

## CS308 Homework 2

Exercises for Algorithm Design and Analysis by Li Jiang, 2016 Autumn Semester

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**Coverage** : Syntax Analysis.

1. (Section 4.2, Exercises 4.2.1) Consider the context-free grammar:

$$S \rightarrow S S + \mid S S * \mid a$$

and the string  $aa + a^*$ .

- (a) Give a leftmost derivation for the string.
- (b) Give a rightmost derivation for the string.
- (c) Give a parse tree for the string.

**Solution.**

(a)  $S \Rightarrow S S * \Rightarrow S S + S * \Rightarrow a S + S * \Rightarrow a a + S * \Rightarrow a a + a *$

(b)  $S \Rightarrow S S * \Rightarrow S a * \Rightarrow S S + a * \Rightarrow S a + a * \Rightarrow a a + a *$

(c) below is the parse tree

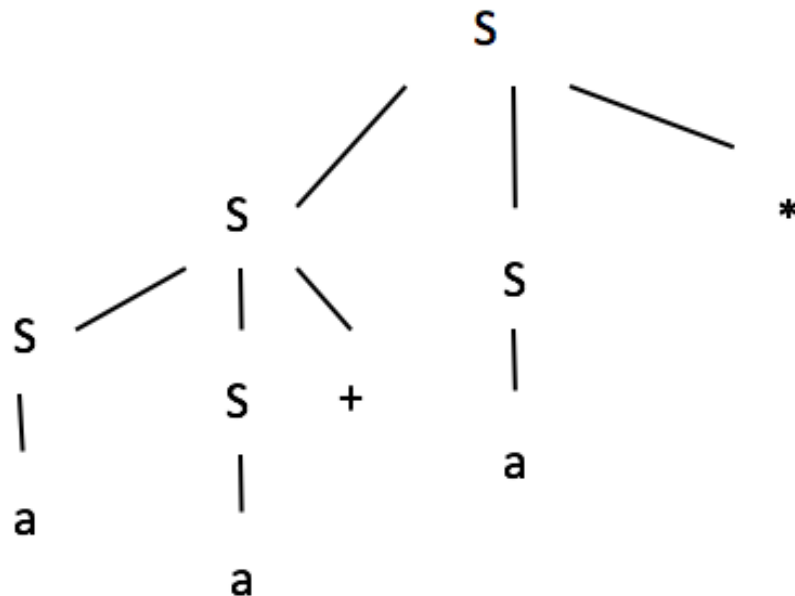


图 1: parse tree

□

2. (Section 4.3, Exercises 4.3.1) The following is a grammar for regular expressions over symbols  $a$  and  $b$  only, using  $+$  in place of  $|$  for union, to avoid conflict with the use of vertical bar as a metasympol in grammars:

$$\begin{aligned} \text{rexpr} &\rightarrow \text{rexpr} + \text{rterm} \mid \text{rterm} \\ \text{rterm} &\rightarrow \text{rterm} \text{rfactor} \mid \text{rfactor} \\ \text{rfactor} &\rightarrow \text{rfactor} * \mid \text{rprimary} \\ \text{rprimary} &\rightarrow a \mid b \end{aligned}$$

- (a) Left factor this grammar.
- (b) Does left factoring make the grammar suitable for top-down parsing?
- (c) In addition to left factoring, eliminate left recursion from the original grammar.
- (d) Is the resulting grammar suitable for top-down parsing?

**Solution.**

- (a) There is no common non-empty prefix, so we can't do left factoring, this grammar stays the same.
- (b) No. Because this is the left-recursive grammar. If we use top-down parsing, it will get into an infinite loop.
- (c) Below is the grammar after left recursion elimination.

$$\begin{aligned} \text{rexpr} &\rightarrow \text{rterm} A \\ A &\rightarrow + \text{rterm} A \mid \epsilon \\ \text{rterm} &\rightarrow \text{rfactor} B \\ B &\rightarrow \text{rfactor} B \mid \epsilon \\ \text{rfactor} &\rightarrow \text{rprimary} C \\ C &\rightarrow * C \mid \epsilon \\ \text{rprimary} &\rightarrow a \mid b \end{aligned}$$

- (d) Yes. Because there is no left recursion.

□

3. (Section 4.4, Exercises 4.4.3) Compute FIRST and FOLLOW for the grammar:

(a)  $S \rightarrow 0 S 1 \mid 0 1$  (with string 000111).

(b)  $S \rightarrow + S S \mid * S S \mid a$  (with string  $+ * aaa$ ).

**Solution.**

(a)  $\text{FIRST}(S) = [ 0 ], \text{FOLLOW}(S) = [ 1, \$ ]$ .

(b)  $\text{FIRST}(S) = [ +, *, a ], \text{FOLLOW}(S) = [ +, *, a, \$ ]$ .

□