

COMPUTER SYSTEMS AND ORGANIZATION

Bitwise Operations

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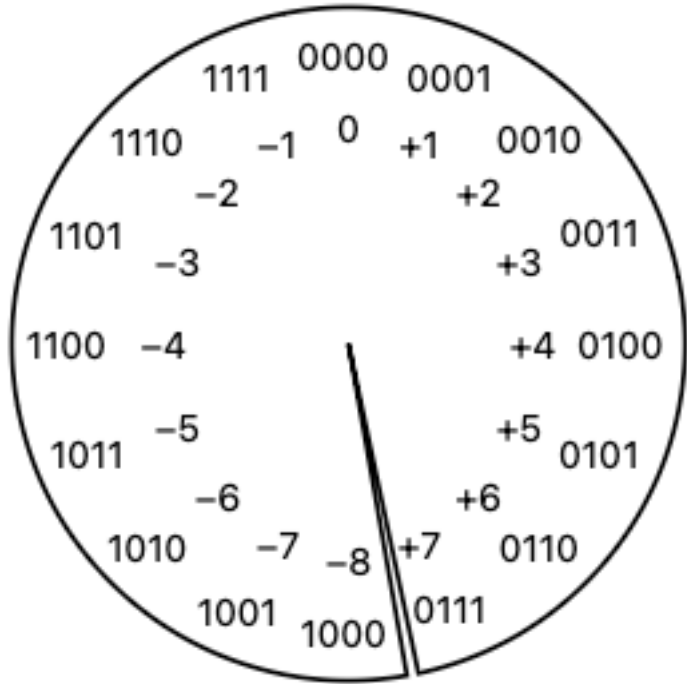


UNIVERSITY
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ENGINEERING

REVIEW

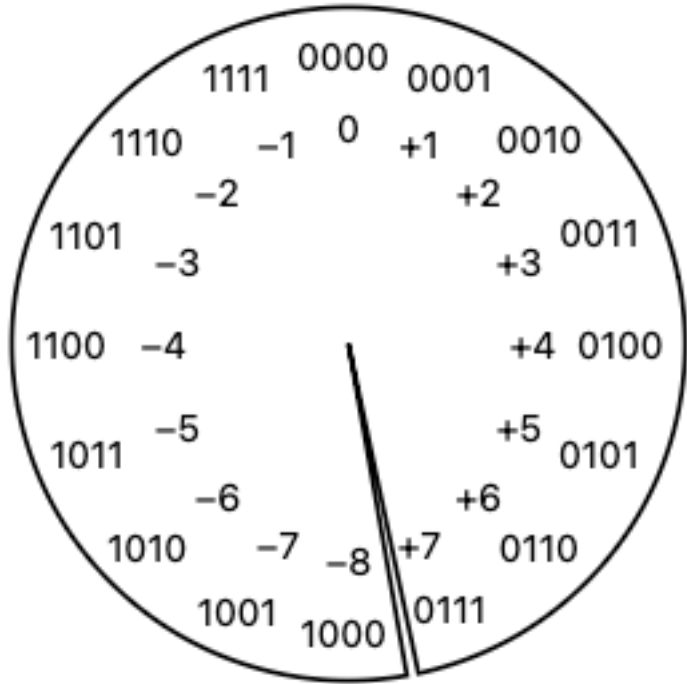
TWO COMPLEMENT



Two's complement picks a number (typically half of the maximum number we can write, rounded up) and decides that that number and all numbers bigger and including it are negative

Flip the bits and Add on

TWO COMPLEMENT



Flip the bits and add on trick for converting between positive and negative numbers?

EXAM REVIEW FALL 2018

The following assume 8-bit 2's-complement numbers. For each number, bit 0 is the low-order bit, bit 7 is the high-order bit.

Question 2 [2 pt]: (see above) Complete the following sum, showing your work (carry bits, etc)

$$\begin{array}{r} 0\ 0\ 1\ 1\ 0\ 0\ 1\ 1 \\ +\ 0\ 1\ 1\ 0\ 0\ 1\ 1\ 0 \\ \hline \end{array}$$

What is the result in base 10? Is it negative or positive? Would you get the same result in decimal if you had more bits 😊 ?

DEFINING OVERFLOW

If the sum of **two positive numbers** yields a **negative result**, the sum has **overflowed**.

If the sum of **two negative numbers** yields a **positive result**, the sum has **overflowed**.

Otherwise, the sum has not overflowed.

Over only exist for operation on signed numbers.

$$\begin{array}{r} 0111_2 \quad +7 \\ \& 0001_2 \quad +1 \\ \hline 1000_2 \quad -8 \end{array}$$


NOT OVERFLOW

If the sum of **two positive numbers** yields a **negative result**, the sum has **overflowed**.

If the sum of **two negative numbers** yields a **positive result**, the sum has **overflowed**.

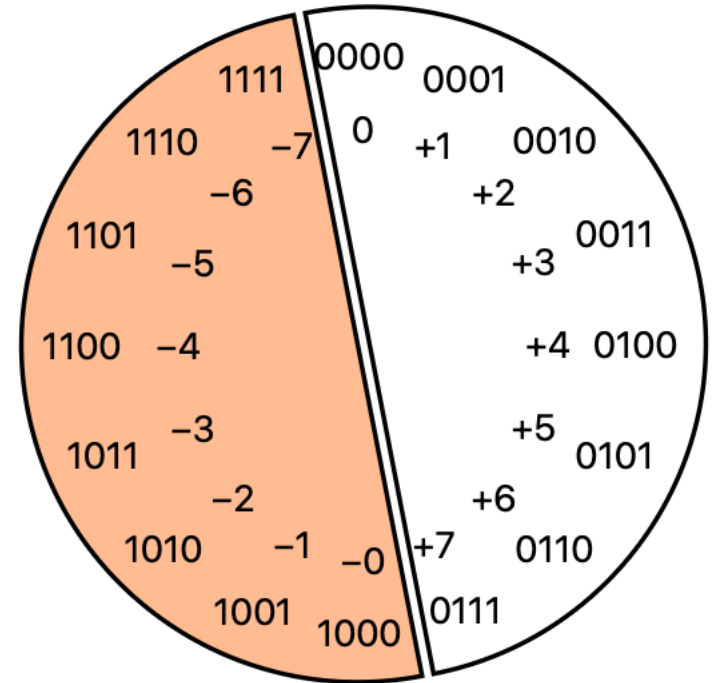
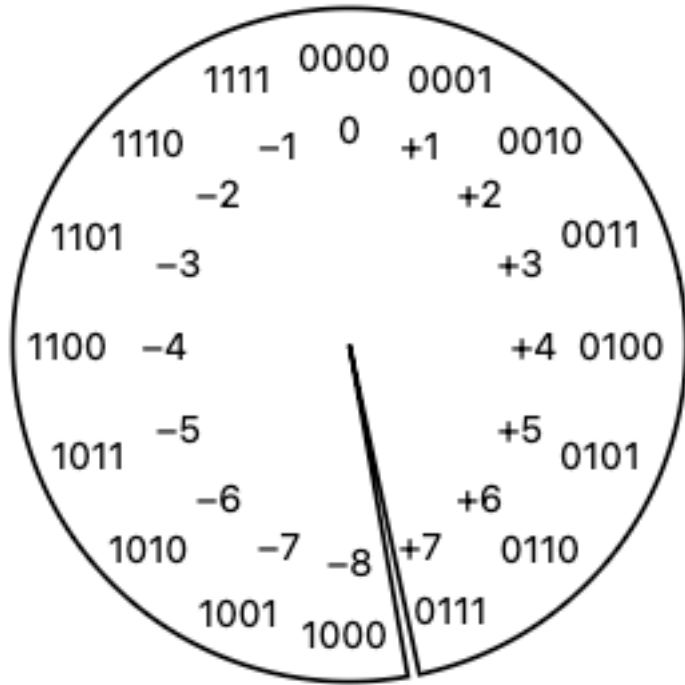
Otherwise, the sum has not overflowed.

Overflow only exists for operation on signed numbers.

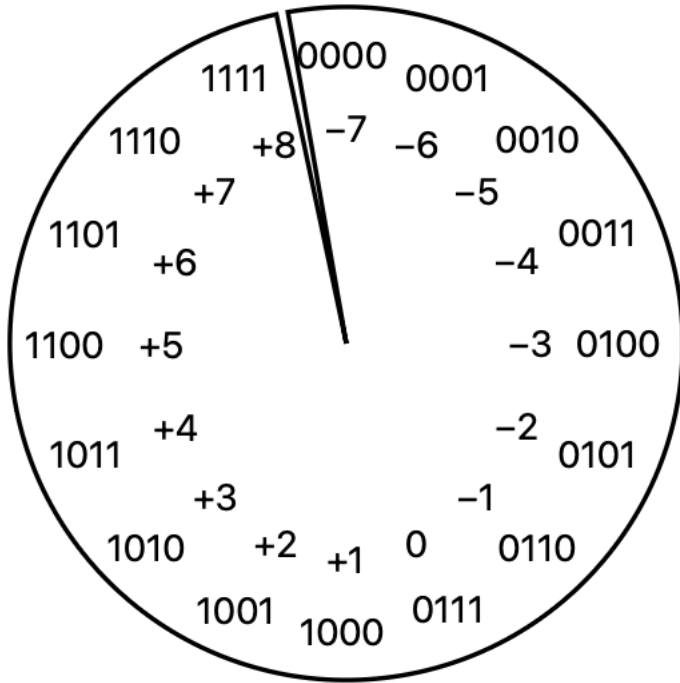
$$\begin{array}{r} 1111_2^{-1} \\ \& 0001_2^{+1} \\ \hline 10000_2 \end{array}$$


Carry Ignored, But **NOT** overflow. The answer is correctly zero

TWOS COMPLEMENT VS SIGN BIT



BIAS



From original number to BIAS

$$\text{BIAS} = \text{FLOOR} (\text{MAX_NUM}/2)$$

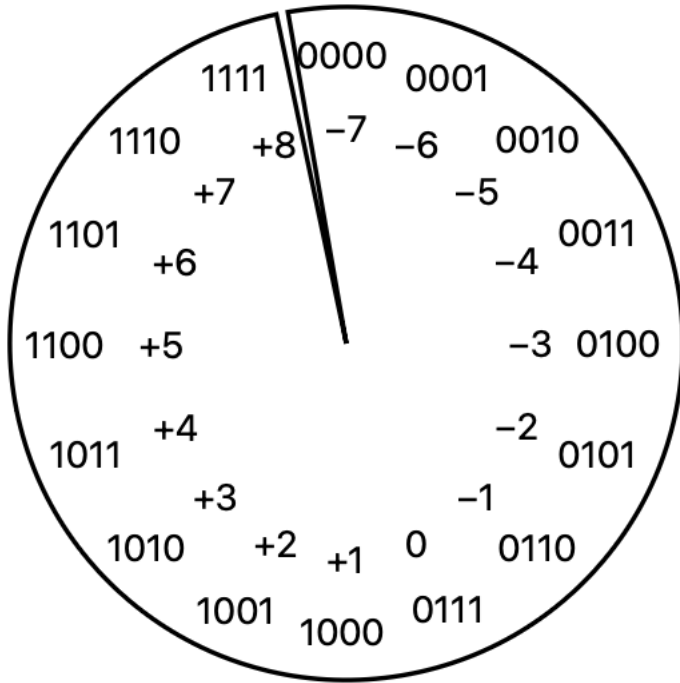
$$\text{REPRESENTATION} = \text{ORIGINAL_NUMBER} + \text{BIAS}$$

From BIAS to Original

$$\text{BIAS} = \text{FLOOR} (\text{MAX_NUM}/2)$$

$$\text{ORIGINAL_NUMBER} = \text{REPRESENTATION} - \text{BIAS}$$

BIAS EXAMPLE



From original number to BIAS

$$\text{BIAS} = \text{FLOOR} (\text{MAX_NUM}/2)$$

$$\text{REPRESENTATION} = \text{ORIGINAL_NUMBER} + \text{BIAS}$$

Example

$$\text{BIAS} = \text{FLOOR} (15/2) = 7$$

$$\text{REPRESENTATION} = -2 + 7 = 5$$

WRITING LONG BINARY IS NO FUN.
LET'S EXPRESS IT IN ANOTHER BASE TO MAKE
EASIER. DEFINITELY CHOOSE SOMETHING
LARGER THAN BASE 10

Hex Digit	Decimal	Binary
0	0	0000
1	1	0001
2	2	0010
3	3	0011
4	4	0100
5	5	0101
6	6	0110
7	7	0111
8	8	1000
9	9	1001
A	10	1010
B	11	1011
C	12	1100
D	13	1101
E	14	1110
F	15	1111

HEXADECIMAL

Convert **00101110** to hexadecimal Answer: **2E**

Group them

0010 = 2

1110 = E

Final **0x2E**

- Some programming languages uses prefixes
 - Hex: 0x
 - $0x23AB = 23AB_{16}$
 - Binary: 0b
 - $0b1101 = 1101_2$

BASE 8 OCTAL

Convert 67 to octal

$$67 \div 8 = 8 \text{ remainder } 3$$

$$8 \div 8 = 1 \text{ remainder } 0$$

$$1 \div 8 = 0 \text{ remainder } 1$$



27 (octal) to decimal

103 (octal) to decimal

(23) Strange write haha

Page 4: Binary

4. [6 points] Convert 219 into binary.

Answer

5. [6 points] What is 0b101100110111 in hexadecimal?

Answer

Page 4: Binary

4. [6 points] Convert 219 into binary.

128	64	32	16	8	4	2	1
1	1	0	1	1	0	1	1

₂

Answer

11011011

5. [6 points] What is 0b101100110111 in hexadecimal?

B 3 7

Answer

0xB37

TODAY'S LECTURE

BITWISE AND &

$$\begin{array}{r} 1100_2 \\ \& 0110_2 \\ \hline 0100_2 \end{array}$$

```
#python example  
x = 12  
y = 6  
z = x & y  
print(z)  
#prints 4
```

BITWISE OR |

$$\begin{array}{r} 1100_2 \\ | \quad 0110_2 \\ \hline 1110_2 \end{array}$$

```
#python example  
x = 12  
y = 6  
z = x | y  
print(z)  
#prints 14
```

BITWISE OR XOR ^

$$\begin{array}{r} 1100_2 \\ \wedge 0110_2 \\ \hline 1010_2 \end{array}$$

```
#python example  
x = 12  
y = 6  
z = x ^ y  
print(z)  
#prints 10
```

BIT-WISE RIGHT SHIFT

$$\begin{array}{r} 1101001_2 \\ \gg 3 \\ \hline 1101_2 \end{array}$$

```
#python example  
x = 105  
y = x >> 3  
print(y)  
#prints 13
```

SIGN EXTENSIONS

$$11000_2 \gg 2 = 11110_2$$

With Sign Extension. (The sign bit is copied)

$$11000_2 \gg 2 = 00110_2$$

Without Sign Extension

LEFT SHIFT

$$\begin{array}{r} 1101_2 \\ \ll 3 \\ \hline 1101000_2 \end{array}$$

```
#python example  
x = 13  
y = x << 3  
print(y)  
#prints 104
```

SHIFTING MULTIPLYING AND DIVIDING BY 2

A left shift is equivalent of multiplying by 2

$0001 \ll 1 = 0010$ (2).

$0001 \ll 2 = 0100$ (4)

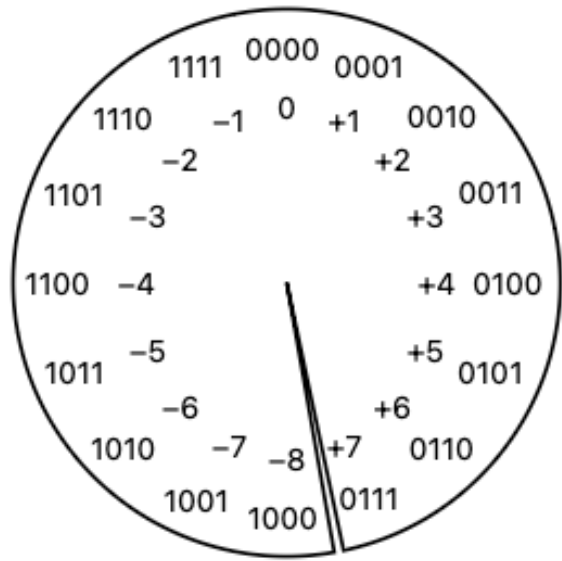
$0001 \ll 3 = 1000$ (8)

A right shift is equivalent to dividing by 2

$01000 \gg 1 = 0100$ (4)

$01000 \gg 2 = 0010$ (2)

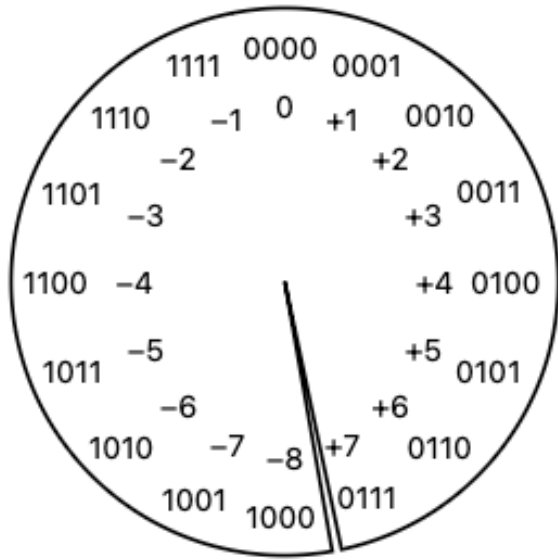
$01000 \gg 3 = 0001$ (2)



$$\begin{array}{r} \sim 0000_2 \\ \hline 1111_2 \end{array}$$

BITWISE INVERT ~

```
#python example
x = 0
z = ~x
print(z)
#prints -1
```

$$\begin{array}{r} \sim 0110_2 \\ \hline 1001_2 \end{array}$$

BITWISE INVERT ~

```
#python example
x = 6
z = ~x
print(z)
#prints -7
```

SETTING BITS TO 1

Set the last bit of this variable 1

$$\begin{array}{r} 0000_2 \\ | \quad 0001_2 \\ \hline 0001_2 \end{array}$$

```
#python example
x = 0
x = x | 0x01
print(x)
#prints 1
```

SETTING BITS TO 1

Set the third bit of x to 1

$$\begin{array}{r} 0000_2 \\ | \quad 0100_2 \\ \hline 0100_2 \end{array}$$

```
#python example
x = 0
x = x | 0x04
print(x)
#prints 4
```

Question: What if it was already one?

SETTING BITS TO 1

Set the n bit of x to 1

$$\begin{array}{r} 0000_2 \\ | \quad 0001_2 \\ \hline 1000_2 \end{array} \ll 3$$

```
#python example
x = 0
n = 3
x = x | (0x01 << n)
print(x)
#prints 8
```

Question: What if it was already one?

FLIPPING BITS

Flip the second bit of x. $1 \Rightarrow 0$ and $0 \Rightarrow 1$

$$\begin{array}{r} 1100_2 \\ \wedge \quad 0010_2 \\ \hline 1110_2 \end{array}$$

What if the nth bit was 1 instead?

FLIPPING BITS

Flip the **n**th bit of x. 1 \Rightarrow 0 and 0 \Rightarrow 1

$$\begin{array}{r} 1100_2 \\ \wedge \quad 0010_2 \\ \hline 1110_2 \end{array}$$

```
#python example
x = 12
n = 1
x = x | (0x01 << n)
print(x)
#prints 14
```

MASKING (EXTRACTING BITS)

The Idea of masking with can extra a certain section of bits by anding.

$$\begin{array}{r} 11011100_2 \\ \& 11110000_2 \\ \hline 11010000_2 \end{array}$$

Upper 4 bits extracted

MASKING (EXTRACTING BITS)

The Idea of masking with can extra a certain section of bits by anding.

$$\begin{array}{r} 11011100_2 \\ \& 00001111_2 \\ \hline 00001100_2 \end{array}$$

Lower 4 bits extracted

```
#python example
x = 220
mask = 0x0F
x = x & mask
print(x)
#prints 12
```


MASKING (EXTRACTING BITS)

The Idea of masking with can extra a certain section of bits by anding.

$$\begin{array}{r} 11011100_2 \\ \& 11110000_2 \\ \hline 11010000_2 \end{array}$$

Upper 4 bits extracted

```
#python example
x = 220
mask = ~0x0F
x = x & mask
print(x)
#prints 208
```

EXAM QUESTION FALL 2018 EXAM 1

Information for questions 6–11

Each question gives two expressions of 32-bit two's-complement integers x and y . If the two are equivalent for all x and y , write “same”; otherwise, write an example x (and y if used in the expressions) for which the two are different.

_____ add example

Question 7 [2 pt]: (see above)

$(x \ll 2) + (x \gg 1)$ and $((x \ll 3) + x) \gg 1$

Question 8 [2 pt]: (see above)

$x \mid (x \gg 1)$ and $x \wedge (x \gg 1)$

COMBINING

We can also set multiple bits simultaneously

$$\begin{array}{r} 10100000_2 \\ | \quad 00001111_2 \\ \hline 10101100_2 \end{array}$$

```
#python example
b = 0x0F
a = 0xA0
x = a | b
print(hex(x))
#prints 0xAF
```

PARITY

Suppose you want to want to calculate the even parity of x.

If the number of one's bit is number is odd the parity value is 1, otherwise it is zero

0010 parity bit is 1

0110 parity bit is 0

parity = 0 repeat 32 times:

parity ^= (x&1)

x >>= 1

PARALLEL EVALUATION

Observe that xor is both transitive and associative; thus we can re-write

$$x_0 \oplus x_1 \oplus x_2 \oplus x_3 \oplus x_4 \oplus x_5 \oplus x_6 \oplus x_7$$

using transitivity as

$$x_0 \oplus x_4 \oplus x_1 \oplus x_5 \oplus x_2 \oplus x_6 \oplus x_3 \oplus x_7$$

and using associativity as

$$(x_0 \oplus x_4) \oplus (x_1 \oplus x_5) \oplus (x_2 \oplus x_6) \oplus (x_3 \oplus x_7)$$

and then compute the contents of all the parentheses at once via

$$x \wedge (x \gg 4).$$

PARALLEL EVALUATION

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and using associativity as

$$(x_0 \oplus x_4) \oplus (x_1 \oplus x_5) \oplus (x_2 \oplus x_6) \oplus (x_3 \oplus x_7)$$

and then compute all at once via

$$x \wedge (x \gg 4).$$

$$x \wedge = (x \gg 16)$$

$$x \wedge = (x \gg 8)$$

$$x \wedge = (x \gg 4)$$

$$x \wedge = (x \gg 2)$$

$$x \wedge = (x \gg 1)$$

$$\text{parity} = (x \& 1)$$

