

# Fast/Discrete Fourier Transform

Group 23

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# Fourier Transform

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## Definition

Fourier transform is a decomposition of a function of time (a signal) into frequencies that make it up.

- The main strategy is to see signals as accumulations of sine waves of different frequencies, and this can be represented as the following

$$f(x) = \int_{-\infty}^{\infty} F(s) e^{2i\pi sx} ds$$

# Discrete Fourier Transform

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## Limitation on FT

In fourier transform, both data in the time domain and frequency domain are **continuous**, which are not feasible to store and process in computers.

Discrete Fourier Transform (DFT) solves this question as it processes a **finite sample sequence** of time domain data and convert it into a same-length sample data for frequency domain.

# Mathematical Formula

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DFT will decompose a sequence  $\{x_n\}$  of length  $N$  into the following style

$$x_n = \frac{1}{N} \sum_{k=0}^{N-1} X_k \cdot e^{i2\pi kn/N}$$

And thus the sequence  $X_k$  which represents how much a certain frequency is in the sample can be written as

$$X_k = \sum_{n=0}^{N-1} x_n \cdot e^{i2\pi kn/N}$$

# Fast Fourier Transform

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If the calculation of  $DFT$  is operated naively, then it is clear that the time complexity is  $\mathcal{O}(n^2)$  for a length- $n$  sequence.

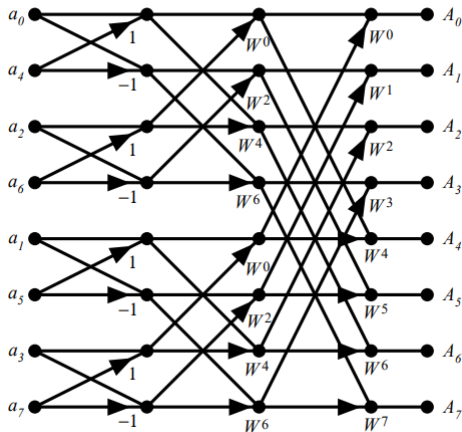
And thus Fast Fourier Transform(FFT) is introduced, which serves to find the internal mathematical basis and simplify the calculation, usually to  $\mathcal{O}(n \log n)$ . And here the algorithm introduced will be [Cooley-Tukey algorithm](#)

The algorithm uses the idea of divide and conquer, so here the case when  $n = 2^p$  is considered.

# Cooley-Tukey Algorithm

$a$ : input sequence,  $A$ : output sequence

$W$ :  $e^{-i2\pi/n}$ ,  $\text{weight}_{j,k} = W^{\frac{8}{2^j} \times k}$



# Discussion on Bitwise Reverse

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The most naive method in bitwise reverse is iteratively divide the number by 2 to get all the bits. In this way reversing one number should be  $\mathcal{O}(\log n)$  complexity.

I found that when the number of bits is a power of 2, this can do it in the way of flipping every adjacent 2,4,...,n/2 bits to reverse it.

So the algorithm I am using now is to find the smallest 2-power number larger than the number of bits now, flipping in that way then right shift the extra 0s in the end.



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References

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- [https://en.wikipedia.org/wiki/Fast\\_Fourier\\_transform](https://en.wikipedia.org/wiki/Fast_Fourier_transform)