

# BU CS320 Assignment 6: Context Free Grammars

November 6, 2023

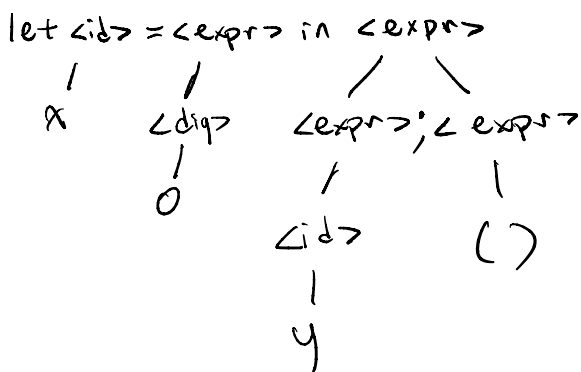
- Given the following grammar where  $\langle expr \rangle$  is the starting symbol:

$\langle id \rangle$	$::= a \mid b \mid c \mid \dots \mid z$
$\langle dig \rangle$	$::= 0 \mid 1 \mid 2 \mid \dots \mid 9$
$\langle expr \rangle$	$::= () \mid \langle dig \rangle \mid \langle id \rangle$ $\quad \mid \text{let } \langle id \rangle = \langle expr \rangle \text{ in } \langle expr \rangle$ $\quad \mid \langle expr \rangle ; \langle expr \rangle$ $\quad \mid \text{begin } \langle expr \rangle \text{ end}$

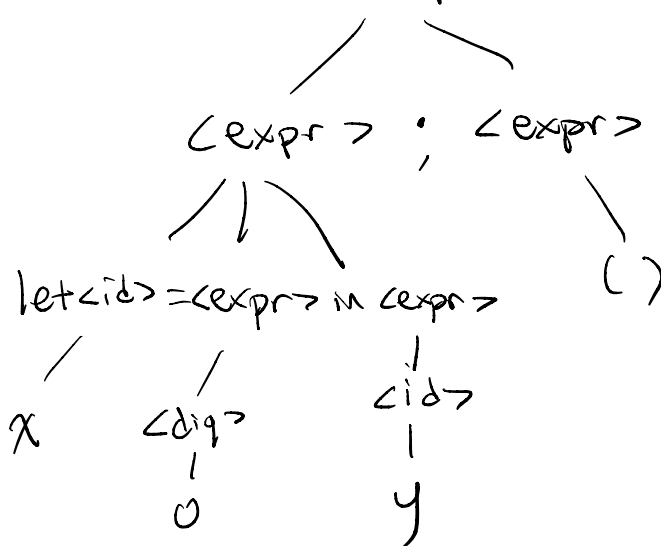
Demonstrate the grammar above is ambiguous.

let  $x = 0$  in  $y; ()$

$\langle expr \rangle$



$\langle expr \rangle$



2 parse trees to evaluate  
expression:

let  $x = 0$  in  $y; ()$

2. Modify the grammar (reproduced below) to be unambiguous. Hint: There is not just one way.

```

<id> ::= a | b | c | ... | z
<dig> ::= 0 | 1 | 2 | ... | 9
<expr> ::= () | <dig> | <id>
          | let <id> = <expr> in <expr>
          | <expr> ; <expr>
          | begin <expr> end

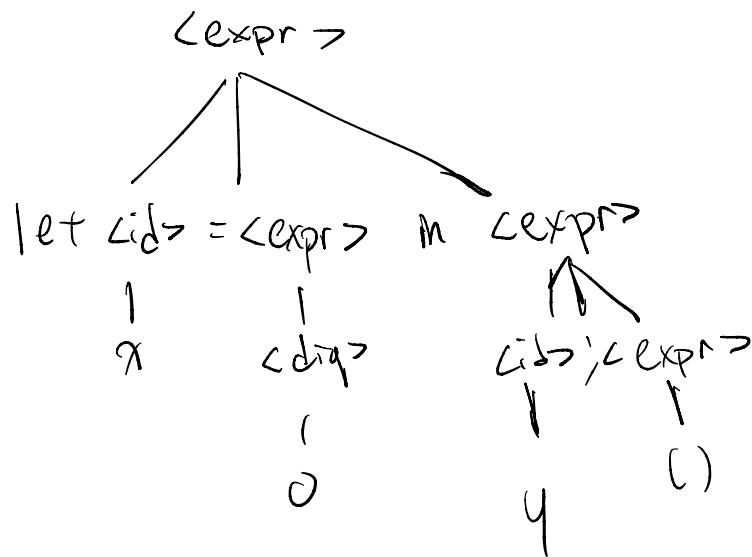
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$\langle id \rangle ::= a | b | c | \dots | z$   
 $\langle dig \rangle ::= 0 | 1 | 2 | \dots | 9$   
 $\langle expr \rangle ::= () | \langle dig \rangle | \langle id \rangle$   
 $\quad | \text{let } \langle id \rangle = \langle expr \rangle \text{ in } \langle expr \rangle$   
 $\quad | \text{begin } \langle expr \rangle \text{ end}$   
 $\quad | ( ) ; \langle expr \rangle$   
 $\quad | \langle dig \rangle ; \langle expr \rangle$   
 $\quad | \langle id \rangle ; \langle expr \rangle$   
 $\quad | \text{begin } \langle expr \rangle \text{ end} ; \langle expr \rangle$

making the  
grammar left associative

3. Demonstrate your modified grammar fixes the previously shown ambiguity.

let  $x = 0$  in  $y$ ;  $()$



Only 1 way to construct the expression

let  $x = 0$  in  $y$ ;  $()$

No more ambiguity.