



# The Audacity of Chapel: Scalable Parallel Programming Done Right (director's cut, with outtakes)

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# Safe Harbor Statement



This presentation may contain forward-looking statements that are based on our current expectations. Forward looking statements may include statements about our financial guidance and expected operating results, our opportunities and future potential, our product development and new product introduction plans, our ability to expand and penetrate our addressable markets and other statements that are not historical facts. These statements are only predictions and actual results may materially vary from those projected. Please refer to Cray's documents filed with the SEC from time to time concerning factors that could affect the Company and these forward-looking statements.



# Fair Warning Statement



This keynote contains many slides and no guitar solos 😞

BUT... it does contain thematically relevant  
**UK rock lyrics trivia!!**





*I am the egg man*

## Introductions

**I am the Walrus**  
The Beatles  
*Magical Mystery Tour*



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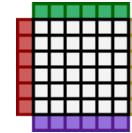
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# ("Pssst... Who is this turkey?")



## Hi, I'm Brad Chamberlain

- graduate of the University of Washington
  - worked on a data-parallel array language, ZPL
- principal engineer at Cray Inc.
  - founding member and technical lead of the Chapel project
- more of a practical parallel computing guy than a PL expert...



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# The Context for My Work



## HPC: High Performance Computing

- parallelism at large scales
  - lots of distributed processors and memories
- performance rules
  - and too often, is **all** that matters
- programmers are virtually living in the dark ages

**CRAY**: The Supercomputer Company



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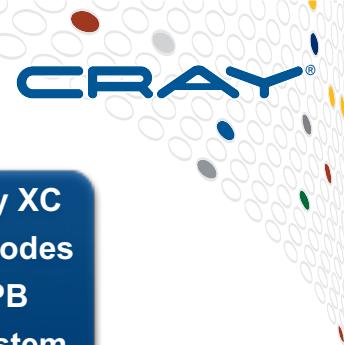
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# Recent Highlighted Cray Systems



Next-Gen Cray XC  
19k compute nodes  
40+ PF, 2+ PB  
80+PB File System  
5200 sq ft

Cray/Intel partnership  
50k+ compute nodes  
180PF, 7+ PB  
150+ PB File System



Petroleum Geo-Services

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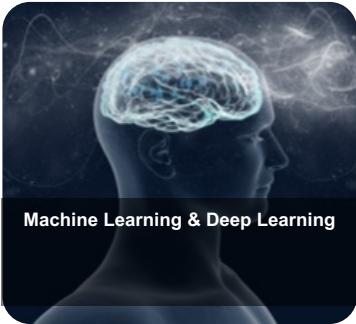
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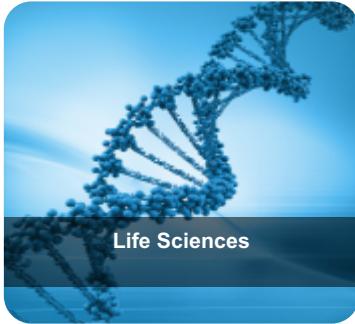
# Cray Market Segments



Manufacturing



Machine Learning & Deep Learning



Life Sciences



Higher Education



Cybersecurity



Earth Sciences



Energy



Financial Services



Government and Defense

**"I don't really care about HPC programming..."** 

- OK, but do you care about parallelism & concurrency?
  - What about performance?
  - What about scaling up your data set sizes?
  - What about targeting next-generation processors?

***Next-generation processors and computations are increasingly resembling traditional HPC.***

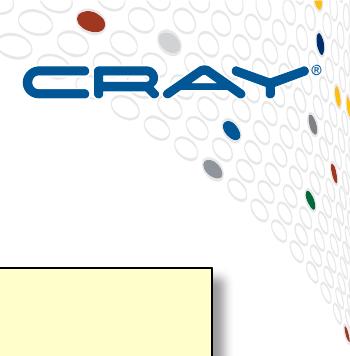


*If you didn't care what happened to me,  
And I didn't care for you,  
We would zig zag our way  
Through the boredom and pain...*

## Motivation for Chapel

**Pigs on the Wing (part one)**  
Pink Floyd  
*Animals*





## Why Consider New Languages at all?

- **Do we need a language? And a compiler?**
  - If higher level syntax is needed for productivity
    - We need a language
  - If static analysis is needed to help with correctness
    - We need a compiler (front-end)
  - If static optimizations are needed to get performance
    - We need a compiler (back-end)



# What is Chapel?



**Chapel:** A productive parallel programming language

- portable
- open-source
- a collaborative effort

## Goals:

- Support general parallel programming
  - “any parallel algorithm on any parallel hardware”
- Make parallel programming at scale far more productive



# What does “Productivity” mean to you?



## Recent Graduates:

“something similar to what I used in school: Python, Matlab, Java, ...”

## Seasoned HPC Programmers:

“that sugary stuff that I don’t need because I ~~was born to suffer~~  
want full control to ensure performance”

## Computational Scientists:

“something that lets me express my parallel computations without having to wrestle  
with architecture-specific details”

## Chapel Team:

“something that lets computational scientists express what they want,  
without taking away the control that HPC programmers want,  
implemented in a language as attractive as recent graduates want.”



# “The Audacity of Chapel”?



***audacity*** (according to Google):

/əd'asɪti/

*noun*

1. a willingness to take bold risks.

“I applaud the *audacity* of the Chapel team in attempting to create a new language given how hard it is for new languages to succeed.”

2. rude or disrespectful behaviour; impudence.

“I can’t believe the Chapel team has the *audacity* to create a new language when we already have [ C++ | Python | ... ]!”



# This Talk's Thesis



Programming language designers have, to date, largely failed the large-scale parallel computing community.



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*This just feels like spinning plates  
I'm living in cloud-cuckoo land*

## The Status Quo in HPC Programming

**Like Spinning Plates**  
Radiohead  
*Amnesiac*



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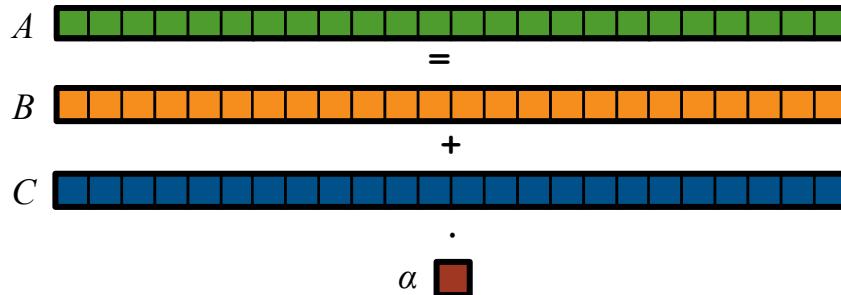
# STREAM Triad: a trivial parallel computation



**Given:**  $m$ -element vectors  $A, B, C$

**Compute:**  $\forall i \in 1..m, A_i = B_i + \alpha \cdot C_i$

**In pictures:**



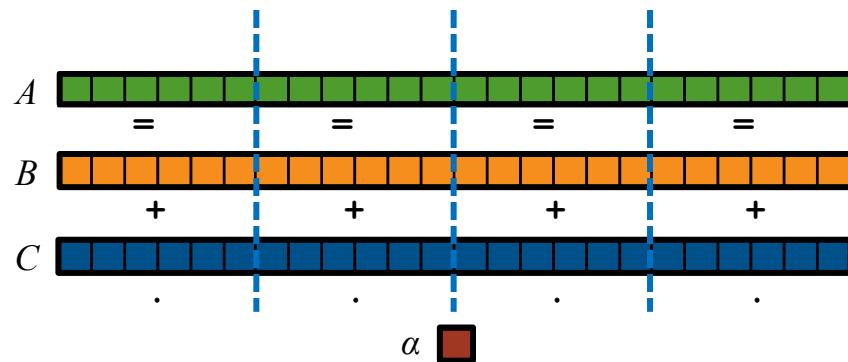
# STREAM Triad: a trivial parallel computation



**Given:**  $m$ -element vectors  $A, B, C$

**Compute:**  $\forall i \in 1..m, A_i = B_i + \alpha \cdot C_i$

**In pictures, in parallel (shared memory / multicore):**



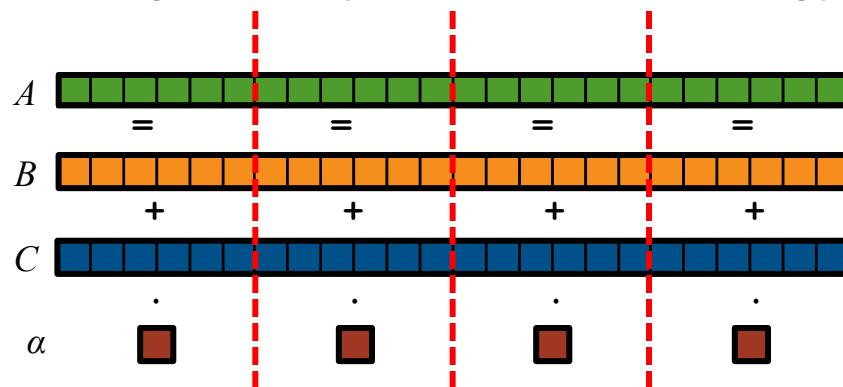
# STREAM Triad: a trivial parallel computation



**Given:**  $m$ -element vectors  $A, B, C$

**Compute:**  $\forall i \in 1..m, A_i = B_i + \alpha \cdot C_i$

**In pictures, in parallel (distributed memory):**



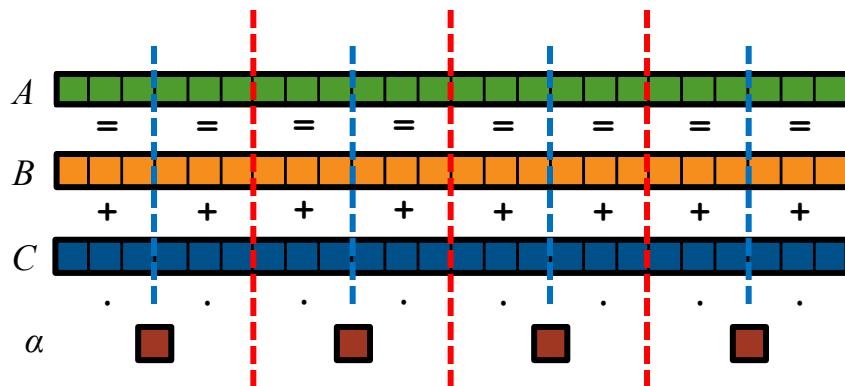
# STREAM Triad: a trivial parallel computation



**Given:**  $m$ -element vectors  $A, B, C$

**Compute:**  $\forall i \in 1..m, A_i = B_i + \alpha \cdot C_i$

**In pictures, in parallel (distributed memory multicore):**



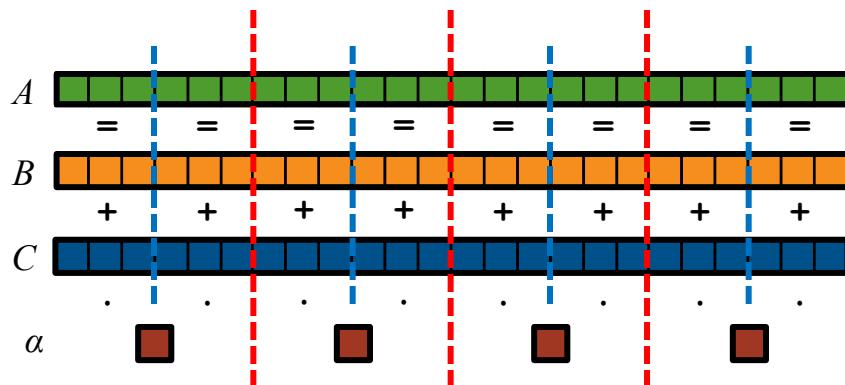
# Scalable Parallel Programming Concerns



**Q:** What should scalable parallel programmers focus on?

**A:** *Parallelism*: What should execute simultaneously?

*Locality*: Where should those tasks execute?



# STREAM Triad: MPI



```
#include <hpcc.h>
MPI
static int VectorSize;
static double *a, *b, *c;

int HPCC_StarStream(HPCC_Params *params) {
    int myRank, commSize;
    int rv, errCount;
    MPI_Comm comm = MPI_COMM_WORLD;

    MPI_Comm_size( comm, &commSize );
    MPI_Comm_rank( comm, &myRank );

    rv = HPCC_Stream( params, 0 == myRank );
    MPI_Reduce( &rv, &errCount, 1, MPI_INT, MPI_SUM,
                0, comm );
}

return errCount;
}

int HPCC_Stream(HPCC_Params *params, int doIO) {
    register int j;
    double scalar;

    VectorSize = HPCC_LocalVectorSize( params, 3,
        sizeof(double), 0 );

    a = HPCC_XMALLOC( double, VectorSize );
    b = HPCC_XMALLOC( double, VectorSize );
    c = HPCC_XMALLOC( double, VectorSize );
```

```
if (!a || !b || !c) {
    if (c) HPCC_free(c);
    if (b) HPCC_free(b);
    if (a) HPCC_free(a);
    if (doIO) {
        fprintf( outFile, "Failed to
            allocate memory (%d).\n",
            VectorSize );
        fclose( outFile );
    }
    return 1;
}

for (j=0; j<VectorSize; j++) {
    b[j] = 2.0;
    c[j] = 1.0;
}
scalar = 3.0;

for (j=0; j<VectorSize; j++)
    a[j] = b[j]+scalar*c[j];

HPCC_free(c);
HPCC_free(b);
HPCC_free(a);

return 0; }
```



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# STREAM Triad: MPI+OpenMP



```
#include <hpcc.h>
#ifndef _OPENMP
#include <omp.h>
#endif
```

```
static int VectorSize;
static double *a, *b, *c;
```

**MPI + OpenMP**

```
int HPCC_StarStream(HPCC_Params *params) {
    int myRank, commSize;
    int rv, errCount;
    MPI_Comm comm = MPI_COMM_WORLD;

    MPI_Comm_size( comm, &commSize );
    MPI_Comm_rank( comm, &myRank );

    rv = HPCC_Stream( params, 0 == myRank );
    MPI_Reduce( &rv, &errCount, 1, MPI_INT, MPI_SUM,
                0, comm );

    return errCount;
}

int HPCC_Stream(HPCC_Params *params, int doIO) {
    register int j;
    double scalar;

    VectorSize = HPCC_LocalVectorSize( params, 3,
                                       sizeof(double), 0 );

    a = HPCC_XMALLOC( double, VectorSize );
    b = HPCC_XMALLOC( double, VectorSize );
    c = HPCC_XMALLOC( double, VectorSize );
```

```
if (!a || !b || !c) {
    if (c) HPCC_free(c);
    if (b) HPCC_free(b);
    if (a) HPCC_free(a);
    if (doIO) {
        fprintf( outFile, "Failed to
allocate memory (%d).\\n",
VectorSize );
        fclose( outFile );
    }
    return 1;
}

#ifndef _OPENMP
#pragma omp parallel for
#endif
for (j=0; j<VectorSize; j++) {
    b[j] = 2.0;
    c[j] = 1.0;
}
scalar = 3.0;

#ifndef _OPENMP
#pragma omp parallel for
#endif
for (j=0; j<VectorSize; j++)
    a[j] = b[j]+scalar*c[j];

HPCC_free(c);
HPCC_free(b);
HPCC_free(a);

return 0; }
```



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# STREAM Triad: MPI+OpenMP

```
#include <hpcc.h>
#ifndef _OPENMP
#include <omp.h>
#endif
```

```
static int VectorSize;
static double *a, *b, *c;

int HPCC_StarStream(HPCC_Params *params) {
    int myRank, commSize;
    int rv, errCount;
    MPI_Comm comm = MPI_COMM_WORLD;

    MPI_Comm_size( comm, &commSize );
    MPI_Comm_rank( comm, &myRank );

    rv = HPCC_Stream( params, 0 == myRank );
    MPI_Reduce( &rv, &errCount, 1, MPI_INT, MPI_SUM,
        0, comm );

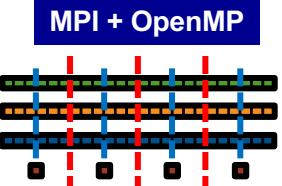
    return errCount;
}

int HPCC_Stream(HPCC_Params *params, int doIO) {
    register int j;
    double scalar;

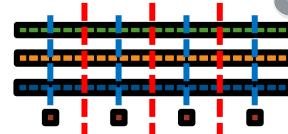
    VectorSize = HPCC_LocalVectorSize( params, 3,
        sizeof(double), 0 );

    a = HPCC_XMALLOC( double, VectorSize );
    b = HPCC_XMALLOC( double, VectorSize );
    c = HPCC_XMALLOC( double, VectorSize );
```

**MPI + OpenMP**



**CUDA**



```
if (!a || !b) #defined N 2000000
    if (c) HPC_free(c);
    if (b) HPC_free(b);
    if (a) HPC_free(a);
    if (doIO) {
        float scalar;
        fprintf( outFile, "Failed to
            allocate cudaMemcpy(&d_a, sizeof(float)*N);
            cudaMemcpy(&d_b, sizeof(float)*N);
            VectorSize cudaMemcpy(&d_c, sizeof(float)*N);
            fclose(outFile );
        }
        return 1;
    }

    dim3 dimBlock(128);
    dim3 dimGrid(N/dimBlock.x );
    if ( N % dimBlock.x != 0 ) dimGrid

    #ifdef _OPENMP
    #pragma omp parallel for
    #endif
    scalar=3.0f;
    for (j=0; j<VectorSize; j++) {
        b[j] = 2.0;
        c[j] = 1.0;
    }
    cudaFree(d_a);
    scalar = 3.0;
    cudaFree(d_b);
    cudaFree(d_c);

    #ifdef _OPENMP
    #pragma omp parallel for
    #endif
    __global__ void set_array(float *a, float value, int len) {
        for (j=0; j<VectorSize; j++)
            if (idx < len) a[idx] = value;
    }

    HPC_free(c);
    HPC_free(b);
    HPC_free(a);
    __global__ void STREAM_Triad( float *a, float *b, float *c,
                                float scalar, int len) {
        int idx = threadIdx.x + blockIdx.x * blockDim.x;
        if (idx < len) c[idx] = a[idx]+scalar*b[idx]; }
```

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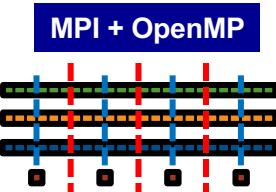
# STREAM Triad: MPI+OpenMP



```
#include <hpcc.h>
#ifndef _OPENMP
#include <omp.h>
#endif
```

```
static int VectorSize;
static double *a, *b, *c;

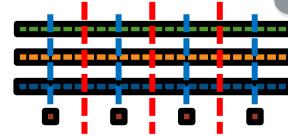
int HPCC_StarStream(HPCC_Params *params) {
    int myRank, commSize;
```



```
if (!a || !b || !c) {
    if (c) HPCC_free(c);
    if (b) HPCC_free(b);
    if (a) HPCC_free(a);
    if (doIO) {
        fprintf( outFile, "Failed to
allocate memory (%d).\n",
VectorSize );
        fclose( outFile );
    }
}
```

```
#define N 2000000
int main() {
    float *d_a, *d_b, *d_c;
    float scalar;
    cudaMalloc((void**)&d_a, sizeof(float)*N);
    cudaMalloc((void**)&d_b, sizeof(float)*N);
    cudaMalloc((void**)&d_c, sizeof(float)*N);

    dim3 dimBlock(128);
```



*HPC suffers from too many distinct notations for expressing parallelism and locality.  
This tends to be a result of bottom-up language design.*

```
    HPCC_SrcComm( params, 0, myRank,
MPI_Reduce( &rv, &errCount, 1, MPI_INT, MPI_SUM,
0, comm );
return errCount;
}
```

```
int HPCC_Stream(HPCC_Params *params, int doIO) {
register int j;
double scalar;

VectorSize = HPCC_LocalVectorSize( params, 3,
sizeof(double), 0 );

a = HPCC_XMALLOC( double, VectorSize );
b = HPCC_XMALLOC( double, VectorSize );
c = HPCC_XMALLOC( double, VectorSize );

for (j=0; j<VectorSize; j++) {
    b[j] = 2.0;
    c[j] = 1.0;
}
scalar = 3.0;

#ifdef _OPENMP
#pragma omp parallel for
#endif
for (j=0; j<VectorSize; j++)
    a[j] = b[j]+scalar*c[j];

HPCC_free(c);
HPCC_free(b);
HPCC_free(a);

return 0;
}
```

```
STREAM_Triad<<<dimGrid, dimBlock>>>(d_b, d_c, d_a, scalar, N);
cudaThreadSynchronize();

cudaFree(d_a);
cudaFree(d_b);
cudaFree(d_c);

_global_ void set_array(float *a, float value, int len) {
int idx = threadIdx.x + blockIdx.x * blockDim.x;
if (idx < len) a[idx] = value;
}

_global_ void STREAM_Triad( float *a, float *b, float *c,
                           float scalar, int len) {
int idx = threadIdx.x + blockIdx.x * blockDim.x;
if (idx < len) c[idx] = a[idx]+scalar*b[idx]; }
```



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# Why so many programming models?



## HPC tends to approach programming models bottom-up:

Given a system and its core capabilities...

...provide features that permit users to access the available performance.

- portability? generality? programmability? These are second- or third-order concerns, if that.

Type of HW Parallelism	Programming Model	Unit of Parallelism
Inter-node	MPI	executable
Intra-node/multicore	OpenMP / pthreads	iteration/task
Instruction-level vectors/threads	pragmas	iteration
GPU/accelerator	CUDA / Open[MP CL ACC]	SIMD function/task

**benefits:** lots of control; decent generality; easy to implement

**downsides:** lots of user-managed detail; brittle to changes



# STREAM Triad: Chapel



```
#include <hpcc.h>
#ifndef _OPENMP
#include <omp.h>
#endif

static int VectorSize;
static double *a, *b, *c;

int HPCC_StarStream(HPCC_Params *params)
    int myRank, commSize;
    int rv, errCount;
    MPI_Comm comm = MPI_COMM_WORLD;
    MPI_Comm_size( comm, &commSize );
    MPI_Comm_rank( comm, &myRank );

    rv = HPCC_Stream( params, 0 == myRank );
    MPI_Reduce( &rv, &errCount, 1, MPI_INT,
        0, comm );
    return errCount;
}

int HPCC_Stream(HPCC_Params *params)
    register int j;
    double scalar;
    VectorSize = HPCC(sizeof(double),
        a = HPCC_XMALLOC(sizeof(double)*VectorSize);
        b = HPCC_XMALLOC(sizeof(double)*VectorSize);
        c = HPCC_XMALLOC(sizeof(double)*VectorSize);

        for(j=0; j<VectorSize; j++) {
            a[j] = 1.0;
            b[j] = 2.0;
            c[j] = 3.0;
        }

        for(j=0; j<VectorSize; j+=16) {
            for(int i=0; i<16; i++) {
                a[j+i] = b[j+i] + alpha * c[j+i];
            }
        }

        MPI_Reduce( &rv, &errCount, 1, MPI_INT,
            0, comm );
    }
    return errCount;
}
```

```
use ...;

config const m = 1000,
          alpha = 3.0;

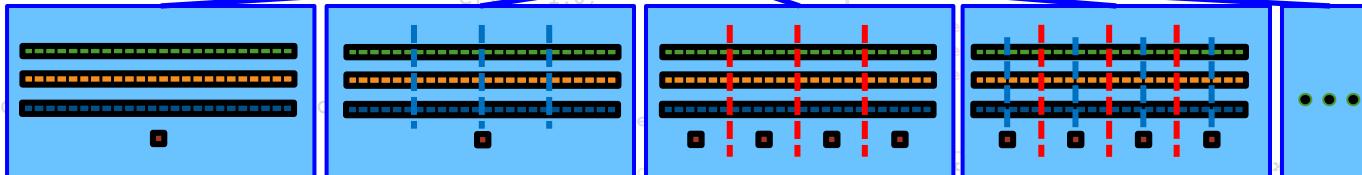
const ProblemSpace = {1..m} dmapped ...;

var A, B, C: [ProblemSpace] real;

B = 2.0;
C = 1.0;

A = B + alpha * C;
```

**The special sauce:**  
How should this index set—and any arrays and computations over it—be mapped to the system?



Philosophy: Good, *top-down* language design can tease system-specific implementation details away from an algorithm, permitting the compiler, runtime, applied scientist, and HPC expert to each focus on their strengths.

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# This Talk's Takeaways



If you design a parallel programming language...

...don't tie yourself to low-level, architecture-specific mechanisms

- yet don't make them inaccessible either...
  - permit interoperating with such mechanisms
  - or support them as the “assembly” to your higher-level features



*It's so easy to laugh, it's so easy to hate  
It takes guts to be gentle and kind*

## SPMD Programming Models like MPI

**I Know It's Over**  
The Smiths  
*The Queen is Dead*



# HPC's Status Quo: SPMD Programming



## ***SPMD:*** Single Program, Multiple Data

- concept: write one program, run multiple copies of it in parallel
- a “bottom-up” programming model design
  - “HPC systems can run lots of programs, so let’s get parallelism that way”
- often clumsy in practice



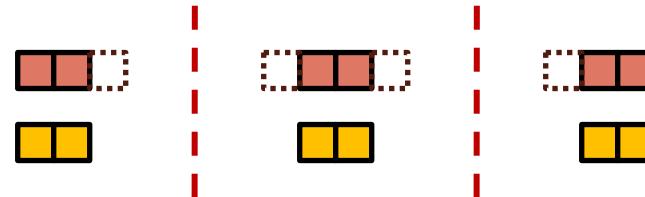
# SPMD by Example (in pictures)

“Apply a 3-Point Stencil to a vector”

*Conceptual View*

$$\begin{aligned} & \left( \begin{array}{cccccc} \text{brown} & \text{brown} & \text{brown} & \text{pink} & \text{pink} \end{array} \right) \\ & + \left( \begin{array}{ccccc} \text{pink} & \text{pink} & \text{brown} & \text{brown} & \text{brown} \end{array} \right) / 2 \\ & = \left( \begin{array}{cccccc} \text{yellow} & \text{orange} & \text{orange} & \text{orange} & \text{yellow} \end{array} \right) \end{aligned}$$

*SPMD View*



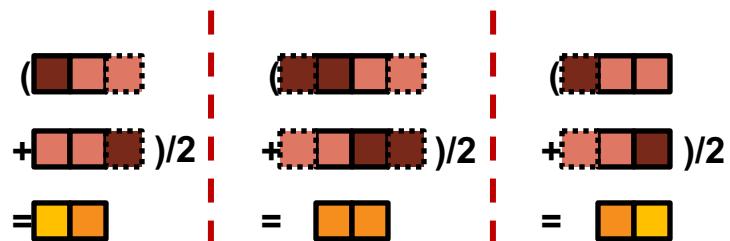
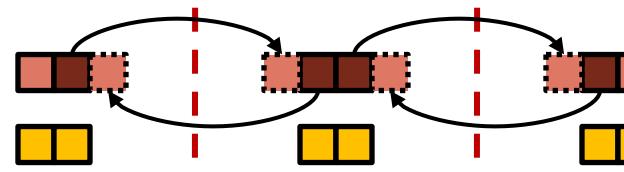
# SPMD by Example (in pictures)

“Apply a 3-Point Stencil to a vector”

*Conceptual View*

$$\begin{aligned}
 & \left( \begin{array}{cccccc} \text{brown} & \text{brown} & \text{brown} & \text{pink} & \text{pink} \end{array} \right) \\
 & + \left( \begin{array}{ccccc} \text{pink} & \text{brown} & \text{brown} & \text{brown} & \text{brown} \end{array} \right) / 2 \\
 & = \left( \begin{array}{cccccc} \text{yellow} & \text{orange} & \text{orange} & \text{orange} & \text{orange} & \text{yellow} \end{array} \right)
 \end{aligned}$$

*SPMD View*



# SPMD by Example (in code)



“Apply a 3-Point Stencil to a vector”



## SPMD pseudo-code

```
proc main() {
    var n = 1000;
    var p = numProcs(),
        me = myProc(),
        myN = n/p,
    var A, B: [0..myN+1] real;

    if (me < p-1) {
        send(me+1, A[myN]);
        recv(me+1, A[myN+1]);
    }
    if (me > 0) {
        send(me-1, A[1]);
        recv(me-1, A[0]);
    }
    forall i in 1..myN do
        B[i] = (A[i-1] + A[i+1])/2;
}
```



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# SPMD by Example (in code)



“Apply a 3-Point Stencil to a vector”

## Global-view code (Chapel)

```
proc main() {
    const n = 1000,
          D = {1..n} dmapped Block(...);
    var A, B: [1..n] real;

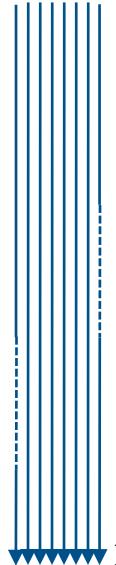
    forall i in 2..n-1 do
        B[i] = (A[i-1] + A[i+1])/2;
}
```



## SPMD pseudo-code

```
proc main() {
    var n = 1000;
    var p = numProcs(),
        me = myProc(),
        myN = n/p,
    var A, B: [0..myN+1] real;

    if (me < p-1) {
        send(me+1, A[myN]);
        recv(me+1, A[myN+1]);
    }
    if (me > 0) {
        send(me-1, A[1]);
        recv(me-1, A[0]);
    }
    forall i in 1..myN do
        B[i] = (A[i-1] + A[i+1])/2;
}
```



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# SPMD by Example (in code)



“Apply a 3-Point Stencil to a vector”

## Global-view code (Chapel)

```
proc main() {
    const n = 1000,
          D = {1..n} dmapped Block(...);
    var A, B: [1..n] real;

    forall i in 2..n-1 do
        }
```

Bug: Refers to uninitialized values at ends of A

## SPMD pseudo-code

```
proc main() {
    var n = 1000;
    var p = numProcs(),
        me = myProc(),
        myN = n/p,
    var A, B: [0..myN+1] real;

    if (me < p-1) {
        send(me+1, A[myN]);
        recv(me+1, A[myN+1]);
    }
    if (me > 0) {
        send(me-1, A[1]);
        recv(me-1, A[0]);
    }
    forall i in 1..myN do
        B[i] = (A[i-1] + A[i+1])/2;
    }
```



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# SPMD by Example (in code)



“Apply a 3-Point Stencil to a vector”

## Global-view code (Chapel)

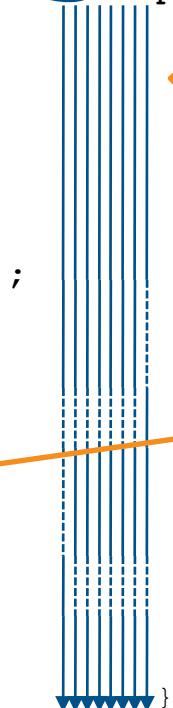
```
proc main() {  
    const n = 1000,  
          D = {1..n} dmapped Block(...);  
    var A, B: [1..n] real;  
  
    forall B[i]
```

Communication becomes  
geometrically more complex  
for higher-dimensional arrays

## SPMD pseudo-code

```
proc main() {  
    var n = 1000;  
    var p = numProcs(),  
        me = myProc(),  
        myN = n/p,  
        myLo = 1,  
        myHi = myN;  
    var A, B: [0..myN+1] real;  
  
    if (me < p-1) {  
        send(me+1, A[myN]);  
        recv(me+1, A[myN+1]);  
    } else  
        myHi = myN-1;  
    if (me > 0) {  
        send(me-1, A[1]);  
        recv(me-1, A[0]);  
    } else  
        myLo = 2;  
    forall i in myLo..myHi do  
        B[i] = (A[i-1] + A[i+1])/2;
```

Assumes p divides n



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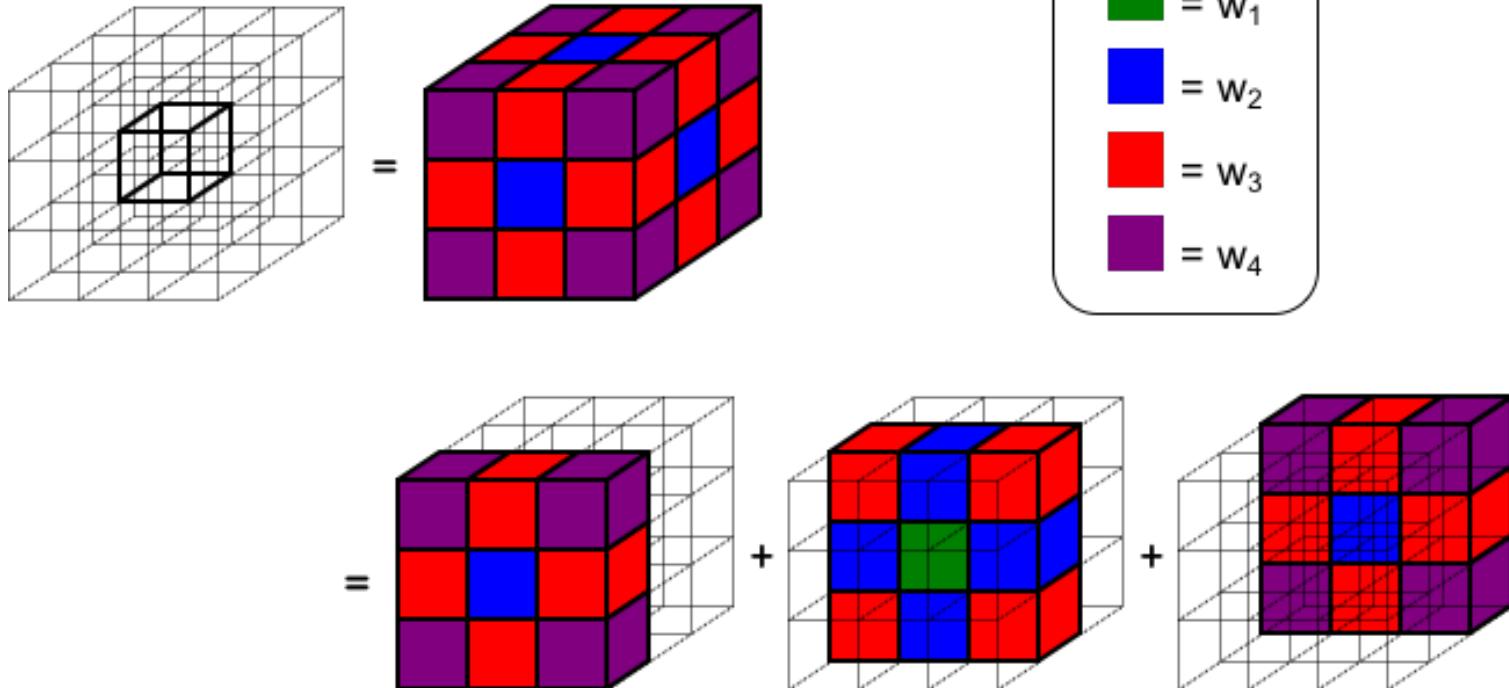


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# 27-point stencils (*rprj3* from NAS MG)



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# rprj3 in Fortran + MPI



```
subroutine rprj3(r,m1k,m2k,m3k,s,m1j,m2j,m3j,k)
implicit none
include 'cafnpb.h'
include 'globals.h'

integer m1k, m2k, m3k, m1j, m2j, m3j, k

double precision r(m1k,m2k,m3k), s(m1j,m2j,m3j)
integer j3, j2, j1, i3, i2, i1, d1, d2, d3, j
double precision x1(m), y1(m), x2,y2

if(m1k.eq.3)then
  d1 = 2
else
  d1 = 1
endif

if(m2k.eq.3)then
  d2 = 2
else
  d2 = 1
endif

if(m3k.eq.3)then
  d3 = 2
else
  d3 = 1
endif
```

```
do j3=2,m3j-1
  i3 = 2*j3-d3
  do j2=2,m2j-1
    i2 = 2*j2-d2
    do j1=2,m1j
      i1 = 2*j1-d1
      x1(i1-1) = r(i1-1,i2-1,i3 ) + r(i1-1,i2+1,i3 )
      >           + r(i1-1,i2, i3-1) + r(i1-1,i2, i3+1)
      y1(i1-1) = r(i1-1,i2-1,i3-1) + r(i1-1,i2-1,i3+1)
      >           + r(i1-1,i2+1,i3-1) + r(i1-1,i2+1,i3+1)
    enddo
    do j1=2,m1j-1
      i1 = 2*j1-d1
      y2 = r(i1, i2-1,i3-1) + r(i1, i2-1,i3+1)
      >           + r(i1, i2+1,i3-1) + r(i1, i2+1,i3+1)
      x2 = r(i1, i2-1,i3 ) + r(i1, i2+1,i3 )
      >           + r(i1, i2, i3-1) + r(i1, i2, i3+1)
      s(j1,j2,j3) =
        > 0.5D0 * r(i1,i2,i3)
        > + 0.25D0 * (r(i1-1,i2,i3) + r(i1+1,i2,i3) + x2)
        > + 0.125D0 * ( x1(i1-1) + x1(i1+1) + y2)
        > + 0.0625D0 * ( y1(i1-1) + y1(i1+1) )
    enddo
  enddo
enddo
j = k-1
call comm3(s,m1j,m2j,m3j,j)
return
end
```



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# rprj3 in Fortran + MPI

```

subroutine rprj3(r,m1k,m2k,m3k,s,m1j,m2j,m3j,k)
implicit none
include 'cafnpb.h'
include 'globals.h'

integer m1k, m2k, m3k, m1j, m2j, m3j, k

double precision r(m1k,m2k,m3k), s(m1j,m2j,m3j)
integer j3, j2, j1, i3, i2, i1, d1, d2, d3, j
double precision x1(m), y1(m), x2,y2

if(m1k.eq.3) then
    d1 = 2
else
    d1 = 1
endif

if(m2k.eq.3) then
    d2 = 2
else
    d2 = 1
endif

if(m3k.eq.3) then
    d3 = 2
else
    d3 = 1
endif

```

```

do j3=2,m3j-1
    i3 = 2*j3-d3
    do j2=2,m2j-1
        i2 = 2*j2-d2
        do j1=2,m1j
            i1 = 2*j1-d1
            x1(i1-1) = r(i1-1,i2-1,i3 ) + r(i1-1,i2+1,i3 )
            >           + r(i1-1,i2, i3-1) + r(i1-1,i2, i3+1)
            y1(i1-1) = r(i1-1,i2-1,i3-1) + r(i1-1,i2-1,i3+1)
            >           + r(i1-1,i2+1,i3-1) + r(i1-1,i2+1,i3+1)
        enddo
        do j1=2,m1j-1
            i1 = 2*j1-d1
            y2 = r(i1, i2-1,i3-1) + r(i1, i2-1,i3+1)
            >           + r(i1, i2+1,i3-1) + r(i1, i2+1,i3+1)
            x2 = r(i1, i2-1,i3 ) + r(i1, i2+1,i3 )
            >           + r(i1, i2, i3-1) + r(i1, i2, i3+1)
            s(j1,j2,j3) =
            >           0.5D0 * r(i1,i2,i3)
            >           + 0.25D0 * (r(i1-1,i2,i3) + r(i1+1,i2,i3) + x2)
            >           + 0.125D0 * ( x1(i1-1) + x1(i1+1) + y2)
            >           + 0.0625D0 * ( y1(i1-1) + y1(i1+1) )
        enddo
        enddo
    j = k-1
call comm3(s,m1j,m2j,m3j,j)
return
end

```



# ***comm3: the communication for rprj3***



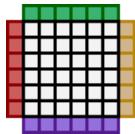
# Being Gutsy, Gentle, and Kind to MPI



- It's enabled the vast majority of HPC results for the past ~20 years
- It's very analogous to assembly programming
  - explicitly move data from memory to registers  
vs.  
explicitly move data between compute nodes' memories
- Just like assembly, it's an important technology
  - for programming at low levels
  - for enabling higher-level technologies
- Yet, as with assembly, we should develop higher-level alternatives



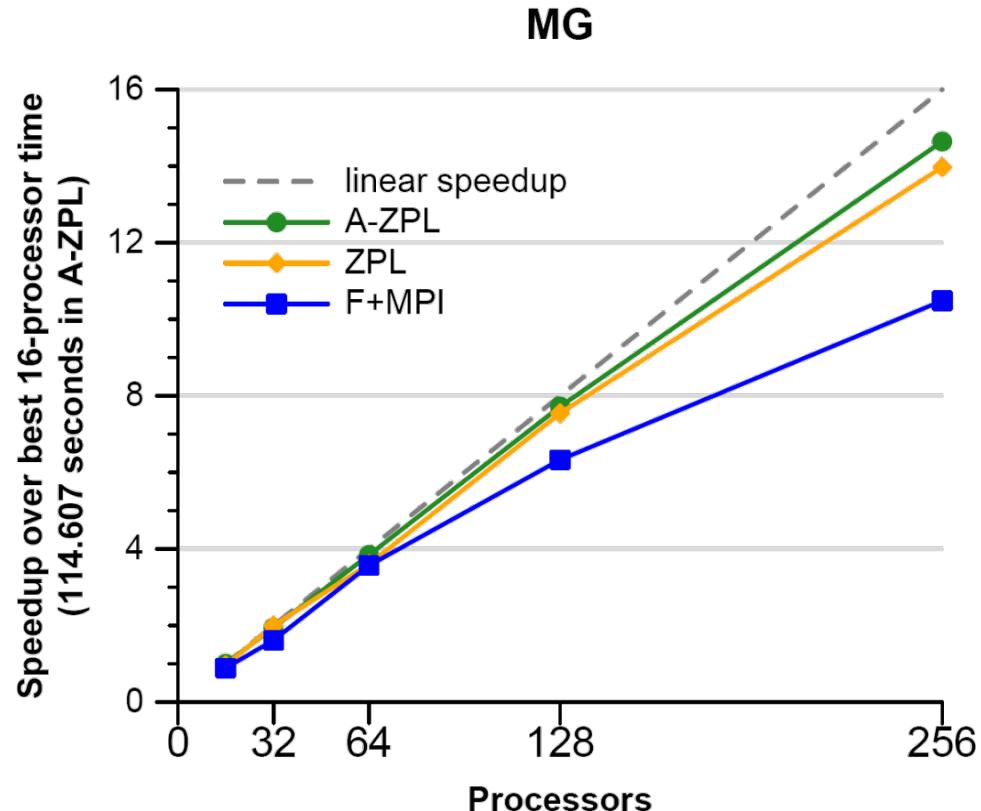
# rprj3 in ZPL



```
procedure rprj3(var S,R: [,,] double;
                 d: array [] of direction);
begin
  S := 0.5000 * R +
        0.2500 * (R@^d[ 1, 0, 0] + R@^d[ 0, 1, 0] + R@^d[ 0, 0, 1] +
                   R@^d[-1, 0, 0] + R@^d[ 0,-1, 0] + R@^d[ 0, 0,-1] +
  0.1250 * (R@^d[ 1, 1, 0] + R@^d[ 1, 0, 1] + R@^d[ 0, 1, 1] +
                   R@^d[ 1,-1, 0] + R@^d[ 1, 0,-1] + R@^d[ 0, 1,-1] +
                   R@^d[-1, 1, 0] + R@^d[-1, 0, 1] + R@^d[ 0,-1, 1] +
                   R@^d[-1,-1, 0] + R@^d[-1, 0,-1] + R@^d[ 0,-1,-1]) +
  0.0625 * (R@^d[ 1, 1, 1] + R@^d[ 1, 1,-1] +
                   R@^d[ 1,-1, 1] + R@^d[ 1,-1,-1] +
                   R@^d[-1, 1, 1] + R@^d[-1, 1,-1] +
                   R@^d[-1,-1, 1] + R@^d[-1,-1,-1]);
end;
```



# NAS MG Speedup: Cray T3E



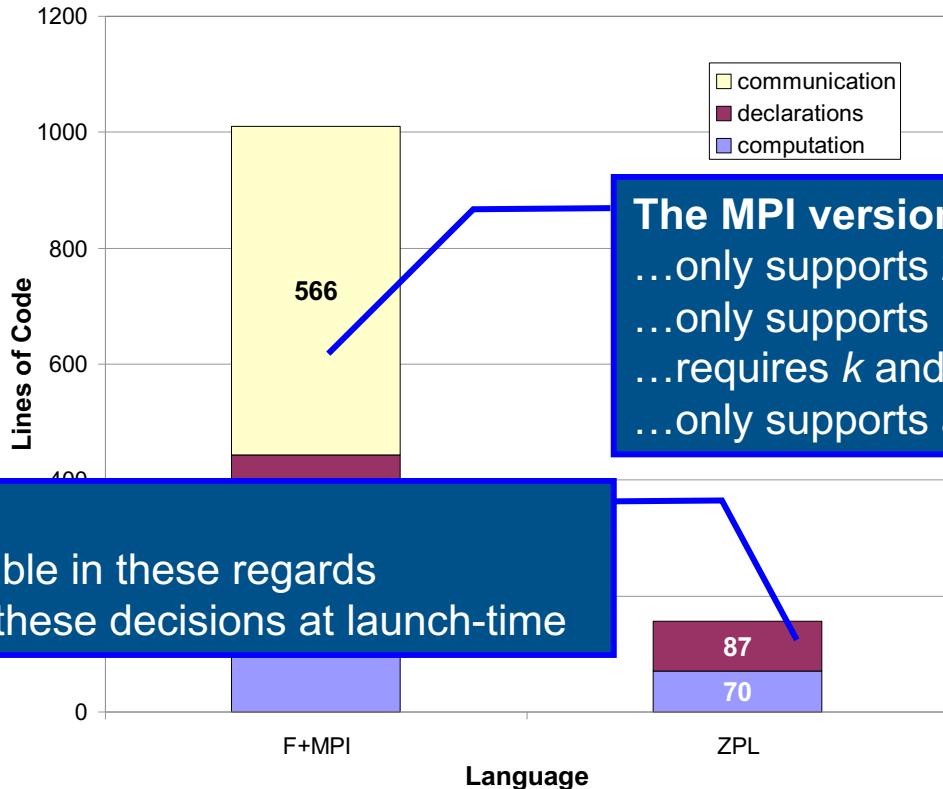
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# Code Size Comparison



# This Talk's Takeaways



If you design a parallel programming language...

...don't base your model for parallelism and locality on SPMD

- instead, support a global view of parallelism and locality (like ZPL)



# Epilogue: So why was ZPL not adopted?



- **Too restricted in terms of generality:**
  - only a single level of array-based data parallelism
  - only a single parallel array type: block-distributed
  - lack of “manual overrides” to drop down closer to the system
  - choices that were dated (no OOP, modula-based syntax, ...)
- **Great academic project, not a great practical language**
  - however, ZPL’s experiences informed Chapel’s design greatly



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A color photograph of Rowan Atkinson as Mr. Bean. He is standing behind a dark wooden podium, wearing a black tuxedo and white shirt. He has his hands clasped in front of him and is looking slightly downwards. On the podium to his left is a silver teapot and a small glass. To his right is a white telephone. The background is a dark, textured wall.

And now for something  
completely different.

## *[Instrumental]*

# Chapel Characteristics

**The Liberty Bell March**  
John Philip Sousa  
(unreleased?)



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# Chapel's Goal



## To create a language that is...

- ...as productive as Python
- ...as fast as Fortran
- ...as portable as C
- ...as scalable as MPI
- ...as fun as *[insert your favorite language here]*



# The Challenge



**Q: So why don't we already have such a language already?**

**A: ~~Technical challenges?~~**

- while they exist, we don't think this is the main issue...

**A: Due to a lack of...**

...long-term efforts

...resources

...community will

...co-design between developers and users

...patience

***Chapel is our attempt to reverse this trend***



# Chapel is Portable



- Chapel's design and implementation are hardware-independent
- The current release requires:
  - a C/C++ compiler
  - a \*NIX environment: Linux, Mac OS X, Windows 10 w/ bash, Cygwin, ...
  - POSIX threads
  - UDP, MPI, or RDMA (if distributed memory execution is desired)
- Chapel runs on...
  - ...laptops and workstations
  - ...commodity clusters
  - ...the cloud
  - ...HPC systems from Cray and other vendors
  - ...modern processors like Intel Xeon Phi, GPUs\*, etc.

\* = academic work only; not yet supported in the official release



# Chapel is Open-Source



- Chapel's development is hosted at GitHub
  - <https://github.com/chapel-lang>
- Chapel is licensed as Apache v2.0 software
- Instructions for download + install are online
  - see <http://chapel.cray.com/download.html> to get started



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# The Chapel Team at Cray (May 2016)



14 full-time employees + 2 summer interns + occasional visiting academics  
(one of each started after photo taken)



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# Chapel Community R&D Efforts



Lawrence Berkeley  
National Laboratory



Yale

(and several others...)

<http://chapel.cray.com/collaborations.html>



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*You say you got a real solution  
Well, you know  
We'd all love to see the plan*

## Chapel in a Nutshell

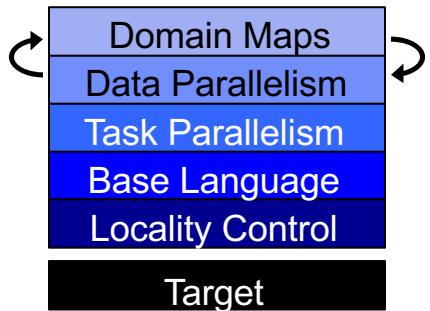
**Revolution**  
The Beatles  
*The Beatles*



# Chapel language feature areas



*Chapel language concepts*



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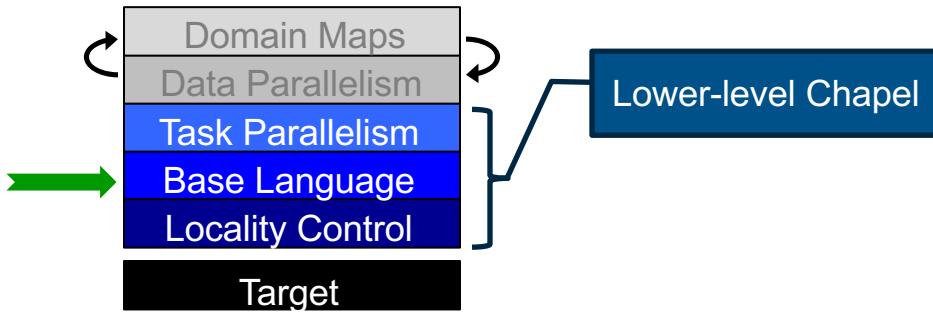
|

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# Base Language



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# Base Language Features, by example



```
iter fib(n) {  
    var current = 0,  
        next = 1;  
  
    for i in 1..n {  
        yield current;  
        current += next;  
        current <=gt; next;  
    }  
}
```

```
config const n = 10;  
  
for f in fib(n) do  
    writeln(f);
```

```
0  
1  
1  
2  
3  
5  
8  
...
```



# Base Language Features, by example



## Modern iterators

```
iter fib(n) {
    var current = 0,
        next = 1;

    for i in 1..n {
        yield current;
        current += next;
        current <= next;
    }
}
```

```
config const n = 10;

for f in fib(n) do
    writeln(f);
```

```
0  
1  
1  
2  
3  
5  
8  
...
```



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# Base Language Features, by example



Configuration declarations  
(to avoid command-line argument parsing)  
./a.out --n=1000000

```
iter fib(n) {
    var current = 0,
        next = 1;

    for i in 1..n {
        yield current;
        current += next;
        current <=> next;
    }
}
```

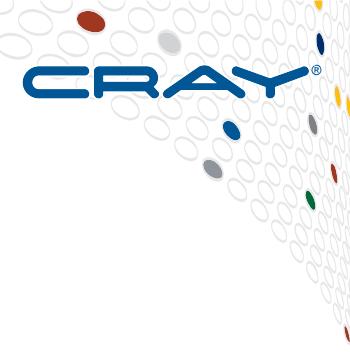
```
config const n = 10;

for f in fib(n) do
    writeln(f);
```

```
0
1
1
2
3
5
8
...
...
```



# Base Language Features, by example



Static type inference for:

- arguments
- return types
- variables

```
iter fib(n) {
    var current = 0,
        next = 1;

    for i in 1..n {
        yield current;
        current += next;
        current <= next;
    }
}
```

```
config const n = 10;

for f in fib(n) do
    writeln(f);
```

```
0  
1  
1  
2  
3  
5  
8  
...
```



# Base Language Features, by example



```
iter fib(n) {
    var current = 0,
        next = 1;

    for i in 1..n {
        yield current;
        current += next;
        current <= next;
    }
}
```

```
config const n = 10;

for (i,f) in zip(0..#n, fib(n)) do
    writeln("fib #", i, " is ", f);
```

Zippered iteration

```
fib #0 is 0
fib #1 is 1
fib #2 is 1
fib #3 is 2
fib #4 is 3
fib #5 is 5
fib #6 is 8
...
```



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# Base Language Features, by example



Range types and operators

```
iter fib(n) {
    var current = 0,
        next = 1;

    for i in 1..n {
        yield current;
        current += next;
        current <= next;
    }
}
```

```
config const n = 10;

for (i,f) in zip(0..#n, fib(n)) do
    writeln("fib #", i, " is ", f);
```

```
fib #0 is 0
fib #1 is 1
fib #2 is 1
fib #3 is 2
fib #4 is 3
fib #5 is 5
fib #6 is 8
...
```



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# Base Language Features, by example



tuples

```
iter fib(n) {
    var current = 0,
        next = 1;

    for i in 1..n {
        yield current;
        current += next;
        current <= next;
    }
}
```

```
config const n = 10;

for (i,f) in zip(0..#n, fib(n)) do
    writeln("fib #", i, " is ", f);
```

```
fib #0 is 0
fib #1 is 1
fib #2 is 1
fib #3 is 2
fib #4 is 3
fib #5 is 5
fib #6 is 8
...
```



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# Base Language Features, by example



```
iter fib(n) {
    var current = 0,
        next = 1;

    for i in 1..n {
        yield current;
        current += next;
        current <= next;
    }
}
```

```
config const n = 10;

for (i,f) in zip(0..#n, fib(n)) do
    writeln("fib #", i, " is ", f);
```

```
fib #0 is 0
fib #1 is 1
fib #2 is 1
fib #3 is 2
fib #4 is 3
fib #5 is 5
fib #6 is 8
```

...

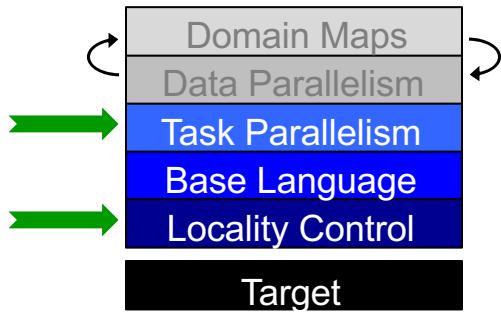


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# Task Parallelism and Locality Control



# Task Parallelism and Locality, by example



taskParallel.chpl

```
coforall loc in Locales do
    on loc {
        const numTasks = here.maxTaskPar;
        coforall tid in 1..numTasks do
            writef("Hello from task %n of %n "+
                "running on %s\n",
                tid, numTasks, here.name);
    }
```

```
prompt> chpl taskParallel.chpl -o taskParallel
prompt> ./taskParallel --numLocales=2
Hello from task 1 of 2 running on n1033
Hello from task 2 of 2 running on n1032
Hello from task 2 of 2 running on n1033
Hello from task 1 of 2 running on n1032
```

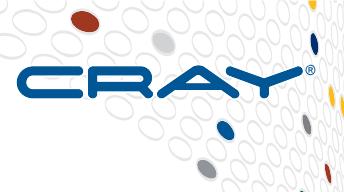


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# Task Parallelism and Locality, by example



Abstraction of  
System Resources

taskParallel.chpl

```
coforall loc in Locales do
    on loc {
        const numTasks = here.maxTaskPar;
        coforall tid in 1..numTasks do
            writef("Hello from task %n of %n "+
                "running on %s\n",
                tid, numTasks, here.name);
    }
```

```
prompt> chpl taskParallel.chpl -o taskParallel
prompt> ./taskParallel --numLocales=2
Hello from task 1 of 2 running on n1033
Hello from task 2 of 2 running on n1032
Hello from task 2 of 2 running on n1033
Hello from task 1 of 2 running on n1032
```



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# Task Parallelism and Locality, by example



High-Level  
Task Parallelism

taskParallel.chpl

```
coforall loc in Locales do
    on loc {
        const numTasks = here.maxTaskPar;
        coforall tid in 1..numTasks do
            writef("Hello from task %n of %n "+
                "running on %s\n",
                tid, numTasks, here.name);
    }
```

```
prompt> chpl taskParallel.chpl -o taskParallel
prompt> ./taskParallel --numLocales=2
Hello from task 1 of 2 running on n1033
Hello from task 2 of 2 running on n1032
Hello from task 2 of 2 running on n1033
Hello from task 1 of 2 running on n1032
```



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# Task Parallelism and Locality, by example



Control of Locality/Affinity

taskParallel.chpl

```
coforall loc in Locales do
    on loc {
        const numTasks = here.maxTaskPar;
        coforall tid in 1..numTasks do
            writef("Hello from task %n of %n "+
                "running on %s\n",
                tid, numTasks, here.name);
    }
```

```
prompt> chpl taskParallel.chpl -o taskParallel
prompt> ./taskParallel --numLocales=2
Hello from task 1 of 2 running on n1033
Hello from task 2 of 2 running on n1032
Hello from task 2 of 2 running on n1033
Hello from task 1 of 2 running on n1032
```



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# Task Parallelism and Locality, by example



Abstraction of  
System Resources

taskParallel.chpl

```
coforall loc in Locales do
    on loc {
        const numTasks = here.maxTaskPar;
        coforall tid in 1..numTasks do
            writef("Hello from task %n of %n "+
                "running on %s\n",
                tid, numTasks, here.name);
    }
```

```
prompt> chpl taskParallel.chpl -o taskParallel
prompt> ./taskParallel --numLocales=2
Hello from task 1 of 2 running on n1033
Hello from task 2 of 2 running on n1032
Hello from task 2 of 2 running on n1033
Hello from task 1 of 2 running on n1032
```



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# Task Parallelism and Locality, by example



High-Level  
Task Parallelism

taskParallel.chpl

```
coforall loc in Locales do
    on loc {
        const numTasks = here.maxTaskPar;
        coforall tid in 1..numTasks do
            writef("Hello from task %n of %n "+
                "running on %s\n",
                tid, numTasks, here.name);
    }
```

```
prompt> chpl taskParallel.chpl -o taskParallel
prompt> ./taskParallel --numLocales=2
Hello from task 1 of 2 running on n1033
Hello from task 2 of 2 running on n1032
Hello from task 2 of 2 running on n1033
Hello from task 1 of 2 running on n1032
```



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# Task Parallelism and Locality, by example



taskParallel.chpl

```
coforall loc in Locales do
    on loc {
        const numTasks = here.maxTaskPar;
        coforall tid in 1..numTasks do
            writef("Hello from task %n of %n "+
                "running on %s\n",
                tid, numTasks, here.name);
    }
```

Not seen here:  
Data-centric task coordination  
via atomic and full/empty vars

```
prompt> chpl taskParallel.chpl -o taskParallel
prompt> ./taskParallel --numLocales=2
Hello from task 1 of 2 running on n1033
Hello from task 2 of 2 running on n1032
Hello from task 2 of 2 running on n1033
Hello from task 1 of 2 running on n1032
```



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# Task Parallelism and Locality, by example



taskParallel.chpl

```
coforall loc in Locales do
    on loc {
        const numTasks = here.maxTaskPar;
        coforall tid in 1..numTasks do
            writef("Hello from task %n of %n "+
                "running on %s\n",
                tid, numTasks, here.name);
    }
```

```
prompt> chpl taskParallel.chpl -o taskParallel
prompt> ./taskParallel --numLocales=2
Hello from task 1 of 2 running on n1033
Hello from task 2 of 2 running on n1032
Hello from task 2 of 2 running on n1033
Hello from task 1 of 2 running on n1032
```



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# Parallelism and Locality: Distinct in Chapel



- This is a **parallel**, but local program:

```
coforall i in 1..msgs do  
    writeln("Hello from task ", i);
```

- This is a **distributed**, but serial program:

```
writeln("Hello from locale 0!");  
on Locales[1] do writeln("Hello from locale 1!");  
on Locales[2] do writeln("Hello from locale 2!");
```

- This is a **distributed parallel** program:

```
coforall i in 1..msgs do  
    on Locales[i%numLocales] do  
        writeln("Hello from task ", i,  
               " running on locale ", here.id);
```



# Partitioned Global Address Space (PGAS) Languages

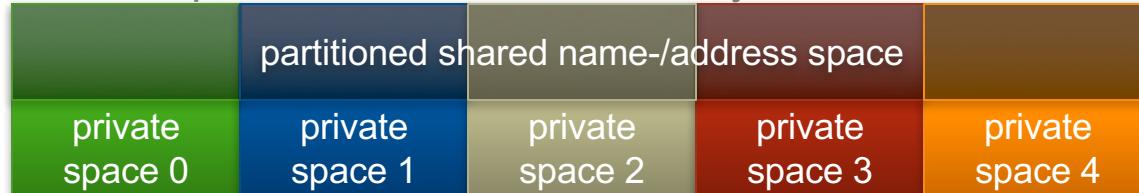
(Or more accurately: partitioned global namespace languages)

- **abstract concept:**

- support a shared namespace on distributed memory
  - permit parallel tasks to access remote variables by naming them
- establish a strong sense of ownership
  - every variable has a well-defined location
  - local variables are cheaper to access than remote ones

- **traditional PGAS languages have been SPMD in nature**

- best-known examples: Fortran 2008's co-arrays, Unified Parallel C (UPC)



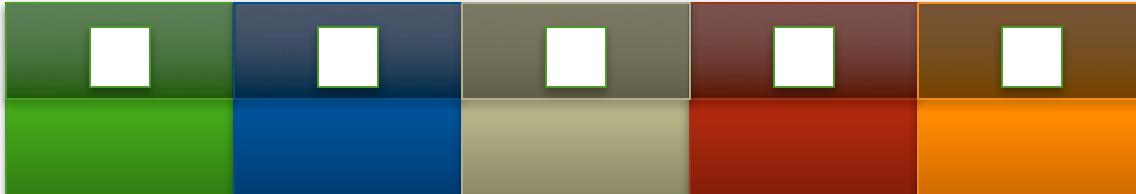
# SPMD PGAS Languages

(using a pseudo-language, not Chapel)



```
shared int i(*) ;           // declare a shared variable i
```

$i =$



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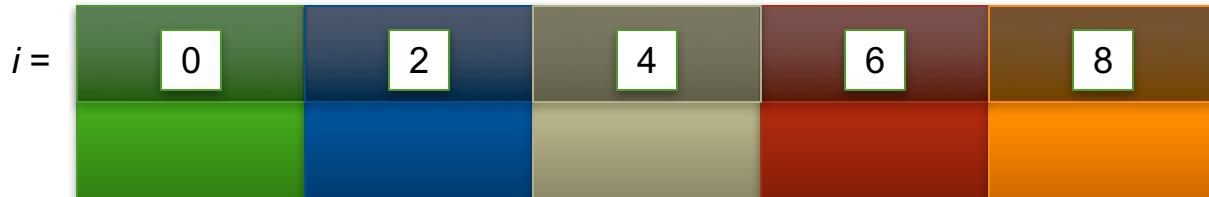
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# SPMD PGAS Languages

(using a pseudo-language, not Chapel)



```
shared int i(*) ;           // declare a shared variable i
function main() {
    i = 2*this_image();    // each image initializes its copy
```



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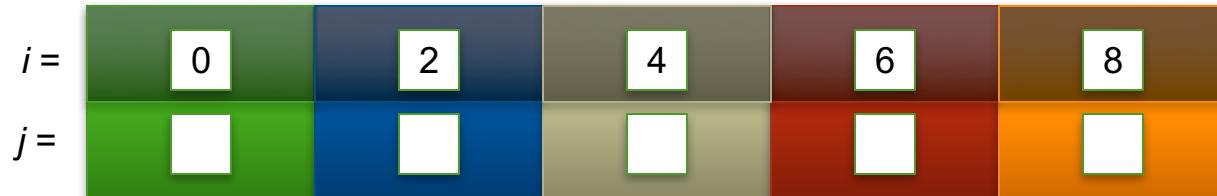
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# SPMD PGAS Languages

(using a pseudo-language, not Chapel)



```
shared int i(*) ;           // declare a shared variable i  
function main() {  
    i = 2*this_image();    // each image initializes its copy  
  
    private int j;          // declare a private variable j
```



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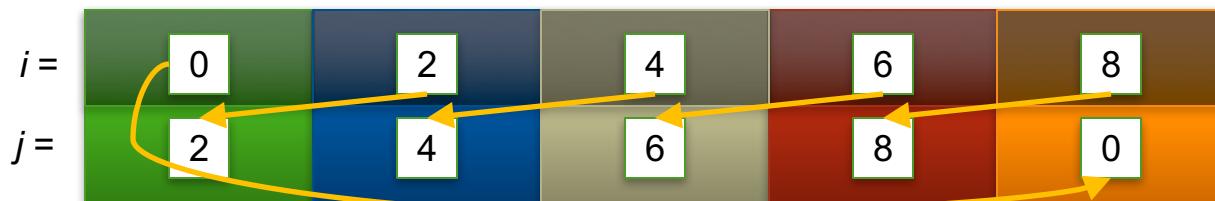
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# SPMD PGAS Languages

(using a pseudo-language, not Chapel)



```
shared int i(*) ;           // declare a shared variable i
function main() {
    i = 2*this_image() ;   // each image initializes its copy
    barrier();
    private int j;          // declare a private variable j
    j = i( (this_image()+1) % num_images() ) ;
    // ^ access our neighbor's copy of i; compiler and runtime implement the communication
    // Q: How did we know our neighbor had an i?
    // A: Because it's SPMD – we're all running the same program so if we have an i, so do they.
```



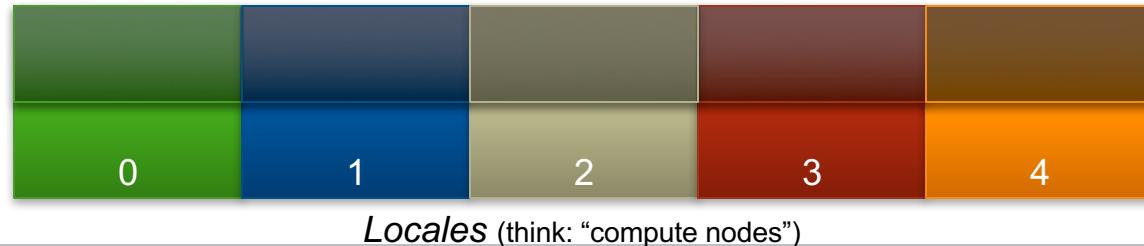
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# Chapel and PGAS

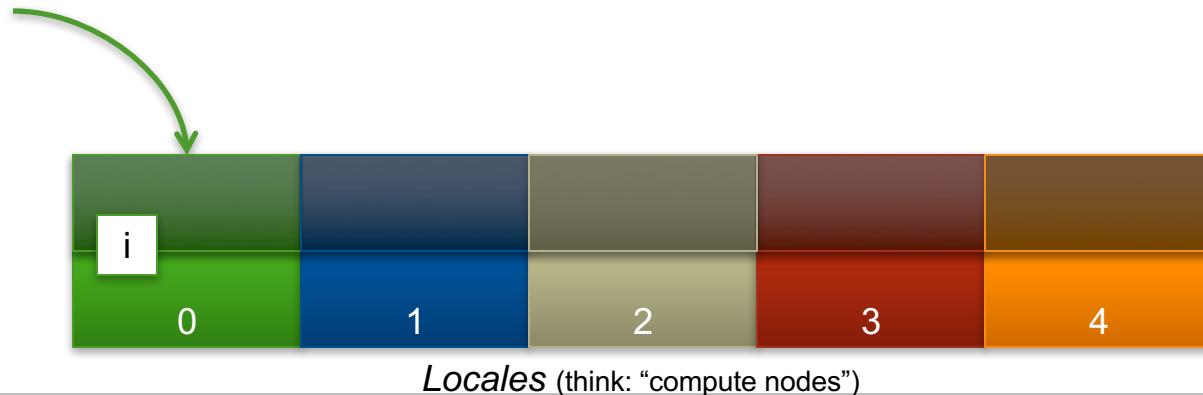
- Chapel is PGAS, but unlike most, it's not inherently SPMD
  - never think about “the other copies of the program”
  - “global name/address space” comes from lexical scoping
    - as in traditional languages, each declaration yields one variable
    - variables are stored on the locale where the task declaring it is executing



# Chapel: Scoping and Locality



```
var i: int;
```



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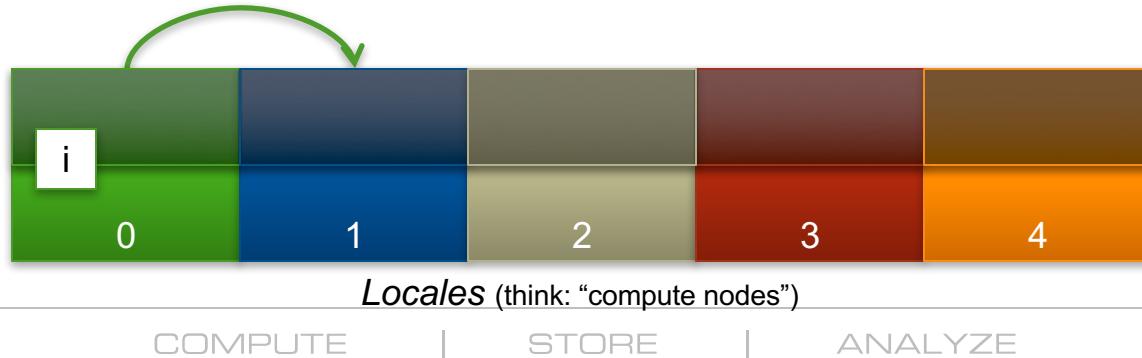
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# Chapel: Scoping and Locality



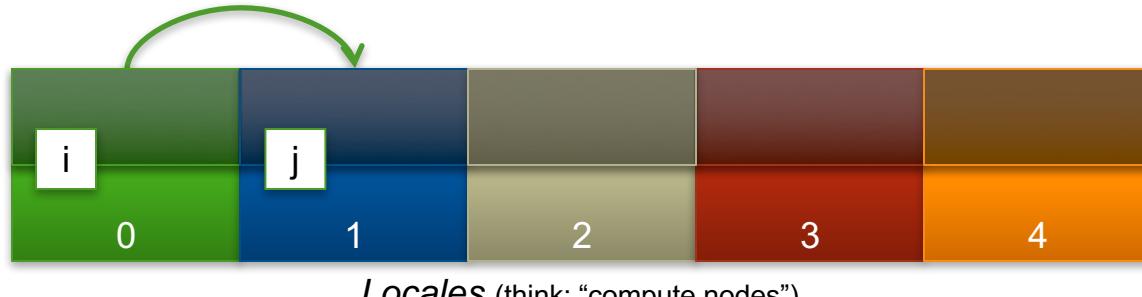
```
var i: int;  
on Locales[1] {
```



# Chapel: Scoping and Locality



```
var i: int;  
on Locales[1] {  
    var j: int;
```



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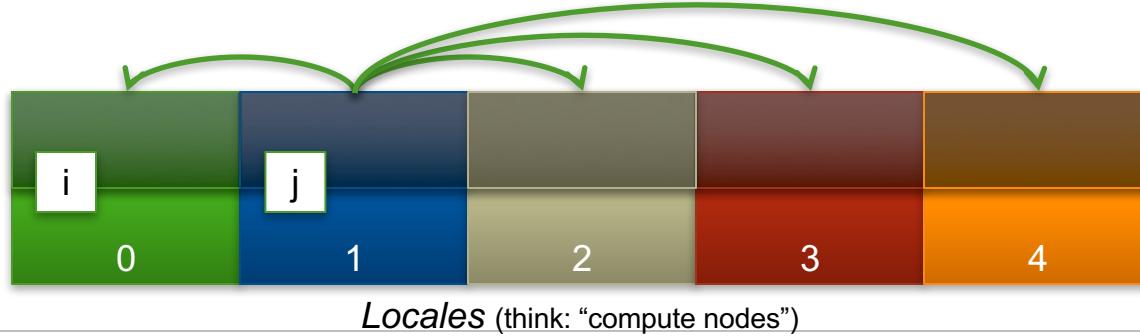
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# Chapel: Scoping and Locality



```
var i: int;  
on Locales[1] {  
    var j: int;  
    coforall loc in Locales {  
        on loc {
```



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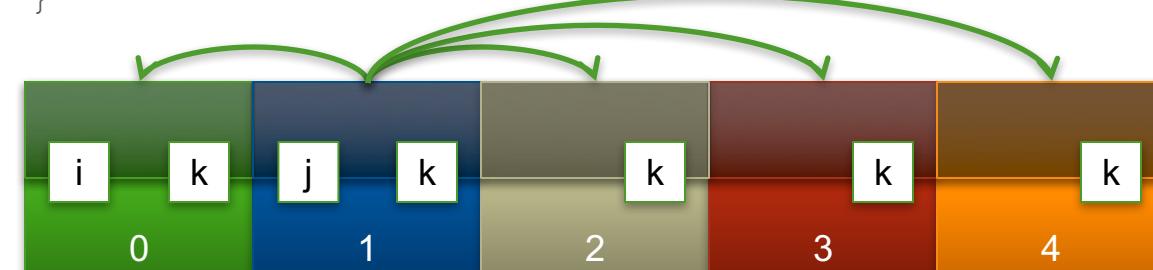
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# Chapel: Scoping and Locality



```
var i: int;  
on Locales[1] {  
    var j: int;  
    coforall loc in Locales {  
        on loc {  
            var k: int;  
            ...  
        }  
    }  
}
```



*Locales* (think: “compute nodes”)

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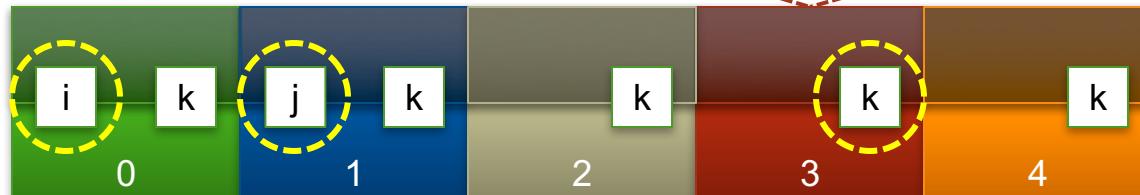
# Chapel: Scoping and Locality



```
var i: int;  
on Locales[1] {  
    var j: int;  
    coforall loc in Locales {  
        on loc {  
            var k: int;  
            k = 2*i + j;  
        }  
    }  
}  
}  
}
```

OK to access  $i$ ,  $j$ , and  $k$   
wherever they live

$k = 2*i + j;$



*Locales* (think: “compute nodes”)

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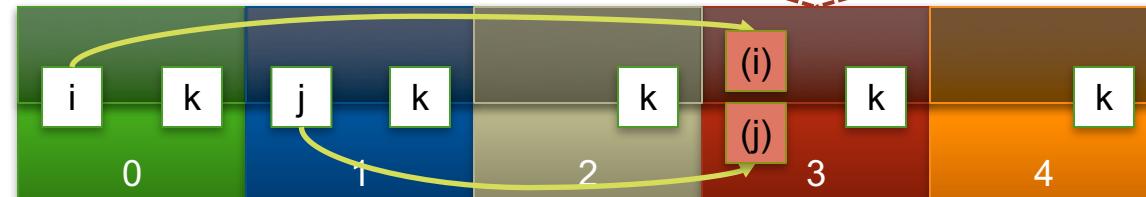


# Chapel: Scoping and Locality



```
var i: int;  
on Locales[1] {  
    var j: int;  
    coforall loc in Locales {  
        on loc {  
            var k: int;  
            k = 2*i + j;  
        }  
    }  
}
```

here, *i* and *j* are remote, so  
the compiler + runtime will  
transfer their values



*Locales* (think: “compute nodes”)

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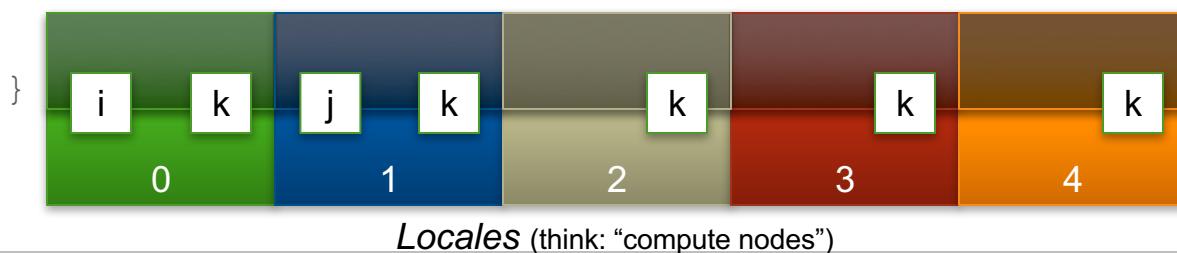
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# Chapel: Locality queries



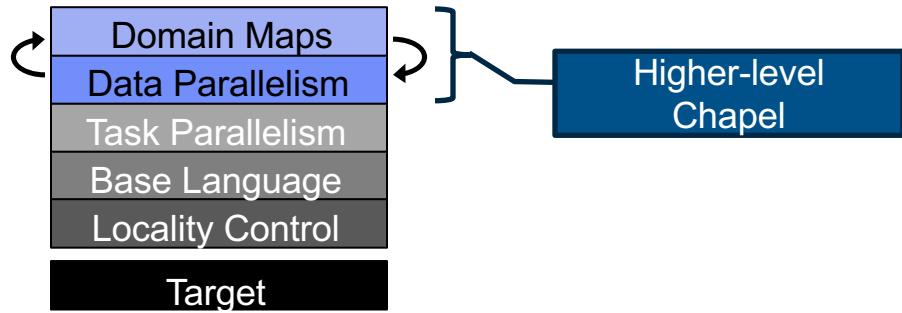
```
var i: int;  
on Locales[1] {  
    var j: int;  
    coforall loc in Locales {  
        on loc {  
            var k: int;  
  
                ...here...          // query the locale on which this task is running  
                ...j.locale...      // query the locale on which j is stored  
        }  
    }  
}
```



# Higher-Level Features



*Chapel language concepts*



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# Data Parallelism, by example



dataParallel.chpl

```
config const n = 1000;
var D = {1..n, 1..n};

var A: [D] real;
forall (i,j) in D do
    A[i,j] = i + (j - 0.5)/n;
writeln(A);
```

```
prompt> chpl dataParallel.chpl -o dataParallel
prompt> ./dataParallel --n=5
1.1 1.3 1.5 1.7 1.9
2.1 2.3 2.5 2.7 2.9
3.1 3.3 3.5 3.7 3.9
4.1 4.3 4.5 4.7 4.9
5.1 5.3 5.5 5.7 5.9
```



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# Data Parallelism, by example



Domains (Index Sets)

dataParallel.chpl

```
config const n = 1000;
var D = {1..n, 1..n};

var A: [D] real;
forall (i,j) in D do
    A[i,j] = i + (j - 0.5)/n;
writeln(A);
```

```
prompt> chpl dataParallel.chpl -o dataParallel
prompt> ./dataParallel --n=5
1.1 1.3 1.5 1.7 1.9
2.1 2.3 2.5 2.7 2.9
3.1 3.3 3.5 3.7 3.9
4.1 4.3 4.5 4.7 4.9
5.1 5.3 5.5 5.7 5.9
```

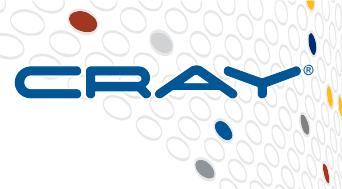


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# Data Parallelism, by example



Arrays

dataParallel.chpl

```
config const n = 1000;
var D = {1..n, 1..n};

var A: [D] real;
forall (i,j) in D do
    A[i,j] = i + (j - 0.5)/n;
writeln(A);
```

```
prompt> chpl dataParallel.chpl -o dataParallel
prompt> ./dataParallel --n=5
1.1 1.3 1.5 1.7 1.9
2.1 2.3 2.5 2.7 2.9
3.1 3.3 3.5 3.7 3.9
4.1 4.3 4.5 4.7 4.9
5.1 5.3 5.5 5.7 5.9
```



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# Data Parallelism, by example



## Data-Parallel Forall Loops

dataParallel.chpl

```
config const n = 1000;
var D = {1..n, 1..n};

var A: [D] real;
forall (i,j) in D do
    A[i,j] = i + (j - 0.5)/n;
writeln(A);
```

```
prompt> chpl dataParallel.chpl -o dataParallel
prompt> ./dataParallel --n=5
1.1 1.3 1.5 1.7 1.9
2.1 2.3 2.5 2.7 2.9
3.1 3.3 3.5 3.7 3.9
4.1 4.3 4.5 4.7 4.9
5.1 5.3 5.5 5.7 5.9
```



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# Distributed Data Parallelism, by example



Domain Maps  
(Map Data Parallelism to the System)

dataParallel.chpl

```
use CyclicDist;
config const n = 1000;
var D = {1..n, 1..n}
    dmapped Cyclic(startIdx = (1,1));
var A: [D] real;
forall (i,j) in D do
    A[i,j] = i + (j - 0.5)/n;
writeln(A);
```

```
prompt> chpl dataParallel.chpl -o dataParallel
prompt> ./dataParallel --n=5 --numLocales=4
1.1 1.3 1.5 1.7 1.9
2.1 2.3 2.5 2.7 2.9
3.1 3.3 3.5 3.7 3.9
4.1 4.3 4.5 4.7 4.9
5.1 5.3 5.5 5.7 5.9
```



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# Distributed Data Parallelism, by example



dataParallel.chpl

```
use CyclicDist;
config const n = 1000;
var D = {1..n, 1..n}
        dmapped Cyclic(startIdx = (1,1));
var A: [D] real;
forall (i,j) in D do
    A[i,j] = i + (j - 0.5)/n;
writeln(A);
```

```
prompt> chpl dataParallel.chpl -o dataParallel
prompt> ./dataParallel --n=5 --numLocales=4
1.1 1.3 1.5 1.7 1.9
2.1 2.3 2.5 2.7 2.9
3.1 3.3 3.5 3.7 3.9
4.1 4.3 4.5 4.7 4.9
5.1 5.3 5.5 5.7 5.9
```



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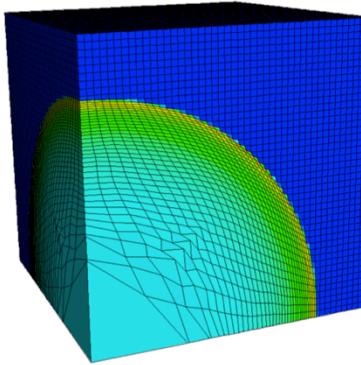
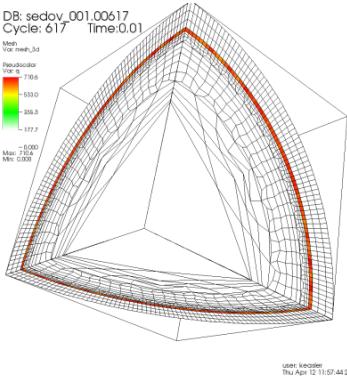
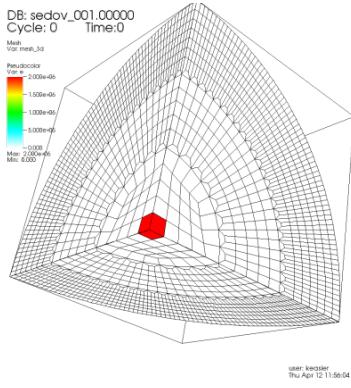
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# LULESH: a DOE Proxy Application



**Goal:** Solve one octant of the spherical Sedov problem (blast wave) using Lagrangian hydrodynamics for a single material



pictures courtesy of Rob Neely, Bert Still, Jeff Keasler, LLNL

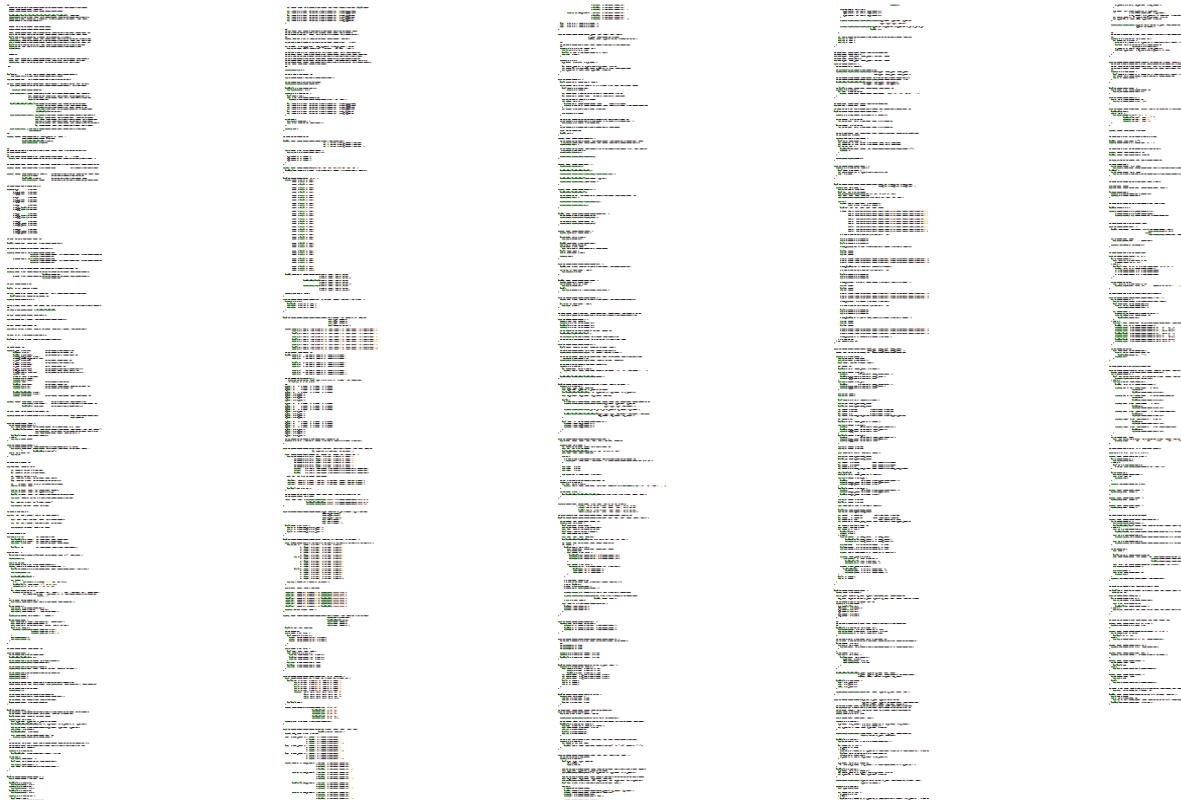


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# LULESH in Chapel



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# LULESH in Chapel



1288 lines of source code  
plus 266 lines of comments  
487 blank lines

(the corresponding C+MPI+OpenMP version is nearly 4x bigger)

This can be found in the Chapel release under examples/benchmarks/lulesh/\*.chpl



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# LULESH in Chapel



This is all of the representation-dependent code.  
It specifies:

- data structure choices
  - structured vs. unstructured mesh
  - local vs. distributed data
  - sparse vs. dense materials arrays
- a few supporting iterators

Small number of changes enabled by domain maps



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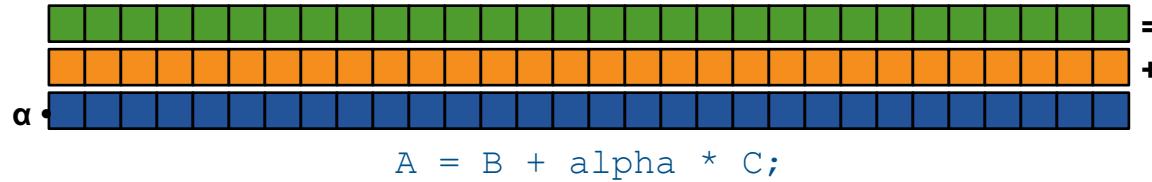
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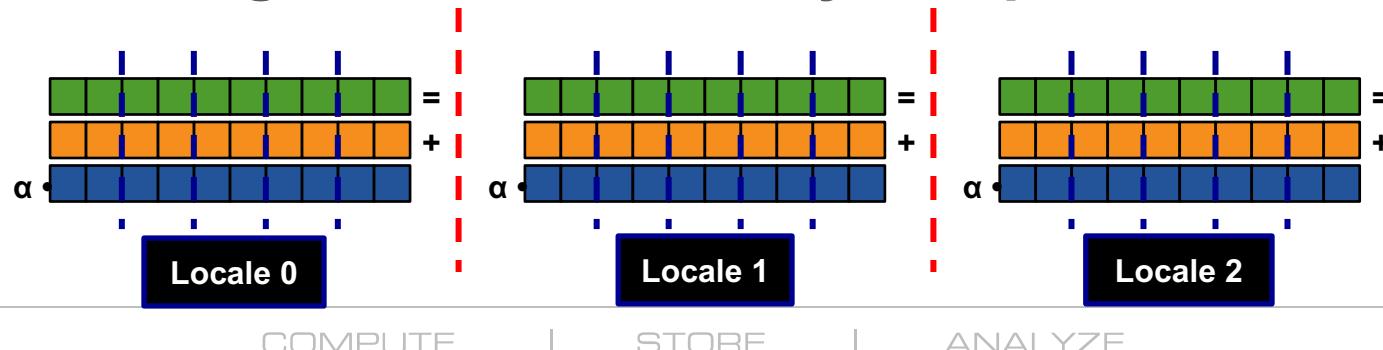
# Domain Maps



Domain maps are “recipes” that instruct the compiler how to map the global view of a computation...



...to the target locales' memory and processors:



# Chapel's Domain Map Philosophy

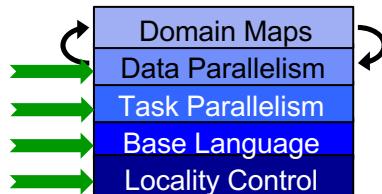


## 1. Chapel provides a library of standard domain maps

- to support common array implementations effortlessly

## 2. Expert users can write their own domain maps in Chapel

- to cope with any shortcomings in our standard library



## 3. Chapel's standard domain maps are written using the end-user framework

- to avoid a performance cliff between “built-in” and user-defined cases
- in fact every Chapel array is implemented using this framework

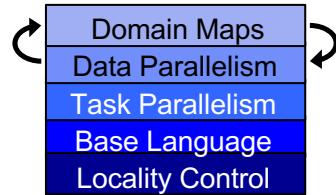


# Chapel's Multiresolution Philosophy



## ***Multiresolution Design:*** Support multiple tiers of features

- higher levels for programmability, productivity
- lower levels for greater degrees of control



- build the higher-level concepts in terms of the lower
- permit users to intermix layers arbitrarily



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# Two Other Multiresolution Features



1) **parallel iterators:** User-specified forall-loop implementations

- how many tasks to use and where they run
- how iterations are divided between the tasks
- how to zipper with other parallel iterators

2) **locale models:** User-specified locale types for new node architectures

- how do I manage memory, create tasks, communicate, ...

Like domain maps, these are...

...written in Chapel by expert users using lower-level features

...available to the end-user via higher-level abstractions



# This Talk's Takeaways



## If you design a parallel programming language...

...create attractive ways of expressing parallelism and locality

- these are key concerns—shouldn't be yet another library call / pragma

...support first-class index sets

- very expressive concept for users (not seen here, today)
- enable the implementation to reason about alignment (ditto)

...support multiresolution features

- give users the ability to develop their own parallel policies



*Puttin' me down for thinking of someone new  
Always the same  
Playin' your game  
Drive me insane...*

## Computer Language Benchmarks Game Results

**Your Time is Gonna Come**  
Led Zeppelin  
Led Zeppelin



# Computer Language Benchmarks Game (CLBG)



## The Computer Language Benchmarks Game

64-bit quad core data set

Will your toy benchmark program be faster if you write it in a different programming language? It depends how you write it!

### Which programs are fast?

Which are succinct? Which are efficient?

<u>Ada</u>	<u>C</u>	<u>Chapel</u>	<u>Clojure</u>	<u>C#</u>	<u>C++</u>
<u>Dart</u>	<u>Erlang</u>	<u>F#</u>	<u>Fortran</u>	<u>Go</u>	<u>Hack</u>
<u>Haskell</u>	<u>Java</u>	<u>JavaScript</u>	<u>Lisp</u>	<u>Lua</u>	
<u>OCaml</u>	<u>Pascal</u>	<u>Perl</u>	<u>PHP</u>	<u>Python</u>	
<u>Racket</u>	<u>Ruby</u>	<u>JRuby</u>	<u>Rust</u>	<u>Scala</u>	
<u>Smalltalk</u>	<u>Swift</u>	<u>TypeScript</u>			

## Website supporting cross-language comparisons

- 13 toy benchmark programs
  - exercise key computational idioms
  - specific approach prescribed

## Take results with a grain of salt

- your mileage may vary

That said, it is one of the only such games in town...



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# Computer Language Benchmarks Game (CLBG)

## The Computer Language Benchmarks Game

64-bit quad core data set

Will your toy benchmark program be faster if you write it in a different programming language? It depends how you write it!

### Which programs are fast?

Which are succinct? Which are efficient?

<u>Ada</u>	<u>C</u>	<u>Chapel</u>	<u>Clojure</u>	<u>C#</u>	<u>C++</u>
<u>Dart</u>	<u>Erlang</u>	<u>F#</u>	<u>Fortran</u>	<u>Go</u>	<u>Hack</u>
<u>Haskell</u>	<u>Java</u>	<u>JavaScript</u>	<u>Lisp</u>	<u>Lua</u>	
<u>OCaml</u>	<u>Pascal</u>	<u>Perl</u>	<u>PHP</u>	<u>Python</u>	
<u>Racket</u>	<u>Ruby</u>	<u>JRuby</u>	<u>Rust</u>	<u>Scala</u>	
<u>Smalltalk</u>	<u>Swift</u>	<u>TypeScript</u>			

## Chapel's approach to the CLBG:

- striving for elegance over heroism
  - ideally: “Want to learn how program xyz works? Read the Chapel version.”

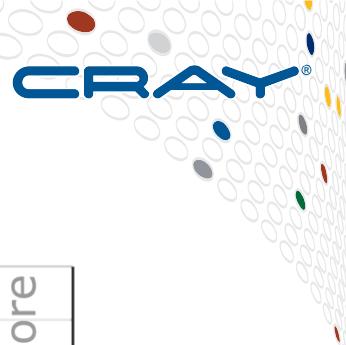


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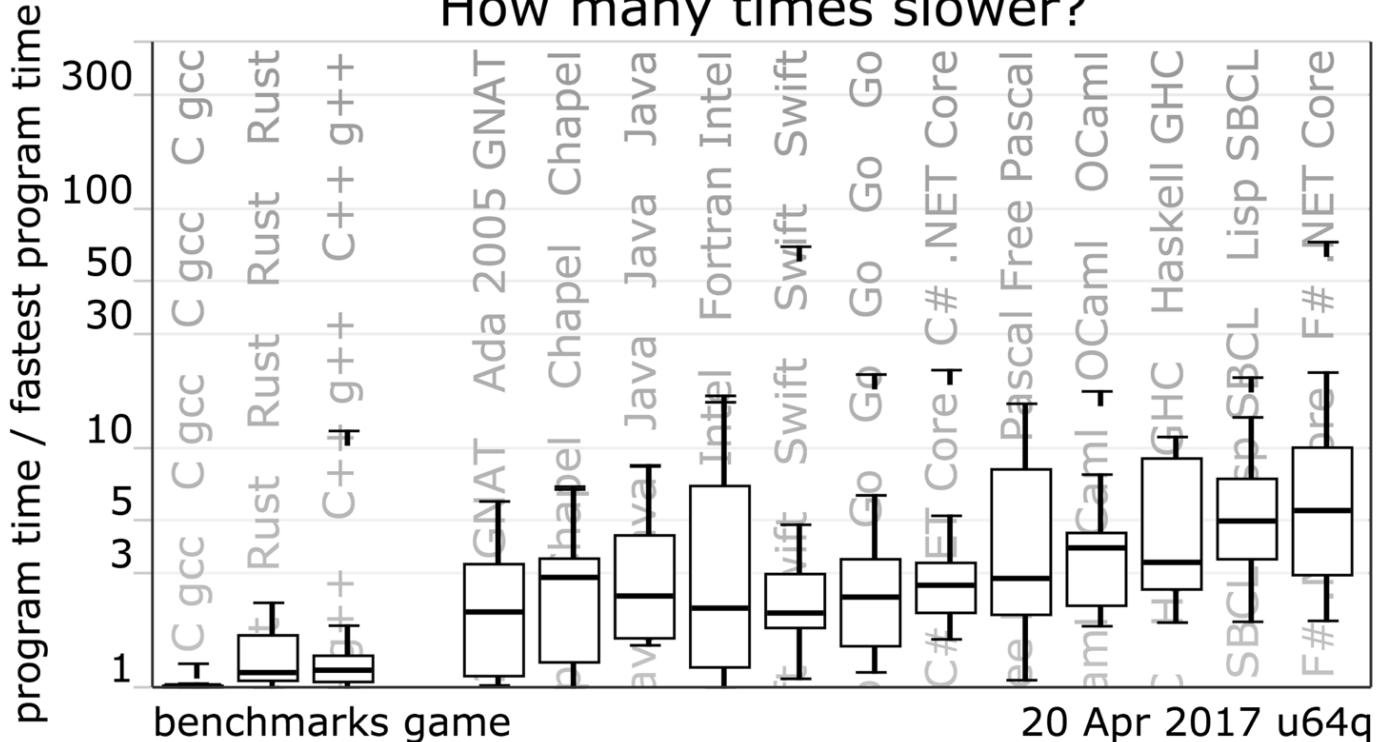
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# CLBG: Relative Performance Summary



How many times slower?

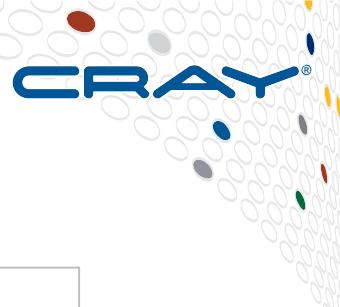


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# CLBG: Website



Can sort results by execution time, code size, memory or CPU use:

The Computer Language Benchmarks Game								
chameneos-redux								
description								
program source code, command-line and measurements								
x	source	secs	mem	gz	cpu	cpu load		
1.0	<a href="#">C gcc #5</a>	<b>0.60</b>	820	2863	2.37	100%	100%	98% 100%
1.2	<a href="#">C++ g++ #5</a>	<b>0.70</b>	3,356	1994	2.65	100%	100%	91% 92%
1.7	<a href="#">Lisp SBCL #3</a>	<b>1.01</b>	55,604	2907	3.93	97%	96%	99% 99%
2.3	<a href="#">Chapel #2</a>	<b>1.39</b>	76,564	1210	5.43	99%	99%	98% 99%
3.3	<a href="#">Rust #2</a>	<b>2.01</b>	56,936	2882	7.81	97%	98%	98% 98%
5.6	<a href="#">C++ g++ #2</a>	3.40	1,880	2016	11.88	100%	51%	100% 100%
6.8	<a href="#">Chapel</a>	4.09	66,584	1199	16.25	100%	100%	100% 100%
8.0	<a href="#">Java #4</a>	<b>4.82</b>	37,132	1607	16.73	98%	98%	54% 99%
8.5	<a href="#">Haskell GHC</a>	<b>5.15</b>	8,596	989	9.26	79%	100%	2% 2%
10	<a href="#">Java</a>	6.13	53,760	1770	8.78	42%	45%	41% 16%
10	<a href="#">Haskell GHC #4</a>	6.34	6,908	989	12.67	99%	100%	2% 1%
11	<a href="#">C# .NET Core</a>	<b>6.59</b>	86,076	1400	22.96	99%	82%	78% 91%
11	<a href="#">Go</a>	<b>6.90</b>	832	1167	24.19	100%	96%	56% 100%
13	<a href="#">Go #2</a>	7.59	1,384	1408	27.65	91%	99%	99% 78%
13	<a href="#">Java #3</a>	7.94	53,232	1267	26.86	54%	96%	98% 94%

The Computer Language Benchmarks Game								
chameneos-redux								
description								
program source code, command-line and measurements								
x	source	secs	mem	gz	cpu	cpu load		
1.0	<a href="#">Erlang</a>	58.90	28,668	<b>734</b>	131.19	62%	60%	51% 53%
1.0	<a href="#">Erlang HIPE</a>	59.39	25,784	<b>734</b>	131.58	60%	56%	56% 54%
1.1	<a href="#">Perl #4</a>	5 min	14,084	<b>785</b>	7 min	40%	40%	29% 28%
1.1	<a href="#">Racket</a>	5 min	132,120	<b>791</b>	5 min	1%	0%	0% 100%
1.1	<a href="#">Racket #2</a>	175.88	116,488	842	175.78	100%	1%	1% 0%
1.2	<a href="#">Python 3 #2</a>	236.84	7,908	<b>866</b>	5 min	24%	48%	27% 45%
1.3	<a href="#">Ruby</a>	90.52	9,396	<b>920</b>	137.53	35%	35%	35% 34%
1.3	<a href="#">Ruby JRuby</a>	48.78	628,968	<b>928</b>	112.15	65%	60%	49% 58%
1.3	<a href="#">Go #5</a>	11.05	832	<b>957</b>	32.48	75%	74%	75% 73%
1.3	<a href="#">Haskell GHC #4</a>	6.34	6,908	<b>989</b>	12.67	99%	100%	2% 1%
1.3	<a href="#">Haskell GHC</a>	5.15	8,596	989	9.26	79%	100%	2% 2%
1.6	<a href="#">OCaml #3</a>	32%	38%	37%	39%			
1.6	<a href="#">Go</a>	100%	96%	56%	100%			
1.6	<a href="#">Chapel</a>	0%	100%	100%	100%			
1.6	<a href="#">Chapel #2</a>	99%	99%	98%	99%			

gz == code size metric  
strip comments and extra whitespace, then gzip

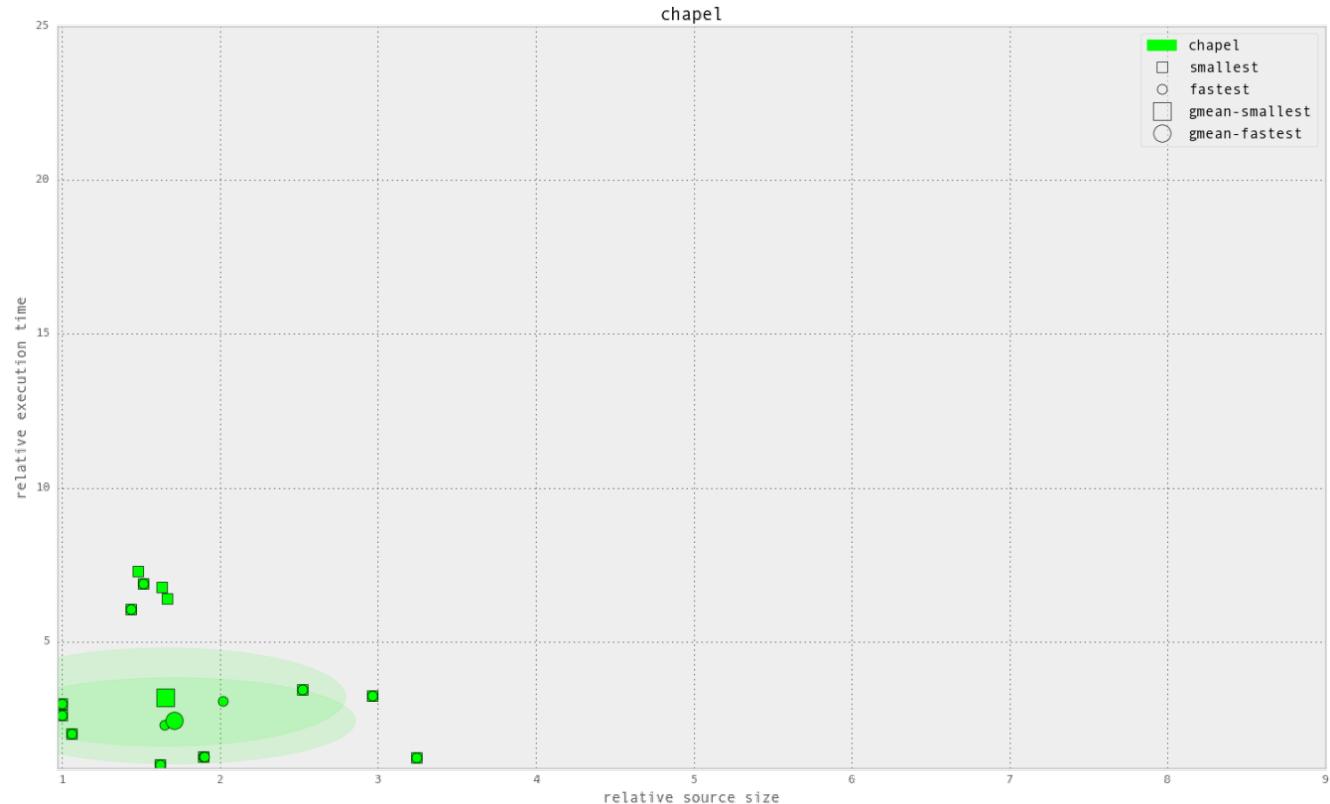


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# CLBG: Chapel entries

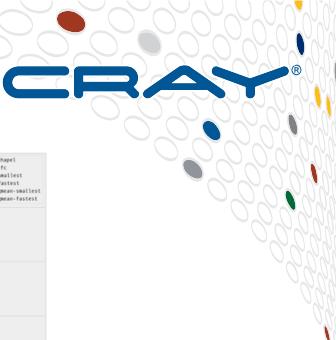


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# CLBG: Chapel vs. 9 key languages

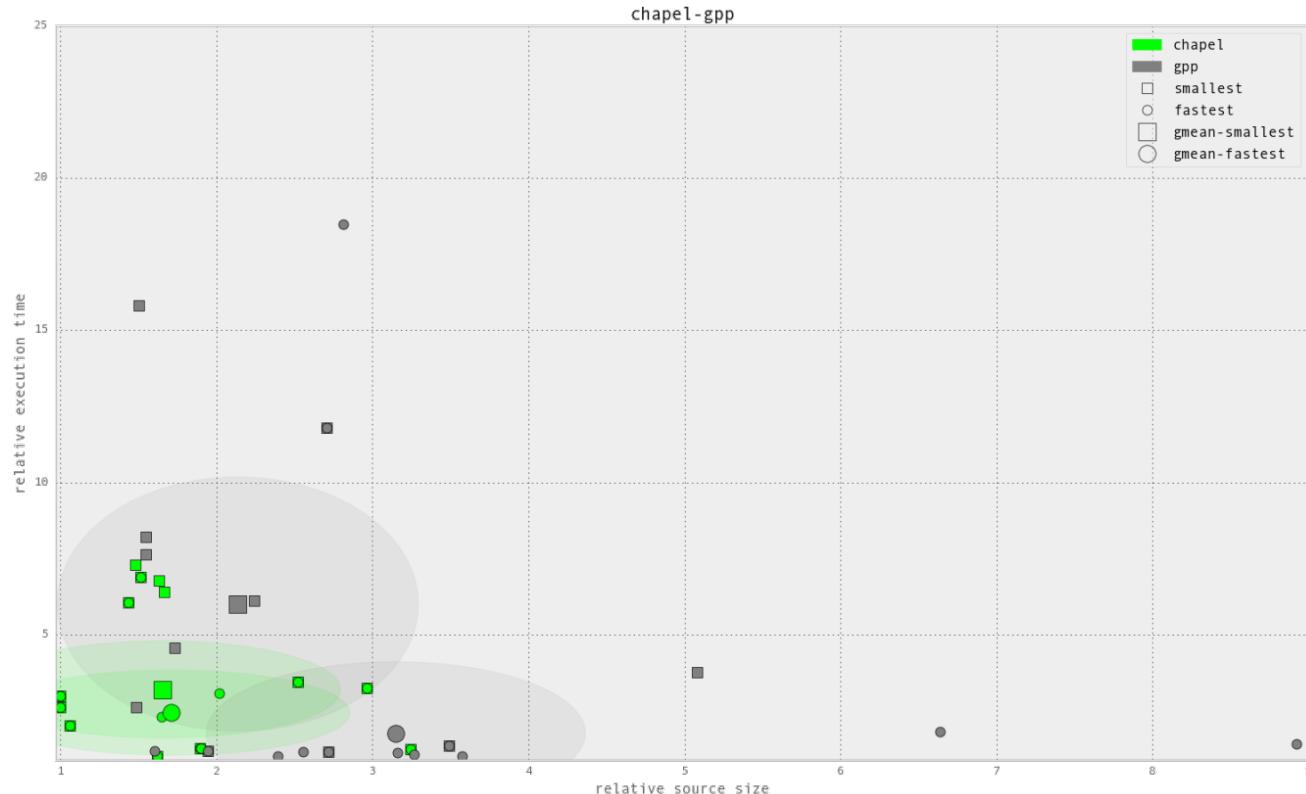


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# Chapel vs. C++

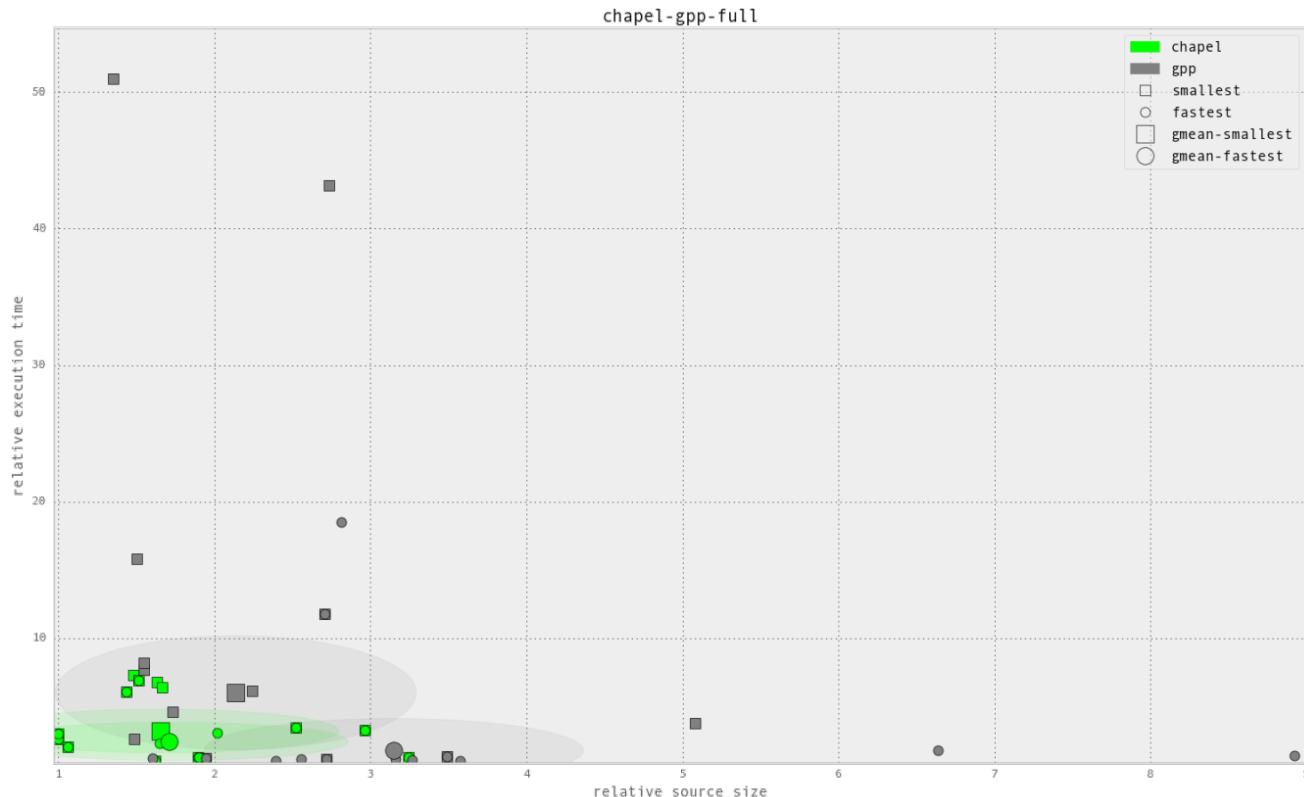


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# Chapel vs. C++ (zoomed out)

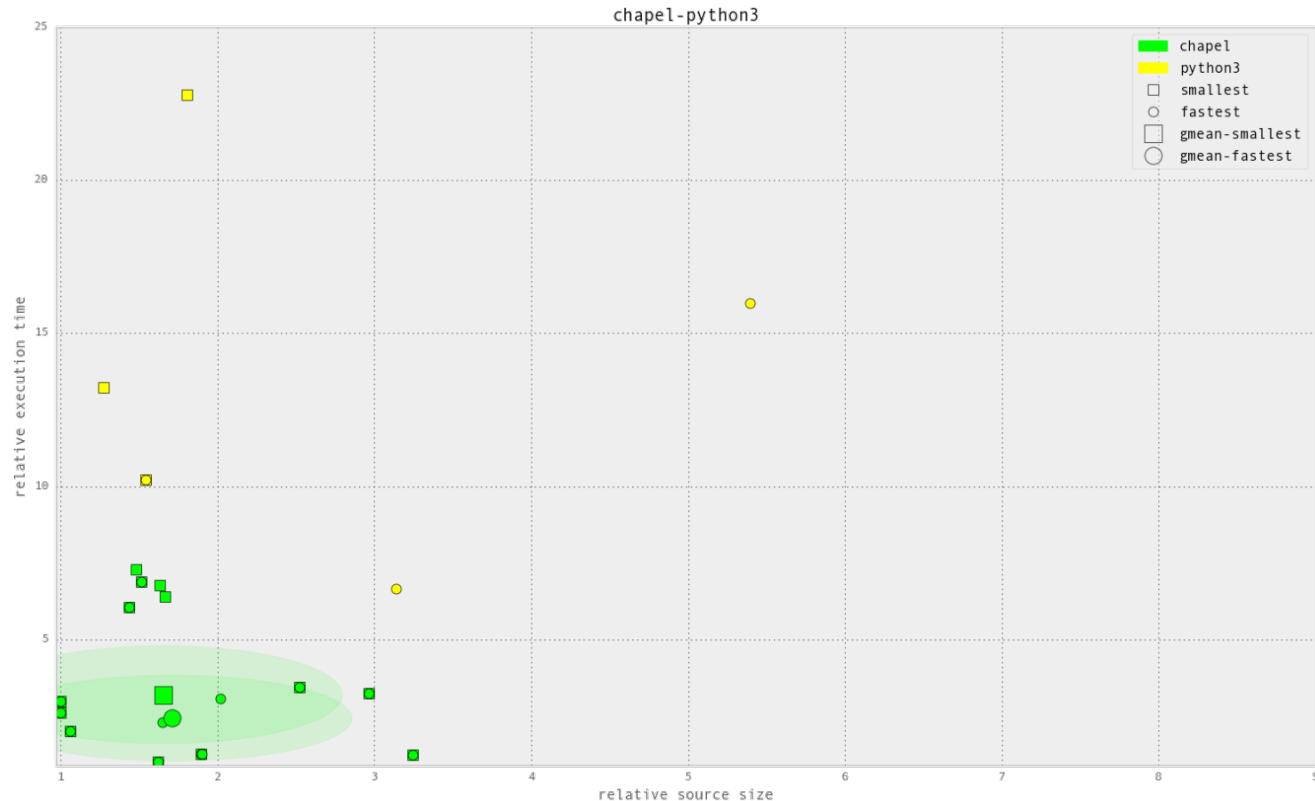


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# Chapel vs. Python

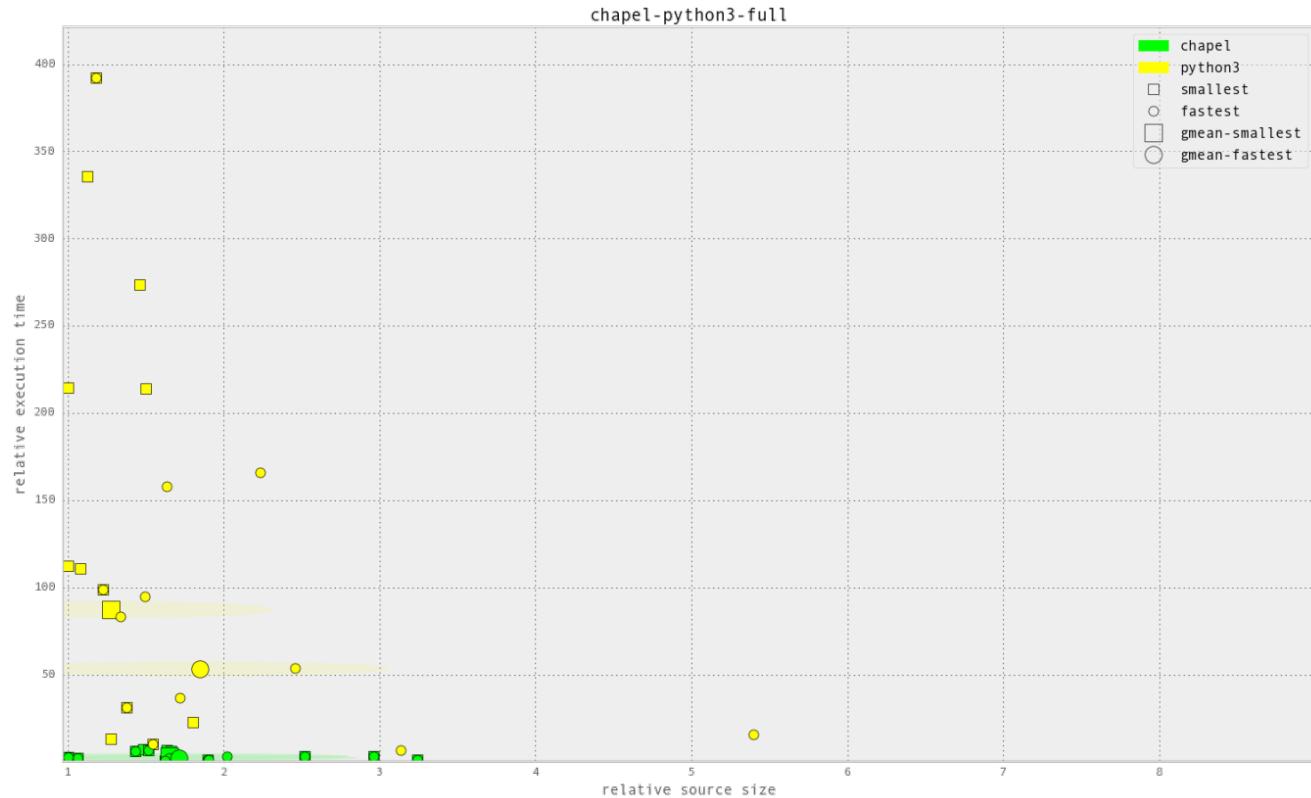


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# Chapel vs. Python (zoomed out)

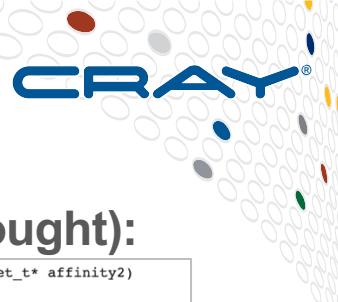


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# CLBG: Qualitative Comparisons



Can also browse program source code (but this requires actual thought):

```
proc main() {
    printColorEquations();

    const group1 = {i in 1..popSize1} new Chameneos(i, ((i-1)%3):Color);
    const group2 = {i in 1..popSize2} new Chameneos(i, colors10[i]);

    cobegin {
        holdMeetings(group1, n);
        holdMeetings(group2, n);
    }

    print(group1);
    print(group2);

    for c in group1 do delete c;
    for c in group2 do delete c;
}

// Print the results of getNewColor() for all color pairs.
//
proc printColorEquations() {
    for c1 in Color do
        for c2 in Color do
            writeln(c1, " + ", c2, " -> ", getNewColor(c1, c2));
    writeln();
}

// Hold meetings among the population by creating a shared meeting
// place, and then creating per-chameneos tasks to have meetings.
//
proc holdMeetings(population, numMeetings) {
    const place = new MeetingPlace(numMeetings);

    coforall c in population do // create a task per chameneos
        c.haveMeetings(place, population);

    delete place;
}
```

*excerpt from 1210.gz Chapel entry*

```
void get_affinity(int* is_smp, cpu_set_t* affinity1, cpu_set_t* affinity2)
{
    cpu_set_t active_cpus;
    FILE* f;
    char buf [2048];
    pos;
    cpu_idx;
    physical_id;
    core_id;
    cpu_cores;
    apic_id;
    cpu_count;
    i;

    char const* processor_str = "processor";
    size_t processor_str_len = strlen(processor_str);
    char const* physical_id_str = "physical id";
    size_t physical_id_str_len = strlen(physical_id_str);
    char const* core_id_str = "core id";
    size_t core_id_str_len = strlen(core_id_str);
    char const* cpu_cores_str = "cpu cores";
    size_t cpu_cores_str_len = strlen(cpu_cores_str);

    CPU_ZERO(&active_cpus);
    sched_getaffinity(0, sizeof(active_cpus), &active_cpus);
    cpu_count = 0;
    for (i = 0; i != CPU_SETSIZE; i += 1)
    {
        if (CPU_ISSET(i, &active_cpus))
        {
            cpu_count += 1;
        }
    }

    if (cpu_count == 1)
    {
        is_smp[0] = 0;
        return;
    }

    is_smp[0] = 1;
    CPU_ZERO(affinity1);
```

*excerpt from 2863.gz C gcc entry*



COMPUTE

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# CLBG: Qualitative Comparisons



Can also browse program source code (but this requires actual thought):

```
proc main() {
    printColorEquations();

    const group1 = [i in 1..popSize1] new Chameneos(i, 0);
    const group2 = [i in 1..popSize2] new Chameneos(i, 0);

    cobegin {
        holdMeetings(group1, n);
        holdMeetings(group2, n);
    }

    print(group1);
    print(group2);

    for c in group1 do delete c;
    for c in group2 do delete c;
}

// Print the results of getNewColor() for all color pairs
// proc printColorEquations() {
//     for c1 in Color do
//         for c2 in Color do
//             writeln(c1, " + ", c2, " = ", getNewColor(c1, c2));
//             writeln();
// }

// Hold meetings among the population by creating a shared
// place, and then creating per-chameneos tasks to have
// them meet
proc holdMeetings(population, numMeetings) {
    const place = new MeetingPlace(numMeetings);

    coforall c in population do // create a task
        c.haveMeetings(place, population);

    delete place;
}
```

excerpt from 1210.gz Chapel entry

```
cobegin {
    holdMeetings(group1, n);
    holdMeetings(group2, n);
}
```

```
void get_affinity(int* is_smp, cpu_set_t* affinity1, cpu_set_t* affinity2)
```

```
size_t
char const*
size_t
char const*
```

```
proc holdMeetings(population, numMeetings) {
    const place = new MeetingPlace(numMeetings);

    coforall c in population do // create a task
        c.haveMeetings(place, population);

    delete place;
}
```

excerpt from 2863.gz C gcc entry

```
active_cpus;
f;
buf [2048];
pos;
cpu_idx;
physical_id;
core_id;
cpu_cores;
apic_id;
cpu_count;
i;

processor_str      = "processor";
processor_str_len = strlen(processor_str);
physical_id_str   = "physical id";
physical_id_str_len = strlen(physical_id_str);
core_id_str        = "core id";
n(core_id_str);
cores';
n(cpu_cores_str);
```

```
is_smp[0] = 1;
CPU_ZERO(affinity1);
```



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# CLBG: Qualitative Comparisons



Can also browse program source code (but this requires actual thought):

```
proc main() {
    char const* core_id_str = "core id";
    size_t core_id_str_len = strlen(core_id_str);
    char const* cpu_cores_str = "cpu cores";
    size_t cpu_cores_str_len = strlen(cpu_cores_str);

    CPU_ZERO(&active_cpus);
    sched_getaffinity(0, sizeof(active_cpus), &active_cpus);
    cpu_count = 0;
    for (i = 0; i != CPU_SETSIZE; i += 1)
    {
        if (CPU_ISSET(i, &active_cpus))
        {
            cpu_count += 1;
        }
    }

    if (cpu_count == 1)
    {
        is_smp[0] = 0;
        return;
    }
}
```

excerpt from 1210.gz Chapel entry

```
void get_affinity(int* is_smp, cpu_set_t* affinity1, cpu_set_t* affinity2)
{
    cpu_set_t active_cpus;
    FILE* f;
    char buf [2048];
    pos;
    cpu_idx;
    physical_id;
    core_id;
    cpu_cores;
    apic_id;
    cpu_count;
    i;

    char const* processor_str = "processor";
    size_t processor_str_len = strlen(processor_str);
    physical_id_str = "physical id";
    physical_id_str_len = strlen(physical_id_str);
    core_id_str = "core id";
    core_id_str_len = strlen(core_id_str);
    cpu_cores_str = "cpu cores";
    cpu_cores_str_len = strlen(cpu_cores_str);

    CPU_ZERO(&active_cpus);
    sched_getaffinity(0, sizeof(active_cpus), &active_cpus);
    cpu_count = 0;
    for (i = 0; i != CPU_SETSIZE; i += 1)
    {
        if (CPU_ISSET(i, &active_cpus))
        {
            cpu_count += 1;
        }
    }

    if (cpu_count == 1)
    {
        is_smp[0] = 1;
        CPU_ZERO(affinity1);
    }
}
```

excerpt from 2863.gz C gcc entry



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*I have run  
I have crawled  
I have scaled...*

## Chapel Scalability Results

U2

I Still Haven't Found What I'm Looking For  
*The Joshua Tree*

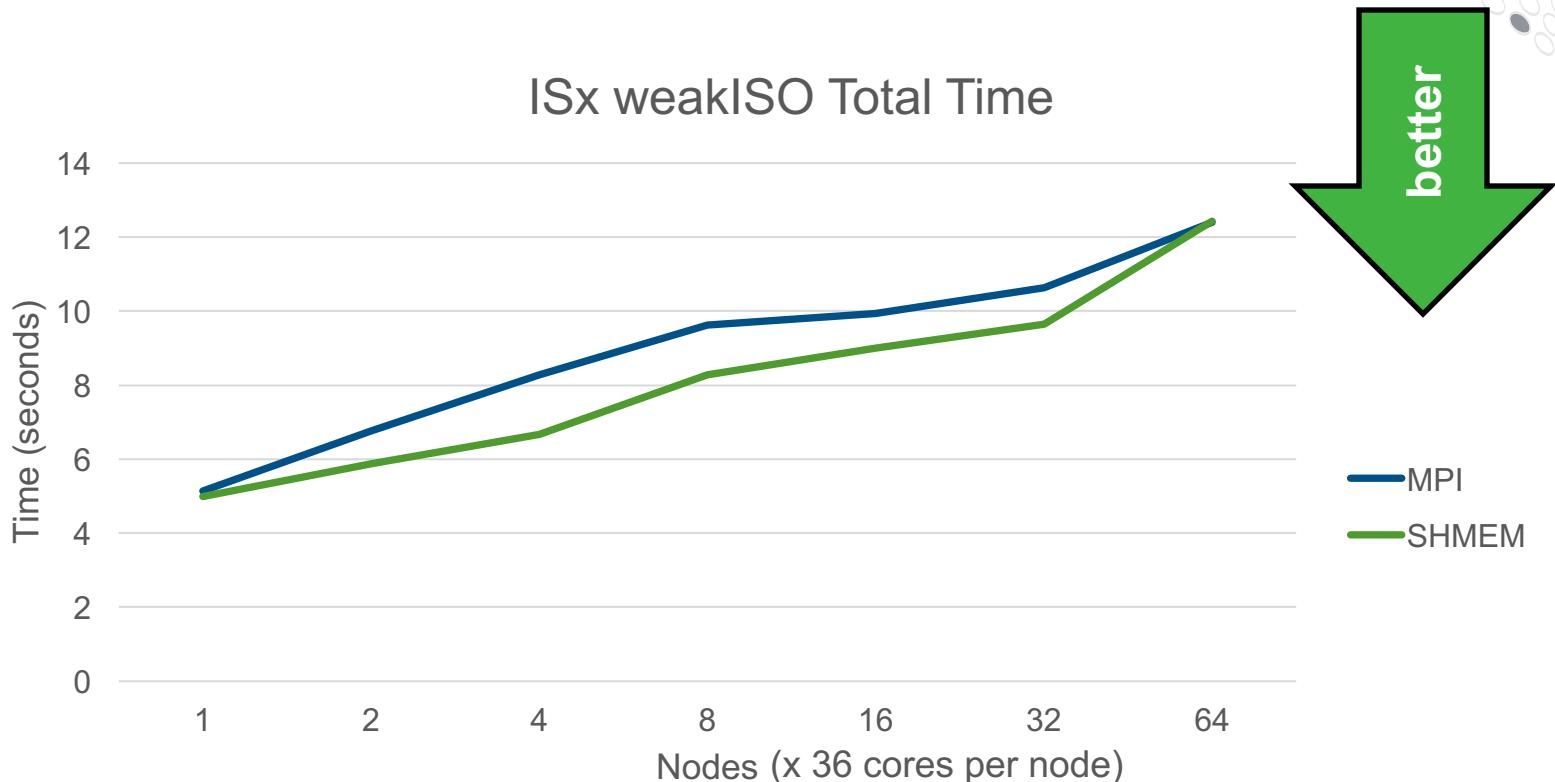


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# ISx Execution Time: MPI, SHMEM

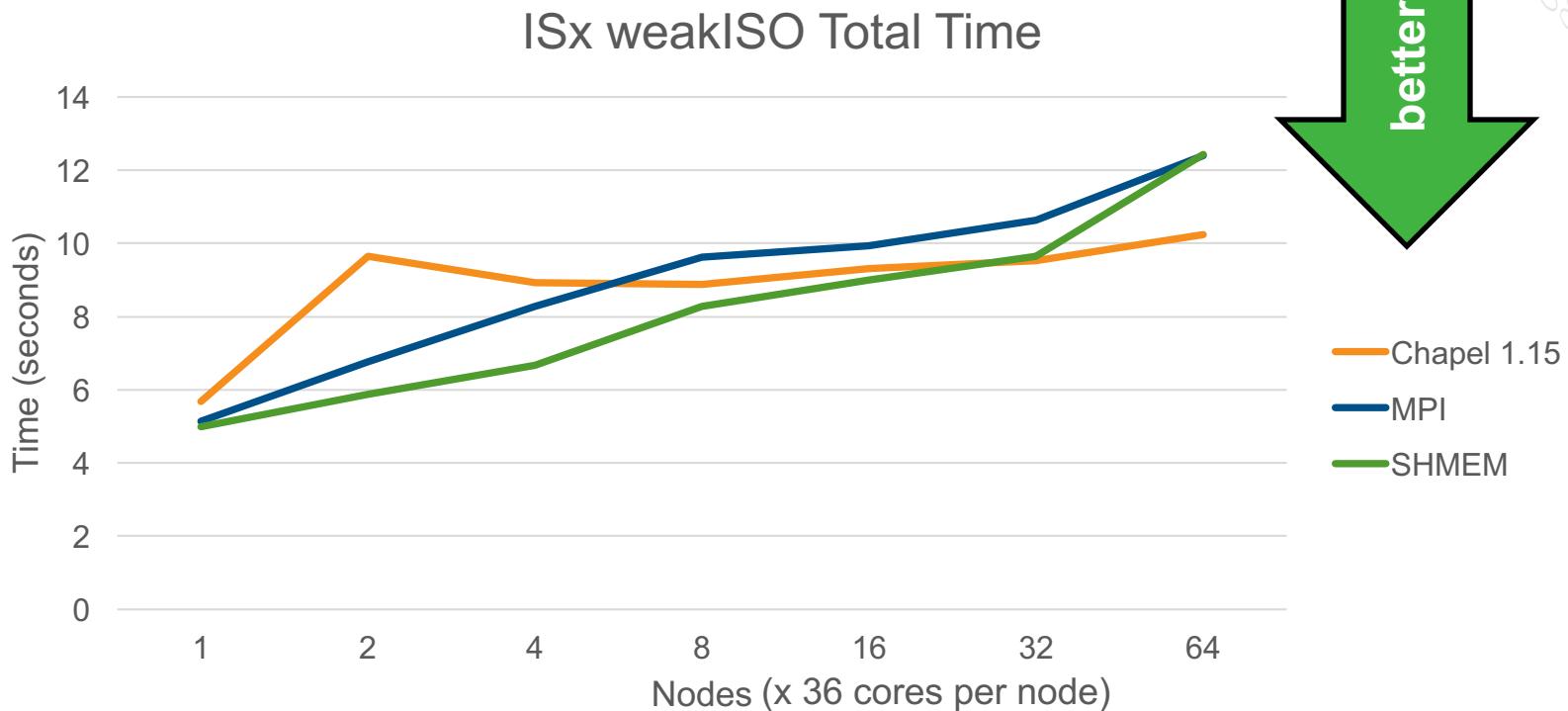


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# ISx Execution Time: MPI, SHMEM, Chapel

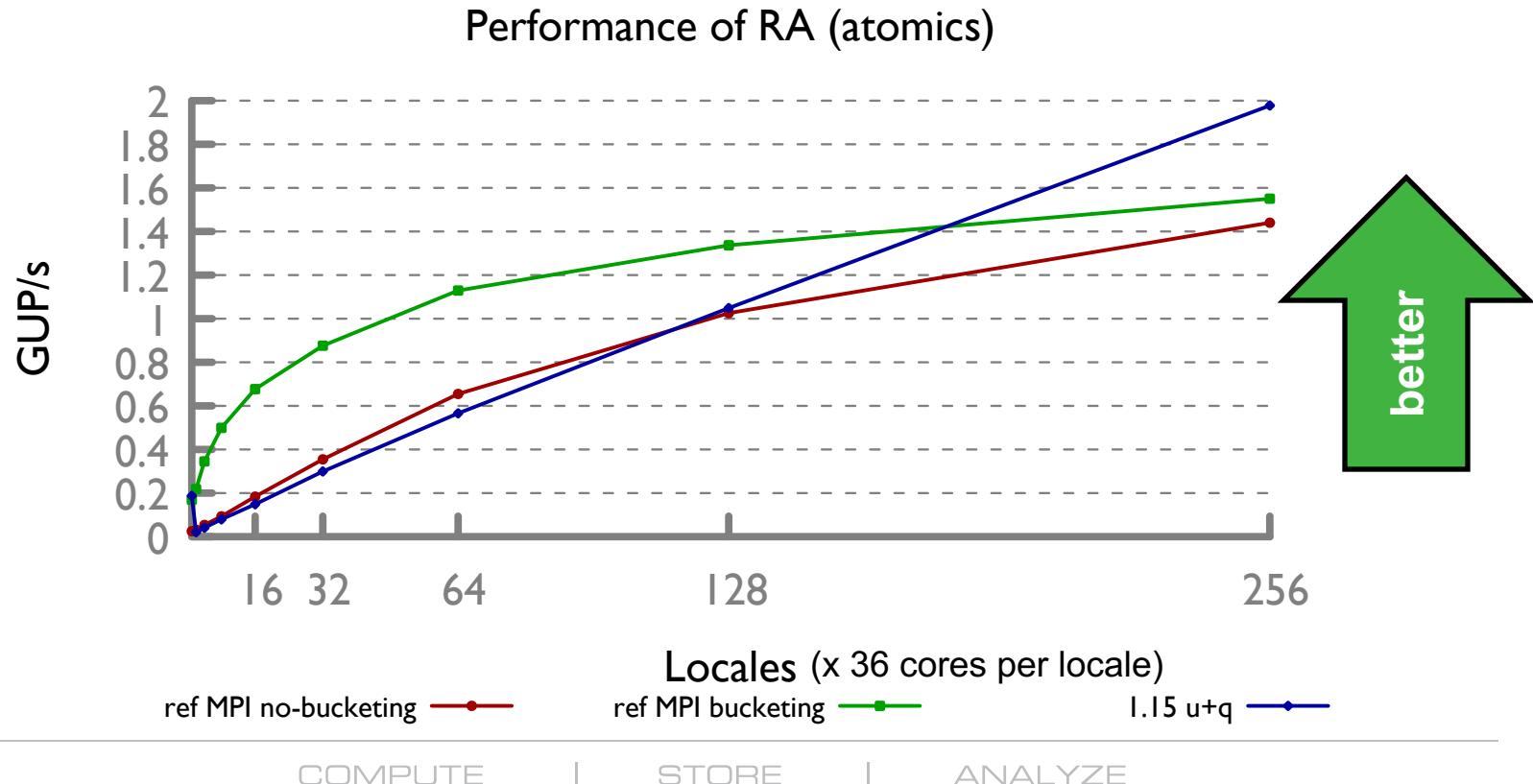


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# RA Performance: Chapel vs. MPI



*[At? And? Ah'm?] fourteen, and you know  
That I've learned the easy way  
Some stupid decisions  
And with a...a broken heart*

## Summary

**That Summer, at Home I Had Become The Invisible Boy**  
The Twilight Sad  
*Fourteen Autumns and Fifteen Winters*



# This Talk's Thesis



**Programming language designers have, to date, largely failed the large-scale parallel computing community.**

- parallel features have historically been an afterthought
  - tacked on through libraries, pragmas, and extensions
  - rarely first-class
- even when languages are designed for parallelism...
  - ...most fail to consider scalable computing (distributed memory)
  - ...others tend to be domain-specific and not very general

**We can do better!**



# We think Chapel is an attractive candidate



- **Attractive language for end-users**
  - modern design
  - separates algorithmic specification from mapping to architecture
- **Able to achieve competitive performance**
  - though a good deal of tuning and optimization work remain...



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# Recap of This Talk's Takeaways



## If you design a parallel programming language...

- ...don't tie yourself to architecture-specific mechanisms
- ...don't base your model for parallelism and locality on SPMD
- ...create attractive ways of expressing parallelism and locality
- ...support first-class index sets
- ...support multiresolution features



# Chapel Challenges: Technical



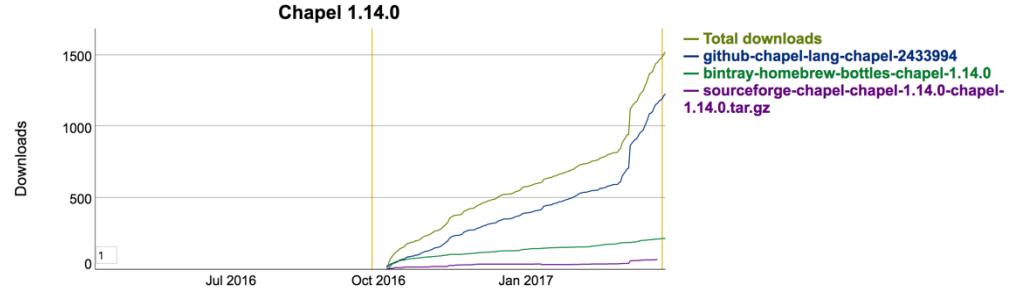
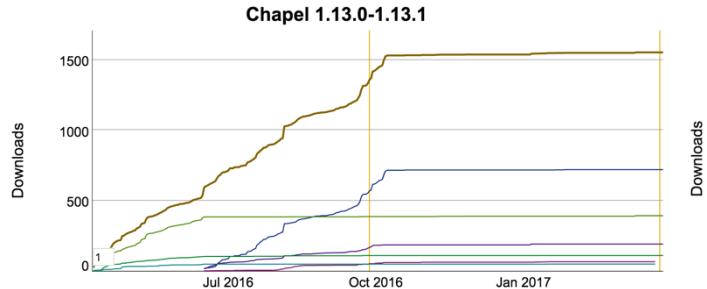
- **design and optimization of multiresolution language features:**
  - domain maps
  - parallel iterators
  - locale models
- **compiler architecture: outgrowing the initial research prototype**
- **tools: classic chicken-and-egg problem**
  - IDE
  - debuggers
  - performance analysis
  - package manager
  - interactive development (interpreter, REPL, iPython, ...)



# Chapel Challenges: Social



- building up momentum in the user community
  - lots of interest, but also lots of fear of being first / only adopter
  - 3000+ downloads / year for 2 releases



- combatting impatience / numbness to our message and progress
- developing an aggressive language in a conservative field (HPC)
- engaging with peers in mainstream computing



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***Scratch my name on your arm with a fountain pen  
This means you really love me***

## Further Resources

**Rusholme Ruffians**  
The Smiths  
*Meat is Murder*



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# How to Stalk Chapel



<http://facebook.com/ChapelLanguage>

<http://twitter.com/ChapelLanguage>

<https://www.youtube.com/channel/UCHmm27bYjhknK5mU7ZzPGsQ/>  
[chapel-announce@lists.sourceforge.net](mailto:chapel-announce@lists.sourceforge.net)



**Chapel Language**  
@ChapelLanguage

Chapel is a productive parallel programming language designed for large-scale computing whose development is being led by @cray\_inc

 [chapel.cray.com](#)  
 Joined March 2016  
  
 115 Photos and videos



**TWEETS 222 FOLLOWING 12 FOLLOWERS 129 LINES 32**

---

**Tweets**   [Tweets & replies](#)   [Media](#)

 **Chapel Language** @ChapelLanguage · 5h  
Doing interesting applications work in Chapel or another PGAS language?  
Submit it to the PAW 2017 workshop at [GSC17](#).  
[sourceryinstitute.github.io/PAW/](#)



## The 2nd Annual PGAS Applications Workshop

Montreal, Quebec, Canada | April 10-11, 2017  
9:00 am - 5:00 pm  
Hosted in conjunction with SOSC17  
The International Conference for High Performance Computing, Networking, Storage and Analysis



In cooperation with



**Chapel Parallel Programming Language**

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**Chapel videos**

**SC16 Chapel Tutorial Promo**  
Chapel Parallel Programming Language  
6 months ago • 277 views

This is a 4-6 minute promotional video for our SC16 Chapel tutorial, and also a good way to get a quick taste of Chapel. All codes shown represent complete Chapel programs, not...

**Chapel Productive, Multiresolution Parallel Programming | Brad Chamberlain, Cray, Inc.**  
ANL Training  
7 months ago • 651 views

Presented at the Argonne Training Program on Extreme-Scale Computing, Summer 2016.

**CHIWU 2016 keynote: "Chapel in the (Cosmological) Wild", Nikhil Padmanabhan**  
Chapel Parallel Programming Language  
10 months ago • 277 views

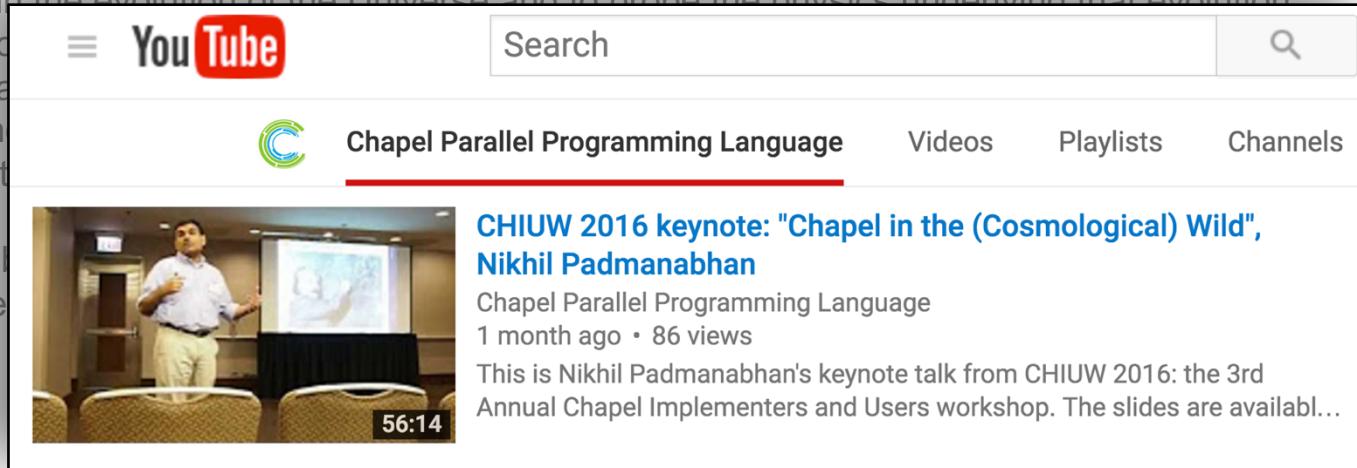
This is Nikhil Padmanabhan's keynote talk from CHIWU 2016: the 3rd Annual Chapel Implementers and Users workshop. The slides are available at...

# A Chapel talk to watch next (user perspective)

## Chapel in the (Cosmological) Wild (CHIUW 2016 keynote)

**Nikhil Padmanabhan**, *Yale University Professor, Physics & Astronomy*

**Abstract:** This talk aims to present my personal experiences using Chapel in my research. My research interests are in observational cosmology; more specifically, I use large surveys of galaxies to constrain the evolution of the Universe and to probe the physics underlying that evolution.



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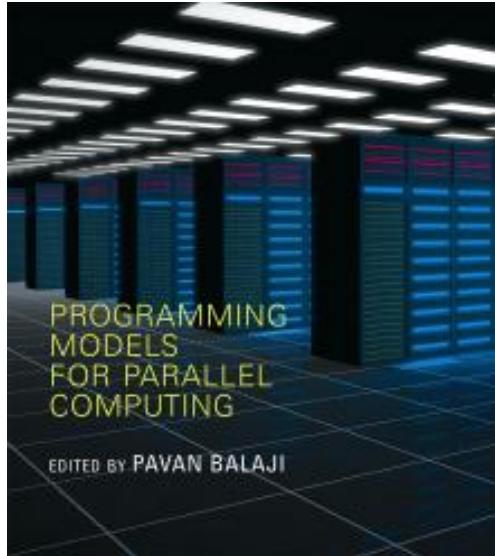
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# Suggested Reading (healthy attention spans)



Chapel chapter from [Programming Models for Parallel Computing](#)

- a detailed overview of Chapel's history, motivating themes, features
- published by MIT Press, November 2015
- edited by Pavan Balaji (Argonne)
- chapter is now also available [online](#)



Other Chapel papers/publications available at <http://chapel.cray.com/papers.html>



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# Suggested Reading (short attention spans)



[Chapel: Productive Parallel Programming](#), [Cray Blog](#), May 2013.

- *a short-and-sweet introduction to Chapel*

[Six Ways to Say “Hello” in Chapel](#) (parts [1](#), [2](#), [3](#)), [Cray Blog](#), Sep-Oct 2015.

- *a series of articles illustrating the basics of parallelism and locality in Chapel*

[Why Chapel?](#) (parts [1](#), [2](#), [3](#)), [Cray Blog](#), Jun-Oct 2014.

- *a series of articles answering common questions about why we are pursuing Chapel in spite of the inherent challenges*

[\[Ten\] Myths About Scalable Programming Languages](#), [IEEE TCSC Blog](#)

([index available on chapel.cray.com “blog articles” page](#)), Apr-Nov 2012.

- *a series of technical opinion pieces designed to argue against standard reasons given for not developing high-level parallel languages*





*All we ever wanted was everything...*

(I have Chapel swag to give you)



*Get up  
Eat jelly  
sandwich bars, ....*

(Have a good break and get some yummy snacks!)

The logo for ACCU 2017 features the word 'ACCU' in a large, orange, blocky font above the year '2017' in a similar style. The 'A' in 'ACCU' is slightly taller than the other letters.

All We Ever Wanted Was Everything  
Bauhaus  
*The Sky's Gone Out*



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