

Fast/Discrete Fourier Transform

Group 23

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Fourier Transform

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Definition

Fourier transform is a decomposition of a function of time (a signal) into frequencies that make it up.

- The main strategy is to see signals as accumulations of sine waves of different frequencies, and this can be represented as the following

$$f(x) = \int_{-\infty}^{\infty} F(s) e^{2i\pi sx} ds$$

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Limitation on FT

In fourier transform, both data in the time domain and frequency domain are **continuous**, which are not feasible to store and process in computers.

Discrete Fourier Transform (DFT) solves this question as it processes a **finite sample sequence** of time domain data and convert it into a same-length sample data for frequency domain.

Mathematical Formula

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DFT will decompose a sequence $\{x_n\}$ of length N into the following style

$$x_n = \frac{1}{N} \sum_{k=0}^{N-1} X_k \cdot e^{i2\pi kn/N}$$

And thus the sequence X_k which represents how much a certain frequency is in the sample can be written as

$$X_k = \sum_{n=0}^{N-1} x_n \cdot e^{i2\pi kn/N}$$

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If the calculation of DFT is operated naively, then it is clear that the time complexity is $\mathcal{O}(n^2)$ for a length- n sequence.

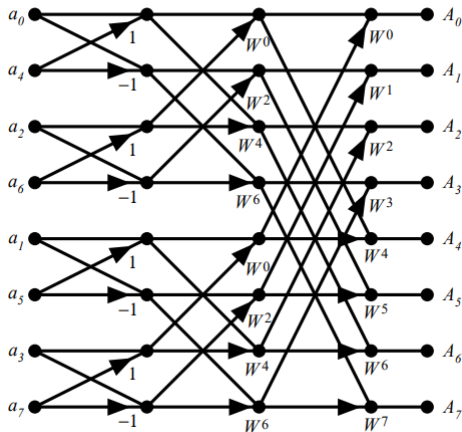
And thus Fast Fourier Transform(FFT) is introduced, which serves to find the internal mathematical basis and simplify the calculation, usually to $\mathcal{O}(n \log n)$. And here the algorithm introduced will be [Cooley-Tukey algorithm](#)

The algorithm uses the idea of divide and conquer, so here the case when $n = 2^p$ is considered.

Cooley-Tukey Algorithm

a : input sequence, A : output sequence

W : $e^{-i2\pi/n}$, $\text{weight}_{j,k} = W^{\frac{8}{2^j} \times k}$



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