

Database technology

TDDD37

Lab 3 : Normalization

Question 1

Considering **R(A,B,C,D,E,F)** with the following FDs:

FD1: **A**→**BC**

FD2: **C**→**AD**

FD3: **DE**→**F**

Use the Armstrong rules to derive the following FD: **C**→**B**

FD4: **C**→**A** (*decomposition of FD2*)

FD5: **A**→**B** (*decomposition of FD1*)

FD6: **C**→**B** (*transitivity of FD4 and FD5*)

Use the Armstrong rules to derive the following FD: **AE**→**F**

FD4: **A**→**C** (*decomposition of FD1*)

FD5: **C**→**D** (*decomposition of FD2*)

FD6: **A**→**D** (*transitivity of FD4 and FD5*)

FD7: **AE**→**DE** (*augmentation of FD6 with E*)

FD7: **AE**→**F** (*transitivity of FD7 and FD3*)

Question 2

Considering **R(A,B,C,D,E,F)** with the following FDs:

FD1: **A**→**BC**

FD2: **C**→**AD**

FD3: **DE**→**F**

Compute the attribute closure of **X = {A}**

Initially **X⁺ = {A}**

By using FD1: **X⁺ = {A,B,C}**

By using FD2: **X⁺ = {A,B,C,D}**

Compute the attribute closure of **X = {C,E}**

Initially **X⁺ = {C,E}**

By using FD2: **X⁺ = {A,C,D,E}**

By using FD1: **X⁺ = {A,B,C,D,E}**

By using FD3: **X⁺ = {A,B,C,D,E,F}**

Question 3

Considering **R(A,B,C,D,E,F)** with the following FDs:

FD1: **AB**→**CDEF**

FD2: **E**→**F**

FD3: **D**→**B**

Candidate keys for R are **{A,B}** and **{A,D}**.

FD2 and FD3 violate the BCNF condition.

Let's decompose R based on FD2:

Relation	FDs	Candidate key(s)
R1(E,F)	FD2	{E}
R2(A,B,C,D,E)	FD3 and FD4: AB → CDE	{AB},{AD}

R1 is in BCNF but R2 isn't. Let's decompose R2 based on FD3:

Relation	FDs	Candidate key(s)
R3(D,B)	FD3	{D}
R4(A,C,D,E)	FD5: AD → CDE Explanation: FD3: D→B by augmentation with A: AD→AB FD4: AB→CDE by transitivity: AD→CDE	{AD}

R3 and R4 are in BCNF.

The decomposition of R into a set of BCNF relations is: **{R1, R3, R4}**.

Question 4

Considering **R(A,B,C,D,E)** with the following FDs:

FD1: **ABC**→**DE**

FD2: **BCD**→**AE**

FD3: **C**→**D**

The candidate key for R is **{B,C}**.

FD3 violate the BCNF condition.

Let's decompose R based on FD3:

Relation	FDs	Candidate key(s)
R1(C,D)	FD3	{C}
R2(A,B,C,E)	FD4: ABC → E FD5: BC → AE	{BC}

R1 and R2 are in BCNF.

The decomposition of R into a set of BCNF relations is: **{R1, R2}**.