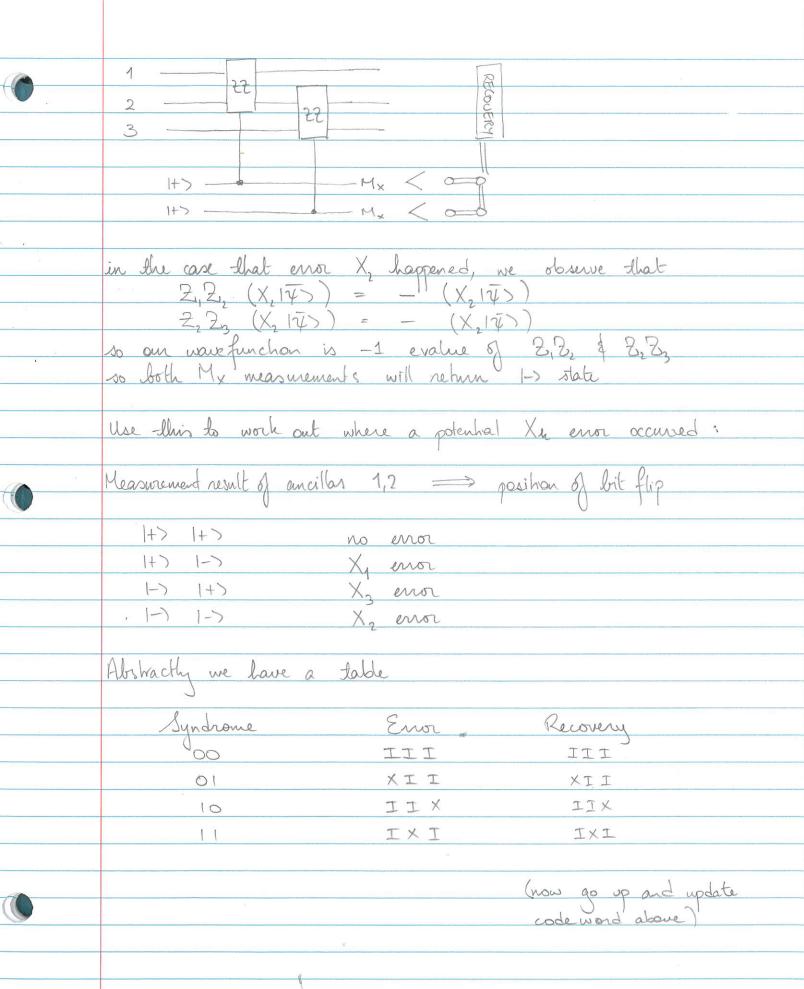
Indirect measurement. Consider I+> Mx where $A^2 = 1$, bottom line is arbitrary number of qubits So A acts on 17> e C2n $A^2 = 1$ => (A+1)(A-1)=0 => eigenvalues of A are ± 1 +1 eigenspace What are the projectors T_{\pm} ? $\frac{I\pm A}{2}$ Proof: Is TI, 14) = 1+A 14) really a. +1 eigenvector of A? $A\left(\Pi_{+}|\psi\rangle\right) = A\left(\frac{1+A}{2}\right)|\psi\rangle = \left(\frac{A+I}{2}\right)|\psi\rangle = \left(\Pi_{+}|\psi\rangle\right)|\Pi$ One annoying thing in that T+14> aren't of unit length... Let's write $|\gamma\rangle = \propto |\gamma+\rangle + \beta |\gamma-\rangle$ where $|\gamma+\rangle = \pi_+ |\gamma-\rangle$ $\langle \gamma | \pi_+ \gamma \rangle$ and $\langle \gamma \rangle = \langle \gamma | \pi_+ \gamma \rangle$ Analyse indirect measurement in three stages

Stage 1 1+> 8 14> = 1/2 (10> + 11>) 8/4> = (2 10) 8 14) + (2 11) 8 14) Stage 2 应100 × I/4> + 应11 × A14> = 12 10) 8 (a14+)+ (14-) + 12 11) (a14+) - (14-) = \f((10) + 11) \omega \alpha | \psi \((10) - 11) \omega \beta | \psi \) = x/+> × 1/+> + B/-> × 1/4-> Measure X will either see H) or I-> and the probabilities are 1012 respectively. if X measures H> then H> is definitely in 14+7 Remark that if I wanted to stay in 14+ > state, but my ancille was read out as 1-> (so I am a chally in 14->) then hopefully I have a procedure / program that I can apply so that 14-> Program 14+>

How do we quantumly error correct? Early ideas in quantum information were just translations of classical ideas into quantum notation. But error correction can't even begin due to no cloning. majority vote 0 -> 000 enor with order p - one with order p2 14> --> 14>14>14> Challenger - no downe theorem - measuring date in order to determine errors will destroy - must correct 4,2 errors, not just X errors - must correct infinite set of unitary errors (and channels which decohere state) (unitary channels, pauli emois, dephasing depolarizing, amplitude damping, photon loss, multiqubit channels) the most difficult ~ errors during computation ~ leading to notion of fault bolerance (only to be louched upon at end)

A quantum code de protect aganst bit flips. 14> = 0 10) + B 11) encoding, 0 1000) + B 1111) = 14> 10) D X 1000) + B M1) (this seems to have excumvented no-cloning problem) Pretend we get a bit flip error on 2nd qubit X, 17) = x (010) + B (101) the subtle point is to realise the middle qubit is different from gulito 1,3 inespective of state 14) rather than doing measurement on qubits 1,23 (which would destroy superposition) we can measure pairwise parity between qubits. measure parity between 1,2 = measure 2,2 eigenvalue superposition 1+) => same parity

(+) => opposite parity do dhis also for parity between gubits 2,3



We could go and book at phase flip errors (thin would be worth while) but I'm more interested in the continuous version of a bit flip error: $R_{O}(X) = \exp(-i\theta X)$ $= \cos \theta I - i \sin \theta X$ Suppose I got this on the first qubit: Codeword 17/5 = 0(500) + BMM) -> Ro(X1) 1/2> = cost 1/2> -2 sm0 X1/2> Note that 2, 2, 17) = +1 17) 2,2, (x, 17) = -1 (x, 17) H) - Mx two opinions: either measure 1+> in prob costo If measure 1+> then wavevector projects to 1+> "
1-> X17> " This has DISCRETIZED the error! So we can return to previous problem where we just lad Pauli-bit-flip errors. seems lo have circumvented problem of infinite set of unitary errors.