## Extended Euclidean Algorithm

An algorithm that we often use in cryptography is the extended Euclidean algorithm. Here we will explain some special issues that let us to construct an efficient algorithm. If we have two integers a and b, and we need to calculate the greatest common divisor we will do the following procedure. For convenience let us consider that  $r_0 = b$  and  $r_1 = a$ .

$$\begin{array}{rcl} r_0 & = & r_1q_1 + r_2 \\ r_1 & = & r_2q_2 + r_3 \\ r_2 & = & r_3q_3 + r_4 \\ \dots & = & \dots \\ r_{i-2} & = & r_{i-1}q_{i-1} + r_i \end{array}$$

This means that in each iteration we calculate the remainder  $r_i$  and the quotient  $q_{i-1}$ . We must remember that our goal is not only to find the gcd, but also the linear combination i.e.  $gcd(r_0, r_1) = sr_0 + tr_1$ . For this purpose, we express the remainder of each ith iteration as a linear combination of  $r_0$  and  $r_1$  i.e.  $r_i = s_i r_0 + t_i r_1$ .

Notice that to construct the linear combination of the ith iteration we need the previous linear combination, thus

$$r_{i-2} = s_{i-2}r_0 + t_{i-2}r_1 \tag{1}$$

$$r_{i-1} = s_{i-1}r_0 + t_{i-1}r_1 \tag{2}$$

Also notice that in the ith iteration we calculate the quotient  $q_{i-1}$  and the remainder  $r_i$  i.e.  $r_{i-2} = r_{i-1}q_{i-1} + r_i$ , if we rearrange this equation we have

$$r_i = r_{i-2} - r_{i-1}q_{i-1}. (3)$$

Now we can use equations 1 and 2 and we use them in 3

$$r_i = s_{i-2}r_0 + t_{i-2}r_1 - (s_{i-1}r_0 + t_{i-1}r_1)q_{i-1}$$

Rearraging we have

$$r_i = (s_{i-2} - s_{i-1}q_{i-1})r_0 + (t_{i-2} - t_{i-1}q_{i-1})r_1$$

This give us a way to calculate  $s_i$  and  $t_i$ :

$$s_i = s_{i-2} - s_{i-1}q_{i-1}$$
  
 $t_i = t_{i-2} - t_{i-1}q_{i-1}$ 

These equations for  $s_i$  and  $t_i$  are valid for  $i \geq 2$ . It is not difficult to see that which are the values for  $s_0, s_1$  and  $t_0, t_1$ . Actually  $s_0 = 1, s_1 = 0, t_0 = 0, t_1 = 1$ . Finally you must notice that to calculate the inverse, we do not require the value of  $s_i$ . All we need is to calculate  $t_i$ . Keeping all these facts we can write down a pseudocode to find the extended Euclidean algorithm.