

ÉCOLE POLYTECHNIQUE DE BRUXELLES

INFO-F521

Graphs and Networks

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Chapter 1

The Basics

1.1 Graphs

A *graph* is a pair $G = (V, E)$ of sets such that $E \subseteq [V]^2$: the elements of E are 2-element subsets of V . The elements of V are the *vertices* (or *nodes*, or *points*) and the elements of E are the *edges* (or *lines*).

The number of vertices of a graph G is its *order*, written as $|G|$. The number of edges is noted $||G||$.

Two vertices x, y of G are *adjacent* or *neighbours* if xy is an edge of G . Two edges $e \neq f$ are adjacent if they have an end in common. If all the vertices of G are pairwise adjacent, then G is complete. A complete graph on n vertices is a K^n - a K^3 is a triangle.

Let $G = (V, E)$ and $G' = (V', E')$ be two graphs. We call G and G' *isomorphic*, and write $G \simeq G'$, if there exists a bijection $\varphi : V \rightarrow V'$ with $xy \in E \iff \varphi(x)\varphi(y) \in E'$ for all $x, y \in V$. Such a map φ is called an *isomorphism*; if $G = G'$, it is called an *automorphism*.

We set $G \cup G' := (V \cup V', E \cup E')$, the intersection is defined similarly with \cap . If $G \cap G' = \emptyset$, then G and G' are *disjoint*. If $V' \subseteq V$ and $E' \subseteq E$, then G' is a *subgraph* of G and G is a supergraph of G' .

If $G' \subseteq G$ and G' contains all the edges $xy \in E$ with $x, y \in V'$, then G' is an *induced subgraph* of G ; we say that V' *induces* or *spans* G' in G , and write $G[V'] := G'$.

A *spanning subgraph* G' of a graph G is a subgraph that includes all of the vertices of G ($V' = V$).

1.2 The degree of a vertex

Let $G = (V, E)$ be a non-empty graph. The set of neighbours of a vertex v in G is denoted $N_G(v)$, or briefly $N(v)$.

The *degree* (or *valency*) $d_G(v) = d(v)$ of a vertex v is the number $|E(v)|$ of edges at v ; this is equal to the number of neighbours of v . A vertex of degree 0 is *isolated*.

The number $\delta(G) := \min \{ d(v) \mid v \in V \}$ is the *minimum degree* of G , the number $\Delta(G) := \max \{ d(v) \mid v \in V \}$ is the *maximum degree*.

If all the vertices of G have the same degree k , then G is *k-regular*, or simply *regular*. A 3-regular graph is called *cubic*.

A *k-factor* is a k -regular spanning subgraph.

1.3 Paths and cycles

A *path* is a non-empty graph $P = (V, E)$ of the form:

$$V = \{x_0, x_1, \dots, x_k\} \quad E = \{x_0x_1, x_1x_2, \dots, x_{k-1}x_k\},$$

where the x_i are all distinct. The vertices x_0 and x_k are *linked* by P and are called its *endvertices* or *ends*. The other vertices are called the *innervertices*.

For P a path from x_0 to x_k , we write with $0 \leq i \leq j \leq k$:

$$\begin{aligned} Px_i &:= x_0 \dots x_i \\ x_i P &:= x_i \dots x_k \\ x_i P x_j &:= x_i \dots x_j \end{aligned}$$

1.4 Connectivity

A non-empty graph G is called *connected* if any two vertices are linked by a path in G . If $U \subseteq V(G)$ and $G[U]$ is connected, we also call U itself connected in G .

Disconnected = not connected.

A graph G is *k-connected* (for $k \in \mathbb{N}$) if $|G| > k$ and $G - X$ is connected for every set $X \subseteq V$ with $|X| < k$. In other words, no two vertices of G are separated by fewer than k other vertices.

1.5 Bipartite graphs

Let $r \geq 2$ be an integer. A graph $G = (V, E)$ is called *r-partite* if V admits a partition into r classes such that every edge has its ends in different classes: vertices in the same partition class must not be adjacent. Instead of "2-partite", one usually says *bipartite*.

An r -partite graph in which every two vertices from different partition classes are adjacent is called *complete*; the complete r -partite graphs for all r together are the *complete multipartite graphs*.

Chapter 2

Matchings

Given the graph $G = (V, E)$, the set $M \subseteq E$ is a *matching* if all edges in M are independent. That is, no two edges share an endpoint. The size of a matching is the number of edges it contains.

For the graph G , $X \subseteq V$ is a *vertex cover* if every edge in E has at least one endpoint in X .

A *perfect matching* (a.k.a. 1-factor) is a matching with $|M| = \frac{|V|}{2}$. That is, every vertex of the graph is incident to exactly one edge of the matching. Every perfect matching is maximum and hence maximal. In some literature, the term complete matching is used. A perfect matching is also a minimum-size edge cover.

Observation: if M is a matching and X a vertex cover, then $|X| \geq |M|$. Corollary: the minimum size of a vertex cover is \geq the maximum size of a matching.

2.1 Matching in bipartite graphs

For this whole section, we let $G = (V, E)$ be a fixed bipartite graph with bipartition $\{A, B\}$. Vertices denoted as a, a' , etc. will be assumed to lie in A , vertices noted as b etc. will lie in B .

An *alternating path* (w.r.t. some matching M) is a path in G that starts at an unmatched vertex in A and alternates between edges in $E - M$ and edges in M (it can be made of only one edge).

An *augmenting path* (w.r.t. some matching M) is an alternating path whose end will be at an unmatched vertex (of B). Note that if P is an augmenting path, then $M \Delta E(P)$ ¹ is again a matching, of size $|M| + 1$.

Alternating paths play an important role in the practical search for large matchings. In fact, if we start with any matching and keep applying augmenting paths until no

¹ Δ , a symmetric difference, defined as $X \Delta Y = (X - Y) \cup (Y - X) = (X \cup Y) - (X \cap Y)$

further such improvement is possible, the matching obtained will always be an optimal one, a matching with the largest possible number of edges.

König's theorem (1931)

The maximum cardinality of a matching in G is equal to the minimum cardinality of a vertex cover.

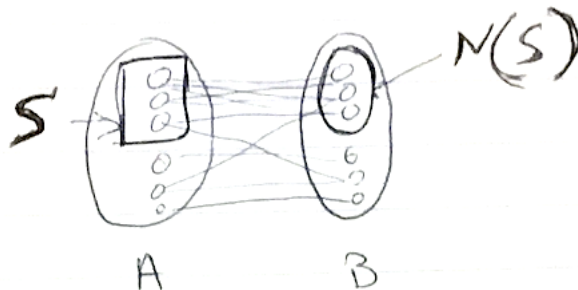
Let $G = (V, E)$ be a bipartite graph with bipartition $V = (A, B)$. Then the maximum size of a matching in G is the minimum size of a vertex cover in G .

1. Let M be a maximum matching. From every edge in M let us choose one of its ends: its end in B if some alternating path ends in that vertex, and its end in A otherwise. We shall prove that this set of these $|M|$ vertices - let us name it U - covers E ; since any vertex cover of E must cover M , there can be none with fewer than $|M|$ vertices, and so the theorem will follow.
2. Let $ab \in E$ be an edge; we show that either a or b lies in U . If $ab \in M$, this holds by definition of U - so we assume that $ab \notin M$. Since M is a maximal matching, there is an edge $a'b' \in M$ with $a = a'$ or $b = b'$. For all $ab \notin M$, we have:
 - (a) First, let's assume $b = b'$: a is the unmatched start of ab , which is then an alternating path, and so the end of $a'b' \in M$ chosen for U was the vertex $b' = b$.
 - (b) Now with $a' = a$. If the start $a = a'$ is not in U , then $b' \in U$ (as $a'b'$ is an edge in M), and some alternating path P ends in b' (by definition of U). But then there is also an alternating path P' ending in b , either $P := Pb$ (if $b \in P$) or $P' := Pb'a'b$. By the maximality of M , however, P' can't be an augmenting path. So b must be matched, and was chosen for U from the edge of M containing it.

Hall's theorem (1935)

If G is bipartite with partitions A, B , it contains a matching of A if and only if $|N(S)| \geq |S|$ for all $S \subseteq A$ (marriage condition).

Here, $N(S) = \{x \in B \text{ s.t. } x \text{ is adjacent to some vertex in } S\}$.



Let $G = (V, E)$ be a bipartite graph with bipartition $V = (A, B)$. Then (keeping in mind that for each partition X , any $x \in X$ isn't connected to any other $x' \in X$):

$$A \text{ can be matched} \iff |N(S)| \geq |S| \quad \forall S \subseteq A \quad (2.1)$$

Proving both implications:

1. \implies : done orally, quite intuitive.
2. \impliedby : we prove the contrapositive²: if A cannot be matched then $\exists S \subseteq A$ such that $|N(S)| < |S|$. So assume A cannot be matched:

$$A \text{ can be matched} \iff |N(S)| \geq |S| \quad \forall S \subseteq A \quad (2.2)$$

$$(2.3)$$

is equivalent to:

$$A \text{ cannot be matched} \implies \exists S \subseteq A : |N(S)| < |S| \quad (2.4)$$

3. If G contains no matching of A , then by König's theorem, it has a cover U with fewer than A vertices:

$$|A \cap U| + |B \cap U| < |A| \quad (2.5)$$

hence:

$$|B \cap U| < |A| - |A \cap U| = |A \setminus (A \cap U)| \quad (2.6)$$

4. By definition of U , G has no edge between $A \setminus (A \cap U)$ and $B \setminus (B \cap U)$. Let $A \setminus (A \cap U)$:

$$|N(S)| \leq |B \cap U| < |A \setminus (A \cap U)| = |S| \quad (2.7)$$

and the marriage condition fails.

$\mu(G) := \mu := |\text{max matching}|$
 $\tau(G) := \tau := |\text{min vertex cover}|$
 $G \text{ bipartite} \Rightarrow \mu(G) = \tau(G)$
 In general $\Rightarrow \mu(G) \leq \tau(G) \leq 2\mu(G)$

2.2 Matching in general graphs

Component (or *connected component*) is a connected subgraph which is inclusion wise maximal with this property.

Observation: every component of G is an induced subgraph of G .

Given a graph G , let us denote by \mathcal{C}_G the set of its components, and by $q(G)$ the number of its odd components, those of odd order.

²The contrapositive of a statement has its antecedent and consequent inverted and flipped: the contrapositive of $P \rightarrow Q$ is thus $\neg Q \rightarrow \neg P$

Tutte's theorem (1947)

A graph G has a perfect matching if and only if $q(G - S) \leq |S|$ for all $S \subseteq V(G)$.

\Rightarrow) Obviously, every odd component of $G - S$ is linked to S by at least one edge from the matching M (because of parity). Hence $|S| \geq q(G - S)$.

\Leftarrow) Let's prove the contrapositive: assume there isn't a perfect matching. We want to show:

$$\exists S \subseteq V : |S| < q(G - S) \quad (2.8)$$

A set such as S will then be said "bad".

We will make two observations:

1. If $|V|$ is odd, then a set $S = \emptyset$ would be bad because $q(G - S) = q(G)$, then at least one of the components of G must be odd since $|V|$ is odd. Hence $q(G - S) = q(G) \geq 1 > |S| = 0$. So we assume $|V|$ is even.
2. If H is a graph and H' is a subgraph obtained by deleting some edge e of H , then $q(H') \geq q(H)$.

Corollary 1 If H' is a spanning subgraph of H then $q(H') \geq q(H)$.

Corollary 2 If S is a bad set for H then S is bad for all spanning subgraphs of H . We have $q(H - S) > |S|$ since S is bad for H . Now if H' is a spanning subgraph of H , then $H' - S$ is a spanning subgraph of $H - S$ and thus by Cor. 1, $q(H' - S) \geq q(H - S)$, thus $q(H' - S) > |S|$.

Trick Let H be obtained from G by iteratively adding edges as long as possible, without creating any perfect matching. By previous observation, if we find a bad set S for H then S is also bad for G . Hence, we may focus on H .

Say that $S \subseteq V$ is "nice". That is, if:

1. Every $v \in S$ is connected to all other vertices from $V - \{v\}$.
2. All components of $H - S$ are complete subgraphs

Observation: if $S \subseteq V$ is nice, then S is bad. In every component, let's cover all but one vertex. For every odd component, connect the vertex we left to S . Indeed, if S were not bad, then we could construct a perfect matching in H . But we know there isn't a perfect matching so S is bad.

So WMA there isn't a nice $S \subseteq V$. We will show that it leads to a contradiction. Let $S := \{v \in V : v \text{ adjacent to all vertices in } V - \{v\}\}$. S is the set of universal vertices (satisfies condition 1). By definition, S satisfies condition 1 from the definition of nice subsets. Thus S must violate condition 2. That is, there is a component X of $G - S$ which is not complete, i.e. $\exists a, a' \in X$ such that $aa' \notin E$. Consider a shortest path

between a and a' , and let a, b, c be the first three vertices on that path. Since $b \notin S$, $\exists d \in V$ such that $bd \notin E$. There is a perfect matching M_1 in $H + ac$, also there's a perfect matching in M_2 in $H + bd$

Note: we know there are at least 3 vertices in this component to have a least 1 connection: components are connected.

Let $H + ac$ have a perfect matching M_1 and $H + bd$ have a perfect matching M_2 . With:

$$M_1 \Delta M_2 = (M_1 - M_2) \cup (M_2 - M_1) \quad (2.9)$$

We get an "alternating" path of pink / blue edges (which are in the symmetric difference), paths alternating between colours, which goes back to the starting point (i.e. a cycle) because the graph is finite and both matchings are perfect. Also, there could be as many cycles as we want.

Now, using this result: remember that ac is an edge of M_1 , as bd is in M_2 , not in the other:

$$bd \in M_1 \Delta M_2 \quad (2.10)$$

Then bd is in a cycle C in $M_1 \Delta M_2$. We have several cases:

1. $ac \notin E(C)$: there's a cycle containing bd . We can switch M_2 on the cycle (i.e. $M_2 \Delta E(C)$): still getting a perfect matching which avoids ac and bd , this is a p.m. in $H \longrightarrow$ contradiction.
2. $ac \in E(C)$: the cycle contains ac . We cannot do the switch because we lose bd and ac appears in M_2 : the new perfect matching... wasn't a perfect matching. The idea is: shortcut c as in bc then we can do the switching and $M_2 \Delta E(C')$ is a p.m. in $H \longrightarrow$ contradiction.

Chapter 3

Connectivity

A graph is *connected* when there is a path between every pair of vertices. In a connected graph, there are no unreachable vertices. A graph that is not connected is disconnected.

A graph with just one vertex is connected. An edgeless graph with two or more vertices is *disconnected*.

A graph G is k -connected if $|V(G)| \geq k + 1$ and $\nexists X \subseteq V$ with $|X| \leq k - 1$ such that $G - X$ is disconnected.

The *vertex connectivity* of G is the maximum k so that G is k -connected.

A vertex $v \in V(G)$ is a cut vertex of a connected graph G if $G - \{v\}$ is not connected.

3.1 2-Connected graphs and subgraphs

There are at least 3 vertices and there isn't any cut vertex.

A *block* of a connected graph G is an induced subgraph B of G such that B is connected and has no cut vertex and B is maximal with this property.

A block B can be:

- a bridge: $B = \{x, y\}$ with $xy \in E(G)$ and $G - xy$ is disconnected.
- a 2-connected induced subgraph: if $|B| \geq 3$, and B is maximal with this property.
- a single vertex: $|B| = 1$, when G is a singleton component.

Let us make an observation: if B_1 and B_2 are two distinct blocks, then they have at most one vertex in common, and if they do then that vertex is a cut vertex of G .

And another: let's make a bipartite graph with, on the one hand, the blocks and on the other hand, the cut vertices. We get a tree.

Theorem: ear decomposition of 2-connected graphs

Let G be 2-connected $\iff G$ can be built starting with a cycle and iteratively adding ears.

Proof:

1. \Leftarrow OK
2. \Rightarrow Let $H \subseteq G$ be a subgraph of G that can be constructed by the procedure and maximal with this property. Note that H exists because G has at least one subgraph that can be constructed, namely a cycle (since G is 2-connected). We want to show $H = G$. Suppose not (pf by contradiction):
 - (a) $V(H) = V(G) \implies \exists e \in E(G) - E(H)$ then we could have added e to H , contradicting the maximality of H .
 - (b) $V(H) \neq V(G)$. We connect a vertex from an ear to a vertex from the cycle H . Since G is connected, $\exists ab \in E(G)$ with $a \in V(H), b \notin V(H)$. Since G is 2-connected, $G - a$ is connected, so there is a path P linking b to $V(H) - \{a\}$ in $G - a$, hence we could have added the ear $P + ab$ to $H \rightarrow$ contradiction.

3.2 The structure of 3-connected graphs**Lemma**

If G is 3-connected and $|V(G)| \geq 5$, then $\exists e \in E(G)$ such that G/e is still 3-connected.

Note that G/e is the graph G following the contraction of the edge e .

Proof: by contradiction, assume there is no such edge.

1. Consider any edge $xy \in E(G)$. Now $H := G/xy$ is not 3-connected, thus there is a subset S of vertices of H with $|S| \leq 2$ so that $H - S$ is not connected.
2. Observation: if v_{xy} denotes the vertex of H resulting from the contraction of xy , then we must have $v_{xy} \in S$, because without otherwise S would also be a cut set of G of size ≤ 3 .
3. On the other hand, if S' is obtained from S by replacing v_{xy} with x, y then S' is a cut set of $G \implies |S'| \geq 3 \implies |S| = 2$.

Any 3-connected graph can be contracted to K_4 by iteratively contracting edges, in such a way that each intermediate state remains 3-connected.

3.3 Menger's theorem

Let $G = (V, E)$ be a graph and $A, B \subseteq V$. Then the minimum number of vertices separating A from B in G is equal to the maximum number of disjoint $A - B$ paths in

G .

That means that G is connected \iff every pair a, b of distinct vertices are linked by k independent paths (no vertex in common except for a and b).

1. Proof \Leftarrow : Take distinct $a, b \in V(G)$. Since there are k independent $a-b$ paths and at least $k-1$ of these have at least one internal vertex, G has $\geq 2+k-1 = k+1$ vertices. Moreover, there is no $Y \subseteq V(G)$ of size less than k such that $G - Y$ is not connected, because o/w: pick a, b in distinct components of $G - Y$. Every $a-b$ path goes through Y , hence there cannot be k independent $a-b$ paths.
2. Proof \Rightarrow : consider two distinct vertices $a, b \in V(G)$.
 - (a) Case 1: $ab \notin E(G)$ Let $A := N(a)$ and $B := N(b)$. If there are k disjoint $A-B$ paths then there are k such paths avoiding a, b and we see that there are k independent paths in G . Observe $|A| \geq k$ and $|B| \geq k$.

If $\nexists k$ disjoint $A-B$ paths then $\exists A-B$ separation Y of size $< k$ but this contradicts the fact that G is k -connected.

- (b) Case 2: $ab \in E(G)$ Let $A := N(a) - \{b\}$ and $B := N(b) - \{a\}$. Then $|A| \geq k-1$ and $|B| \geq k-1$. As always when we have an edge that's annoying, we try to contract it or to delete it: consider $G - ab$. This graph is $(k-1)$ -connected (this always stands for k -connected graphs). And we can find $(k-1)$ disjoint $A-B$ paths in $G - ab$ avoiding a, b . So adding a, b and using the edge ab we get k independent $a-b$ paths in G .

Chapter 4

Planar graphs

A *plane graph* is a pair $G = (V, E)$ where $V \subseteq \mathbb{R}^2$ such that:

1. Every edge is an arc linking two distinct vertices
2. $\forall e \in E$, the interior $\text{int}(e)$ (the arc, minus the two vertices) contains no vertex of G nor any point from edges $f \in E \setminus \{e\}$.
3. Every pair of vertices is linked by at most one edge.

Without the third condition, this could be a plane multi-graph (possibly involving "parallel" edges).

A plane graph defines a corresponding abstract graph G . With a slight abuse of notation, we will write G both for the abstract graph and for the corresponding subset of the plane: $V \cup \bigcup_{e \in E} e \subseteq \mathbb{R}^2$.

The *faces* of G are the regions of $\mathbb{R}^2 \setminus G$. Exactly one of these faces is unbounded and is called the *outer face*. The other faces are called *inner faces*.

All acyclic plane graphs will have only an outer face. As soon as there is a cycle, there will be at least one inner face. In other words: a plane graph has only one face (the outer face) \iff it is a forest.

Lemma 1

Let $F \in E(G)$. Then:

1. $\partial F \subseteq G$
2. $\forall e \in E(G)$, either $e \in \partial F$, or $\partial F \cap \text{int}(e) = \emptyset$
3. If $e \in E(G)$ is on a cycle $C \in G$, then e is contained in the boundary of exactly two faces of G , one inside C and the other outside C .
4. If $e \in E(G)$ is not included in any cycle, then e appears in the boundary of exactly one face.

Lemma 2

If G is a 2-connected plane graph, then the boundary of every face is a cycle of G .

Some other properties:

- A graph is planar if it can be realised as a plane graph.
- A plane graph is *maximally plane* if one cannot add any edge to it (while staying a plane graph).
- A plane graph is a *plane triangulation* if every face is bounded by a triangle of the graph.

Lemma 3

A plane graph on at least 3 vertices is maximally plane \iff it is a plane triangulation

Theorem 1: Euler's formula

Let G be a connected plane graph with n vertices, m edges and f faces. Then:

$$n - m + f = 2 \quad (4.1)$$

Note that this formula works only for connected graphs.

Let's prove it with induction on m .

1. Base case: $m = n - 1$ (as it's connected, it contains a spanning tree, it has $n - 1$ edges). In this case G is a tree and $f = 1$. So $n - m + f = 2$.
2. Inductive step: $m \geq n$. The connected graph is not a tree, then it has a cycle. Consider a cycle of G and let $e \in E(C)$. The edge is on the cycle, so it's on the boundary of exactly two distinct faces of the graph $F_1, F_2 \in F(G)$. If we remove this edge from the graph, the graph stays connected but two faces merge so f decreases by 1. So $f \rightarrow f - 1$, $m \rightarrow m - 1$ and n remains the same. Then $-m + f$ is constant and $n - m + f = 2$.

Corollary 1

If G is an n -vertex plane graph with n being at least 3, then G has at most $3n - 6$ edges.

With my assumption WLOG that G is maximally plane, i.e. G is a plane triangulation. Let $m := |E(G)|$ and $f := |F(G)|$.

$$3f = \sum_{F \in F(G)} 3 = 2m \quad (4.2)$$

If we look at every face in the triangulation, each edge will be counted two times as it is boundary to two faces. Hence:

$$m = \frac{3}{2}f \quad (4.3)$$

And by Euler's formula (which we will use as we're considering a plane triangulation):

$$n - m + f = 2 \quad (4.4)$$

Thus

$$n - m + \frac{2}{3}m = 2 \quad (4.5)$$

$$\implies n - \frac{1}{3}m = 2 \quad (4.6)$$

$$\implies m = 3n - 6 \quad (4.7)$$

Corollary 2

K_5 is not planar.

K_5 has $n = 5$ vertices and $m = 10$ edges but $3n - 6 = 9 < 10$.

Lemma 3

If G is an n -vertex plane graph ($n \geq 3$) and has no triangle (i.e cycle of length 3 or 3 vertices that are pairwise adjacent), then G has at most $2n - 4$ edges.

Corollary 3

$K_{3,3}$ is not planar.

$K_{3,3}$ has $n = 6$ vertices and $m = 9$ edges: $2n - 4 = 8 < 9$.

Corollary 4

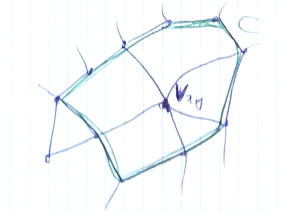
No planar graph contains K_5 or $K_{3,3}$ as a minor.

Theorem 2: Kuratowski (1930)

A graph G is planar $\iff G$ it has neither K_5 nor $K_{3,3}$ as minor.

1. \implies : already done.
2. \impliedby : We first prove the statement in the case that G is 3-connected. That is, we will show: " G 3-connected and has neither K_5 or $K_{3,3}$ as minor $\implies G$ is planar", by reduction on $|V(G)|$.
 - (a) Base case: $|V(G)| = 4$. Then $G = K_4$ and is planar (every k -connected graph has at least $k + 1$ vertices).

- (b) Inductive step: $|V(G)| \geq 5$. By Tutte's theorem, there is an edge $xy \in E(G)$ such that when we contract (we get $G \setminus xy$), the graph is 3-connected. Note that $G \setminus xy$ has no K_5 nor $K_{3,3}$ minor (by transitivity of the minor relation), thus $G \setminus xy$ is planar by induction. Let v_{xy} be the vertex resulting from the edge contraction. Consider a drawing of $G \setminus xy$ in the plane. The plane graph $(G \setminus xy) - v_{xy}$ is 2-connected. In particular, each face is bounded by a cycle of that graph. Let C denote the cycle bounding the face where v_{xy} was drawn.

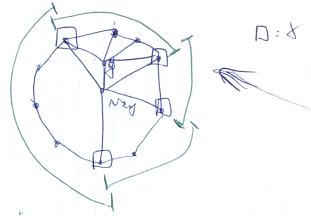


Then all neighbours of v_{xy} in $G \setminus xy$ are on C . Let:

$$X := \{v \in V(C) : vx \in E(G)\} \quad (4.8)$$

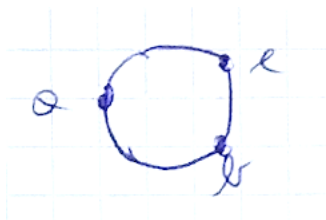
$$Y := \{v \in V(C) : vy \in E(G)\} \quad (4.9)$$

Note that $|X| \geq 2$ and $|Y| \geq 2$, since G is 3-connected.

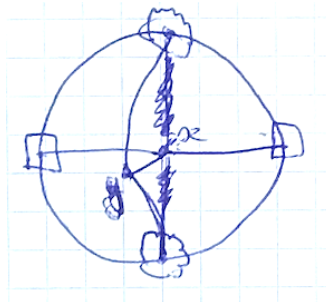


Consider the plane graph $H := (G \setminus xy) - \{vy : v \in Y - X\}$. Note that all the remaining edges are at least adjacent to X .

- i. Vertices of X define intervals of C . One can see H as a plane drawing of $G - y$ (where v_{xy} plays the role of x). Our aim is to add y to obtain a drawing of G .
- ii. If there exists a single interval containing all of Y then we can extend this drawing to a drawing of G by adding y in the face made by the vertices of Y .
- iii. Now let's assume this is not the case. We will show that this leads to a contradiction (and this Y must be contained in a single interval).
 - A. $|X \cap Y| \geq 3$, then there is a K_5 minor in G .
 Let $a, b, c \in X \cap Y$ be distinct. Consider the minor of G obtained by contracting each of these three paths into edges.
 Then in that minor $\{a, b, c, x, y\}$ are pairwise linked (a, b, c are a triangle, x and y are both connected to all a, b, c). Hence K_5 is a minor of G . Contradiction.



B. $|X \cap Y| \leq 2$, then there is a $K_{3,3}$ minor in G :



Note this usual approach in graph theory: assume the graph has some connectivity, prove the theorem, then extend

Chapter 5

Colouring

A vertex colouring of a graph $G = (V, E)$ is a map $\phi : V \rightarrow S$ such that $\phi(u) \neq \phi(v) \quad \forall uv \in E(G)$.

If ϕ is a k -coloring, $|S| \leq k$. Also, we define the *chromatic number* χ as:

$$\chi(G) := \min\{k : \exists k\text{-colouring of } G\} \quad (5.1)$$

A graph G with $\chi_G = k$ is called k -chromatic; if $\chi_G \leq q$, we call G k -colourable.

Note that k -colouring is nothing but a vertex partition into k independent sets, now called *colour classes*. The non-trivial 2-colourable graphs, for example, are precisely the bipartite graphs.

5.1 Colouring maps and planar graphs

The four colour conjecture (1852)

If a graph G is planar, then $\chi_G \leq 4$.

This has been "proved" by Appel and Haken in 1977 using 2000 graph configurations (resulting in 200 pages). but none has really accepted their proof.

Then it has been reproved by Robertson, Sanders, Seymour, Thomas (1997, 600 configurations, 160 pages), who tried to understand the first proof but couldn't make sense of it.

In 2005, Gonthiers (Microsoft research) made a complete formal proof so that it can be verified by computers, which it has.

The five colour theorem (Kempe and Heawood, 1890)

If G is planar, $\chi_G \leq 5$.

Proof: by induction on $|V(G)|$, valid if $|V(G)| \leq 5$ so let's assume $|V(G)| \geq 6$.

1. The average degree of G is:

$$d(G) = \frac{\sum_{v \in V(G)} d(v)}{|V(G)|} = \frac{2|E(G)|}{|V(G)|} \leq \frac{2(3|V(G)| - 6)}{|V(G)|} < 6 \quad (5.2)$$

Now if this average is strictly less than 6, there is at least one vertex v with $d(v) \leq 5$.

2. If $d(v) = 4$, then consider a 5-colouring of $G - v$, one colour is available at v , done.
3. If $d(v) = 5$, showing that χ_G could be 6 is now quite straightforward, think of removing v and then putting it back in G). Now we have a potential problem if we consider a 5-colouring: there must be two distinct non-adjacent vertices x, y in the neighbours of v such that $xy \notin E(G)$. We will first contract xv and yv (thus keeping the planar property, so we can apply induction). By induction, $G/\{vx, vy\}$ has a 5-colouring ϕ' . Let:

$$\phi(w) := \phi'(w) \quad \forall w \in V(G) - \{vx, vy\} \quad (5.3)$$

$$\phi(x) := \phi(v) := \phi(y) = \phi'(\text{contracted vertex}) \quad (5.4)$$

ϕ is not a valid colouring because of the edges vx, vy . However, these are the only two problematic edges. Moreover, since $\phi(x) = \phi(y)$, there is some colour c that is not used on any neighbour of v , thus setting $\phi(v) := c$ makes ϕ a valid 5-colouring.

Discharging method (Weinicke, 1904)

One example to prove above's theorem. If G is a plane triangulation and G has minimum degree, then $\exists vw \in E(G)$ such that $d(v) = 5$ and $d(w) \in \{5, 6\}$.

Put a charge of $6 - d(v)$ on each v in G with:

$$\sum_{v \in V(G)} \text{charge}(v) = 6|V(G)| - \sum_{v \in V(G)} d(v) = 6|V(G)| - 2|E(G)| \quad (5.5)$$

$$= 6|V(G)| - 2(3|V(G)| - 6) \quad (5.6)$$

$$= 12 > 0 \quad (5.7)$$

As we're in a plane triangulation, we can use Euler's formula.

One discharging rule: every degree-5 vertex sends a charge of $1/5$ to each of its neighbours.

Let $\text{charge}'(v)$ denote the new charge of v .

$$\sum_{v \in V(G)} \text{charge}'(v) = \sum_{v \in V(G)} \text{charge}(v) = 12 > 0 \quad (5.8)$$

Let $v \in V(G)$ be such that $\text{charge}'(v) > 0$. Now we need to do a simple case analysis:

1. $d(v) \in \{5, 6\}$. Then v received some charge from some neighbour w , which thus has degree 5 and hence vw is an edge as desired.
2. $d(v) \geq 8$. This cannot happen because:

$$\text{charge}'(v) \leq \text{charge}(v) + d(v)/5 \quad (5.9)$$

$$= 6 - d(v) + d(v)/5 \quad (5.10)$$

$$= 6 - \frac{4}{5}d(v) \quad (5.11)$$

$$\leq 6 - \frac{4 \cdot 8}{5} \leq 0 \quad (5.12)$$

3. $d(v) = 7$. By the same calculation, v receives charges from at least 6 from its neighbours. With the condition that the new charge must be strictly positive, we can find two neighbours of v that have degree 5 and are adjacent.

A *clique* is a set of pairwise adjacent vertices. A clique of size 3 is a triangle. We have:

$$w(G) := \text{maximum size of a clique in } G \quad (5.13)$$

Clearly, $\chi(G) \geq w(G) \forall G$. For example, a cycle of 5 vertices has a clique of size 2 but requires 3 colours. Formally, we have $\chi(C_{2k+1}) = 3$ but $w(C_{2k+1}) = 2$, for $k \geq 2$.

In fact, there are triangle-free graphs (i.e. $w \leq 2$) with arbitrarily high χ .

The *girth* of a graph G is the length of a shortest cycle contained in G . Its value is considered infinite if G has no cycle.

Theorem (Erdo's, 1950)

$$\forall k \geq 1, \exists G : \chi(G) \geq k \wedge \text{girth}(G) \geq k \quad (5.14)$$

One can always 2-colour a tree. If locally, up until some extent, the graph is a tree but contains a cycle globally, the cycle will enforce the value of χ .

The chromatic number is a global phenomenon.

So the clique size is a lower bound, but at times this can be a very bad bound. For bipartite graphs, we have equality, but that's not always so.

5.2 Perfect graphs

A graph G is *perfect* if $\chi(H) = w(H)$ for all induced subgraphs H of G . For example, bipartite graphs are perfect.

Chapter 6

Random graphs

6.1 The notion of a random graph

Let V be a fixed set of n elements, say $V = \{0, \dots, n-1\}$

Theorem

Every graph G has a bipartite subgraph H with $|E(H)| \geq |E(G)|/2$.

Let H be a bipartite subgraph of G with bipartition A, B defined as follows:

$$\forall v \in V(G), \begin{cases} P(v \in A) = 1/2 \\ P(v \in B) = 1/2 \end{cases} \quad (6.1)$$

(independently at random).

Then $E(H)$ consists of all the edges of G with one endpoint in A and the other in B .
 $\forall e \in E(G)$, define X_e as the random variable with:

$$X_e = \begin{cases} 1 & \text{if } e \in E(H) \\ 0 & \text{otherwise} \end{cases} \quad (6.2)$$

Let $X = |E(H)|$. Then:

$$E(X) = E\left(\sum_{e \in E(G)} X_e\right) = \sum_{e \in E(G)} E(X_e) = \sum_{e \in E(G)} \frac{2}{4} = \frac{|E(G)|}{2} \quad (6.3)$$

Then in particular, $\exists H$ with $|E(H)| \geq |E(G)|/2$.

Quick remark on the number of edges in a bipartite graph: $K_{a,b}$, number of edges ab , maximal when $a = b = n/2$, then $K_{n/2, n/2}$ has $n^2/4$ edges.

Theorem (Erdős's, 1956)

$\forall k \geq 0, \exists \mathcal{G}$ with $\text{girth}(G) > k$ and $\chi(G) > k$.

The idea is to use the Erdős's-Renyi model of random graphs $\mathcal{G}(n, p)$. There are n vertices and for every vertex pair we put an edge at probability p , independently at random.

Lemma

For $G \in \mathcal{G}(n, p)$, let X_k be the random variable counting the number of cycles of length k in G . Then:

$$E(X_k) = \frac{n!}{2k(n-k)!} p^k \quad (6.4)$$

Proof:

1. The number of k -cycles in K_n is (considering that both orientations $/2$ and any shifting leads to the same cycle $(/k)$):

$$\binom{n}{k} k! \cdot \frac{1}{2k} = \frac{n!}{2k(n-k)!} \quad (6.5)$$

2. For C a k -cycle in K_n , we let X_C be the independent random variable encoding whether $C \subset \mathcal{G}$. Then:

$$E(X_k) = E\left(\sum_C X_C\right) = \sum_C E(X_C) = \sum_C p^k = \frac{n!}{2k(n-k)!} p^k \quad (6.6)$$

Lemma

Let $n \geq k \geq 3$ and $p \in [0, 1]$ with $p \geq 6k \ln n$. Then, for $G \in \mathcal{G}(n, p)$:

$$P(\alpha(G)) \geq \frac{n}{2k} \rightarrow 0 \quad (6.7)$$

First proof: let $t := \text{ceil}(n/2k)$. We have:

$$P(\alpha(G) \geq t) \stackrel{\text{union bound}}{\leq} \binom{n}{t} (1-p)^{t/2} \quad (6.8)$$

$$\leq n^t e^{-pt(t-1)/2} \quad (6.9)$$

$$= (ne^{-pt/2} e^{p/2})^t \quad (6.10)$$

If that last expression goes towards zero, clearly it goes quicker to zero as t grows. So we write:

$$(ne^{-pt/2} e^{p/2})^t \leq (ne^{-\frac{6k \ln n}{n} \frac{n}{2k} \frac{1}{2} e^{p/2}})^t \quad (6.11)$$

$$= (ne^{-\frac{3}{2} \ln n} e^{p/2})^t \quad (6.12)$$

And $t \geq 1/2$. So:

$$\underbrace{(n^{-\frac{1}{2}})}_{\rightarrow 0} \underbrace{e^{p/2}}_{\in [1, e]} t \quad (6.13)$$

clearly goes to 0.

Hence $P(\alpha(G) \geq n/2k) \xrightarrow{n \rightarrow \infty} 0$.

Proof of the theorem: we want to show $\exists G$ with $\text{girth}(G) > k$ and $\chi(G) > k$. Let $\varepsilon > 0$ be fixed with $\varepsilon < 1/k$. In the rest of the proof, we take n "large enough" compared to k and ε . Let $p := \frac{n^\varepsilon}{n}$.

For $G \in \mathcal{G}(n, p)$, let $X_{\leq k}$ count the number of cycles of length $\leq k$. Then $X_{\leq k} = \sum_{i=3}^k X_i$. Now:

$$E(X_{\leq k}) = \sum_{i=3}^k E(X_i) = \sum_{i=3}^k \frac{n!}{2i(n-i)!p^i} \leq \frac{1}{2} \sum_{i=3}^k n^i p^i \leq \frac{1}{2} (k-2)(np)^k \quad (6.14)$$

For this that inequality, we take n large enough to ensure $np = n^\varepsilon \geq 1$. Finally, we have:

$$\frac{1}{2} (k-2)(np)^k \leq \frac{k}{2} n^{\varepsilon k} \quad (6.15)$$

And $\varepsilon k < 1$.

Quick recap on Markov's inequality If X is a random variable with $X \geq 0$ and $a > 0$, then $P(x \geq a) \leq \frac{E(X)}{a}$. Proof (discrete case): say X can take values in $\{v_1, v_2, \dots, v_n\}$ in increasing order and $v_1 \geq 0$. Then $E(X) = \sum_{i \geq 1} v_i P(x = v_i)$. But this cannot have any negative value:

$$\sum_{i \geq 1} v_i P(x = v_i) \geq \sum_{i \geq 1, v_i \geq 0} v_i P(x = v_i) \quad (6.16)$$

$$\geq \sum_{i \geq 1, v_i \geq 0} a P(x = v_i) \quad (6.17)$$

$$= a \sum_{i \geq 1, v_i \geq a} P(x = v_i) \quad (6.18)$$

$$= P(x \geq a) \quad (6.19)$$

So $E(X) \geq aP(x \geq a)$.

Back to the proof. By Markov's inequality:

$$P(X_{\leq k} \geq n/2) \leq \frac{E(X_{\leq k})}{n/2} \quad (6.20)$$

$$\leq k n^{\varepsilon k - 1} \quad (6.21)$$

As $\varepsilon k - 1 < 0$, this whole last expression tends to 0. Also, since $p = \frac{n^\varepsilon}{n} \geq \frac{6k \ln n}{n}$, we can use previous lemma and get:

$$P(\alpha(G) \geq \frac{n}{2k}) \xrightarrow{n \rightarrow \infty} 0. \quad (6.22)$$

. Take n large enough so that:

$$P(X_{\leq k} \geq \frac{n}{2}) < 1/2, \quad P(\alpha(G) \geq \frac{n}{2k} < 1/2 \quad (6.23)$$

By the union bound:

$$P(X_{\leq k} \geq \frac{n}{2} | \alpha(G) \geq \frac{1}{2k}) < \frac{1}{2} + \frac{1}{2} < 1 \quad (6.24)$$

So $\exists G$ with $< n/2$ cycles of length $\leq k$ and $\alpha(G) < \frac{n}{2k}$. For each cycle of length in G , choose a vertex on the cycle. Let X be the set of chosen vertices and let $H := G - S$.

Then $\text{girth}(H) > k$. Also:

$$\chi(H) \geq \frac{|V(H)|}{\alpha(H)} \geq \frac{|V(G)| - |S|}{\alpha(H)} \geq \frac{n/2}{\alpha(H)} \quad (6.25)$$

But we notice something about the stability number of H : it is at most the stability number of G . By dividing by $\alpha(G)$, we get another lower bound:

$$\geq \frac{n/2}{\alpha(G)} \quad (6.26)$$

$$> \frac{n/2}{n/(2k)} \quad (6.27)$$

$$= k \quad (6.28)$$

So the chromatic number of H is strictly more than k , and so is its girth.

First moment method

- Markov's inequality to bind to probability of a RV non negative, smaller than its expectation divided by some constant
- Union bound

Local lemma

A_1, \dots, A_n events in a probability space "bad events"

Union bound

$$P(A \cup \dots \cup A_n) \leq \sum_i P(A_i) \quad (6.29)$$

So if $\sum_i P(A_i) < 1$, we have that $P(\bar{A}_1, \dots, \bar{A}_n) > 0$. Note that this is usually a weak bound.

A_i 's independent: if $P(A_i) < 1 \forall i$: then $P(\cap \bar{A}_i) = \prod_i P(\bar{A}_i) > 0$.

With local dependencies: A is independent from B_1, \dots, B_k if :

$$\forall J \subseteq \{1, \dots, k\} : P(A \cap \bigcap_{j \in J} B_j) = P(A)P(\bigcap_{j \in J} B_j) \quad (6.30)$$

Remark: it is not enough to just check pairs. For $S = \{1, 2, 3, 4, 5, 6\}$ and $A(1, 2, 3, 4), B_1(2, 4, 5), B_2(3, 4, 5)$, we have:

$$P(A) = 2/3, P(B_1) = P(B_2) = 1/2, P(A \cap B_1) = 1/2, P(A \cap B_2) = 1/3 \quad (6.31)$$

$$\implies P(A \cap B_1 \cap B_2) = 1/6 \neq P(A)P(B_1 \cap B_2) = 2/3 \cdot 1/3 = 2/9 \quad (6.32)$$

Say A_1, \dots, A_n are events. A dependency digraph on A_1, \dots, A_n is a digraph with vertex set $\{1, \dots, n\}$ such that A_i is independent from the collection $\{A_j : j \neq i, (i, j) \notin E(D)\}$, $\forall i$.

Lovász local lemma (symmetric version)

Given events A_1, \dots, A_n and a dependency digraph with Δ , outdegree d , if $P(A_i) \leq p \forall i$ and $eP(D+1) \leq 1$, then $P(\bigcap_i \bar{A}_i) > 0$.

Lovász local lemma (asymmetric version)

Given events A_1, \dots, A_n and a dependency digraph D . Suppose that $\exists x_1, \dots, x_n \in [0, 1]$ such that $P(A_i) \leq \prod_{(i,j) \in E(D)} (1 - x_j) \forall i$. Then $P(\bigcap_i \bar{A}_i) \geq \prod_i (1 - x_i) > 0$.

Hypergraph 2-coloring

Note that a hypergraph is a set system. Hypergraph H is k -uniform if each edge has size k .

The *2-coloring* of H is a coloring of its vertices such that no edge is monochromatic.

Theorem

If H is a k -uniform hypergraph and each edge intersects $\leq 2^{k-1}/2$ others, then \exists 2-coloring of H .

Proof:

- Color each vertex red (1/2) or blue (1/2) independently at random
- Bad events: for each $e \in E(H) : A_e := "$ e is monochromatic $"$.
- $P(A_e) = 2 \cdot \frac{1}{2^k} = 2^{-(k-1)}$
- Dependency d -graph: put an edge (e, f) whenever $e \cap f \neq \emptyset \forall e, f \in E(H)$

By the LLL: since D has max outdegree $\leq \frac{2^{k-1}}{e} - 1$ and $P(A_e) = 2^{-(k-1)} \forall e \in E(H)$ and $e2^{-(k-1)} \left(\frac{2^{k-1}}{e} \right) \leq 1$, $P(\bigcap_{e \in E(H)} \bar{A}_e) > 0$ and thus \exists 2-coloring.

Constructive proof

in case of 2-colorings of hypergraphs.

Moser 08-09:

- Start with an arbitrary 2-coloring (all red for instance)
- For each $e \in E(H)$: if e is monochromatic: $\text{fix}(e)$

$\text{fix}(e)$ is defined as:

- Resample the colors of all vertices in e
- For each $f \in E(H)$ with $f \cap e \neq \emptyset$: if f is monochromatic, $\text{fix}(f)$

Analysis (not for exam)

Consider a call of $\text{fix}(e)$ (ie this edge was not chromatic). We get a recursive tree.