

Solution for PS3-1.

- a. $(2m) + (2n) = 2(m + n)$, which is even.
- b. $(2m + 1) + (2n + 1) = 2(m + n + 1)$, which is even.
- c. $(2m + 1) + (2n) = 2(m + n) + 1$, which is odd.
- d. $(2m) \cdot (n) = 2(mn)$, which is even.
- e. $(2m + 1) \cdot (2n + 1) = 2(2mn + m + n) + 1$, which is odd.

Solution for PS3-2.

- a. Since $3^b = 3^d$, we can simplify the equation to $2^a = 2^c$, which makes $a = c$. Thus, $(a, b) = (c, d)$.
- b. Rearranging the equation, $2^{a-c} = 3^{d-b}$. So let $p = a - c, q = d - b$. By assumption, $q > 0$, so $q \in \mathbb{N}$.
- c. Each time we multiply an odd integer by three, the result must be odd. Since 3^q is obtained by starting with one (an odd integer) and multiplying by three q times, the final result must be odd.
- d. We have that $2^p = 3^q$, which is odd. If $p < 0$, $2p$ is not even an integer. If $p > 0$, $2^p = 2 \times 2^{p-1}$, which is two times an integer, so it is even. The only way that 2^p can be odd is if $p = 0$.
- e. Since $p = a - c = 0$, we have $a = c$. Thus, $2^a = 2^c$ and the original equation simplifies to $3^b = 3^d$, which implies $b = d$. Thus, $(a, b) = (c, d)$.
- f. We have shown that $f(a, b) = f(c, d) \implies (a, b) = (c, d)$, so by definition, f is injective.

Solution for PS3-3. If A is countable, then we know there exists an injection $g : A \rightarrow \mathbb{N}$ by the definition of countability. If B is a subset of A we can define a function $f : B \rightarrow \mathbb{N}$ with $f(x) = g(x)$. To prove that B is countable we will show that f is an injection.

Let x_1 and x_2 be any two elements of B . If $f(x_1) = f(x_2)$, then $g(x_1) = g(x_2)$. But g is injective, so this implies $x_1 = x_2$. Therefore f is also injective.

Solution for PS3-4.

- a. If A and B are both countable then there are injections $g : A \rightarrow \mathbb{N}$ and $h : B \rightarrow \mathbb{N}$. We can define a function $f : (A \cup B) \rightarrow \mathbb{N}$ by the following:

$$f(x) = \begin{cases} 2g(x) + 1, & \text{if } x \in A, \\ 2h(x), & \text{otherwise (i.e., } x \in (B - A)). \end{cases}$$

We can show this function is an injection. Let x_1, x_2 be elements of $A \cup B$ with $f(x_1) = f(x_2)$. Either $f(x_1)$ will be odd, or it will be even.

If $f(x_1)$ is odd, then $f(x_1) = 2g(x_1) + 1$ (note that $2h(x)$ can never be odd because $h(x)$ is a natural number). Similarly, $f(x_2)$ is odd (it is equal to $f(x_1)$ after all) so we have $f(x_2) = 2g(x_2) + 1$. This means $2g(x_1) + 1 = 2g(x_2) + 1$ which implies $g(x_1) = g(x_2)$. But we know g is an injective function, so $x_1 = x_2$.

If $f(x_1)$ is even, then $f(x_1) = 2h(x_1)$ (again note that $2g(x) + 1$ can never be even because $g(x)$ is a natural number). We also know $f(x_2)$ is even so it equals $2h(x_2)$, therefore $2h(x_1) = 2h(x_2)$. This in turn implies $h(x_1) = h(x_2)$. We know h is an injection so $x_1 = x_2$.

So in either case we know if $f(x_1) = f(x_2)$ then $x_1 = x_2$. Therefore f is an injective function from $A \cup B$ to \mathbb{N} and there $A \cup B$ must be countable.

- b. If A and B are both countable then there are injections $g : A \rightarrow \mathbb{N}$ and $h : B \rightarrow \mathbb{N}$. As in part a, we can define a function $f : (A \times B) \rightarrow \mathbb{N}$ by $f(a, b) = 2^{g(a)}3^{h(b)}$. As before we can prove this is an injection. Let $f(a_1, b_1) = f(a_2, b_2)$, so $2^{g(a_1)}3^{h(b_1)} = 2^{g(a_2)}3^{h(b_2)}$. From the uniqueness of prime factorization we know that this implies $g(a_1) = g(a_2)$ and $h(a_1) = h(a_2)$. But g and h are both injections so $a_1 = a_2$ and $b_1 = b_2$. This means the ordered pairs (a_1, b_1) and (a_2, b_2) are equal to each other. Therefore f is an injective function and $A \times B$ is countable.

Solution for PS3-5. Let's repeatedly invoke the fact " A and B countable $\implies A \times B$ countable".

Using it with $A = \mathbb{N}$ and $B = \mathbb{N}$ tells us that $\mathbb{N} \times \mathbb{N}$ is countable.

Now, using it with $A = \mathbb{N} \times \mathbb{N}$ and $B = \mathbb{N}$ tells us that $\mathbb{N} \times \mathbb{N} \times \mathbb{N}$ is countable.

Next, using it with $A = \mathbb{N} \times \mathbb{N} \times \mathbb{N}$ and $B = \mathbb{N}$ tells us that $\mathbb{N} \times \mathbb{N} \times \mathbb{N} \times \mathbb{N}$ is countable.

Finally, using it with $A = \mathbb{N} \times \mathbb{N} \times \mathbb{N} \times \mathbb{N}$ and $B = \mathbb{N}$ tells us that $\mathbb{N} \times \mathbb{N} \times \mathbb{N} \times \mathbb{N} \times \mathbb{N}$ is countable.

Solution for PS3-6. Since A is infinite, there exists some element $a_0 \in A$.

Since A is infinite, it has at least two elements, so there exists some element $a_1 \in A$ distinct from a_0 .

Since A is infinite, it has at least three elements, so there exists some element $a_3 \in A$ distinct from a_0, a_1 .

Since A is infinite, it has at least four elements, so there exists some element $a_4 \in A$ distinct from a_0, a_1, a_2 .

Proceeding in this fashion, for each $n \in \mathbb{N}$ we have an element a_n distinct from all elements a_m where $m < n$.

Now define a function $f : A \rightarrow \mathbb{N}$ as follows.

$$f(x) = \begin{cases} n, & \text{if } x = a_n \text{ for some } n \in \mathbb{N}, \\ 0, & \text{otherwise (i.e., } x \text{ is not in the list } (a_0, a_1, a_2, \dots)). \end{cases}$$

This function is surjective because, given any $n \in \mathbb{N}$, there exists the element $a_n \in A$ for which $f(a_n) = n$.

Solution for PS3-7. Since S is finite, given any particular length ℓ , there are only a finite number of ℓ -length strings in S^* . (To be precise, there are $|S|^\ell$ such strings, though we don't need this fact.) Therefore, we can list all the elements of S^* as follows:

empty string, followed by
all strings of length 1 in some arbitrary order, followed by
all strings of length 2 in some arbitrary order, followed by
all strings of length 3 in some arbitrary order, followed by

This listing implicitly defines a bijection $f : \mathbb{N} \rightarrow S^*$, proving that S^* is countable.

Solution for PS3-8. Every Python program is just a string over a certain alphabet: say, the alphabet of all Unicode characters.

Thus, the set of all Python programs is a subset of a countable set, so it is itself a countable set.

Solution for PS3-9. Let $I = (0, 1)$. We define $f(x, y)$ as follows, for $(x, y) \in I \times I$. Let

$$\begin{aligned} x &= 0.a_1a_2a_3\cdots, \\ y &= 0.b_1b_2b_3\cdots \end{aligned}$$

be the unique decimal representations of x and y , as defined in class. Now construct the number

$$z = 0.a_1b_1a_2b_2a_3b_3\cdots.$$

The sequence of digits in this definition of z has infinitely many non-9s, so it is a legit decimal representation of a real number in I . Set $f(x, y) = z$.

Students: You should write up a formal proof that f is indeed an injection.

Solution for PS3-10. Consider the function $g : (0, 1] \rightarrow (0, 1)$ defined by:

$$g(x) = \begin{cases} \frac{x}{2}, & \text{if } \exists n \in \mathbb{Z} \text{ such that } x = 2^{-n}, \\ x, & \text{otherwise.} \end{cases}$$

To prove this is a bijection we will construct an inverse function. Define $h(x) : (0, 1) \rightarrow (0, 1]$ to be:

$$h(x) = \begin{cases} 2x, & \text{if } \exists n \in \mathbb{Z} \text{ such that } x = 2^{-n}, \\ x, & \text{otherwise.} \end{cases}$$

To prove that g and h are inverses of each other (which, in turn, shows that g is a bijection) we must show that $g \circ h = \text{id}_{(0,1)}$ and $h \circ g = \text{id}_{(0,1]}$.

To show $g \circ h = \text{id}_{(0,1)}$, let x be any element of $(0, 1)$. We must show $g(h(x)) = x$. If $x = 2^{-n}$ for some integer n , then $g(h(x)) = g(2x)$, but $2x = 2^{1-n}$ and $(n-1)$ is an integer, so $g(2x) = \frac{2x}{2} = x$. If instead $x \neq 2^{-n}$ for every integer n , we have $g(h(x)) = g(x) = x$. In either case $g(h(x)) = x$ so $g \circ h = \text{id}_{(0,1)}$.

Similarly To show $h \circ g = \text{id}_{(0,1]}$, we let $x \in (0, 1]$. Once again we have two cases either $x = 2^{-n}$ for some integer n , or $x \neq 2^{-n}$ for every integer n . In the former case, $h(g(x)) = h(\frac{x}{2})$. However $\frac{x}{2} = 2^{-(n+1)}$ and $(n+1)$ is an integer so $h(\frac{x}{2}) = 2 \cdot \frac{x}{2} = x$. In the latter case $h(g(x)) = h(x) = x$. Therefore $h(g(x))$ is always x and $h \circ g = \text{id}_{(0,1]}$.

So g and h are inverse functions of each other. Because g is a function with an inverse function it must be a bijection.