs. Scala (zoomed out)



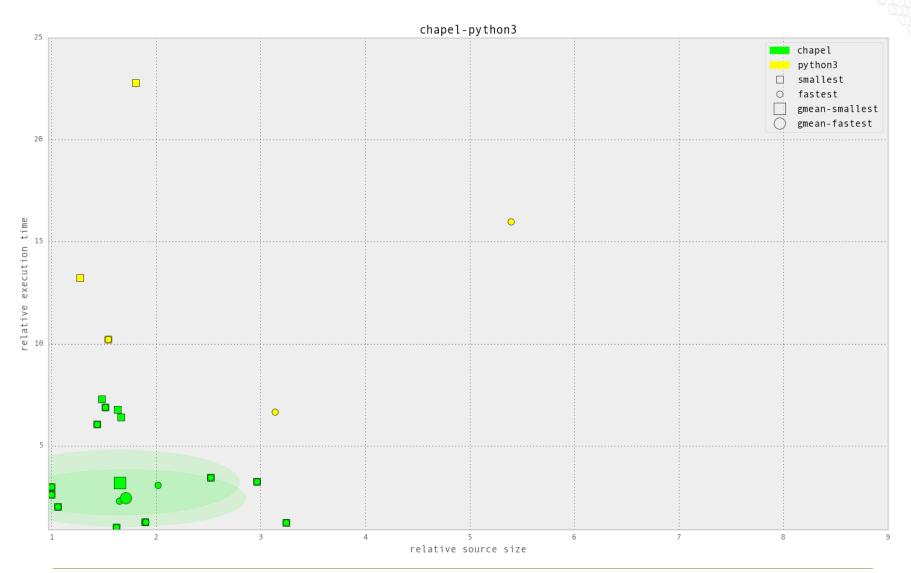




COMPUTE

# **Chapel vs. Python**

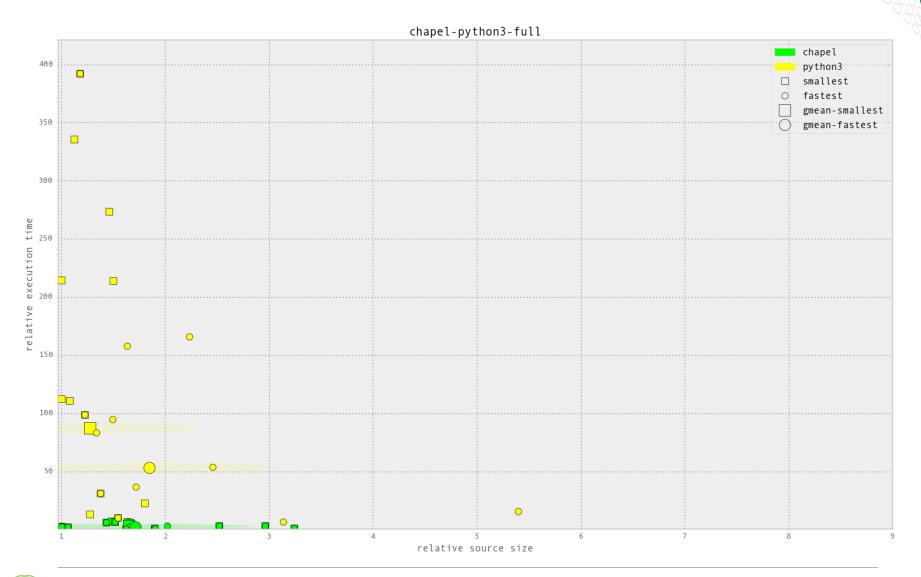






# **Chapel vs. Python (zoomed out)**







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