

Math 206B Introduction to Numerical Analysis  
notes based off "Numerical Analysis by R.L.  
Burden and J.D. Faires" Section 1 which is  
Numerical Solution of Nonlinear Equations

Charlie Seager

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**2.1 The Bisection Method**

**Bisection Technique** The first technique based on the Intermediate Value Theorem, is called the Bisection or Binary Search method.

**Theorem 2.1** Suppose that  $f \in C[a, b]$  and  $f(a) \cdot f(b) < 0$ . The Bisection method generates a sequence  $\{p_n\}_{n=1}^{\infty}$  approximating a zero  $p$  of  $f$  with

$$|p_n - p| \leq \frac{b-a}{2^n}, \text{ when } n \geq 1$$

**Chapter 2.2 Fixed Point Iteration** A fixed point for the function is a number at which the value of the function does not change when the function is applied.

**Definition 2.2** The number  $p$  is a fixed point for a given function  $g$  if  $g(p) = p$ .

**Theorem 2.3** (i) If  $g \in C[a, b]$  and  $g(x) \in [a, b]$  for all  $x \in [a, b]$ , then  $g$  has at least one fixed point in  $[a, b]$ .

(ii) If, in addition  $g'(x)$  exists on  $(a, b)$  and a positive constant  $k < 1$  exists with

$$|g'(x)| \leq k, \text{ for all } x \in (a, b)$$

then there is exactly one fixed point in  $[a, b]$ .

**Theorem 2.4 (Fixed Point Theorem)** Let  $g \in C[a, b]$  be such that  $g(x) \in [a, b]$  for all  $x$  in  $[a, b]$ . Suppose in addition that  $g'$  exists on  $(a, b)$  and that a constant  $0 < k < 1$  exists with

$$|g'(x)| \leq k, \text{ for all } x \in (a, b)$$

Then for any number  $p_0$  in  $[a, b]$ , the sequence defined by

$$p_n = g(p_{n-1}), n \geq 1.$$

converges to the unique fixed point  $p$  in  $[a, b]$ .

**Corollary 2.5** If  $g$  satisfies the hypotheses of Th. 2.4, then bounds for the error involved in using  $p_n$  to approximate  $p$  are given by

$$|p_n - p| \leq k^n \max\{p_0 - a, b - p_0\}$$

and

$$|p_n - p| \leq \frac{k^n}{1-k} |p_1 - p_0|, \text{ for all } n \geq 1$$

### Chapter 2.3 Newtons Method and Its Extensions

Newtons (or the Newton-Raphson) method is one of the most powerful and well known numerical methods for solving a root-finding problem. There are many ways of introducing Newton's Method.

**Theorem 2.6** Let  $f \in C^2[a, b]$ . If  $p \in (a, b)$  is such that  $f(p) = 0$  and  $f'(p) \neq 0$  then there exists a  $\delta > 0$  such that Newton's Method generates a sequence  $\{p_n\}_{n=1}^{\infty}$  converging to  $p$  for any initial approximations  $p_0 \in [p-\delta, p+\delta]$ .

**Chapter 2.4 Error Analysis for Iterative Methods** In this section we investigate the order of convergence of functional iteration schemes and as a means of obtaining rapid convergence, rediscover Newton's method. We also consider ways of accelerating the convergence of Newton's method in special circumstances. First, however, we need a new procedure for measuring how rapidly a sequence converges.

**Defintion 2.7 Order of convergence** Suppose  $\{p_n\}_{n=0}^{\infty}$  is a sequence that converges to  $p$ , with  $p_n \neq p$  for all  $n$ . If positive constants  $\lambda$  and  $\alpha$  exist with

$$\lim_{n \rightarrow \infty} \frac{|p_{n+1} - p|}{|p_n - p|^\alpha} = \lambda$$

then  $\{p_n\}_{n=0}^{\infty}$  converges to  $p$  of order  $\alpha$  with asymptotic error constant  $\lambda$ .

**Theorem 2.8** Let  $g \in C[a, b]$  be such that  $g(x) \in [a, b]$  for all  $x \in [a, b]$ . Suppose, in addition that  $g'$  is continuous on  $(a, b)$  and a positive constant  $k < 1$  exists with

$$|g'(x)| \leq k, \text{ for all } x \in (a, b)$$

If  $g'(p) \neq 0$ , then for any number  $p_0 \neq p$  in  $[a, b]$ , the sequence

$$p_n = g(p_{n-1}) \text{ for } n \geq 1$$

converges only linearly to the unique fixed point  $p$  in  $[a, b]$

**Theorem 2.9** Let  $p$  be a solution of the equation  $x = g(x)$ . Suppose that  $g'(p) = 0$  and  $g''$  is continuous with  $|g''(x)| < M$  on an open interval  $I$  containing  $p$ . Then there exists a  $\delta > 0$  such that for  $p_0 \in [p-\delta, p+\delta]$ , the sequence defined by  $p_n = g(p_{n-1})$  when  $n \geq 1$ , converges at least quadratically to  $p$ . Moreover for sufficiently large values of  $n$ .

$$|p_{n+1} - p| < \frac{M}{2} |p_n - p|^2$$

**Defintion 2.10** A solution  $p$  of  $f(x) = 0$  is a zero of multiplicity  $m$  of  $f$  if for  $x \neq p$  we can write  $f(x) = (x - p)^m q(x)$ , where  $\lim_{x \rightarrow p} q(x) \neq 0$ .

**Theorem 2.11** The function  $f \in C^1[a, b]$  has a simple zero at  $p$  in  $(a, b)$  if and only if  $f(p) = 0$ , but  $f'(p) \neq 0$

**Theorem 2.12** The function  $f \in C^m[a, b]$  has a zero of multiplicity  $m$  at  $p$  in  $(a, b)$  if and only if

$$0 = f(p) = f'(p) = f''(p) = \dots = f^{(m-1)}(p), \text{ but } f^{(m)}(p) \neq 0$$

**Chapter 2.5 Accelerating Convergence** Theorem 2.8 indicates that it is rare to have the luxury of quadratic convergence. We now consider a technique called Aitken's  $\Delta^2$  method that can be used to accelerate the convergence of a sequence that is linearly convergent, regardless of its origin or application.

**Definition 2.13** For a given sequence  $\{p_n\}_{n=0}^\infty$  the forward difference  $\Delta p_n$  (read "deltap<sub>n</sub>") is defined by

$$\Delta p_n = p_{n+1} - p_n \text{ for } n \geq 0$$

Higher powers of the operator  $\Delta$  are defined recursively by

$$\Delta^k p_n = \Delta(\Delta^{k-1} p_n) \text{ for } k \geq 2$$

**Theorem 2.14** Suppose that  $\{p_n\}_{n=0}^\infty$  is a sequence that converges linearly to the limit  $p$  and that

$$\lim_{n \rightarrow \infty} \frac{p_{n+1} - p}{p_n - p} < 1$$

Then the Aitken's  $\Delta^2$  sequence  $\{\hat{p}_n\}_{n=0}^\infty$  converges to  $p$  faster than  $\{p_n\}_{n=0}^\infty$  in the sense that

$$\lim_{n \rightarrow \infty} \frac{\hat{p}_n - p}{p_n - p} = 0.$$

**Theorem 2.15** Suppose that  $x = g(x)$  has the solution  $p$  with  $g'(p) \neq 1$ . If there exists a  $\delta > 0$  such that  $g \in C^3[p - \delta, p + \delta]$ , then Steffensen's method gives quadratic convergence for any  $p_0 \in [p - \delta, p + \delta]$

## Chapter 2.6 Zeros of Polynomials and Mullers Method

### Theorem 2.16 (Fundamental Theorem of Algebra)

If  $P(x)$  is a polynomial of degree  $n \geq 1$  with real or complex coefficients, the  $P(x) = 0$  has at least one (possibly complex) root.

**Corollary 2.17** If  $P(x)$  is a polynomial of degree  $n \geq 1$  with real or complex coefficients, then there exist unique constants  $x_1, x_2, \dots, x_k$  possibly complex and unique positive integers  $m_1, m_2, \dots, m_k$  such that  $\sum_{i=1}^k m_i = n$  and

$$P(x) = a_n(x - x_1)^{m_1}(x - x_2)^{m_2} \dots (x - x_k)^{m_k}$$

**Corollary 2.18** Let  $P(x)$  and  $Q(x)$  be polynomials of degree at most  $n$ . If  $x_1, x_2, \dots, x_k$  with  $k \leq n$ , are distinct numbers with  $P(x_i) = Q(x_i)$  for  $i = 1, 2, \dots, k$  then  $P(x) = Q(x)$  for all values.

### Theorem 2.19 (Horner's Method)

Let

$$P(x) = a_n X^n + a_{n-1} X^{n-1} + \dots + a_1 x + a_0$$

Define  $b_n = a_n$  and

$$b_k = a_k + b_{k+1}x_0. \text{ for } k = n-1, n-2, \dots, 1, 0.$$

Then  $b_0 = P(x_0)$ . Moreover if

$$Q(x) = b_n x^{n-1} + b_{n-1} x^{n-2} + \dots + b_2 x + b_1.$$

then

$$P(x) = (x - x_0)Q(x) + b_0$$

**Theorem 2.20** If  $z = bi$  is a complex zero of multiplicity  $m$  of the polynomial  $P(x)$  with real coefficients then  $\bar{z} = a - bi$  is also a zero of multiplicity  $m$  of the polynomial  $P(x)$  and  $(x^2 - 2ax + a^2 + b^2)^m$  is a factor of  $P(x)$ .

#### **Chapter 2.7 Survey of Methods and Software**