

# Database technology

## TDDD37

### Lab 3 : Normalization

#### Question 1

Considering **R(A,B,C,D,E,F)** with the following FDs:

FD1: **A**→**BC**

FD2: **C**→**AD**

FD3: **DE**→**F**

Use the Armstrong rules to derive the following FD: **C**→**B**

FD4: **C**→**A**                      (*decomposition of FD2*)

FD5: **A**→**B**                      (*decomposition of FD1*)

FD6: **C**→**B**                      (*transitivity of FD4 and FD5*)

Use the Armstrong rules to derive the following FD: **AE**→**F**

FD4: **A**→**C**                      (*decomposition of FD1*)

FD5: **C**→**D**                      (*decomposition of FD2*)

FD6: **A**→**D**                      (*transitivity of FD4 and FD5*)

FD7: **AE**→**DE**                  (*augmentation of FD6 with E*)

FD7: **AE**→**F**                      (*transitivity of FD7 and FD3*)

#### Question 2

Considering **R(A,B,C,D,E,F)** with the following FDs:

FD1: **A**→**BC**

FD2: **C**→**AD**

FD3: **DE**→**F**

Compute the attribute closure of **X = {A}**

*Initially*                      **X<sup>+</sup> = {A}**

*By using FD1:*              **X<sup>+</sup> = {A,B,C}**

*By using FD2:*              **X<sup>+</sup> = {A,B,C,D}**

Compute the attribute closure of **X = {C,E}**

*Initially*                      **X<sup>+</sup> = {C,E}**

*By using FD2:*              **X<sup>+</sup> = {A,C,D,E}**

*By using FD1:*              **X<sup>+</sup> = {A,B,C,D,E}**

*By using FD3:*              **X<sup>+</sup> = {A,B,C,D,E,F}**

### Question 3

Considering **R(A,B,C,D,E,F)** with the following FDs:

FD1: **AB**→**CDEF**

FD2: **E**→**F**

FD3: **D**→**B**

Candidate keys for R are **{A,B}** and **{A,D}**.

FD2 and FD3 violate the BCNF condition.

Let's decompose R based on FD2:

Relation	FDs	Candidate key(s)
<b>R1(E,F)</b>	FD2	<b>{E}</b>
<b>R2(A,B,C,D,E)</b>	FD3 and FD4: <b>AB</b> → <b>CDE</b>	<b>{AB}</b>

R1 is in BCNF but R2 isn't. Let's decompose R2 based on FD3:

Relation	FDs	Candidate key(s)
<b>R3(D,B)</b>	FD3	<b>{D}</b>
<b>R4(A,C,D,E)</b>	FD5: <b>A</b> → <b>CDE</b>	<b>{A}</b>

R3 and R4 are in BCNF.

The decomposition of R into a set of BCNF relations is: **{R1, R3, R4}**.

### Question 4

Considering **R(A,B,C,D,E)** with the following FDs:

FD1: **ABC**→**DE**

FD2: **BCD**→**AE**

FD3: **C**→**D**

The candidate key for R is **{B,C}**.

FD3 violate the BCNF condition.

Let's decompose R based on FD3:

Relation	FDs	Candidate key(s)
<b>R1(C,D)</b>	FD3	<b>{C}</b>
<b>R2(A,B,C,E)</b>	FD4: <b>ABC</b> → <b>E</b> FD5: <b>BC</b> → <b>AE</b>	<b>{BC}</b>

R1 and R2 are in BCNF.

The decomposition of R into a set of BCNF relations is: **{R1, R2}**.