last time

course intro / logistics

building C programs (cc = clang/gcc, etc.)

`cc -c file.c` (makes `file.o`) — compile+assemble
reads `file.c` and all the files it #includes

`cc file1.o file2.o ... -o executable` — link
reads .o files + some system files

`...-Lpath -lname` — libraries:
static (`libname.a`): included in executable itself
dynamic (`libname.so`): found + loaded at program start — one copy on system

runtime search paths for dynamic libraries

anonymous feedback (1)

Holding class in Mcleod is very hard for most of this class as I know at least DMT2, which a lot of us are in , is all the way in gilmer and is a more than 15 min to get here. I ask to please consider holding class somewhere else on grounds or starting class a few minutes later every day so that everyone has ample time to get here and be prepared to learn.

I'm pretty sure there's not an alternative room (I didn't volunteer for a long walk from Rice...)

disappointed to lose some lecture time, ...

quiz demo

anonymous feedback (2)

I was hoping you could do some introductions to some concepts before diving into the slides. I think it would help clarify things for us before we learn new content due to the large gap we have had since last talking about this subject. Also if you could continue some in class exercises and add examples to the slides that would be very helpful.

Is there a way we could get a C refresher. Are there any good resources to learn memory allocation etc? We did not have good practice with that in cso1

warmup assignment

C exercise

```
int array[4] = {10,20,30,40};
int *p;
p = &array[0];
p += 2;
p[1] += 1;
```

array =

A. compile or runtime error
B. {10,20,30,41}
C. {10,20,32,41}
D. {10,21,30,40}
E. {12,21,30,40}
F. none of these

some avenues for review

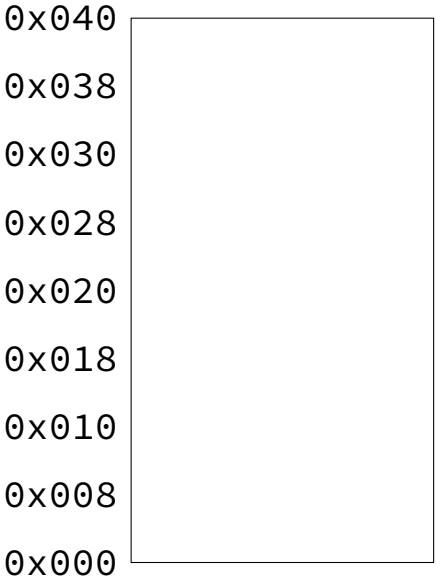
review CSO1 stuff

- labs 9–12 (of last Spring)
  `https://www.cs.virginia.edu/~jh2jf/courses/cs2130/spring2023/`

exercises we've used in the past:

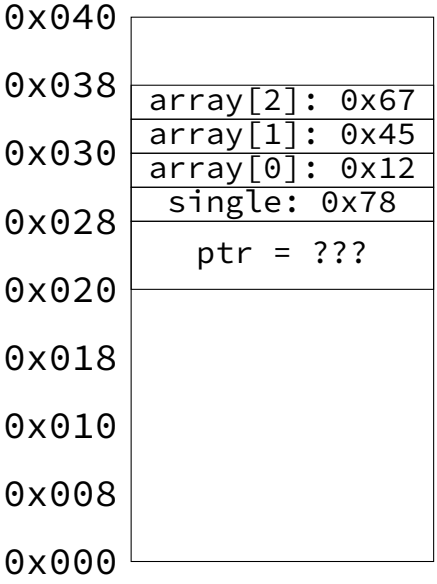
- implement strsep library function
- implement conversion from dynamic array to linked list

some pointer stuff



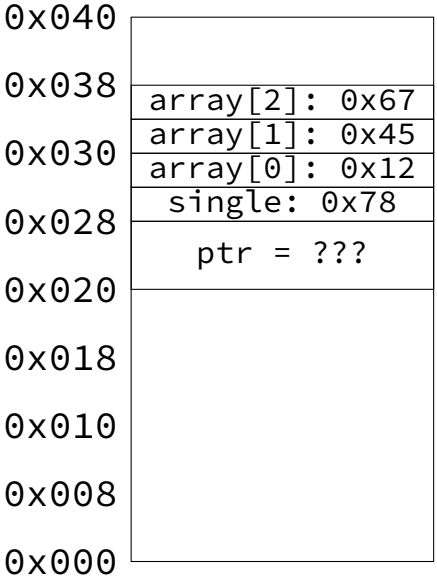

```
int array[3]={0x12,0x45,0x67};
int single = 0x78;
int *ptr;
```

some pointer stuff




```
int array[3]={0x12,0x45,0x67};
int single = 0x78;
int *ptr;
```

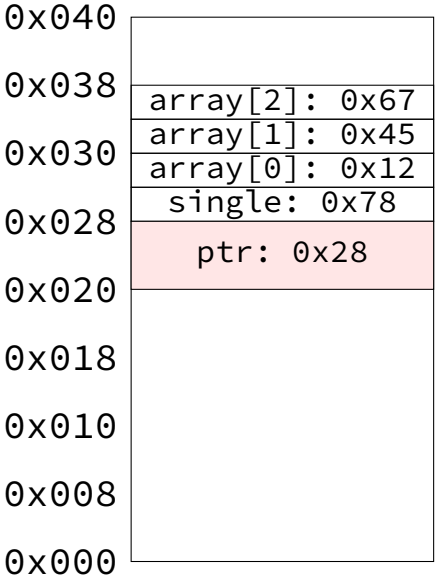
some pointer stuff




```
int array[3]={0x12,0x45,0x67};
int single = 0x78;
int *ptr;
```

~~`*ptr = 0xAB;`~~ compile error

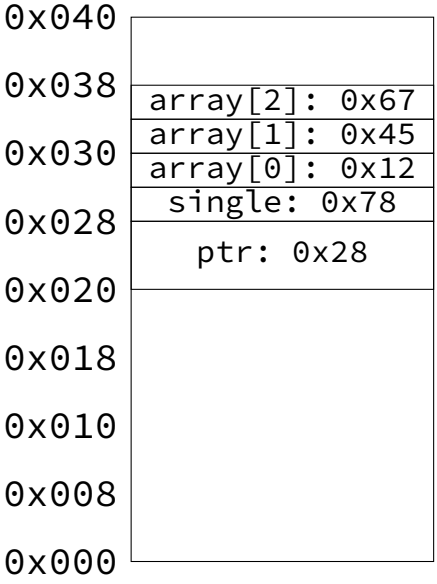
some pointer stuff




```
int array[3]={0x12,0x45,0x67};
int single = 0x78;
int *ptr;

ptr = &single;
ptr = (int*) 0x28;
```

addr. of single

some pointer stuff



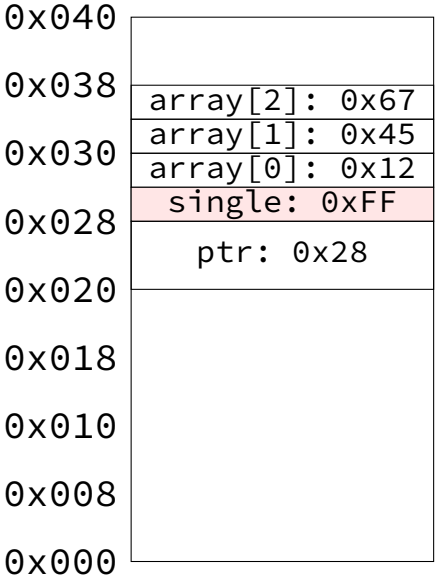

```
int array[3]={0x12,0x45,0x67};
int single = 0x78;
int *ptr;

ptr = &single;
ptr = (int*) 0x28;
```
addr. of single

~~`ptr = 0x28;`~~ compile error

~~`ptr = (int*) single;`~~

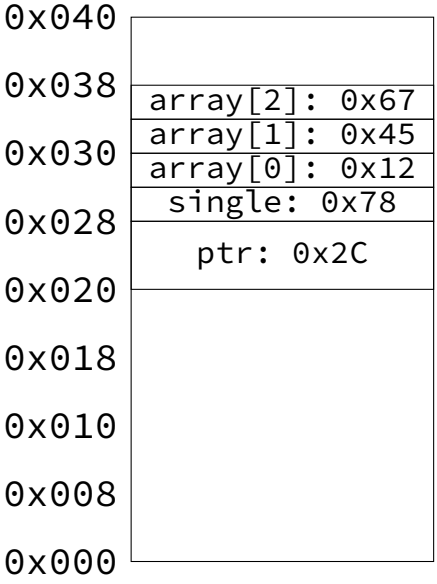
pointer to unknown place

some pointer stuff



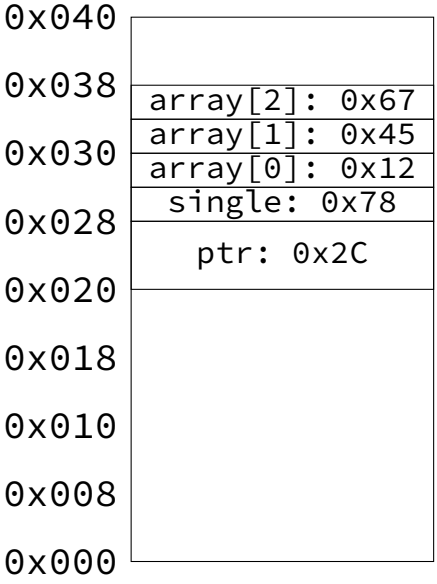

```
int array[3]={0x12,0x45,0x67};
int single = 0x78;
int *ptr;
ptr = &single;

*ptr = 0xFF;
```

some pointer stuff




```
int array[3]={0x12,0x45,0x67};
int single = 0x78;
int *ptr;

ptr = array;
ptr = &array[0];
ptr = (int*) 0x2C;
```

some pointer stuff



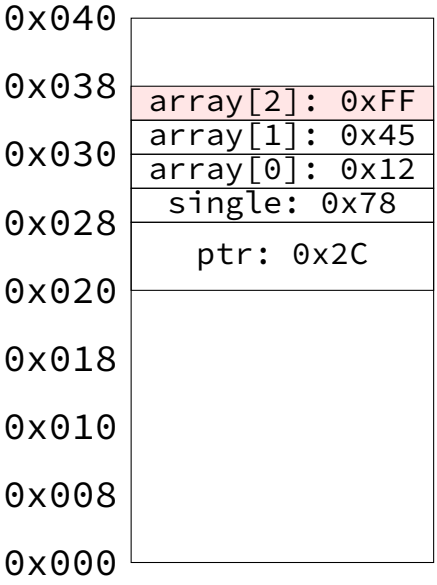

```
int array[3]={0x12,0x45,0x67};
int single = 0x78;
int *ptr;

ptr = array;
ptr = &array[0];
ptr = (int*) 0x2C;
```

~~`ptr = array[0];`~~ compile error

~~`ptr = (int*) array[0];`~~

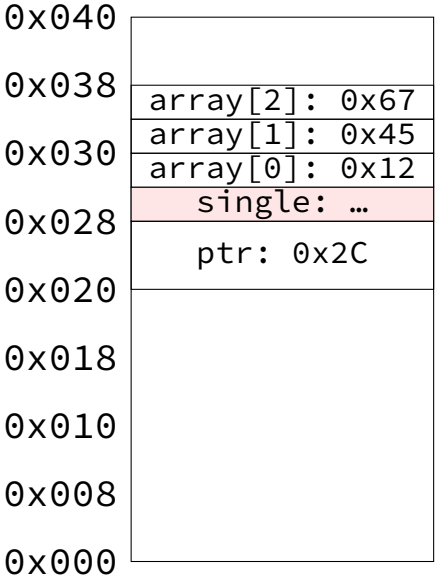
pointer to unknown place

some pointer stuff




```
int array[3]={0x12,0x45,0x67};
int single = 0x78;
int *ptr;
ptr = &array[0];

ptr[2] = 0xFF;
*(ptr + 2) = 0xFF;

int *temp1; temp1 = ptr + 2;
*temp1 = 0xFF;

int *temp2; temp2 = &ptr[2];
*temp2 = 0xFF;
```

some pointer stuff




```
int array[3]={0x12,0x45,0x67};
int single = 0x78;
int *ptr;

void change_arg(int *x) {
    *x = compute_some_value();
}
…
change_arg(&single);
```

make

`make` — Unix program for "making" things...

...by running commands based on what's changed

what commands? based on *rules* in *makefile*

make rules

```
main.o: main.c main.h extra.h
▶           clang -c main.c
```

before colon: target(s) (file(s) generated/updated)

after colon: prerequisite(s)

following lines prefixed by a tab character: command(s) to run

`make` runs commands if any prereq modified date after target

make rules

```
main.o: main.c main.h extra.h
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make rules

```
main.o: main.c main.h extra.h
▶       clang -c main.c
```

before colon: target(s) (file(s) generated/updated)

after colon: prerequisite(s)

following lines prefixed by a tab character: command(s) to run

`make` runs commands if any prereq modified date after target

...after making sure prerequisites up to date

make rule chains

```
program: main.o extra.o
▶           clang -o program main.o extra.o

extra.o: extra.c extra.h
▶           clang -c extra.c

main.o: main.c main.h extra.h
▶           clang -c main.c
```

to *make* `program`, first...

update main.o and extra.o if they aren't

running make

"`make` *`target`*"

look in `Makefile` in current directory for rules
check if *`target`* is up-to-date
if not, rebuild it (and dependencies, if needed) so it is

"`make` *`target1 target2`*"

check if both target1 and target2 are up-to-date
if not, rebuild it as needed so they are

"`make`"

if "*`firstTarget`*" is the first rule in Makefile,
same as '`make` *`firstTarget`*"

exercise: what will run?

```
W: X Y
►    buildW
X: Q
►    buildX
Y: X Z
►    buildY
```

W	modified 1 minute ago
X	modified 3 hours ago
Y	does not exist
Z	modified 1 hour ago
Q	modified 2 hours ago

exercise: "make W" will run what commands?

A. none
B. `buildY` only
C. `buildW` then `buildY`
D. `buildY` then `buildW`
E. `buildX` then `buildY` then `buildW`
F. `buildX` then `buildW`
G. something else

‘phony’ targets (1)

common to have Makefile targets that aren't files

```
all: program1 program2 libfoo.a
```

“`make all`” effectively shorthand for “`make program1 program2 libfoo.a`”

no actual file called “all”

‘phony’ targets (2)

sometimes want targets that don’t actually build file

example: “`make clean`” to remove generated files

```
clean:
▶	rm --force main.o extra.o
```

but what if I create...

```
clean:
▶	rm --force main.o extra.o

all: program1 program2 libfoo.a
```

Q: if I make a file called “all” and then “make all” what happens?

Q: same with “clean” and “make clean”?

marking phony targets

```
clean:
▶       rm --force main.o extra.o

all: program1 program2 libfoo.a

.PHONY: all clean
```

special .PHONY rule says " 'all' and 'clean' not real files"

(not required by POSIX, but in every make version I know)

conventional targets

common convention:

target name	purpose
(default), `all`	build everything
`install`	install to standard location
`test`	run tests
`clean`	remove generated files

redundancy (1)

```
program: main.o extra.o
▶       clang -o program main.o extra.o

extra.o: extra.c extra.h
▶       clang -o extra.o -c extra.c
main.o: main.c main.h extra.h

▶       clang -o main.o -c main.c
```

what if I want to run `clang` with `-Wall`?

what if I want to change to `gcc`?

variables/macros (1)

```
CC = gcc
CFLAGS = -Wall -pedantic -std=c11 -fsanitize=address
LDFLAGS = -Wall -pedantic -fsanitize=address
LDLIBS = -lm

program: main.o extra.o
▸       $(CC) $(LDFLAGS) -o program main.o extra.o $(LDLIBS)

extra.o: extra.c extra.h
▸       $(CC) $(CFLAGS) -o extra.o -c extra.c

main.o: main.c main.h extra.h
▸       $(CC) $(CFLAGS) -o main.o -c main.c
```

variables/macros (2)

```
CC = gcc
CFLAGS = -Wall
LDFLAGS = -Wall
LDLIBS = -lm

program: main.o extra.o
▶       $(CC) $(LDFLAGS) -o $@ $^ $(LDLIBS)

extra.o: extra.c extra.h
▶       $(CC) $(CFLAGS) -o $@ -c $<

main.o: main.c main.h extra.h
▶       $(CC) $(CFLAGS) -o $@ -c $<
```

aside: `$^` works on GNU make (usual on Linux), but not portable.

suffix rules

```
CC = gcc
CFLAGS = -Wall
LDFLAGS = -Wall

program: main.o extra.o
▶	$(CC) $(LDFLAGS) -o $@ $^

.c.o:
▶	$(CC) $(CFLAGS) -o $@ -c $<

extra.o: extra.c extra.h
main.o: main.c main.h extra.h
```

aside: `$^` works on GNU make (usual on Linux), but not portable.

pattern rules

```
CC = gcc
CFLAGS = -Wall
LDFLAGS = -Wall
LDLIBS = -lm

program: main.o extra.o
▶       $(CC) $(LDFLAGS) -o $@ $^ $(LDLIBS)

%.o: %.c
▶       $(CC) $(CFLAGS) -o $@ -c $<

extra.o: extra.c extra.h
main.o: main.c main.h extra.h
```

aside: these rules work on GNU make (usual on Linux), but less portable than suffix rules.

built-in rules

'make' has the 'make .o from .c' rule built-in already, so:

```
CC = gcc
CFLAGS = -Wall
LDFLAGS = -Wall
LDLIBS = -lm

program: main.o extra.o
▶       $(CC) $(LDFLAGS) -o $@ $^ $(LDLIBS)

extra.o: extra.c extra.h
main.o: main.c main.h extra.h
```

(don't actually need to write supplied rule!)

writing Makefiles?

error-prone to automatically all .h dependencies

`-M` option to `gcc` or `clang`

- outputs Make rule
- ways of having make run this

Makefile generators

- other programs that write Makefiles

other build systems

alternatives to writing Makefiles:

other make-ish build systems

- ninja, scons, bazel, maven, xcodebuild, msbuild, ...

tools that generate inputs for make-ish build systems

- cmake, autotools, qmake, ...

things programs on portal shouldn't do

read other user's files

modify OS's memory

read other user's data in memory

hang the entire system

things programs on portal shouldn't do

read other user's files

modify OS's memory

read other user's data in memory

hang the entire system

privileged operation: problem

how can hardware (HW) plus operating system (OS) allow:

read your own files from hard drive

but disallow:

read others files from hard drive

some ideas

OS tells HW ‘okay’ parts of hard drive before running program code

- complex for hardware and for OS

some ideas

OS tells HW 'okay' parts of hard drive before running program code

- complex for hardware and for OS

OS verifies your program's code can't do bad hard drive access

- no work for HW, but complex for OS
- may require compiling differently to allow analysis

some ideas

OS tells HW 'okay' parts of hard drive before running program code

- complex for hardware and for OS

OS verifies your program's code can't do bad hard drive access

- no work for HW, but complex for OS
- may require compiling differently to allow analysis

OS tells HW to only allow OS-written code to access hard drive

- that code can enforce only 'good' accesses
- requires program code to call OS routines to access hard drive
- relatively simple for hardware

kernel mode

extra one-bit register: "are we in *kernel mode*"

other names: privileged mode, supervisor mode, ...

not in kernel mode = *user mode*

certain operations only allowed in kernel mode

privileged instructions

example: talking to any I/O device

what runs in kernel mode?

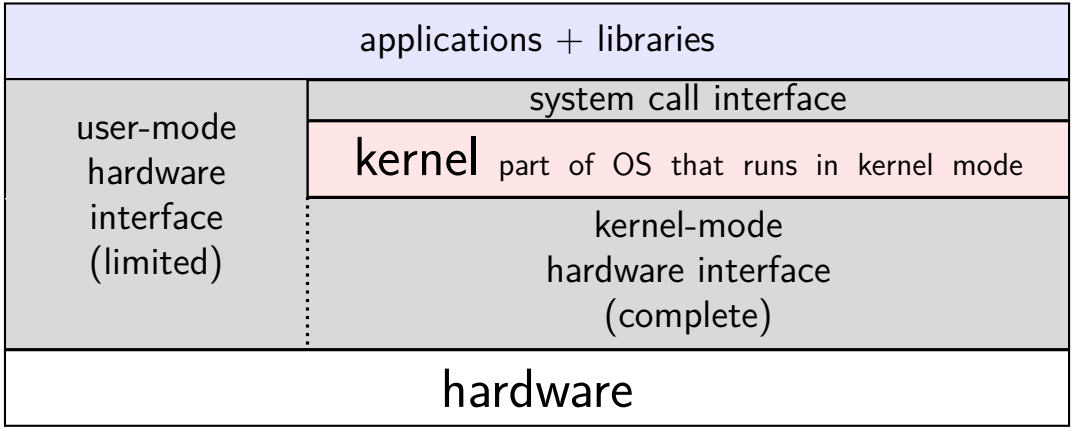
system boots in kernel mode

OS switches to user mode to run program code

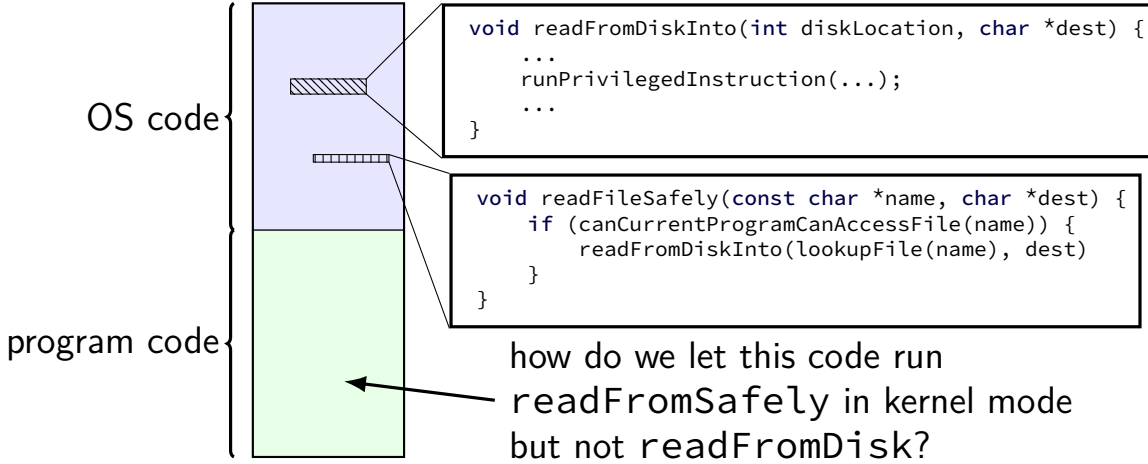
next topic: when does system switch back to kernel mode?

- how does OS tell HW where the (trusted) OS code is?

hardware + system call interface



calling the OS?



controlled entry to kernel mode (1)

special instruction: “system call”

runs OS code in kernel mode at location specified earlier

OS sets up at boot

location can’t be changed without privilieged instrution

controlled entry to kernel mode (2)

OS needs to make specified location:

figure out what operation the program wants

- calling convention, similar to function arguments + return value

be "safe" — not allow the program to do 'bad' things

- example: checks whether current program is allowed to read file before reading it
- requires exceptional care — program can try weird things

Linux x86-64 system calls

special instruction: `syscall`

runs OS specified code in kernel mode

Linux syscall calling convention

before `syscall`:

`%rax` — system call number

`%rdi`, `%rsi`, `%rdx`, `%r10`, `%r8`, `%r9` — args

after `syscall`:

`%rax` — return value

on error: `%rax` contains -1 times "error number"

almost the same as normal function calls

Linux x86-64 hello world

```
.globl _start
.data
hello_str: .asciz "Hello, World!\n"
.text
_start:
  movq $1, %rax # 1 = "write"
  movq $1, %rdi # file descriptor 1 = stdout
  movq $hello_str, %rsi
  movq $15, %rdx # 15 = strlen("Hello, World!\n")
  syscall

  movq $60, %rax # 60 = exit
  movq $0, %rdi
  syscall
```

approx. system call handler

```
sys_call_table:
    .quad handle_read_syscall
    .quad handle_write_syscall
    // ...

handle_syscall:
    ... // save old PC, etc.
    pushq %rcx // save registers
    pushq %rdi
    ...
    call *sys_call_table(,%rax,8)
    ...
    popq %rdi
    popq %rcx
    return_from_exception
```

Linux system call examples

`mmap`, `brk` — allocate memory

`fork` — create new process

`execve` — run a program in the current process

`_exit` — terminate a process

`open`, `read`, `write` — access files

`socket`, `accept`, `getpeername` — socket-related

system call wrappers

library functions to not write assembly:

```
open:
    movq $2, %rax // 2 = sys_open
    // 2 arguments happen to use same registers
    syscall
    // return value in %eax
    cmp $0, %rax
    jl has_error
    ret
has_error:
    neg %rax
    movq %rax, errno
    movq $-1, %rax
    ret
```

system call wrappers

library functions to not write assembly:

```
open:
    movq $2, %rax // 2 = sys_open
    // 2 arguments happen to use same registers
    syscall
    // return value in %eax
    cmp $0, %rax
    jl has_error
    ret
has_error:
    neg %rax
    movq %rax, errno
    movq $-1, %rax
    ret
```

system call wrapper: usage

```
/* unistd.h contains definitions of:
    O_RDONLY (integer constant), open() */
#include <unistd.h>
int main(void) {
  int file_descriptor;
  file_descriptor = open("input.txt", O_RDONLY);
  if (file_descriptor < 0) {
      printf("error: %s\n", strerror(errno));
      exit(1);
  }
  ...
  result = read(file_descriptor, ...);
  ...
}
```

system call wrapper: usage

```
/* unistd.h contains definitions of:
    O_RDONLY (integer constant), open() */
#include <unistd.h>
int main(void) {
  int file_descriptor;
  file_descriptor = open("input.txt", O_RDONLY);
  if (file_descriptor < 0) {
      printf("error: %s\n", strerror(errno));
      exit(1);
  }
  ...
  result = read(file_descriptor, ...);
  ...
}
```

strace hello_world (1)

`strace` — Linux tool to trace system calls

run on assembly program we saw earlier:

```
$ strace -o trace.txt ./hello_world
$ cat trace.txt
execve("./hello_world", ["./hello_world"],
        0x7ffeedafd0a0 /* 28 vars */) = 0
write(1, "Hello, World!\n\0", 14)       = 14
exit(0)                                 = ?
+++ exited with 0 +++
```

strace hello_world (2)

```
#include <stdio.h>
int main() { puts("Hello, World!"); }
```

---

when statically linked:

```
execve("./hello_world", ["./hello_world"], 0x7ffeb4127f70 /* 28 vars */)
                                        = 0
brk(NULL)                               = 0x22f8000
brk(0x22f91c0)                          = 0x22f91c0
arch_prctl(ARCH_SET_FS, 0x22f8880)      = 0
uname({sysname="Linux", nodename="reiss-t3620", ...}) = 0
readlink("/proc/self/exe", "/u/cr4bd/spring2023/cs3130/slide"..., 4096)
                                        = 57
brk(0x231a1c0)                          = 0x231a1c0
brk(0x231b000)                          = 0x231b000
access("/etc/ld.so.nohwcap", F_OK)      = -1 ENOENT (No such file or
                                                     directory)
fstat(1, {st_mode=S_IFCHR|0620, st_rdev=makedev(136, 4), ...}) = 0
write(1, "Hello, World!\n", 14)         = 14
exit_group(0)                           = ?
+++ exited with 0 +++
```

aside: what are those syscalls?

execve: run program

brk: allocate heap space

arch_prctl(ARCH_SET_FS, ...): thread local storage pointer

may make more sense when we cover concurrency/parallelism later

uname: get system information

readlink of /proc/self/exe: get name of this program

access: can we access this file [in this case, a config file]?

fstat: get information about open file

exit_group: variant of exit

strace hello_world (2)

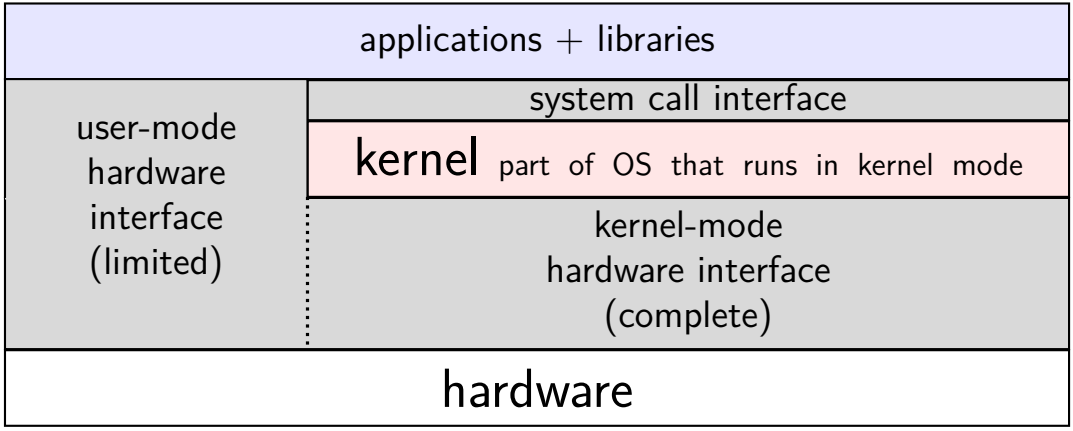
```
#include <stdio.h>
int main() { puts("Hello, World!"); }
```

---

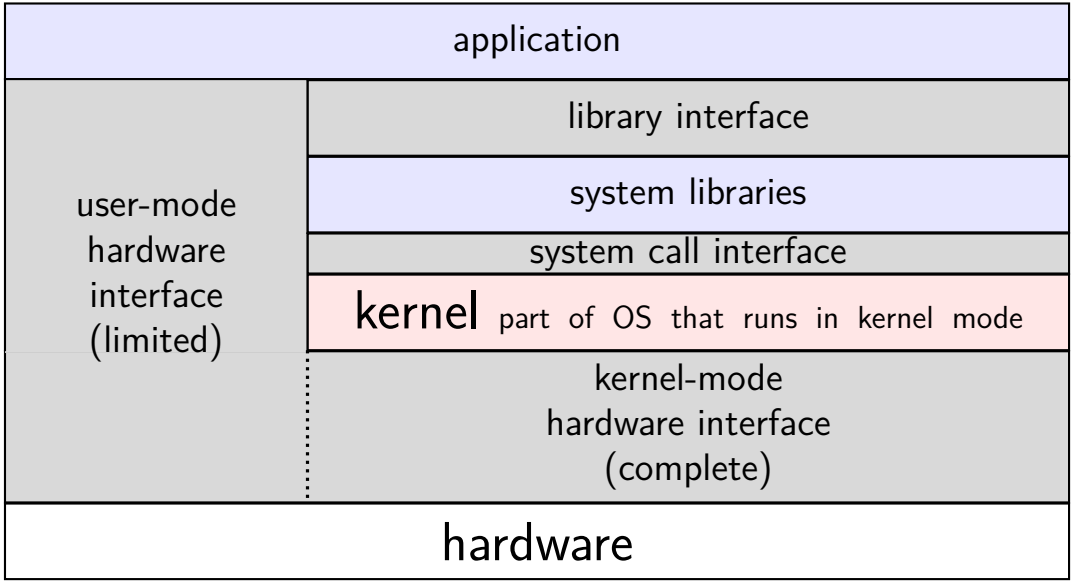
when dynamically linked:

```
execve("./hello_world", ["./hello_world"], 0x7ffcfe91d540 /* 28 vars */)
                                        = 0
brk(NULL)                               = 0x55d6c351b000
...
openat(AT_FDCWD, "/etc/ld.so.cache", O_RDONLY|O_CLOEXEC) = 3
fstat(3, {st_mode=S_IFREG|0644, st_size=196684, ...}) = 0
mmap(NULL, 196684, PROT_READ, MAP_PRIVATE, 3, 0) = 0x7f7a62dd3000
close(3)                                = 0
access("/etc/ld.so.nohwcap", F_OK)      = -1 ENOENT (No such file or directory
openat(AT_FDCWD, "/lib/x86_64-linux-gnu/libc.so.6", O_RDONLY|O_CLOEXEC) = 3
read(3, "\177ELF\2\1\1\3\0\0\0\0\0\0\0\0\3\0>\0\1\0\0\0"..., 832) = 832
...
close(3)                                = 0
write(1, "Hello, World!\n", 14)         = 14
exit_group(0)                           = ?
+++ exited with 0 +++
```

hardware + system call interface



hardware + system call + library interface



things programs on portal shouldn't do

read other user's files

modify OS's memory

read other user's data in memory

hang the entire system

memory protection

modifying another program's memory?

Program A	Program B

`0x10000: .long 42`
`// ...`
`// do work`
`// ...`
`movq 0x10000, %rax`
	`// while A is working:`
	`movq $99, %rax`
	`movq %rax, 0x10000`
	`...`

memory protection

modifying another program's memory?

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`// ...`
`movq 0x10000, %rax`
	`// while A is working:`
	`movq $99, %rax`
	`movq %rax, 0x10000`
	`...`

result: `%rax` (in A) is ...

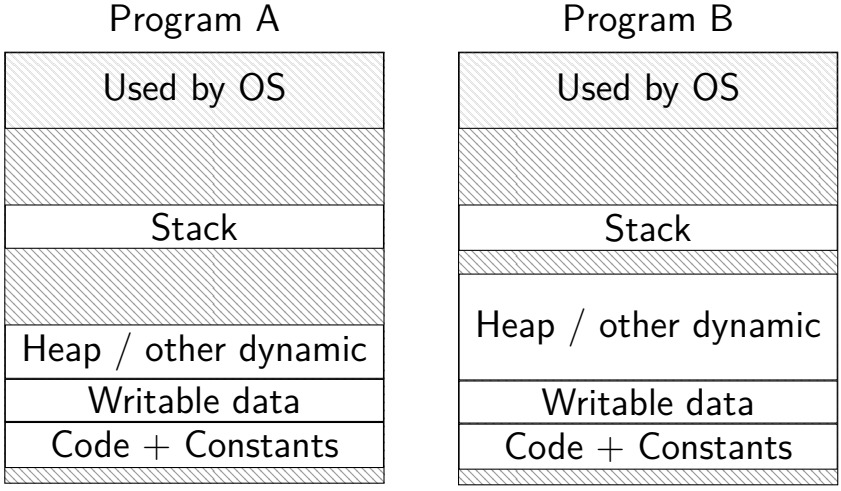
A. 42 B. 99 C. 0x10000

D. 42 or 99 (depending on timing/program layout/etc)

E. 42 or 99 or program might crash (depending on ...)

F. something else

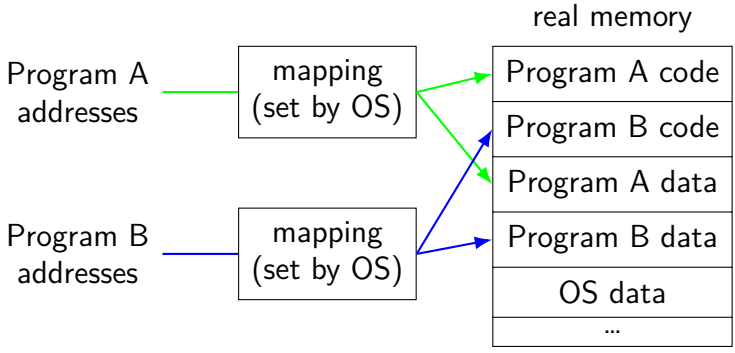
program memory (two programs)



address space

programs have illusion of own memory

called a program's address space



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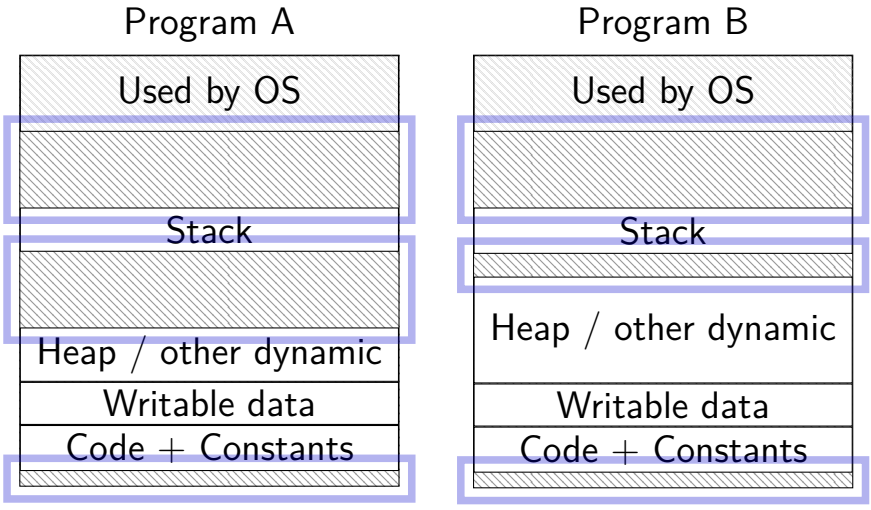

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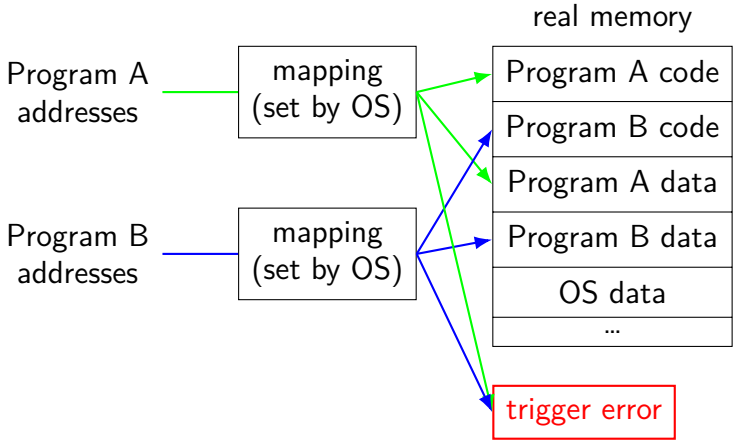


address space mechanisms

topic after exceptions

called virtual memory

mapping called page tables

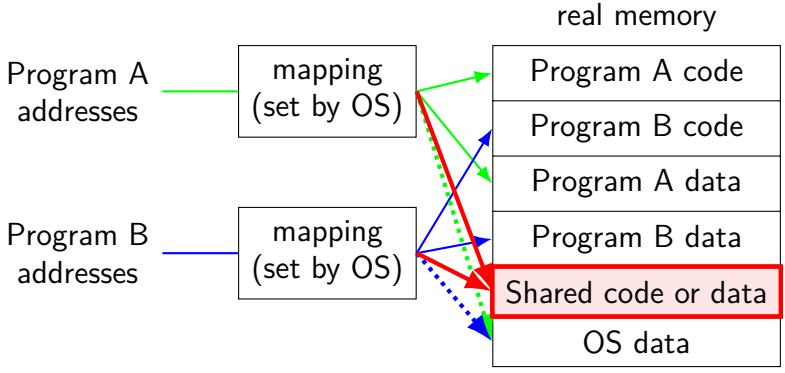
mapping part of what is changed in context switch

shared memory

recall: dynamically linked libraries

would be nice not to duplicate code/data...

we can!



one way to set shared memory on Linux

```
/* regular file, OR: */
int fd = open("/tmp/somefile.dat", O_RDWR);
/* special in-memory file */
int fd = shm_open("/name", O_RDWR);
...
/* make file's data accessible as memory */
void *memory = mmap(NULL, size, PROT_READ | PROT_WRITE,
                    MAP_SHARED, fd, 0);
```

mmap: "map" a file's data into your memory

will discuss a bit more when we talk about virtual memory

part of how Linux loads dynamically linked libraries

memory protection

modifying another program's memory?

Program A	Program B

`0x10000: .long 42`
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`// ...`
`movq 0x10000, %rax`
	`// while A is working:`
	`movq $99, %rax`
	`movq %rax, 0x10000`
	`...`

result: `%rax` (in A) is 42 (always)	result: might crash

A. 42 B. 99 C. 0x10000

D. 42 or 99 (depending on timing/program layout/etc)

E. 42 or 99 or program might crash (depending on ...)

F. something else

program crashing?

what happens on processor when program crashes?

other program informed of crash to display message

use processor to run some other program

program crashing?

what happens on processor when program crashes?

other program informed of crash to display message

use processor to run some other program

how does hardware do this?

would be complicated to tell about other programs, etc.

instead: hardware runs designated OS routine

exceptions

recall: system calls — software asks OS for help

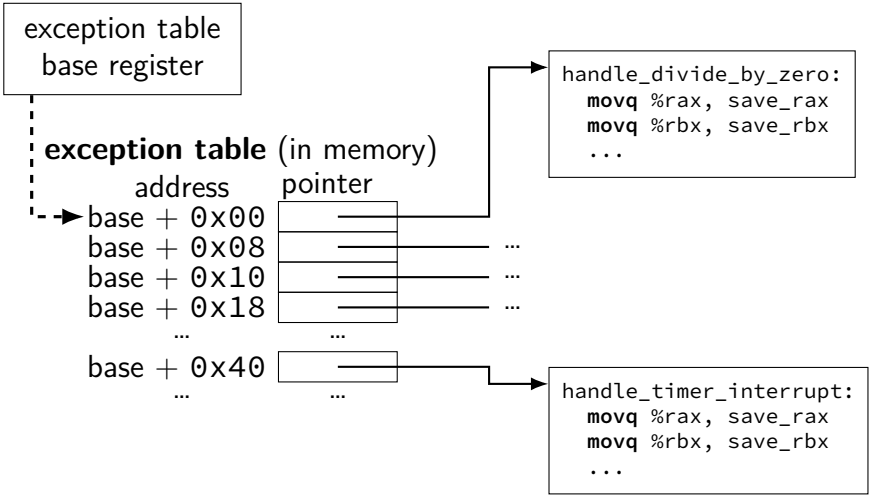
also cases where hardware asks OS for help

different triggers than system calls

but same mechanism as system calls:

- switch to kernel mode (if not already)
- call OS-designated function

locating exception handlers (one strategy)



running the exception handler

hardware saves the old program counter (and maybe more)

identifies location of exception handler via table

then jumps to that location

OS code can save anything else it wants to , etc.

types of exceptions

system calls

- intentional — ask OS to do something

errors/events in programs

- memory not in address space ("Segmentation fault")
- privileged instruction
- divide by zero, invalid instruction
- ...

(and more we'll talk about later)

types of exceptions

system calls

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(and more we'll talk about later)

synchronous
triggered by current program

things programs on portal shouldn't do

read other user's files

modify OS's memory

read other user's data in memory

hang the entire system

an infinite loop

```
int main(void) {
    while (1) {
        /* waste CPU time */
    }
}
```

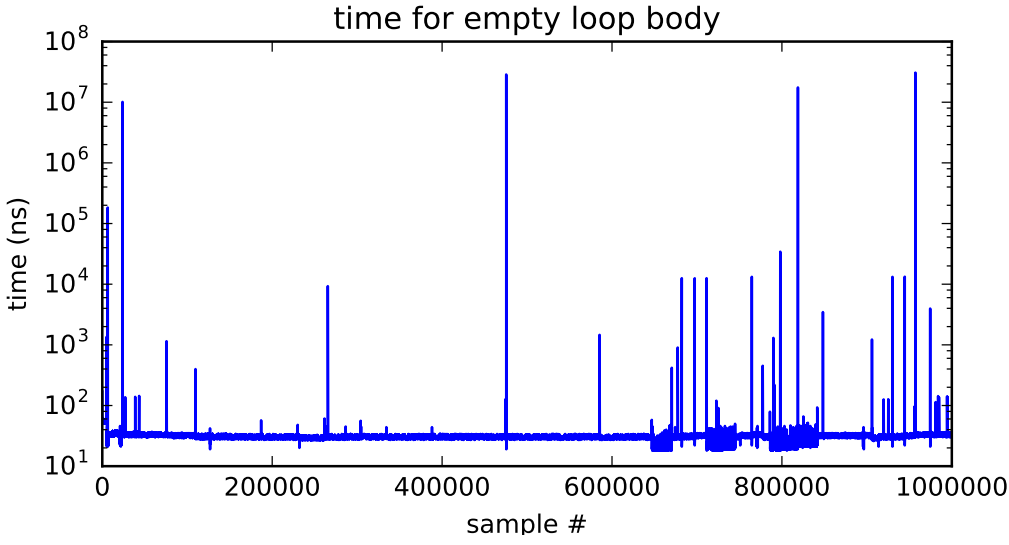
If I run this on a shared department machine, can you still use it?
...if the machine only has one core?

timing nothing

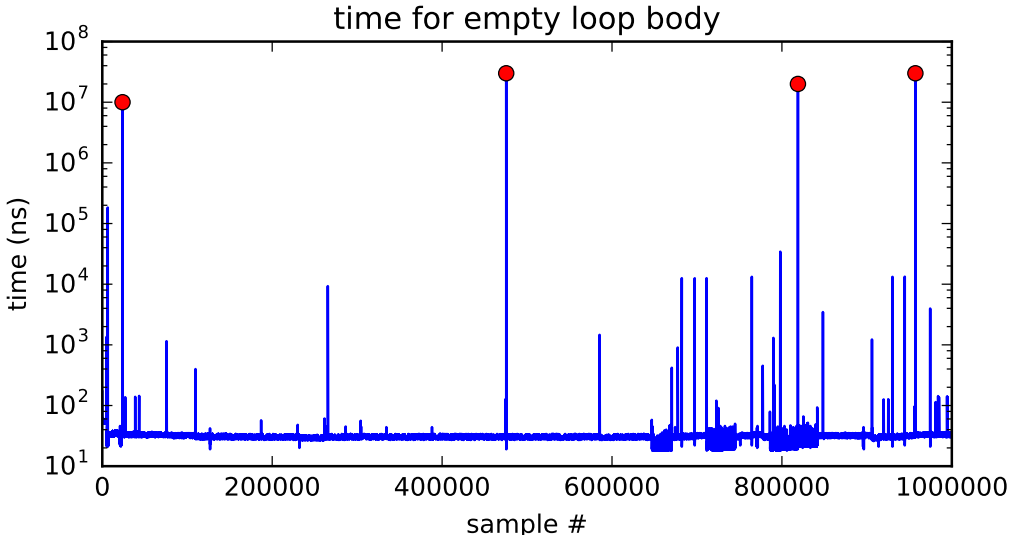
```
long times[NUM_TIMINGS];
int main(void) {
    for (int i = 0; i < N; ++i) {
        long start, end;
        start = get_time();
        /* do nothing */
        end = get_time();
        times[i] = end - start;
    }
    output_timings(times);
}
```

same instructions — same difference each time?

doing nothing on a busy system



doing nothing on a busy system



types of exceptions

system calls

intentional — ask OS to do something

errors/events in programs

memory not in address space ("Segmentation fault")
privileged instruction
divide by zero, invalid instruction
...

synchronous
triggered by current program

external — I/O, etc.

timer — configured by OS to run OS at certain time
I/O devices — key presses, hard drives, networks, ...
hardware is broken (e.g. memory parity error)

asynchronous
not triggered by running program

time multiplexing

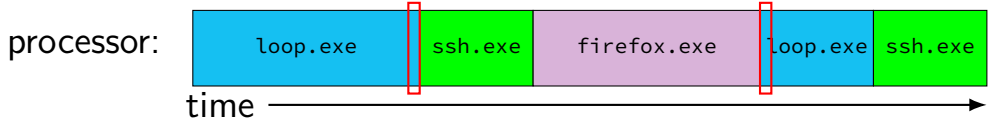
time multiplexing


```
...
call get_time
    // whatever get_time does
movq %rax, %rbp
```

million cycle delay

```
call get_time
    // whatever get_time does
subq %rbp, %rax
...
```

time multiplexing

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    // whatever get_time does
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```

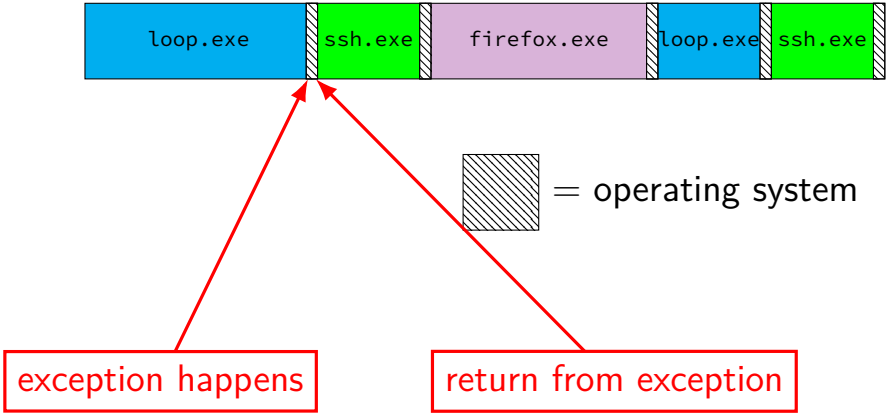
million cycle delay

```
call get_time
    // whatever get_time does
subq %rbp, %rax
...
```

time multiplexing really

 = operating system

time multiplexing really



types of exceptions

system calls

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} synchronous
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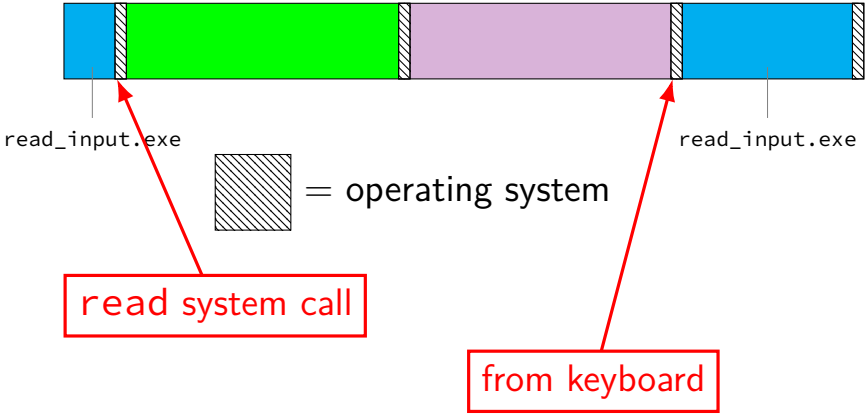
external — I/O, etc.

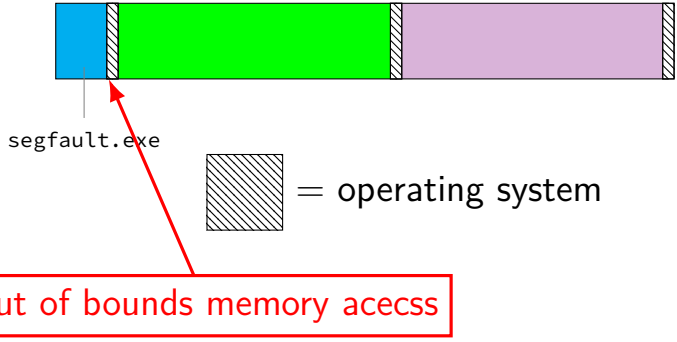
timer — configured by OS to run OS at certain time
I/O devices — key presses, hard drives, networks, ...
hardware is broken (e.g. memory parity error)

} asynchronous
not triggered by running program

keyboard input timeline



crash timeline timeline



threads

thread = illusion of own processor

own register values

own program counter value

threads

thread = illusion of own processor

own register values

own program counter value

actual implementation:
many threads sharing one processor

> problem: where are register/program counter values
> when thread not active on processor?

switching programs

OS starts running somehow

- some sort of exception

saves old registers + program counter

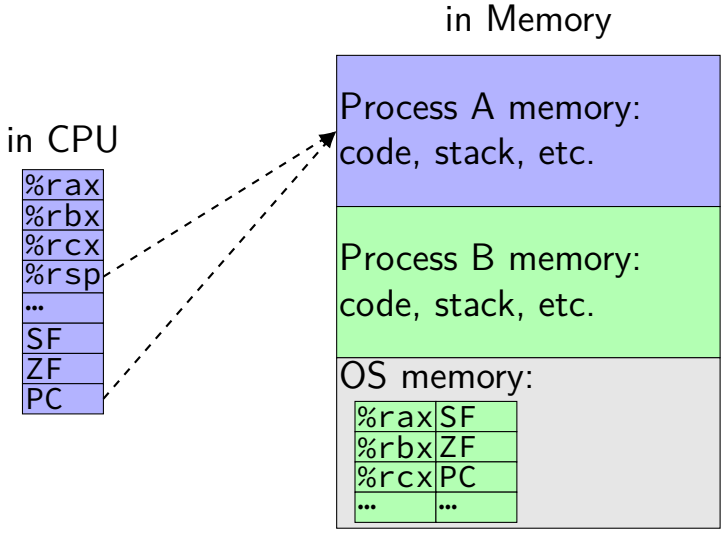
- (optimization: could omit when program crashing/exiting)

sets new registers, jumps to new program counter

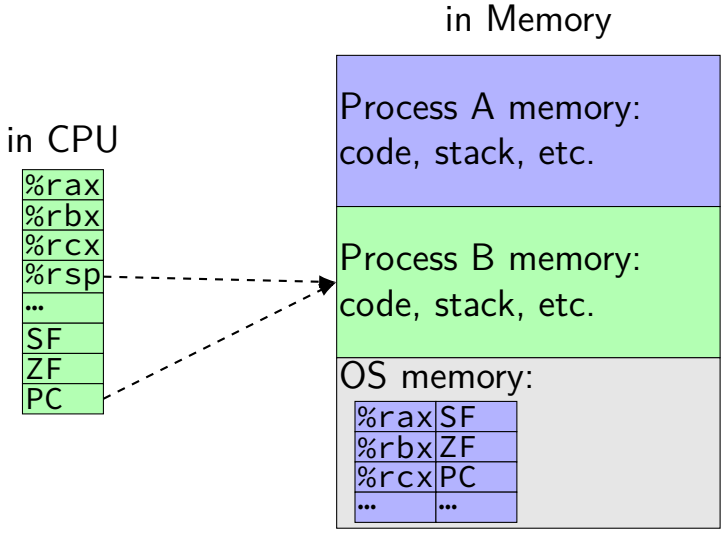
called context switch

- saved information called context

contexts (A running)



contexts (B running)



review: definitions

exception: hardware calls OS specified routine

- many possible reasons
- system calls: type of exception

context switch: OS switches to another thread

- by saving old register values + loading new ones
- part of OS routine run by exception

which of these require exceptions? context switches?

A. program calls a function in the standard library

B. program writes a file to disk

C. program A goes to sleep, letting program B run

D. program exits

E. program returns from one function to another function

F. program pops a value from the stack

terms for exceptions

terms for exceptions aren't standardized

our readings use one set of terms

interrupts = externally-triggered
faults = error/event in program
trap = intentionally triggered

all these terms appear differently elsewhere

The Process

process = thread(s) + address space

illusion of dedicated machine:

- thread = illusion of own CPU
- address space = illusion of own memory